



**Q2. (20 pts.)** Consider a relation R with attributes {A, B, C, D, E, F, G, H} and the set of functional dependencies  $F = \{G \rightarrow FD, E \rightarrow D, GD \rightarrow CE, BD \rightarrow A\}$

- a) Find the minimal cover of the given set of functional dependencies F. **(10 points)**
- b) Decompose the relation into a collection of 3NF relations using the minimal cover found in part a). **(5 points)**
- c) Is your decomposition also BCNF, lossless-join and dependency preserving as well? Explain. **(3 points)**
- d) In designing a relational database schema why might we choose a non BCNF but 3NF design? **(2 points)**

**Q3. (8 pts.)** Consider the relation R(A,B,C,D,E) with multi-valued dependencies  $MD = \{A \twoheadrightarrow B, AB \twoheadrightarrow C\}$  and functional dependencies  $FD = \{A \rightarrow D, AB \rightarrow E\}$ .

- a) Find key(s) for the relation? **(3 points)**
- b) Decompose the relation into a collection of relation schemas in 4NF. **(5 points)**

**Q4. (12 pts.)** For each of the following schedules determine whether it is conflict-serializable and/or view-serializable. Justify your answers.

a)  $W_3(C) W_2(B) R_2(B) R_3(C) W_1(C) R_2(C)$

Is it conflict-serializable? Why or why not? If it is give the equivalent serial schedule.

Is it view-serializable? Why or why not? If it is give the equivalent serial schedule.

b)  $R_2(C) W_1(C) W_2(C) R_1(B) W_3(C)$

Is it conflict-serializable? Why or why not? If it is, give the equivalent serial schedule.

Is it view-serializable? Why or why not? If it is, give the equivalent serial schedule.

**Q5. (20 pts.)** Consider the following schedule involving three transactions T1, T2 and T3:

Time	T1	T2	T3
1	R(C)		
2			R(B)
3	R(A)		
4		R(C)	
5			W(B)
6	W(A)		
7		R(B)	
8			R(A)
9		W(C)	
10		W(B)	
11	Commit		
12		Commit	
13			Commit

- Is this schedule a recoverable schedule? Justify your answer. **(5 points)**
- Describe how the strict two-phase locking with deadlock detection would handle the schedule. **(10 points)**
- Assuming T1 is the oldest and T3 is the youngest transaction, describe how the **strict two-phase locking with wait-die deadlock prevention scheme** would handle the schedule. **(5 points)**

**Q6 (20 pts.)** Consider the following relational table:

**Enroll**(student\_id, course\_no, status)

- [illegible]

**Q7. (25 pts.)** Consider a database with data items {A, B, C, D}. The system has the following log records. Note that an entry <T, X, old, new> means transaction T changes the value of X from old to new. In the following log, DPT represents the Dirty Page Table and TT represents the Transaction Table.

1. <T2, C, 10, 11>
2. <T2 commit>
3. <T1, B, 20, 21>
4. <begin checkpoint>  
DPT = (B, 3)  
TT = (T1, running, 3)  
<end checkpoint>
5. <T3, D, 30, 31>
6. <T4, C, 11, 41>
7. <T4 commit>
8. <T3, A, 50, 51>
9. <T1, D, 31, 32>
10. <T5, A, 51, 52>
11. <T3 commit>
12. <T6, B, 21, 33>
13. <T5 commit>

**CRASH**

- a) What is the smallest Log sequence number (LSN) accessed during the Analysis phase? (2 pts.)

Answer: \_\_\_\_\_

- b) Fill in the contents of the **Dirty Page Table** and the **Transaction Table** at the end of the Analysis phase. (8 pts.)

**DPT**

PageID	RecLSN

**TT**

XID	Status	LastLSN

- c) What is the first Log sequence number (LSN) to be redone during the Redo phase? (2 pts.)

Answer: \_\_\_\_\_

- d) List all possible values of A, B, C and D on disk at the time of crash (after action 13)? Assume 3 buffer pages are available. (4 points)

- e) What log entries are written to the log during Undo phase? (5 points)

- f) What are the values of A, B, C, D on disk at the end of crash recovery? (4 points)