Содержание

1	.1 CMake	1 1 1
2 I	m Bugs	1
	·	2
4 N	Numbers	2
	•	3
6 F	Push-free segment tree	5
7	7.1 Chinese remainder theorem without overflows	6 6 7
8 5	Suffix Automaton	7
1 1.1	Setup & Scripts CMake	
	ke_minimum_required(VERSION 3.14) ject(olymp)	
add #se	(CMAKE_CXX_STANDARD 17) _compile_definitions(LOCAL) t(CMAKE_CXX_FLAGS "\${CMAKE_CXX_FLAGS} -fsanitize=undefined -fno- → sanitize-recover") nitizers: address, leak, thread, undefined, memory	_
	_executable(olymp f.cpp) wipe.sh	
tou	ch {al}.cpp	
for don	<pre>file in ?.cpp ; do cat template.cpp > \$file ; e</pre>	

2 Bugs

- powmod :)
- Всегда чекать Куна дважды, особенно на количество итераций

- uniform int distribution от одного параметра
- for (char c : "NEWS")
- Порядок верхних и нижних границ в случае, когда задача двумерна $t-b \neq b-t$
- static с мультитестами

3 Geometry

3.1 Пересечение прямых

$$AB = A - B; CD = C - D$$
$$(A \times B \cdot CD.x - C \times D \cdot AB.x : A \times B \cdot CD.y - C \times D \cdot AB.y : AB \times CD)$$

4 Numbers

• A lot of divisors

$$- \leq 20 : d(12) = 6$$

$$- \leq 50 : d(48) = 10$$

$$- \leq 100 : d(60) = 12$$

$$- \leq 1000 : d(840) = 32$$

$$- \leq 10^4 : d(9240) = 64$$

$$- \leq 10^5 : d(83160) = 128$$

$$- \leq 10^6 : d(720720) = 240$$

$$- \leq 10^7 : d(8648640) = 448$$

$$- \leq 10^8 : d(91891800) = 768$$

$$- \leq 10^9 : d(931170240) = 1344$$

$$- \leq 10^{11} : d(97772875200) = 4032$$

$$- \leq 10^{12} : d(963761198400) = 6720$$

$$- \leq 10^{15} : d(866421317361600) = 26880$$

$$- \leq 10^{18} : d(897612484786617600) = 103680$$

• Numeric integration

- simple:
$$F(0)$$

- simpson: $\frac{F(-1)+4\cdot F(0)+F(1)}{6}$
- runge2: $\frac{F(-\sqrt{\frac{1}{3}})+F(\sqrt{\frac{1}{3}})}{2}$
- runge3: $\frac{F(-\sqrt{\frac{3}{5}})\cdot 5+F(0)\cdot 8+F(\sqrt{\frac{3}{5}})\cdot 5}{18}$

5 Graphs

5.1 Weighted matroid intersection

```
1 // here we use T = __int128 to store the independent set
2 // calling expand k times to an empty set finds the maximum
3 // cost of the set with size exactly k,
4 // that is independent in blue and red matroids
5 // ver is the number of the elements in the matroid,
6 // e[i].w is the cost of the i-th element
7 // first return value is new independent set
8 // second return value is difference between
9 // new and old costs
10 // oracle(set, red) and oracle(set, blue) check whether
11 // or not the set lies in red or blue matroid respectively
12
13 auto expand = [\delta] (T cur_set) \rightarrow pair<T, int>
14 {
       vector<int> in(ver);
15
       for (int i = 0; i < ver; i++)
16
           in[i] = ((cur set >> i) & 1);
17
18
       const int red = 1;
19
       const int blue = 2;
20
21
       vector<vector<int>>> g(ver);
22
       for (int i = 0; i < ver; i++)</pre>
23
       for (int j = 0; j < ver; j++)
24
25
           T swp_mask = (cur_set ^(T(1) \ll i) ^(T(1) \ll j));
26
           if (!in[i] & in[j])
27
28
                if (oracle(swp_mask, red))
29
                    g[i].push back(j);
30
                if (oracle(swp_mask, blue))
31
                    g[j].push_back(i);
32
           }
33
       }
34
35
       vector<int> from, to;
36
       for (int i = 0; i < ver; i++) if (!in[i])</pre>
37
38
           T add_mask = cur_set ^(T(1) \ll i);
39
           if (oracle(add mask, blue))
40
                from.push_back(i);
41
           if (oracle(add_mask, red))
42
                to.push back(i);
43
       }
44
45
       auto get_cost = [&] (int x)
46
```

```
{
47
            const int cost = (!in[x] ? e[x].w : -e[x].w);
48
            return (ver + 1) * cost - 1;
49
       };
50
51
       const int inf = int(1e9);
52
       vector<int> dist(ver, -inf), prev(ver, -1);
53
       for (int x : from)
54
            dist[x] = get_cost(x);
55
56
       queue<int> q;
57
58
59
       vector<int> used(ver);
       for (int x : from)
60
61
62
            q.push(x);
            used[x] = 1;
63
       }
64
65
       while (!q.empty())
66
67
       {
            int cur = q.front(); used[cur] = 0; q.pop();
68
69
            for (int to : g[cur])
70
71
72
                int cost = get cost(to);
                if (dist[to] < dist[cur] + cost)</pre>
73
74
                     dist[to] = dist[cur] + cost;
75
76
                     prev[to] = cur;
                     if (!used[to])
77
78
                         used[to] = 1;
79
                         q.push(to);
80
                     }
81
                }
82
            }
83
       }
84
85
       int best = -inf, where = -1;
86
       for (int x : to)
87
88
       {
            if (dist[x] > best)
89
90
                best = dist[x];
91
                where = x;
92
93
            }
       }
94
95
       if (best = -inf)
96
```

```
return pair<T, int>(cur_set, best);
97
98
        while (where \neq -1)
99
100
             cur_set ^= (T(1) \ll where);
101
             where = prev[where];
102
         }
103
104
        while (best % (ver + 1))
105
             best++;
106
         best \neq (ver + 1);
107
108
109
         assert(oracle(cur_set, red) & oracle(cur_set, blue));
110
         return pair<T, int>(cur_set, best);
111 };
        Push-free segment tree
    6
 1 class pushfreesegtree
 2 {
        vector<modulo♦> pushed, unpushed;
 3
 4
        modulo ⇔ add(int l, int r, int cl, int cr, int v, const modulo
 5
            \leftrightarrow \Leftrightarrow \&x)
         {
 6
             if (r \le cl || cr \le l)
 7
 8
                  return 0;
             if (l ≤ cl & cr ≤ r)
 9
             {
 10
                  unpushed[v] += x;
 11
12
                  return x * (cr - cl);
13
             }
 14
15
             int ct = (cl + cr) / 2;
16
17
             auto tmp = add(l, r, cl, ct, 2 * v, x) + add(l, r, ct, cr,
18
                \hookrightarrow 2 * v + 1, x);
19
             pushed[v] += tmp;
20
21
             return tmp;
22
         }
23
24
25
        modulo ⇔ sum(int l, int r, int cl, int cr, int v)
26
27
             if (r \leq cl || cr \leq l)
28
29
                  return 0;
             if (l \leq cl \& cr \leq r)
30
```

```
return pushed[v] + unpushed[v] * (cr - cl);
31
32
           int ct = (cl + cr) / 2;
33
34
           return sum(l, r, cl, ct, 2 * v) + unpushed[v] * (min(r, cr)
35
               \rightarrow - max(l, cl)) + sum(l, r, ct, cr, 2 * v + 1);
       }
36
37
   public:
38
       pushfreesegtree(int n): pushed(2 * up(n)), unpushed(2 * up(n))
39
       {}
40
41
42
43
       modulo ⇔ sum(int l, int r)
44
           return sum(l, r, 0, pushed.size() / 2, 1);
45
       }
46
47
48
       void add(int l, int r, const modulo ♦ &x)
49
50
           add(l, r, 0, pushed.size() / 2, 1, x);
51
52
       }
53 };
```

7 Number theory

7.1 Chinese remainder theorem without overflows

```
1 // Replace T with an appropriate type!
2 using T = long long;
3
4 // Finds x, y such that ax + by = gcd(a, b).
5 T gcdext (T a, T b, T &x, T &y)
6 {
       if (b = 0)
7
8
       {
9
           x = 1, y = 0;
10
           return a;
       }
11
12
       T res = gcdext(b, a \% b, y, x);
13
       y -= x * (a / b);
14
15
       return res;
16 }
17
18 // Returns true if system x = r1 \pmod{m1}, x = r2 \pmod{m2} has
      → solutions
19 // false otherwise. In first case we know exactly that x = r \pmod{m}
      \rightarrow )
```

```
20
21 bool crt (T r1, T m1, T r2, T m2, T &r, T &m)
22 {
       if (m2 > m1)
23
24
       {
25
            swap(r1, r2);
            swap(m1, m2);
26
27
       }
28
       T g = \underline{gcd(m1, m2)};
29
       if ((r2 - r1) \% g \neq 0)
30
31
            return false;
32
       T c1, c2;
33
       auto nrem = gcdext(m1 / g, m2 / g, c1, c2);
34
35
       assert(nrem = 1);
       assert(c1 * (m1 / g) + c2 * (m2 / g) = 1);
36
       T a = c1;
37
       a *= (r2 - r1) / g;
38
       a \% = (m2 / g);
39
40
       m = m1 / g * m2;
41
       r = a * m1 + r1;
       r = r \% m;
42
       if (r < 0)
43
44
            r += m;
45
46
       assert(r % m1 = r1 \& r % m2 = r2);
47
       return true;
48 }
        Integer points under a rational line
1 // integer (x, y) : 0 \le x < n, 0 < y \le (kx + b) / d
2 // (real division)
3 // In other words, sum_{x=0}^{n-1} [(kx+b)/d]
4 ll trapezoid (ll n, ll k, ll b, ll d)
  {
5
       if (k = 0)
6
            return (b / d) * n;
7
       if (k \ge d \parallel b \ge d)
8
            return (k / d) * n * (n - 1) / 2 + (b / d) * n + trapezoid(
9
               \hookrightarrow n, k % d, b % d, d);
       return trapezoid((k * n + b) / d, d, (k * n + b) % d, k);
10
11 }
       Suffix Automaton
1 struct Vx{
2
       static const int AL = 26;
3
       int len, suf;
4
       int next[AL];
```

```
Vx(){}
5
       Vx(int l, int s):len(l), suf(s){}
6
  };
7
8
9 struct SA{
       static const int MAX LEN = 1e5 + 100, MAX V = 2 * MAX LEN;
10
       int last, vcnt;
11
       Vx v[MAX_V];
12
13
       SA(){
14
           vcnt = 1;
15
           last = newV(0, 0); // root = vertex with number 1
16
17
       int newV(int len, int suf){
18
           v[vcnt] = Vx(len, suf);
19
           return vcnt++;
20
       }
21
22
23
       int add(char ch){
24
           int p = last, c = ch - 'a';
           last = newV(v[last].len + 1, 0);
25
           while(p & !v[p].next[c]) //added p &
26
                v[p].next[c] = last, p = v[p].suf;
27
           if(!p)
28
                v[last].suf = 1;
29
           else{
30
31
                int q = v[p].next[c];
                if (v[q].len = v[p].len + 1)
32
                    v[last].suf = q;
33
34
                else{
                    int r = newV(v[p].len + 1, v[q].suf);
35
                    v[last].suf = v[q].suf = r;
36
                    memcpy(v[r].next, v[q].next, sizeof(v[r].next));
37
                    while(p & v[p].next[c] = q)
38
                        v[p].next[c] = r, p = v[p].suf;
39
                }
40
41
42
           return last;
       }
43
44 };
```















