Team Reference

SPb SU LOUD Enough 2019–2020

1 Setup & Scripts

1.1 CMake

2 Bugs

- powmod:)
- Всегда чекать Куна дважды, особенно на количество итераций
- uniform int distribution от одного параметра
- for (char c : "NEWS")
- Порядок верхних и нижних границ в случае, когда задача двумерна $t-b \neq b-t$
- static с мультитестами

3 Geometry

3.1 Пересечение прямых

$$AB = A - B; CD = C - D$$
$$(A \times B \cdot CD.x - C \times D \cdot AB.x : A \times B \cdot CD.y - C \times D \cdot AB.y : AB \times CD)$$

4 Numbers

• A lot of divisors

$$- \leq 20 : d(12) = 6$$

$$- \leq 50 : d(48) = 10$$

$$- \leq 100 : d(60) = 12$$

$$- \leq 1000 : d(840) = 32$$

$$- \leq 10^4 : d(9240) = 64$$

$$- \leq 10^5 : d(83160) = 128$$

$$- \leq 10^6 : d(720720) = 240$$

$$- \leq 10^7 : d(8648640) = 448$$

$$- \leq 10^8 : d(91891800) = 768$$

$$- \leq 10^9 : d(931170240) = 1344$$

$$- \leq 10^{11} : d(97772875200) = 4032$$

$$- \leq 10^{12} : d(963761198400) = 6720$$

$$- \leq 10^{15} : d(866421317361600) = 26880$$

$$- \leq 10^{18} : d(897612484786617600) = 103680$$

• Numeric integration

- simple:
$$F(0)$$

- simpson: $\frac{F(-1)+4\cdot F(0)+F(1)}{6}$
- runge2: $\frac{F(-\sqrt{\frac{1}{3}})+F(\sqrt{\frac{1}{3}})}{2}$
- runge3: $\frac{F(-\sqrt{\frac{3}{5}})\cdot 5+F(0)\cdot 8+F(\sqrt{\frac{3}{5}})\cdot 5}{18}$

5 Graphs

5.1 Weighted matroid intersection

```
// here we use T = __int128 to store the independent set
// calling expand k times to an empty set finds the maximum
// cost of the set with size exactly k,
// that is independent in blue and red matroids
// ver is the number of the elements in the matroid,
// e[i].w is the cost of the i-th element
```

```
// first return value is new independent set
// second return value is difference between
// new and old costs
// oracle(set, red) and oracle(set, blue) check whether
// or not the set lies in red or blue matroid respectively
auto expand = [\&B] (T cur_set) \rightarrow pair<T, int>
{
    vector<int> in(ver);
    for (int i = 0; i < ver; i++)
        in[i] = ((cur set >> i) & 1);
    const int red = 1;
    const int blue = 2;
    vector<vector<int>>> g(ver);
    for (int i = 0; i < ver; i++)
    for (int j = 0; j < ver; j++)
        T swp_mask = (cur_set ^ (T(1) << i) ^ (T(1) << j));
        if (!in[i] & in[j])
            if (oracle(swp mask, red))
                g[i].push_back(j);
            if (oracle(swp mask, blue))
                g[j].push back(i);
        }
    }
    vector<int> from, to;
    for (int i = 0; i < ver; i++) if (!in[i])</pre>
        T add_mask = cur_set ^(T(1) \ll i);
        if (oracle(add mask, blue))
            from.push back(i);
        if (oracle(add_mask, red))
            to.push back(i);
    }
    auto get_cost = [\delta] (int x)
        const int cost = (!in[x] ? e[x].w : -e[x].w);
        return (ver + 1) * cost - 1;
    };
    const int inf = int(1e9);
    vector<int> dist(ver, -inf), prev(ver, -1);
    for (int x : from)
        dist[x] = get_cost(x);
```

```
queue<int> q;
vector<int> used(ver);
for (int x : from)
    q.push(x);
    used[x] = 1;
}
while (!q.empty())
    int cur = q.front(); used[cur] = 0; q.pop();
    for (int to : g[cur])
        int cost = get_cost(to);
        if (dist[to] < dist[cur] + cost)</pre>
            dist[to] = dist[cur] + cost;
            prev[to] = cur;
            if (!used[to])
            {
                 used[to] = 1;
                 q.push(to);
            }
        }
    }
}
int best = -inf, where = -1;
for (int x : to)
    if (dist[x] > best)
    {
        best = dist[x];
        where = x;
    }
}
if (best = -inf)
    return pair<T, int>(cur_set, best);
while (where \neq -1)
    cur_set ^= (T(1) \ll where);
    where = prev[where];
}
while (best % (ver + 1))
    best++;
```

```
best \neq (ver + 1);
    assert(oracle(cur_set, red) & oracle(cur_set, blue));
    return pair<T, int>(cur set, best);
};
   Push-free segment tree
6
class pushfreesegtree
    vector<modulo♦> pushed, unpushed;
    modulo ⇔ add(int l, int r, int cl, int cr, int v, const modulo
       \leftrightarrow \Leftrightarrow \&x)
    {
        if (r \le cl || cr \le l)
             return 0;
        if (l \le cl \& cr \le r)
             unpushed[v] += x;
            return x * (cr - cl);
        }
        int ct = (cl + cr) / 2;
        auto tmp = add(l, r, cl, ct, 2 * v, x) + add(l, r, ct, cr,
           \rightarrow 2 * v + 1, x);
        pushed[v] += tmp;
        return tmp;
    }
    modulo ⇔ sum(int l, int r, int cl, int cr, int v)
    {
        if (r \le cl || cr \le l)
             return 0;
        if (l \leq cl \& cr \leq r)
             return pushed[v] + unpushed[v] * (cr - cl);
        int ct = (cl + cr) / 2;
        return sum(l, r, cl, ct, 2 * v) + unpushed[v] * (min(r, cr))
           \rightarrow - max(l, cl)) + sum(l, r, ct, cr, 2 * v + 1);
    }
public:
    pushfreesegtree(int n): pushed(2 * up(n)), unpushed(2 * up(n))
```

```
{}
    modulo ⇔ sum(int l, int r)
        return sum(l, r, 0, pushed.size() / 2, 1);
    }
    void add(int l, int r, const modulo ♦ &x)
        add(l, r, 0, pushed.size() / 2, 1, x);
};
7
    Number theory
     Chinese remainder theorem without overflows
// Replace T with an appropriate type!
using T = long long;
// Finds x, y such that ax + by = gcd(a, b).
T gcdext (T a, T b, T &x, T &y)
    if (b = 0)
        x = 1, y = 0;
        return a;
    T res = gcdext(b, a \% b, y, x);
    y -= x * (a / b);
    return res;
}
// Returns true if system x = r1 \pmod{m1}, x = r2 \pmod{m2} has
   → solutions
// false otherwise. In first case we know exactly that x = r \pmod{m}
   \hookrightarrow )
bool crt (T r1, T m1, T r2, T m2, T &r, T &m)
    if (m2 > m1)
    {
        swap(r1, r2);
        swap(m1, m2);
    }
    T g = \underline{gcd(m1, m2)};
```

```
if ((r2 - r1) \% g \neq 0)
        return false;
    T c1, c2;
    auto nrem = gcdext(m1 / g, m2 / g, c1, c2);
    assert(nrem = 1);
    assert(c1 * (m1 / g) + c2 * (m2 / g) = 1);
    T a = c1;
    a *= (r2 - r1) / g;
    a \% = (m2 / g);
    m = m1 / g * m2;
    r = a * m1 + r1;
    r = r \% m;
    if (r < 0)
        r += m;
    assert(r % m1 = r1 & r % m2 = r2);
    return true;
}
7.2 Integer points under a rational line
// integer (x, y) : 0 \le x < n, 0 < y \le (kx + b) / d
// (real division)
// In other words, sum \{x=0\}^{n-1} [(kx+b)/d]
ll trapezoid (ll n, ll k, ll b, ll d)
{
    if (k = 0)
        return (b / d) * n;
    if (k \ge d || b \ge d)
        return (k / d) * n * (n - 1) / 2 + (b / d) * n + trapezoid(
           \rightarrow n, k % d, b % d, d);
    return trapezoid((k * n + b) / d, d, (k * n + b) % d, k);
}
```