



Formal Method in Software Engineering (SE-313)

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Introduction to Petri Net

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Petri Net

Petri Net Applications

- **Software: design, specification, simulation, validation, and implementation**
- **Manufacturing, production, and scheduling systems**
- **Sequence controllers (Programmable Logic Controller, PLC)**
- **Communication protocols and networks**

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Petri Net

Petri Nets

- **A Petri net is an abstract, formal model(graphical and mathematical) of information flow.**
- **Petri nets make it possible to model and visualize behaviors with parallelism, concurrency, synchronization and resource sharing.**
- **Powerful methods for describing and analyzing the flow of information and control in systems**
- **Petri nets represent computer systems by**
 - **providing a means to abstract the basic elements of the system**
 - **and its informational flow using only four fundamental components.**

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Petri Net

A Petri net is a bipartite directed graph composed of places and transitions. Petri nets are usually represented graphically according to the following conventions:-

- Places are represented by circles,(resources, conditions, buffers, locations)
- transitions by bars,(events, actions)
- input function by arcs directed from places to transitions,
- output function by arcs directed from transitions to places, and
- Tokens are dynamic objects used to track information and appear as solid dots within the circle of a place.
- The state of a Petri net, called the marking, is determined by the distribution of tokens over the places.

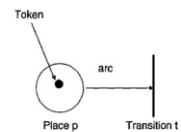
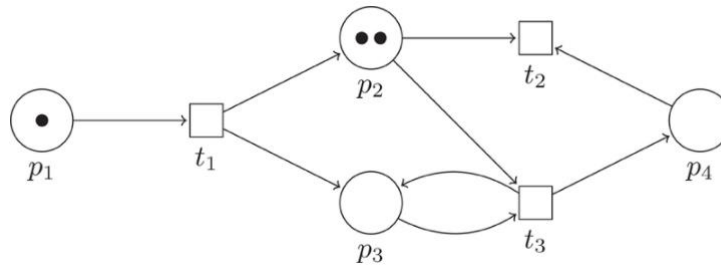


Fig 1: Basic Petri net components

Typical Interpretations of Transitions and Places

| Input Places | Transitions | Output Places |
|--------------------|------------------|--------------------|
| Precondition | Event | Postcondition |
| Input Data | Computation Step | Output Data |
| Input Signal | Signal Processor | Output Signals |
| Required Resources | Task or Job | Released Resources |
| Conditions | Logic Clause | Conclusion |
| Buffers | Process | Buffers |

- Example:



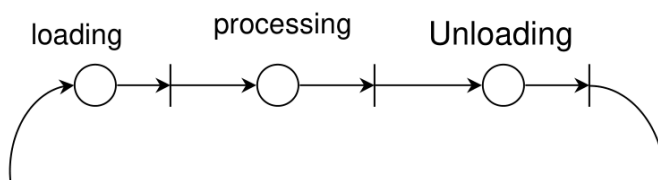
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- Enabling Rule:

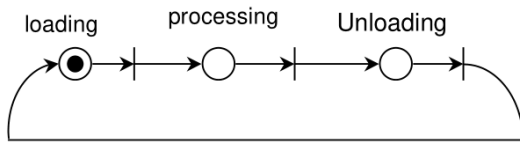
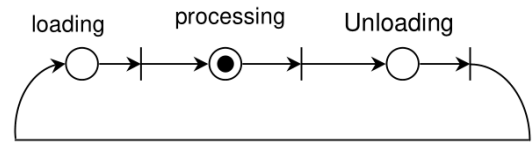
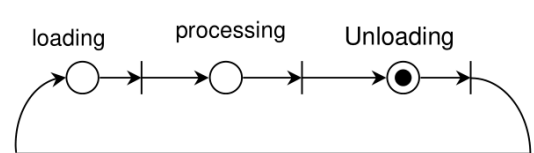
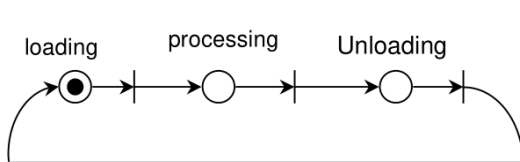
- » A transition t is enabled if every input place contains at least one token

- Firing Rule:

- » Firing an enabled transition
 - removes one token from each input place of the transition
 - adds one token to each output place of the transition



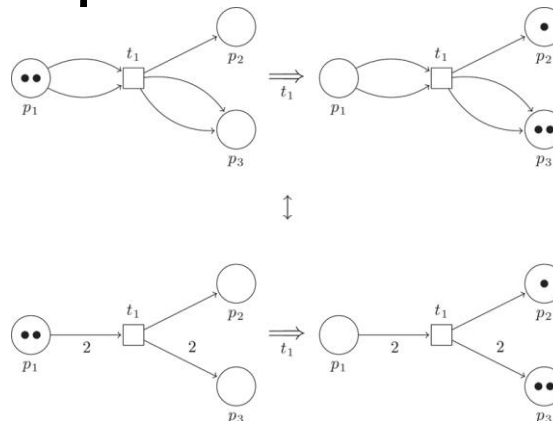
Initial State:


 t_1
State after t_1 is fired:
 t_2
State after t_2 is fired:State after t_3 is fired:
 t_3

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Multiplicity of Arcs:

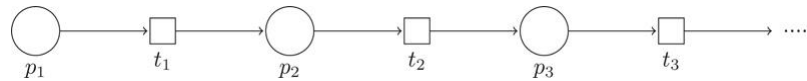
- Firing a transition may require consuming multiple tokens from an input place and may generate multiple tokens to an output place.



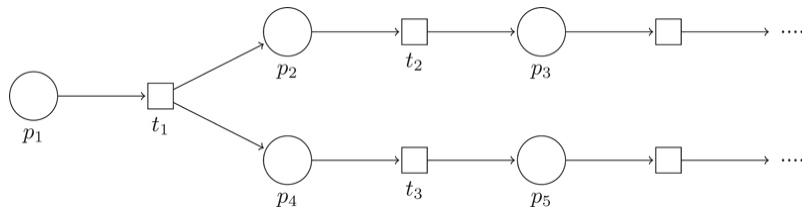
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Common Petri Net Structures

- Sequential Execution



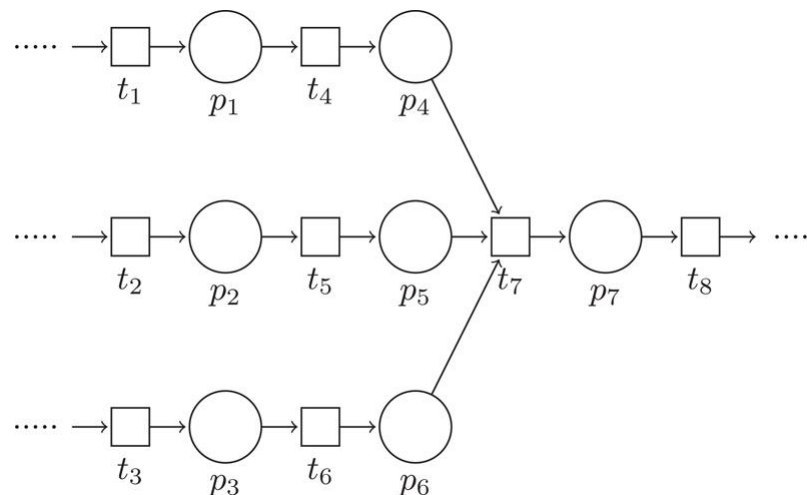
- Concurrent Execution



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Common Petri Net Structures

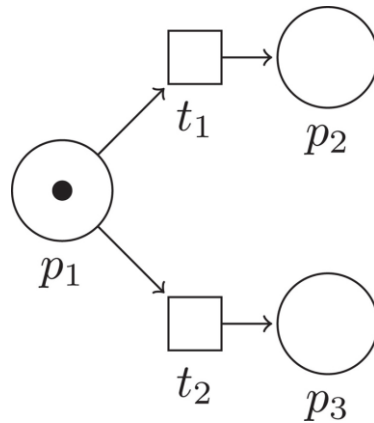
- Synchronization



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Common Petri Net Structures

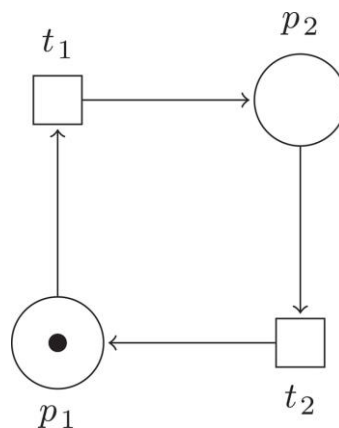
- Nondeterminism



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Common Petri Net Structures

- Loop

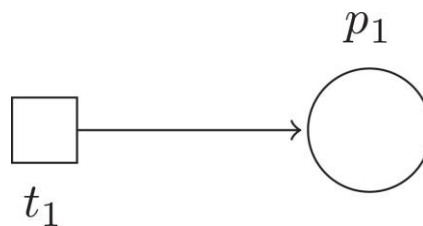


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Common Petri Net Structures

Source

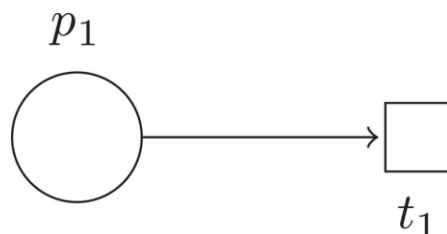
- A source is a transition that has no input place but does have an output place.
- This can be used to model an infinite source of resources entering into the system.



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Common Petri Net Structures

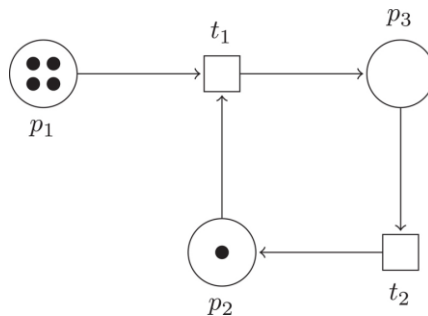
- Consumer
- The opposite of a source is a consumer. A consumer is a transition that has an input place with no output places.
- This consumes the tokens in the input place for the transition.



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Common Petri Net Structures

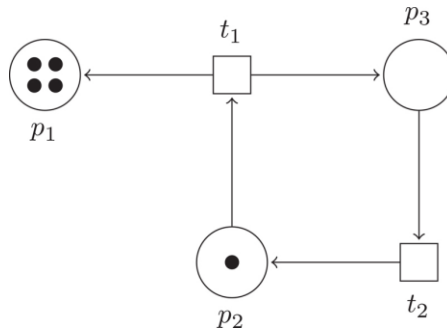
- **Control**
- Modeling some kind of fixed resource can be accomplished rather easily. In this case, a control place is initialized with the available resources.
- That place can be incorporated into the Petri net as a limiter.
- It can also be used as a break point for the loop.



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Common Petri Net Structures

- **Accumulator**
- The opposite of a control is an accumulator. Rather than a place that holds tokens which are to be consumed, an accumulator is a place where tokens are created and never consumed.
- It effectively holds a count of how many times t_1 has been fired.



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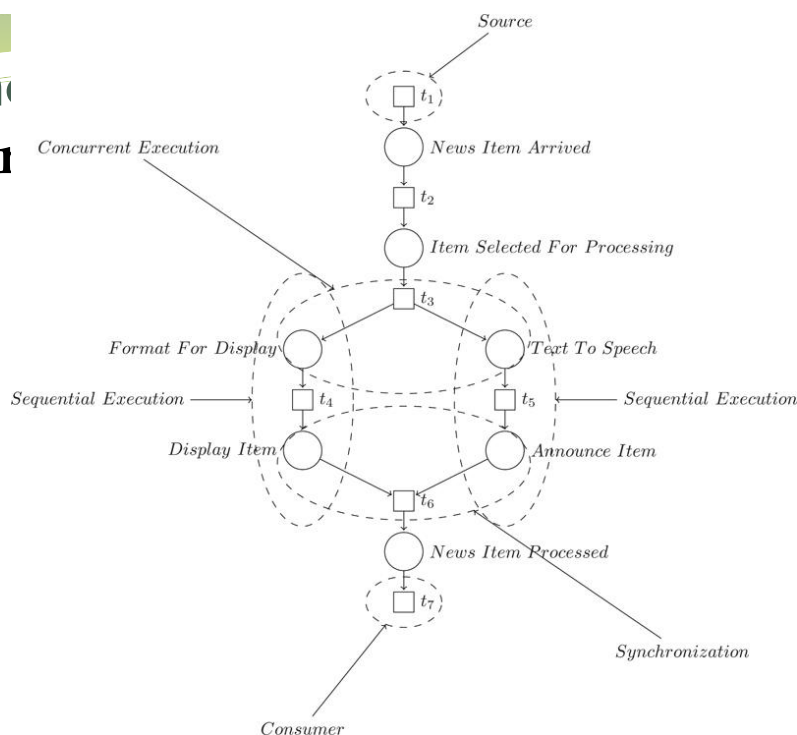
• Example

- A Petri net representing a general announcement program that displays news items on a monitor and speaks the news item aloud over a loudspeaker. It is easy to analyze this program by recognizing the various structures and substructures within it.

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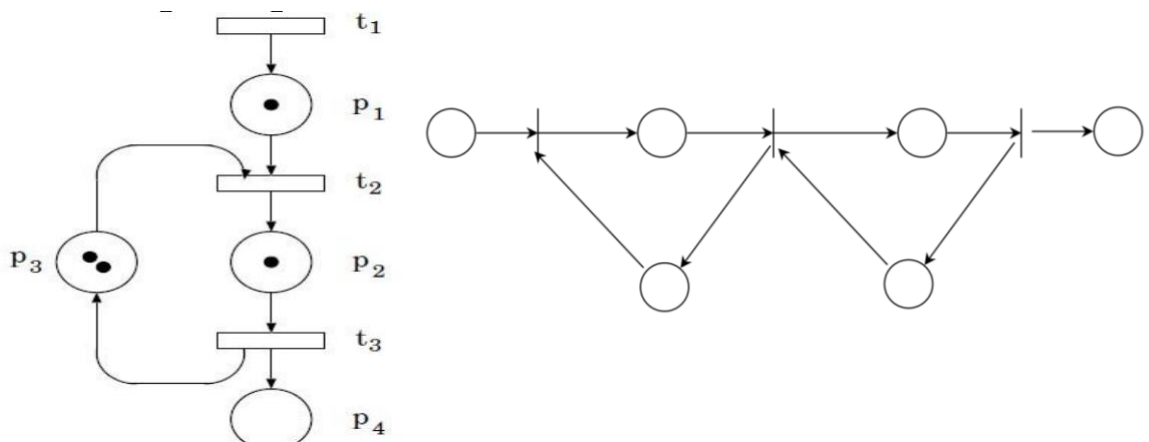
Comm

• Synch



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Buffer (Queue)



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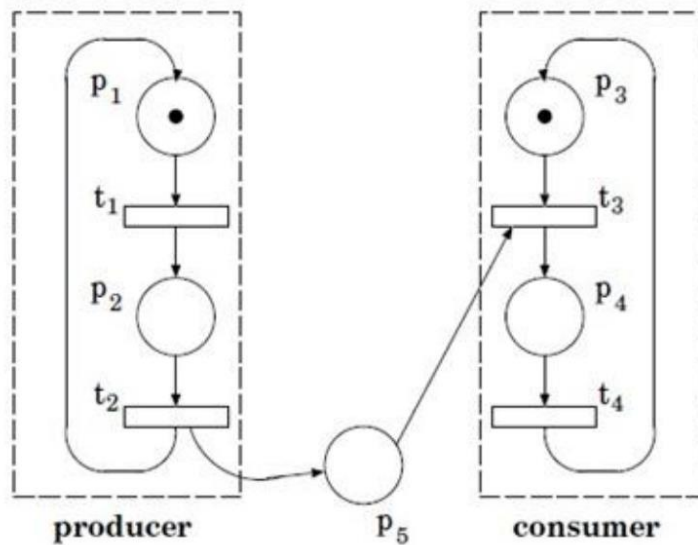
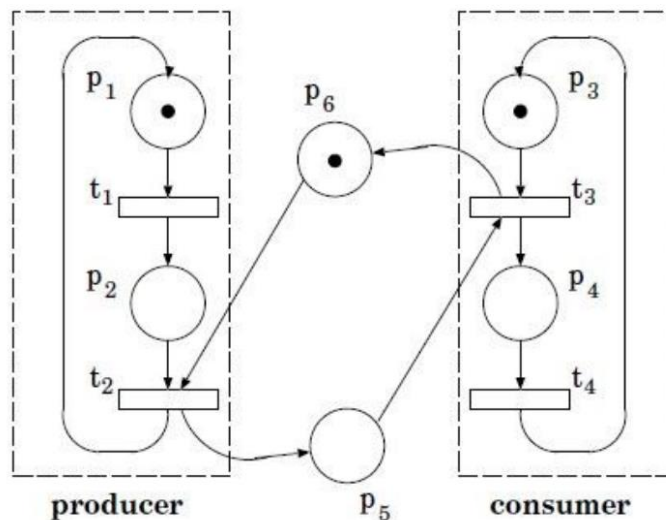


Fig 15: The producer/consumer problem with unbounded buffer

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: The producer/consumer problem with finite buffer

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Resource Sharing

Mutual Exclusion (Conflict)

- Places p_1 and p_5 represent C_1 and C_2 working independently;
- p_2 and p_6 represent C_1 and C_2 requesting access to C_s ;
- Place p_4 determines availability of resource C_s ; prevents p_3 & p_7 to be marked at the same time;
- p_3 and p_7 represent C_s busy with C_1 and C_2 respectively.
- Firing of t_3 or t_6 models the release of the common resource.

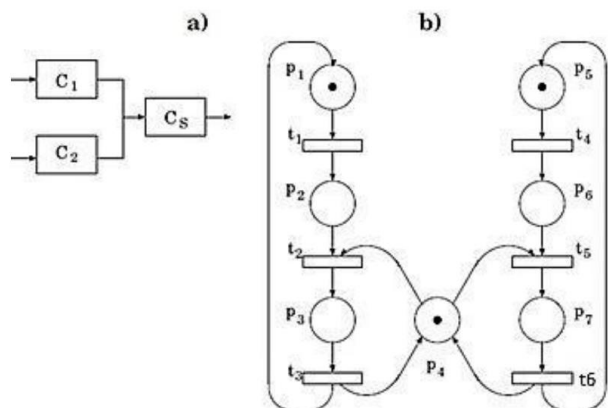
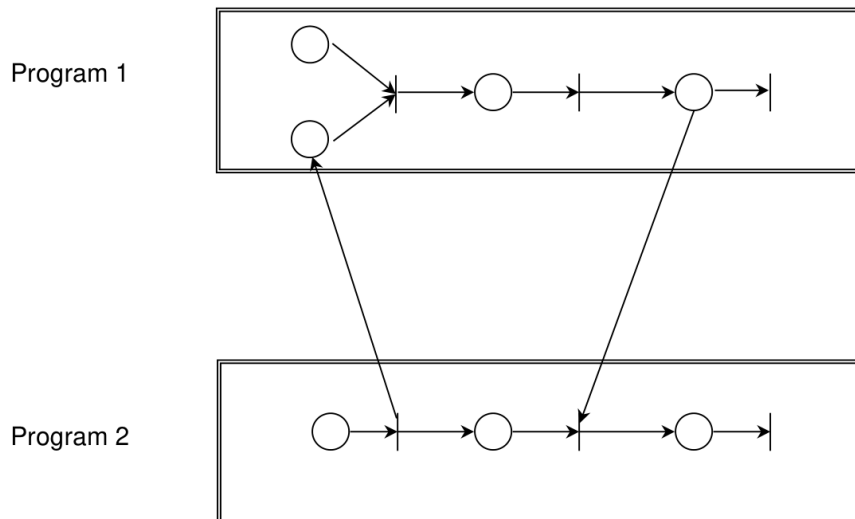


Fig 17: The mutual exclusion problem

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Communication



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Petri Net

Example

Consider for example the following description of a computer system:

- Jobs appear and are put on an input list. When the processor is free, and there is a job on the input list, the processor starts to process the job.
- When the job is complete, it is placed on an output list, and if there are more jobs on the input list, the processor continues with another job; otherwise it waits for another job.

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Petri Net

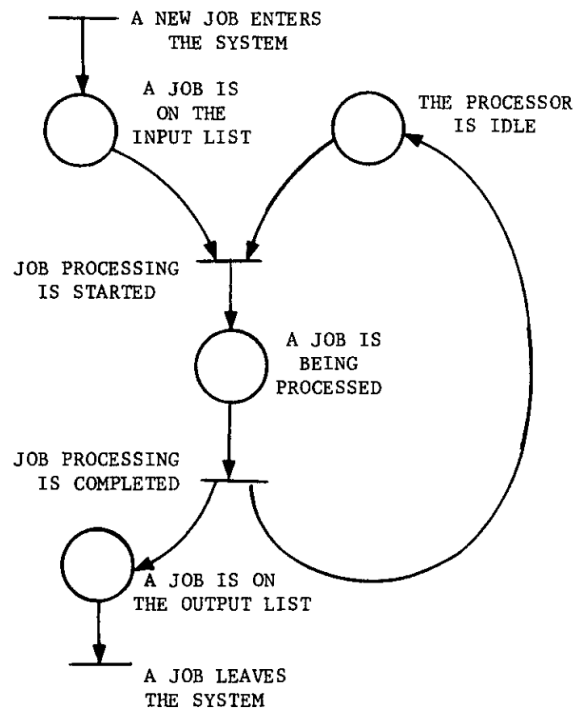
Petri Net Example

- **conditions:**

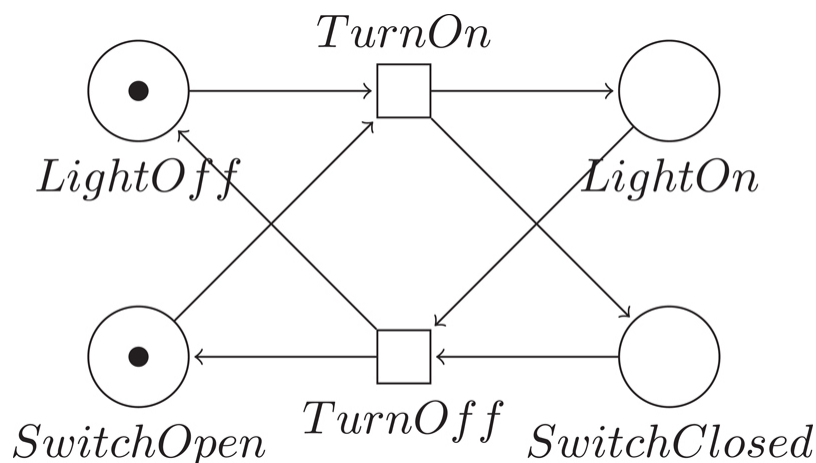
- The processor is idle;
- A job is on the input list;
- A job is being processed;
- A job is on the output list;

- **Events:**

- A new job enters the system;
- Job processing is started;
- Job processing is completed;
- A job leaves the system.

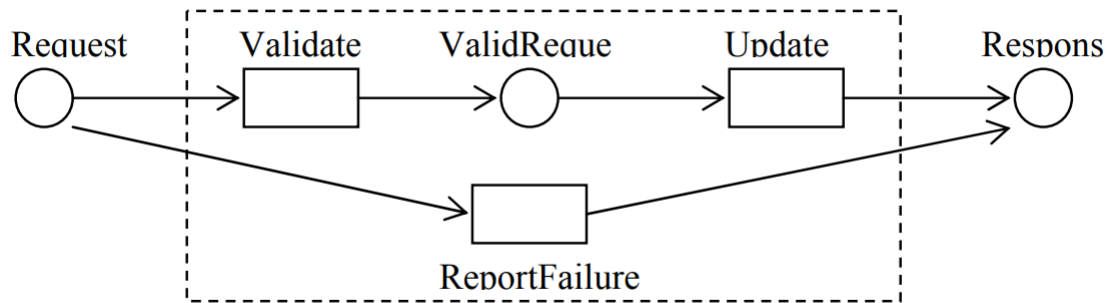


Example Model Light bulb and switch circuit using petri Net.



Petri Net

Petri Net Example



Model of Transactions

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Petri Net

Mathematical Description of Petri Nets

A graphical depiction of a Petri net allows one to understand it structurally and intuitively.

However, a picture doesn't help answer a number of questions that are of interest when specifying a system, including such questions as

- Can the system reach a specific state from the current state?
- What states of the system are accessible?
- Is the system alive?

To answer these questions, one must adopt a mathematical model of a Petri net.

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Petri Net

Mathematical Description of Petri Nets

A Petri net N is a tuple $N = \{P, T, I, O, M_0\}$, where

- P is a finite set of places, graphically represented by circles.
- T is a finite set of transitions, graphically represented by boxes.
- $I: P \times T \rightarrow (\{0, 1, 2, \dots\})$ is the pre-incidence function representing input arcs.
- $O: T \times P \rightarrow (\{0, 1, 2, \dots\})$ is the post-incidence function representing output arcs.
- $M_0: P \rightarrow N$ is the initial marking representing the initial distribution of tokens.

Places P and transitions T are disjoint ($P \cap T = \emptyset$).

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Petri Net

Mathematical Description of Petri Nets

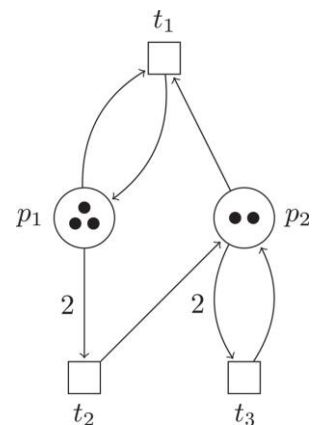
Example

$$P = \{p_1, p_2\}$$

$$T = \{t_1, t_2, t_3\}$$

$$I(p_1, t_1) = 1 \quad I(p_1, t_2) = 2 \quad I(p_1, t_3) = 0$$

$$I(p_2, t_1) = 1 \quad I(p_2, t_2) = 0 \quad I(p_2, t_3) = 2$$



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Petri Net

Mathematical Description of Petri Nets

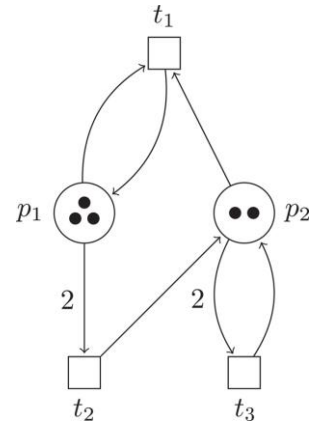
Example

$$O(t_1, p_1) = 1 \quad O(t_1, p_2) = 0$$

$$O(t_2, p_1) = 0 \quad O(t_2, p_2) = 1$$

$$O(t_3, p_1) = 0 \quad O(t_3, p_2) = 1$$

$$M_0 = (3, 2)$$



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Petri Net

Extensions of Petri Nets

- **Deterministic Timed Petri Nets**
 - Deterministic time delays with transitions
- **Stochastic Timed Petri Nets**
 - Stochastic time delays with transitions
- **Color Petri Nets**
 - Tokens with different colors
- **Hybrid Nets**
 - Combine object-oriented concept into Petri nets

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