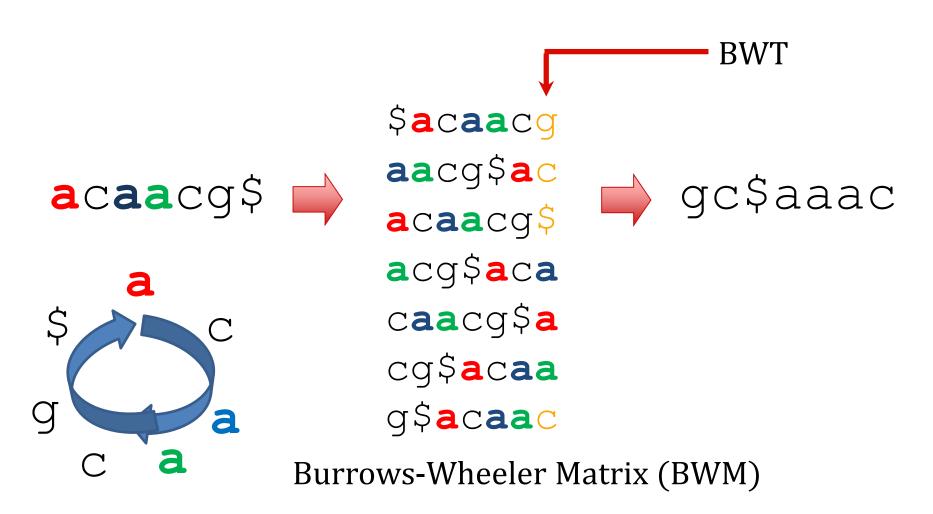
Alignment of Short Reads: Burrows-Wheeler Transform (BWT)

Ahmet Sacan

Burrows-Wheeler Transform (BWT)



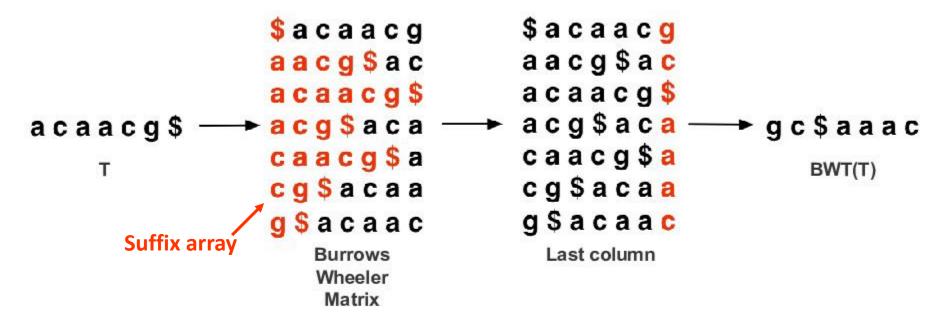
Applying BWT to the Double Helix Paper by Watson&Crick

```
nd Corey (1). They kindly made their manuscript availa ..... a
nd criticism, especially on interatomic distances. We ..... a
nd cytosine. The sequence of bases on a single chain d ..... a
nd experimentally (3,4) that the ratio of the amounts o ..... u
nd for this reason we shall not comment on it. We wish ..... a
nd guanine (purine) with cytosine (pyrimidine). In oth ..... a
nd ideas of Dr. M. H. F. Wilkins, Dr. R. E. Franklin ..... a
nd its water content is rather high. At lower water co ..... a
nd pyrimidine bases. The planes of the bases are perpe ..... a
nd stereochemical arguments. It has not escaped our no ..... a
nd that only specific pairs of bases can bond together ..... u
nd the atoms near it is close to Furberg's 'standard co ..... a
nd the bases on the inside, linked together by hydrogen ..... a
nd the bases on the outside. In our opinion, this stru ..... a
nd the other a pyrimidine for bonding to occur. The hy ..... a
nd the phosphates on the outside. The configuration of ..... a
nd the ration of guanine to cytosine, are always very c ..... a
nd the same axis (see diagram). We have made the usual ..... u
nd their co-workers at King's College, London. One of ..... a
```

"and" is a frequent repeat in English texts

The Burrows-Wheeler Transform is a reversible representation with handy properties

Sort all the possible rotations of original string



- Once BWT(T) is built, all else shown here is discarded
 - Matrix will be shown for illustration only

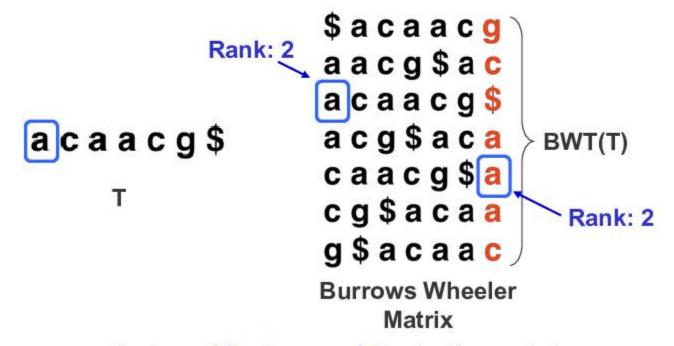
Burrows M, Wheeler DJ: A block sorting lossless data compression algorithm. Digital Equipment Corporation, Palo Alto, CA 1994, Technical Report 124; 1994

Recovering the original string

• If the BWT(s) is: "gc\$aaac", what is the original string s?

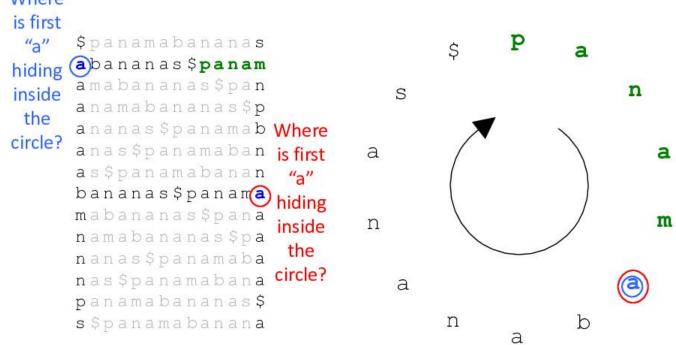
A text occurrence has the same rank in the first and last columns

 When we rotate left and sort, the first character retains its rank. Thus the same text occurrence of a character has the same rank in the Last and First columns.

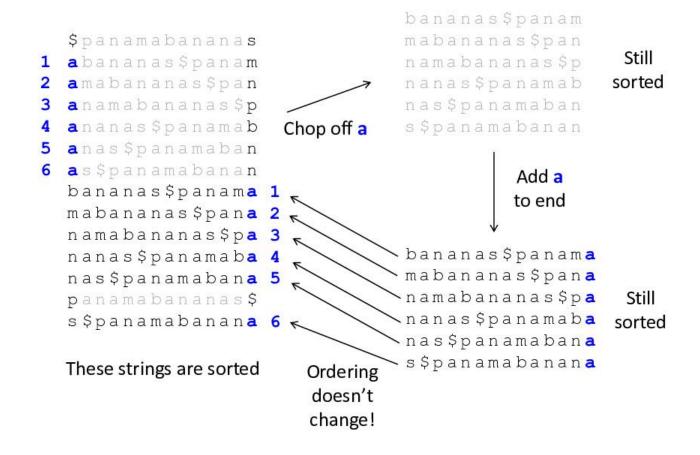


Courtesy of Ben Langmead. Used with permission.

A Strange Observation



Is It True in General?



The Last to First (LF) function matches character and rank

- Occ(qc): First occurrence of the character type qc.
- Count(ind, qc): How many characters of type qc are before ind?

Walk Left to Invert BWT (recover original)

```
i = 0
     while bwt[i] != `$':
       t = bwt[i] + t
       i = LF(i, bwt[i])
                                                        Final t
                                    aacg
                                             caacq
                                                       acaacq
$acaacg
          $acaacg
                     $acaacg
                                $acaacg
                                           $acaacg
                                                     $acaacg
                     aacg$ac
                                aacg$ac
aacq$ac
          aacg$/c
                                                     aacq Sac
acaacg$
          acaa / g $
                    acaacg$
                                acaacg$
acg$aca
          a c g 🌠 a c a
                     acg$aca
                                           acg$aca
caacg$a
          caacq$a
                     caacg$a
                                caacq Sa
                                           caacg$a
cg$acaa
                                cg$acaa
                                           cg$acaa
                                                     cg$acaa
g $ a c a a c
                     q $ a c a a c
                                q$acaac
                                           a Sacaac
                                                     q$acaac
```

Find ana in panamabananas

Lets Start by Matching the Last Symbol (a)

```
$1panamabananas1
a1bananas$panam1
a2mabananas$pan1
a3namabananas$p1
a4nanas$panamab1
a5nas$panamaban2
a6s$panamabanan3
b1ananas$panama1
m1abananas$pana2
n1amabananas$pana2
n1amabananas$pana5
p1anamabananas$1
s1$panamabanana6
```

Matching the Last Two Symbols (na)

```
$1panamabananas1
a1bananas$panam1
a2mabananas$pan1
a3namabananas$p1
a4nanas$panamab1
a5nas$panamaban2
a6s$panamabanan3
b1ananas$panama1
m1abananas$pana2
n1amabananas$pana2
n1amabananas$pana5
p1anamabananas$1
s1$panamabanana6
```

Three Matches of **na** Found!

```
$1panamabananas1
a1bananas$panam1
a2mabananas$pana
a3namabananas$p1
a4nanas$panamab1
a5nas$panamaban2
a6s$panamabananas$pana2
n1amabananas$pana2
n1amabananas$pana2
n1amabananas$pana2
n1amabananas$pana2
n1amabananas$pana2
```

Three Matches of **na** Found!

```
$1panamabananas1
a1bananas$panam1
a2mabananas$pan1
a3namabananas$p1
a4nanas$panamab1
a5nas$panamaban2
a6s$panamabanan3
b1ananas$panama1
m1abananas$pana2
n1amabananas$pana2
n1amabananas$pana5
p1anamabananas$1
s1$panamabanana6
```

Matching **ana**

```
$1panamabananas1
a1bananas$panam1
a2mabananas$pan1
a3namabananas$p1
a4nanas$panamab1
a5nas$panamaban2
a6s$panamabanan3
b1ananas$panama1
m1abananas$pana2
n1amabananas$pana2
n1amabananas$pana2
n1amabananas$pana6
p1anamabananas$1
s1$panamabanana6
```

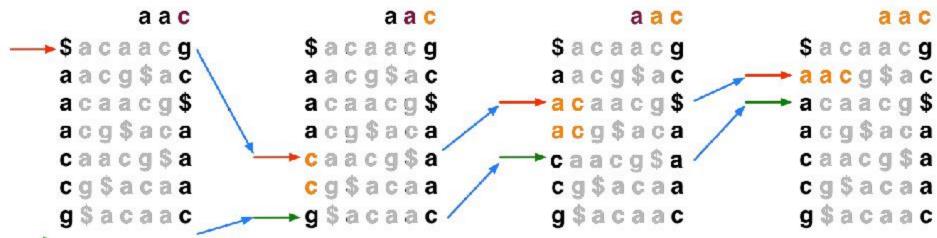
Three Matches of ana Found!

```
$1panamabananas1
a1bananas$panam1
a2mabananas$pan1
a3namabananas$p1
a4nanas$panamab1
a5nas$panamaban2
a6s$panamabanan3
b1ananas$panama1
m1abananas$pana2
n1amabananas$pana2
n1amabananas$pana6
p1anamabananas$1
s1$panamabanana6
```

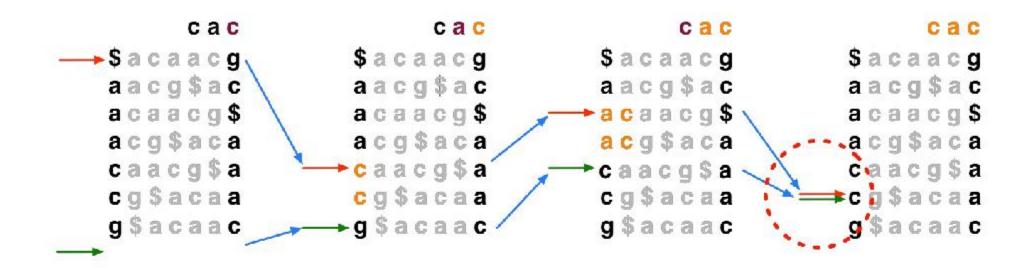
Exact Matching with FM Index

```
q = "aac"
top = 0
bot = len(bwt)
for qc in reverse(q):
  top = LF(top, qc)
  bot = LF(bot, qc)
```

In each iteration top & bot delimit the range of rows beginning with progressively longer suffixes of q



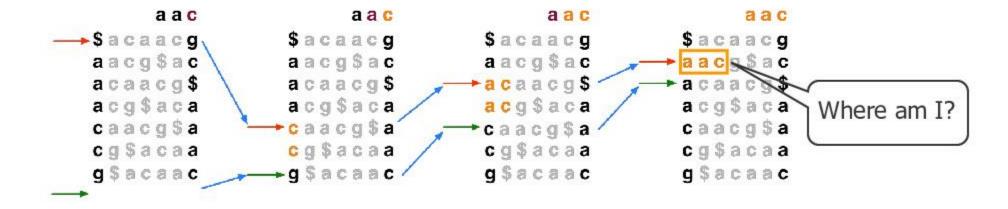
Exact Matching with FM Index



 If range becomes empty (top = bot) the query suffix (and therefore the query) does not occur in the text

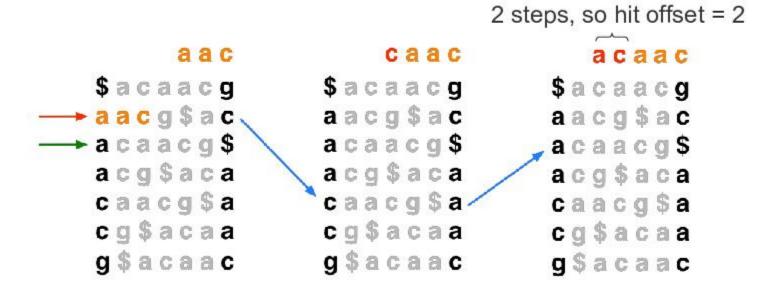
Courtesy of Ben Langmead. Used with permission.

 Once we know a row contains a legal alignment, how do we determine its position in the reference?



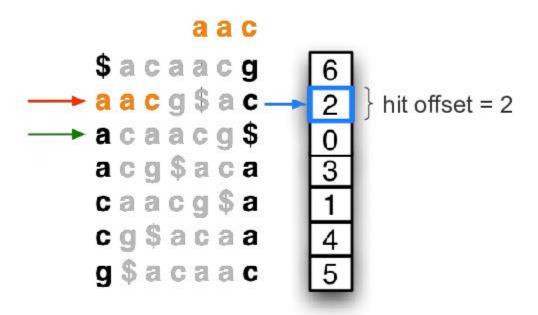
Courtesy of Ben Langmead. Used with permission.

 Naïve solution 1: Use "walk-left" to walk back to the beginning of the text; number of steps = offset of hit



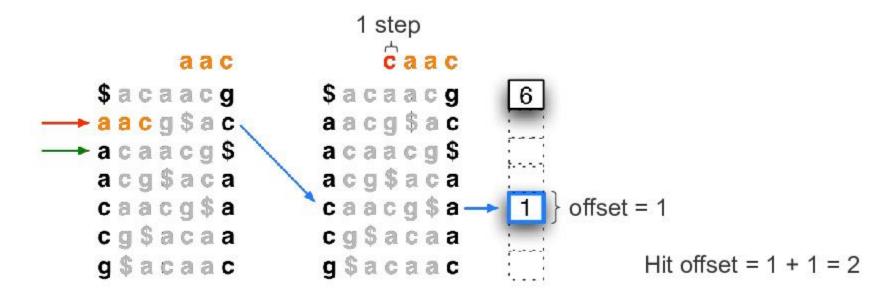
Linear in length of text in general – too slow

Naïve solution 2: Keep whole suffix array in memory.
 Finding reference position is a lookup in the array.



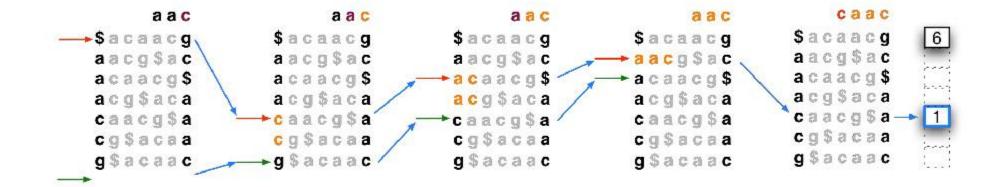
Suffix array is ~12 gigabytes for human – too big

- Hybrid solution: Store sample of suffix array; "walk left" to next sampled ("marked") row to the left
 - Due to Ferragina and Manzini



Bowtie marks every 32nd row by default (configurable)

Put It All Together

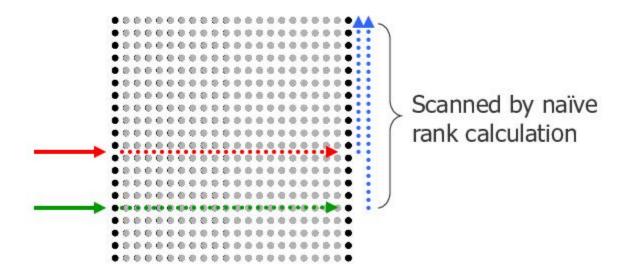


 Algorithm concludes: "aac" occurs at offset 2 in "acaacg"

Courtesy of Ben Langmead. Used with permission.

The FM index makes LF fast

- LF(i, qc) must determine the rank of qc in row i
- Naïve way: count occurrences of qc in all previous rows
 - This LF(i, qc) is linear in length of text too slow



A Full-text Minute-size (FM) index makes LF constant time

 Solution: pre-calculate cumulative counts for A/C/G/T up to periodic checkpoints in BWT [A: 240, C: 225, G: 226, T: 267] Rank: 242 Rank: 309 [A: 310, C: 299, G: 308, T: 332] LF(i, qc) is now constant-time (if space between checkpoints is considered constant)

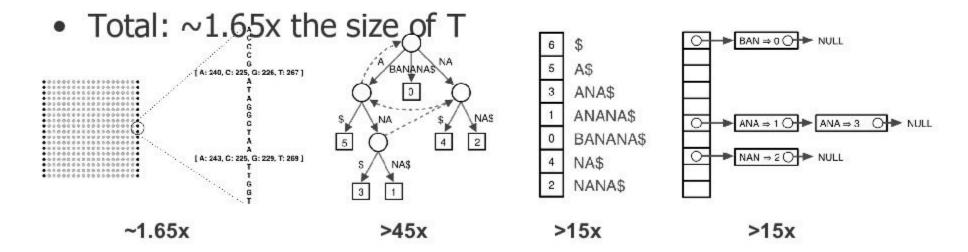
An FM Index is Small

- Entire FM Index on DNA reference consists of:
 - BWT (same size as T)
 - Checkpoints (~15% size of T)
 - Suffix array sample
 (~50% size of T)

Assuming 2-bit-per-base encoding and no compression, as in Bowtie

Assuming a 16-byte checkpoint every 448 characters, as in Bowtie

Assuming Bowtie defaults for suffixarray sampling rate, etc



How about Approximate Matching?

• searching for ana in panamabananas

```
$1panamabananas1
a1bananas$panam1
a2mabananas$pan1
a3namabananas$p1
a4nanas$panamab1
a5nas$panamabanan2
a6s$panamabanan3
b1ananas$panama1
m1abananas$pana2
n1amabananas$pana2
n1amabananas$pana5
p1anamabananas$1
s1$panamabanana6
```

Exact matching

BWT Pattern Matching with 1 Mismatch

• searching for ana in panamabananas

To allow for 1 mismatch, we need to analyze the rows ending in red letters as well.

```
$1panamabananas1

a1bananas$panam1
a2mabananas$pan1
a3namabananas$p1
a4nanas$panamab1
a5nas$panamabanan2
a6s$panamabanan3
b1ananas$panama1
m1abananas$pana2
n1amabananas$pana2
n1amabananas$pana2
n1amabananas$pana6
p1anamabananas$1
s1$panamabanana6
```

Approximate matching with at most 1 mismatch

BWT Pattern Matching with 1 Mismatch

• searching for ana in panamabananas

To allow for 1 mismatch, we need to analyze the rows ending in red letters as well.

```
# Mismatches
$ 1 panamabananas 1
a 1 bananas $ panam 1
a 2 mabananas $ pan 1
a 3 namabananas $ pan 1
a 4 nanas $ panamab 1
a 5 nas $ panamaban 2
a 6 s $ panamabanan 3
b 1 ananas $ panama 1
m 1 abananas $ panama 1
m 1 abananas $ panama 2
n 1 amabananas $ pana 3
n 2 anas $ panamabana 4
n 3 as $ panamabana 5
p 1 anamabananas $ 1
s 1 $ panamabanana 6
```

BWT Pattern Matching with 1 Mismatch

• searching for ana in panamabananas

Now we analyze all rows with at most 1 mismatch using the First-Last property.

This row results in a 2nd mismatch (the \$), so we discard it.

```
#Mismatches
$\(^1\)panamabananas\(^1\)
a\(^1\)bananas\(^2\)panama\(^1\)
a\(^2\)mabananas\(^2\)panamabananas\(^2\)
a\(^3\)panamabanamaban\(^2\)
a\(^3\)panamabanana\(^3\)
b\(^1\)ananas\(^3\)panamaban\(^2\)
a\(^3\)panamabananas\(^3\)panamabana\(^3\)
\(^1\)anamabananas\(^3\)panamabanana\(^3\)
\(^1\)mamabanamabanamaban\(^3\)
\(^1\)mamabanamabanamaban\(^3\)
\(^1\)mamabanamabanamabana\(^3\)
\(^1\)mamabanamabanamabana\(^3\)
\(^1\)mamabanamabanamabana\(^3\)
\(^1\)mamabanamabanamabana\(^3\)
\(^1\)mamabanamabanamabanana\(^3\)
```

Five Approximate Matches Found!

```
# Mismatches
$,panamabananas,
a<sub>1</sub>bananas Spanam<sub>1</sub>
a<sub>2</sub>mabananas$pan<sub>1</sub>
a<sub>3</sub>namabananas$p<sub>1</sub>
a<sub>4</sub>nanas$panamab<sub>1</sub>
                                          0
a<sub>5</sub>nas$panamaban<sub>2</sub>
                                          0
a<sub>6</sub>s$panamabanan<sub>3</sub>
b<sub>1</sub>ananas$panama<sub>1</sub>
m<sub>1</sub>abananas $pana<sub>2</sub>
n_1amabananaspa_3
n<sub>2</sub>anas$panamaba<sub>4</sub>
n<sub>3</sub>as$panamabana<sub>5</sub>
p<sub>1</sub>anamabananas $<sub>1</sub>
s<sub>1</sub>$panamabanana<sub>6</sub>
```