

# Alignment of Short Reads: Suffix Trees

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Slides adapted from:

<https://www.coursera.org/learn/algorithms-on-strings>

<https://ocw.mit.edu/courses/biology/7-91j-foundations-of-computational-and-systems-biology-spring-2014/video-lectures/lecture-5-library-complexity-and-short-read-alignment-mapping/>

[http://www.cs.utsa.edu/~jrjuan/teaching/cs5263\\_spring\\_2015/slides/slide5\\_shortreadmapping.ppt](http://www.cs.utsa.edu/~jrjuan/teaching/cs5263_spring_2015/slides/slide5_shortreadmapping.ppt)

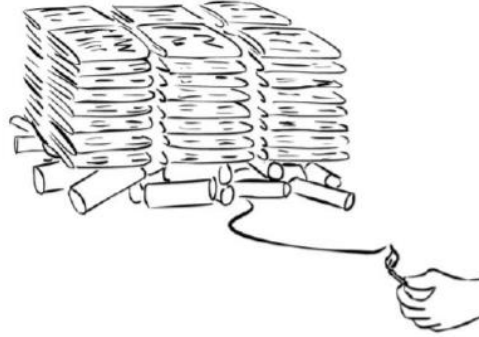
# The Newspaper Problem



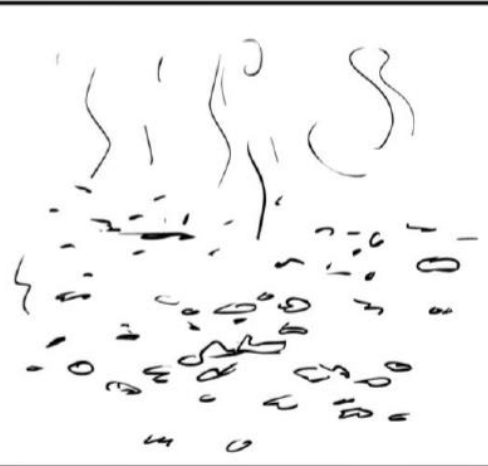
stack of NY Times, June 27, 2000



stack of NY Times, June 27, 2000  
on a pile of dynamite

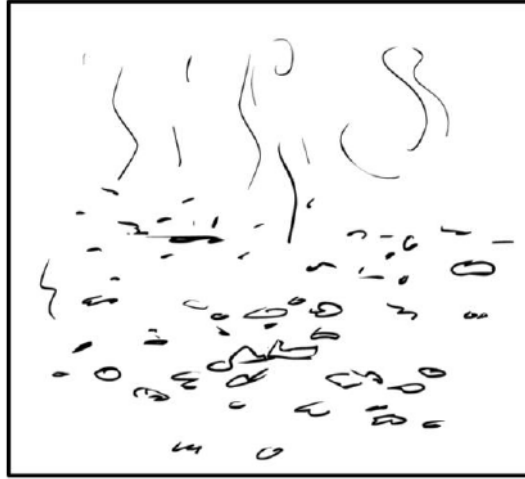


this is just hypothetical



so, what did the June 27, 2000 NY  
Times say?

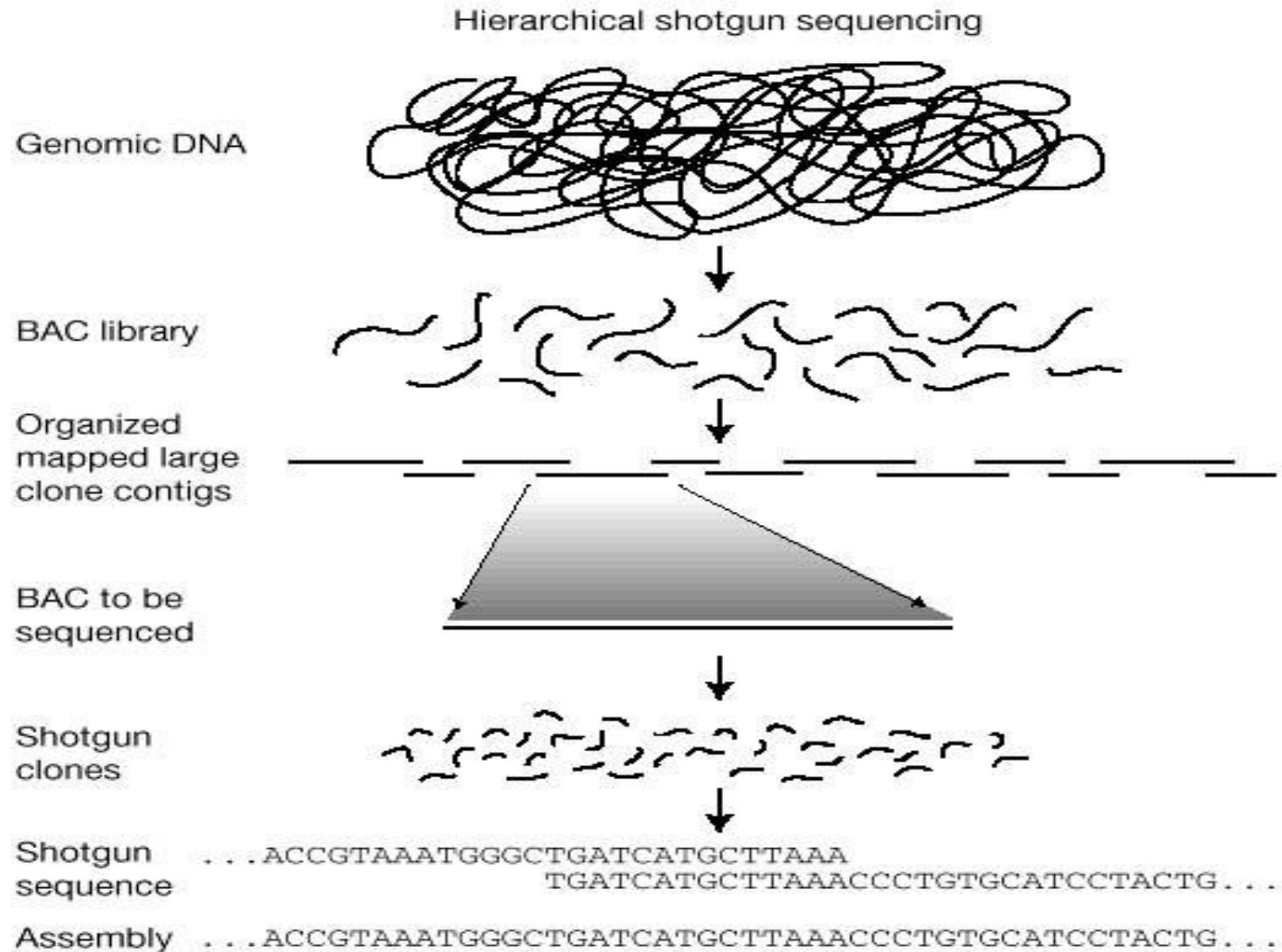
# The Newspaper Problem as an Overlapping Puzzle

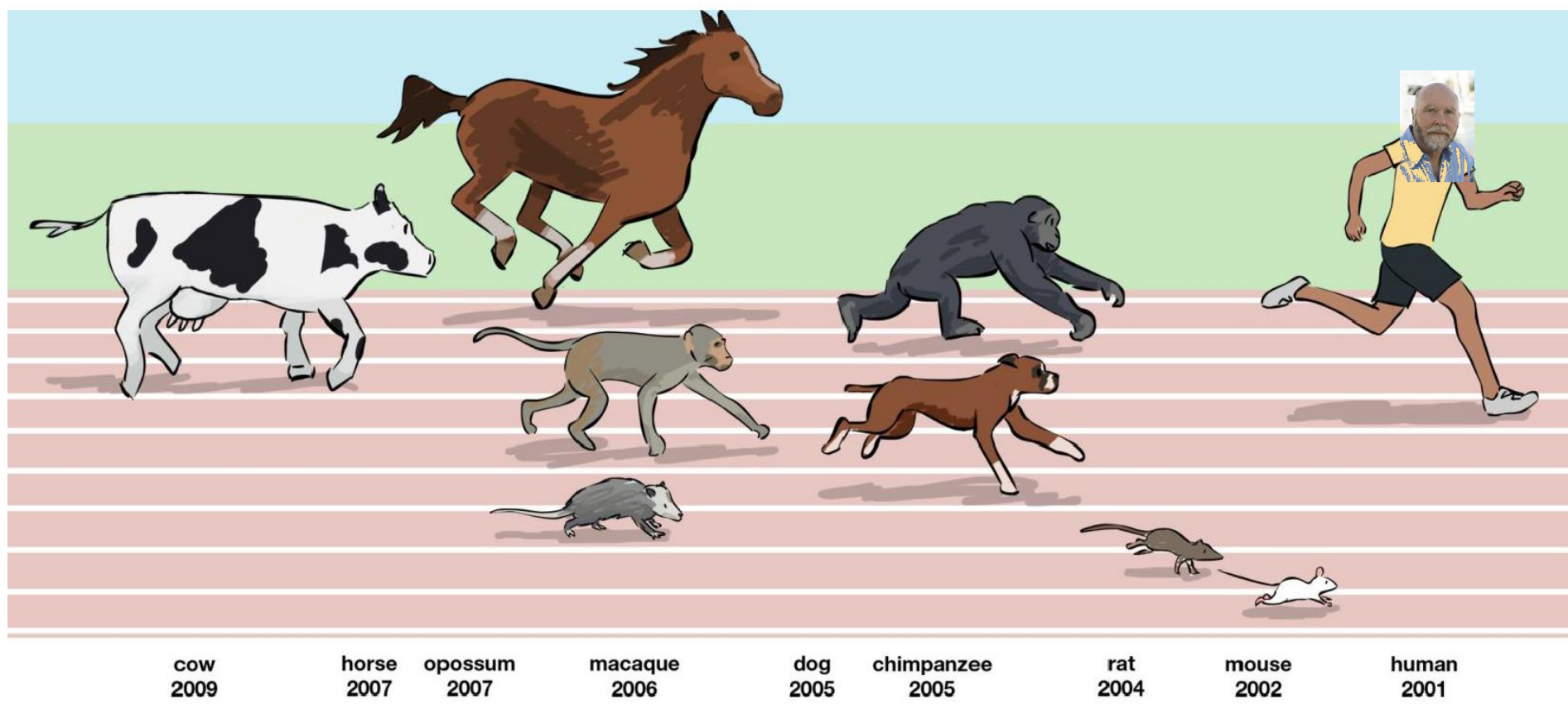


...noodie, appri  
...e have not yet named  
...mation is welc

...lie, appri  
...yet named any suspects, alt  
...is welc  
...2'  
...e ca

# Whole-Genome Sequencing





# From Reference Genome to Personal Genomes

- **Reference human genome** assembled in 2000.

CTGAGGATGGACTACGCTACTACTGATAGCTGTTT      reference  
genome

GAGGA            CACG            TGA-AG  
CTGA            GGAC            ACTAC    A-AGCT      reads  
          GATGG    ACGCT                    TGTTT



CTGAGGATGGACACGCTACTACTGA-AGCTGTTT      personal  
genome

# Reference-based sequencing

- Map short reads to the reference genome.
- Exact "pattern matching"
  - Locate the occurrence of short read exactly.
- Approximate "pattern matching"
  - Allow mismatches, insertions, deletions.

# Searching short read (pattern) in genome (text)

*Pattern* drives along *Text*

p a n a m a b a n a n a s

*Text*

n a n a

*Pattern*





# Searching short read (pattern) in genome (text)

*Pattern* drives along *Text*

p a n a m a b a n a n a s

*Text*

n a n a

*Pattern*



# Searching for multiple patterns

Genome

CTGATGATGGACTACGCTACTACTGCTAGCTGTATTACGATCAGCTAC**C**ACATCGTAGCTAC



# Searching for multiple patterns

Genome

CTGATGATGGACTACGCTACTACTGCTAGCTGTATTACGATCAGCTAC**C**ACATCGTAGCTAC



# Computational Complexity

- How long does it take to find a single pattern?  
 $O(|\text{text}| * |\text{pattern}|)$
- Multiple patterns?  
 $O(|\text{pattern1}| * |\text{text}|) + O(|\text{pattern2}| * |\text{text}|) + \dots$   
 $= O(|\text{text}| * |\text{patterns}|)$
- Human genome:
  - $|\text{Text}| \approx 10^9$
  - $|\text{Patterns}| \approx 10^{12}$

# Pack Patterns onto a bus

CTGATGATGGACTACGCTACTACTGCTAGCTGTATTACGATCAGCTAC**C**ACATCGTAGCTAC





## ***Patterns***

banana

pan

nab

antenna

bandana

ananas

nana



## ***Patterns***

**banana**

pan

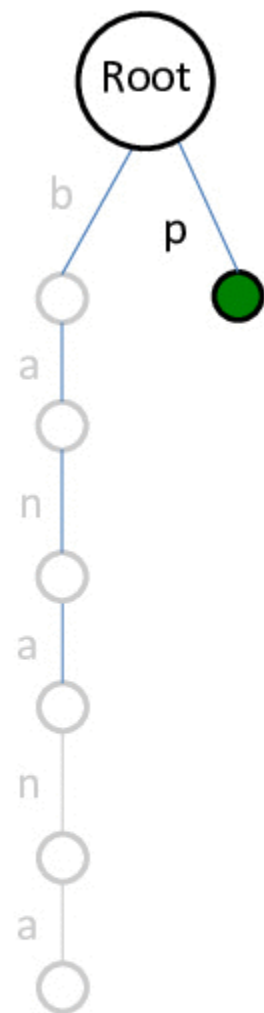
nab

antenna

bandana

ananas

nana



## ***Patterns***

banana

**pan**

and

nab

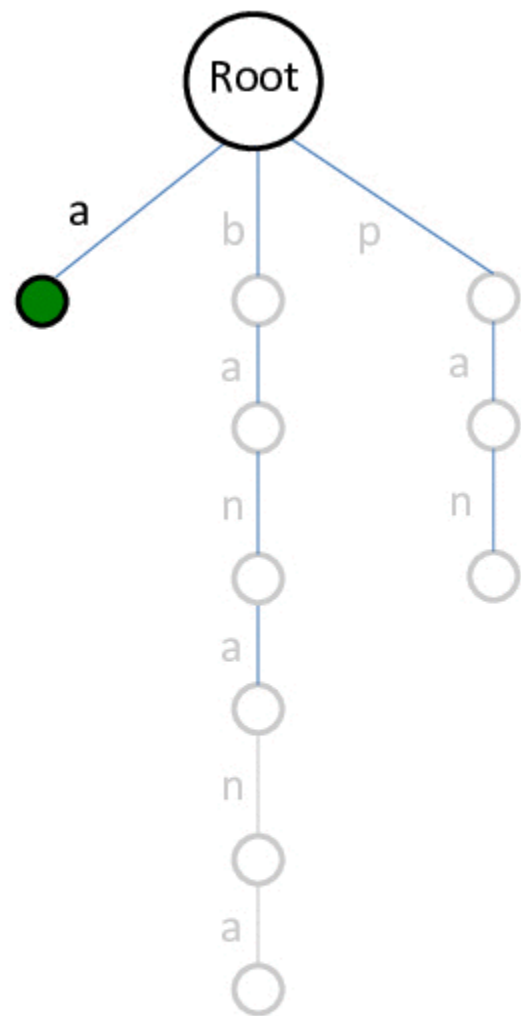
antenna

bandana

ananas

nana





## ***Patterns***

banana

pan

**and**

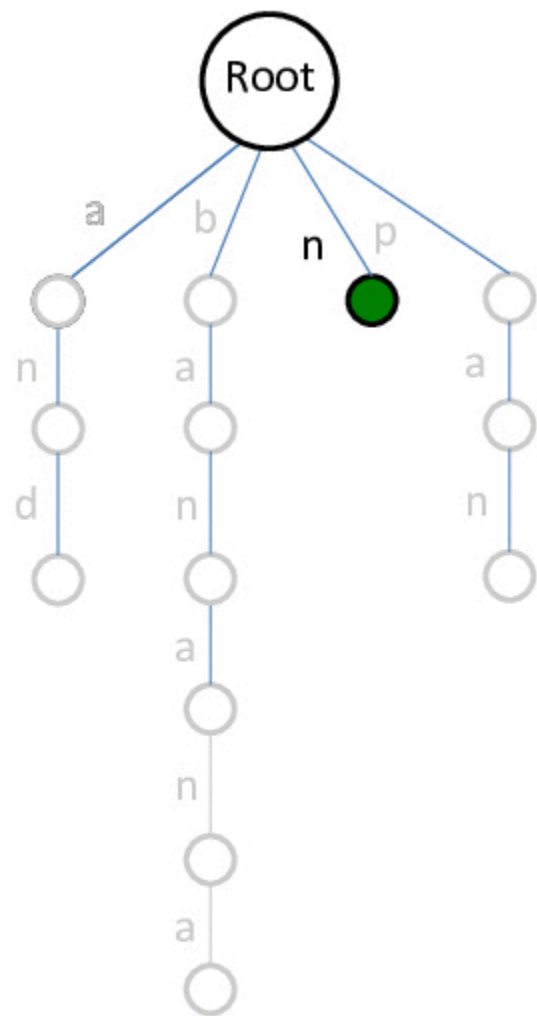
nab

antenna

bandana

ananas

nana



## Patterns

banana

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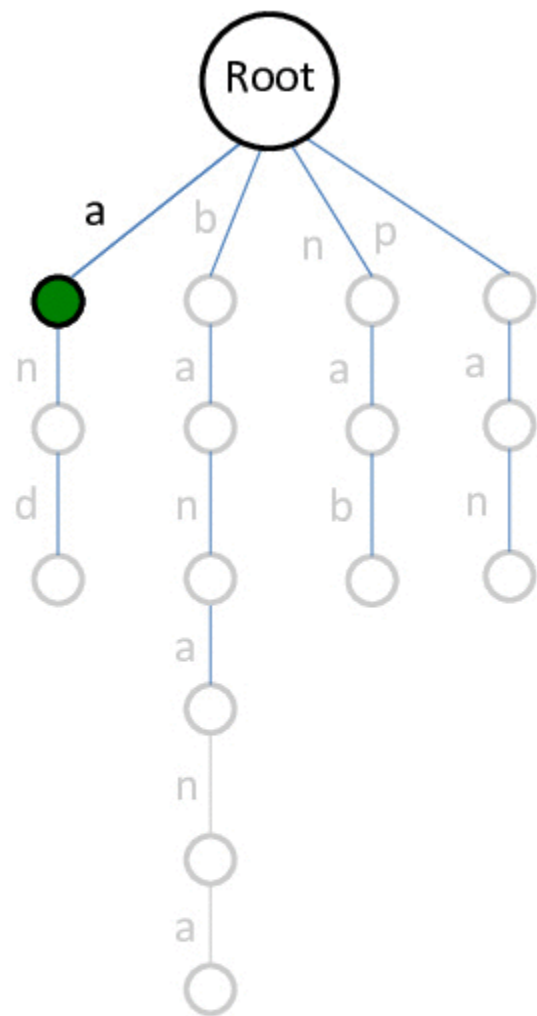
**nab**

antenna

bandana

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nana



## ***Patterns***

banana

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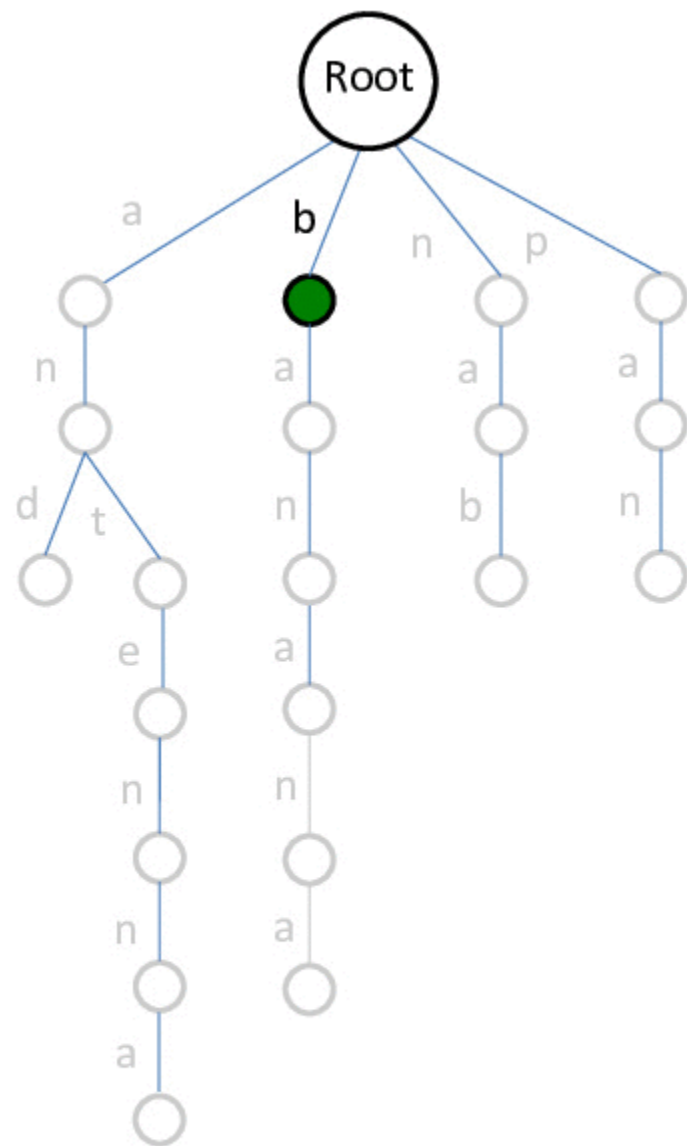
nab

**antenna**

bandana

ananas

nana



## Patterns

banana

pan

and

nab

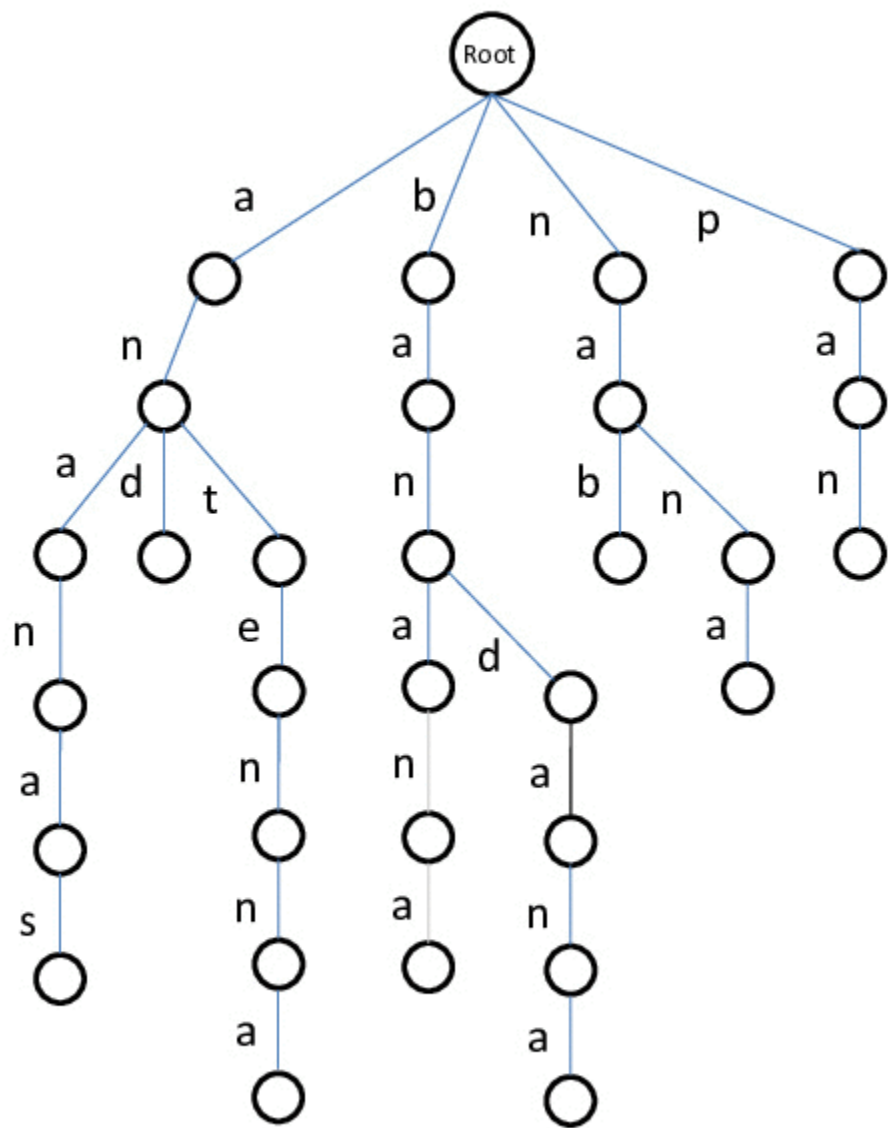
antenna

**bandana**

ananas

nana

panamabananas

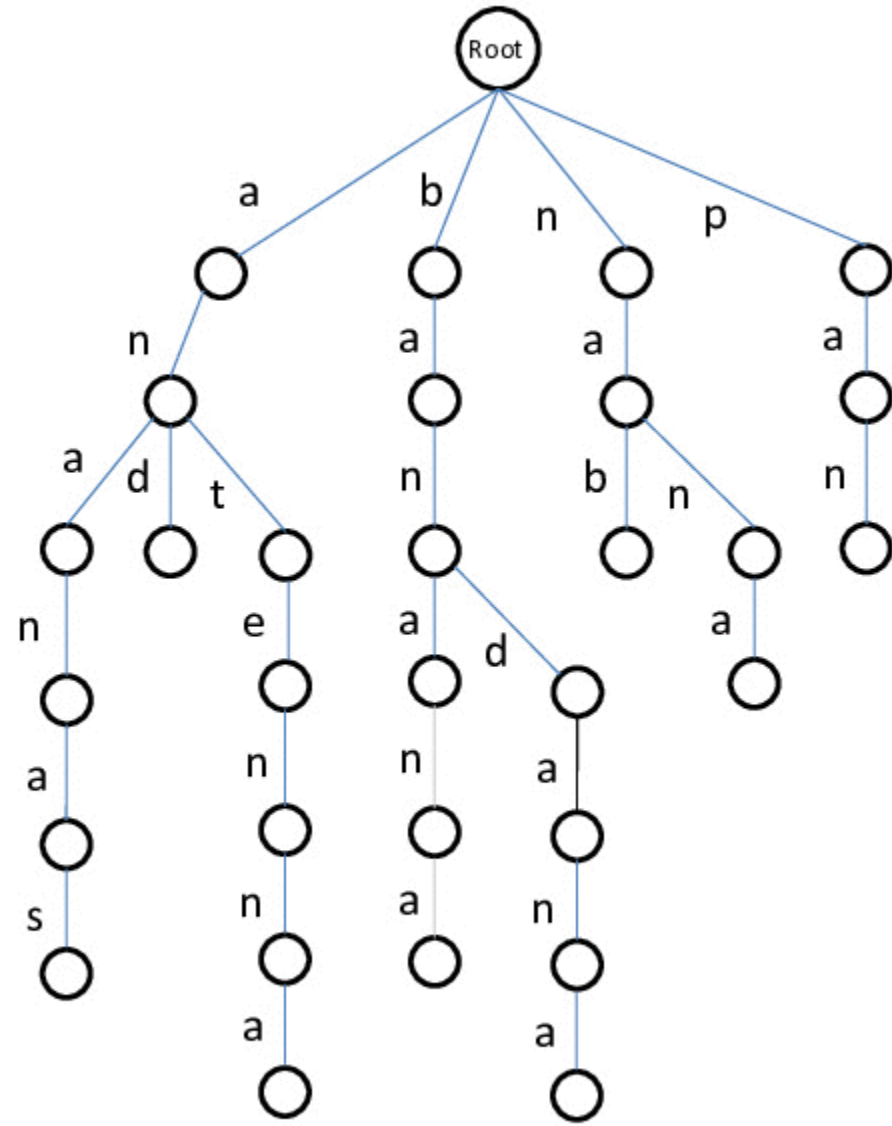


**TrieMatching**(*Text*, *Patterns*):  
drive Trie(*Patterns*) along *Text*  
at each position of *Text*

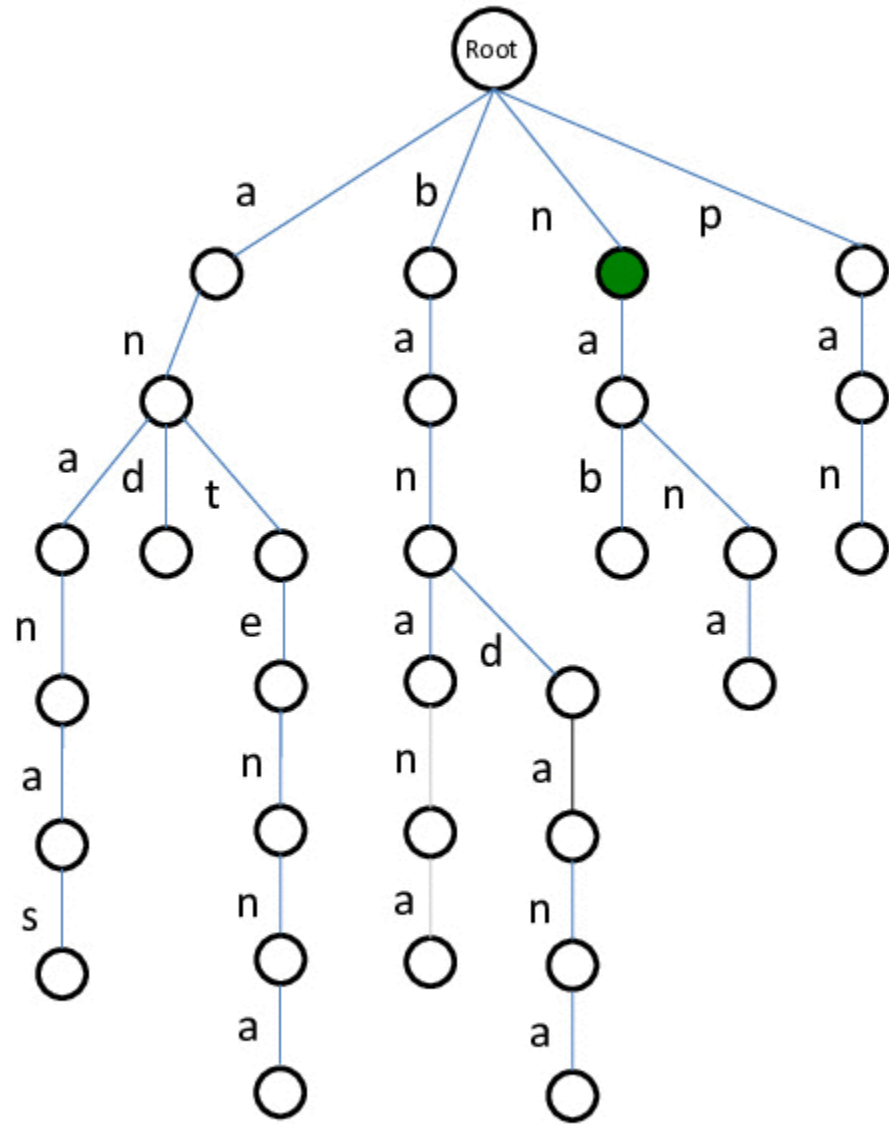
- walk down  $\text{Trie}(\text{Patterns})$  by spelling symbols of  $\text{Text}$
- a pattern from  $\text{Patterns}$  matches  $\text{Text}$  each time you reach a leaf!

For simplicity, we assume that no pattern is a substring of another pattern

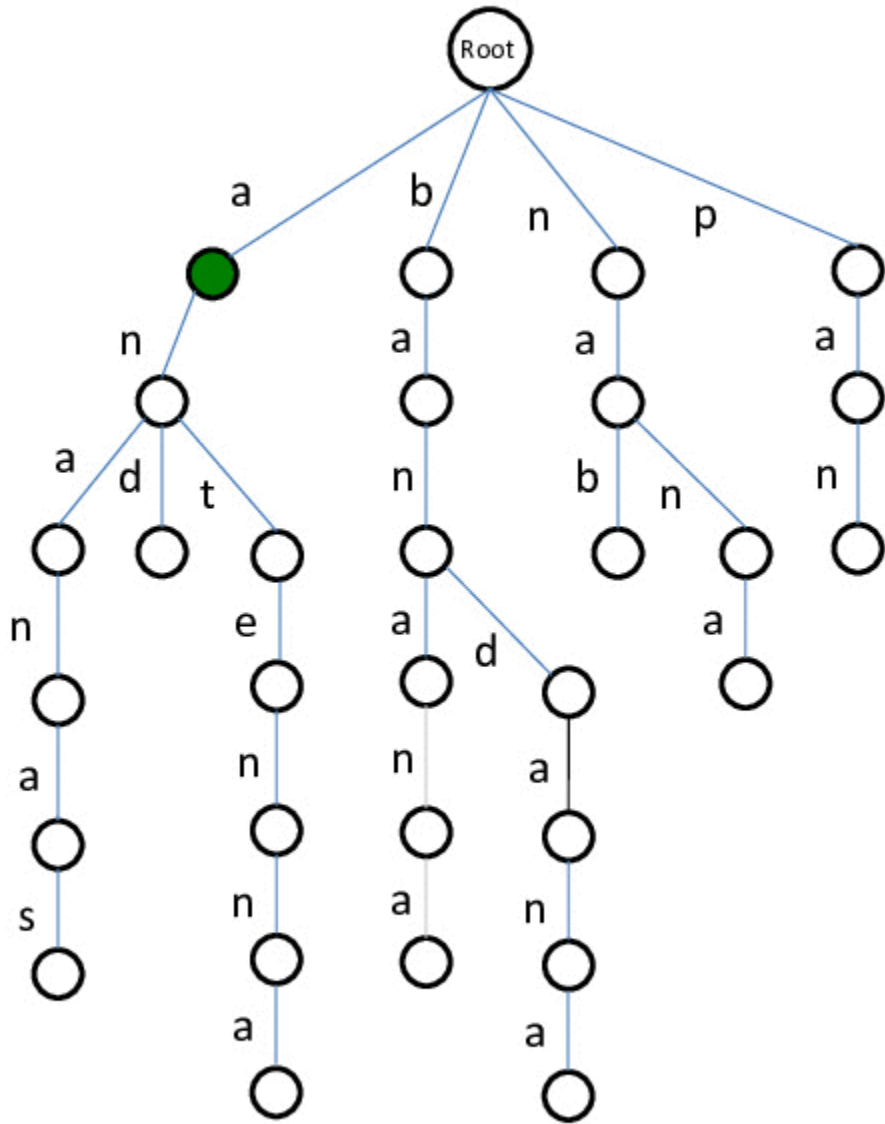
panamabanas



pa**n**amabananas



pan**a**mabanas





# Computational Complexity

- How long does it take to find multiple patterns?
  - Brute force:  $O(|\text{text}| * |\text{patterns}|)$
  - Pattern trie:  $O(|\text{text}| * |\text{LongestPattern}|)$
- Space complexity
  - Number of edges in pattern trie:  $O(|\text{Patterns}|)$
  - Human genome:  $|\text{Patterns}| \approx 10^{12}$

# Alternative: Pack *Text* onto a bus

- Generate all suffixes of *Text*
  - The suffixes represent all possible places a pattern can match
- Form a trie from these suffixes (**suffix trie**)
- For each pattern, walk down from the root of the trie to see if there is a match.

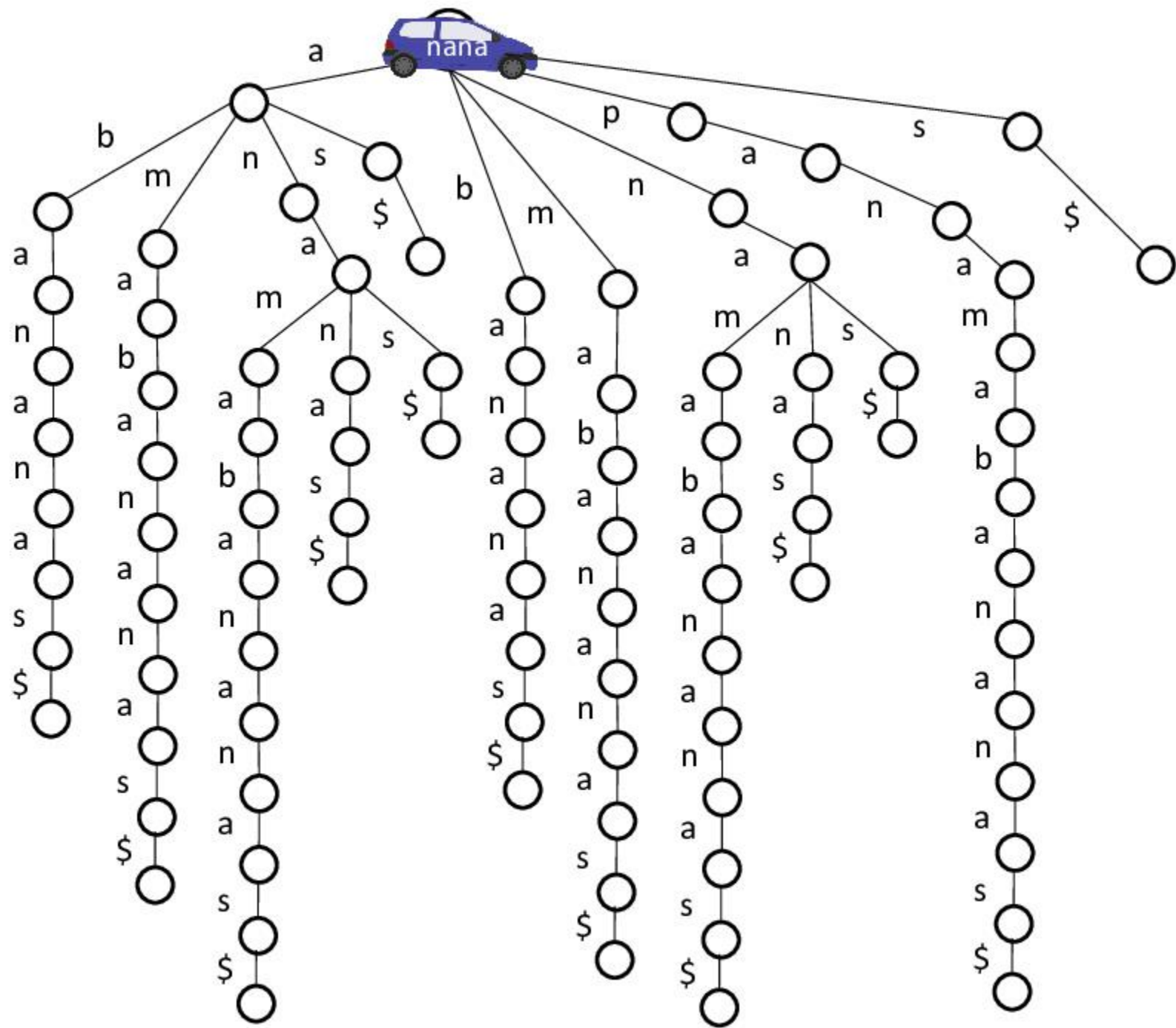
Root

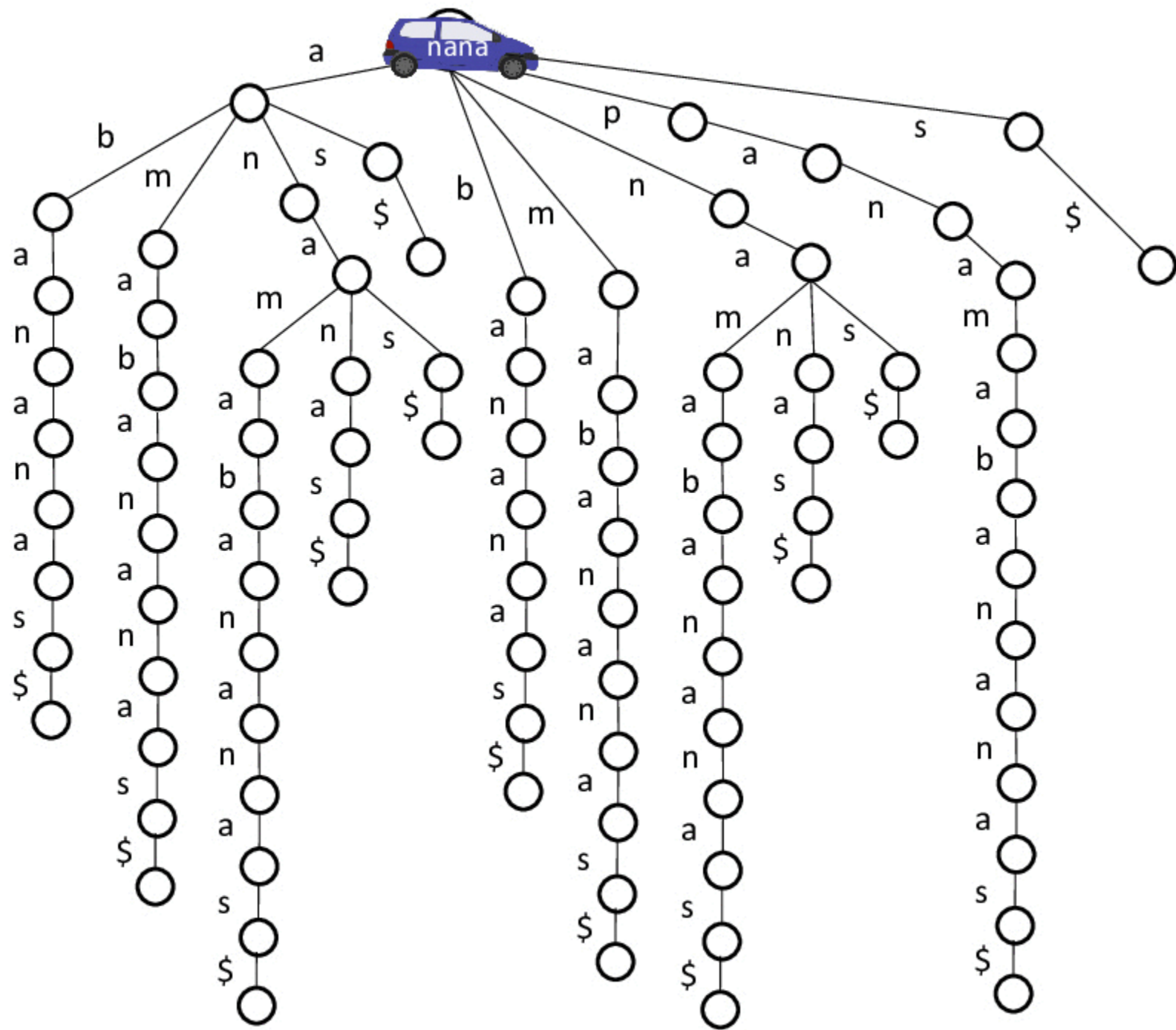
panamabanas

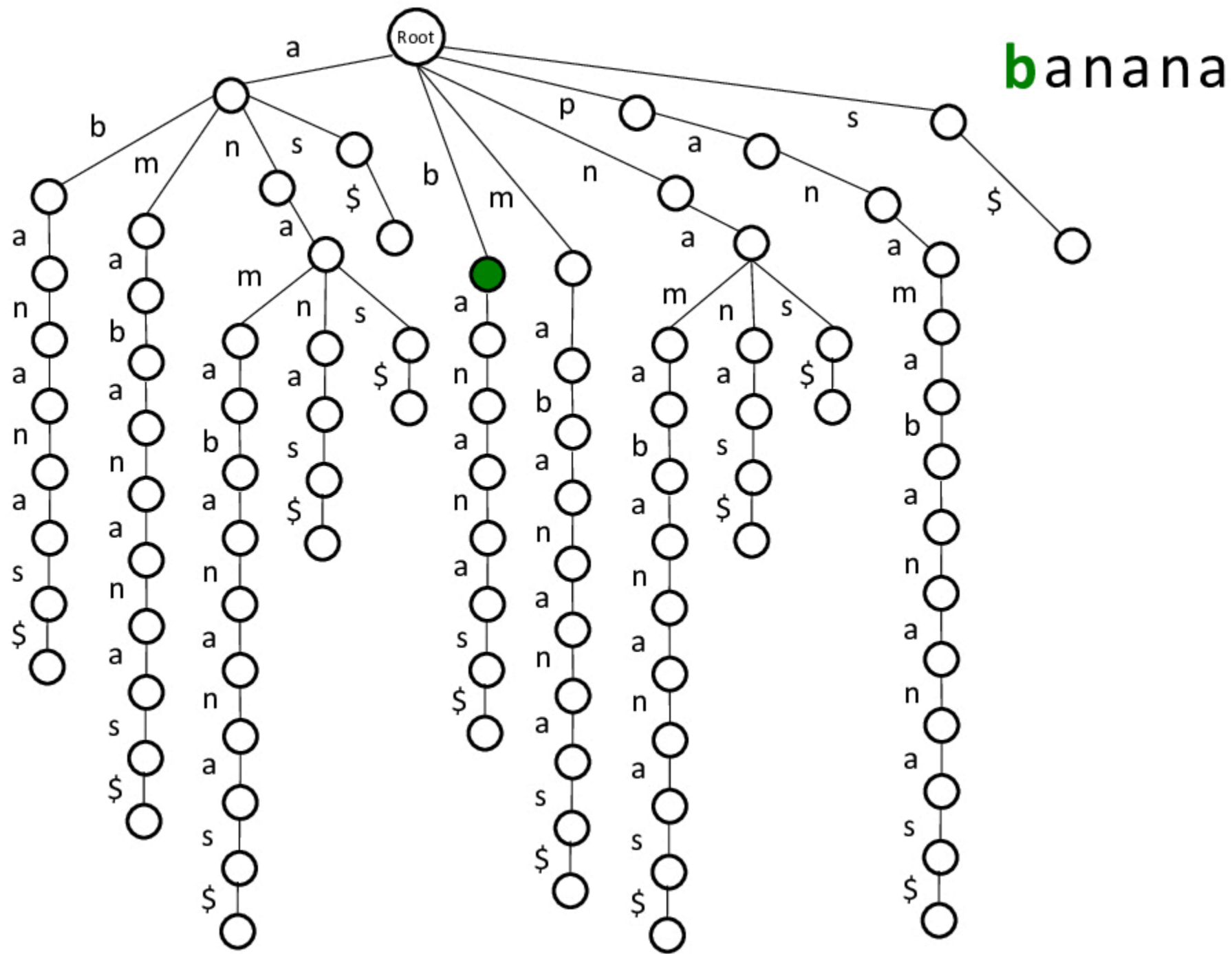
Root

panamabanas\$

Adding “\$” sign in the end (we’ll explain later why)

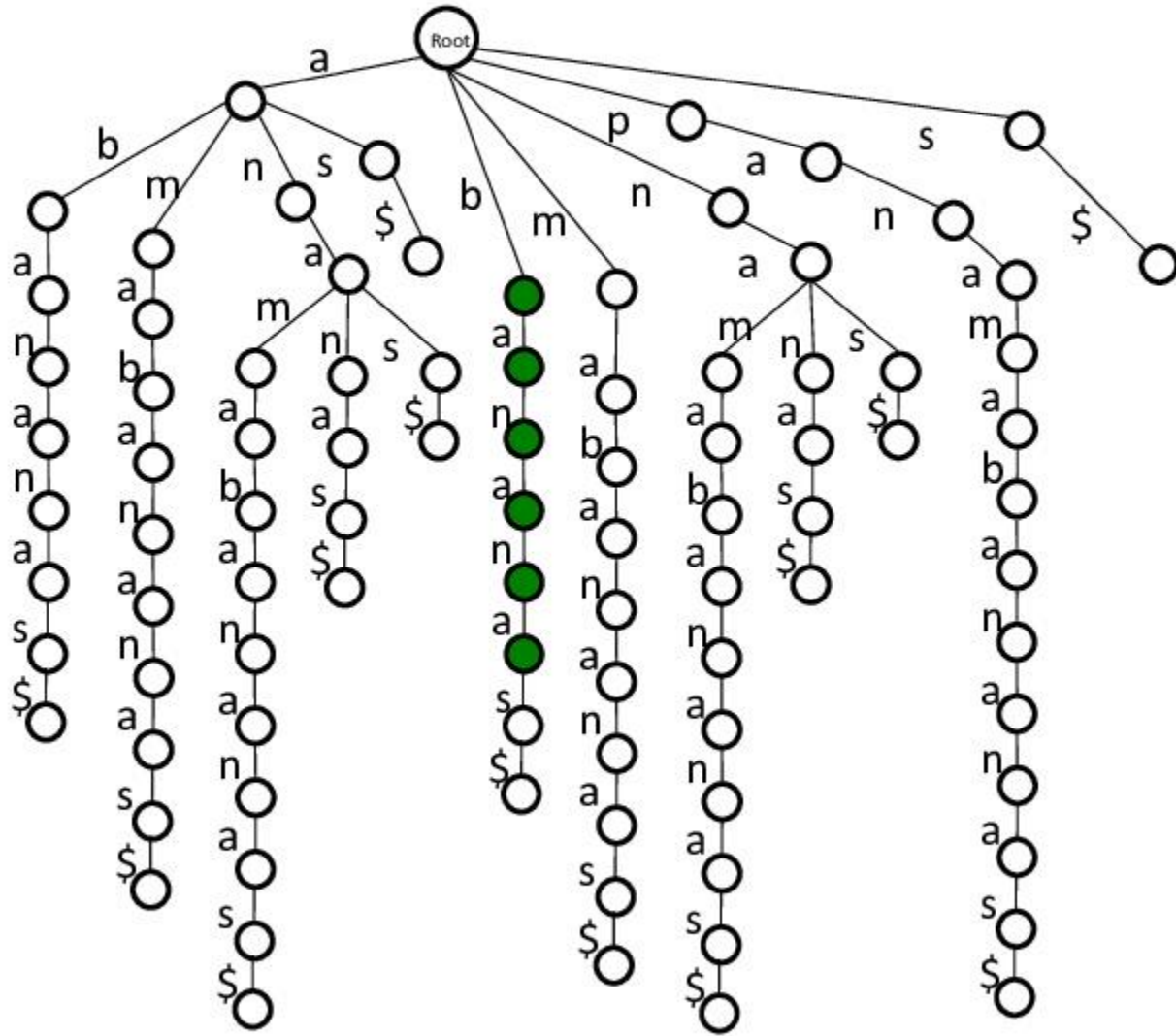






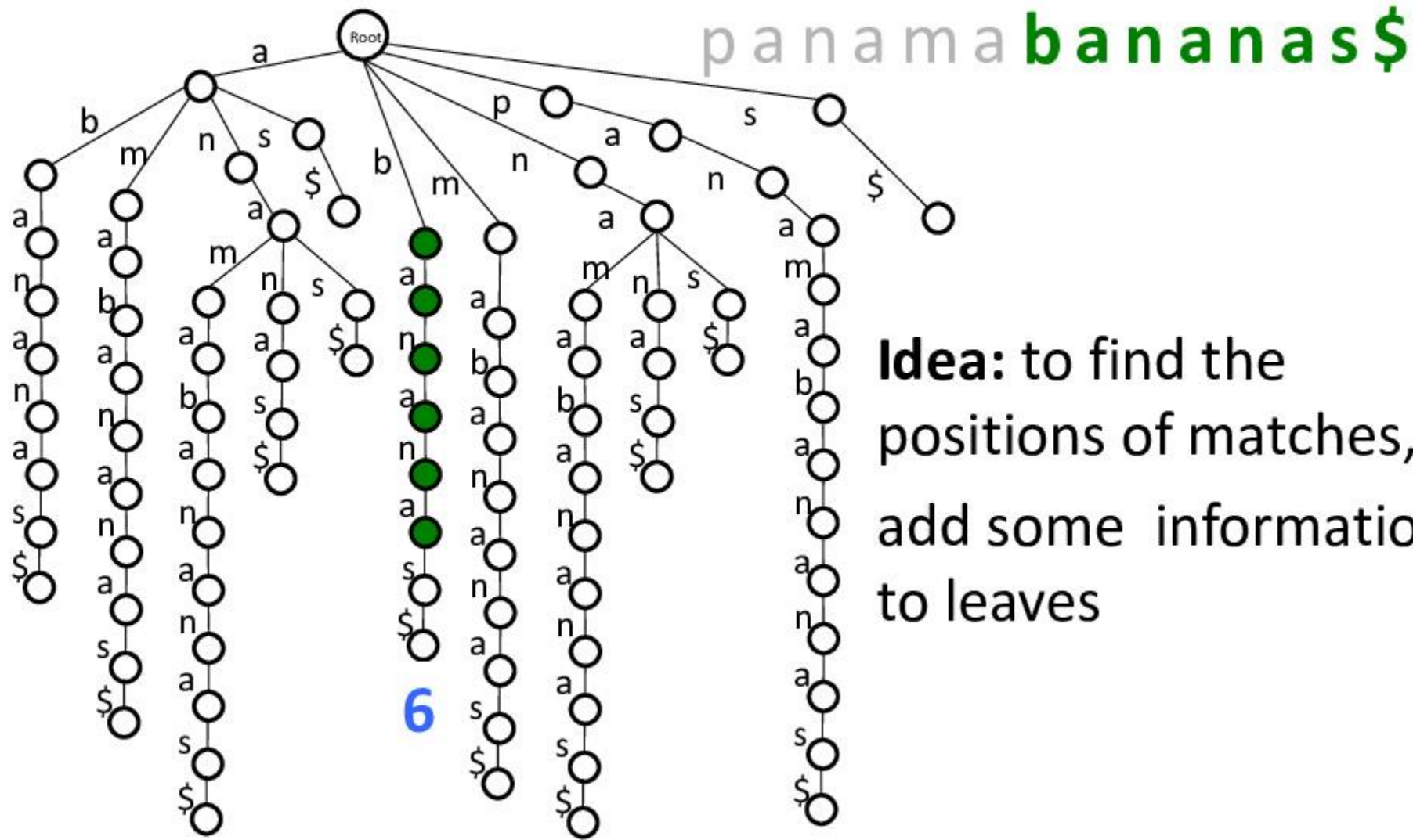
# Where Are the Matches???

bananas\$

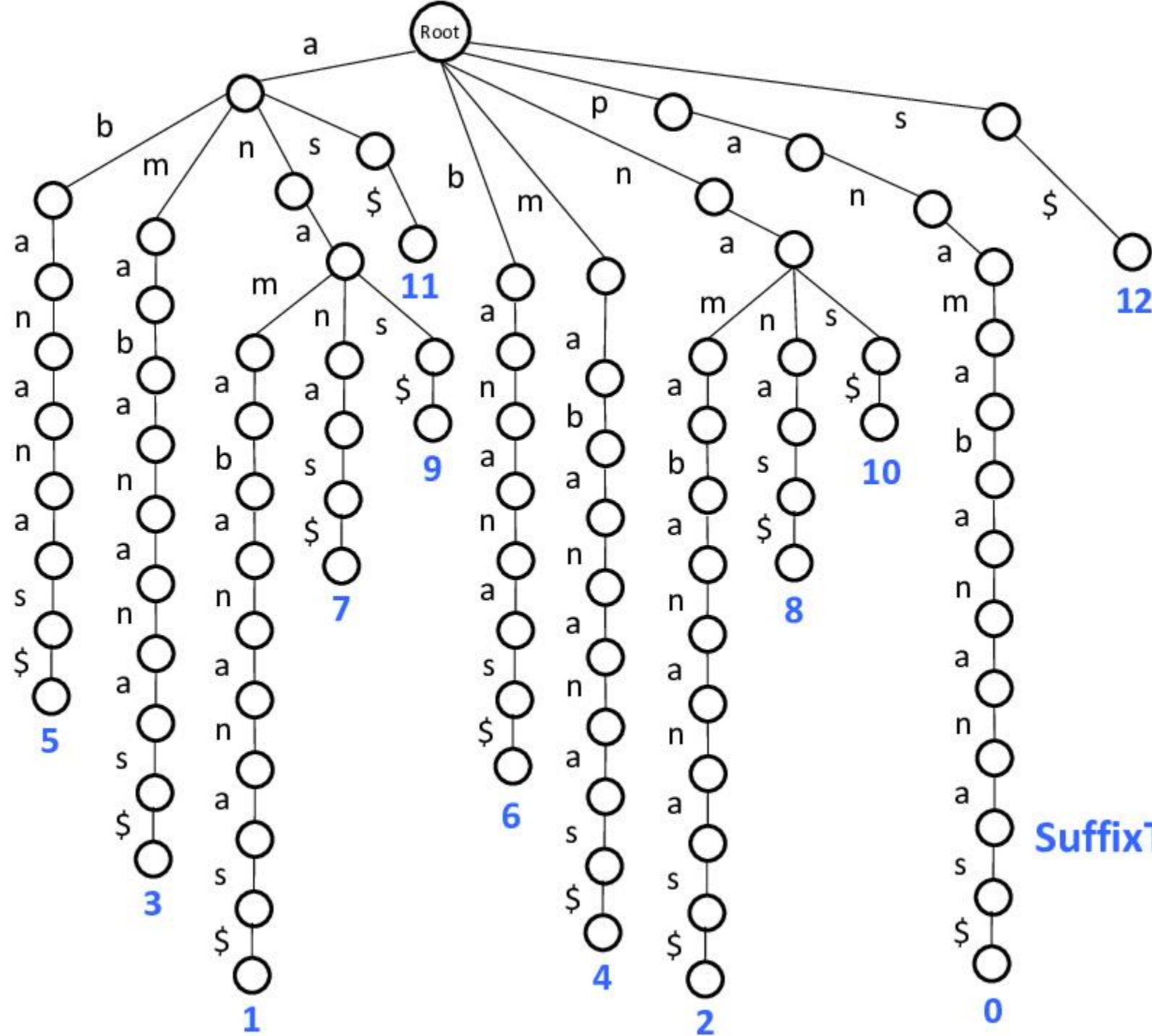




# Where Are the Matches???



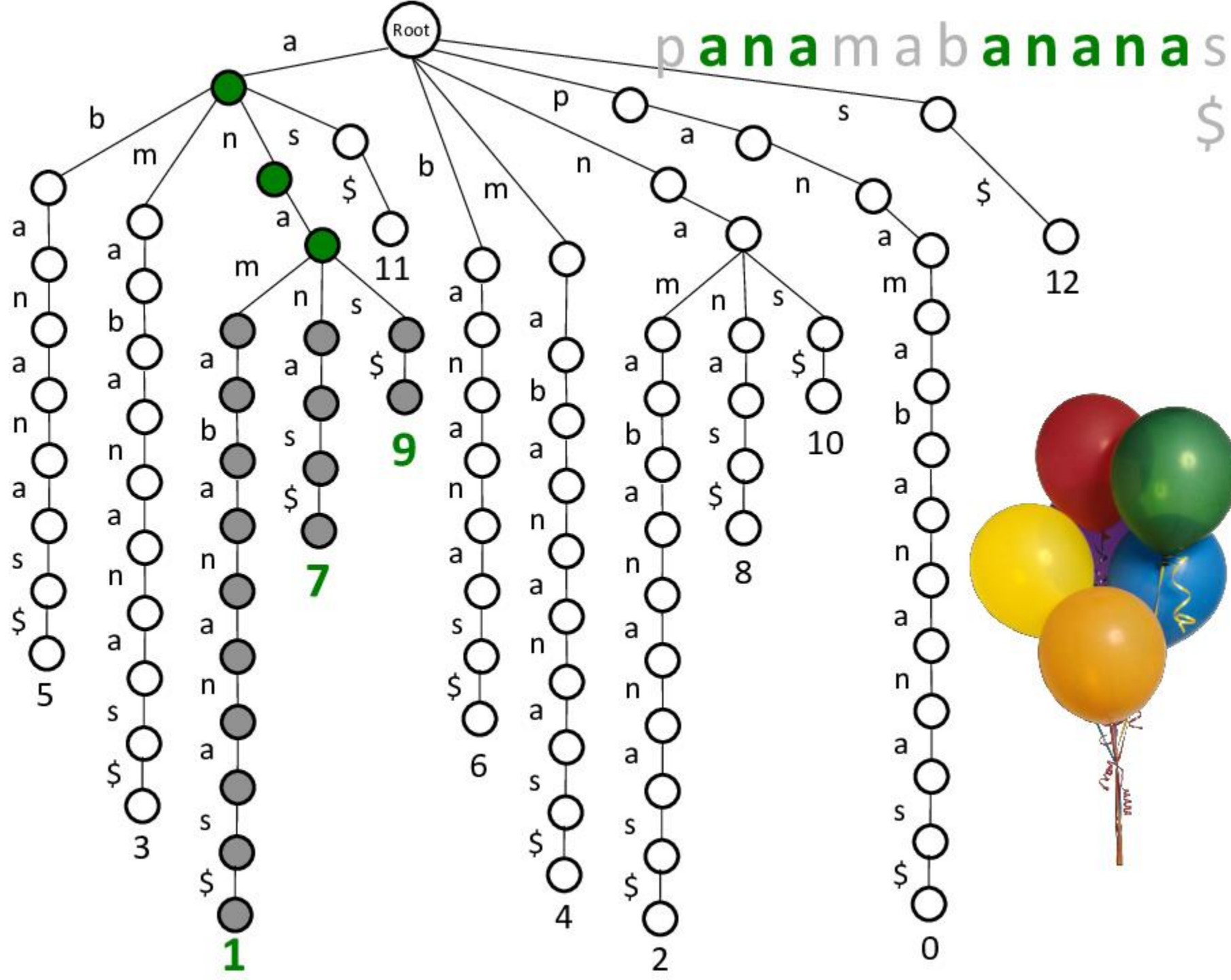
**Idea:** to find the positions of matches, add some information to leaves



SuffixTrie(Text)

# Identifying position of match in text

- Once we find a match, walk down to the leaves to get the position(s) of the matches.



## Memory Footprint of Suffix Trie

The suffix trie is formed from  $|Text|$  suffixes with total length:

$$|Text| * (|Text| - 1) / 2$$

For human genome:

- $|Text| \approx 3 \cdot 10^9$

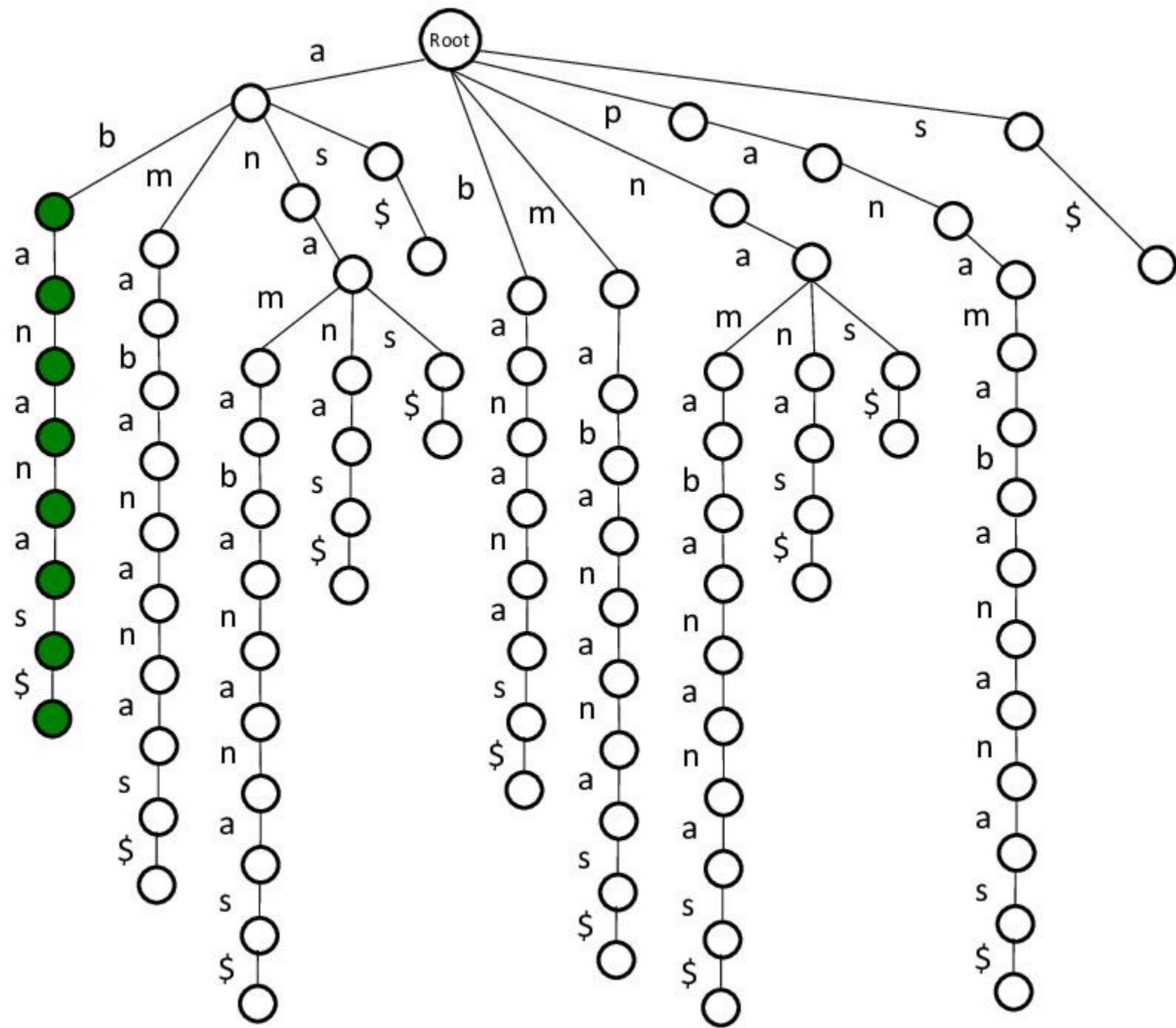


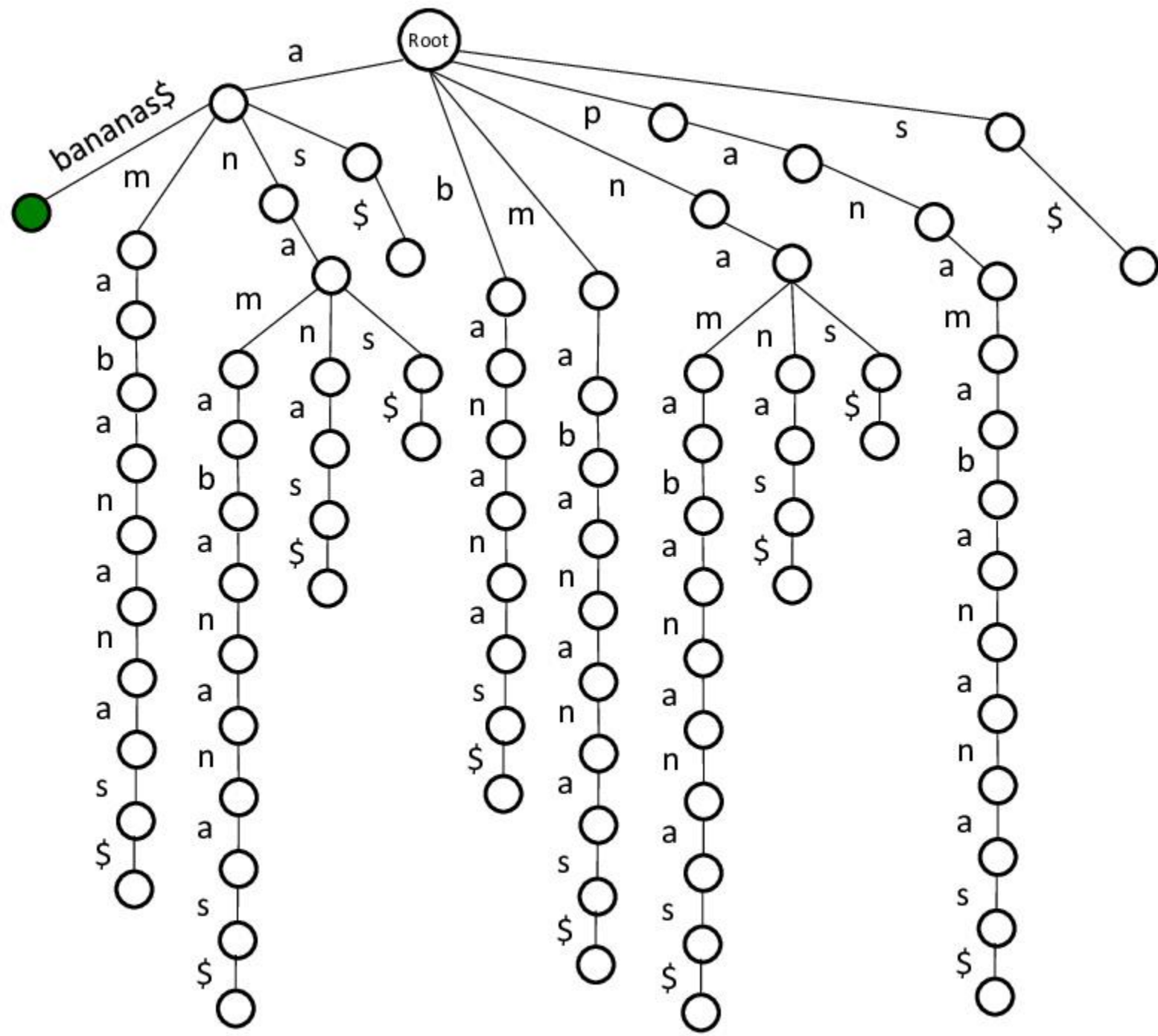
$| \textit{Text} |$  symbols

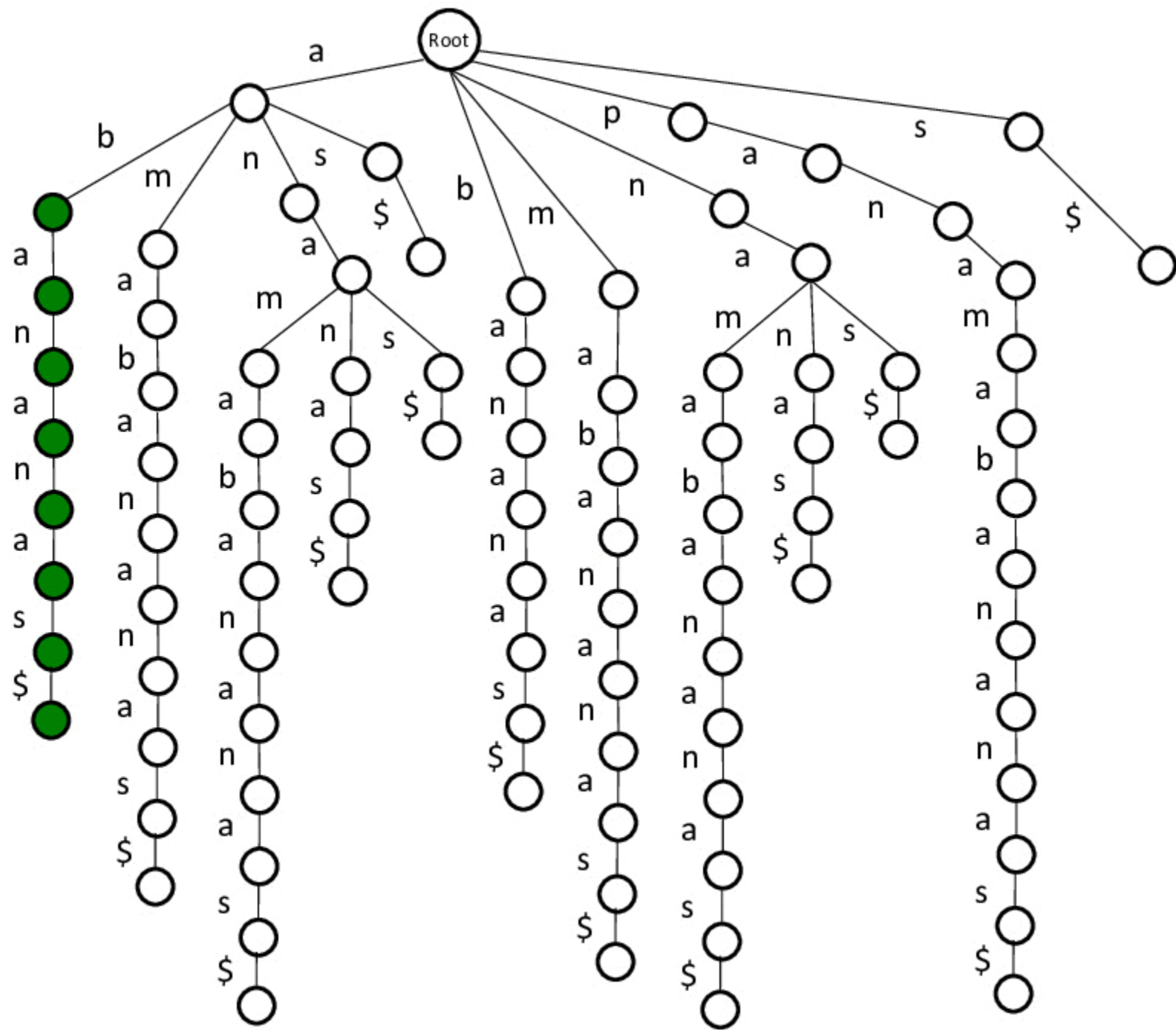
p	a	n	a	m	a	b	a	n	a	n	a	s	\$
	a	n	a	m	a	b	a	n	a	n	a	s	\$
		n	a	m	a	b	a	n	a	n	a	s	\$
			a	m	a	b	a	n	a	n	a	s	\$
				m	a	b	a	n	a	n	a	s	\$
					a	b	a	n	a	n	a	s	\$
						b	a	n	a	n	a	s	\$
							a	n	a	n	a	s	\$
								n	a	n	a	s	\$
									a	n	a	s	\$
										n	a	s	\$
											a	s	\$
												s	\$
													\$

$| \textit{Text} |$  suffixes

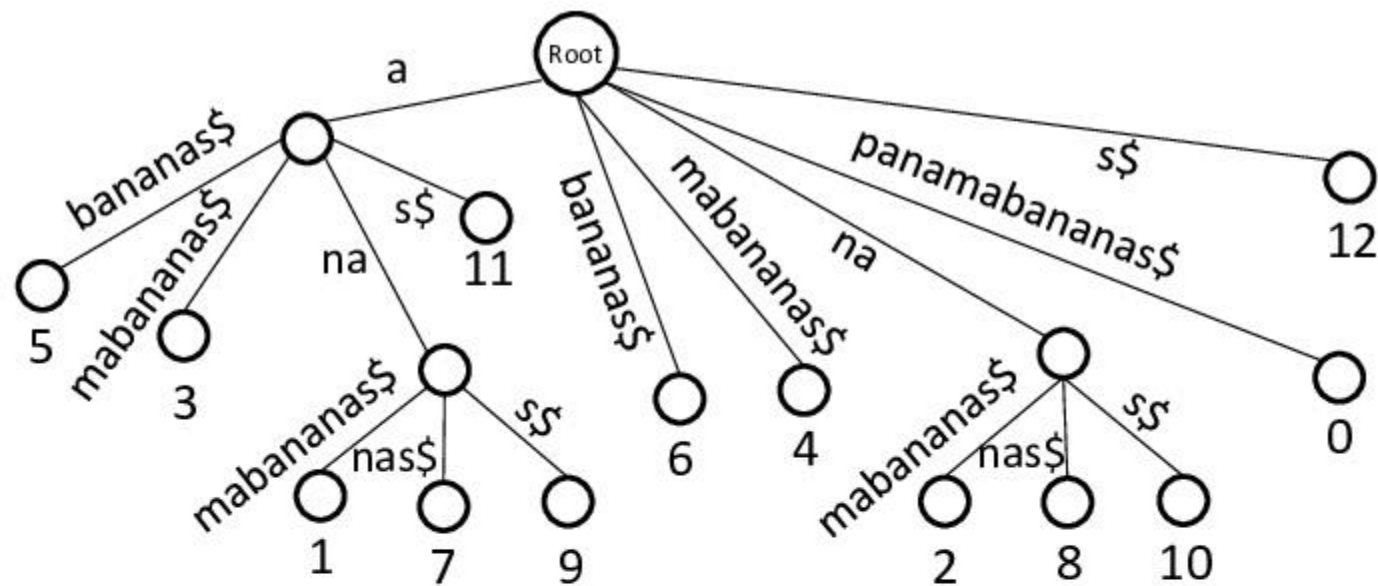






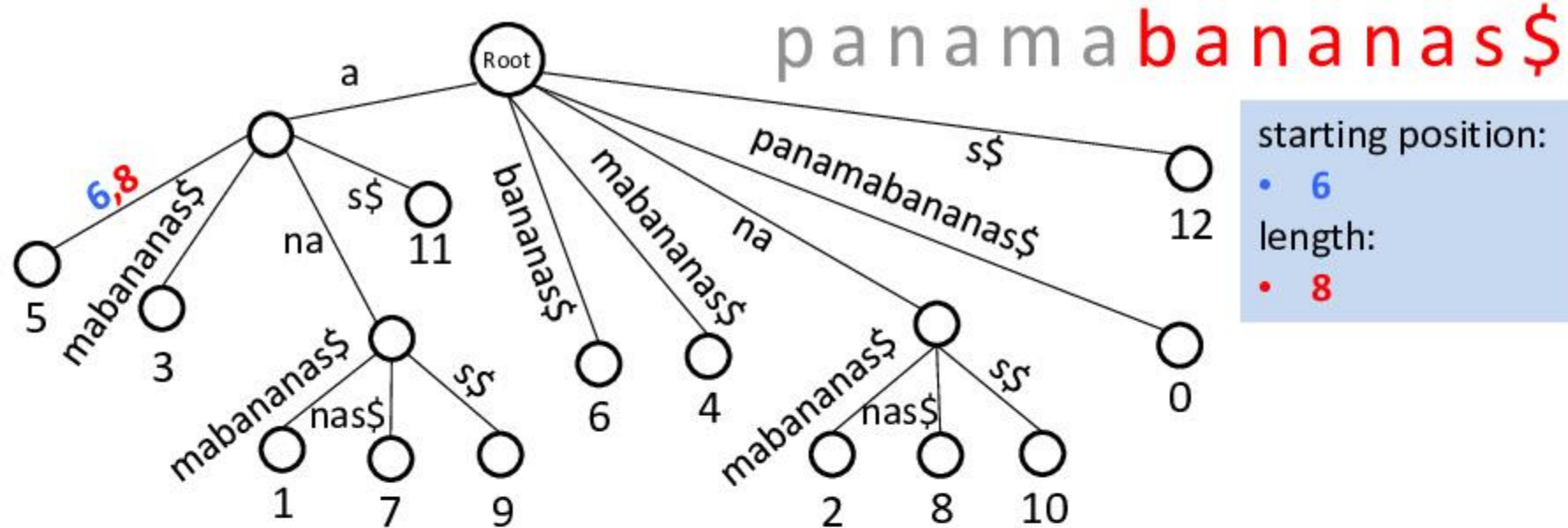






Since each suffix adds one leaf and at most one internal vertex to the suffix tree:

- $\# \text{ vertices} < 2 | \text{Text} |$
- memory footprint of the suffix tree:  $O(| \text{Text} |)$



Since each suffix adds one leaf and at most one internal vertex to the suffix tree:

- $\# \text{ vertices} < 2 | \text{Text} |$
- memory footprint of the suffix tree:  $O(| \text{Text} |)$
- storing edge labels

# Overview of Suffix Tree

- Fast **Exact** Multiple Pattern Matching
  - Time:  $O(|\text{text}| + |\text{patterns}|)$
  - Memory:  $O(|\text{text}|)$ 
    - Actual implementation still too demanding in memory requirements.  
 $\sim 20 * |\text{text}|$
- Construction:  $O(|\text{text}|)$
- Need better method that:
  - Can handle mutations (approximate matching)
  - Has smaller memory footprint