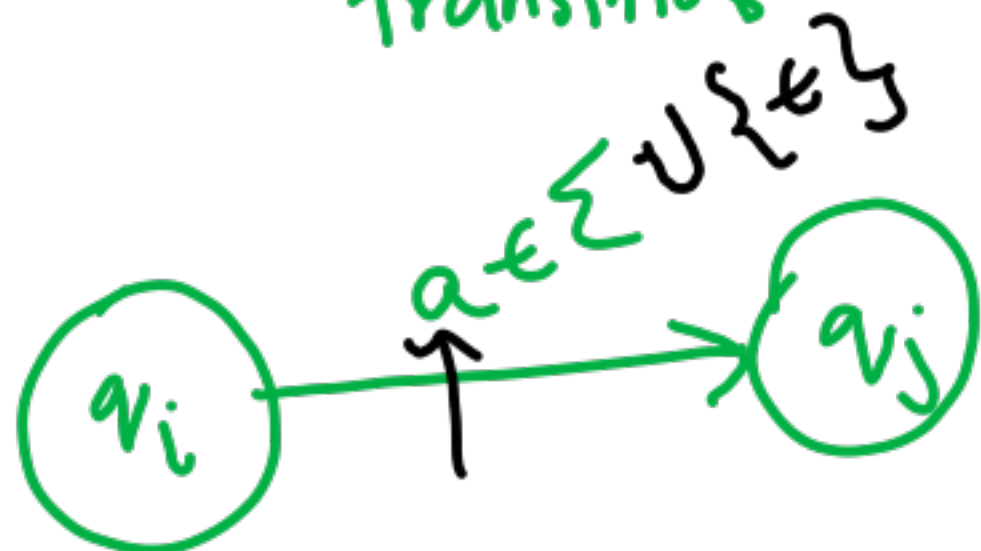


22.8.2025

FSA

$$\langle \Sigma, Q, \delta, q_0, F \rangle$$

transitions
Start



Systematic

NFA + ϵ

$$\delta: Q \times \Sigma \cup \{\epsilon\} \rightarrow Q$$

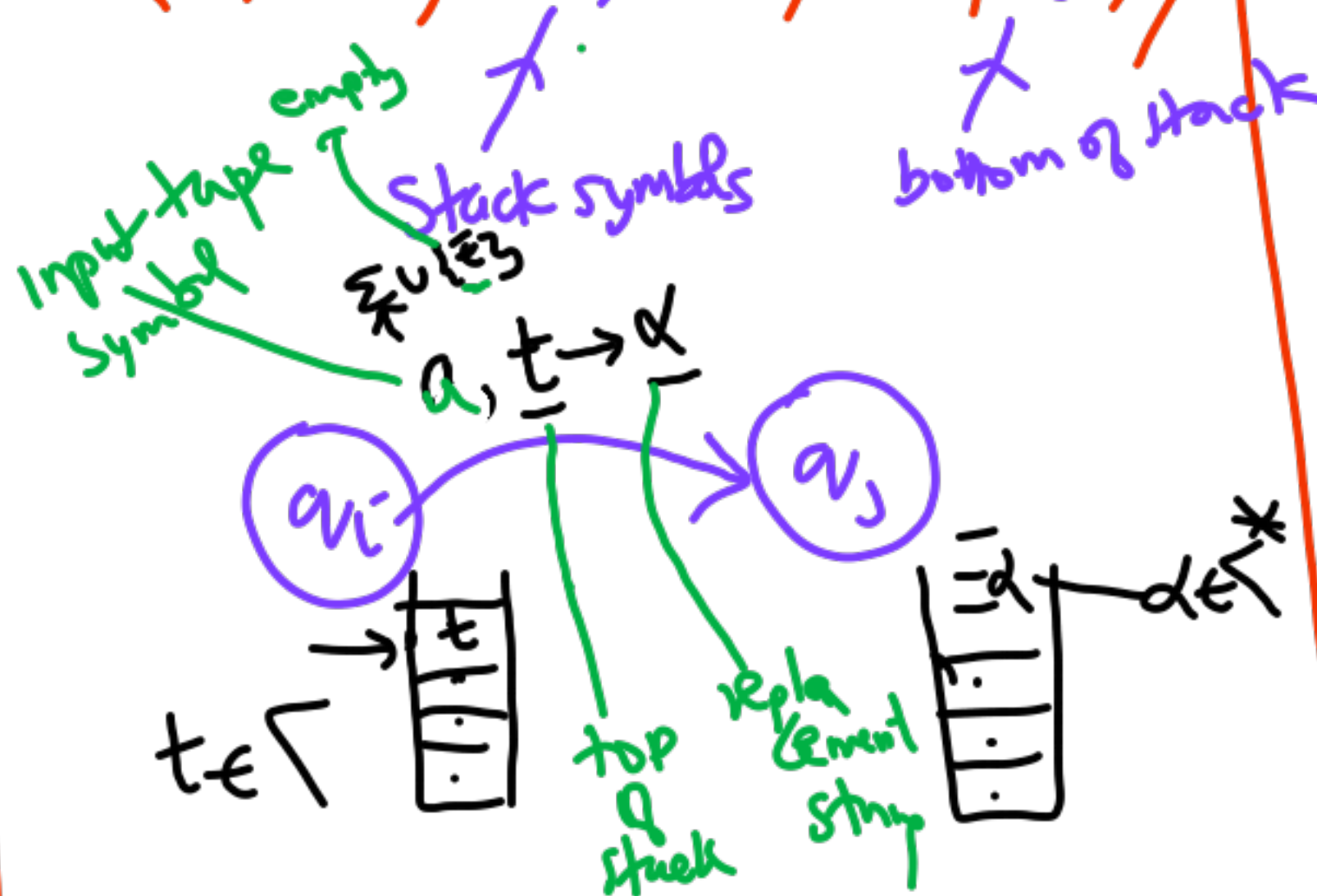
No ϵ -transitions incl. \emptyset
any subset

function $Q \times \Sigma \rightarrow Q$

exactly one next state

PDA

$$\langle Q, \Sigma, \Gamma, \delta, q_0, z_0, F \rangle$$



$$\delta: Q \times \Sigma \cup \{\epsilon\} \times \Gamma \rightarrow 2^{Q \times \Gamma^*}$$

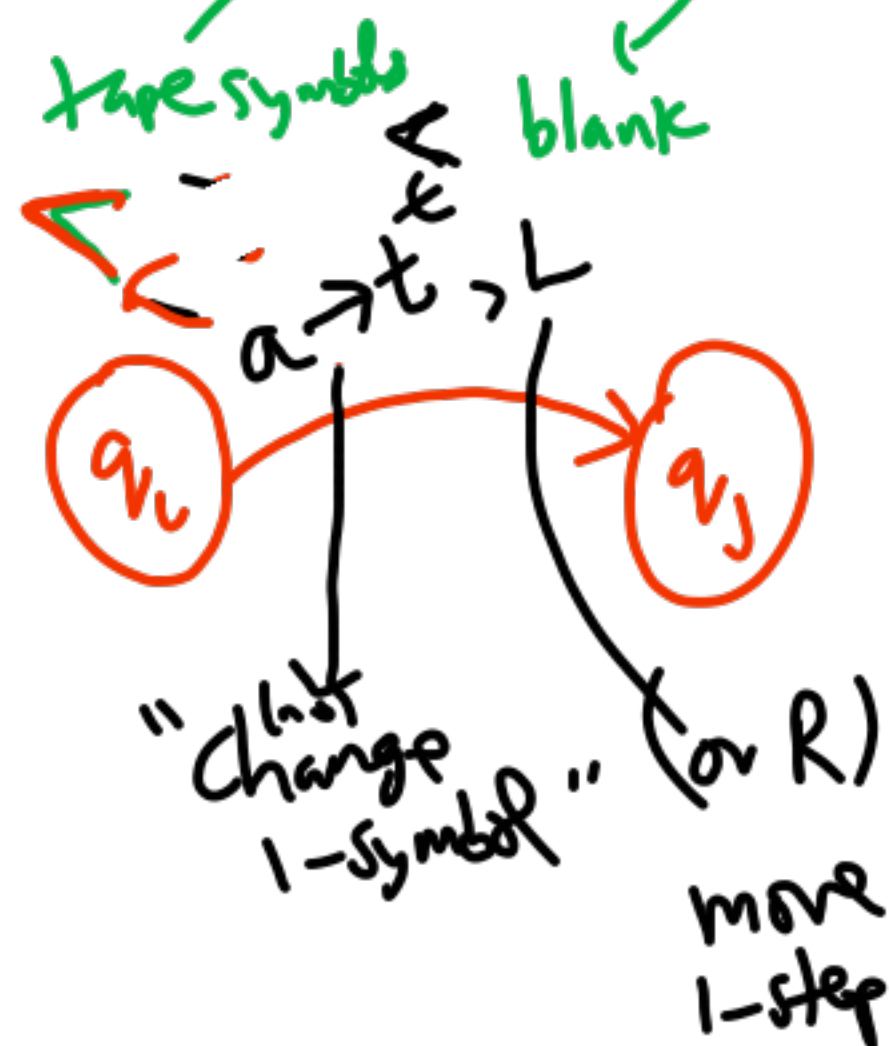
at most one next state

no choice \Rightarrow deterministic
no arrow for a, t not

Can replace only top symbol

TM

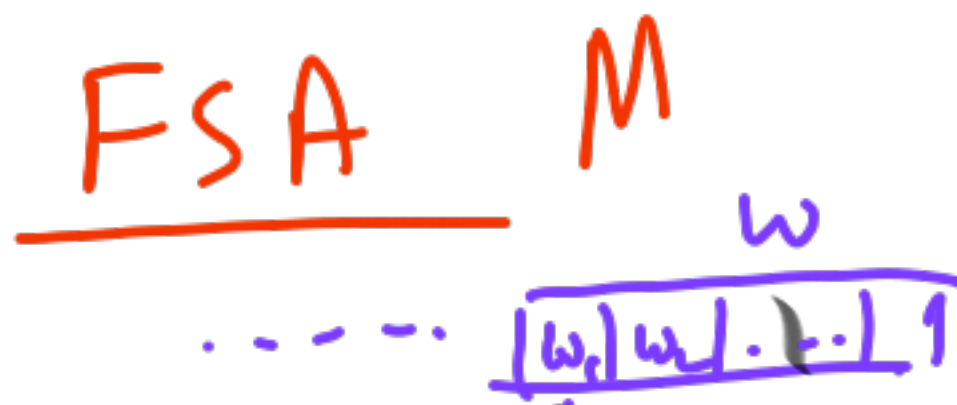
$$\langle Q, \Sigma, \Gamma, \delta, q_0, \emptyset, F \rangle$$



$$\delta: Q \times \Gamma \rightarrow Q \times \Gamma \cup \{\emptyset\}$$

Crash if \emptyset no move!
 $\emptyset =$ possible

Input tape
b | a | ...



ID $\langle q_i, w \rangle$

Run is \vdash^* Any number of steps

$$L(M) = \{ w \mid M \text{ accepts } w \}$$

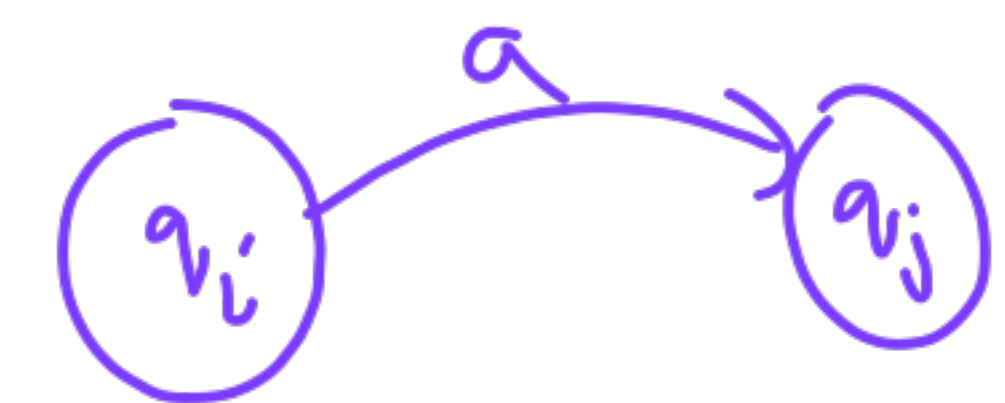
try all for non-det M

Move \vdash ID₁ \vdash ID₂

There is a run

$$\langle q_0, w \rangle \vdash^* \langle q_f, \epsilon \rangle$$

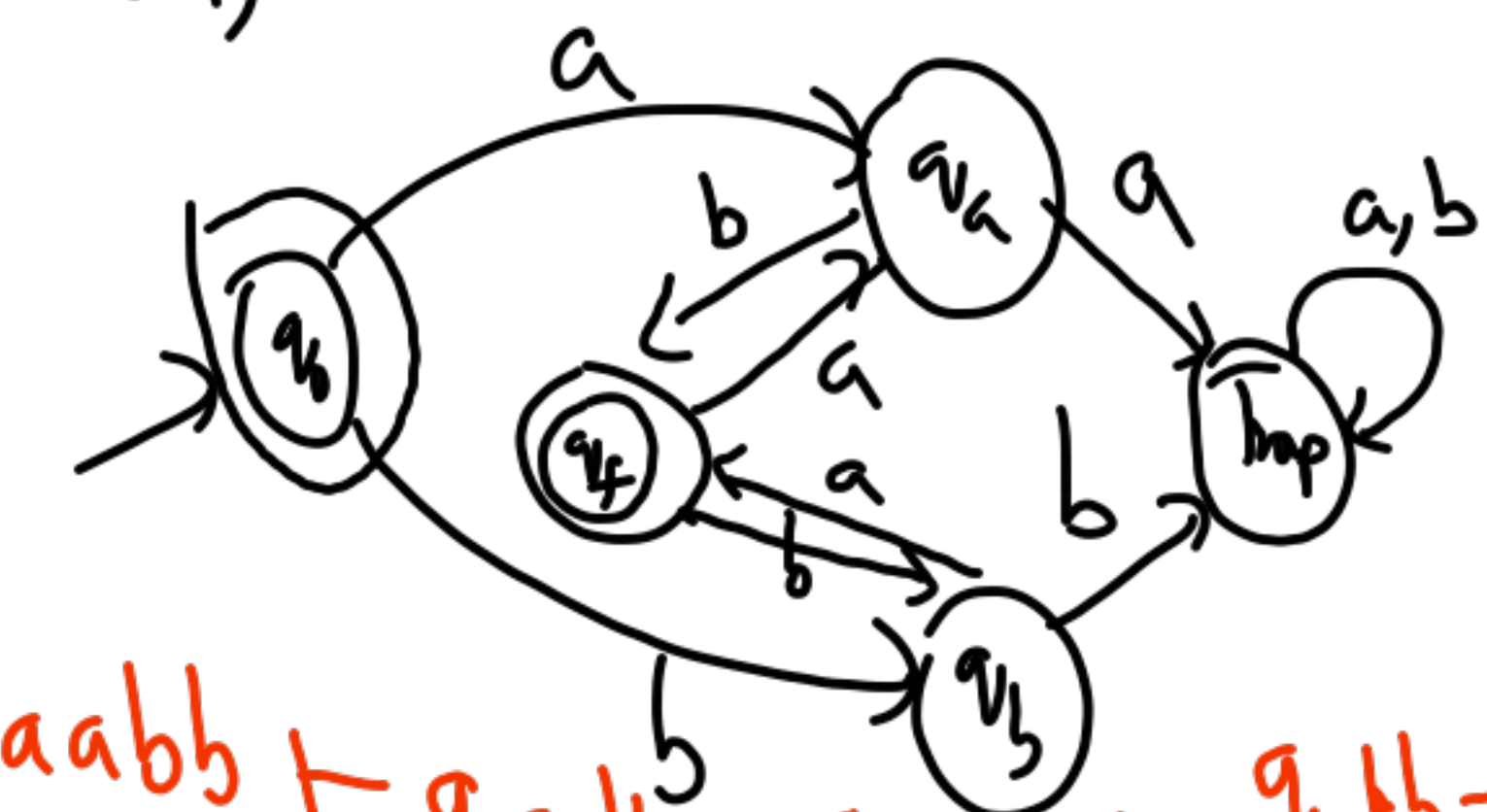
① $\langle q_i, \underline{a}w \rangle \vdash \langle q_j, w \rangle$



② $\langle q_i, \underline{a}w \rangle \vdash \langle q_j, \underline{a}w \rangle$



$(ab+ba)^*$ $q_f \in F$



$q_0 baabb \vdash q_1 aabb \vdash q_f abb \vdash q_a bb \rightarrow q_f b \vdash q_b$

PDA M

ID

$\langle q, w, \alpha \rangle$

$\epsilon \epsilon^*$

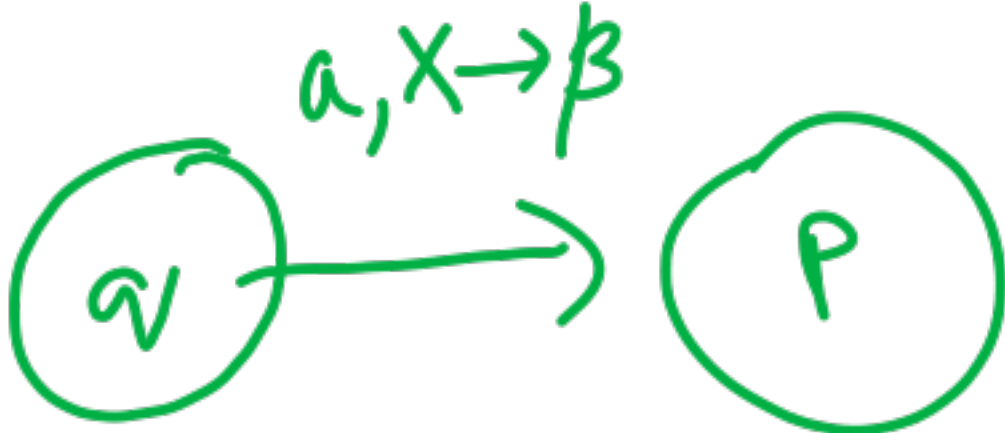
$\epsilon \epsilon^*$

input symbol

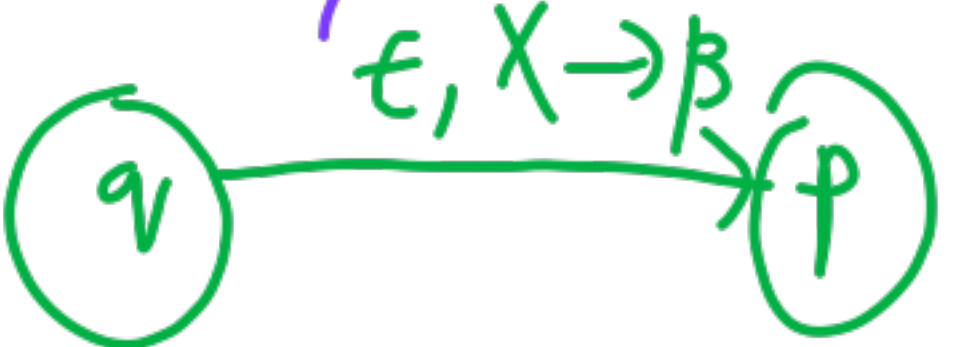
top of stack

(1) $\langle q, aw, X\alpha \rangle$

$\vdash \langle p, w, \beta\alpha \rangle$



(2) $\langle q, aw, X\alpha \rangle \vdash \langle p, aw, \beta\alpha \rangle$



$L(M) = \{w \mid M \text{ accepts } w\}$

by final state

$\langle q_0, w, z_0 \rangle$

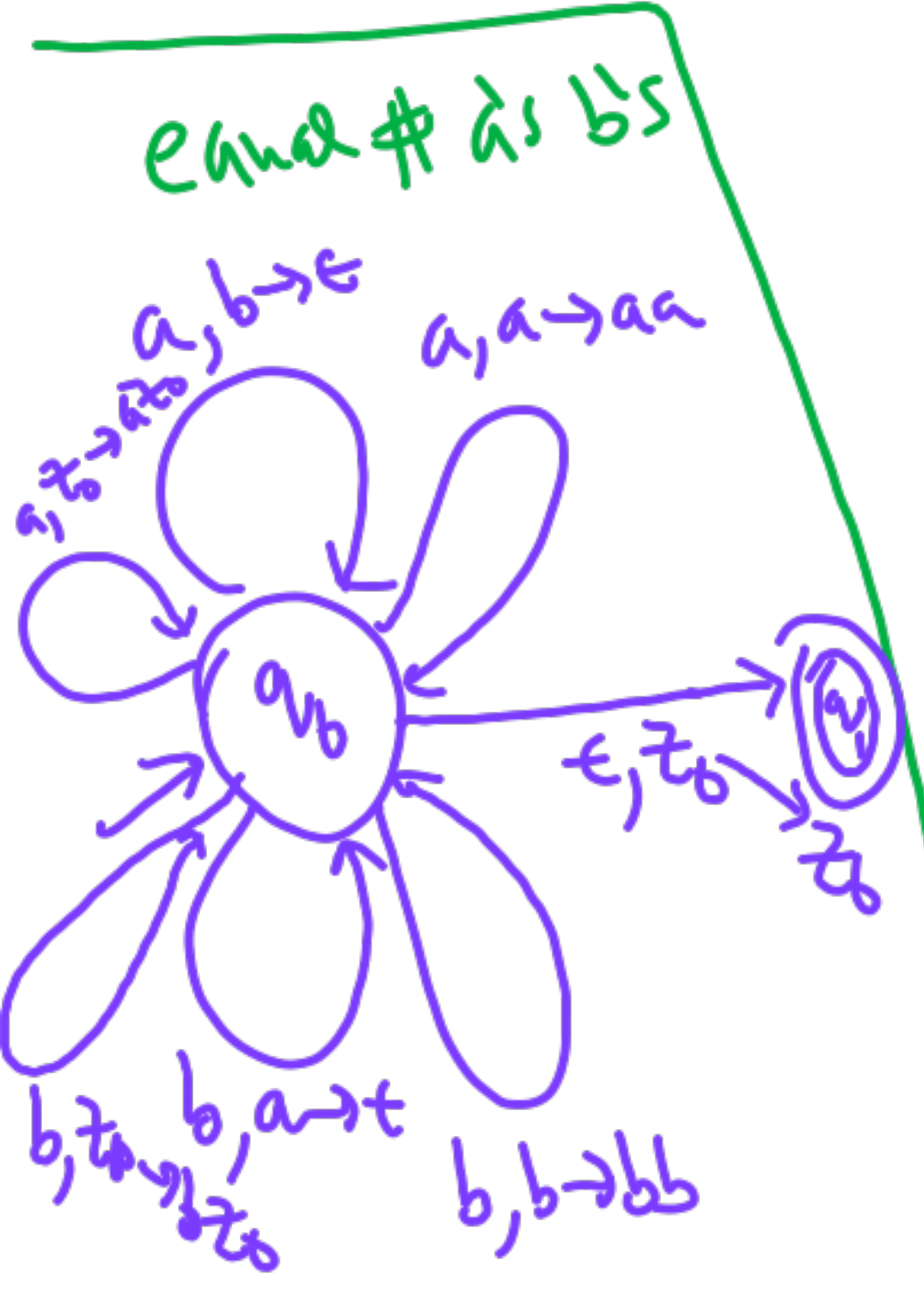
\vdash^*

$\langle q_f, \epsilon, \alpha \rangle$

in Γ^*

empties (all input processed)

anything on stack



$\{a^i b^j c^k \mid i=j \text{ or } j=k\}$

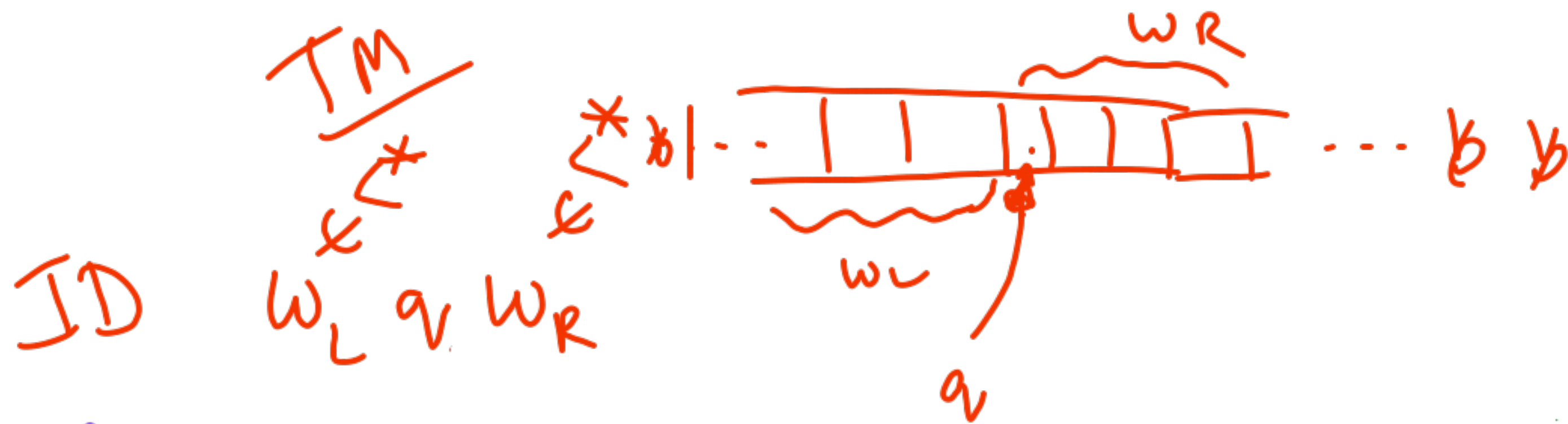
$a^i b^j c^k$

$n=bl$

q_0

q_1

q_2

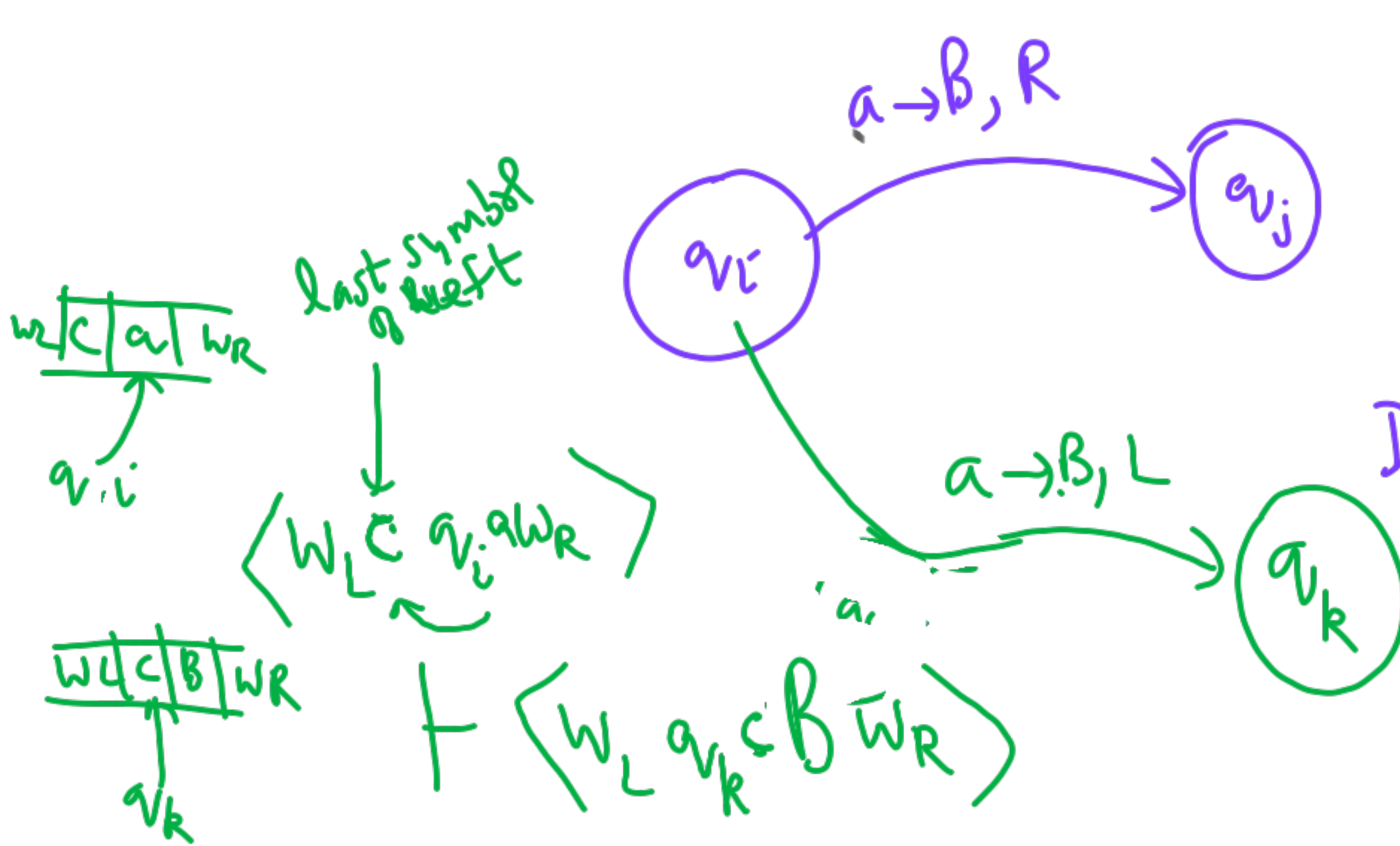


accepts

move

$$w_L q_i a w_R \vdash w_L B q_j w_R$$

$$q_0 w \vdash^* \underbrace{x_1}_{\text{any string}} q_f \underbrace{x_2}_{\text{any string}}$$



turing machine	
	a, b
$w w$	
$w \# w$	given

DO

without #

Final state

(Halt)

(no more possible)

