Machine-Level Programming IV: Data

15-213: Introduction to Computer Systems

8th Lecture, Sep. 24, 2015

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Today

Arrays

- One-dimensional
- Multi-dimensional (nested)
- Multi-level

Structures

- Allocation
- Access
- Alignment

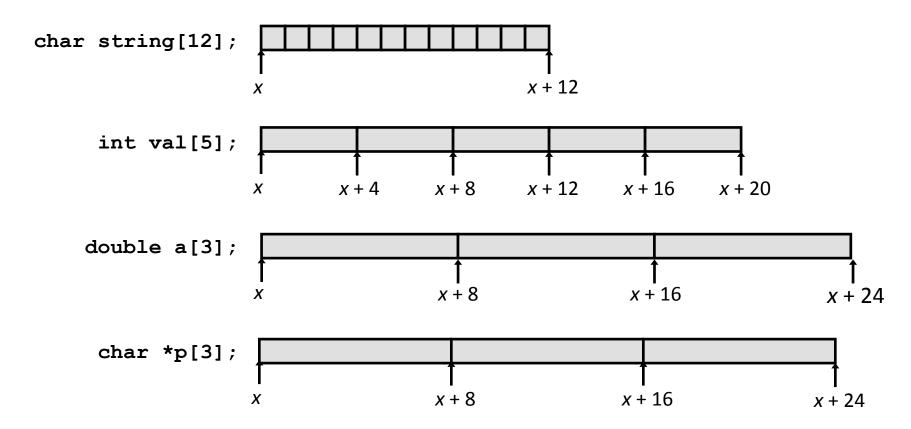
Floating Point

Array Allocation

Basic Principle

 $T \mathbf{A}[L];$

- Array of data type T and length L
- Contiguously allocated region of L * sizeof (T) bytes in memory

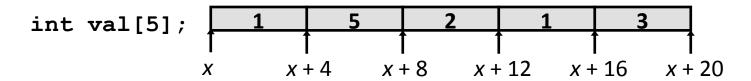


Array Access

■ Basic Principle

```
T \mathbf{A}[L];
```

- Array of data type T and length L
- Identifier **A** can be used as a pointer to array element 0: Type *T**

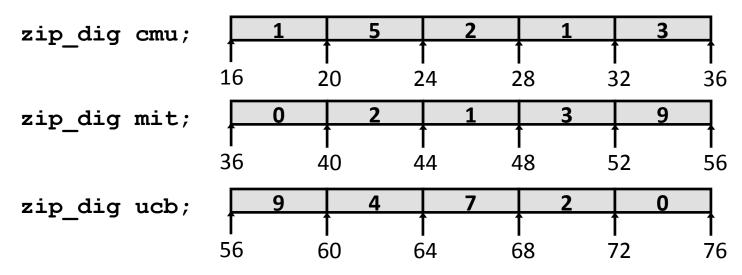


Reference	Type	Value
val[4]	int	3
val	int *	X
val+1	int *	x + 4
&val[2]	int *	<i>x</i> + 8
val [5]	int	5 5
*(val+1)	int	5
val + <i>i</i>	int *	x + 4i

Array Example

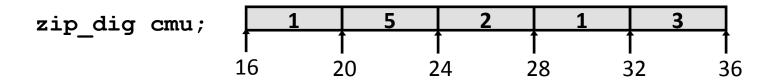
```
#define ZLEN 5
typedef int zip_dig[ZLEN];

zip_dig cmu = { 1, 5, 2, 1, 3 };
zip_dig mit = { 0, 2, 1, 3, 9 };
zip_dig ucb = { 9, 4, 7, 2, 0 };
```



- Declaration "zip_dig cmu" equivalent to "int cmu[5]"
- Example arrays were allocated in successive 20 byte blocks
 - Not guaranteed to happen in general

Array Accessing Example



```
int get_digit
  (zip_dig z, int digit)
{
  return z[digit];
}
```

IA32

```
# %rdi = z
# %rsi = digit
movl (%rdi, %rsi, 4), %eax # z[digit]
```

- Register %rdi contains starting address of array
- Register %rsi contains array index
- Desired digit at %rdi + 4*%rsi
- Use memory reference (%rdi,%rsi,4)

Array Loop Example

```
void zincr(zip_dig z) {
   size_t i;
   for (i = 0; i < ZLEN; i++)
      z[i]++;
}</pre>
```

```
# %rdi = z
 movl $0, %eax
                         \# i = 0
                         # goto middle
 jmp .L3
.L4:
                         # loop:
 addl $1, (%rdi,%rax,4) # z[i]++
 addq $1, %rax
                         # i++
.L3:
                         # middle
                         # i:4
 cmpq $4, %rax
                         # if <=, goto loop</pre>
 jbe
     .L4
 rep; ret
```

Multidimensional (Nested) Arrays

Declaration

 $T \mathbf{A}[R][C];$

- 2D array of data type T
- R rows, C columns
- Type T element requires K bytes

Array Size

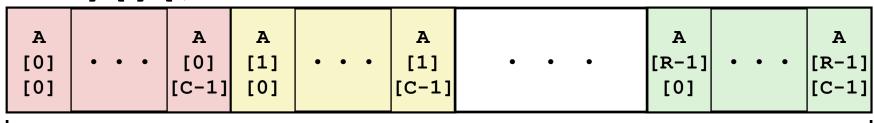
• *R* * *C* * *K* bytes

Arrangement

Row-Major Ordering

A[0][0] • • • A[0][C-1] • • • A[R-1][0] • • • A[R-1][C-1]

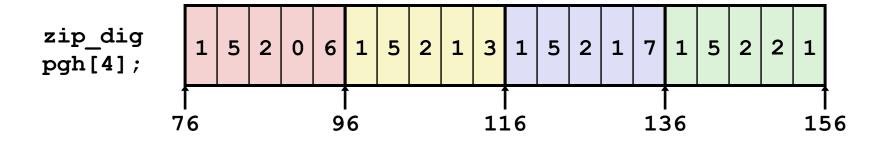
int A[R][C];



4*R*C Bytes

Nested Array Example

```
#define PCOUNT 4
zip_dig pgh[PCOUNT] =
   {{1, 5, 2, 0, 6},
    {1, 5, 2, 1, 3},
   {1, 5, 2, 1, 7},
   {1, 5, 2, 2, 1 }};
```



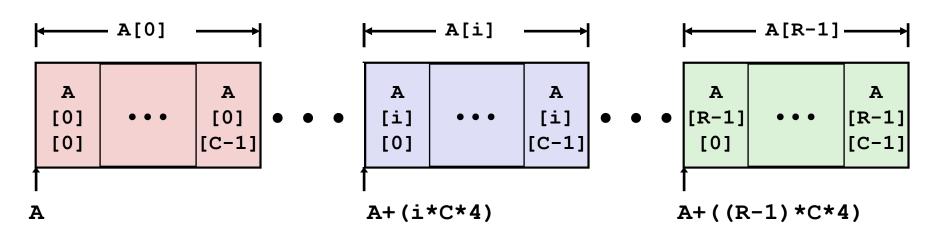
- "zip_dig pgh[4]" equivalent to "int pgh[4][5]"
 - Variable pgh: array of 4 elements, allocated contiguously
 - Each element is an array of 5 int's, allocated contiguously
- "Row-Major" ordering of all elements in memory

Nested Array Row Access

Row Vectors

- A[i] is array of C elements
- Each element of type T requires K bytes
- Starting address A + i * (C * K)

int A[R][C];



Nested Array Row Access Code

```
1 5 2 0 6 1 5 2 1 3 1 5 2 1 7 1 5 2 2 1

pgh

int *get_pgh_zip(int index)
{
    return pgh[index];
}
```

```
# %rdi = index
leaq (%rdi,%rdi,4),%rax # 5 * index
leaq pgh(,%rax,4),%rax # pgh + (20 * index)
```

Row Vector

- pgh[index] is array of 5 int's
- Starting address pgh+20*index

Machine Code

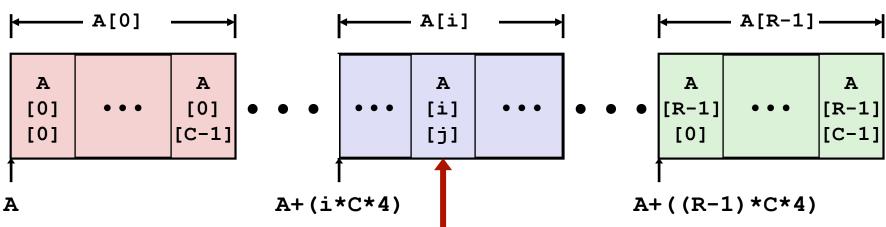
- Computes and returns address
- Compute as pgh + 4*(index+4*index)

Nested Array Element Access

Array Elements

- **A**[i][j] is element of type *T*, which requires *K* bytes
- Address **A** + i * (C * K) + j * K = A + (i * C + j) * K





$$A+(i*C*4)+(j*4)$$

Nested Array Element Access Code

```
leaq (%rdi,%rdi,4), %rax # 5*index
addl %rax, %rsi # 5*index+dig
movl pgh(,%rsi,4), %eax # M[pgh + 4*(5*index+dig)]
```

Array Elements

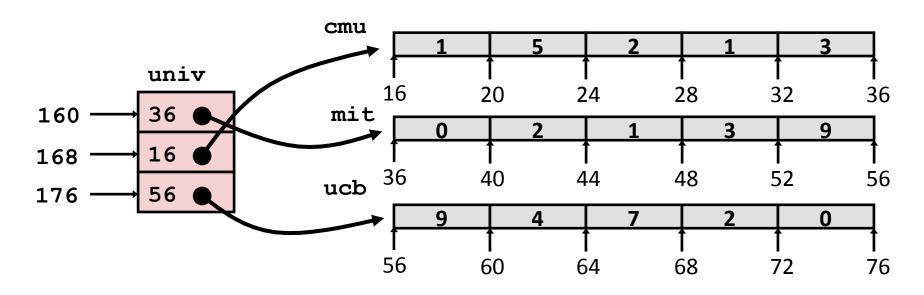
- pgh[index][dig] is int
- Address: pgh + 20*index + 4*dig
 - = pgh + 4*(5*index + dig)

Multi-Level Array Example

```
zip_dig cmu = { 1, 5, 2, 1, 3 };
zip_dig mit = { 0, 2, 1, 3, 9 };
zip_dig ucb = { 9, 4, 7, 2, 0 };
```

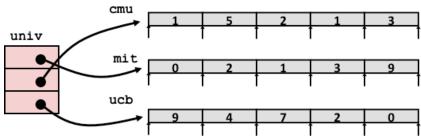
```
#define UCOUNT 3
int *univ[UCOUNT] = {mit, cmu, ucb};
```

- Variable univ denotes array of 3 elements
- Each element is a pointer
 - 8 bytes
- Each pointer points to array of int's



Element Access in Multi-Level Array

```
int get_univ_digit
  (size_t index, size_t digit)
{
  return univ[index][digit];
}
```



```
salq $2, %rsi  # 4*digit
addq univ(,%rdi,8), %rsi # p = univ[index] + 4*digit
movl (%rsi), %eax  # return *p
ret
```

Computation

- Element access Mem [Mem [univ+8*index]+4*digit]
- Must do two memory reads
 - First get pointer to row array
 - Then access element within array

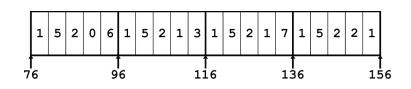
Array Element Accesses

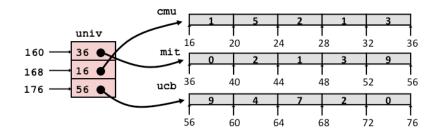
Nested array

```
int get_pgh_digit
  (size_t index, size_t digit)
{
  return pgh[index][digit];
}
```

Multi-level array

```
int get_univ_digit
   (size_t index, size_t digit)
{
   return univ[index][digit];
}
```





Accesses looks similar in C, but address computations very different:

Mem[pgh+20*index+4*digit] Mem[Mem[univ+8*index]+4*digit]

N X N Matrix Code

- Fixed dimensions
 - Know value of N at compile time
- Variable dimensions, explicit indexing
 - Traditional way to implement dynamic arrays
- Variable dimensions, implicit indexing
 - Now supported by gcc

16 X 16 Matrix Access

Array Elements

- Address **A** + i * (C * K) + j * K
- C = 16, K = 4

```
/* Get element a[i][j] */
int fix_ele(fix_matrix a, size_t i, size_t j) {
  return a[i][j];
}
```

```
# a in %rdi, i in %rsi, j in %rdx
salq $6, %rsi # 64*i 16.44
addq %rsi, %rdi # a + 64*i
movl (%rdi,%rdx,4), %eax # M[a + 64*i + 4*j]
ret
```

n X n Matrix Access

Array Elements

- Address **A** + i * (C * K) + j * K
- C = n, K = 4
- Must perform integer multiplication

```
/* Get element a[i][j] */
int var_ele(size_t n, int a[n][n], size_t i, size_t j)
{
  return a[i][j];
}
```

```
# n in %rdi, a in %rsi, i in %rdx, j in %rcx
imulq %rdx, %rdi  # n*i
leaq (%rsi,%rdi,4), %rax # a + 4*n*i
movl (%rax,%rcx,4), %eax # a + 4*n*i + 4*j
ret
```

Today

Arrays

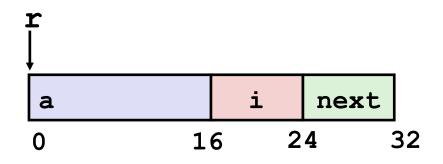
- One-dimensional
- Multi-dimensional (nested)
- Multi-level

Structures

- Allocation
- Access
- Alignment
- Floating Point

Structure Representation

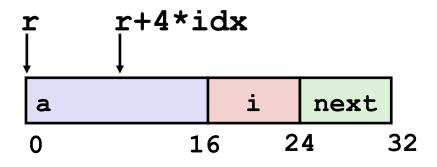
```
struct rec {
   int a[4];
   size_t i;
   struct rec *next;
};
```



- Structure represented as block of memory
 - Big enough to hold all of the fields
- Fields ordered according to declaration
 - Even if another ordering could yield a more compact representation
- Compiler determines overall size + positions of fields
 - Machine-level program has no understanding of the structures in the source code

Generating Pointer to Structure Member

```
struct rec {
   int a[4];
   size_t i;
   struct rec *next;
};
```



Generating Pointer to Array Element

- Offset of each structure member determined at compile time
- Compute as r + 4*idx

```
int *get_ap
  (struct rec *r, size_t idx)
{
   return &r->a[idx];
}
```

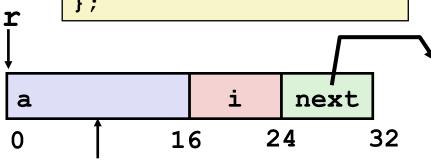
```
# r in %rdi, idx in %rsi
leaq (%rdi,%rsi,4), %rax
ret
```

Following Linked List

C Code

```
void set_val
  (struct rec *r, int val)
{
  while (r) {
    int i = r->i;
    r->a[i] = val;
    r = r->next;
  }
}
```

```
struct rec {
   int a[4];
   int i;
   struct rec *next;
};
```

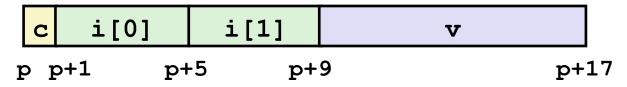


Element i

Register	Value
%rdi	r
%rsi	val

Structures & Alignment

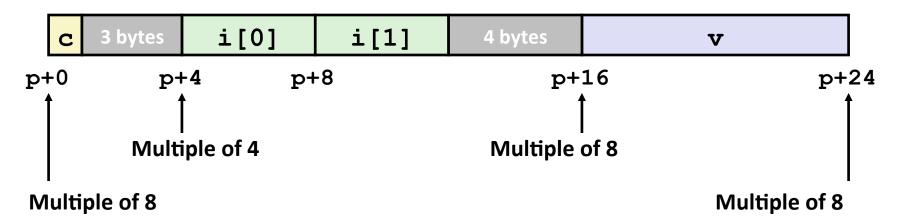
Unaligned Data



```
struct S1 {
  char c;
  int i[2];
  double v;
} *p;
```

Aligned Data

- Primitive data type requires K bytes
- Address must be multiple of K



Alignment Principles

Aligned Data

- Primitive data type requires K bytes
- Address must be multiple of K
- Required on some machines; advised on x86-64

Motivation for Aligning Data

- Memory accessed by (aligned) chunks of 4 or 8 bytes (system dependent)
 - Inefficient to load or store datum that spans quad word boundaries
 - Virtual memory trickier when datum spans 2 pages

Compiler

Inserts gaps in structure to ensure correct alignment of fields

Specific Cases of Alignment (x86-64)

- 1 byte: char, ...
 - no restrictions on address
- 2 bytes: short, ...
 - lowest 1 bit of address must be 02
- 4 bytes: int, float, ...
 - lowest 2 bits of address must be 002
- 8 bytes: double, long, char *, ...
 - lowest 3 bits of address must be 0002
- 16 bytes: long double (GCC on Linux)
 - lowest 4 bits of address must be 0000?

Satisfying Alignment with Structures

Within structure:

• Must satisfy each element's alignment requirement

Overall structure placement

- Each structure has alignment requirement K
 - K = Largest alignment of any element
- Initial address & structure length must be multiples of K

Example:

K = 8, due to double element

```
        c
        3 bytes
        i [0]
        i [1]
        4 bytes
        v

        p+0
        p+4
        p+8
        p+16
        p+24

        Multiple of 4
        Multiple of 8
        Multiple of 8

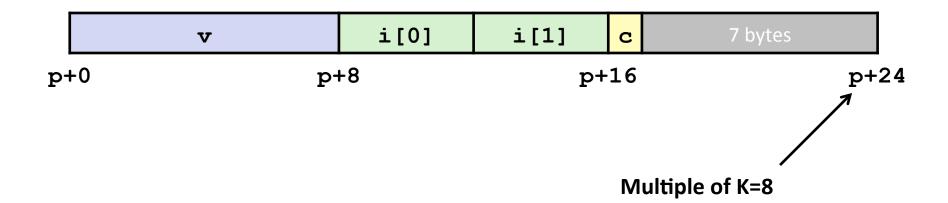
Multiple of 8
```

```
struct S1 {
  char c;
  int i[2];
  double v;
} *p;
```

Meeting Overall Alignment Requirement

- For largest alignment requirement K
- Overall structure must be multiple of K

```
struct S2 {
  double v;
  int i[2];
  char c;
} *p;
```

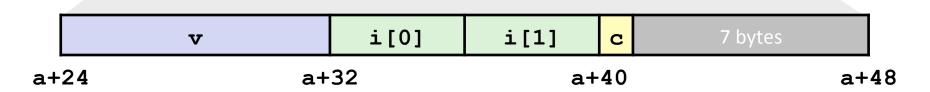


Arrays of Structures

- Overall structure length multiple of K
- Satisfy alignment requirement for every element

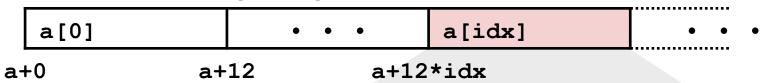
```
struct S2 {
  double v;
  int i[2];
  char c;
} a[10];
```





Accessing Array Elements

- Compute array offset 12*idx
 - sizeof(S3), including alignment spacers
- Element j is at offset 8 within structure
- Assembler gives offset a+8
 - Resolved during linking





```
short get_j(int idx)
{
  return a[idx].j;
}
```

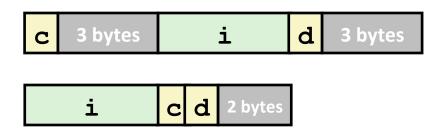
```
# %rdi = idx
leaq (%rdi,%rdi,2),%rax # 3*idx
movzwl a+8(,%rax,4),%eax
```

Saving Space

Put large data types first

```
struct S4 {
  char c;
  int i;
  char d;
} *p;
struct S5 {
  int i;
  char c;
  char d;
} *p;
```

■ Effect (K=4)



Today

Arrays

- One-dimensional
- Multi-dimensional (nested)
- Multi-level

Structures

- Allocation
- Access
- Alignment

Floating Point

Background

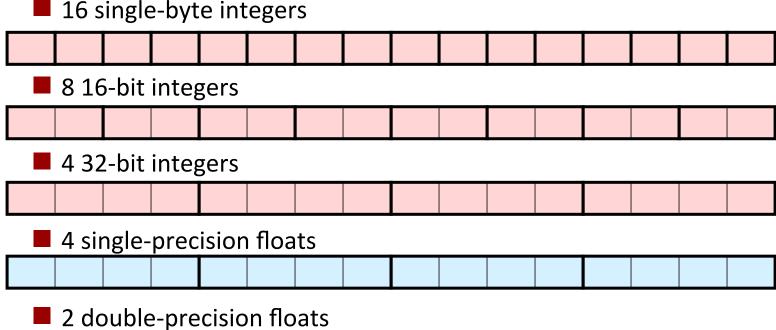
History

- x87 FP
 - Legacy, very ugly
- SSE FP
 - Supported by Shark machines
 - Special case use of vector instructions
- AVX FP
 - Newest version
 - Similar to SSE
 - Documented in book

Programming with SSE3

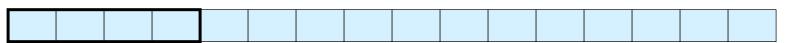
XMM Registers

- 16 total, each 16 bytes
- 16 single-byte integers



2 double-precision floats





1 double-precision float



Scalar & SIMD Operations

Scalar Operations: Single Precision addss %xmm0,%xmm1 %xmm0 %xmm1 SIMD Operations: Single Precision addps %xmm0,%xmm1 %xmm0 %xmm1 Scalar Operations: Double Precision addsd %xmm0,%xmm1 %xmm0 %xmm1 35

FP Basics

- Arguments passed in %xmm0, %xmm1, ...
- Result returned in %xmm0
- All XMM registers caller-saved

```
float fadd(float x, float y)
{
    return x + y;
}
```

```
double dadd(double x, double y)
{
    return x + y;
}
```

```
# x in %xmm0, y in %xmm1
addss %xmm1, %xmm0
ret
```

```
# x in %xmm0, y in %xmm1
addsd %xmm1, %xmm0
ret
```

FP Memory Referencing

- Integer (and pointer) arguments passed in regular registers
- **■** FP values passed in XMM registers
- Different mov instructions to move between XMM registers, and between memory and XMM registers

```
double dincr(double *p, double v)
{
    double x = *p;
    *p = x + v;
    return x;
}
```

```
# p in %rdi, v in %xmm0
movapd %xmm0, %xmm1  # Copy v
movsd (%rdi), %xmm0  # x = *p
addsd %xmm0, %xmm1  # t = x + v
movsd %xmm1, (%rdi) # *p = t
ret
```

Other Aspects of FP Code

Lots of instructions

Different operations, different formats, ...

Floating-point comparisons

- Instructions ucomiss and ucomisd
- Set condition codes CF, ZF, and PF

Using constant values

- Set XMM0 register to 0 with instruction xorpd %xmm0, %xmm0
- Others loaded from memory

Summary

Arrays

- Elements packed into contiguous region of memory
- Use index arithmetic to locate individual elements

Structures

- Elements packed into single region of memory
- Access using offsets determined by compiler
- Possible require internal and external padding to ensure alignment

Combinations

Can nest structure and array code arbitrarily

Floating Point

Data held and operated on in XMM registers

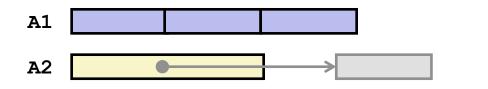
Decl	A <i>n</i>				*An	
	Cmp	Bad	Size	Cmp	Bad	Size
int A1[3]						
int *A2						

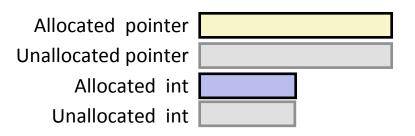
Cmp: Compiles (Y/N)

Bad: Possible bad pointer reference (Y/N)

Size: Value returned by sizeof

Decl		An			*An	
	Cmp	Bad	Size	Cmp	Bad	Size
int A1[3]	Y	N	12	Y	N	4
int *A2	Y	N	8	Y	Y	4





Cmp: Compiles (Y/N)

Bad: Possible bad pointer reference (Y/N)

■ Size: Value returned by sizeof

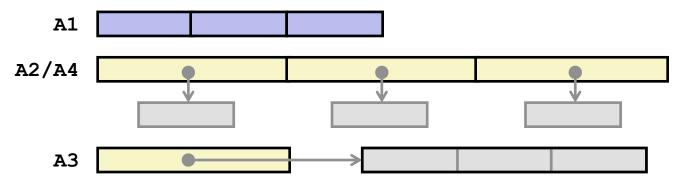
Decl		An			*An		**An		
	Cmp	Bad	Size	Cmp	Bad	Size	Cmp	Bad	Size
int A1[3]									
int *A2[3]									
int (*A3)[3]									
int (*A4[3])									

Cmp: Compiles (Y/N)

Bad: Possible bad pointer reference (Y/N)

Size: Value returned by sizeof

Decl		An *An **An				*A <i>n</i>			
	Cmp	Bad	Size	Cmp	Bad	Size	Cmp	Bad	Size
int A1[3]	Y	N	12	Y	N	4	N	-	-
int *A2[3]	Y	N	24	Y	N	8	Y	Y	4
int (*A3)[3]	Y	N	8	Y	Y	12	Y	Y	4
int (*A4[3])	Y	N	24	Y	N	8	Y	Y	4

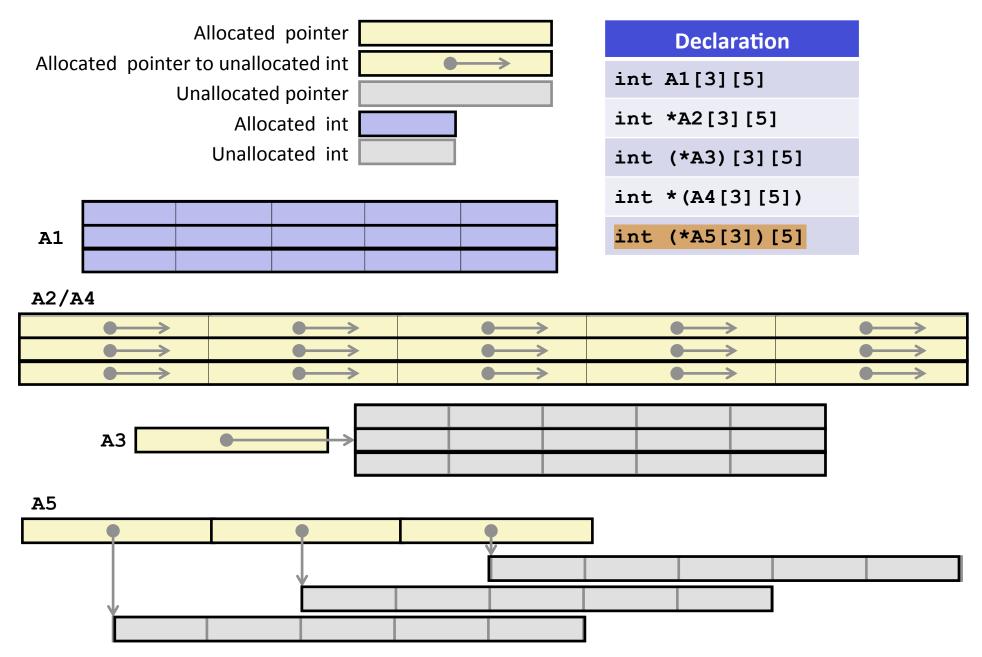


Allocated pointer
Unallocated pointer
Allocated int
Unallocated int

Decl	A <i>n</i>		*An			**An			
	Cmp	Bad	Size	Cmp	Bad	Size	Cmp	Bad	Size
int A1[3][5]									
int *A2[3][5]									
int (*A3)[3][5]									
int *(A4[3][5])									
int (*A5[3])[5]									

- Cmp: Compiles (Y/N)
- Bad: Possible bad pointer reference (Y/N)
- Size: Value returned by sizeof

Decl	***An					
	Cmp	Bad	Size			
int A1[3][5]						
int *A2[3][5]						
int (*A3)[3][5]						
int *(A4[3][5])						
int (*A5[3])[5]						



Decl	A <i>n</i>			*An			**An		
	Cmp	Bad	Size	Cmp	Bad	Size	Cmp	Bad	Size
int A1[3][5]	Y	N	60	Y	N	20	Y	N	4
int *A2[3][5]	Y	N	120	Y	N	40	Y	N	8
int (*A3)[3][5]	Y	N	8	Y	Y	60	Y	Y	20
int *(A4[3][5])	Y	N	120	Y	N	40	Y	N	8
int (*A5[3])[5]	Y	N	24	Y	N	8	Y	Y	20

- Cmp: Compiles (Y/N)
- Bad: Possible bad pointer reference (Y/N)
- Size: Value returned by sizeof

Decl	***An				
	Cmp Bad				
int A1[3][5]	N	-	-		
int *A2[3][5]	Y	Y	4		
int (*A3)[3][5]	Y	Y	4		
int *(A4[3][5])	Y	Y	4		
int (*A5[3])[5]	Y	Y	4		