# Classical IPC Problems

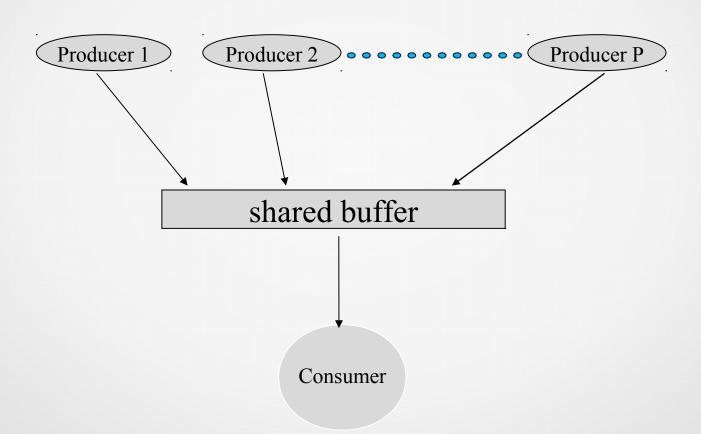
Computer Operating Systems
BLG 312E

2016-2017 Spring

#### **Problems**

- producer consumer
- readers writers
- dining philosophers
- sleeping barber

#### Producer - Consumer



#### Producer - Consumer

- access to shared buffer through mutual exclusion
- circular buffer
- if buffer empty → consumer waits (synchronization)

#### Producer – Consumer

- use counting semaphores
  - takes on ≥ 0 integers
  - used when resource capacity > 1
  - initial value = initial free resource capacity
  - P: one more unit of capacity in use
  - V: one unit of capacity freed

#### Producer – Consumer

- shared buffer implemented through a shared array of size N
  - array[N]
- binary semaphore:
  - mutex ← 1
- counting semaphores:
  - full ← 0 : number of full buffer locations
  - empty ← N : number of free **buffer locations**

```
int in=0, out=0;
          item data;
process producer(){
  while (true) {
    -- produce data -
    p(empty);
    P(mutex);
      array[in]=data;
      in=(in+1) %N;
    v(mutex);
    v(full);
```

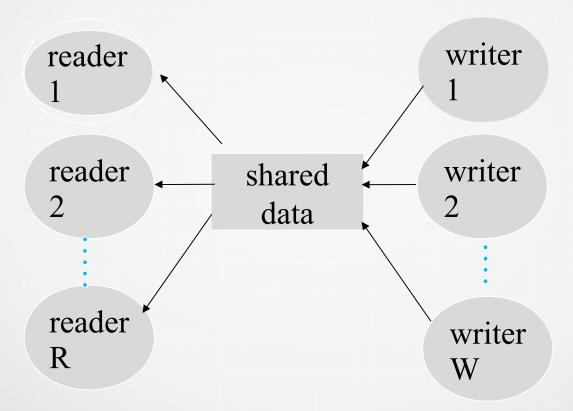
constant N=100;

item array[N];

semaphore full=0, empty=N, mutex=1;

```
process consumer(){
  while (true) {
    p(full);
    p(mutex);
      data=array[out];
      out=(out+1) %N;
    v(mutex);
    v (empty);
    -- use data --
```

### Readers - Writers



#### Readers - Writers

- more than one reader may read shared data (no writers)
- when a writer uses shared data, all other writers and readers must be excluded

```
process reader() {
  while (true) {
    p(read mutex);
      reader=reader+1;
      if (reader==1)
        p(data mutex);
    v(read mutex);
    -- read data --
    p(read mutex);
      reader=reader-1;
      if (reader==0)
        v(data mutex);
    v(read mutex);
```

int reader=0;

semaphore read mutex=1, data mutex=1;

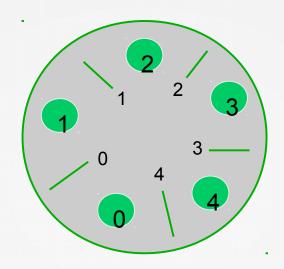
```
process writer() {
  while (true) {
    -- execute --
    p(data_mutex);
    -- write data --
    v(data_mutex);
  }
}
```

readers have priority over writers!

Possible indefinite postponement!

#### Readers - Writers

- must find a fair solution
- apply rules for access order:
  - if a writer is waiting for readers to be finished, do not allow any more readers
  - if a reader is waiting for a writer to finish, give reader priority



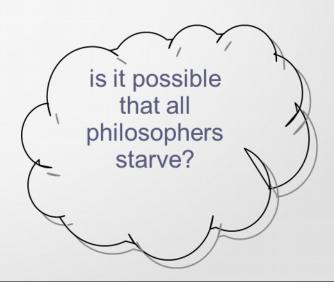
**Problem:** share resources (forks) among philosophers without causing deadlock or starvation

- philosophers
  - eat pasta
  - think
- philosophers need two forks to eat

- fact: two philosophers sitting side by side cannot eat at the same time
  - e.g. for N=5, at most 2 philosophers can eat at the same time
- solution must provide maximum amount of parallelism

```
philosopher(i) {
  while (true) {
                                         what happens if
                                          all philosophers
    think();
                                           take their left
    take fork(i); //left fork
                                             forks?
    take fork((i+4)%5);//right fork
     -- eat
    leave fork(i);
    leave fork ((i+4)\%5);
```

```
philosopher(i) {
  while (true) {
    think();
    take fork(i); //left fork
    if (fork free((i+4)%5)==FALSE)
       leave fork(i);
    else {
      take fork((i+4)\%5);//right fork
      --- eat ----
      leave fork(i);
      leave fork ((i+4)\%5);
```



```
philosopher(i) {
  while (true) {
    P(mutex); //binary semaphore
       think();
       take fork(i); //left fork
       take fork((i+4)%5);//right fork
       --- eat ----
       leave fork(i);
                                             at most how
       leave fork ((i+4)\%5);
                                                many
                                            philosophers can
    V(mutex);
                                             eat together?
```

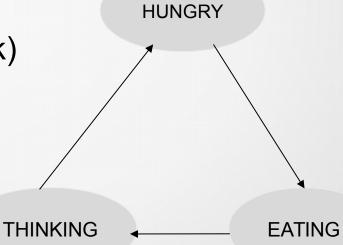
## Dining Philosophers (Correct Solution)

state[i] : state of ith philosopher

- 0: THINKING

- 1: HUNGRY (wait for fork)

- 2 : EATING



## Dining Philosophers (Correct Solution)

- a philosopher can be "EATING" only if both neighbors are not "EATING"
- use a binary semaphore per philosopher
  - blocks on semaphore if a fork is not available when requested

#### Variables:

• semaphores:

- N=5 philosophers
  states:
   THINKING = 0
   HUNGRY = 1
   EATING = 2
- state[5]: array of size 5
- $mutex \leftarrow 1$   $s[5] \leftarrow 0 \text{ array of size 5}$

```
process philosopher(i) {
  while (true) {
    think();
    take fork(i);
    --- eat ---
    leave fork(i);
take fork(i) {
  P(mutex);
    state[i] = HUNGRY; //request to eat
    try[i]; //try to take forks
  V(mutex);
  P(S[i]); //blocks if can't take forks
```

```
leave fork(i) {
  left=(i+1) %5;
  right=(i+4) %5;
 P(mutex);
    state[i]=THINK;
    try(left);
    try(right);
 V(mutex);
try(i) {
  left=(i+1) %5;
 right=(i+4) %5;
  if ((state[i]=HUNGRY) ^
      (state[left]≠EATING) ^
      (state[right]≠EATING))
    state[i]=EATING;
    v(s[i]);
```

### Sleeping Barber

- in a barber shop
  - 1 barber
  - 1 customer seat
  - N waiting seats
- barber sleeps if there are no customers
- arriving customer wakes barber up
- if barber is busy when customer arrives
  - waits if waiting seats available
  - leaves if no waiting seats available

### Sleeping Barber

- 3 semaphores needed for the solution
  - customers : number of customers waiting (excluding the one in the customer seat)
  - barbers : number of available barbers (0/1 in this problem)
  - mutex : for mutual exclusion

```
constant CHAIRS=5;
int waiting=0;
semaphore customers=0,barber=0,mutex=1;
```

```
process barber() {
   while(true) {
      P(customers); //sleep if no customers
      P(mutex);
      waiting--; //remove customer
      V(barber); //barber ready to cut hair
      V(mutex);
      -- cut hair -
   }
}
```

```
process customer() {
   P(mutex);
   if (waiting<CHAIRS) { //shop full?</pre>
      waiting=++;//admite customer
      V (customers); //wake-up barber (possibly)
      V(mutex);
      P (barber); //sleep if barber busy
      -- cut hair -
   else
      V (mutex); //shop is full, so leave
```