

# BLG336E - Analysis of Algorithms II

## 2017-2018 Spring, Project 1 Report

Submission Deadline: 18.03.2018, 23:59

1) Present your problem formulation, state and action representations in detail.

I used a special well-known method which is called **bit masking** to represent the states with only numbers. In this method the bits of the numbers are used to store the state, and it provides mainly two advantages: reducing the used memory, and making easier to control states.

In this problem, there are just 16 different states which are specified below, and we can express them with numbers from 0 to 15. Let's consider all these numbers in **binary form**, they have 4 bits and each bit is responsible for one character (Farmer, Fox, Rabbit, Carrot respectively). If a bit is 0, it means the corresponding character is at the east side, otherwise it is at the west side.

Number	Binary Representation	State
0	0000	Farmer Fox Rabbit Carrot
1	0001	Farmer Fox Rabbit    Carrot
2	0010	Farmer Fox Carrot    Rabbit
3	0011	Farmer Fox    Rabbit Carrot
4	0100	Farmer Rabbit Carrot    Fox
5	0101	Farmer Rabbit    Fox Carrot
6	0110	Farmer Carrot    Fox Rabbit
7	0111	Farmer    Fox Rabbit Carrot
8	1000	Fox Rabbit Carrot    Farmer
9	1001	Fox Rabbit    Farmer Carrot
10	1010	Fox Carrot    Farmer Rabbit
11	1011	Fox    Farmer Rabbit Carrot
12	1100	Rabbit Carrot    Farmer Fox
13	1101	Rabbit    Farmer Fox Carrot
14	1110	Carrot    Farmer Fox Rabbit
15	1111	Farmer Fox Rabbit Carrot

Table 1. Number-State Transitions  
(Invalid states are red-colored)

According to the number-state transition table, we can easily construct the state diagram: **red** cells represent invalid states, and **green** cells represent valid states. The source cell is colored with **yellow** and sink cell is colored with **blue**. Actually all state actions are bi-directional, however turning back to the same state would not advantageous, hence I just use directional edges.

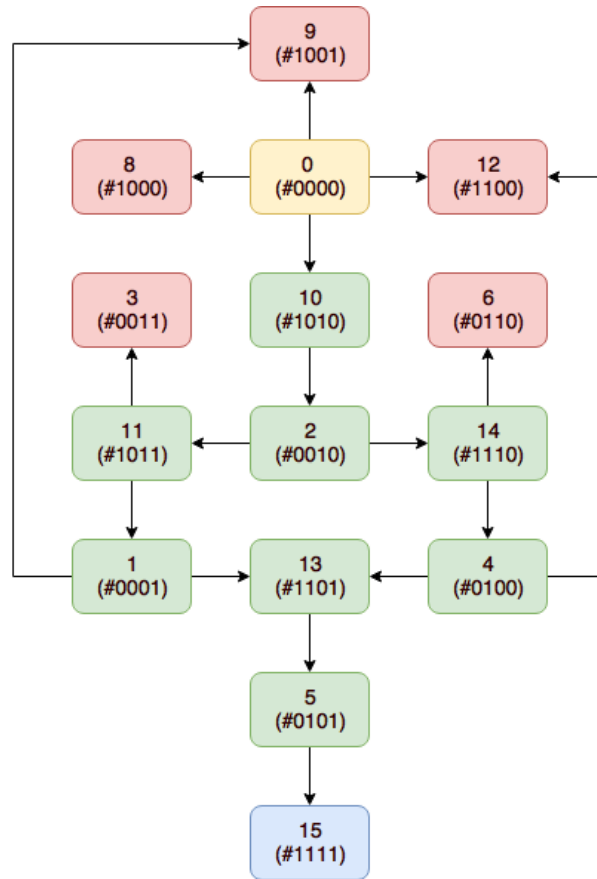


Figure 1. State Diagram

Now all states and actions are determined, and we can just use these transitions instead constructing a graph explicitly.

- 2) How does your algorithm work?
- Write your pseudo-code.
  - Show your complexity of your algorithm on pseudo-code.

<b>function</b> BFS():	O(N)
Q = new Queue() // holds current_state, previous_state pairs	O(1)
add initial pair (0, 0) to Q	O(1)
<b>while</b> Q is not empty:	O(N)
increase number_of_visited_nodes by 1	O(1)
maximum_node_count = max(maximum_node_count, Q.size())	O(1)
current_state, previous_state = Q.pop()	O(1)
mark the current state visited	O(1)
trace[current_state] = previous_state // we can reach the current state from previous state	O(1)
<b>if</b> (current_state is destination (15)):	O(1)
<b>return</b> true	O(1)
<b>for each</b> valid movement from current_state:	O(3)
<b>if</b> next_state is not visited:	O(1)
add (next_state, current_state) pair to Q	O(1)
<b>return</b> false	

<b>function</b> DFS (int state, int previous, int depth)	O(N)
increase number_of_visited_counts by 1	O(1)
max_node_cnt = max(max_node_cnt, depth);	O(1)
trace[state] = previous; // it means to reach the state, first you have to reach "previous" state	O(1)
<b>if</b> state is destination (15) :	O(1)
<b>return</b> true	O(1)
mark the current state visited	O(1)
<b>for each</b> valid movement from current_state:	O(3)
<b>if</b> next_state is not visited <b>and</b> DFS (next_state, current_state, depth+1) is <b>true</b> :	O(N)
<b>return</b> true	O(1)
<b>return</b> false	O(1)

3) In Depth-First Search algorithm, why should we maintain a list of discovered nodes?

Because if the algorithms encounter a node which was discovered already, it causes an infinite loop and it can never move out the loop.

4) Is the graph constructed by the given problem formulation a bipartite graph? Why or why not? (Implementation is not required for this part.)

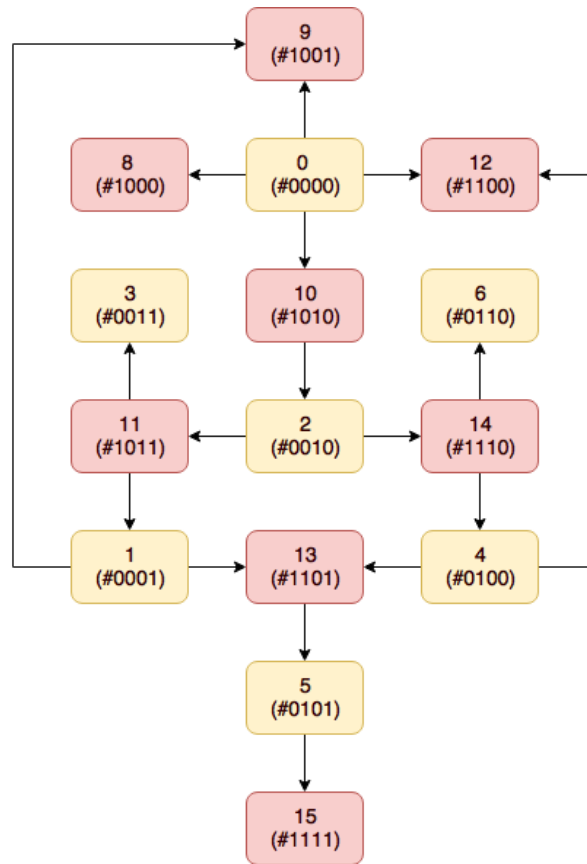


Figure 2. Bipartite Graph

As we can see the state diagram in Figure 1 the graph is a bipartite graph, we can decompose the graph into two disjoint set such that no two graph nodes within the same set are adjacent. Colors represent the set I composed.

5) Analyze and explain the algorithm results in terms of:

- **The number of visited nodes:**

According to the program outputs, Depth-First Search algorithm visits **8** nodes (states) and Breadth-First Search algorithm visits **12** nodes (states).

- **The maximum number of nodes kept in memory:**

In DFS, the maximum number of nodes kept in memory is the maximum depth of the recursion, and the maximum depth is **8** for this problem.

In BFS, the maximum number of nodes kept in memory is the maximum size of queue structure that is reached at any moment, and it is just **2** for given problem.

- The running time: 0.001 s

DFS and BFS algorithms have  $O(E + V)$  time complexity and because  $E$  and  $V$  are both very small (16), the program finds the solution rapidly.

- The number of move steps on the found solution: 7

There is only two solutions in this problem as we can see the state diagram in Figure 1, and both of them has 7 move steps without repetition. DFS and BFS algorithms that I implemented for this problem can find these solutions too.

**Compilation Command:** `g++ kod.cpp -o kod -std=c++11`

**Running Command:** `./kod bfs`

```

[otok@ssh ~]$ g++ kod.cpp -o kod -std=c++11
[otok@ssh ~]$ ./kod dfs
Algorithm: dfs
Number of the visited nodes: 8
Maximum number of nodes kept in the memory: 8
Running time: 0 milliseconds
Solution move count: 7
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[otok@ssh ~]$ ./kod bfs
Algorithm: bfs
Number of the visited nodes: 12
Maximum number of nodes kept in the memory: 2
Running time: 0 milliseconds
Solution move count: 7
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[otok@ssh ~]$

```

Figure 2. The funny output of program with emojis