

1.7: INTRODUCTION TO PROOFS

Basic Definitions Needed for Proofs In The Remainder of This Section:

- (1) Given an integer n .
 - n is *even* if there exists an integer k such that

$$n = 2k.$$
 - n is *odd* if there exists an integer k such that

$$n = 2k+1.$$
- (2) A number x is *rational* if it can be written in the form $x = m/n$, where m and n are integers (with $n \neq 0$). x is *irrational* if it is not a rational number.

I. Direct Proof of $P \rightarrow Q$

Suppose P .

⋮

Therefore Q .

Thus, $P \rightarrow Q$. ■

Example 1: Prove that if n is odd, then $n + 1$ is even.

Your strategy for developing a direct proof of an implication should involve the following steps:

1. Determine precisely the hypothesis (P) and the conclusion (Q).
2. If necessary, replace P (and/or Q) with a more usable, but still equivalent, proposition.
3. Develop a chain of statements, each deducible from its predecessors or other known results, that leads from P to Q .

Example 2: Suppose n is an integer. Prove that n is even if n^2 is even.

Recall that $(P \rightarrow Q) \equiv (\neg Q \rightarrow \neg P)$.

Proof by Contrapositive of $P \rightarrow Q$

Suppose $\neg Q$.

\vdots

Therefore $\neg P$.

Thus, $P \rightarrow Q$. ■

Example 3: Prove that $\sqrt{2}$ is irrational.

Proof by Contradiction of Q	Proof by Contradiction of $P \rightarrow Q$
<p>Suppose $\neg Q$.</p> <p>\vdots</p> <p>Therefore R.</p> <p>\vdots</p> <p>Therefore $\neg R$.</p> <p>Hence, $R \wedge \neg R$, a contradiction.</p> <p>Thus, Q. ■</p>	<p>Suppose P and $\neg Q$.</p> <p>\vdots</p> <p>Therefore $\neg P$.</p> <p>Hence, $P \wedge \neg P$, a contradiction.</p> <p>Thus, Q. ■</p>

Example 4: Let m and n be positive real numbers. Prove if $m \geq n$, then $m^0 \geq n^0$.

Trivial proof: $P \rightarrow Q$ is automatically true whenever Q is always true.

Example 5: Let m and n be positive real numbers. Prove if $mn < 0$, then $m < 0$ or $n < 0$.

Vacuous proof: $P \rightarrow Q$ is automatically true if P is always false.

SECTION 1.7 – INTRODUCTION TO PROOFS

METHODS OF PROVING THEOREMS

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This is a *trivial proof*. $P \rightarrow Q$ is automatically true whenever Q is always true.

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This is a *vacuous proof*. $P \rightarrow Q$ is automatically true if P is always false.