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Base Case

$count\ NilT \le 2^{(height\ NilT)} - 1$	
$0 \le 2^0 - 1$	ctb, htb
$0 \le 0$	

Base case holds

Assume

$$count (Node \ n \ x \ y) \le 2^{\left(height \ (Node \ n \ x \ y)\right)} - 1$$

Prove

$$count \ (Node \ n \ x \ y) \leq 2^{\left(height \ (Node \ n \ x \ y)\right)} - 1$$

$$(count \ x) + (count \ y) + 1 \leq 2^{\left(\max \ (height \ x) \ (height \ y)\right) + 1} - 1 \qquad ctr, htr$$

The following step can be done because it is the maximum case.

$$2^{height\,x} - 1 + 2^{height\,y} - 1 + 1 \le 2^{\left(\max\left(height\,x\right)\left(height\,y\right)\right) + 1} - 1 \qquad IA$$

$$2^{height\,x} + 2^{height\,y} \le 2^{\left(\max\left(height\,x\right)\left(height\,y\right)\right) + 1}$$

Case mc1

$$2^{height x} + 2^{height x} \le 2^{(height x)+1}$$
$$2^{(height x)+1} \le 2^{(height x)+1}$$

Case mc2

$$2^{height y} + 2^{height y} \le 2^{(height y)+1}$$
$$2^{(height y)+1} \le 2^{(height y)+1}$$