PEARSON

ELEVENTH EDITION

Digital Systems

Principles and Applications

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- *Selected areas covered in this chapter*:
  - Terminology associated with memory systems.
  - Difference between read/write memory and read-only memory.
  - Difference between volatile and nonvolatile memory.
  - Capacity of a memory device from inputs & outputs.
  - Steps that occur when the CPU reads from or writes to memory.
  - Various types of ROMs & common applications.
  - Organization/operation of static and dynamic RAMs.
  - Relative advantages/disadvantages of EPROM, EEPROM, and flash memory.
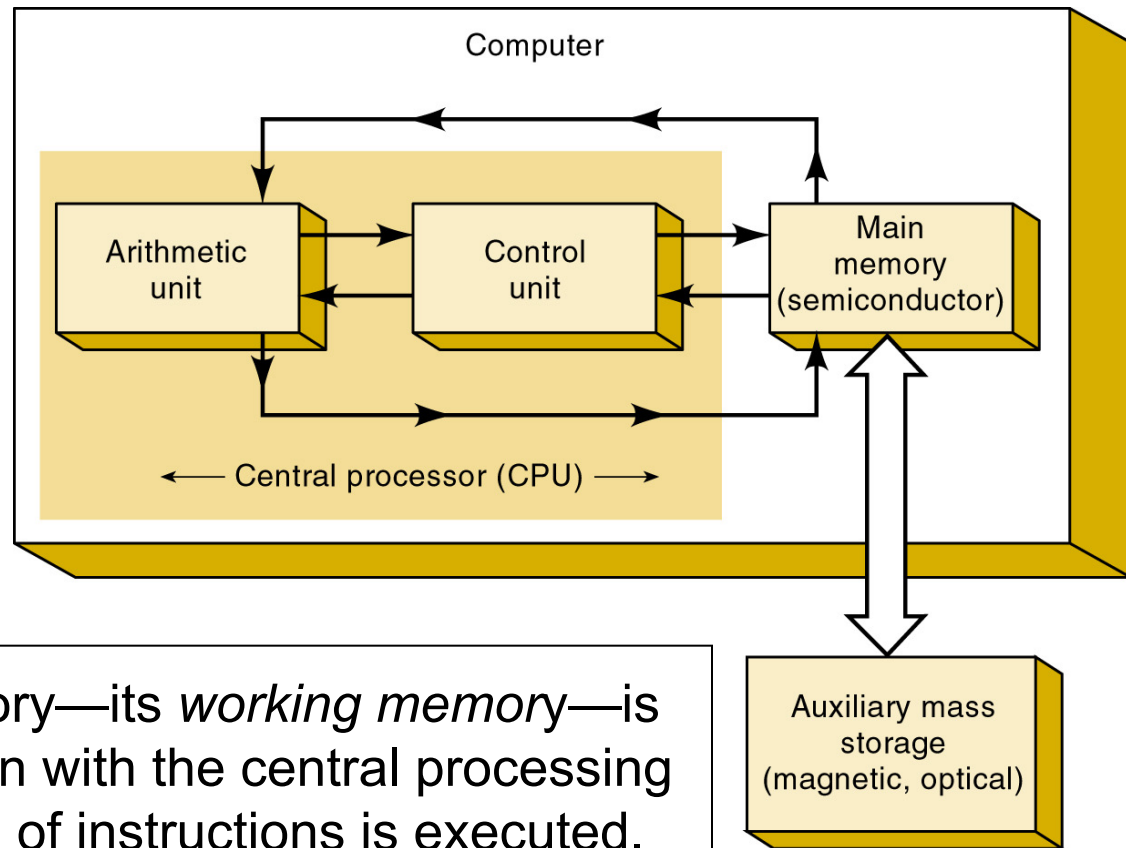 - Using test results to determine memory faults.

- Digital data can be stored as charges on capacitors.
  - An important semiconductor memory type does this for high-density storage, at low power-requirements.

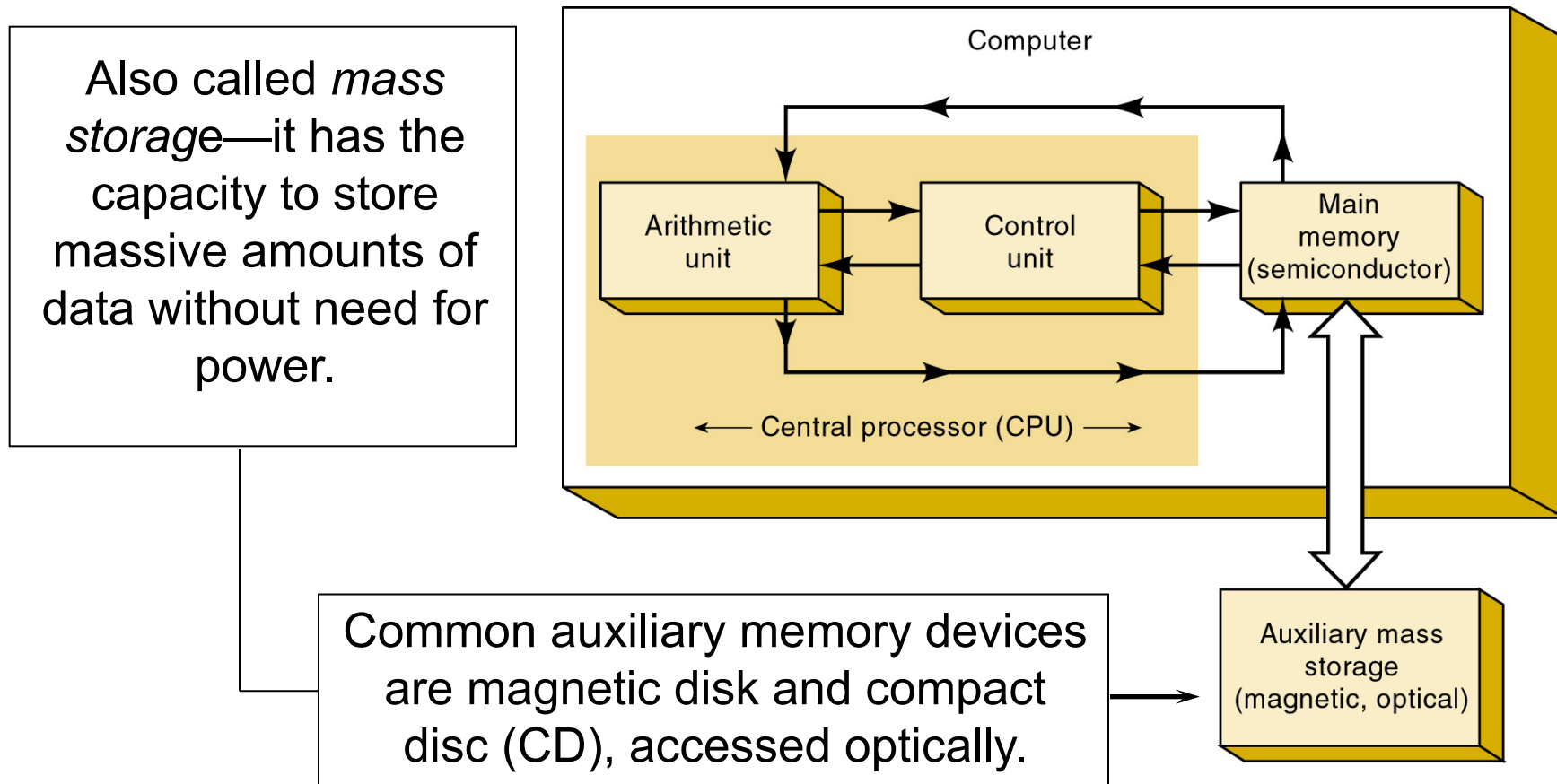- Semiconductor memories are used as the **main memory** of a computer where fast operation is important.

RAM and ROM make up main memory.

Computer

Arithmetic unit

Control unit

Main memory (semiconductor)

← Central processor (CPU) →

Auxiliary mass storage (magnetic, optical)

A computer's main memory—its *working memory*—is in constant communication with the central processing unit (CPU) as a program of instructions is executed.

- ## Another form of storage is performed by **auxiliary memory,** separate from the main memory.

Also called *mass storage*—it has the capacity to store massive amounts of data without need for power.

Common auxiliary memory devices are magnetic disk and compact disc (CD), accessed optically.

- **Memory Cell**—a device or an electrical circuit used to store a single bit (0 or 1).
  - Examples include a flip-flop, charged capacitor, or a single spot on magnetic tape or disk.

- **Memory Word**—a group of bits (cells) in a memory that represents instructions or data.
  - Word sizes in computers typically range from 8 to 64 bits, depending on the size of the computer.

- **Byte**—a special term used for a group of eight bits.
  - Always consists of eight bits.

- **Capacity**—a way of specifying how many bits can be stored in a memory device or system.

- **Density**—another term for *capacity.*
  - A memory device with greater density can store more bits in a given amount of space.

- **Address**—a number that identifies the location of a word in memory.
  - Addresses always exist in a digital system as a binary number, although octal, hex & decimal numbers are often used to represent the address for convenience.

A small memory consisting of eight words.

Each of these eight words has a specific address represented as a three-bit number—from 000 to 111.

To refer to a specific word location in memory, use its address code to identify it.

Addresses	
000	Word 0
001	Word 1
010	Word 2
011	Word 3
100	Word 4
101	Word 5
110	Word 6
111	Word 7

- **Read Operation**—the binary word stored in a specific memory location (address) is sensed and then transferred to another device.
  - Often called a *fetch* operation because a word is being fetched from memory.

- **Write Operation**—operation whereby a new word is placed into a particular memory location.
  - Also referred to as a *store* operation, it replaces the word that was previously stored there.

- **Access Time**—measure of memory device speed, it is the time between the memory receiving a new address input & data is available at the output.

- **Volatile Memory**—any memory that requires the application of electrical power to store information.

  – If the electrical power is removed, all information stored in the memory will be lost.

- **Random-Access Memory (RAM)**—memory in which actual physical location of a memory word has no effect on how long it takes to read from, or write into, that location.

  – Access time is the same for any address in memory.

    • Most semiconductor memories are RAMs.

- **Sequential-Access Memory (SAM)**—type of memory in which the access time is not constant but varies depending on the address location.

  – A stored word is found by sequencing through all address locations until the desired address is reached.

    • Access times far longer than random-access memory.

- **Read/Write Memory (RWM)**—any memory that can be read from or written into with equal ease.

- **Read-Only Memory (ROM)**—broad class of semiconductor memories designed for applications with a high ratio of read- to write- operations.

- **Static Memory Devices**—semiconductor memory devices in which stored data remains permanently stored as long as power is applied.
  - Without need for periodically rewriting data to memory.

- **Dynamic Memory Devices**—semiconductor memory in which stored data will *not* remain permanently stored, even with power applied.
  - Unless the data are periodically rewritten into memory.
    - A *refresh* operation.

- **Main Memory**—a computer's *working memory.*
  - Stores instructions and data the CPU is currently working on.

- **Cache Memory**—high-speed block of memory that operates between slower main memory and the CPU to optimize the speed of the computer.
  - Physically located in the CPU, mother board, or both.

- **Auxiliary Memory**—referred to as *mass storage* because it stores massive amounts of information external to the main memory.
  - Slower than main memory, always nonvolatile.

- Every memory system requires I/O lines to:
  - Apply the binary address of the memory location that is to be accessed.
  - Enable memory devices to respond to control inputs.
  - Place data stored in the specified address.
  - In a read operation, enable the tristate outputs.
    - Which applies the data to the output pins.
  - In a write operation, apply the data to be stored to the data input pins.
  - Enable the write operation, which causes the data to be stored at the specified location.
  - Deactivate the read or write controls when done reading or writing and disable the memory IC.
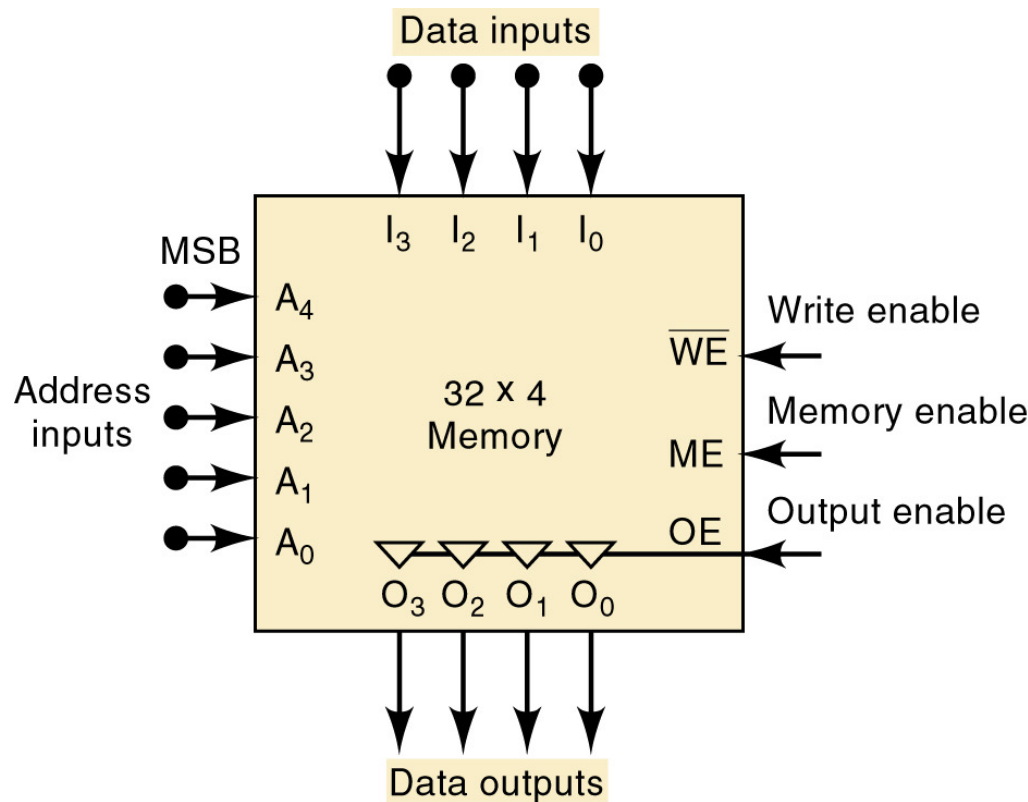
Diagram of a 32 x 4 memory, and the virtual arrangement of memory cells into 32 four-bit words.
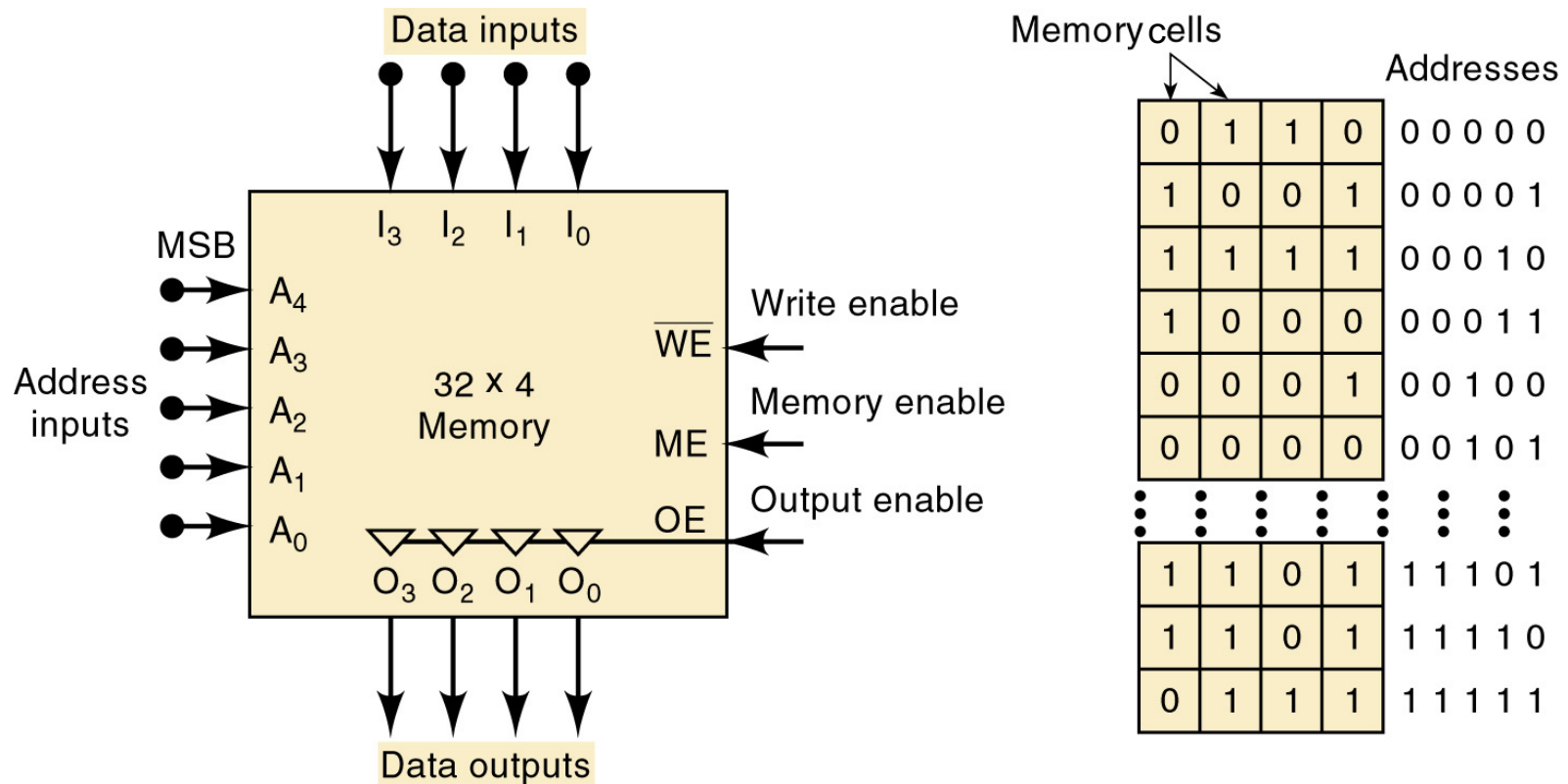
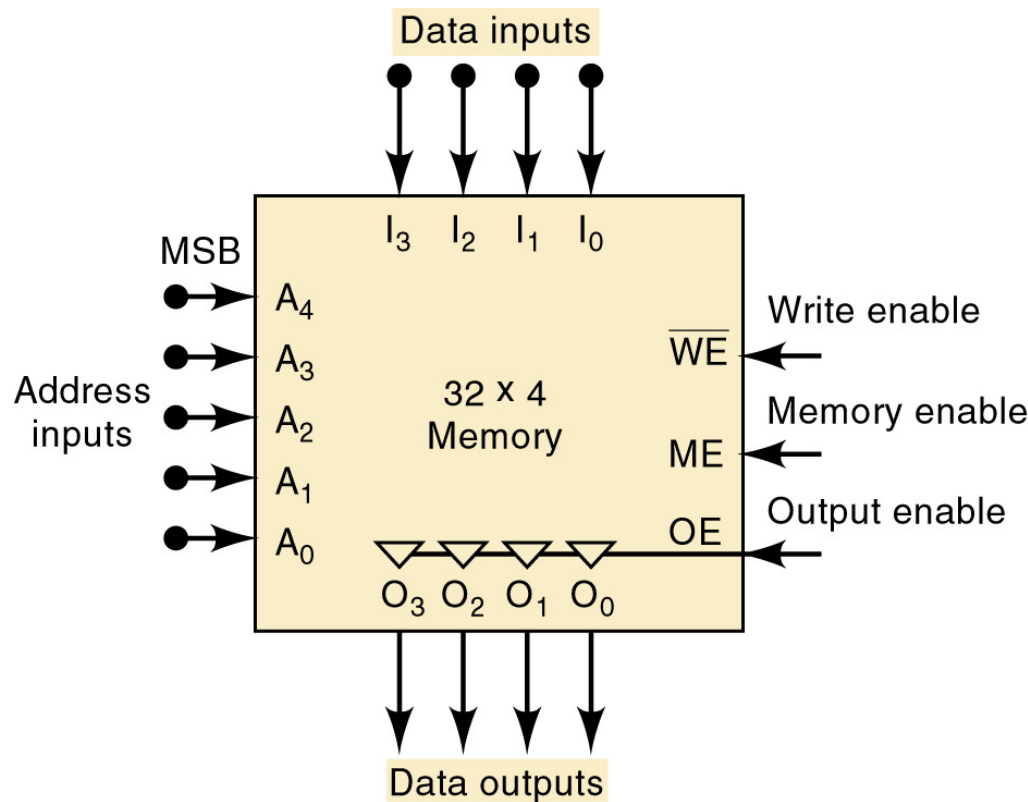## Diagram of a 32 x 4 memory, and the virtual arrangement of memory cells into 32 four-bit words.

Because this memory stores 32 words, it has 32 different storage locations & 32 different binary addresses, from 00000 to 11111 (0 to 31 in decimal).

There are five address inputs—A_0 to A_4.

Diagram of a 32 x 4 memory, and the virtual arrangement of memory cells into 32 four-bit words.
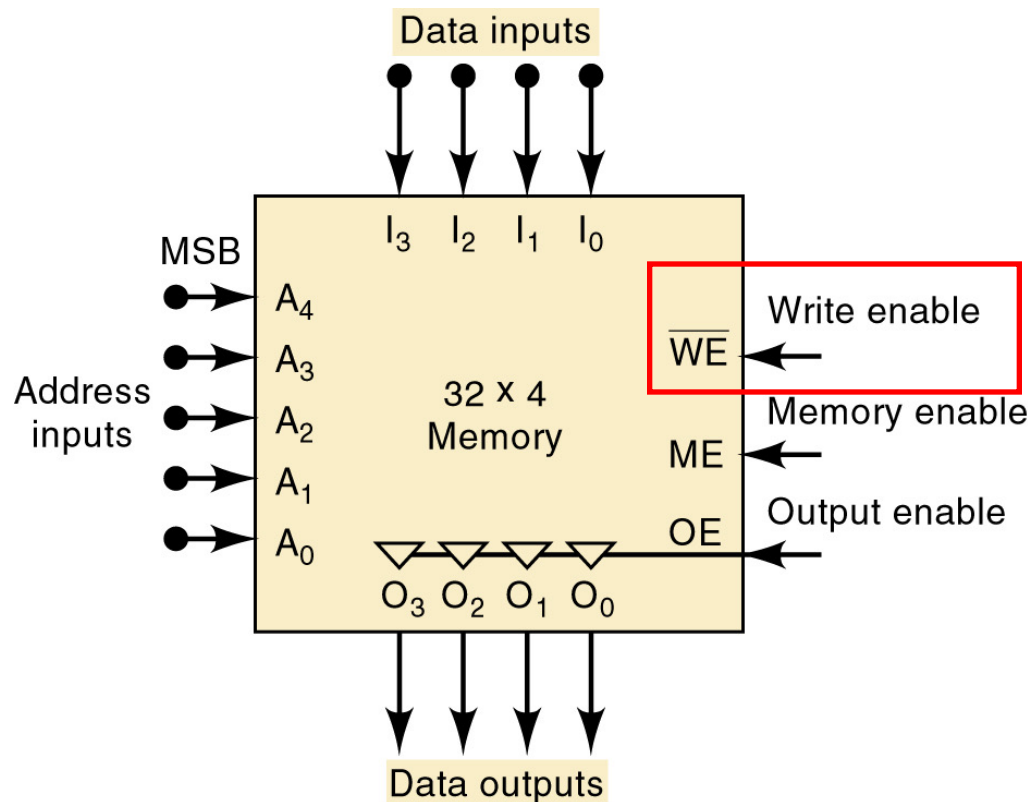

To access one a memory location for read or write, the five-bit address code is applied to the address inputs.

In general, N address inputs are required for a memory with a capacity of 2^N words.

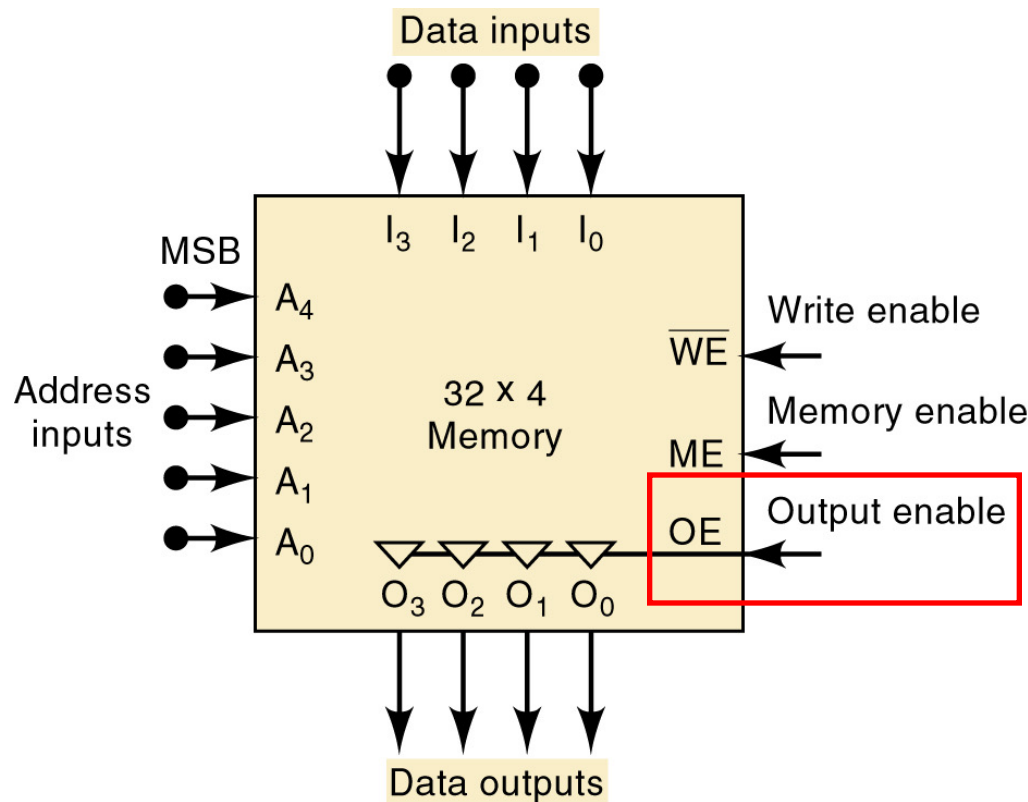## Diagram of a 32 x 4 memory, and the virtual arrangement of memory cells into 32 four-bit words.

The $\overline{WE}$ (write enable) input is activated to allow the memory to store data.

The overbar indicates that the write operation takes place when $\overline{WE} = 0$.

Diagram of a 32 x 4 memory, and the virtual arrangement of memory cells into 32 four-bit words.
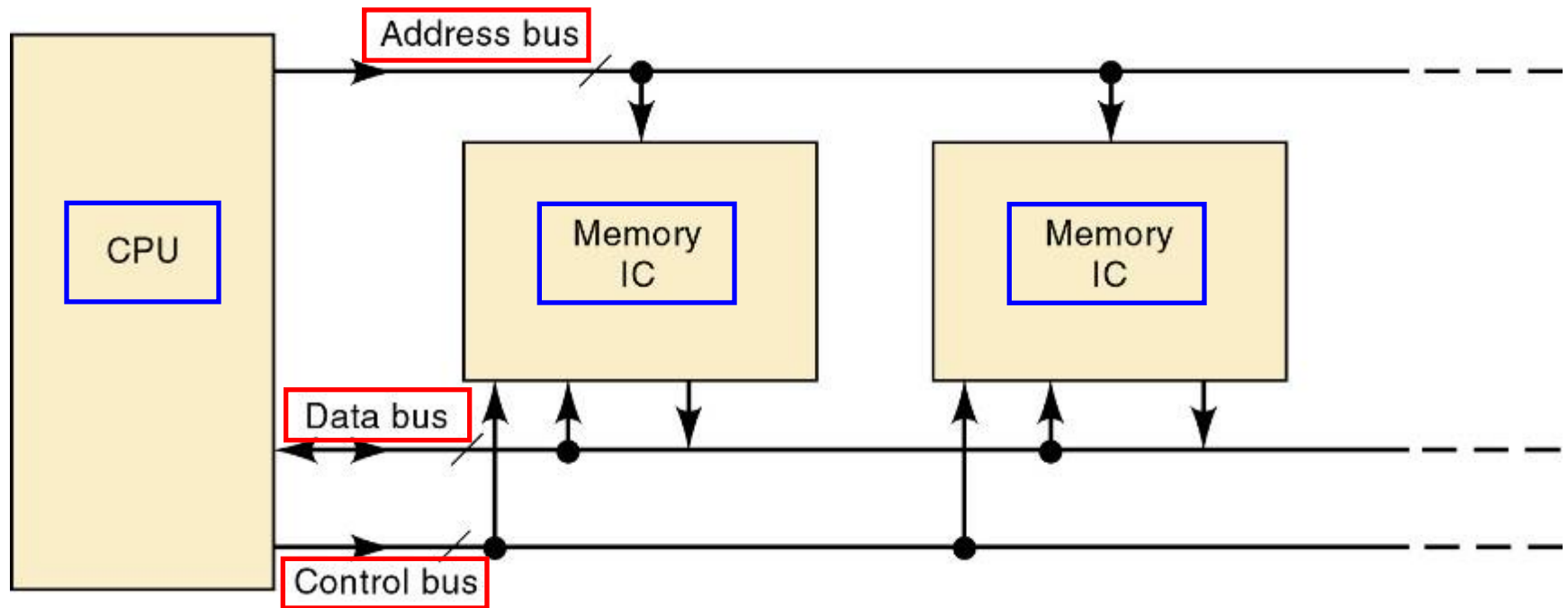

The *OE* pin is activated to enable the tristate buffer and deactivated to place the buffers in the high impedance (hi-Z) state.

A control signal connected to *OE* is active only when the bus is ready to receive data from the memory.

- ## Main memory is interfaced to the CPU through:
  - Address bus; Data bus; Control bus.

The three buses play a necessary part in allowing the CPU to write data into memory and to read data from memory.

- Write operation process:
  - The CPU supplies the binary address of the memory location where the data are to be stored.
    - It places this address on the address bus lines.
  - An address decoder activates the memory device's enable input *(CE or CS)*.
  - CPU places data to be stored on the data bus lines.
  - CPU activates appropriate control signal lines for the memory write operation $\overline{WR}$ *or R/$\overline{W}$.*
  - Memory ICs internally decode the binary address to determine the location selected for the store operation.
  - The data on the data bus are transferred to the selected memory location.

- Read operation process:
  - The CPU supplies the binary address of the memory location from which data are to be retrieved.
    - It places this address on the address bus lines.
  - An address decoder activates the memory device's enable input *(CE* or *CS)*.
  - The CPU activates the appropriate control signal lines for the memory read operation, which is normally connected to on the memory IC.
    Example: $\overline{WR}$ *or* $R/\overline{W}$.
  - Memory ICs internally decode the binary address to determine the location is being selected to read.
    - They place data from the memory location onto the data bus, from which they are transferred to the CPU.

- Function of each of the system buses:
  - **Address Bus**—*unidirectional* bus that carries binary address outputs from the CPU to the memory ICs.
    - To select one memory location.
  - **Data Bus**—*bidirectional* bus that carries data between the CPU and the memory ICs.
  - **Control Bus**—carries control signals ($\overline{RD}$ *or* $\overline{WR}$) from the CPU to the memory ICs.

- Read-only memory is semiconductor memory designed to hold data that are permanent or will not change frequently.

- Some ROMs cannot have their data changed once they have been programmed—others can be *erased* & reprogrammed as often as desired.
  - The process of entering data is called **programming** or *burning* the ROM.

- A major use of ROMs storage of programs in microcomputers.
  - As ROMs are *nonvolatile,* programs are not lost when electrical power is turned off.

Typical ROM
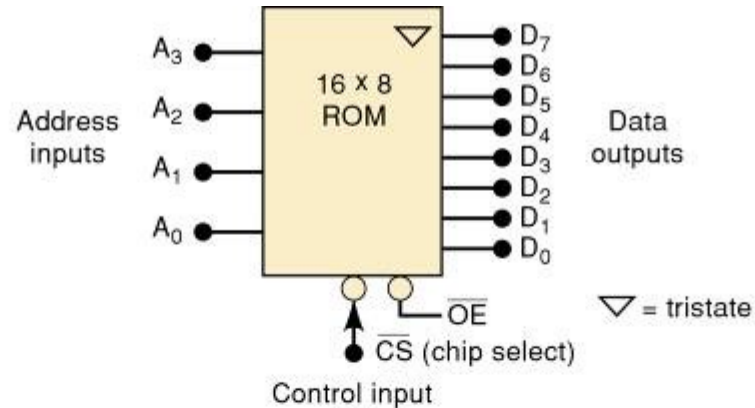
ord	A₃ A₂ A₁ A₀	D₇ D₆ D₅ D₄ D₃ D₂ D₁ D₀
	Address	Data
ord	A_3 A_2 A_1 A_0	D_7 D_6 D_5 D_4 D_3 D_2 D_1 D_0
0	0 0 0 0	1 1 0 1 1 1 1 0
1	0 0 0 1	0 0 1 1 1 0 1 0
2	0 0 1 0	1 0 0 0 0 1 0 1
3	0 0 1 1	1 0 1 0 1 1 1 1
4	0 1 0 0	0 0 0 1 1 0 0 1
5	0 1 0 1	0 1 1 1 1 0 1 1
6	0 1 1 0	0 0 0 0 0 0 0 0
7	0 1 1 1	1 1 1 0 1 1 0 1
8	1 0 0 0	0 0 1 1 1 1 0 0
	1 0 0 1	1 1 1 1 1 1 1 1
10	1 0 1 0	1 0 1 1 1 0 0 0
11	1 0 1 1	1 1 0 0 0 1 1 1
12	1 1 0 0	0 0 1 0 0 1 1 1
13	1 1 0 1	0 1 1 0 1 0 1 0
14	1 1 1 0	1 1 0 1 0 0 1 0
15	1 1 1 1	0 1 0 1 1 0 1 1

| Word	A_3 A_2 A_1 A_0	D_7 D_0
	Address	Data
0	0	DE
1	1	3A
2	2	85
3	3	A
4	4	19
5	5	7B
6	6	00
7	7	ED
8	8	3C
9	9	
10	A	B8
11	B	C7
12	C	27
13	D	6A
14	E	D2
15		5B

- To read a data word from ROM requires:
  - Applying the appropriate address inputs.
  - Activating the control inputs.

- The internal architecture (structure) of a ROM IC is complex—but has four basic parts:
  - ***Register array****—stores data programmed into ROM.*
    - *Each register contains several memory cells equal to the word size.*
  - ***Address decoders****—Row & Column decoders.*
    - Only one register will be in both row & column selected by the address inputs, and this one will be enabled.
  - **Output buffers**—pass data to external data outputs.
    - The register that is enabled by the address inputs will place its data on the data bus.
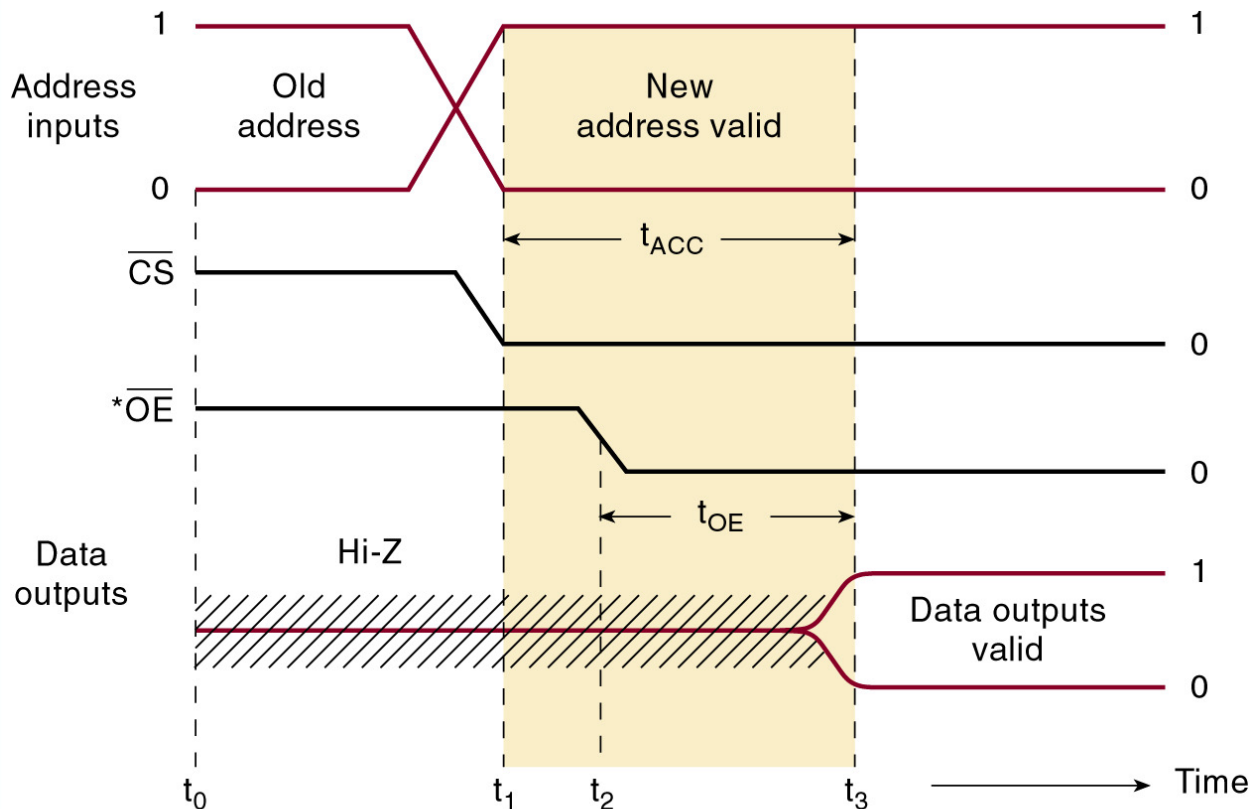   - These data feed into the output buffers, which will pass the data to the external data outputs.

Architecture of a 16 x 8 ROM.
Each register stores one eight-bit word.
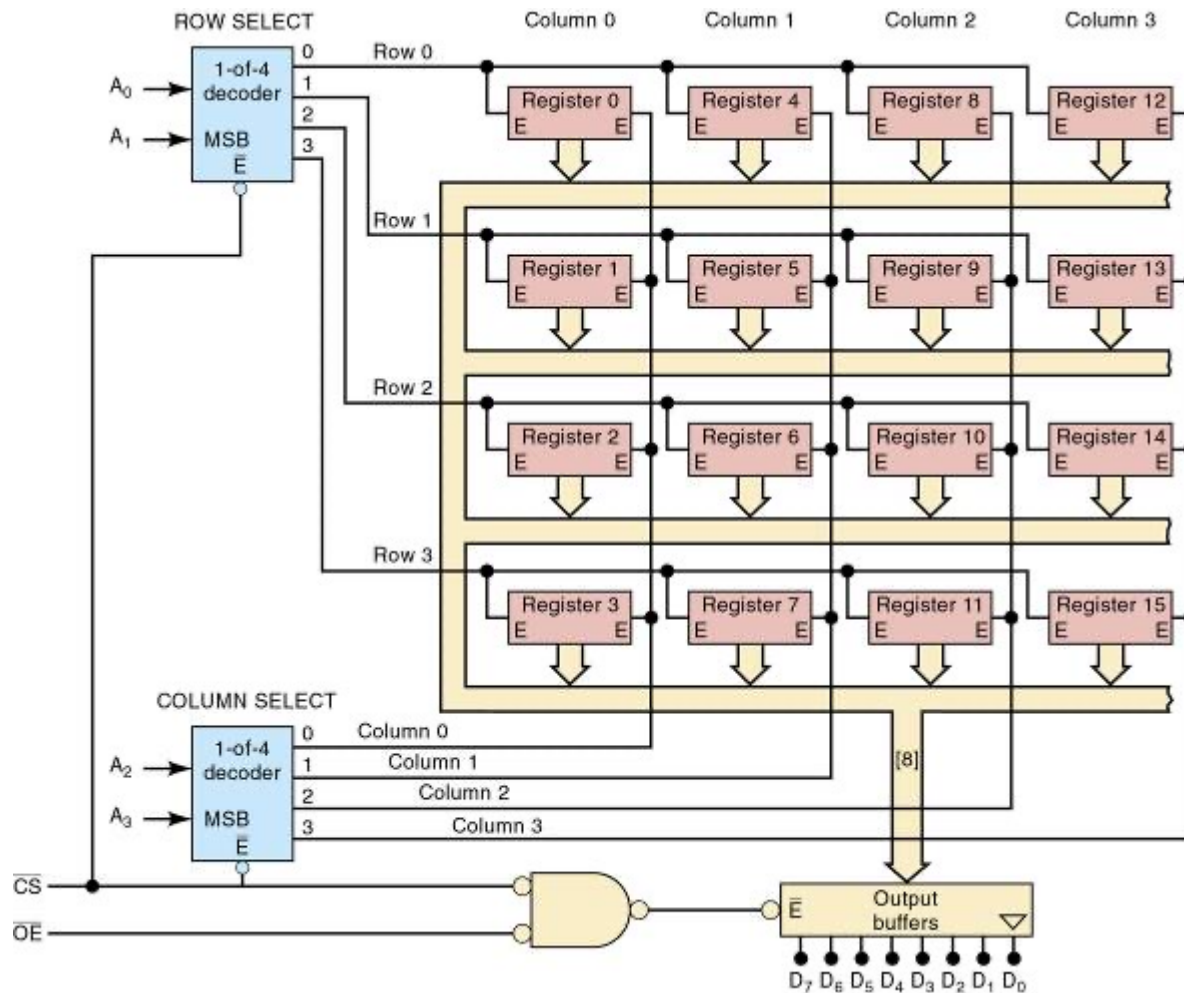
- There will be propagation delay between the application of a ROM's inputs and appearance of the data outputs during a read operation.
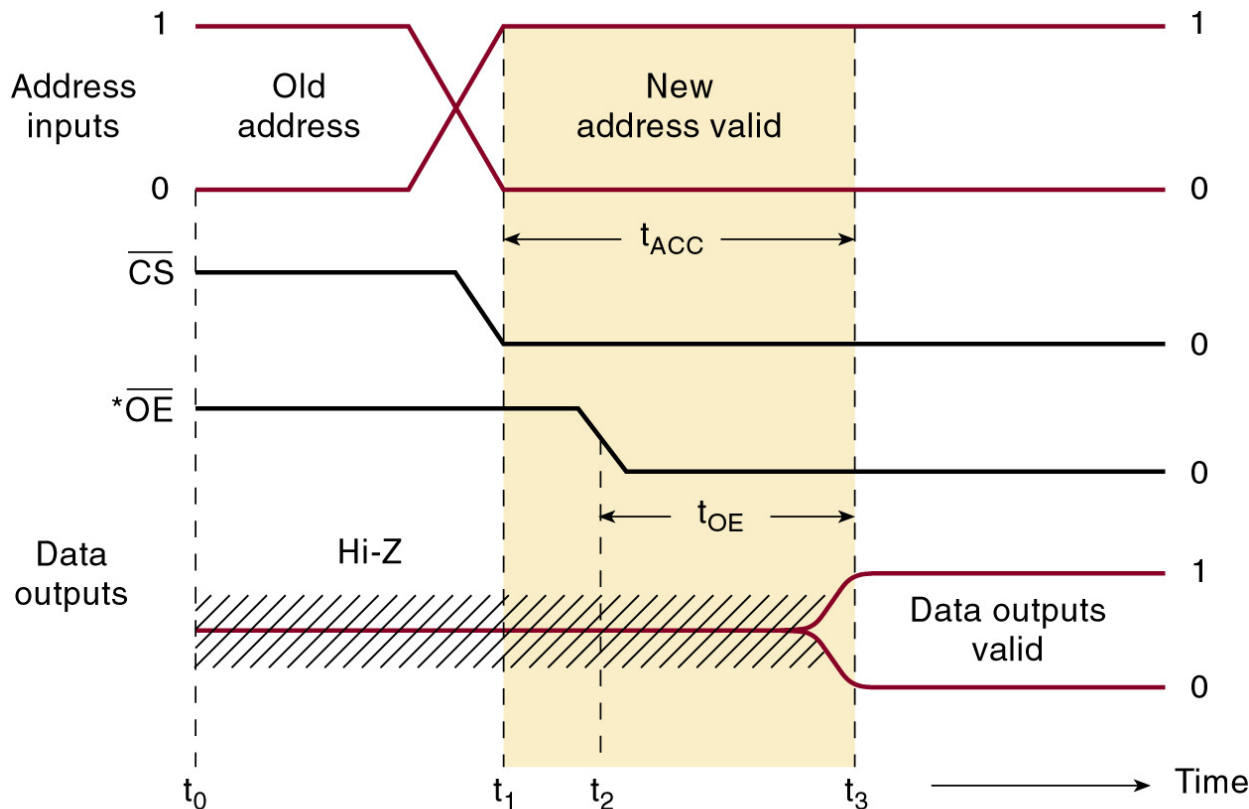

Called access time (t_{ACC}), the delay is a measure of ROM operating speed.

*t_{OE} is measured from the time $\overline{CS}$ and $\overline{OE}$ have both been asserted.

- There will be propagation delay between the application of a ROM's inputs and appearance of the data outputs during a read operation.


Another important timing parameter is *output enable time* (t_{OE}), the delay between input and valid data output.

Values for t_{OE} are always shorter than access time.

*t_{OE} is measured from the time $\overline{CS}$ and $\overline{OE}$ have both been asserted.

- **Mask-programmed ROM** (**MROM**) has data stored at the time the IC is manufactured.

ROMs are made up of a rectangular array of transistors.

Information is stored by either connecting or disconnecting the source of a transistor to the output column.
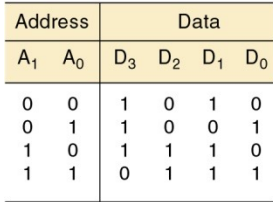
Address		Data			
A_1	A_0	D_3	D_2	D_1	D_0
0	0	1	0	1	0
0	1	1	0	0	1
1	0	1	1	1	0
1	1	0	1	1	1

- **Mask-programmed ROM (MROM)** has data stored at the time the IC is manufactured.

The last step in the manufacturing is to form all these conducting paths or connections.

The process uses a "mask" to deposit metals on the silicon that determine where connections form.
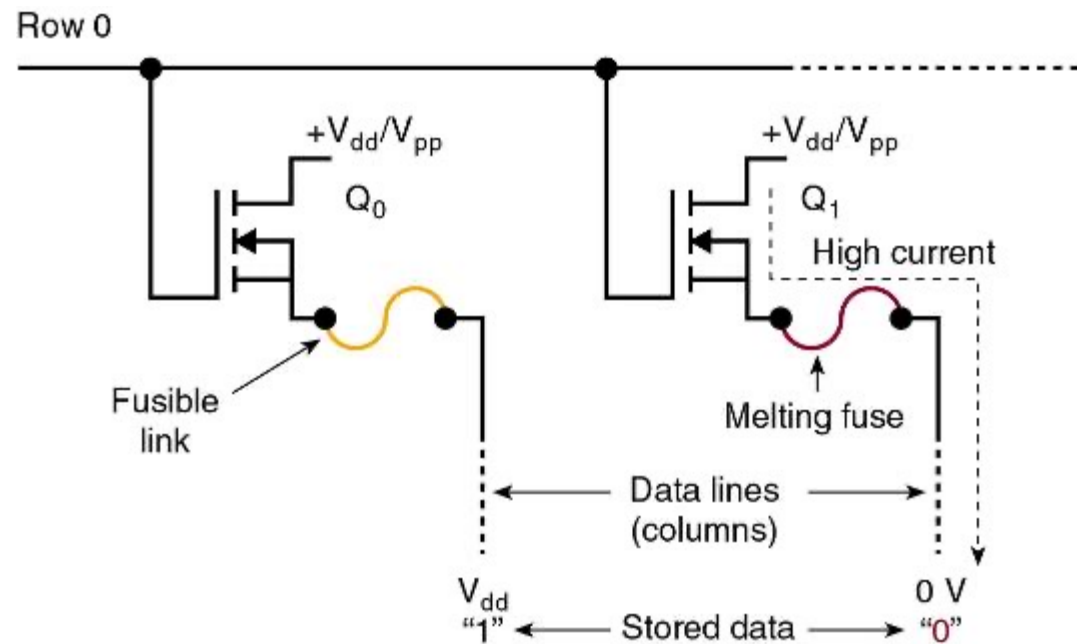

Address		Data			
A_1	A_0	D_3	D_2	D_1	D_0
0	0	1	0	1	0
0	1	1	0	0	1
1	0	1	1	1	0
1	1	0	1	1	1
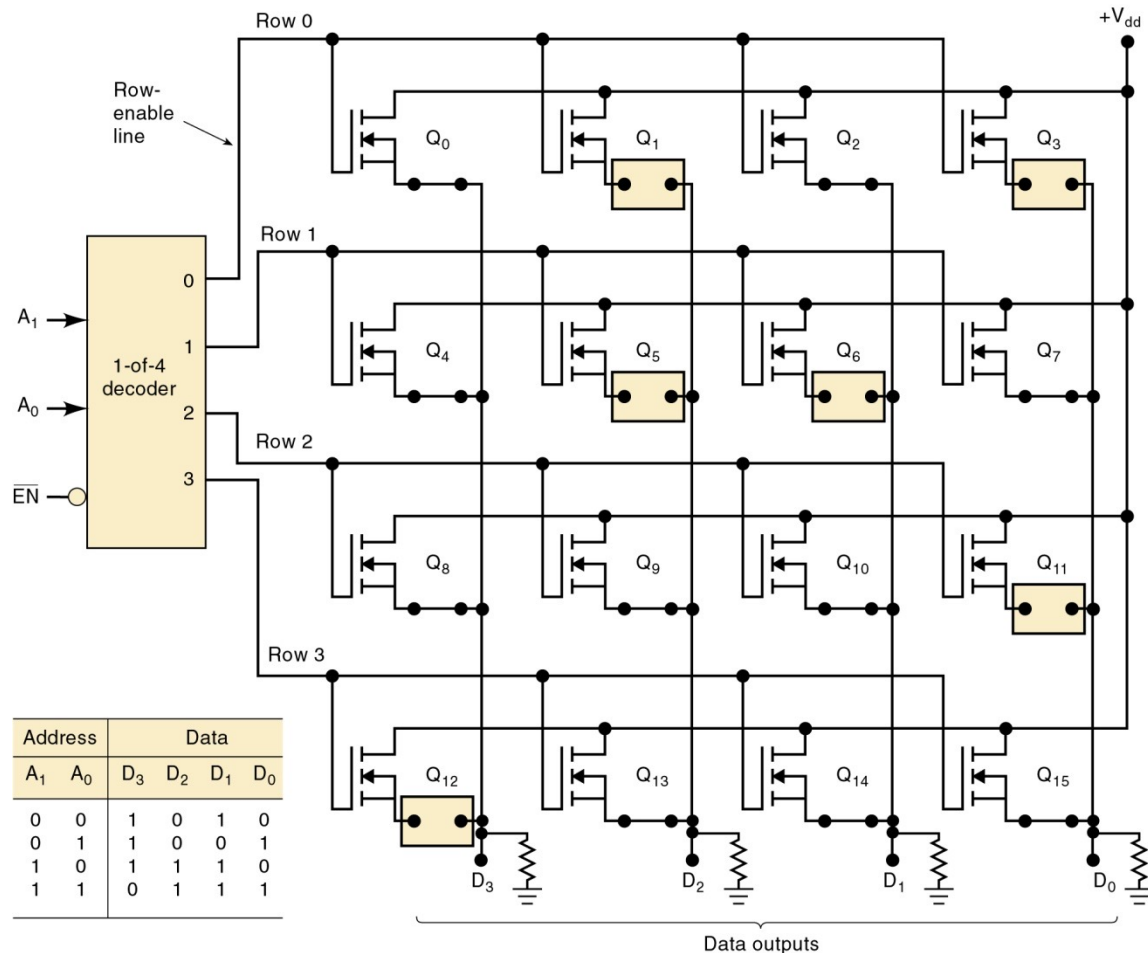
- The mask is very precise, expensive and must be made specifically for the customer, with the correct binary information.
  - Economical only when many ROMs are being made with exactly the same information.

- ## For lower-volume applications, user-programmable **fusible-link PROMs** are available.
   - Custom-programmed by the user, it cannot be erased and reprogrammed.

If information in the PROM is faulty or must be changed, it must be discarded.

Often referred to as "one-time programmable" (OTP) ROMs.

Row 0

$+V_{dd}/V_{pp}$ Q_0

$+V_{dd}/V_{pp}$ Q_1

High current

Fusible link

Melting fuse

Data lines (columns)

V_{dd} "1"   0 V "0"

Stored data

- An EPROM can be user-programmed, *erased* and reprogrammed as often as desired.
  - Once programmed, it is a *nonvolatile* memory that will hold its stored data indefinitely.

A **UV** light is used to clear the device.
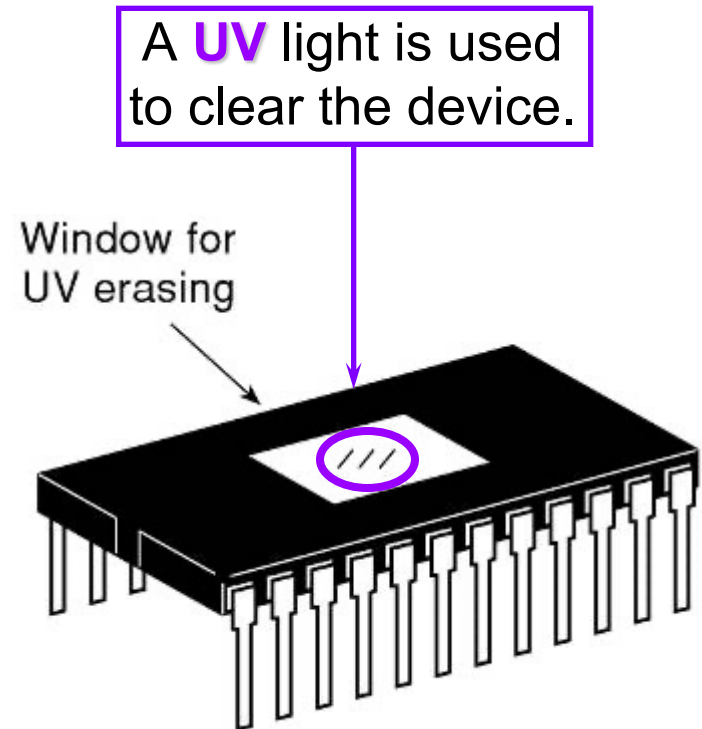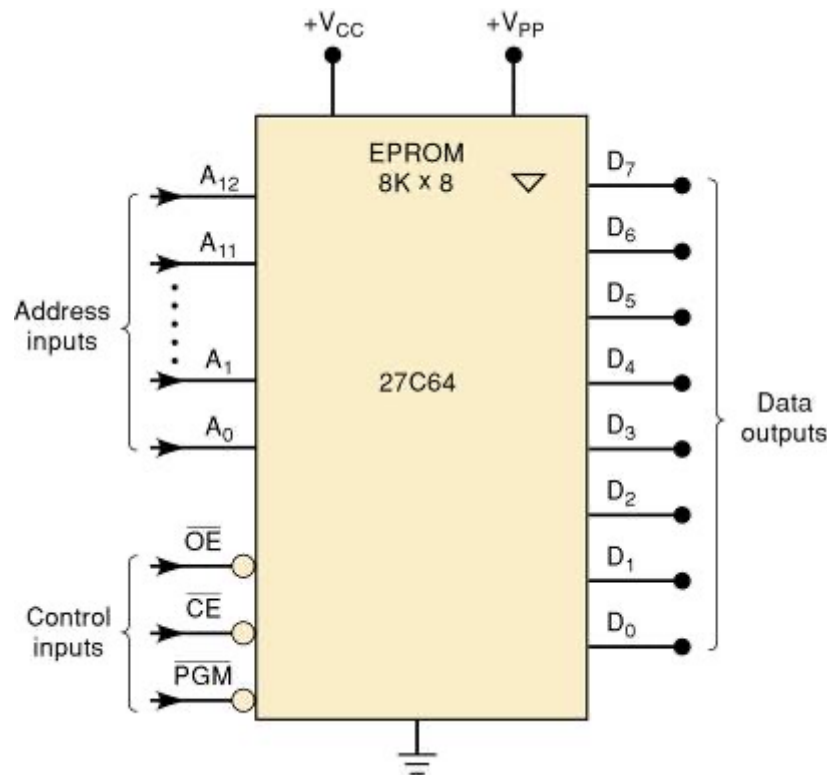
- Major disadvantages of UVEPROMs:
  - They must be removed from the circuit to be programmed and erased.
  - The erase operation erases the entire chip.
  - The erase operation takes up to 20 minutes.

- The major characteristic of **electrically erasable PROM (EEPROM)** is electrical erasability.
  - Also the ability to erase and rewrite *individual* bytes in the memory array.

- Because the internal process of storing a data value in an EEPROM is quite slow, the speed of the data transfer operation can also be slower.
  - EEPROM devices are available in eight-pin packages interfaced to a two- or three-wire *serial* bus.

- A flash memory cell is like the simple single-transistor EPROM cell, with a cost considerably less than for EEPROM.

In-circuit, electrically erasable byte-by-byte

In-circuit, electrically erasable by sector or in bulk (all cells)

UV erasable in bulk; erased and reprogrammed out of circuit

Cannot be erased and reprogrammed

EEPROM

Flash

EPROM

MROM and PROM

Device complexity and cost

- The unique feature of the 28F256A CMOS flash memory IC is the *command register.*
  - Command codes are written into this register to control which operations take place inside the chip.
  - State control logic examines the contents of the command register and generates logic and control signals.
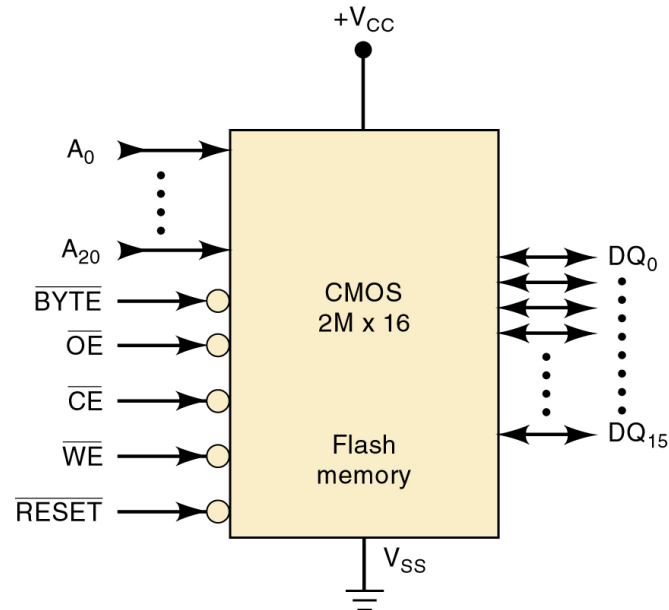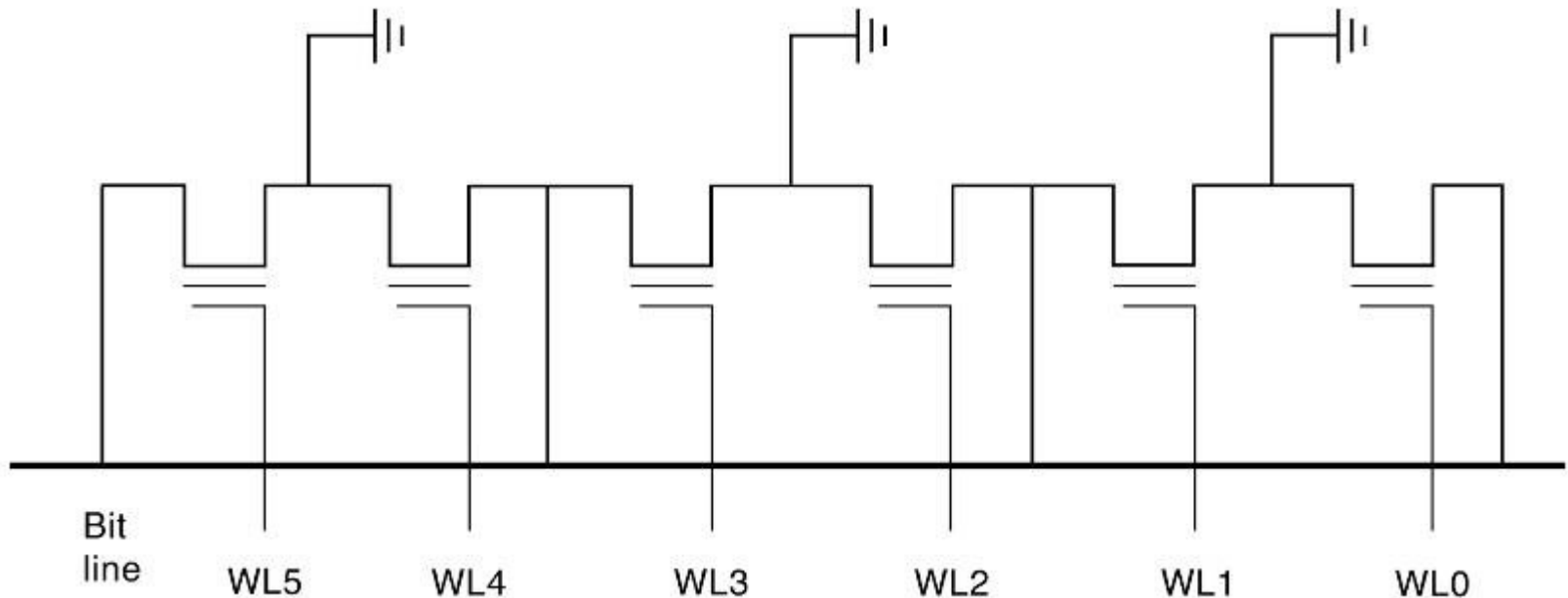
Functional diagram of a flash memory chip.
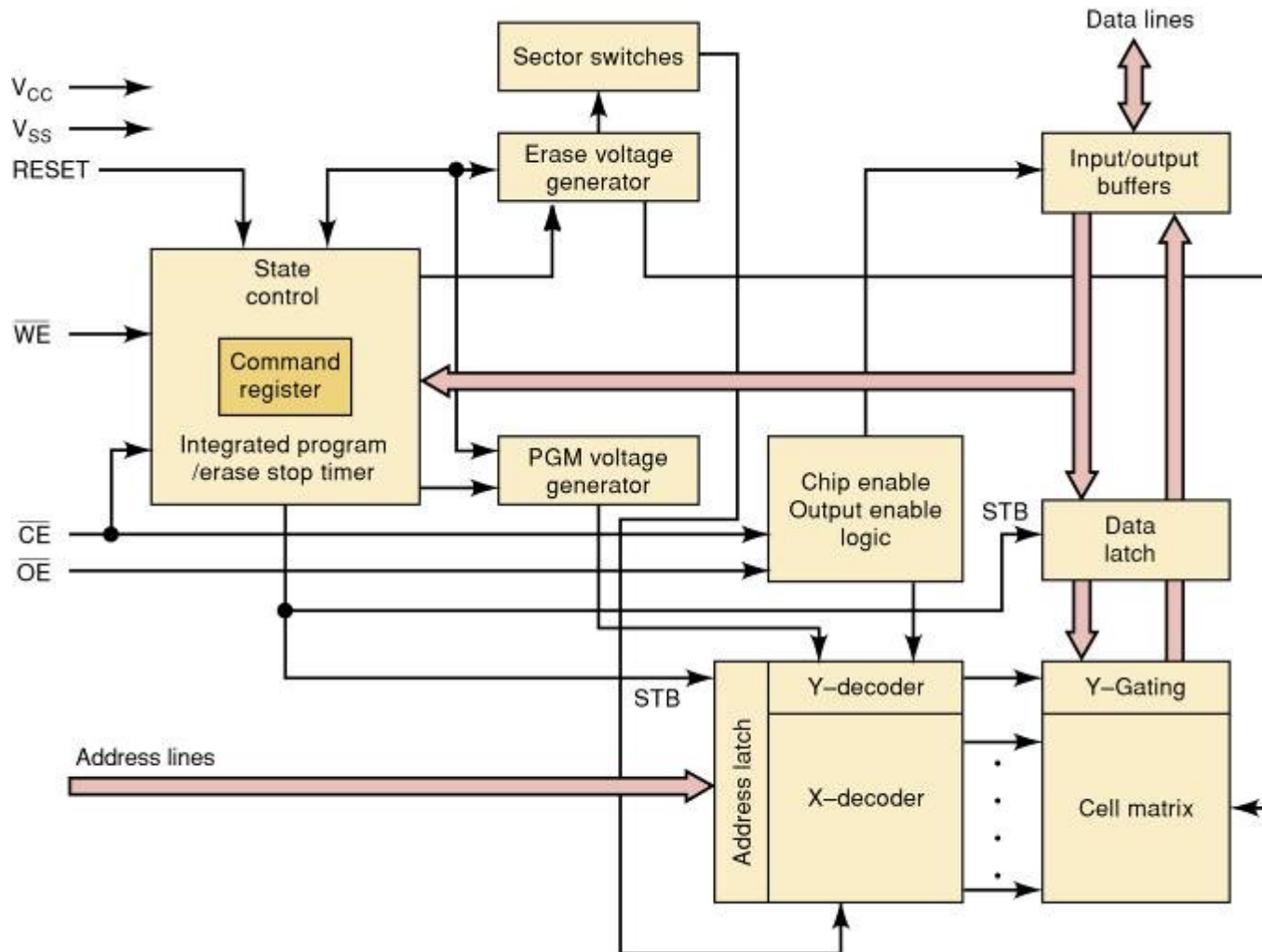
- The first flash devices, created to improve on EEPROM, used *NOR flash* technology.
  - Circuit functions logically like a **NOR** gate.
    - Each transistor can be read or written independent of the status of the other transistors in the group.


Bit line   WL5   WL4   WL3   WL2   WL1   WL0

- **NOR** flash offers quick read access time and random access.
  - Usually used for things like storing program instructions for the microcontroller in your cell phone or PDA.

- Attempts were made to improve density of mass-storage flash devices, resulting in the ***NAND flash***.
  - Data must be accessed in conjunction with the other *word* lines being activated by a control gate voltage.
  - Activated by a control gate voltage large enough to turn on the other transistors—regardless of the amount of charge on the floating gate.
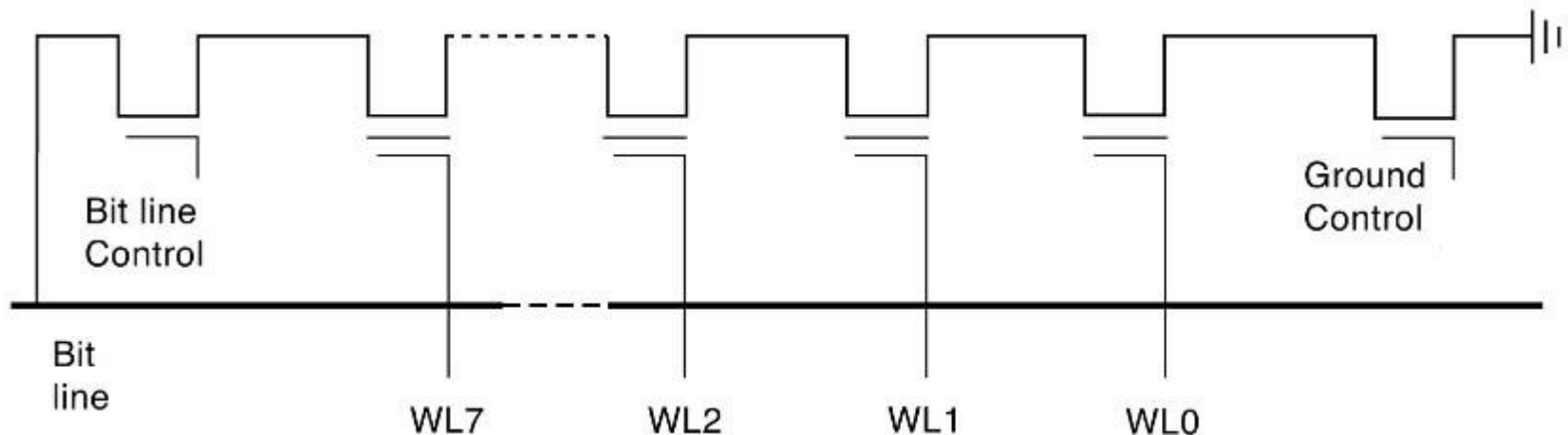
- **NAND** flash circuits offer fast erase & program time—but the data must be dealt with in blocks.
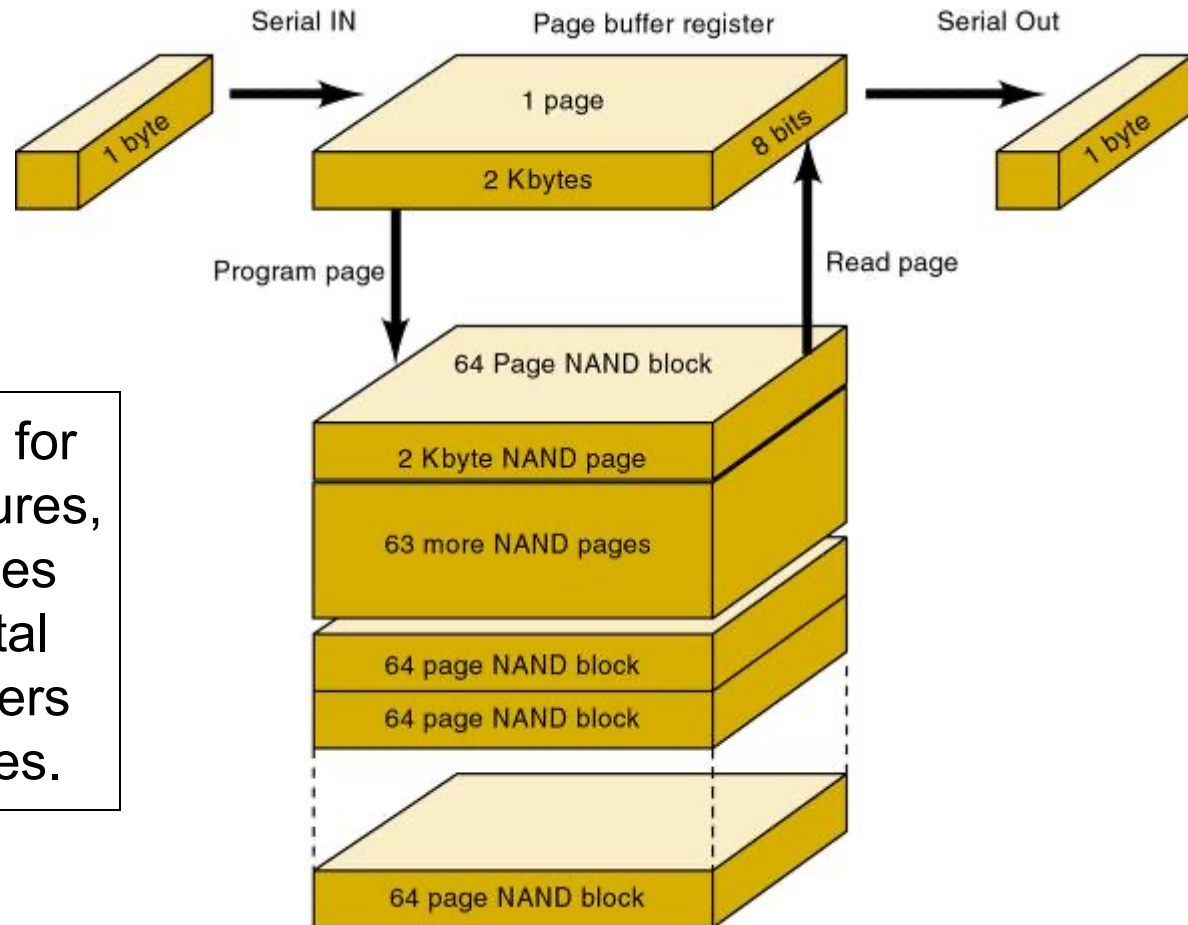
NAND flash is used for mass storage of pictures, music, and other files in devices like digital cameras, MP3 players and USB flash drives.

- Here are some of the most common application areas in which ROMs are used:
  - Embedded microcontroller program memory.
    - Automotive automatic braking systems, cell phones, digital camcorder, microwave ovens, etc.
  - Data transfer and portability.
    - Cell phones, digital cameras, flash drives, MP-3 players.
  - Bootstrap memory.
    - A relatively small program, stored in ROM loads the operating system programs from mass storage (disk) into a computer's main internal memory.

- Here are some of the most common application areas in which ROMs are used:
  - Data tables.
    - Tables of data that do not change like trigonometric and code-conversion tables.
  - Data converter.
    - Data expressed in one type of code is converted to an output expressed in another type, such as BCD to 7-segment LED readouts.

- Here are some of the most common application areas in which ROMs are used:
  - Function generator.
    - Produces waveforms such as sine waves, sawtooth waves, triangle waves, and square waves.
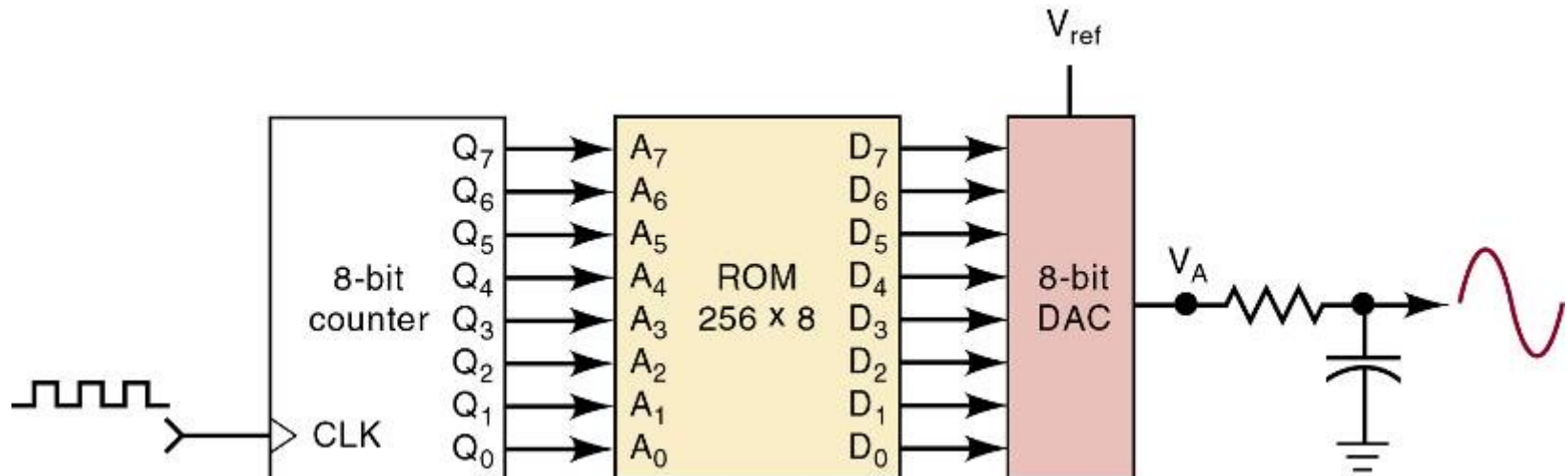
A ROM look-up table and a DAC are used
to generate a sine-wave output signal.

- *RAM—random-access memory*—means any memory address location is as easily accessible as any other.

- Used in computers for *temporary* storage of programs and data—requires fast read/write cycle times to avoid slowing computer operation.
  – RAM can be written into and read from rapidly with equal ease.

- RAM is volatile and will lose all stored information if power is interrupted or turned off.
  – Some CMOS RAMs can be powered from batteries when main power is interrupted.

- It is helpful to think of RAM as consisting of a number of registers.

  – Each storing a single data word, and each having a unique address.

- Most memory chips have one or more CHIP SELECT (*CS*) inputs, used to enable the entire chip or disable it completely.

  – In disabled mode, all data inputs outputs are disabled (Hi-Z)—neither a read nor a write can take place.

- In order to conserve pins on an IC package, manufacturers often combine data input and output functions using common input/output pins.
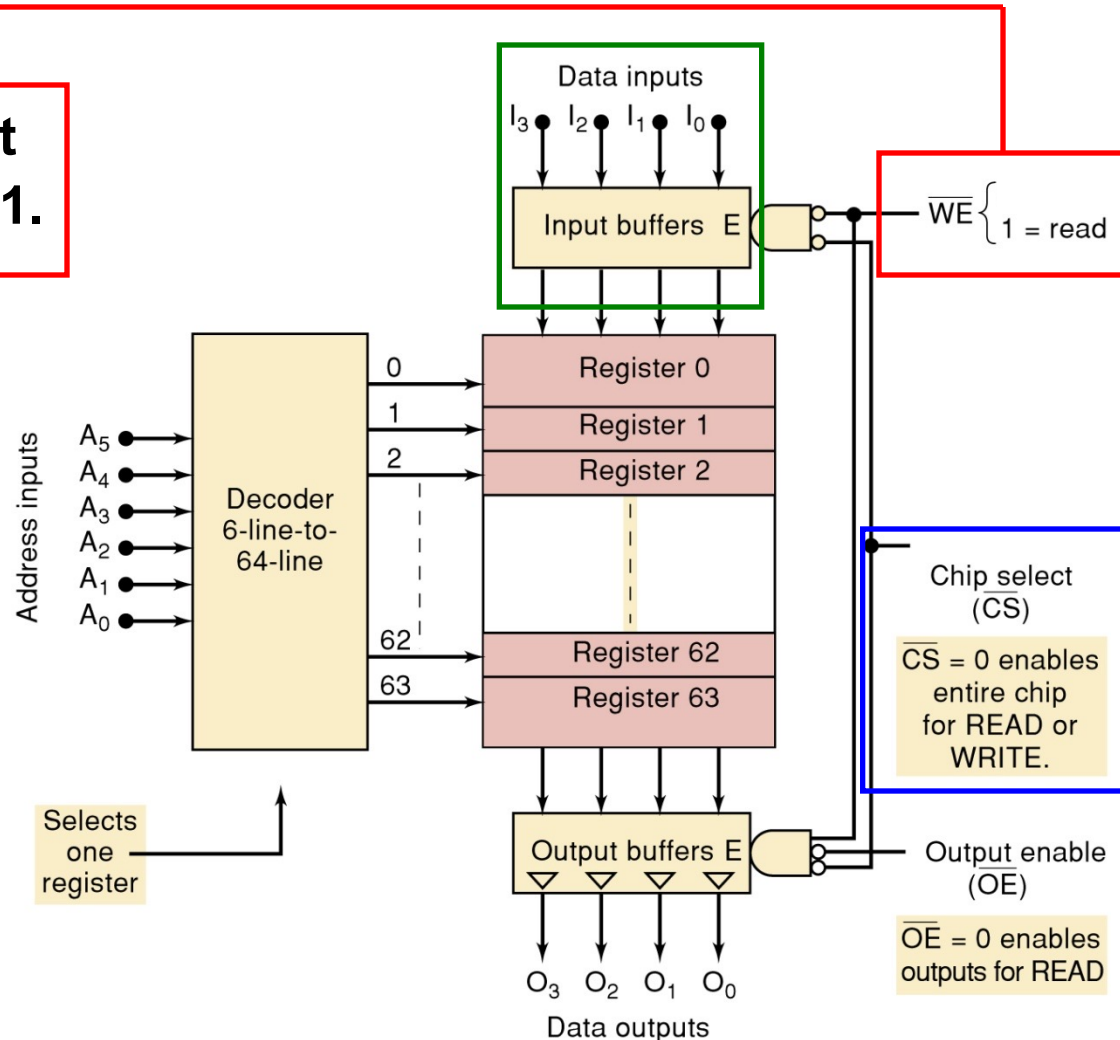
To *READ* the contents of the selected register:

The write enable input $\overline{WE}$ or $R/\overline{W}$ must be a 1.

The CHIP SELECT input must also be activated.
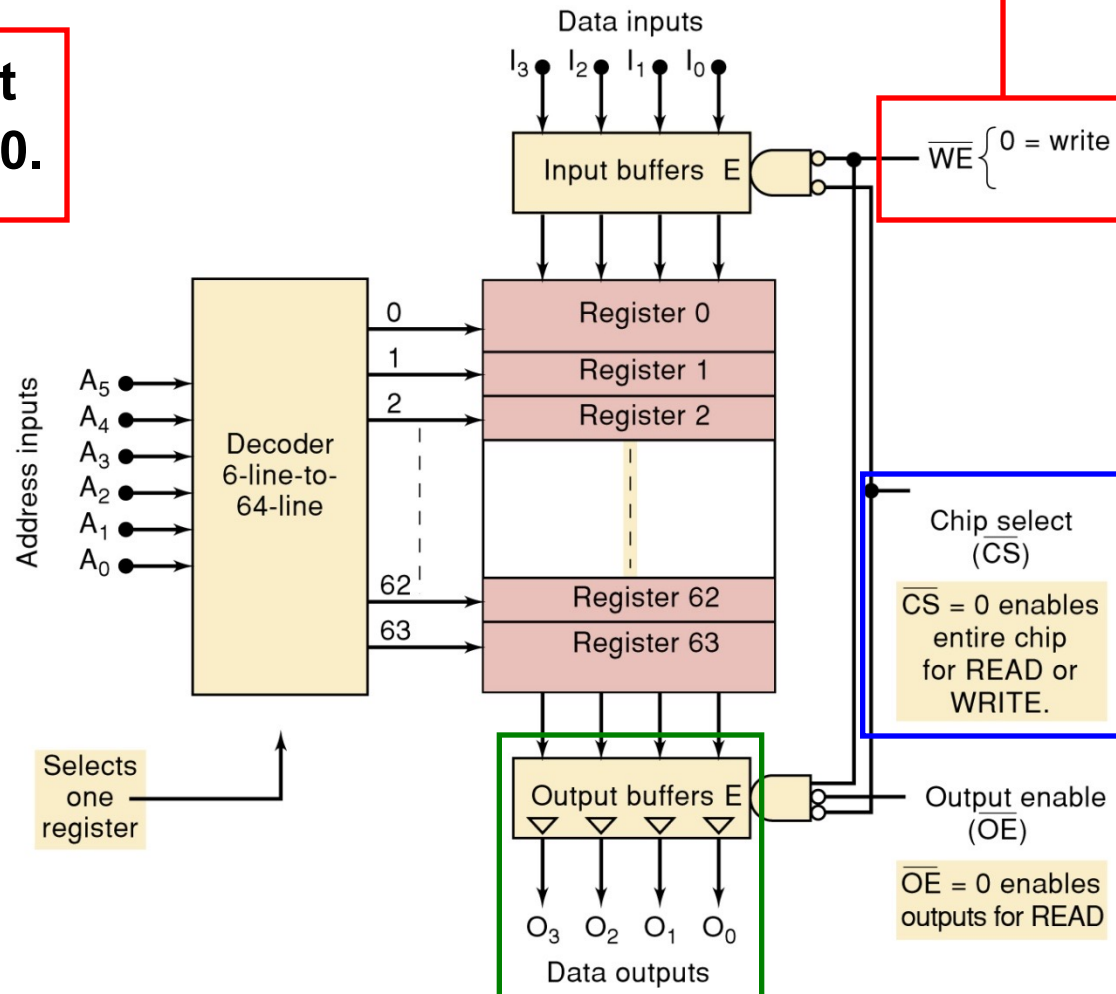
Input buffers are disabled during a data read.

Data inputs

I_3 I_2 I_1 I_0

Input buffers E

$\overline{WE} \begin{cases} 1 = \text{read} \end{cases}$

Address inputs

A_5
A_4
A_3
A_2
A_1
A_0

Decoder
6-line-to-
64-line

0 Register 0
1 Register 1
2 Register 2

62 Register 62
63 Register 63

Selects
one
register

Output buffers E

O_3 O_2 O_1 O_0

Data outputs

Chip select
($\overline{CS}$)

$\overline{CS}$ = 0 enables
entire chip
for READ or
WRITE.

Output enable
($\overline{OE}$)

$\overline{OE}$ = 0 enables
outputs for READ

To *WRITE* a four-bit word to the selected register:

The write enable input $\overline{WE}$ or $R/\overline{W}$ must be a 0.

The CHIP SELECT input must also be a 0.

Tristate output buffers are in Hi-*Z* state during a data write.

Data inputs

I_3 I_2 I_1 I_0

Input buffers   E

$\overline{WE} \left\{ 0 = \text{write} \right.$

Address inputs

A_5
A_4
A_3
A_2
A_1
A_0

Decoder
6-line-to-
64-line

0   Register 0
1   Register 1
2   Register 2

62   Register 62
63   Register 63

Selects
one
register

Chip select
($\overline{CS}$)

$\overline{CS}$ = 0 enables
entire chip
for READ or
WRITE.

Output buffers E

O_3 O_2 O_1 O_0

Data outputs

Output enable
($\overline{OE}$)

$\overline{OE}$ = 0 enables
outputs for READ

- Static-RAM memory cells are essentially flip-flops that stay in a given state (store a bit) indefinitely.
  - Provided power to the circuit is not interrupted.

- Available in bipolar, MOS and BiCMOS variations
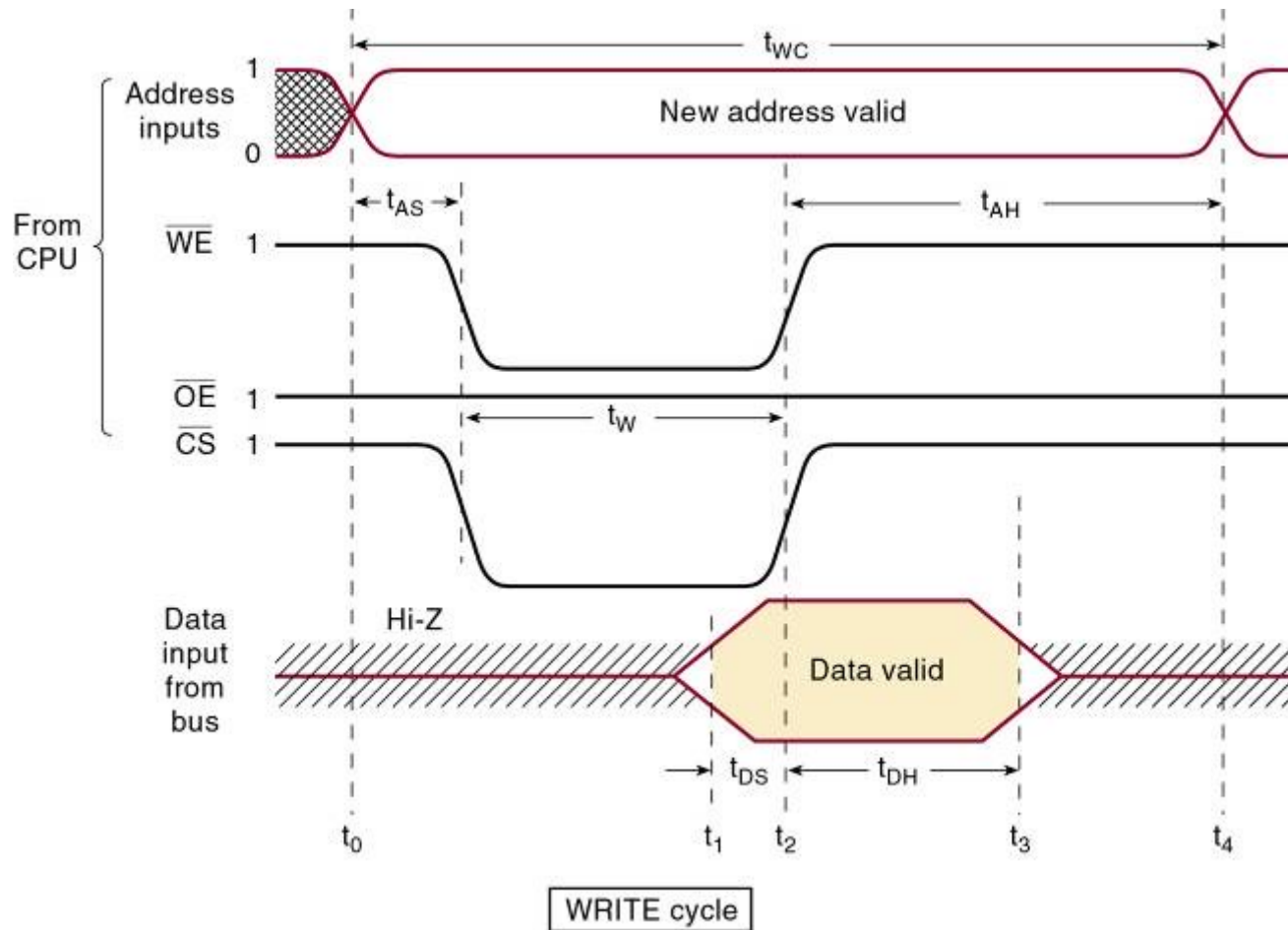 - The majority of applications today use CMOS RAMs.

Bipolar cell   NMOS cell   CMOS cell
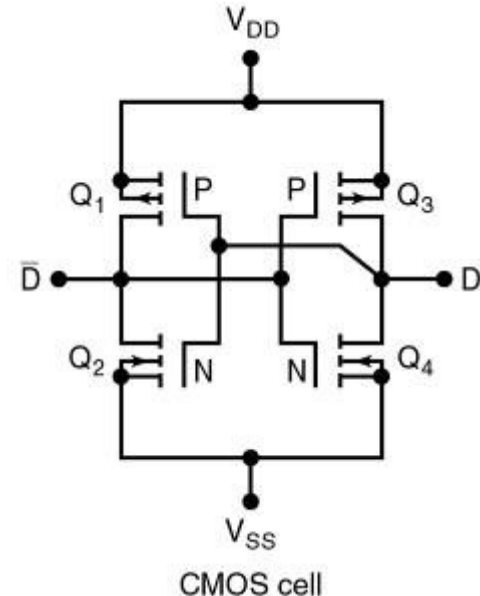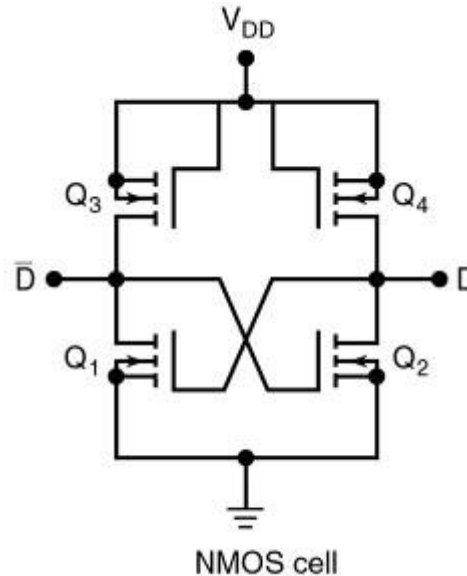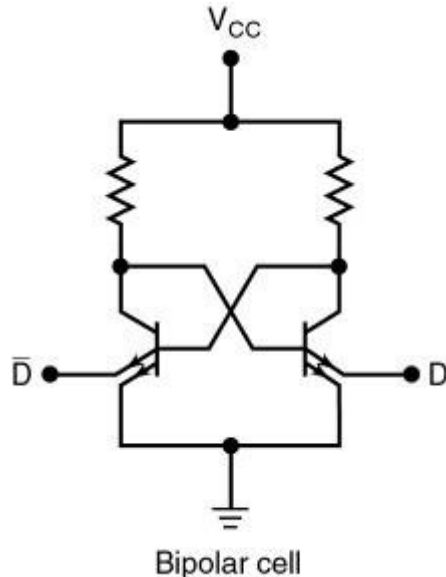
Timing diagram for a complete READ cycle for a typical RAM chip.
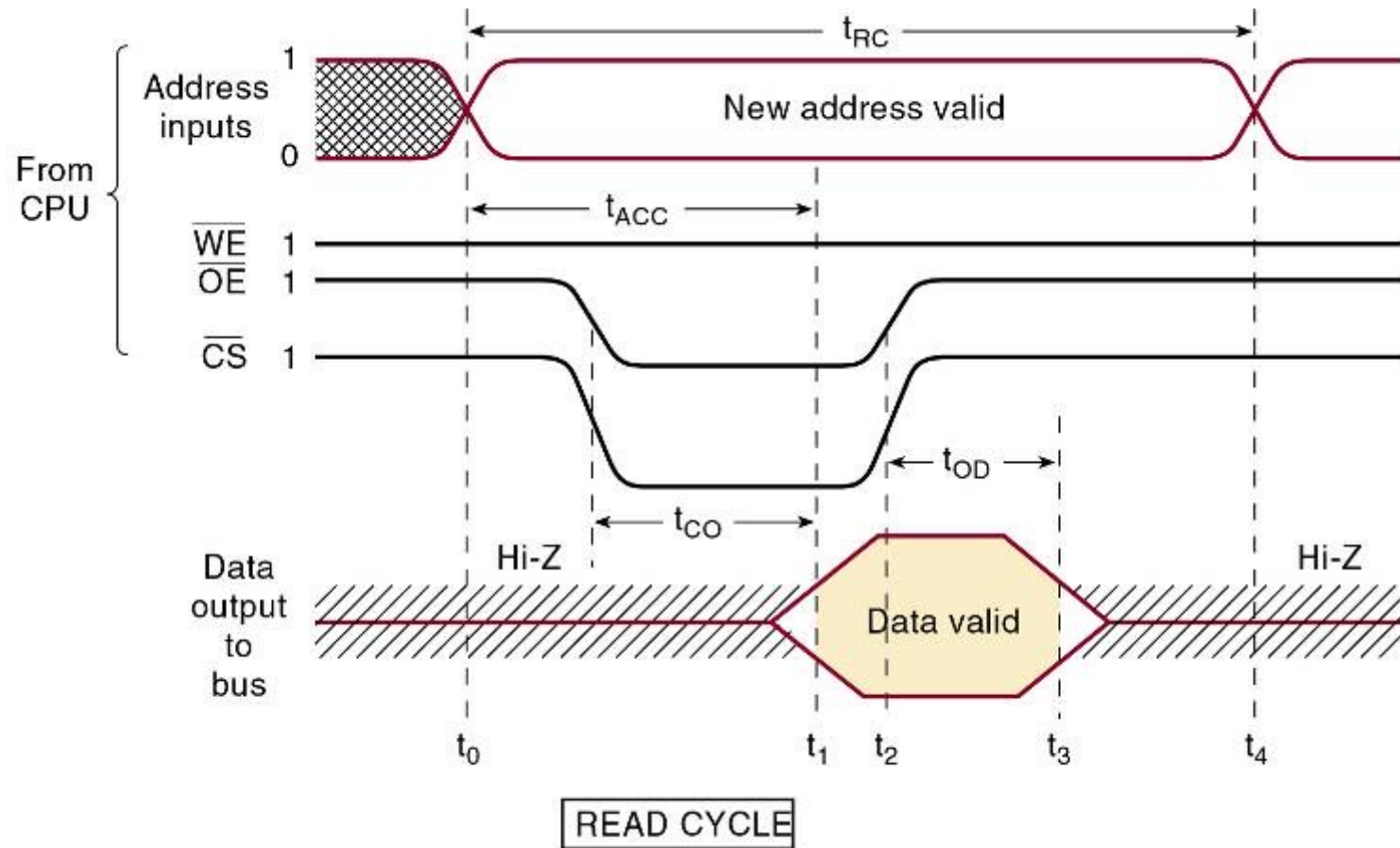
Timing diagram for a complete WRITE cycle for a typical RAM chip.

WRITE cycle

- An example of an actual SRAM IC is the MCM6264C CMOS 8K x 8 RAM,
  - Read- and write-cycle times of 12 ns.
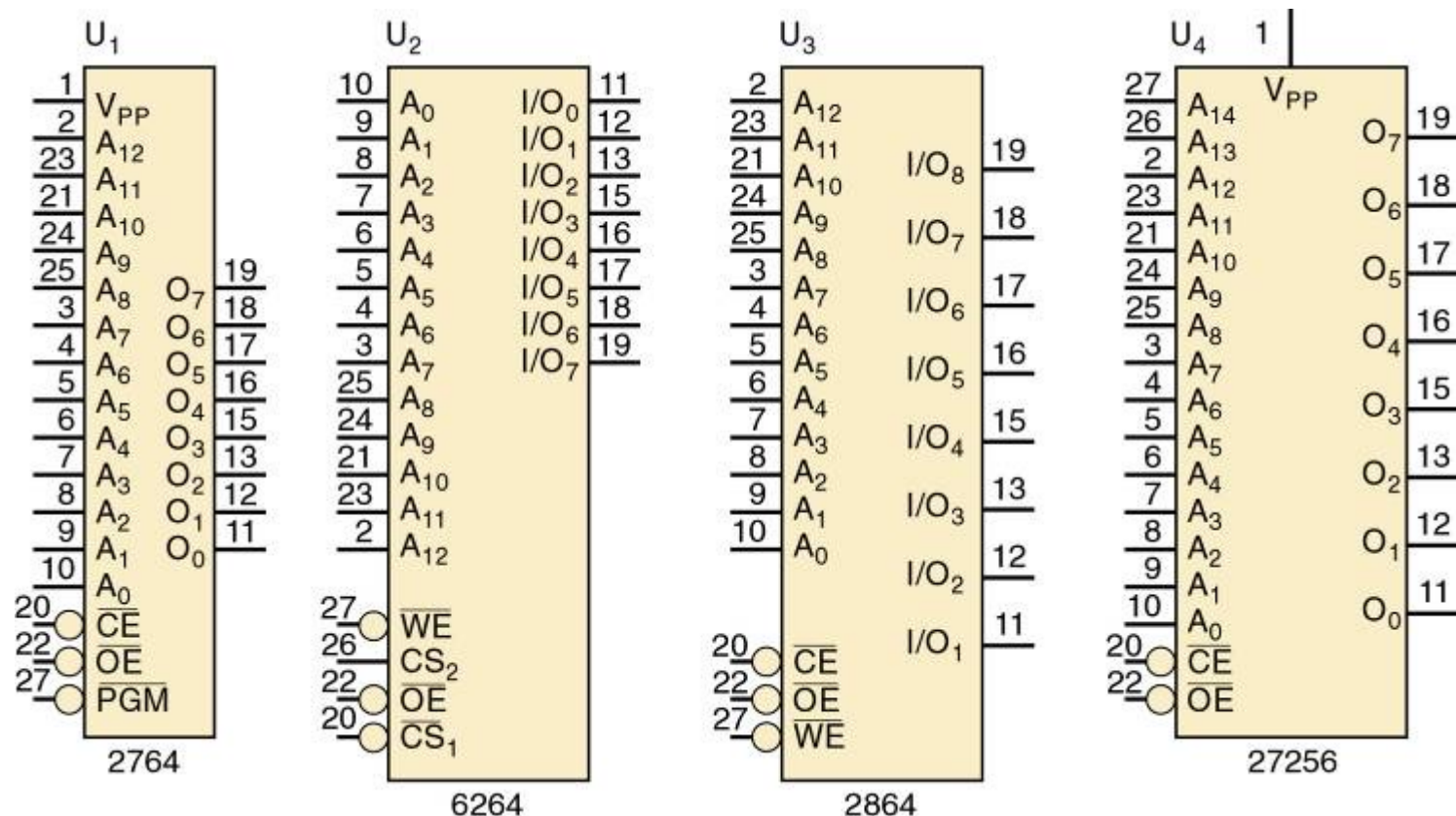 - Standby power consumption of only 100 mW.
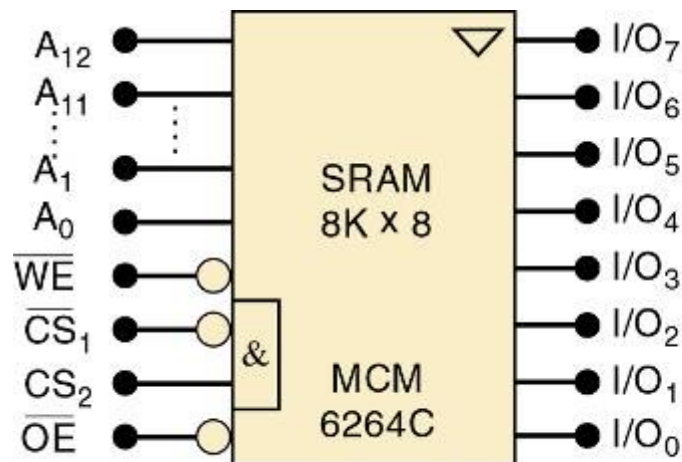
Mode	Inputs				I/O pins
	$\overline{WE}$	$\overline{CS_1}$	CS_2	$\overline{OE}$	
READ	1	0	1	0	DATA$_{OUT}$
WRITE	0	0	1	X	DATA$_{IN}$
Output disable	1	X	X	1	High Z
Not selected	X	1	X	X	High Z
po er do n	X	X	0	X	

X = don't care

- Industry standards created by the Joint Electronic Device Engineering Council **(JEDEC)** have led to memory devices that are interchangeable.

- **Dynamic RAM** stores data as charges on capacitors, which gradually disappear due to capacitor discharge.
  - It is necessary to **refresh** the data periodically by recharging capacitors—typically every 2, 4, or 8 ms.

- Much larger capacities and much lower power consumption make DRAMs the memory of choice.
  - Where the most important design considerations are keeping down size, cost, and power.

- ## The dynamic RAM's internal architecture can be visualized as an array of single-bit cells.
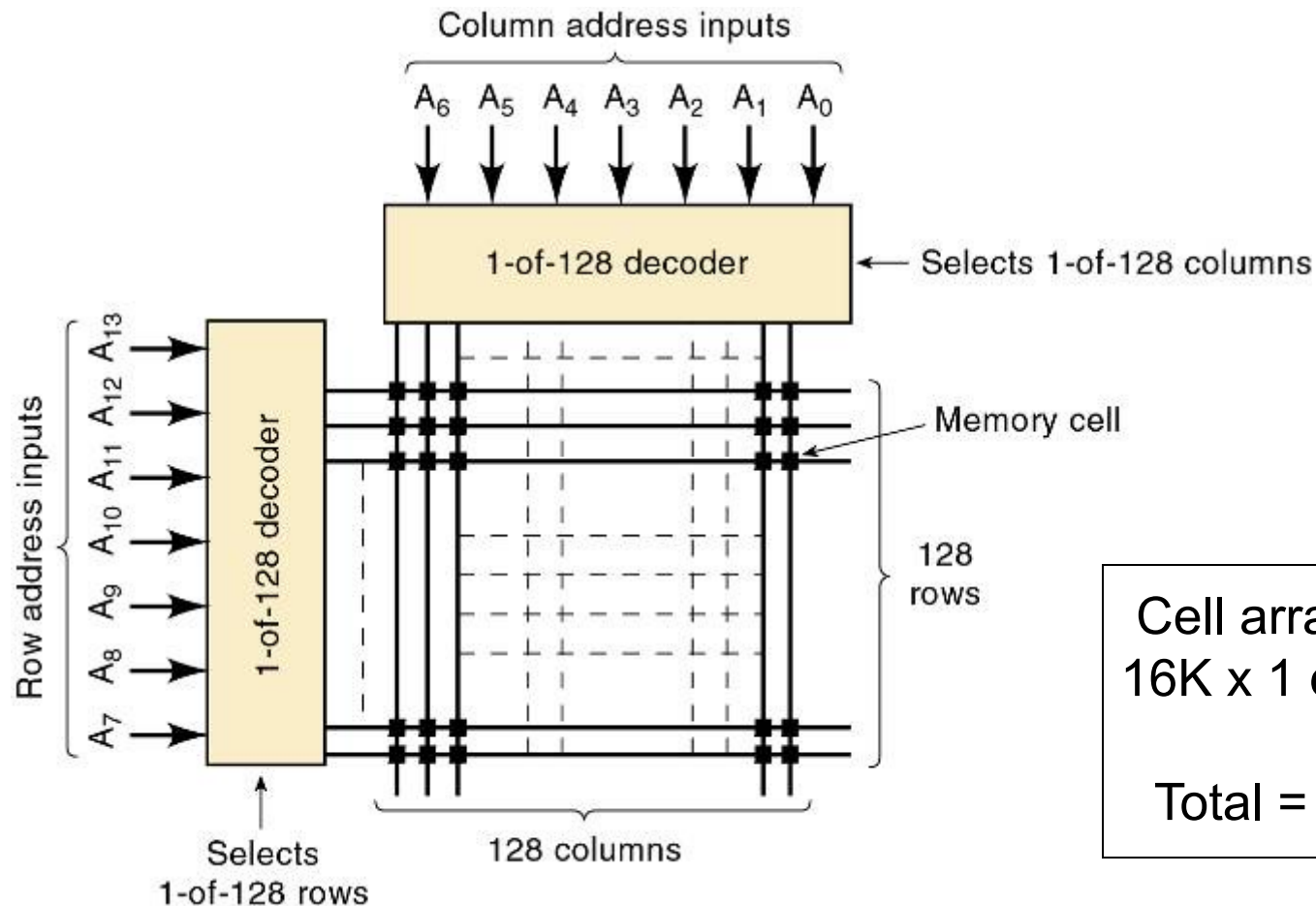
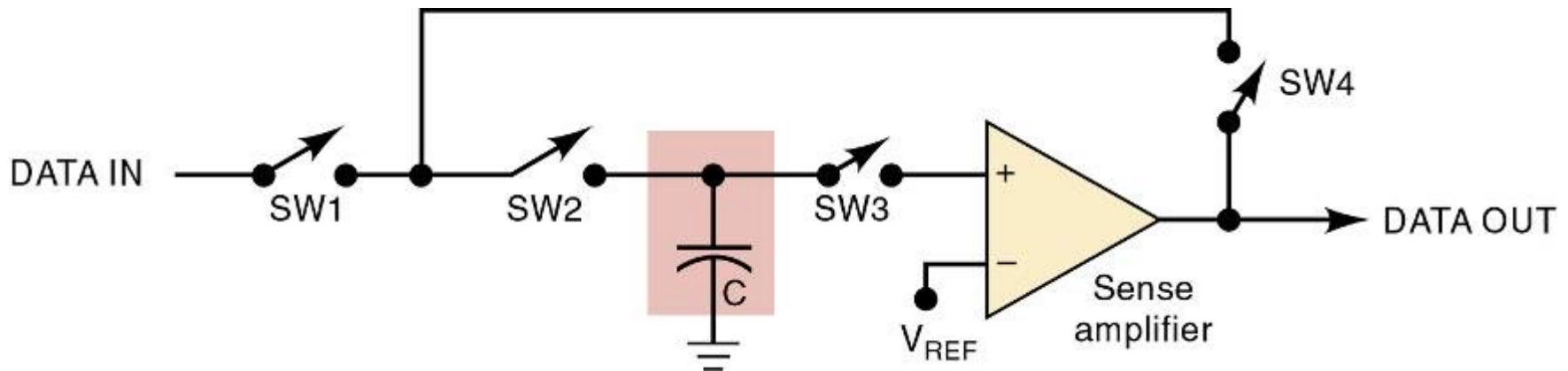Cell arrangement in a 16K x 1 dynamic RAM.

Total = 16,384 cells.

Simplified architecture of a typical DRAM.
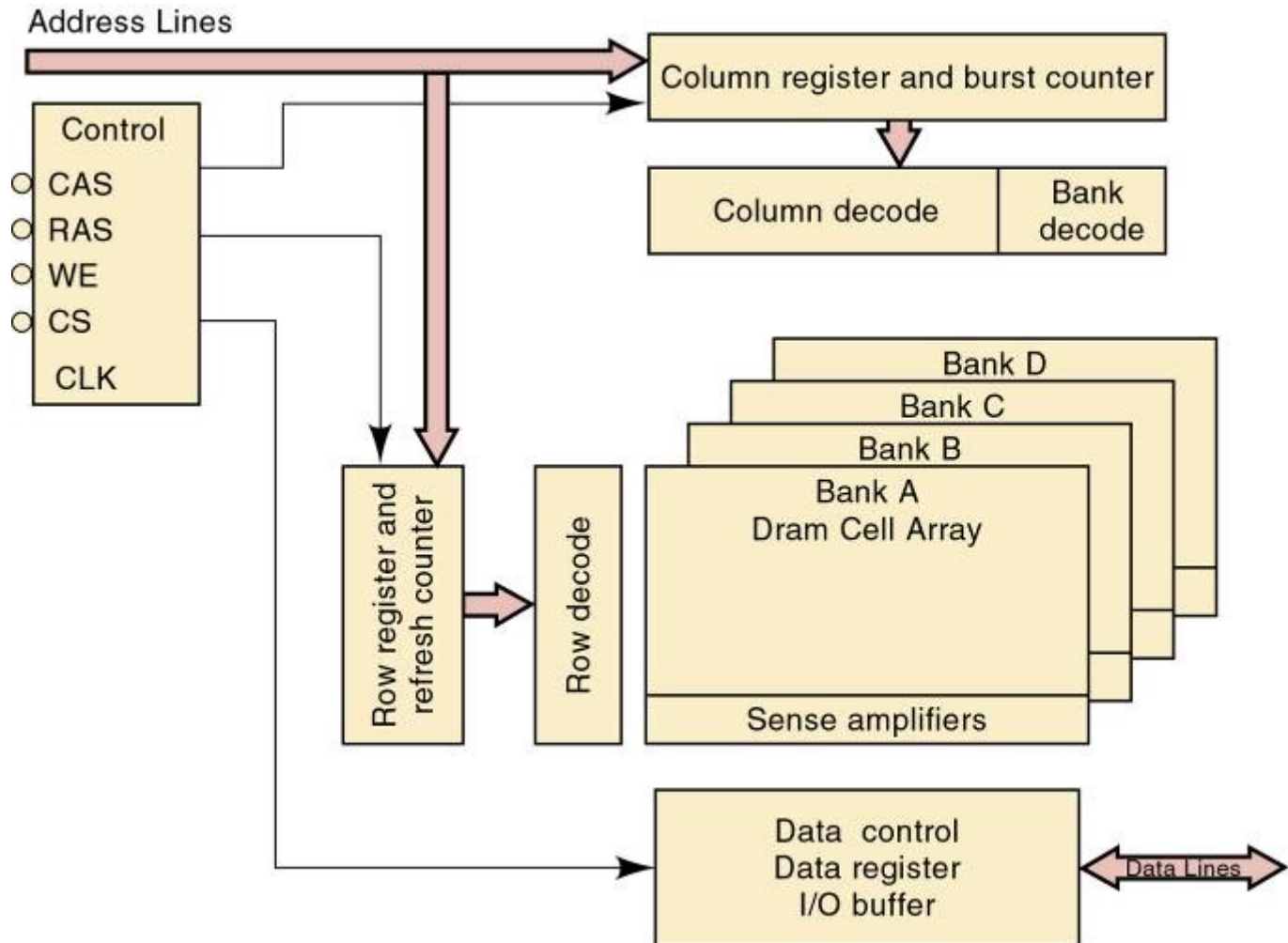
Essential ideas involved in writing to, and reading from a DRAM.

During a WRITE operation, switches SW1 and SW2 are closed. During a read operation, all switches are closed *except* SW1.

- In order to reduce the number of pins, high-capacity DRAM utilize **address multiplexing**.
  - Each address input pin can accommodate two different address bits.

- In multiplexed addressing, the address is applied in two parts—the row, and then column address.

- The address lines are connected directly to both the row & column address registers.
  - The row register stores the upper part of the address, and the column register stores the lower part.

- The **row address strobe** ($\overline{RAS}$) stores the contents of the address inputs into the row address register.

- The **column address strobe** ($\overline{CAS}$) stores the contents of the address inputs into the column address register.

When the CPU wants to access a particular memory location, it generates the complete address and places it on address lines that make up an address bus.
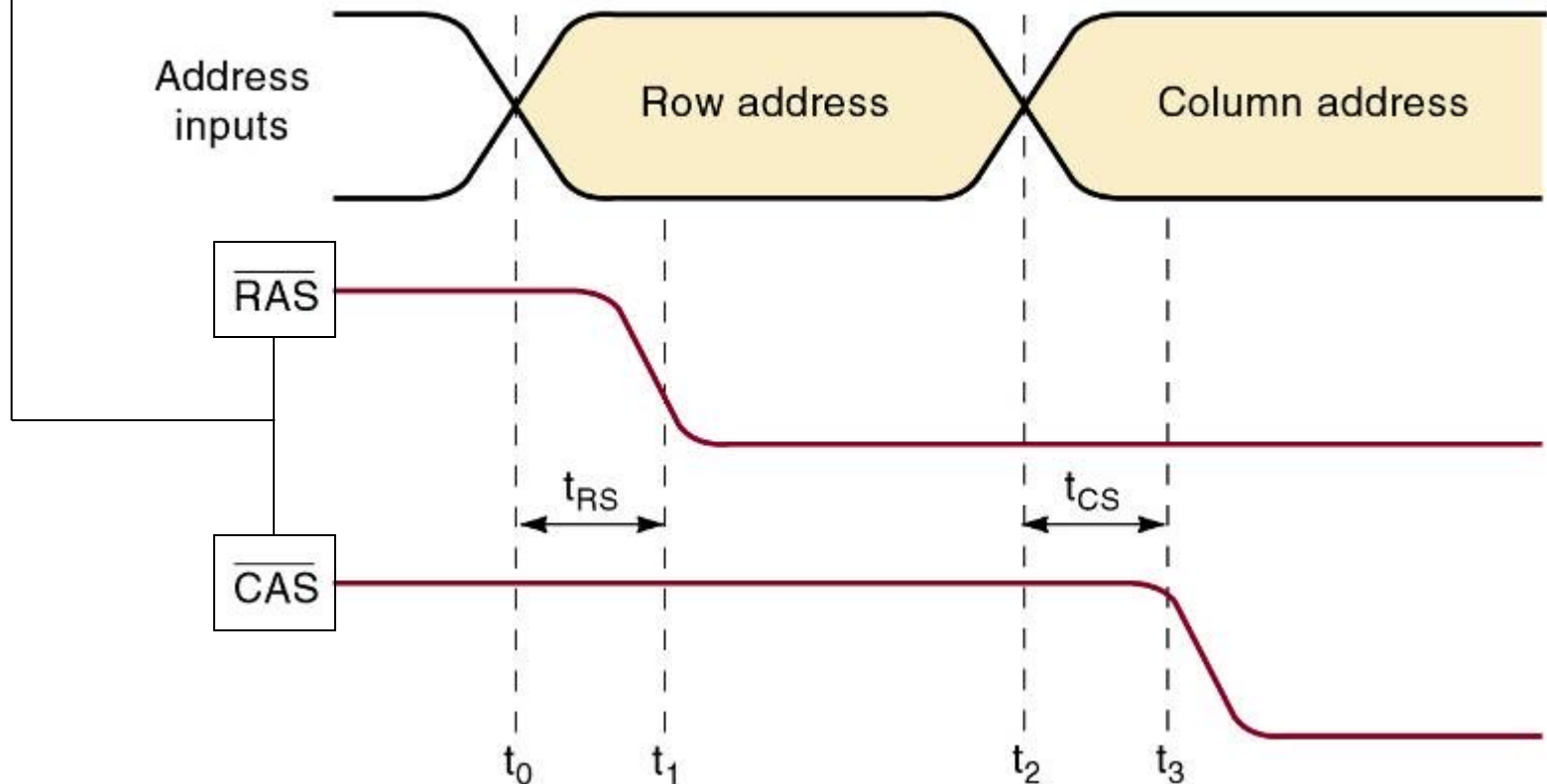
Signal activity for a READ operation on a dynamic RAM.
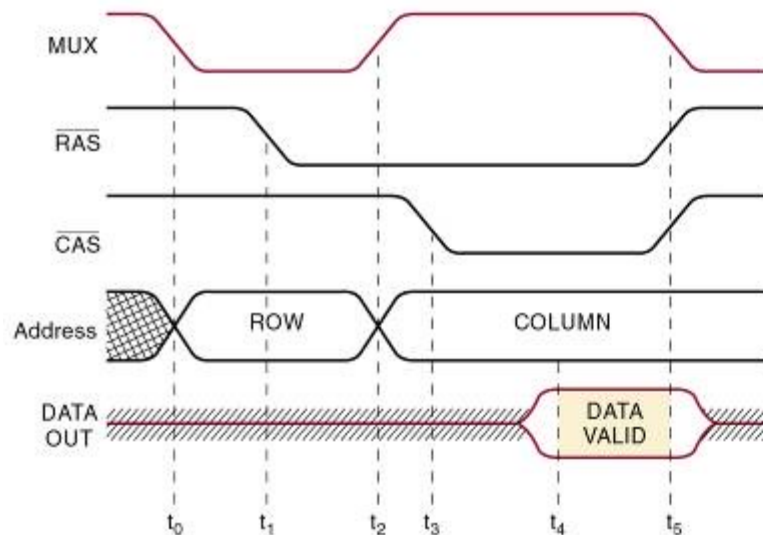
t_0: *MUX* is driven LOW to apply the row address bits (A_8 to A_{15}) to the DRAM address inputs.

t_1: $\overline{RAS}$ is driven LOW to load the row address into the DRAM.

t_2: *MUX* goes HIGH to place the column address (A_0 to A_7) at the DRAM address inputs.

t_3: $\overline{CAS}$ goes LOW to load the column address into the DRAM.

t_4: The DRAM responds by placing valid data from the selected memory cell onto the DATA OUT line.

t_5: *MUX*, $\overline{RAS}$, $\overline{CAS}$, and DATA OUT return to their initial states.

Signal activity for a WRITE operation on a dynamic RAM.

t_0: The LOW at MUX places the row address at the DRAM inputs.

t_1: The NGT at $\overline{RAS}$ loads the row address into the DRAM.

t_2: MUX goes HIGH to place the column address at the DRAM inputs.

t_3: The NGT at $\overline{CAS}$ loads the column address into the DRAM.

t_4: Data to be written are placed on the DATA IN line.

t_5: $\overline{WE}$ is pulsed LOW to write the data into the selected cell.

t_6: Input data are removed from DATA IN.

t_7: MUX, $\overline{RAS}$, $\overline{CAS}$, and $\overline{WE}$ are returned to their initial states.
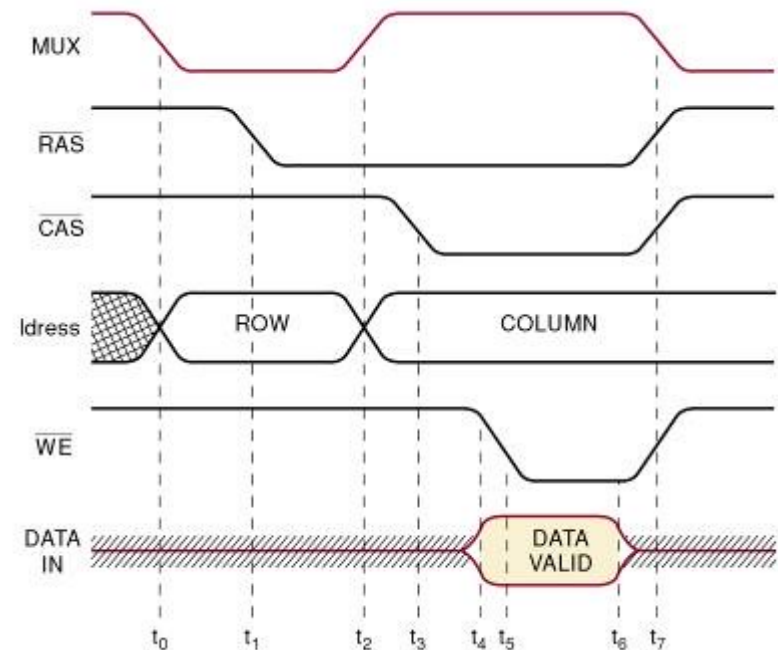
- When a read operation is performed on a cell, all of the cells in the row will be refreshed.

- Refresh control logic is used to make sure each row is refreshed within the time limit.

  – In *burst refresh* mode, normal memory operation is suspended, and each row is refreshed in succession until all rows have been refreshed.

  – In a *distributed refresh*, row refreshing is interspersed with the normal operations of the memory.

- The most universal method for refreshing a DRAM is the $\overline{RAS}$-only refresh.
  - It is performed by strobing in a row address with $\overline{RAS}$ while $\overline{CAS}$ and $\overline{WE}$ remain HIGH.

* R/$\overline{W}$ and $\overline{CAS}$ lines held HIGH
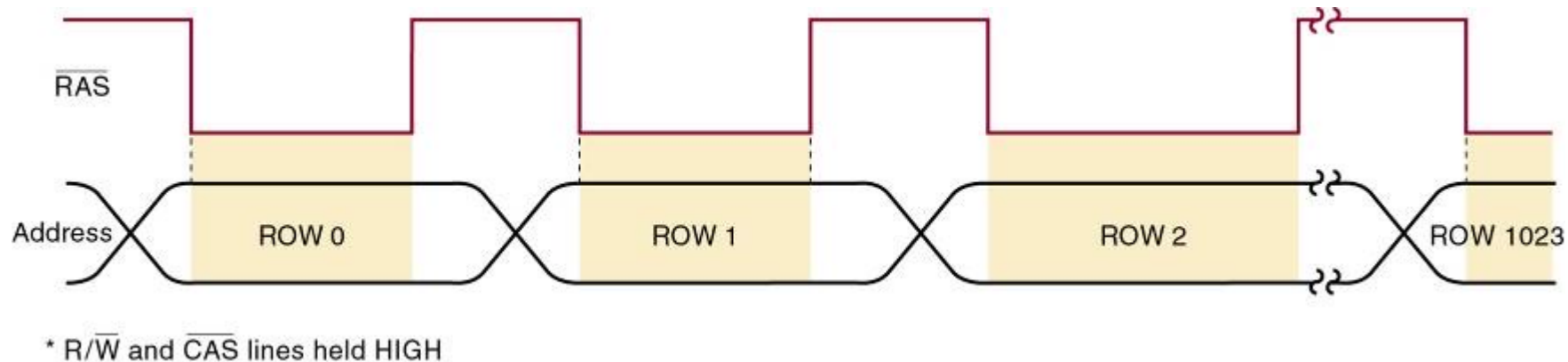
A **dynamic RAM (DRAM) controller** is often used to perform address multiplexing and refresh count sequence generation.

- Standard memory interface connectors are used.
  - The connectors receive a small printed circuit card with contact points on both sides of the card edge.
    - Allow easy memory component installation/replacement.

- Memory modules:
  - **SIMM**—single in-line memory module—a circuit card with 72 functionally equivalent contacts on both sides.
  - **DIMM**—dual-in-line memory module (DIMM) has from 168 to 240 functionally unique pins on each side.
  - **SODIMM**—small-outline, dual-in-line memory module for compact applications, such as laptop computers.
  - **RIMM**—Rambus In-line Memory Module, a proprietary package that holds Direct Rambus DRAM (DRDRAM).

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- **FPM DRAM**—fast page mode (FPM) allows quicker access to random memory locations within the current "page."
  - A page is a range of memory addresses that have identical upper address bit values.
    - Only the lower address lines must be changed.

- **EDO DRAM**—Extended data output (EDO) DRAM offers a minor improvement to FPM.
  - For accesses on a given page, the data value at the current memory location is sensed and latched onto the output pins.
  - While these data are present on the outputs, a new address on the current page can be decoded, and data path circuitry can be reset for the next access.
    - This allows the memory controller to be outputting the next address at the *same* time the current word is being read.

- **SDRAM**—synchronous DRAM is designed to transfer data in rapid-fire bursts of several sequential memory locations.
  - The first location accessed is the slowest due to the overhead (latency) of latching row & column address.
    - After this data values are clocked by the bus system.

- SDRAMs are organized in two (or more) banks.
  - Allows data to be read out at a very fast rate by alternately accessing each of the two banks.

- Self-refresh mode allows the memory device to perform all of the necessary functions to keep its cells refreshed.

- **DDRSDRAM**—Double Data Rate SDRAM refers to the memory module's interface to the PC bus.
  - Achieves higher data rates by transferring data on the rising and falling edge of the system clock.
    - Burst transfer rates twice as fast as the SDRAM ICs.

- DDR2 uses buffering techniques to produce I/O data rates *four times* faster than the SDRAMs.
  - DDR3 transfers data *eight times* faster.

- Speeding up the system clock offers marginal improvement in performance—given that the SDRAM latency is the ultimate limiting factor of maximum speed.

- The first method of magnetic storage of digital information involved reels of magnetic tape for long-term storage/retrieval of programs and data.

- The next improvement involved coating rigid (hard) disks with magnetic media and rotating the disks while moving a magnetic read/write head radially across the disk.
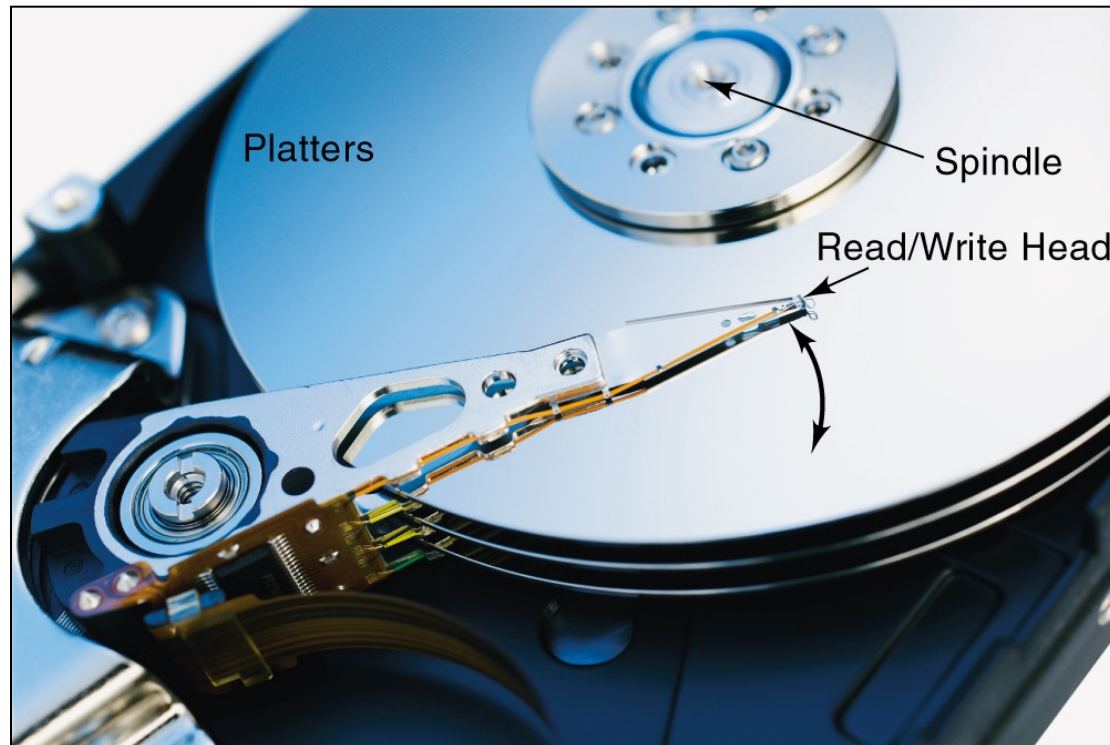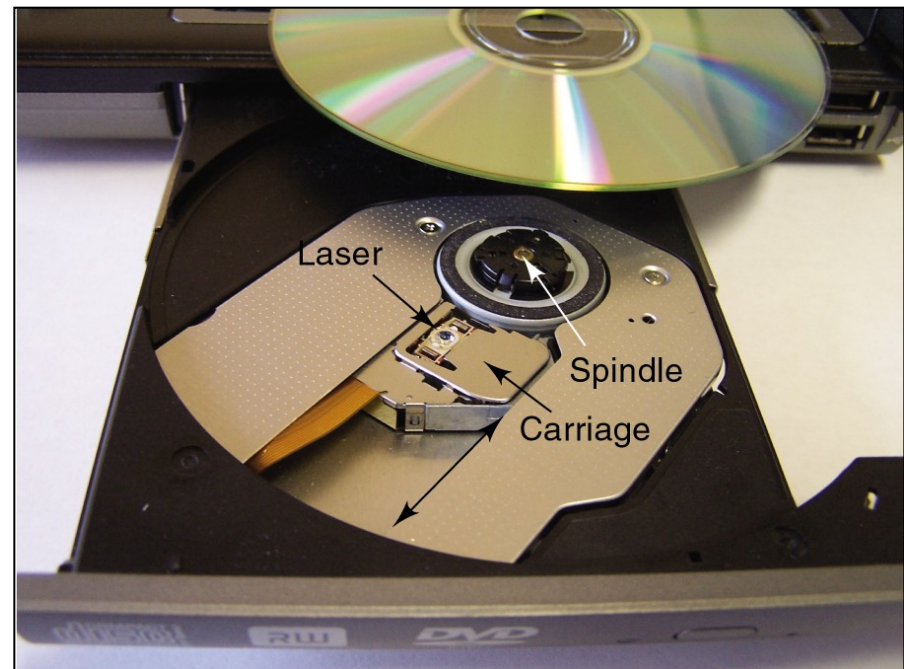
- High-speed, random-access, nonvolatile magnetic storage was also tried in the early days of computers using "magnetic core" technology.
  - Rows and columns of little electromagnets that could be polarized in either direction.

- This basic technology has been brought back recently in the form of **magnetoresistive random access memory (MRAM**).

- ## The optical disc is a very significant digital storage memory technology.
  - Digital audio compact discs (CDs) became available in the early 1980s,
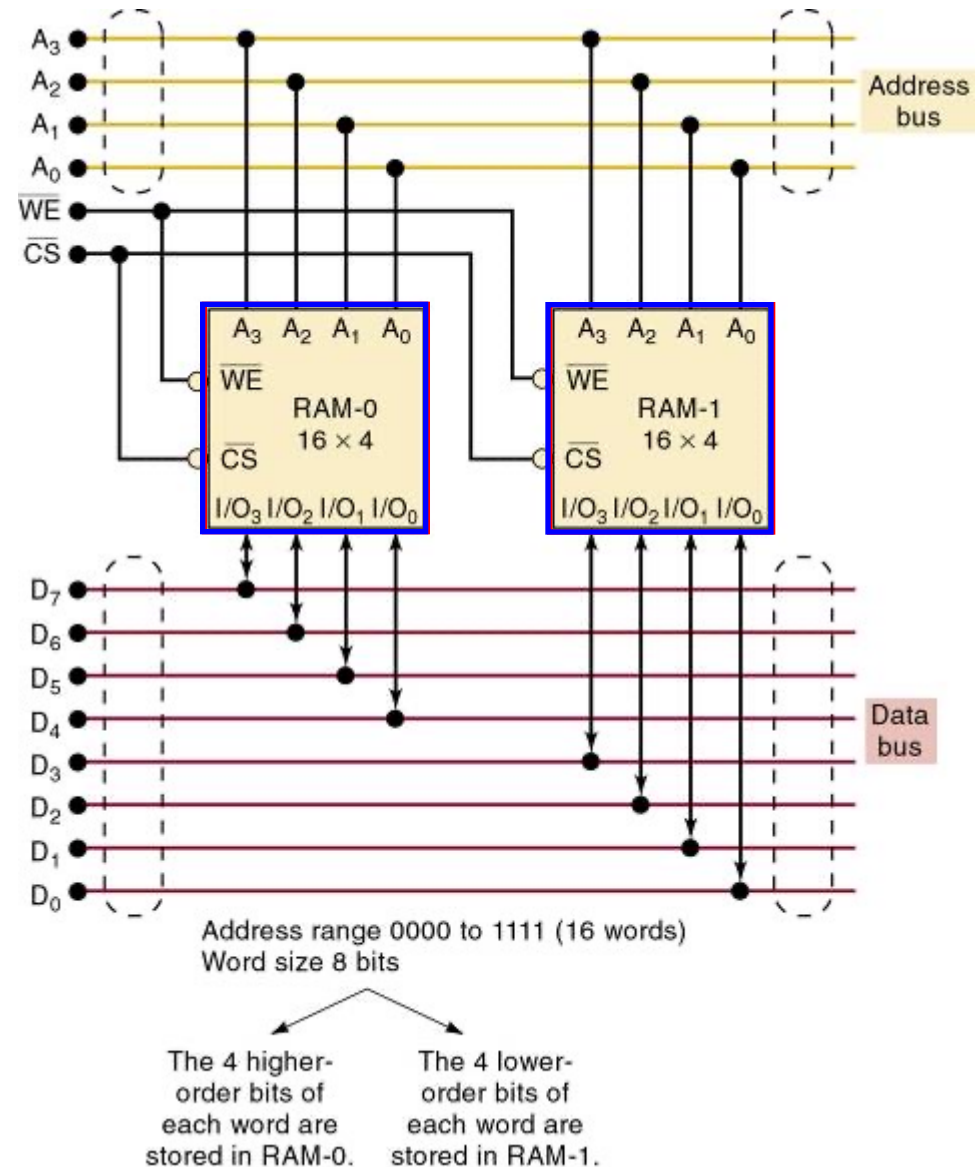 - Digital video (DVD), and Blu-Ray Discs (BD).

All optical storage formats use essentially the same technology, differing largely in format & density.
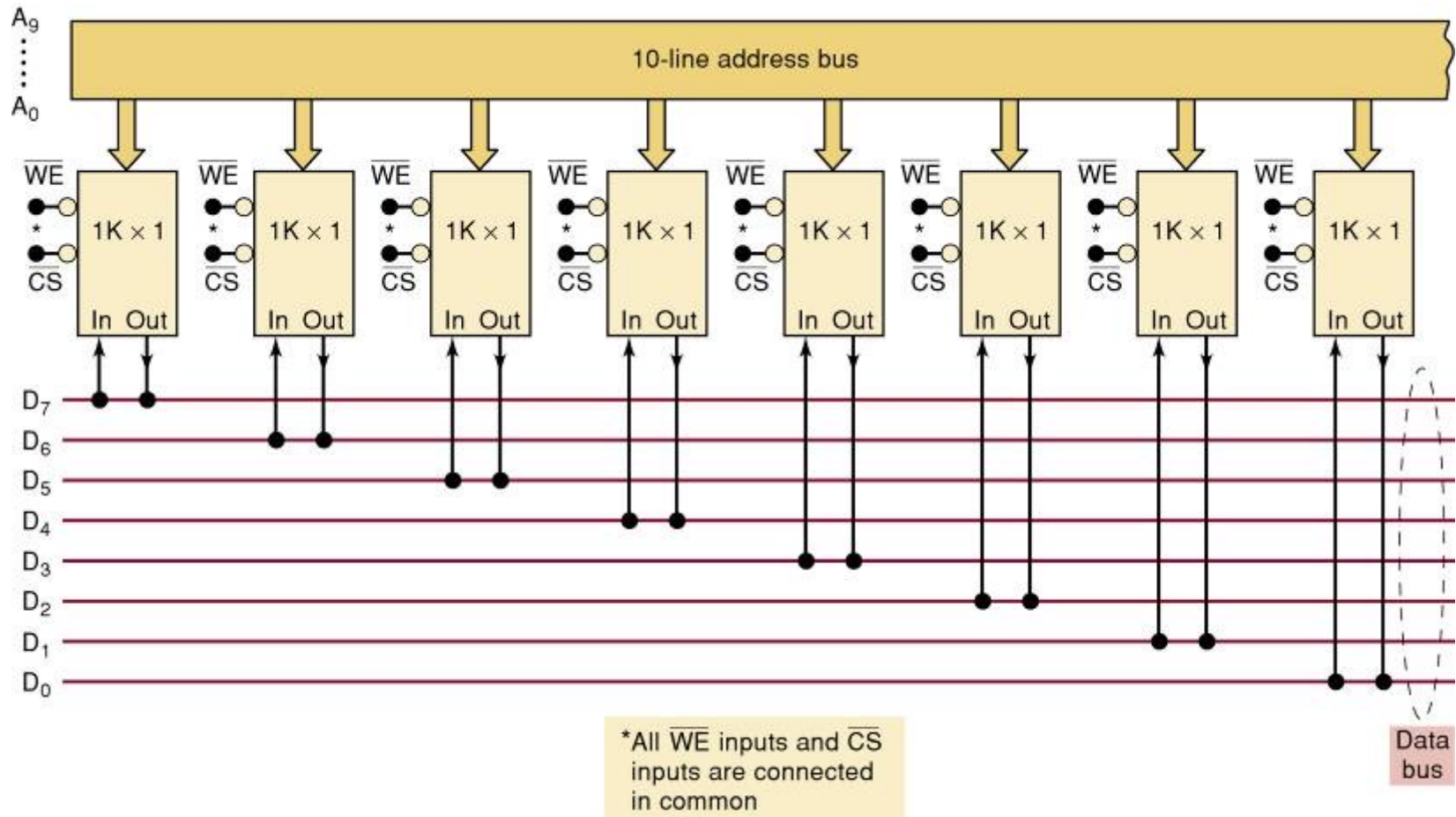
In many applications, required RAM or ROM memory capacity or word size cannot be satisfied by one memory chip.

Several chips must be combined to provide the capacity and/or the word size.

The combination of the two RAM chips acts like a single 16 x 8 memory, and is referred to as a 16 x 8 *memory module.*

Address range 0000 to 1111 (16 words)
Word size 8 bits

The 4 higher-order bits of each word are stored in RAM-0.

The 4 lower-order bits of each word are stored in RAM-1.

Eight 2125A 1K ✕ 1 chips arranged as a 1K ✕ 8 memory.

Incomplete address decoding is useful when different memory devices are used in the same system
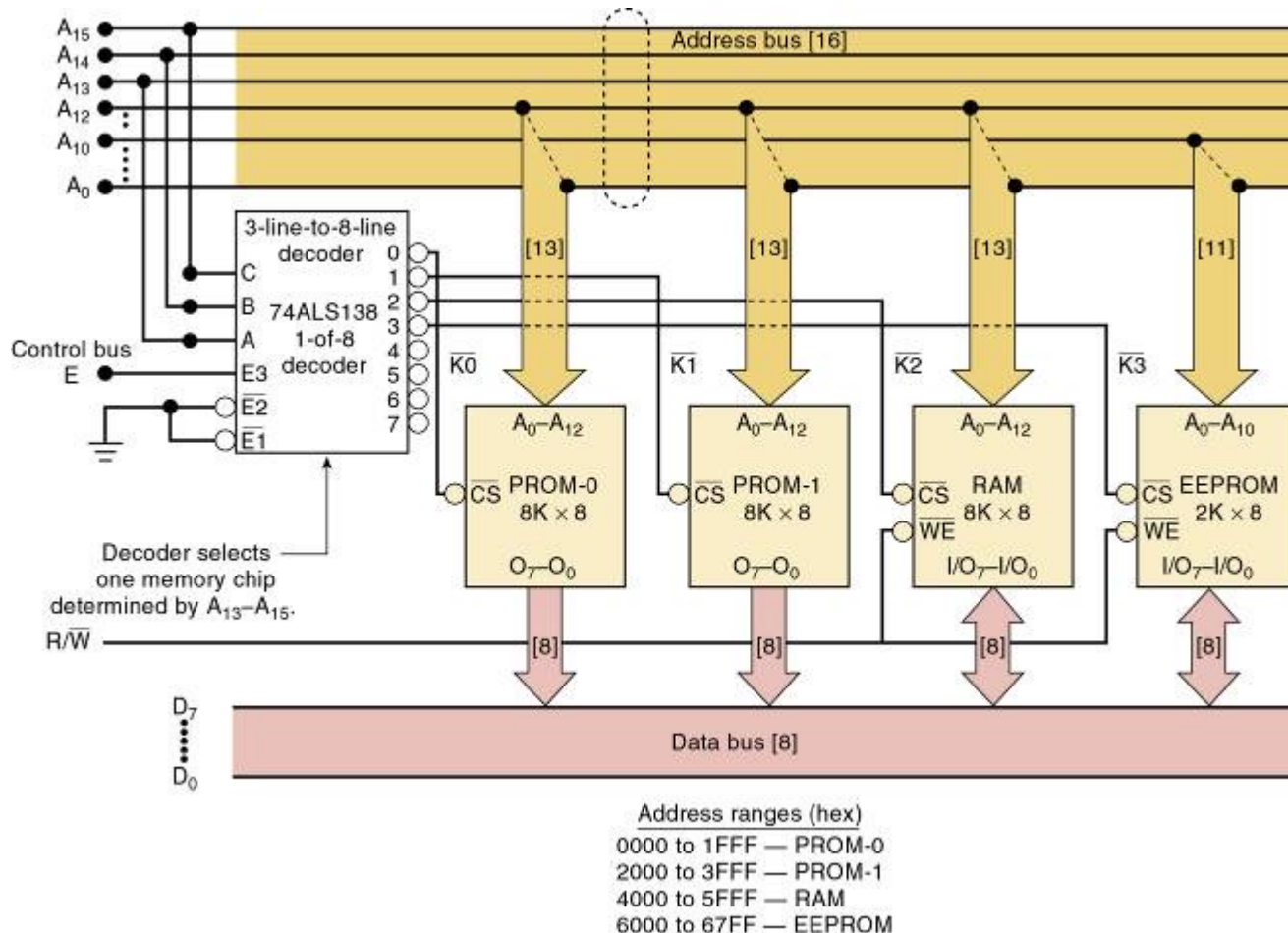
DRAM ICs with word sizes of 1-4 bits must be combined to form larger word size modules
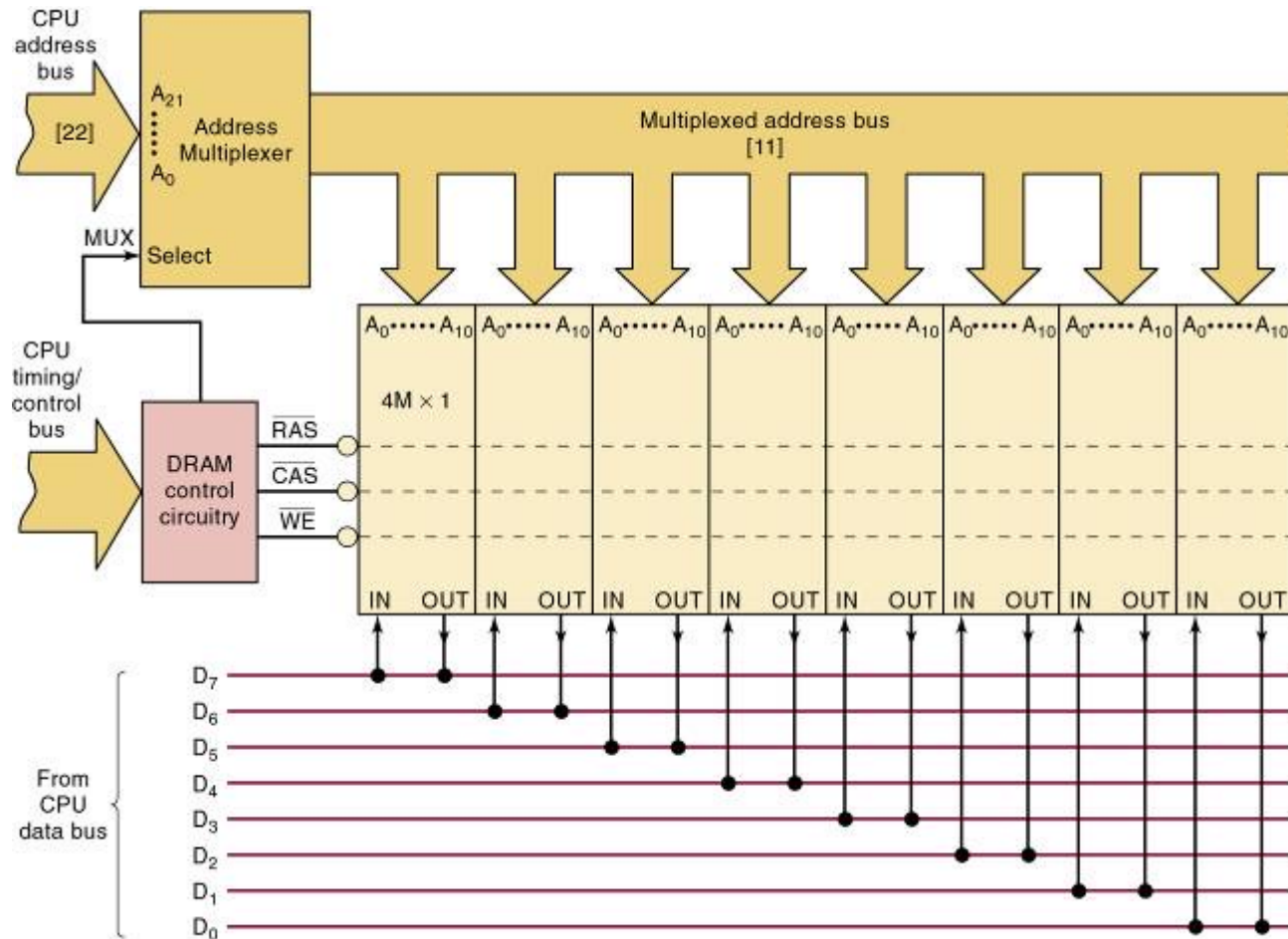
- ## Special functions performed by memory devices in various applications.
    - Power down storage.
        - Critical operating parameters applied when a system is powered up.
        - Industrial process control systems that must retain memory of where they are in a process under all conditions.
    - Cache memory.
        - High speed memory that communicates directly with the CPU
            - Level 1 cache is on CPU.
            - Level 2 cache is SRAM external to CPU.

First in first out (FIFO) memory.

Data written into RAM storage are read out in the same order that they were written in.

Useful as a buffer between systems with different data rates.

First data byte written

1	0 1 1 0 1 1 1 0
2	
3	
4	
5	
6	

Second data byte written

1	1 1 1 0 0 1 1 0
2	0 1 1 0 1 1 1 0
3	
4	
5	
6	

Third data byte written

1	0 0 0 0 0 0 0 1
2	1 1 1 0 0 1 1 0
3	0 1 1 0 1 1 1 0
4	
5	
6	

1	0 0 0 0 0 0 0 1
2	1 1 1 0 0 1 1 0
3	0 1 1 0 1 1 1 0
4	
5	
6	

1	0 0 0 0 0 0 0 1
2	1 1 1 0 0 1 1 0
3	0 1 1 0 1 1 1 0
4	
5	
6	

1	0 0 0 0 0 0 0 1
2	1 1 1 0 0 1 1 0
3	0 1 1 0 1 1 1 0
4	
5	
6	

First data byte
read = 01101110

Second data byte
read = 11100110

Third data byte
read = 00000001
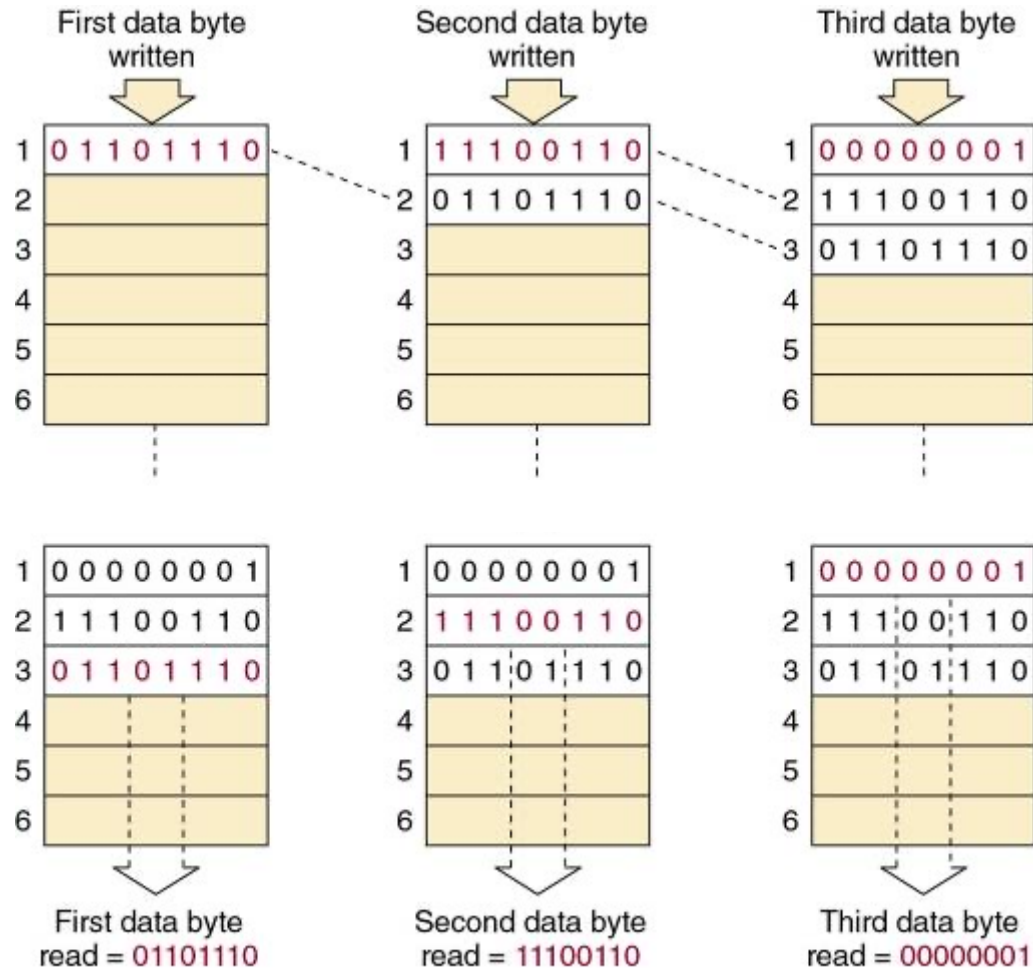
- Data rate buffers (FIFOs) are often referred to as **linear buffers**.
  - As soon as all the locations in the buffer are full, no more entries are made until the buffer is emptied.

- A similar memory system is a **circular buffer**, used to store the last *n* values entered.
  - Where *n* is the number of buffer memory locations.

- Each time a new value is written to a circular buffer, it overwrites (replaces) the oldest value.

- When the highest address is reached, the address counter will "wrap around" and the next location will be the lowest address.

- Faults in RAM can cause unreliable system performance or "crashes."

- In order to determine if RAM is working properly you must know how it normally operates.

- The decoding logic can be tested using signal injection, or by forcing a certain address onto the bus to obtain a known decoder output.

- To test the complete RAM system, patterns of 1s and 0s are written and read from each memory location.
  - By alternating the patterns each bit can be checked for R/W of both 1s and 0s.

- Pattern checking does not catch all errors.
  - There may be errors that occur only in *certain* patterns.

- A memory check is commonly run when a system is powered up.

- Test options include:
  - Printing out a listing of the memory contents.
    - Memory contents are compared to a reference ROM.
  - Checksums.