

Normalization for Relational Databases

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資訊科學系

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Outline

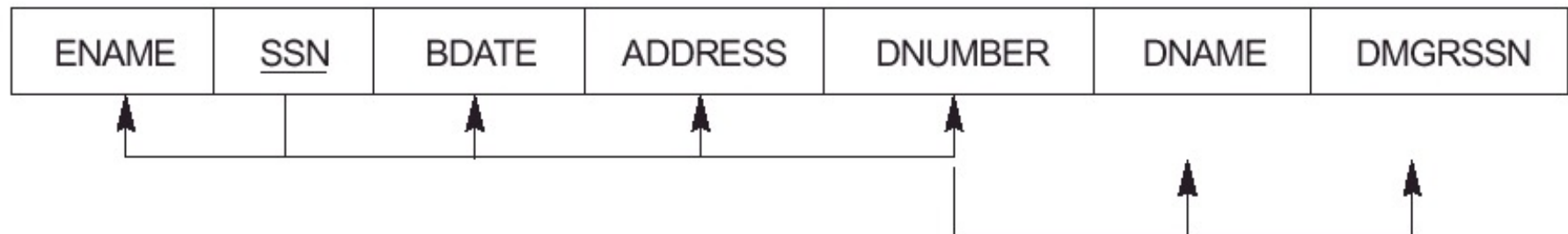
- ◆ Normal Forms based on Primary Keys
- ◆ General Definitions of Second & Third Normal Forms
- ◆ Boyce-Codd Normal Form

Normal Forms Based on Primary Keys

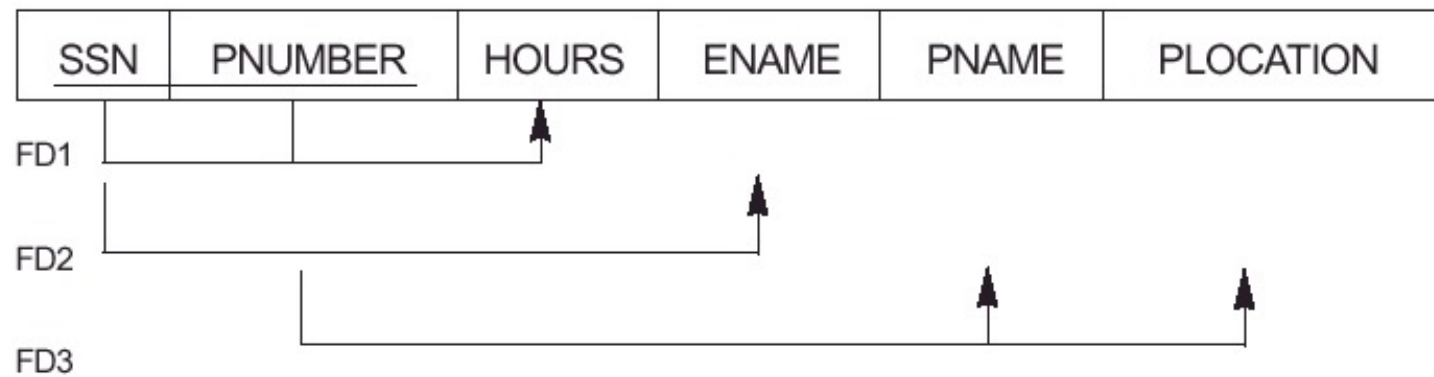
Informal Design Guidelines for Relation Schemas

- ◆ 4 Informal measures of quality for relation schema design
 - Semantics of the attributes
 - Reducing the redundant values in tuples
 - Reducing the null values in tuples
 - Disallowing the possibility of generating spurious tuples

(a)



(b)



EMP_DEPT

ENAME	<u>SSN</u>	BDATE	ADDRESS	DNUMBER	DNAME	DMGRSSN
Smith,John B.	123456789	1965-01-09	731 Fondren,Houston,TX	5	Research	333445555
Wong,Franklin T.	333445555	1955-12-08	638 Voss,Houston,TX	5	Research	333445555
Zelaya, Alicia J.	999887777	1968-07-19	3321 Castle,Spring,TX	4	Administration	987654321
Wallace,Jennifer S.	987654321	1941-06-20	291 Berry,Bellaire,TX	4	Administration	987654321
Narayan,Ramesh K.	666884444	1962-09-15	975 FireOak,Humble,TX	5	Research	333445555
English,Joyce A.	453453453	1972-07-31	5631 Rice,Houston,TX	5	Research	333445555
Jabbar,Ahmad V.	987987987	1969-03-29	980 Dallas,Houston,TX	4	Administration	987654321
Borg,James E.	888665555	1937-11-10	450 Stone,Houston,TX	1	Headquarters	888665555

- To insert a new department that has no employee yet ?
- To delete an employee that happens to be the last employee working for a particular department ?
- To modify information of a department ?

Null Values in Tuples

◆ Guideline 3

- Avoiding placing attributes in a base relation whose values may frequently be **null**
- If nulls are unavoidable, make sure that they apply in exceptional cases only & do not apply to a **majority** of tuples in the relation

- ◆ e.g.: if only 10% employees have individual offices, there is little justification for including an attribute **Office_Number** in the **Employee** relation

◆ e.g.

Dname	<u>Dnumber</u>	Mgr_ssn	Mgr_start_date
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Fname	Minit	Lname	<u>Ssn</u>	Bdate	Address	Sex	Salary	Super_ssn	Dno
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EMP_LOCS

ENAME	PLOCATION
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Smith, John B.	Bellaire
Smith, John B.	Sugarland
Narayan, Ramesh K.	Houston
English, Joyce A.	Bellaire
English, Joyce A.	Sugarland
Wong, Franklin T.	Sugarland
Wong, Franklin T.	Houston
Wong, Franklin T.	Stafford

Zelaya, Alicia J.	Stafford
Jabbar, Ahmad V.	Stafford
Wallace, Jennifer S.	Stafford
Wallace, Jennifer S.	Houston
Borg,James E.	Houston

EMP_PROJ1

SSN	PNUMBER	HOURS	PNAME	PLOCATION
123456789	1	32.5	Product X	Bellaire
123456789	2	7.5	Product Y	Sugarland
666884444	3	40.0	Product Z	Houston
453453453	1	20.0	Product X	Bellaire
453453453	2	20.0	Product Y	Sugarland
333445555	2	10.0	Product Y	Sugarland
333445555	3	10.0	Product Z	Houston
333445555	10	10.0	Computerization	Stafford
333445555	20	10.0	Reorganization	Houston
999887777	30	30.0	Newbenefits	Stafford
999887777	10	10.0	Computerization	Stafford
987987987	10	35.0	Computerization	Stafford
987987987	30	5.0	Newbenefits	Stafford
987654321	30	20.0	Newbenefits	Stafford
987654321	20	15.0	Reorganization	Houston
888665555	20	null	Reorganization	Houston

SSN	PNUMBER	HOURS	PNAME	PLOCATION	
123456789	1	32.5	ProductX	Bellaire	Smith,John B.
* 123456789	1	32.5	ProductX	Bellaire	English,Joyce A.
123456789	2	7.5	ProductY	Sugarland	Smith,John B.
* 123456789	2	7.5	ProductY	Sugarland	English,Joyce A.
* 123456789	2	7.5	ProductY	Sugarland	Wong,Franklin T.
666884444	3	40.0	ProductZ	Houston	Narayan,Ramesh K.
* 666884444	3	40.0	ProductZ	Houston	Wong,Franklin T.
* 453453453	1	20.0	ProductX	Bellaire	Smith,John B.
453453453	1	20.0	ProductX	Bellaire	English,Joyce A.
* 453453453	2	20.0	ProductY	Sugarland	Smith,John B.
453453453	2	20.0	ProductY	Sugarland	English,Joyce A.
* 453453453	2	20.0	ProductY	Sugarland	Wong,Franklin T.
* 333445555	2	10.0	ProductY	Sugarland	Smith,John B.
* 333445555	2	10.0	ProductY	Sugarland	English,Joyce A.
333445555	2	10.0	ProductY	Sugarland	Wong,Franklin T.
* 333445555	3	10.0	ProductZ	Houston	Narayan,Ramesh K.
333445555	3	10.0	ProductZ	Houston	Wong,Franklin T.
333445555	10	10.0	Computerization	Stafford	Wong,Franklin T.
* 333445555	20	10.0	Reorganization	Houston	Narayan,Ramesh K.
333445555	20	10.0	Reorganization	Houston	Wong,Franklin T.

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Approaches of Database Design

◆ Approaches

– Bottom up design

- Consider the basic relationships among individual attributes as the starting point
- Not very popular in practice
- Suffers from the problem of collecting a large number of binary attribute relationships
- Design by **synthesis**

– Top-down design

- Start with a number of groupings of attributes into relations obtained from conceptual design
- Further **decomposition** until all desirable properties are met
- Design by **analysis**

Normalization

◆ Normalization

- Proposed by Codd
- Proceeds in a **top-down** fashion by evaluating each relation against the **criteria** for normal forms & **decomposing relations** as necessary
- Relation design by analysis

◆ Normal forms

- 1NF (First normal form)
 - 2NF (Second normal form)
 - 3NF (Third normal form)
 - BCNF (Boyce-Codd normal form)
 - 4NF (Fourth normal form): **multi-value dependency**
 - 5NF (Fifth normal form): **join-dependency**
- } **functional dependencies**
among attributes

Normalization (cont.)

◆ Normalization procedure

- A formal framework for analyzing relation schemas based on their **keys** & on the **functional dependencies** among their attributes
- To achieve
 - Minimizing redundancy
 - Minimizing the insertion, deletion & update anomalies
- Unsatisfactory relation schemas are decomposed into smaller relation schemas that meet the normal form test

Normalization (cont.)

- ◆ **Superkey**
 - ◆ **Key**: minimal superkey
(removal of any attribute will not be a superkey)
 - ◆ **Candidate key**
 - ◆ **Primary key**
 - ◆ **Prime attribute** of R
 - a member of some **candidate key** of R
- choose one (primary key is also a candidate key)

EMPLOYEE

ENAME	<u>SSN</u>	BDATE	ADDRESS	DNUMBER
-------	------------	-------	---------	---------

p.k.

f.k.

DEPARTMENT

DNAME	<u>DNUMBER</u>	DMGRSSN
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p.k.

f.k.

DEPT_LOCATIONS

<u>DNUMBER</u>	<u>DLOCATION</u>
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p.k.

f.k.

PROJECT

PNAME	<u>PNUMBER</u>	PLOCATION	DNUM
-------	----------------	-----------	------

p.k.

f.k.

WORKS_ON

<u>SSN</u>	<u>PNUMBER</u>	HOURS
------------	----------------	-------

p.k.

f.k.

f.k.

First Normal Form

◆ First normal form

~~(3) identical rows~~

- (1) domain of an attribute must include only **atomic** values ~~composite~~
- (2) value of any attribute in a tuple must be **a single value** from the domain of that attribute ~~multivalued~~
- Disallow multi-value attribute, composite attributes & their combinations
- Be considered to be part of the formal definition of a relation in the basic relational model

relational model \Rightarrow 1st NF

First Normal Form (cont.)

DEPARTMENT

DNAME	<u>DNUMBER</u>	DMGRSSN	DLOCATIONS
-------	----------------	---------	------------

DEPARTMENT

DNAME	<u>DNUMBER</u>	DMGRSSN	DLOCATIONS
Research	5	333445555	{Bellaire, Sugarland, Houston}
Administration	4	987654321	{Stafford}
Headquarters	1	888665555	{Houston}

multivalued
attribute



Approach 1

DEPARTMENT

DNAME	<u>DNUMBER</u>	DMGRSSN	DLOCATIONS
Research	5	333445555	{Bellaire, Sugarland, Houston}
Administration	4	987654321	{Stafford}
Headquarters	1	888665555	{Houston}



DEPARTMENT

DNAME	<u>DNUMBER</u>	DMGRSSN
Research	5	333445555
Administration	4	987654321
Headquarters	1	888665555

DEPT_LOCATIONS

<u>DNUMBER</u>	<u>DLOCATION</u>
1	Houston
4	Stafford
5	Bellaire
5	Sugarland
5	Houston

Approach 2

DEPARTMENT

DNAME	<u>DNUMBER</u>	DMGRSSN	DLOCATIONS
Research	5	333445555	{Bellaire, Sugarland, Houston}
Administration	4	987654321	{Stafford}
Headquarters	1	888665555	{Houston}



DEPARTMENT

DNAME	<u>DNUMBER</u>	DMGRSSN	<u>DLOCATION</u>
Research	5	333445555	Bellaire
Research	5	333445555	Sugarland
Research	5	333445555	Houston
Administration	4	987654321	Stafford
Headquarters	1	888665555	Houston

Approach 3

DEPARTMENT

DNAME	<u>DNUMBER</u>	DMGRSSN	DLOCATIONS
Research	5	333445555	{Bellaire, Sugarland, Houston}
Administration	4	987654321	{Stafford}
Headquarters	1	888665555	{Houston}



DName	DNumber	DMGRSSN	DLocation1	DLocation2	Dlocation3
Research	5	333445555	Bellaire	Sugarland	Houston
Administration	4	987654321	Stafford		
Headquarters	1	888665555	Houston		

EMP_PROJ

SSN	ENAME	PROJS	
		PNUMBER	HOURS

composite attribute



EMP_PROJ

SSN	ENAME	PNUMBER	HOURS
123456789	Smith, John B.	1	32.5
		2	7.5
666884444	Narayan, Ramesh K.	3	40.0
453453453	English, Joyce A.	1	20.0
		2	20.0
333445555	Wong, Franklin T.	2	10.0
		3	10.0
		10	10.0
		20	10.0
999887777	Zelaya, Alicia J.	30	30.0
		10	10.0
987987987	Jabbar, Ahmad V.	10	35.0
		30	5.0
987654321	Wallace, Jennifer S.	30	20.0
		20	15.0
888665555	Borg, James E.	20	null

EMP_PROJ

SSN	ENAME	PROJS	
		PNUMBER	HOURS



EMP_PROJ1

<u>SSN</u>	ENAME
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EMP_PROJ2

<u>SSN</u>	<u>PNUMBER</u>	HOURS
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Second Normal Form

◆ Full functional dependency

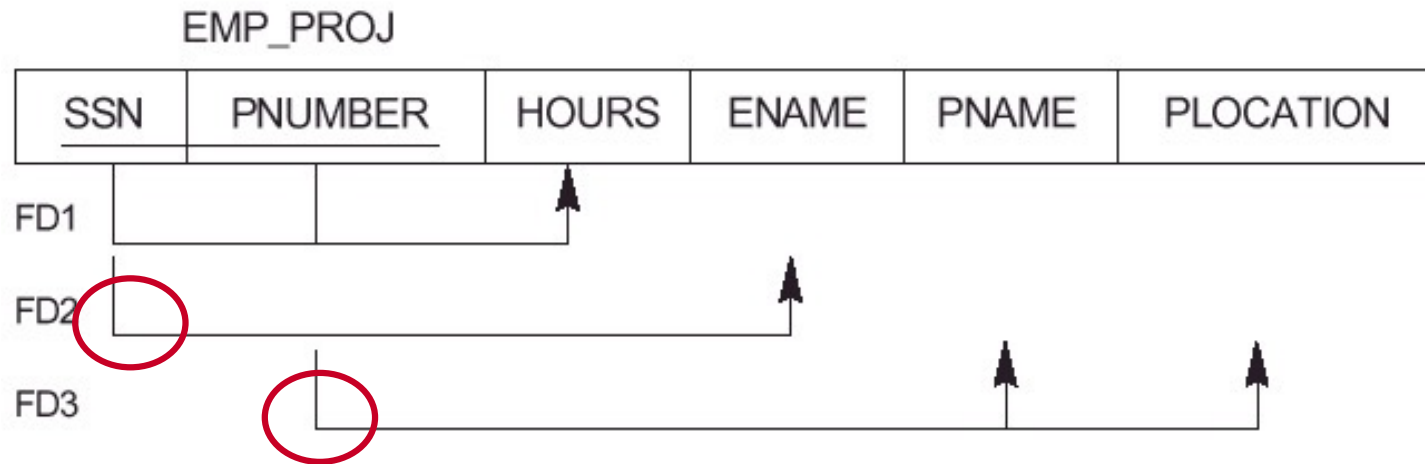
- $X \rightarrow Y$ is a **full functional dependency**

if removal of any attribute A from X means that the dependency does not hold any more

- $X \rightarrow Y$ is a **Partial functional dependency**

if removal of any attribute A from X means that the dependency still holds

- e.g. $\{\text{SSN}, \text{PNumber}\} \rightarrow \{\text{Hours}\}$ is a full FD
- e.g. $\{\text{SSN}, \text{PNumber}\} \rightarrow \{\text{Ename}\}$ is a partial FD
- e.g. $\{\text{SSN}, \text{Pnumber}\} \rightarrow \{\text{Pname}, \text{PLocation}\}$ is a partial FD

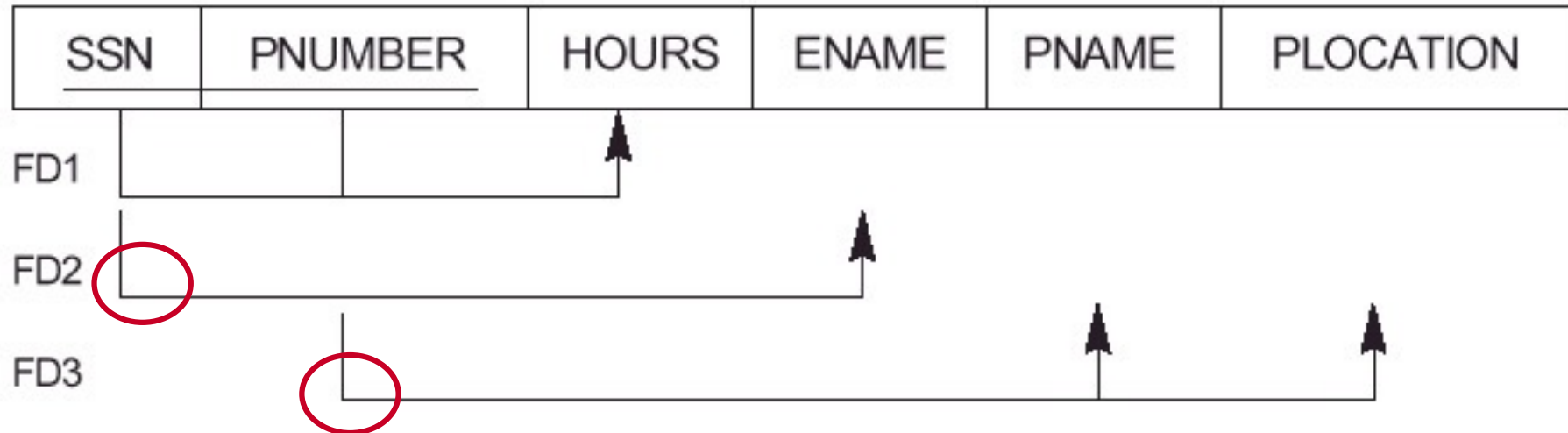


HOURS is full FD on {SSN, PNUMBER}

ENAME is partial FD on {SSN, PNUMBER}

PNAME, PLOCATION is partial FD on {SSN, PNUMBER}

EMP_PROJ



EMP_PROJ1

ENAME	SSN	PNUMBER	HOURS	PNAME	PLOCATION
Smith,John B.	123456789	1	32.5	Product X	Bellaire
Smith,John B.	123456789	2	7.5	Product Y	Sugarland
Narayan,Ramesh K.	666884444	3	40.0	Product Z	Houston
English,Joyce A.	453453453	1	20.0	Product X	Bellaire
English,Joyce A.	453453453	2	20.0	Product Y	Sugarland
Wong,Franklin T.	333445555	2	10.0	Product Y	Sugarland
Wong,Franklin T.	333445555	3	10.0	Product Z	Houston
Wong,Franklin T.	333445555	10	10.0	Computerization	Stafford
Wong,Franklin T.	333445555	20	10.0	Reorganization	Houston
Zelaya, Alicia J.	999887777	30	30.0	Newbenefits	Stafford
Zelaya, Alicia J.	999887777	10	10.0	Computerization	Stafford
Jabbar,Ahmad V.	987987987	10	35.0	Computerization	Stafford
Jabbar,Ahmad V.	987987987	30	5.0	Newbenefits	Stafford
Wallace,Jennifer S.	987654321	30	20.0	Newbenefits	Stafford
Wallace,Jennifer S.	987654321	20	15.0	Reorganization	Houston
Borg,James E.	888665555	20	null	Reorganization	Houston

Redundancy
appear in
ENAME,
PNAME,
PLOCATION

Second Normal Form (cont.)

◆ 2NF

- A relation schema R is in 2NF

 - if every **nonprime attribute** A in R

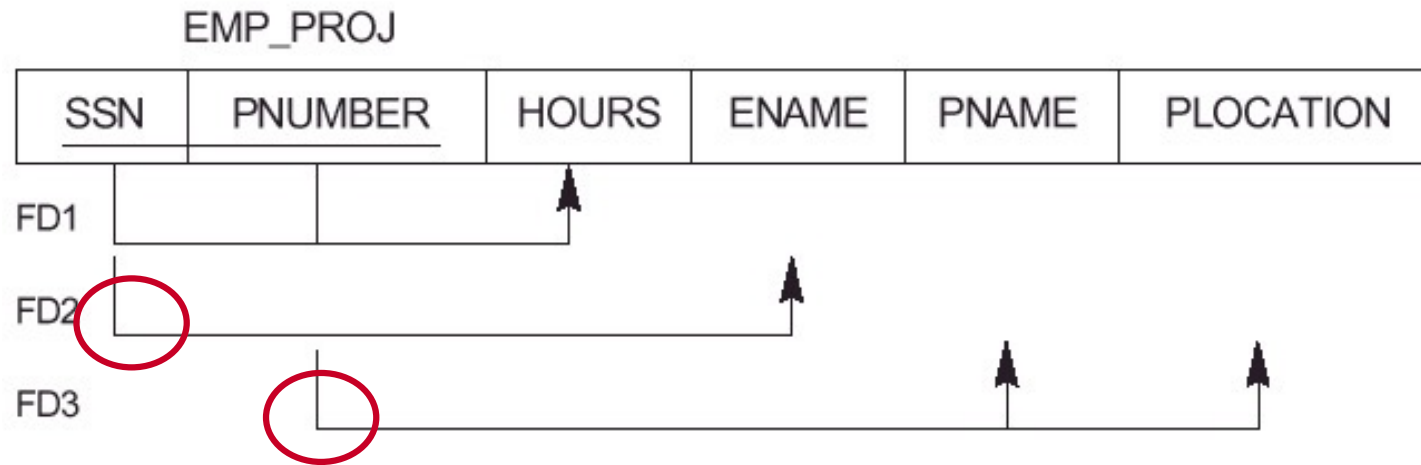
 - is **fully functional dependent** on the **primary key** of R

- If R is not in 2NF,

 - it can be **decomposed** into a number of 2NF relations
in which

 - nonprime attributes are associated only with
the part of the primary key

 - on which they are fully functional dependent



{SSN, PNUMBER} is primary key

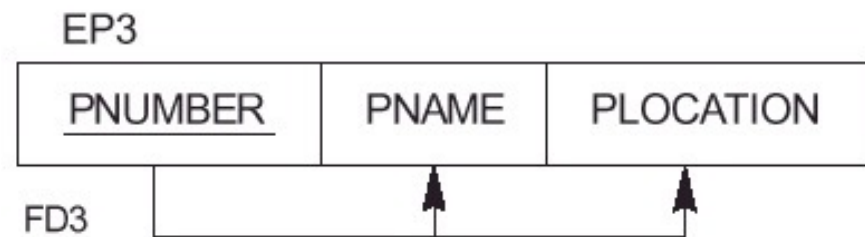
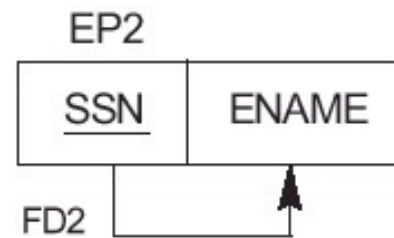
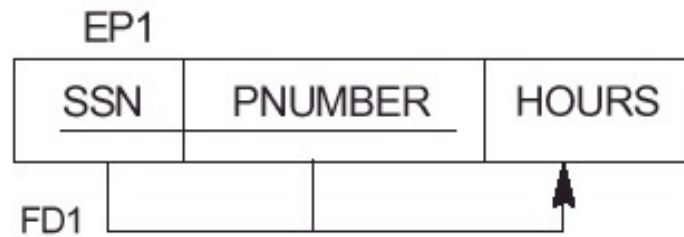
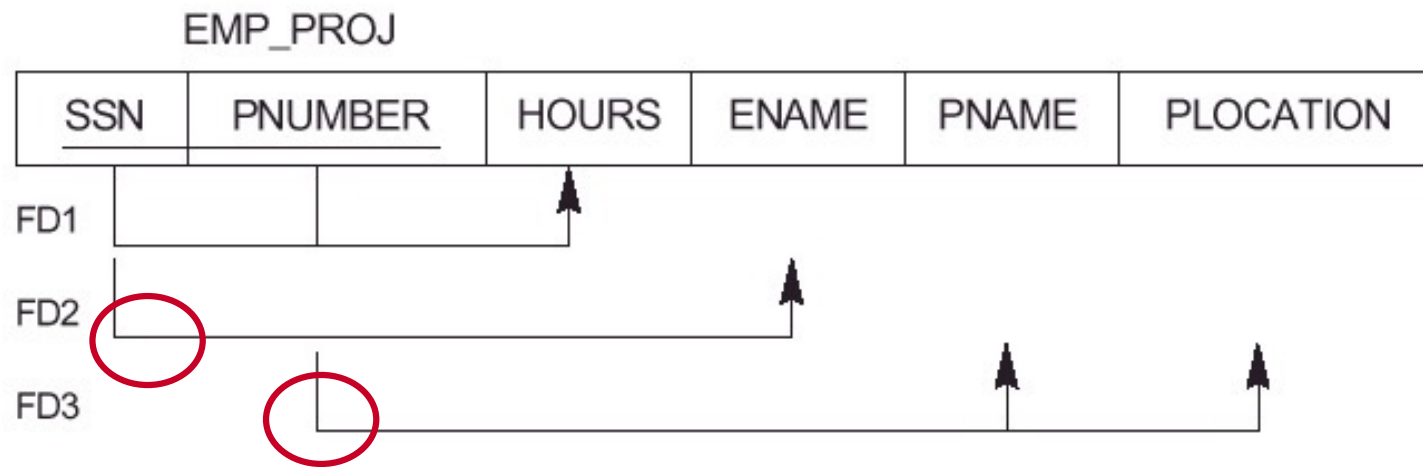
HOURS is full FD on {SSN, PNUMBER}

ENAME is partial FD on {SSN, PNUMBER}

PNAME, PLOCATION is partial FD on {SSN, PNUMBER}



EMP_PROJ is not in 2NF



EMPLOYEE

Fname	Minit	Lname	Ssn	Bdate	Address	Sex	Salary	Super_ssn	Dno
John	B	Smith	123456789	1965-01-09	731 Fondren, Houston, TX	M	30000	333445555	5
Franklin	T	Wong	333445555	1955-12-08	638 Voss, Houston, TX	M	40000	888665555	5
Alicia	J	Zelaya	999887777	1968-01-19	3321 Castle, Spring, TX	F	25000	987654321	4
Jennifer	S	Wallace	987654321	1941-06-20	291 Berry, Bellaire, TX	F	43000	888665555	4
Ramesh	K	Narayan	666884444	1962-09-15	975 Fire Oak, Humble, TX	M	38000	333445555	5
Joyce	A	English	453453453	1972-07-31	5631 Rice, Houston, TX	F	25000	333445555	5
Ahmad	V	Jabbar	987987987	1969-03-29	980 Dallas, Houston, TX	M	25000	987654321	4
James	E	Borg	888665555	1937-11-10	450 Stone, Houston, TX	M	55000	NULL	1

WORKS_ON

Essn	Pno	Hours
123456789	1	32.5
123456789	2	7.5
666884444	3	40.0
453453453	1	20.0
453453453	2	20.0
333445555	2	10.0
333445555	3	10.0
333445555	10	10.0
333445555	20	10.0
999887777	30	30.0
999887777	10	10.0
987987987	10	35.0
987987987	30	5.0
987654321	30	20.0
987654321	20	15.0
888665555	20	NULL

DEPARTMENT

Dname	Dnumber	Mgr_ssn	Mgr_start_date
Research	5	333445555	1988-05-22
Administration	4	987654321	1995-01-01
Headquarters	1	888665555	1981-06-19

DEPT_LOCATIONS

Dnumber	Dlocation
1	Houston
4	Stafford
5	Bellaire
5	Sugarland
5	Houston

DEPENDENT

Essn	Dependent_name	Sex	Bdate	Relationship
333445555	Alice	F	1986-04-05	Daughter
333445555	Theodore	M	1983-10-25	Son
333445555	Joy	F	1958-05-03	Spouse
987654321	Abner	M	1942-02-28	Spouse
123456789	Michael	M	1988-01-04	Son
123456789	Alice	F	1988-12-30	Daughter
123456789	Elizabeth	F	1967-05-05	Spouse

PROJECT

Pname	Pnumber	Plocation	Dnum
ProductX	1	Bellaire	5
ProductY	2	Sugarland	5
ProductZ	3	Houston	5
Computerization	10	Stafford	4
Reorganization	20	Houston	1
Newbenefits	30	Stafford	4

Third Normal Form

- ◆ Transitive dependency

- $X \rightarrow Y$ is a **transitive dependency**

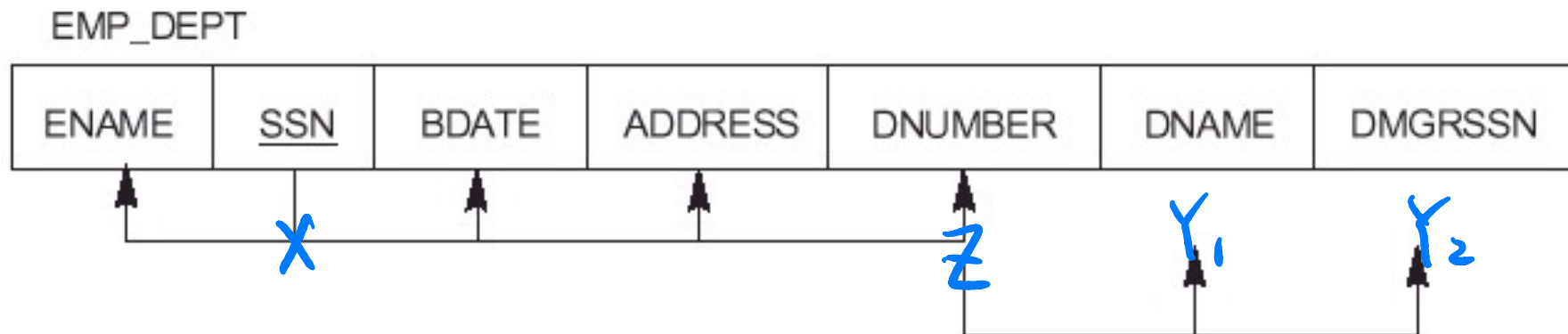
- if there is a set of attributes Z that is

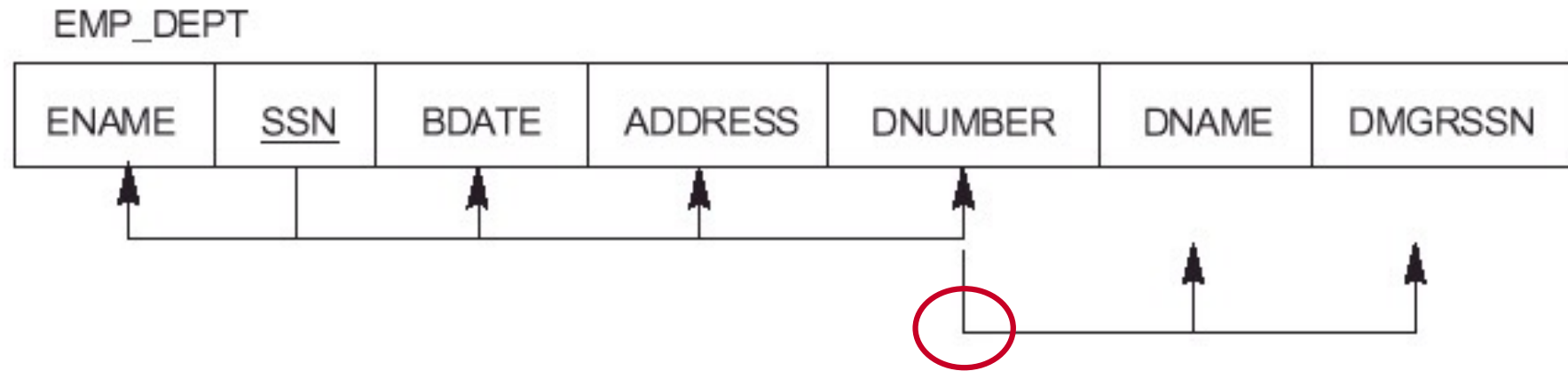
- neither a candidate key

- nor a prime attribute of R &

- $X \rightarrow Z$ and $Z \rightarrow Y$

- e.g. $\{SSN\} \rightarrow \{DNumber\} \rightarrow \{Dname, DMgrSSN\}$





$\{Dname, DMgrSSN\}$ is Transitively Dependent on $\{SSN\}$

EMP_DEPT isn't in 3NF

* $\{SSN\} \rightarrow \{DNumber\} \rightarrow \{Dname, DMgrSSN\}$

X

Z

Y

EMP_DEPT



EMP_DEPT

ENAME	<u>SSN</u>	BDATE	ADDRESS	DNUMBER	DNAME	DMGRSSN
Smith,John B.	123456789	1965-01-09	731 Fondren,Houston,TX	5	Research	333445555
Wong,Franklin T.	333445555	1955-12-08	638 Voss,Houston,TX	5	Research	333445555
Zelaya, Alicia J.	999887777	1968-07-19	3321 Castle,Spring,TX	4	Administration	987654321
Wallace,Jennifer S.	987654321	1941-06-20	291 Berry,Bellaire,TX	4	Administration	987654321
Narayan,Ramesh K.	666884444	1962-09-15	975 FireOak,Humble,TX	5	Research	333445555
English,Joyce A.	453453453	1972-07-31	5631 Rice,Houston,TX	5	Research	333445555
Jabbar,Ahmad V.	987987987	1969-03-29	980 Dallas,Houston,TX	4	Administration	987654321
Borg,James E.	888665555	1937-11-10	450 Stone,Houston,TX	1	Headquarters	888665555

Third Normal Form (cont.)

- ◆ 3NF

- A relation schema R is in 3NF

- if it satisfies 2NF &

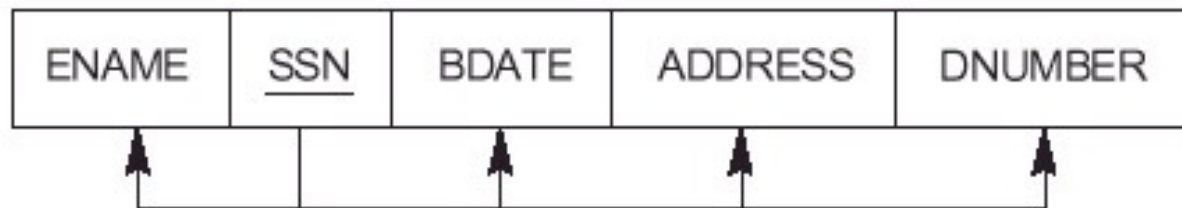
- no nonprime attribute of R

- is transitively dependent on the primary key

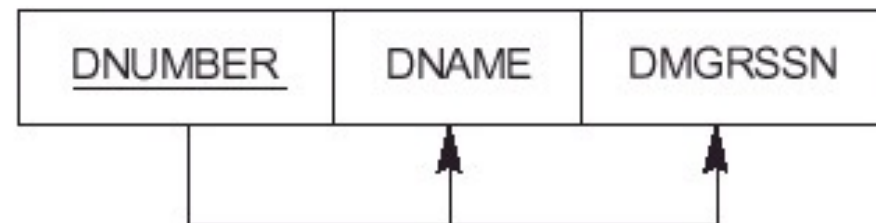
EMP_DEPT



ED1



ED2



EMPLOYEE

Fname	Minit	Lname	Ssn	Bdate	Address	Sex	Salary	Super_ssn	Dno
John	B	Smith	123456789	1965-01-09	731 Fondren, Houston, TX	M	30000	333445555	5
Franklin	T	Wong	333445555	1955-12-08	638 Voss, Houston, TX	M	40000	888665555	5
Alicia	J	Zelaya	999887777	1968-01-19	3321 Castle, Spring, TX	F	25000	987654321	4
Jennifer	S	Wallace	987654321	1941-06-20	291 Berry, Bellaire, TX	F	43000	888665555	4
Ramesh	K	Narayan	666884444	1962-09-15	975 Fire Oak, Humble, TX	M	38000	333445555	5
Joyce	A	English	453453453	1972-07-31	5631 Rice, Houston, TX	F	25000	333445555	5
Ahmad	V	Jabbar	987987987	1969-03-29	980 Dallas, Houston, TX	M	25000	987654321	4
James	E	Borg	888665555	1937-11-10	450 Stone, Houston, TX	M	55000	NULL	1

WORKS_ON

Essn	Pno	Hours
123456789	1	32.5
123456789	2	7.5
666884444	3	40.0
453453453	1	20.0
453453453	2	20.0
333445555	2	10.0
333445555	3	10.0
333445555	10	10.0
333445555	20	10.0
999887777	30	30.0
999887777	10	10.0
987987987	10	35.0
987987987	30	5.0
987654321	30	20.0
987654321	20	15.0
888665555	20	NULL

DEPARTMENT

Dname	Dnumber	Mgr_ssn	Mgr_start_date
Research	5	333445555	1988-05-22
Administration	4	987654321	1995-01-01
Headquarters	1	888665555	1981-06-19

DEPT_LOCATIONS

Dnumber	Dlocation
1	Houston
4	Stafford
5	Bellaire
5	Sugarland
5	Houston

DEPENDENT

Essn	Dependent_name	Sex	Bdate	Relationship
333445555	Alice	F	1986-04-05	Daughter
333445555	Theodore	M	1983-10-25	Son
333445555	Joy	F	1958-05-03	Spouse
987654321	Abner	M	1942-02-28	Spouse
123456789	Michael	M	1988-01-04	Son
123456789	Alice	F	1988-12-30	Daughter
123456789	Elizabeth	F	1967-05-05	Spouse

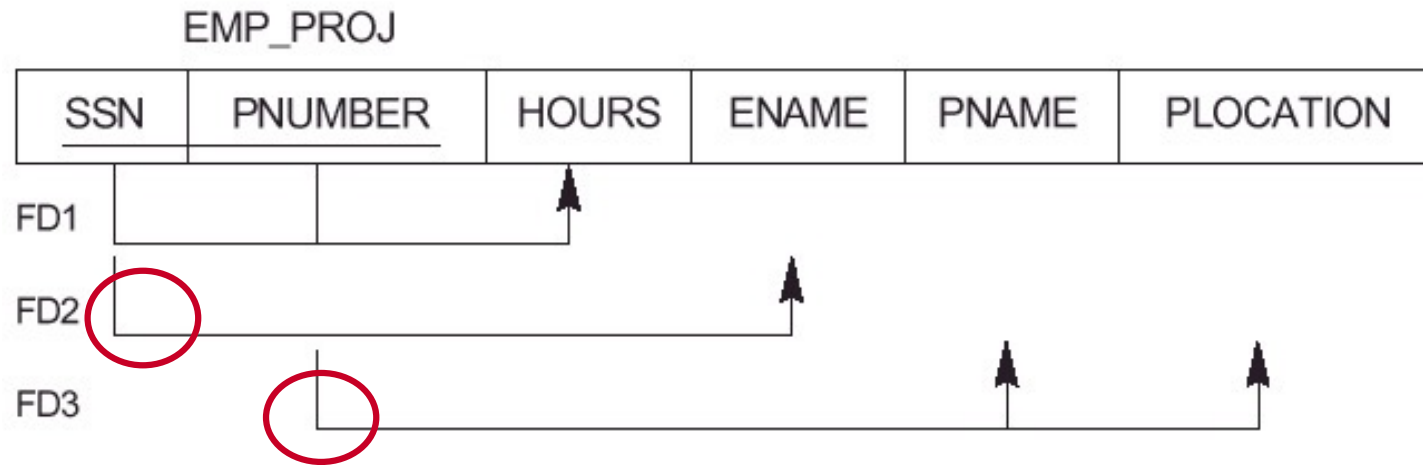
PROJECT

Pname	Pnumber	Plocation	Dnum
ProductX	1	Bellaire	5
ProductY	2	Sugarland	5
ProductZ	3	Houston	5
Computerization	10	Stafford	4
Reorganization	20	Houston	1
Newbenefits	30	Stafford	4

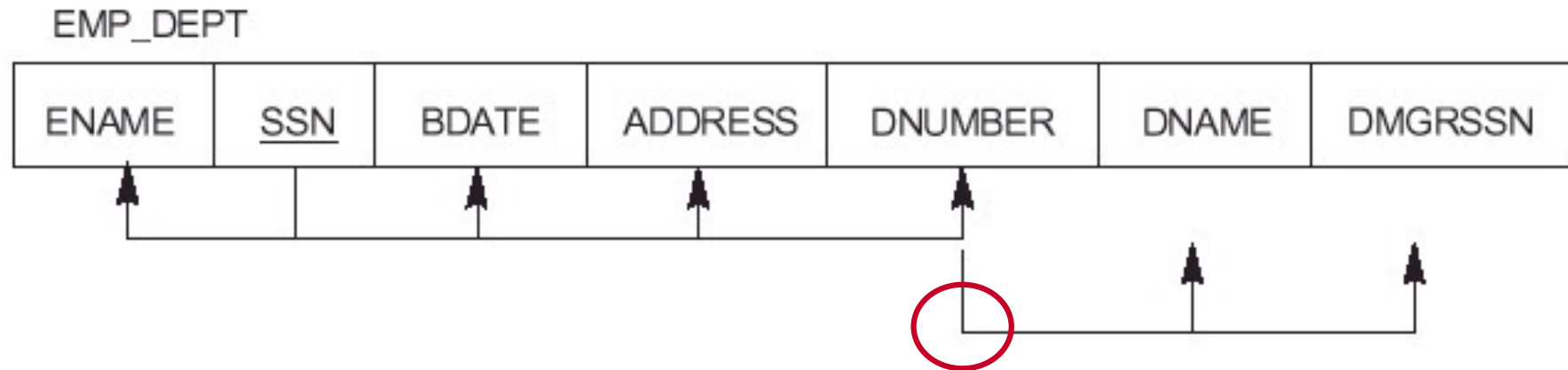
Intuitively about Normalization

- ◆ Any functional dependency in which **left-hand side** is
 - **Proper subset** of the **primary key** (breaks 2nd NF)
 - **Nonkey** attribute (breaks 3rd NF)is a **problematic** functional dependency.

- * Proper subset of primary key: **partial FD**
- * Nonkey attribute: **transitive dependency**



Left hand side of FD is proper subset of the primary key



Left hand side of FD is nonkey attribute

General Definitions of Second & Third Normal Forms

2nd & 3rd Normal Form

建立在與 **Primary Key** 的 Functional Dependency 上，
那如果是 **Candidate Key** 的 Functional Dependency 呢？



General Definition of 2NF

◆ 2NF

- A relation schema R is in 2NF

if every nonprime attribute A in R is

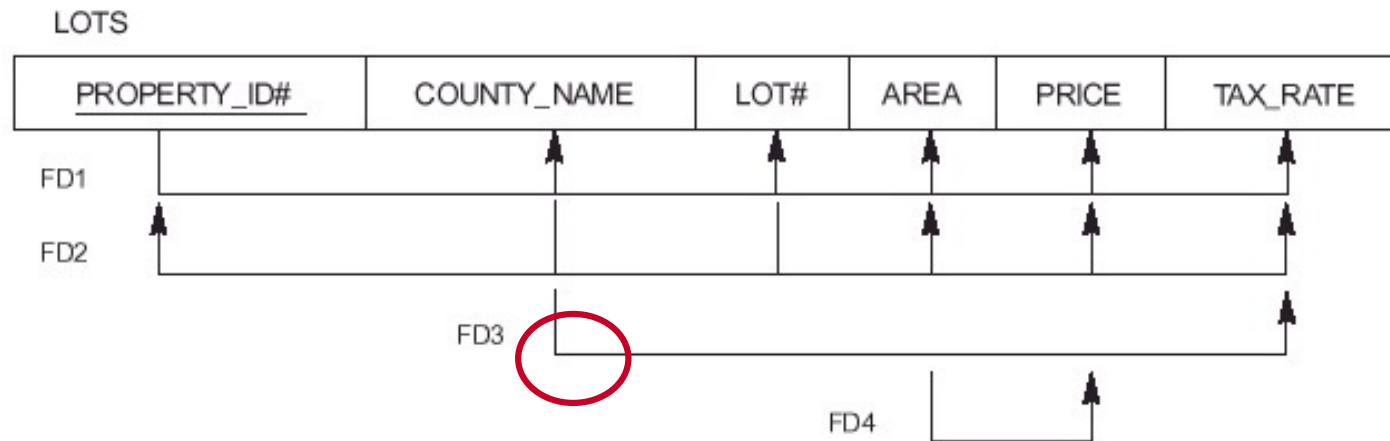
fully functional dependent on any key of R

* Original definition

- A relation schema R is in 2NF

if every nonprime attribute A in R is

fully functional dependent on the primary key of R



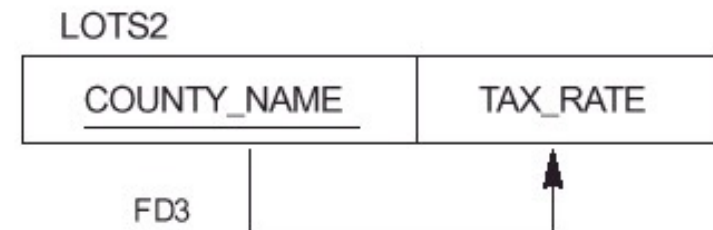
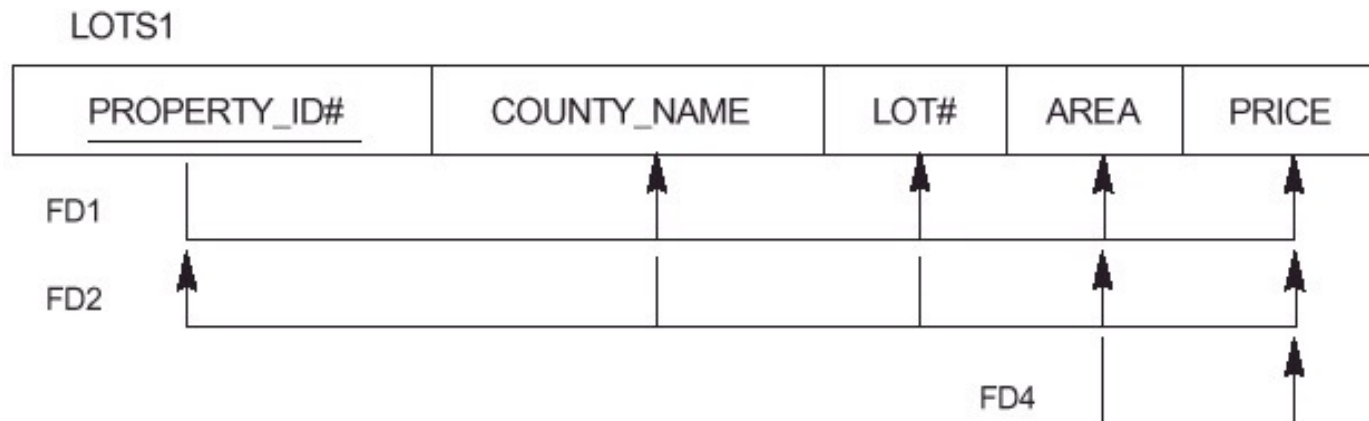
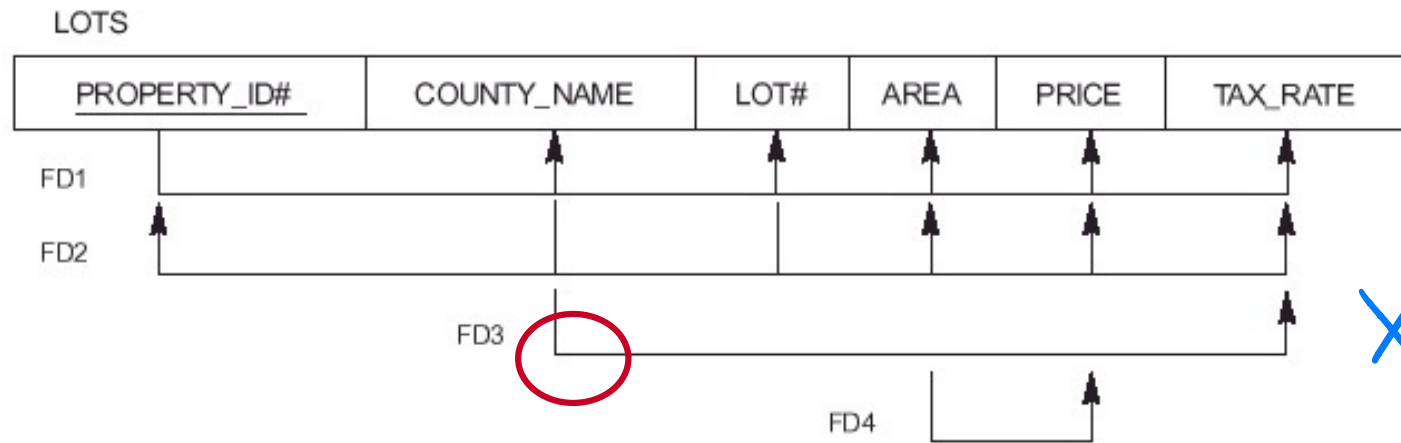
{RPROPERTY_ID#} is primary key

{COUNTY_NAME, LOT#} is a candidate key

TAX_RATE is partial FD on {COUNTY_NAME, LOT#}



LOTS is not in 2NF



General Definition of 3NF

◆ 3NF

- A relation schema R is in 3NF

- if it satisfies **2NF** &

- no nonprime attribute of R

- is transitively dependent on **any key** of R

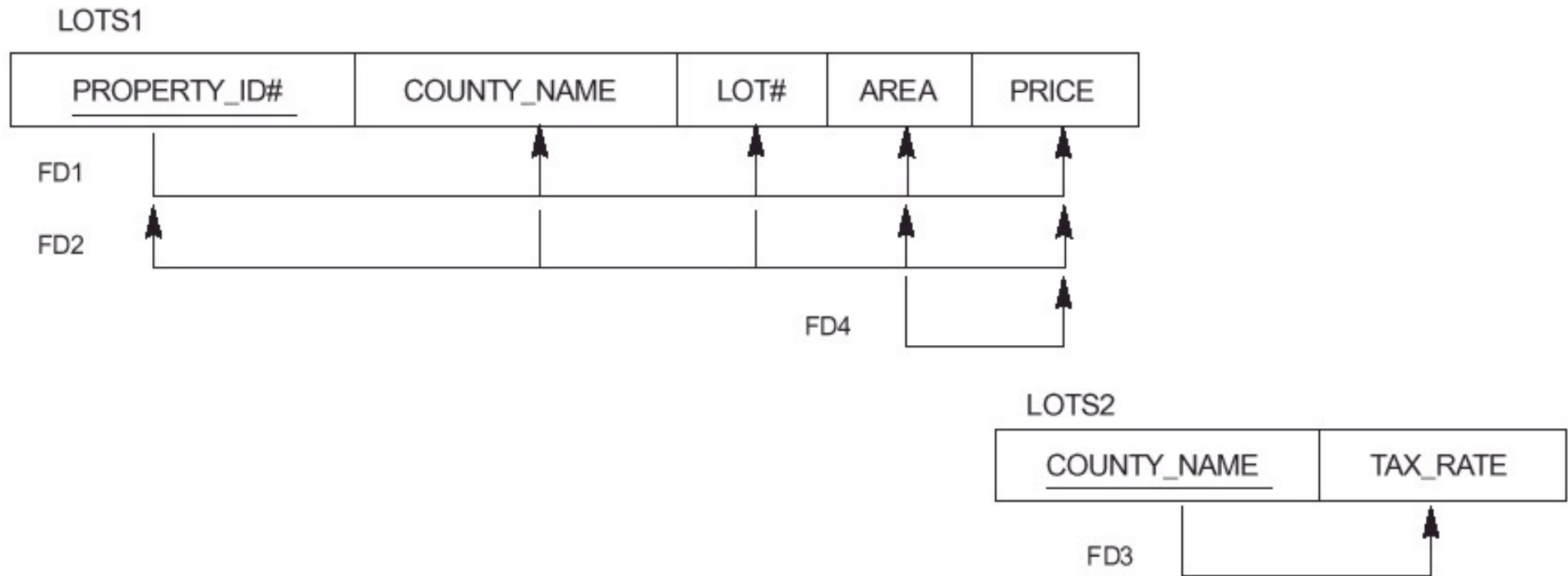
- * Original definition

- A relation schema R is in 3NF

- if it satisfies 2NF &

- no nonprime attribute of R

- is transitively dependent on **primary key** of R



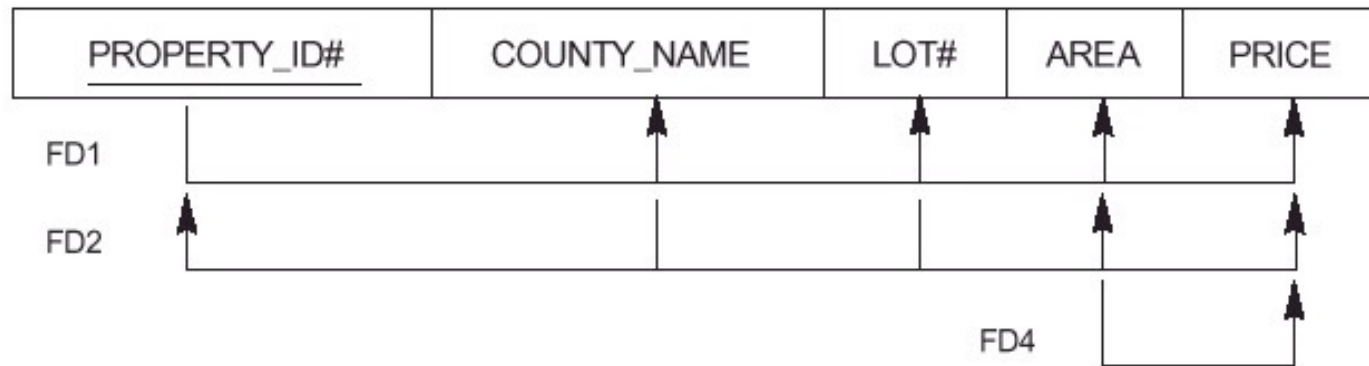
$\{COUNTY_NAME, LOT\#\}$ is a candidate key

PRICE is transitive depend on $\{COUNTY_NAME, LOT\#\}$

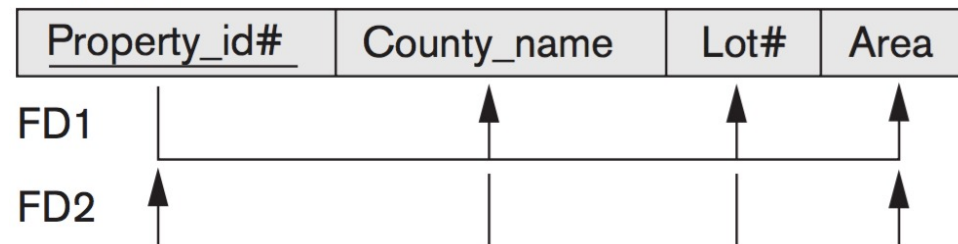


LOTS1 isn't in 3 NF

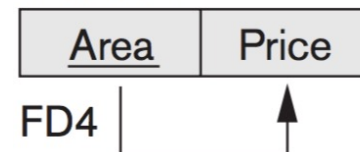
LOTS1



LOTS1A



LOTS1B



Alternative Definition of 3NF

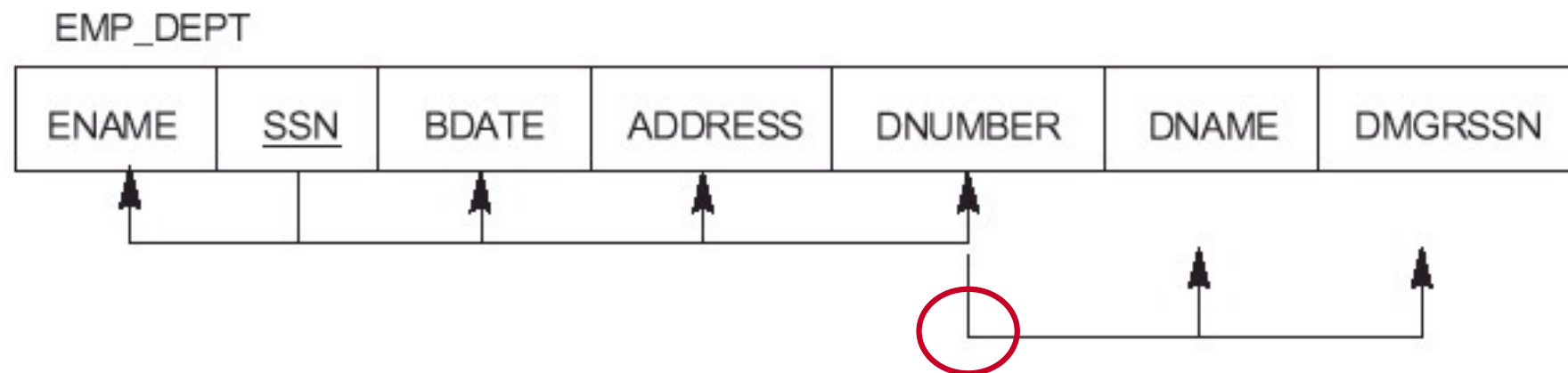
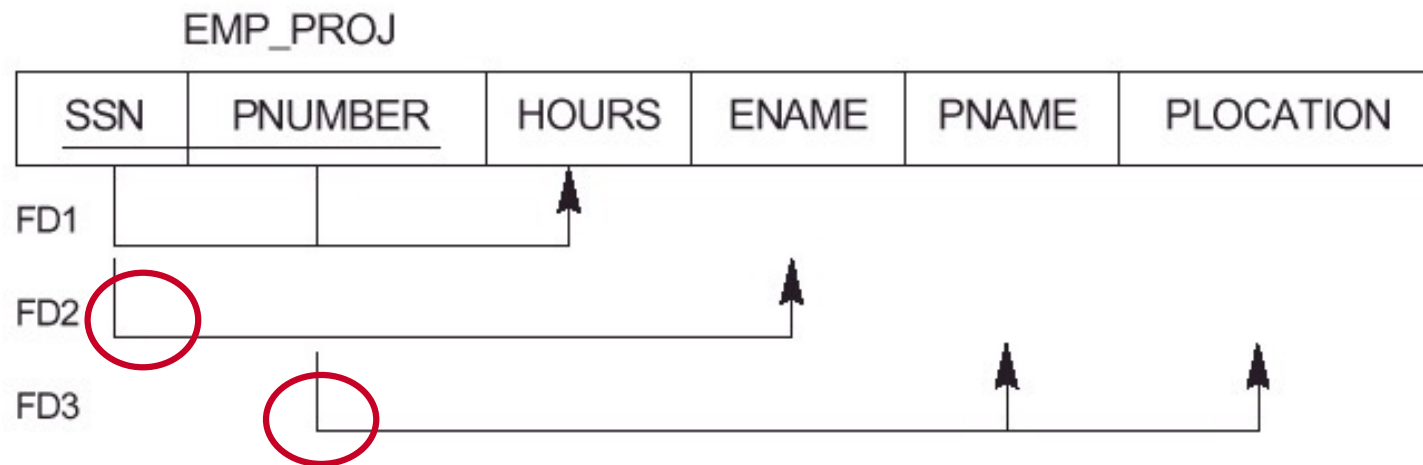
- ◆ Alternative definition of 3NF
 - A relation schema R is in 3NF if every **nonprime** attribute of R meets both of the following conditions
 - It is **fully functionally dependent** on **every key** of R (2nd NF)
 - It is **non-transitively dependent** on **every key** of R (3rd NF)

Intuitively about Normalization



- ◆ Any functional dependency in which **left-hand** side is
 - **Proper subset** of the **key**
 - **Nonkey** attributeis a **problematic** functional dependency.

- * Proper subset of key: **partial FD**
- * Nonkey attribute: **transitive dependency**



Boyce-Codd Normal Form

Boyce-Codd Normal Form

◆ BCNF

- Simpler form of 3NF
- Stricter than 3NF
- Every relation in BCNF is also in 3NF
- A relation in 3NF is not necessarily in BCNF
- A relation schema R is in BCNF
if whenever a non-trivial functional dependency $X \rightarrow A$
holds in R , then
 X is a **superkey** of R

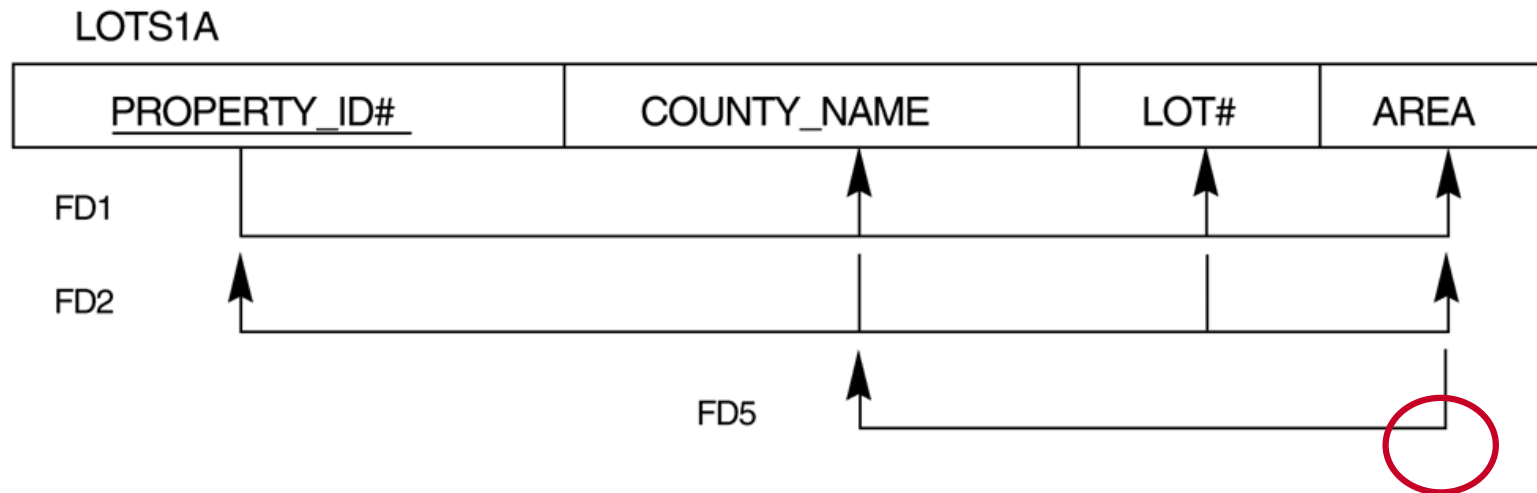
Boyce-Codd Normal Form (cont.)

◆ BCNF

- A relation schema R is in BCNF

if whenever a non-trivial functional dependency $X \rightarrow A$ holds in R , then

X is a **superkey** of R



Assume that {AREA} → {COUNTY_NAME}

Area {1, 3, 5}: King county

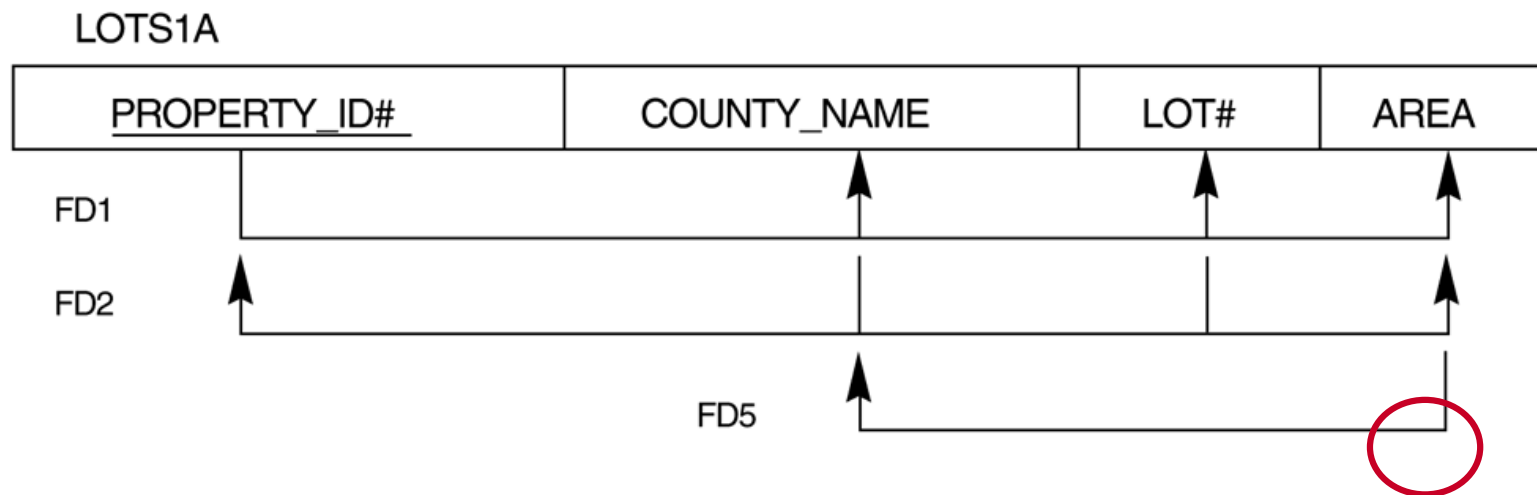
Area {2, 4, 6}: Queen county

LOTS1A is in 3 NF

(FD1, FD2, FD5 do not violate general definition of 3NF)
but the redundancy appear in COUNTY_NAME

Property_ID#	County_Name	Lot#	Area
123456	Queen	123	6
123457	Queen	125	6
123458	Queen	136	2
123459	Queen	236	6
123460	King	125	1
123461	King	288	3
123462	King	369	1

(a)



BCNF Normalization

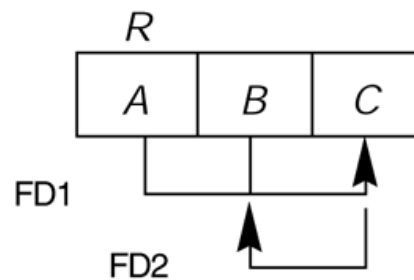
LOTS1AX

<u>PROPERTY_ID#</u>	AREA	LOT#
---------------------	------	------

LOTS1AY

<u>AREA</u>	COUNTY_NAME
-------------	-------------

(b)



Property_ID#	Lot#	Area
123456	123	6
123457	125	6
123458	136	2
123459	236	6
123460	125	1
123461	288	3
123462	369	1

Area	County_Name
6	Queen
4	Queen
2	Queen
1	King
3	King
5	King

Relationship between BCNF & 3NF

◆ BCNF

- A relation schema R is in BCNF

if whenever a non-trivial functional dependency $X \rightarrow A$ holds in R , then

X is a **superkey** of R

* 3NF

- A relation schema is in 3NF

if whenever a nontrivial functional dependency $X \rightarrow A$ holds in R , either

(a) X is a superkey of R **or**

~~(b) A is a prime attribute of R~~

General Definition of 3NF

◆ 3NF

– A relation schema is in 3NF if whenever a nontrivial functional dependency $X \rightarrow A$ holds in R, either

(a) X is a **superkey** of R

(b) A is a **prime attribute** of R

* Prime attribute: an attribute that is part of any candidate key

* General Definition

– A relation schema R is in 3NF

if every **nonprime** attribute of R meets both of the following conditions

- It is **fully functionally dependent** on **every key** of R
- It is **non-transitively dependent** on **every key** of R

General Definition of 3NF (cont.)

◆ 3NF

- A relation schema is in 3NF

if whenever a nontrivial functional dependency $X \rightarrow A$ holds in R , either

- (a) X is a superkey of R **or**
- (b) A is a prime attribute of R

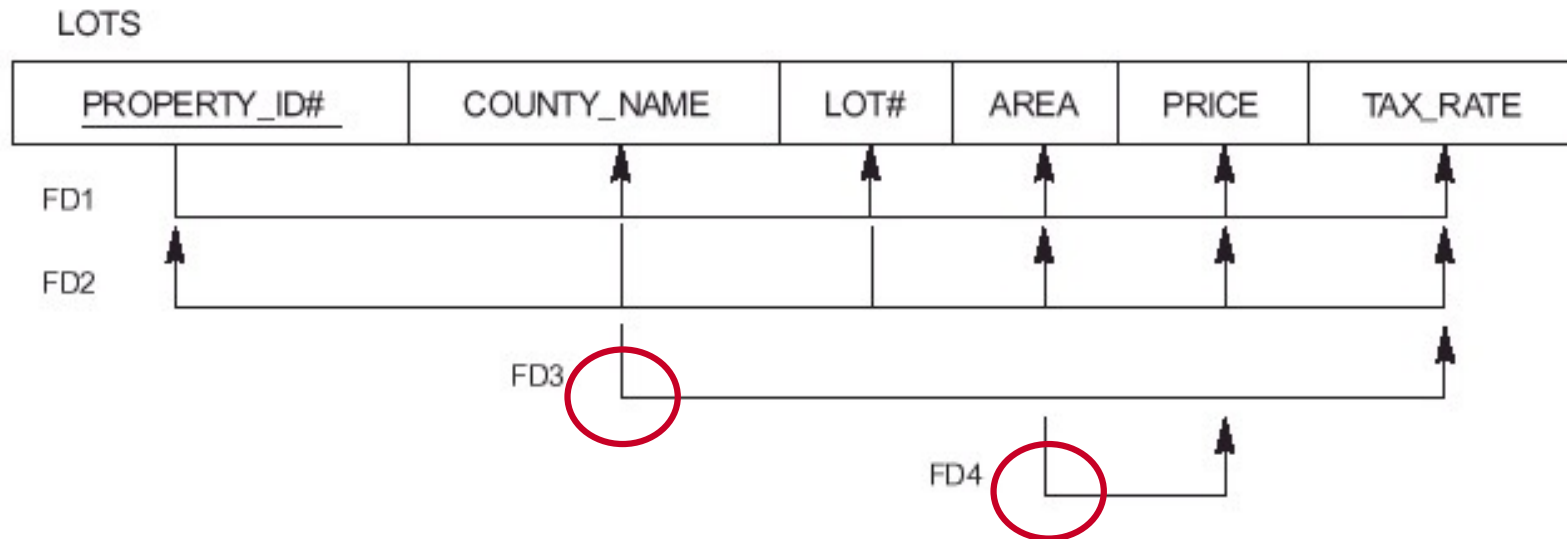
- A relation schema R **violates** 3NF

if a functional dependency $X \rightarrow A$ holds in R that **violates** both

- (a) X is not superset of any key of R **and**
- (b) A is a nonprime attribute

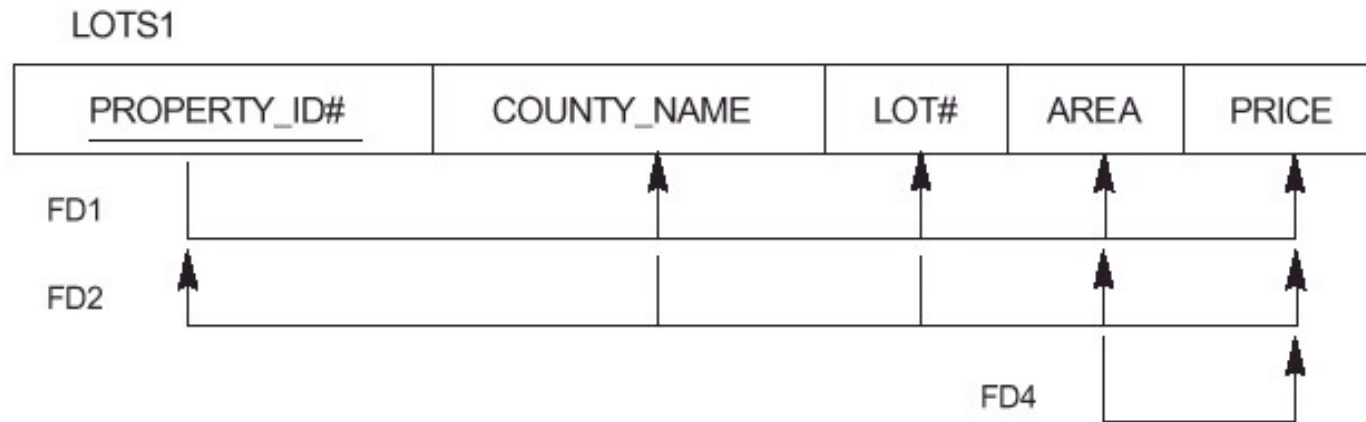
General Definition of 3NF (cont.)

- A relation schema R **violates** 3NF
if a functional dependency $X \rightarrow A$ holds in R that both
 - (a) X is **not superkey** of R and
 - (b) A is a **nonprime** attribute
- * X is not a superkey of R
 - => X could be
 - a **proper subset** of a key of R (**partial dependency**) or
 - **Nonprime** (**transitive dependency**)



$\{\text{COUNTY_NAME}\} \rightarrow \{\text{TAX_RATE}\}$
 COUNTY_NAME is not superkey of LOTS,
 TAX_RATE is a nonprime attribute of LOTS
 LOTS is not in 3NF

$\{\text{AREA}\} \rightarrow \{\text{PRICE}\}$
 AREA is not superkey of LOTS,
 PRICE is a nonprime attribute of LOTS
 LOTS is not in 3NF



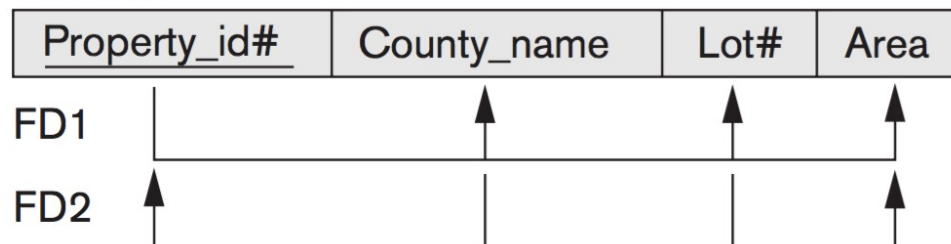
$\{\text{Area}\} \rightarrow \{\text{Price}\}$

Area is not superkey of LOTS1

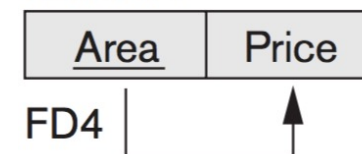
Price is a nonprime attribute of LOTS1

LOTS1 isn't in 3 NF

LOTS1A

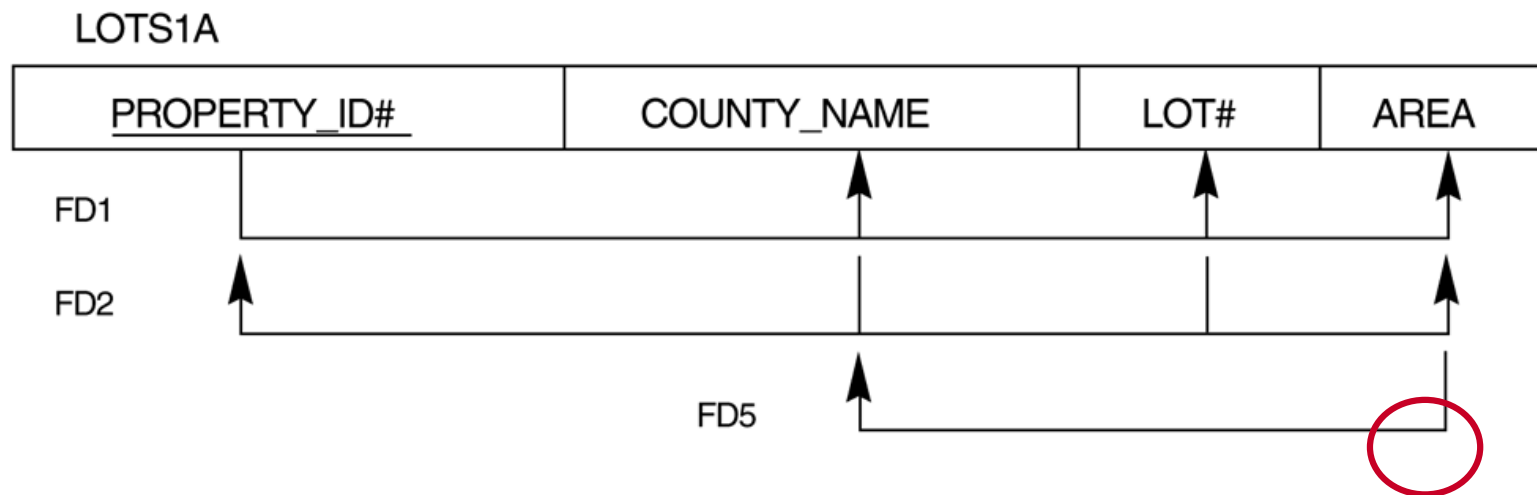


LOTS1B



LOTS1A, LOTS1B are in 3 NF

(a)



BCNF Normalization

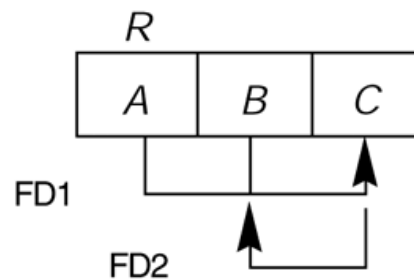
LOTS1AX

<u>PROPERTY_ID#</u>	AREA	LOT#
---------------------	------	------

LOTS1AY

<u>AREA</u>	COUNTY_NAME
-------------	-------------

(b)



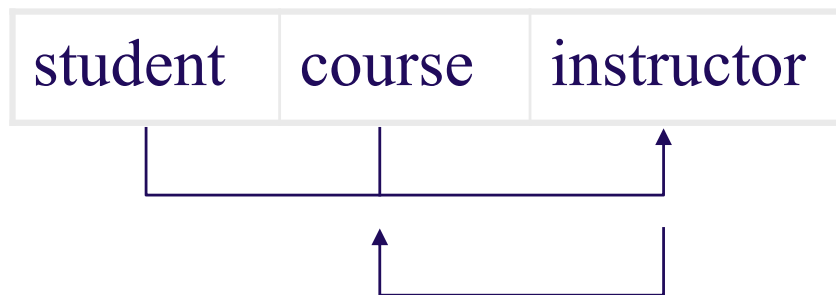
Property_ID#	Lot#	Area
123456	123	6
123457	125	6
123458	136	2
123459	236	6
123460	125	1
123461	288	3
123462	369	1

Area	County_Name
6	Queen
4	Queen
2	Queen
1	King
3	King
5	King

Another Example

◆ Given

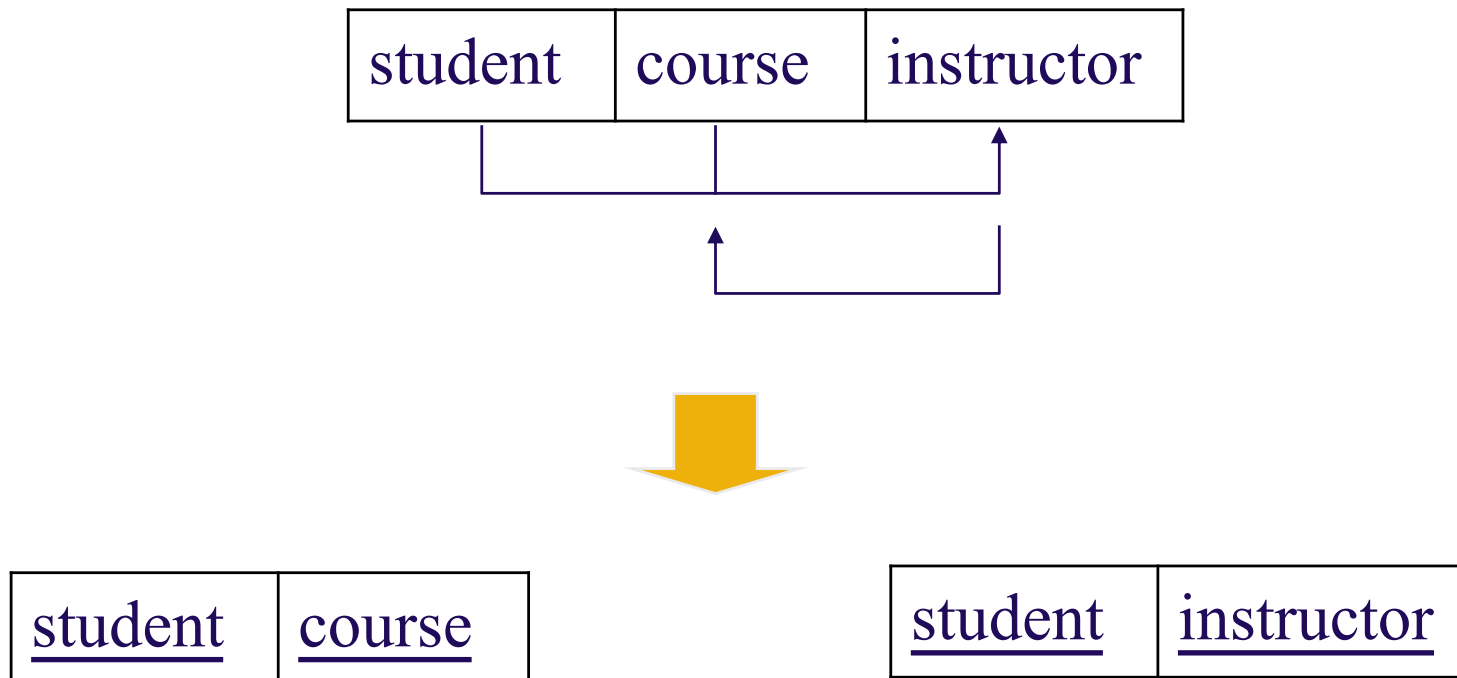
- FD1: {student, course} -> {instructor}
- FD2: {instructor} -> {course}



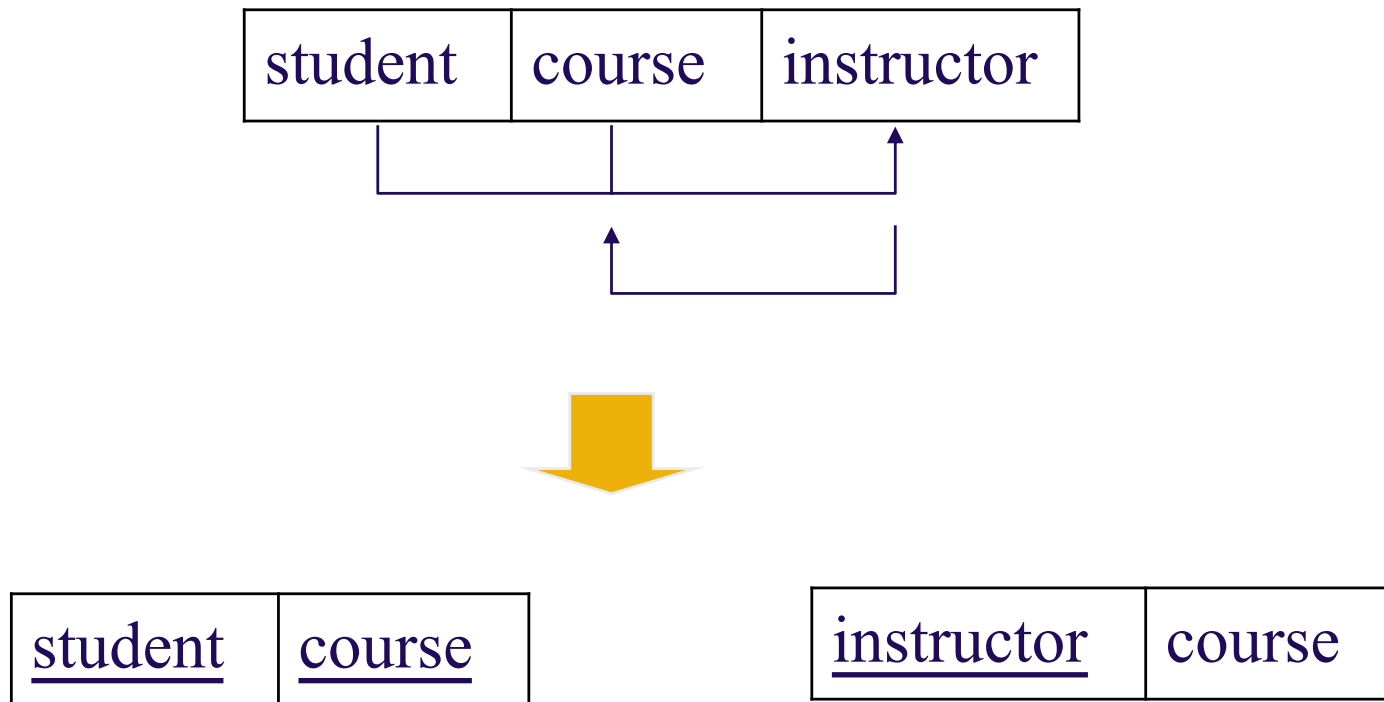
TEACH

STUDENT	COURSE	INSTRUCTOR
Narayan	Database	Mark
Smith	Database	Navathe
Smith	Operating Systems	Ammar
Smith	Theory	Schulman
Wallace	Database	Mark
Wallace	Operating Systems	Ahamad
Wong	Database	Omiecinski
Zelaya	Database	Navathe

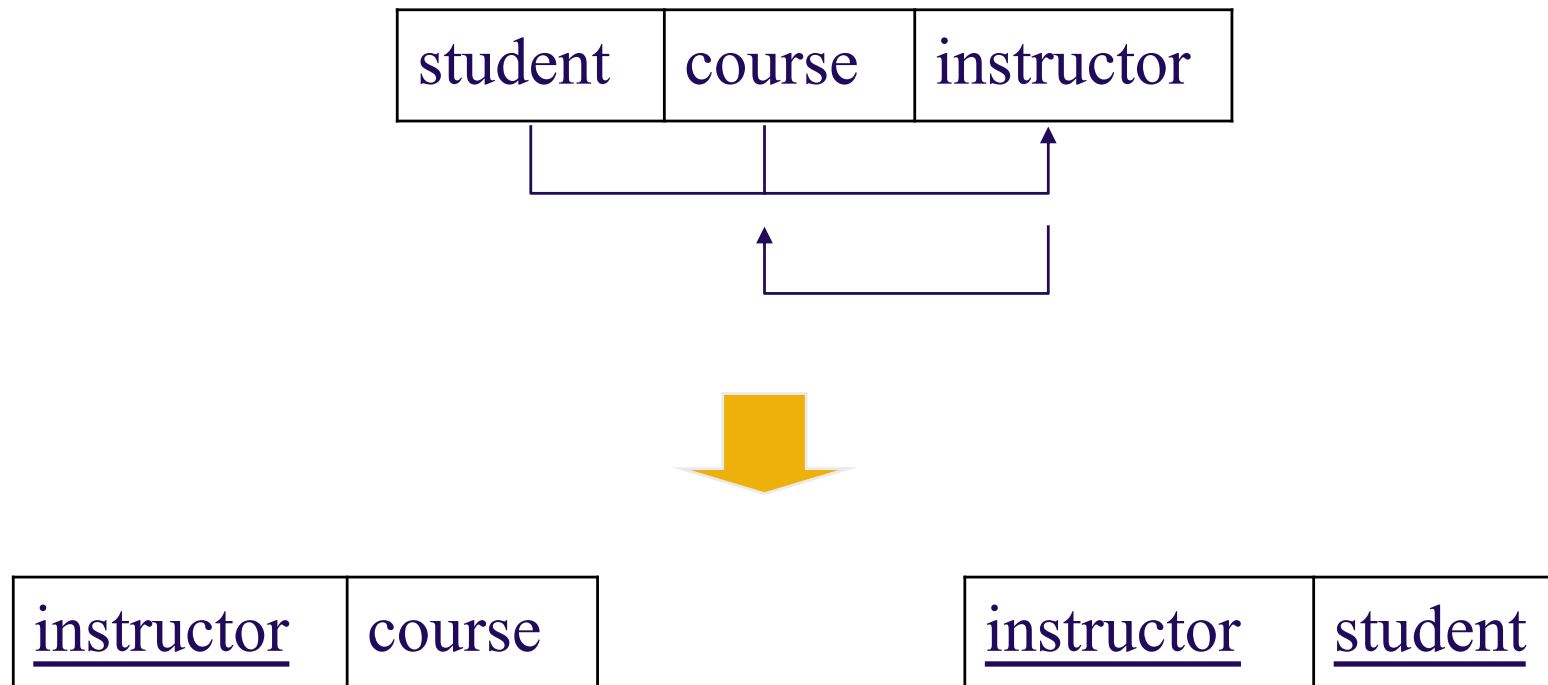
Decomposition Approach 1



Decomposition Approach 2



Decomposition Approach 3



Conclusions

- ◆ Redundancy
 - Insertion anomaly
 - Deletion anomaly
 - Modification anomaly
- ◆ Normalization: decomposition based on functional dependency
 - First Normal Form: Relational Schema
 - Second Normal Form: full functional dependency
 - Third Normal Form: non-transitive functional dependency
 - Boyce-Codd Normal Form

Normalization (cont.)

- ◆ Relations may be left in a lower normalization status for **performance** reasons
- ◆ **Denormalization**: process of storing the join of higher normal form relations as a base relation
(which is in a lower normal form)
- ◆ In general, it is advisable to
 - use normalized base relations
 - specify **views** that include the **JOINS** for placing together the attributes frequently referenced in important queries