

AMATH 301 Final Cheat Sheet that I cannot even use on the final :)

Root Finding

Methods (find roots of a function $f(x) = 0$)

MATLAB

fzero(fun,x0) tries to find a point x where $fun(x) = 0$. This solution is where $fun(x)$ changes sign. **fzero** cannot find a root of a function such as x^2 .

Bisection Method: use midpoints

- Steps: make a left and right point, find the **midpoint**, use for loop and find $f(\text{current_mid})$, the value of the function at the current x value.
- if $(\text{abs}(f_mid) < \text{tol})$ is approximately equal to zero \rightarrow root found, break;
- else if $(f_mid * \text{sin}(\text{left}) < 0) \rightarrow$ set right = mid (the root is in between the left and mid), and vice versa
- **Bisection Cons:** only works if there is **one root within the given starting range**, is somewhat slower than Newton's Method.
- **Pro:** **will always find root** if exists in between starting points

Newton's Method: use derivatives

Steps: choose an initial point, set functions and find derivative of function

$f = @(x) x + 0.5 * \sin(2 * x) - 3$; $fprime = @(x) 1 + \cos(2 * x)$;

Declare $x(k+1) = x(k) - f(x(k))/fprime(x(k))$;

if $\text{abs}(f(x(k+1))) < \text{tolerance}$, break from loop

$$x_{n+1} = x_n - \frac{f(x_n)}{f'(x_n)}$$

Cons: need to choose a **good starting point**, **need to know the derivative of the function**, **will NOT converge if on a critical point!!!**

Pros: **faster than Bisection Method** if converges

Linear Systems

Code for Jacobi:

D = diag(diag(A));

T = A-D;

M = -D\T; c = D\b;

for $k = 1:\text{max}$

$x_next = M * x + c$;

if **error < tol** \rightarrow

break

eigenvalue $|A| < 1$, the linear solving method will be **guaranteed to converge**.

Choose to use **backslash** over inverse ($\text{inv}(A)$) in most cases.

condition number:

This tells you **how much solving a linear system will magnify any noise** in your data. Will never fall below 1. An orthogonal matrix will not

Gaussian Elimination, a slow way of calculating solution for linear system, has an order of growth **$O(N^3)$** ("full price"). Ideal runtime for any iterative method: $O(N \log N)$

Jacobi Iteration: isolate all variables in the linear system

$x_{k+1} = (-15 + y_k + 5z_k)/2$

$y_{k+1} = (4x_k + z_k + 21)/8$

$z_{k+1} = (7 - 4x_k + y_k)/5$

Gauss-Seidel: enhancement of Jacobi Iteration. Use if you want a method that takes one input and does the whole process instead of individually filling out each variable like Jacobi.

A = S+T; % S is the lower triangle of A
% T is everything else

$$M = -S^{-1}T, c = S^{-1}b,$$

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while norm(x_next - x_current, Inf) >= tol && iter < maxIters
x_current = x_next;
x_next = M*x_current + c;
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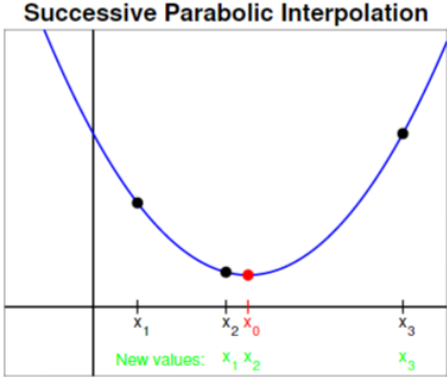
- **Gauss-Seidel Cons:** need solutions from earlier lines, **cannot be ran as parallel process (like computer with many processing cores)**
- **Pros:** much faster than Jacobi Iteration

Strictly Diagonally Dominant: all diagonal elements are strictly larger than the sum of the elements in the same row (of the matrix A). **Jacobi and Gauss-Seidel methods will be guaranteed to converge if A is SDD.** SOR method may converge regardless.

Error checking condition for iterative linear solving methods:

$$\bullet \|Ax_k - b\| \quad \bullet \|x_k - \bar{x}\| \quad \bullet \|x_{k+1} - x_k\|$$

LUP factorization: $A = L * U$; Therefore, $L * U * x = b \rightarrow y = L \backslash b$; $x = y \backslash b$;

<p>amplify any noise in your data.</p> <p>Should I use an iterative method over a direct one? → Large and sparse matrix A</p>	<ul style="list-style-type: none"> - Why use LUP (Lower Triangle, Upper Triangle, Permutation) factorization over Gaussian Elimination? In many engineering applications, when you solve $Ax=b$, the matrix A remains unchanged, while the <i>right hand side vector b keeps changing</i>. - This decouples the factorization phase (usually computationally expensive) from the actual solving phase.
<p>Min & Max Derivative-Free Methods (Optimization)</p> <p>feasibility (of regions) constrained = narrow feasibility, unconstrained = can take in any number</p> <p>fminbnd: a combination of Golden Section Search and SOR optimized to find the min value</p> <ul style="list-style-type: none"> - Only for one-dimensional function over a specific domain 	<p>Golden Section Search: does not use any derivatives, function MUST BE UNIMODAL!!! Is a lot like Bisection Method. Has variables a, b, x1, x2.</p> <p>Successive Parabolic Interpolation: pick three points, draw a curve through them.</p>  $x_0 = \frac{x_1 + x_2}{2} - \frac{(f(x_2) - f(x_1))(x_3 - x_1)(x_3 - x_2)}{2[(x_2 - x_1)(f(x_3) - f(x_2)) - (f(x_2) - f(x_1))(x_3 - x_2)]}$ <p>Get rid of side that is greater than x_0. Ex: If $x_0 < x_2 \rightarrow x_{1\text{new}} = x_{1\text{old}}; x_{2\text{new}} = x_0; x_{3\text{new}} = x_{2\text{old}};$</p> <ul style="list-style-type: none"> - Cons: Sometimes does not converge. Less reliable than Golden Section Search. - Pros: SOR can converge in fewer iterations, is faster than GSS.
<p>Gradient Descent: method that can find the minimum of functions of multiple variables.</p> <p>fminsearch(fun,x0): finds a single vector of values that will minimize a multi-dimensional function given some initial guess</p>	<p>Steps for Gradient Descent</p> <p>Outside the loop: $f = @(p) f_{xy}(p(1), p(2));$ Define the objective function with one vector input. $f_{grad} = @(p) f_{gradxy}(p(1), p(2));$ Define gradient with one vector input. $p = [-2; 0];$ Initial guess</p> <p>Inside the loop: $\mathbf{phi} = @(t) \mathbf{p} - t * \mathbf{f}_{grad}(\mathbf{p});$ Define line that points in the direction of steepest descent. $\mathbf{f_of_phi} = @(t) f(\mathbf{phi}(t));$ Function handle defining value along path. $\mathbf{grad} = \mathbf{f}_{grad}(\mathbf{p});$ Create a vector pointing in the direction, f increases the fastest. $\mathbf{tmin} = \mathbf{fminbnd}(\mathbf{f_of_phi}, 0, 1);$ Find the time it takes to reach min height on path. $\mathbf{p} = \mathbf{phi}(\mathbf{tmin});$ Gives the point on the path that is the updated guess for the minimum.</p> <p>Exam Question: What is domain and range of the minimum finding process of gradient descent? § $\mathbf{R} \rightarrow \mathbf{R}^2$, \mathbf{phi}'s codomain must be \mathbf{R}^2, \mathbf{f}_{grad} is \mathbf{R}^2 to \mathbf{R}^2 Which of the following are valid operations? $\mathbf{f}_{grad}(\mathbf{phi}(t))$ is okay (\mathbf{f}_{grad} needs an \mathbf{R}^2 input, \mathbf{phi} outputs an \mathbf{R}^2) $\mathbf{phi}(\mathbf{f}(\mathbf{p}))$ \mathbf{f} spits out something in \mathbf{R}, \mathbf{phi} wants \mathbf{R} as domain</p>

Curve Fitting

mu = center of a graph (change mu to move right or left)
- bigger mu = more to the right
sigma2 = height of peak, change to move up or down)
ex: sigma2 = 0.2 (very high), sigma2 = 1 (flatter curve)

Interpolation:

getting data point where you don't have a data point (within range of current data)

Polynomial Wiggle:

middle part is accurate, but left side and right side are very inaccurate)

Find value of polynomial at $x = 2$: $y = \text{polyval}(\text{polynomial}, 2)$;

An n -degree polynomial will go through $n+1$ points; Therefore, the error for a 2nd degree polynomial fit of 3 points will always be zero.

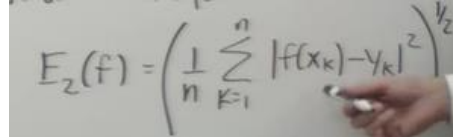
Maximal Error

- Weakness: *can't apply to huge outliers* (trying to minimize length of outlier to line), usually inaccurate in real life data
- Strengths: **fast** if points generally not outliers

Average Error

- Strength: can handle outliers very well, generally the safest method.
- Notes: **best accuracy if worried about worst case scenario, want to match that very well (try to ignore outliers)**

Root-Mean Square



- Strength: *can be thrown off by outliers. but not as much as maximal error*

Least Square Fit

The minimum is from the inside summation. **No max error, only min error!**

Sum of Squared Errors

$y = @ (x)$ blah using mu and sigma

err = sum(abs(y(X)-Y).^2);

- **this is generally well-fitted but ok with including outliers**

Exam Question: To fit a polynomial through $n+1$ data points, what degree must the polynomial be? **N degree (see below polynomial equation)**

$$y = a_n x^n + a_{n-1} x^{n-1} + \dots + a_1 x + a_0.$$

Numerical

Difference Schemes

(solving an ODE)

- error always increases over time

Error Order of

Growth: **Forward**

and Backward have

an error growth of

O(time step)

- Central and Leapfrog are second order accurate.
- **RK4 is fourth order accurate**

Second derivative

error order of growth

= **O(step^2)**

Taylor Series Expansion: The following equations are valid.

$$\checkmark f(x-h) = f(x) - hf'(x) + \frac{h^2}{2!} f''(x) - \frac{h^3}{3!} f'''(x) + \dots$$

$$f(x+h) = f(x) + hf'(x) + \frac{h^2}{2!} f''(x) + \frac{h^3}{3!} f'''(x) + \dots$$

Forward Euler

- **Explicitly solves for solution.** If you know the current point, just plug in X_k .
- UNSTABLE, if the step size is too big, the error will suddenly explode
- Otherwise usually **over-approximating** the actual solution

Backward Euler: **$(f(x) - f(x-dx))/dx$; % %** (rise over run but backwards)

- $y_{k+1} = y_k + \Delta t * f(y_{k+1})$

- **implicit** method value being determined is a function of unknown variables, need to solve an algebraic equation for the unknown, **is stable**

- can also be written as: **$x_{k+1} = \text{inv}(I - \Delta t * A) * x_k$** ; when solving $v = x'$ problem.

Central Difference: **$(f(x+dx) - (f(x-dx)))/(2*dx)$; % %** middle derivative

- is the slope between the two lines representing forward and backward

- only has odd deltas, the even ones cancel out

Leap Frog

$A = [0 \ 1; -g/L \ -r_0]$; % matrix representing change to x and v every step

$x_0 = [\theta_0; v_0]$;

<p>Local Error vs Global Error: If solving for $0 < t < T$, you take order 1/step steps. Then local error = $O(\text{step}^n)$, then global error = $O(\text{step}^{(n-1)})$</p> <p>Phase Portraits: lets you see how state variables depend on each other.</p> <ul style="list-style-type: none"> - May show “limit cycle” (predator prey reliance) 	<pre> x1 = x0 + step*A*x0; % one step of forward to calculate x at first step for t = 2*step: step: 60 x_new = x0 + 2*step*A*x1; % x0 is two before, x1 is previous x0 = x1; x1 = x_new; end </pre> <p>Using ode45 ts = 0:step:T; %% a range of times Z0 = [theta0 ; v0]; odefun2 = @(t,Z) [dxdt2(Z(1),Z(2)); dvdt2(Z(1),Z(2))]; %% t is ts, Z is initial values [t,isol] = ode45(odefun2,ts,Z0);</p> <p>- Note: ode45 becomes unstable under “stiff differential equations” situation</p> <p>Runge-Kutta ODE: RK2 (midpoint): 2nd order of growth error → more accurate than forward and backward. for k = 1: N-1; %% from first point to almost last point k1 = f(x(k)); %% x value at k k2 = f(x(k) + dt/2 * k1); x(k+1) = x(k) + dt*k2; end</p> <p>RK4: even more accurate: Do half step forward, do another half step of forward from first point, use a weighted average of the 4 slopes and take a full step x(k+1) = x(k) + dt/6 * (k1 + 2*k2 + 2*k3 + *k4); %% average of all k's</p>
<p>Numerical Integration</p> <p>MATLAB's built-in integral function: works only for function handles integral(f, 0, 8);</p>	<p>Left and Right rectangle rule (Global error: $O(\text{step})$) left_rect = dt*sum(c(1:end-1)) right_rect = dt*sum(c(2:end))</p> <p>Midpoint Rule - Global error: $O(\text{step}^2)$</p> <p>Trapezoidal rule: Global error = $O(\text{step}^2)$ trap = dt/2*(c(1) + c(end) + 2*sum(c(2:end-1)));</p> <p>Simpson's Rule → closest to true solution, Global error = $O(\text{step}^4)$ simpson = dt/3 * (c(1) + c(end) + 4*sum(c(2:2:end-1)) + 2*sum(c(3:2:end-2))); - The first and last endpoints, and 4*the even points, 2*odd points - must have EVEN NUMBER OF INTERVALS!!!</p>