

# DU\_Kronos

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```
}; 
CHT cht;
cht.init();
1.2 DC_Optimization
void compute(int L, int R, int optL, int optR){
  if(L > R) return;
  int M = L + R >> 1
  pair<ll, int> best(1LL << 60, -1);</pre>
  for(int k = optL; k \le min(M, optR); k++)
    best = min(best, \{dp[prv][k] + C[k + 1][M], k\});
  dp[now][M] = best.ff;
  compute(L, M - 1, optL, best.ss);
  compute(M + 1, R, best.ss, optR);
1.3 SOS_DP
for(int mask = 0; mask < (1<<N); ++mask){
    dp[mask][-1] = A[mask];</pre>
         for(int i = 0; i < N; ++i){
              if(mask \& (1 << i)) dp[mask][i] =
               \rightarrow dp[mask][i-1] + dp[mask ^(1<< i)][i-1];
              else dp[mask][i] = dp[mask][i-1];
         f[mask] = dp[mask][N-1];
    for(int i = 0; i<(1<<N); ++i) F[i] = A[i];
for(int i = 0;i < N; ++i)
         for(int mask = 0; mask < (1<<N); ++mask){
              if(mask & (1<<i))
    F[mask] += F[mask^(1<<i)];</pre>
1.4 dynamic cht
```

```
//add lines with -m and -b and return -ans to
//make this code work for minimums.(not -x)
const ll inf = -(1LL << 62);
struct line {
  ll m, b;
  mutable function<const line*() > succ;
  bool operator < (const line& rhs) const {</pre>
    if (rhs.b != inf) return m < rhs.m;</pre>
    const line* s = succ();
    if (!s) return 0;
    ll \dot{x} = rhs.m;
    return b - s->b < (s->m - m) * x;
struct CHT : public multiset<line> {
  bool bad(iterator y) {
    auto z = next(y);
    if (y == begin()) {
      if (z == end()) return 0;
      return y \rightarrow m == z \rightarrow m \&\& y \rightarrow b <= z \rightarrow b;
    auto x = prev(y);
    if (z == end()) return y -> m == x -> m \&\& y -> b
    return 1.0 * (x -> b - y -> b) * (z -> m - y -> m)
    \Rightarrow >= 1.0 * (y -> b - z -> b) * (y -> m - x -> m);
  void add(ll m, ll b) {
    auto y = insert({ m, b });
    y->succ = [ = ] { return next(y) == end() ? 0 :
    if (bad(y)) {
      erase(v);
      return;
```

```
while (next(y) != end() \&\& bad(next(y)))
     → erase(next(y));
    while (y != begin() \&\& bad(prev(y)))

    erase(prev(y));

  il query(ll x) {
    auto 1 = *lower_bound((line) {
      x, inf
    return l.m * x + l.b;
CHT* cht;
[ll a[N], b[N]
ios_base::sync_with_stdio(0);
  cin_tie(0):
  int n;
  cin >> n;
  for(int i = 0; i < n; i++) cin >> a[i];
for(int i = 0; i < n; i++) cin >> b[i];
cht = new CHT();
  cht -> add(-b[0], 0);
  ll ans = 0;
  for(int i = 1; i < n; i++)
    ans = -cht -> query(a[i]);
    cht -> add(-b[i], -ans);
  cout << ans << nl;
  return 0;
```

# $\begin{array}{c|c} \mathbf{2} & \mathbf{Data}_S tructure \\ \mathbf{2.1} & \mathbf{2D}_{\mathbf{-segtree}} \end{array}$

```
struct Point {
    int x, y, mx;
Point() {}
    Point(int x, int y, int mx) : x(x), y(y), mx(mx) {}
    bool operator < (const Point& other) const {</pre>
        return mx < other.mx;</pre>
struct Segtree2d {
    // I didn't calculate the exact size needed in
        terms of 2D container size.
    // If anyone, please edit the answer.
    // It's just a safe size to store nodes for MAX *
        MAX 2D grids which won't cause stack overflow
    Point T[500000]; // TODO: calculate the accurate

→ space needed
    int n, m;
    // initialize and construct segment tree
    void init(int n, int m) {
        this -> n = n;
        this -> m = m;
        build(1, 1, 1, n, m);
    // build a 2D segment tree from data [ (a1, b1),
    → (a2, b2) ]
// Time: O(n logn)
    Point build(int node, int al, int bl, int a2, int
     → b2) {
         // out of range
        if (a1 > a2 \text{ or } b1 > b2)
             return def();
        // if it is only a single index, assign value
        if (a1 == a2 and b1 == b2)
             return T[node] = Point(a1, b1, P[a1][b1]);
        // split the tree into four segments
```

```
T[node] = def();
    T[node] = maxNode(T[node], build(4 * node - 2,
    \rightarrow a1, b1, (a1 + a2) / 2, (b1 + b2) / 2) );
T[node] = maxNode(T[node], build(4 * node - 1,
         (a1 + a2) / 2 + 1, b1, a2, (b1 + b2) / 2
    T[node] = maxNode(T[node], build(4 * node + 0,
         a1, (b1 + b2) / 2 + 1, (a1 + a2) / 2, b2)
    T[node] = maxNode(T[node], build(4 * node + 1,
         (a1 + a2) / 2 + 1, (b1 + b2) / 2 + 1, a2,
     \stackrel{\square}{=} \stackrel{\backprime}{b2} );
     return T[node];
// helper function for query(int, int, int, int);
Point query(int node, int al, int bl, int a2, int
\rightarrow b2, int x1, int y1, int x2, int y2) {
    // if we out of range, return dummy
    if (x1 > a2 \text{ or } y1 > b2 \text{ or } x2 < a1 \text{ or } y2 < b1
     \rightarrow or a1 > a2 or b1 > b2)
         return def();
     // if it is within range, return the node
    if (x1 \le a1 \text{ and } y1 \le b1 \text{ and } a2 \le x2 \text{ and } b2
     \rightarrow <= v2)
         return T[node];
     // split into four segments
    Point mx = def();
    mx = maxNode(mx), query(4 * node - 2, a1, b1,
         (a1 + a2) / 2, (b1 + b2) / 2, x1, y1, x2,
     mx = maxNode(mx, query(4 * node - 1, (a1 + a2))
         /2 + 1, b1, a2, (b1 + b2) /2, x1, y1,
     \stackrel{\sim}{\rightarrow} x2, y2));
     mx = maxNode(mx, query(4 * node + 0, a1, (b1 +
         b2) / 2 + 1, (a1 + a2) / 2, b2, x1, y1,
     \stackrel{\sim}{\rightarrow} x2, y2));
     mx = maxNode(mx, query(4 * node + 1, (a1 + a2))
         /2 + 1, (b1 + b2) / 2 + 1, a2, b2, x1,
     \Rightarrow y1, x2, y2));
     // return the maximum value
     return mx;
// query from range [ (x1, y1), (x2, y2) ]
// Time: 0(logn)
Point query(int x1, int y1, int x2, int y2) {
     return query(1, 1, 1, n, m, x1, y1, x2, y2);
// helper function for update(int, int, int);
Point update(int node, int al, int bl, int a2, int

    b2, int x, int y, int value) {
     if (a1 > a2 \text{ or } b1 > b2)
         return def();
    if (x > a2 \text{ or } y > b2 \text{ or } x < a1 \text{ or } y < b1)
         return T[node];
    if (x == a1 \text{ and } y == b1 \text{ and } x == a2 \text{ and } y ==
         return T[node] = Point(x, y, value);
     Point mx = def();
     mx = maxNode(mx, update(4 * node - 2, a1, b1,
         (a1 + a2) / 2, (b1 + b2) / 2, x, y, value)
     ≒ );
     mx = maxNode(mx, update(4 * node - 1, (a1 +
         a2) / 2 + 1, b1, a2, (b1 + b2) / 2, x, y,

    value));

     mx = maxNode(mx, update(4 * node + 0, a1, (b1))
         + b2) / 2 + 1, (a1 + a2) / 2, b2, x, y,

    value)):
```

```
mx = maxNode(mx, update(4 * node + 1, (a1 +
             a2) / 2 + 1, (b1 + b2) / 2 + 1, a2, b2, x,
         return T[node] = mx;
    // update the value of (x, y) index to 'value'
    // Time: 0(logn)
    Point update(int x, int y, int value) {
   return update(1, 1, 1, n, m, x, y, value);
    // utility functions; these functions are virtual
     → because they will be overridden in child class
    virtual Point maxNode(Point a, Point b) {
         return max(a, b);
    // dummy node
    virtual Point def() {
    return Point(0, 0, -INF);
/* 2D Segment Tree for range minimum query; a override

→ of Segtree2d class */
struct Segtree2dMin : Segtree2d {
    // overload maxNode() function to return minimum
    Point maxNode(Point a, Point b) {
         return min(a, b);
    Point def() {
        return Point(0, 0, INF);
```

#### 2.2 Segtree\_beats\_desc

```
/*
Description: For update Ai Ai mod x and similar, keep
- range min,
max in node and lazily update whenever min = max. For
- update
Ai min(Ai, x) and similar, keep range max, second max
- in node and
lazily update whenever x > second max.
Time: O(log^2 N), (log N)
*/
```

#### 2.3 gp\_hash\_table

```
using namespace __gnu_pbds;
const int RANDOM = chrono::high_resolution_clock::now(_
→ ).time since epoch().count();
using namespace gnu pbds;
struct chash -
  const int RANDOM = (long
      long)(make unique<char>().get()) ^
      chrono::high resolution clock::now().time since
     epoch().count();
  static unsigned long long hash f(unsigned long long
  x x) {
x += 0x9e3779b97f4a7c15;
    x = (x \land (x >> 30)) * 0xbf58476d1ce4e5b9;

x = (x \land (x >> 27)) * 0x94d049bb133111eb;
    return x ^ (x >> 31);
  static unsigned hash combine(unsigned a, unsigned b)
     { return a * 31 + b; }
  ll operator()(ll x) const { return hash f(x)^RANDOM;
qp hash table<key, long long, chash> table;
```

```
2.4 iterative_segtree
const int N = 5000010;
int n, a[N], tree[N << 1];
void init() {
    for (int i = 0; i < n; ++i) tree[n + i] = a[i];
    for (int i = n - 1; i >= 0; --i) {
        tree[i] = min(tree[i << 1], tree[i << 1 | 1]);
    }
}
void update(int p, int v) {
    for (tree[p += n] = v; p > 1; p >>= 1) {
        tree[p >> 1] = min(tree[p], tree[p ^ 1]);
    }
}
int query(int l, int r) {
    int ret = INT_MAX;
    for (l += n, r += n; l < r; l >>= 1, r >>= 1) {
        if (l & 1) ret = min(ret, tree[l++]);
        if (r & 1) ret = min(ret, tree[--r]);
    }
} return ret;
}
```

## 2.5 mos\_algo

```
struct data{
  int l, r, id, bn;
data() {}
  data(int _l, int _r, int id){
    l = _l, _r = _r, _id = _id;
bn = l / block_sz;
  bool operator < (const data& other) const{</pre>
    if (bn != other.bn) return (bn < other.bn);</pre>
    return ((bn & 1) ? (r < other.r) : (r > other.r));
int curL = 0, curR = -1;
for(int i = 0; i < 0.sz; i++){
  while(curL > Q[i].L){
    curl--; add(curl);
  while(curR < Q[i].R){</pre>
    curR++; add(curR);
  while(curL < Q[i].L){</pre>
    remove(curL); curL++;
  while(curR > 0[i].R){
    remove(curR); curŔ--;
```

#### 2.6 ordered\_set

```
using namespace std;
using namespace __gnu_pbds;
typedef tree<
   int,
   null_type ,
   less < int > , // "less_equal<int>," for multiset
   rb_tree_tag,
   tree_order_statistics_node_update > ordered_set;
ordered_set OS;
```

#### 2.7 persistant\_segtree

```
const int MAX = 100010;
int ncnt = 0;
struct node {
  int sum;
  int left,right;
  node() {}
```

```
node(int val) {
    sum = val;
left = right = -1;
} tree[ ? ];
int ara[MAX];
int version[MAX];
void build(int n,int st,int ed) {
  if (st==ed) {
    tree[n] = node(ara[st]);
    return;
  int mid = (st+ed) / 2;
  tree[n].left = ++ncnt;
  tree[n].right = ++ncnt;
  build(tree[n].left, st, mid);
  build(tree[n].right, mid+1, ed);
  tree[n].sum = tree[tree[n].left].sum +

    tree[tree[n].right].sum;

void update(int prev,int cur,int st,int ed,int id, int
  if (id > ed or id < st) return;</pre>
  if (st == ed) {
    tree[cur] = node(val);
    return;
  int mid = (st+ed) / 2;
  if (id <= mid) ·
    tree[curl.right = tree[prevl.right:
    tree[cur].left = ++ncnt;
    update(tree[prev].left,tree[cur].left, st, mid,
    → id. val):
  else {
    tree[cur].left = tree[prev].left;
    tree[cur].right = ++ncnt;
    update(tree[prev].right, tree[cur].right, mid+1,
     → ed. id. val):
  tree[cur].sum = tree[tree[cur].left].sum +

    tree[tree[cur].right].sum;

int query(int n,int st,int ed,int i,int j){
 if(st>=i && ed<=j) return tree[n].sum;</pre>
 int mid = (st+ed)/2;
if(mid<i) return query(tree[n].right,mid+1,ed,i,j);</pre>
  else if(mid>=j) return
     query(tree[n].left,st,mid,i,j);
  else return query(tree[n].left,st,mid,i,j) +

    query(tree[n].right,mid+1,ed,i,j);

int main() {
 int n,q,l,r,k;
  sii(n,q);
  version[0] = ++ncnt;
  build(version[0],1,n);
  version[1] = ++ncnt;
  update(version[0], version[1], 1, n, id, val);
 query(version[0],1,n,id,id);
query(version[1],1,n,id,id);
  return 0;
```

```
2.8 segment_tree
int ara[MAX];
struct node {
  int sum;
}    tree[4 * MAX];
int lazy[4 * MAX];
node Merge(node a, node b) {
  node ret;
  ret.sum = a.sum + b.sum;
  return ret;
void lazyUpdate(int n, int st, int ed) {
  if(lazy[n] != 0){
    tree[n].sum += ((ed - st + 1) * lazy[n]);
   if(st != ed) {
    lazy[2 * n] += lazy[n];
    lazy[2 * n + 1] += lazy[n];
    lazv[n] = 0;
void build(int n, int st, int ed) {
  lazy[n] = 0;
  if(st == ed){
    tree[n].sum = ara[st];
    return;
  int mid = (st + ed) / 2;
  build(2 * n, st, mid);
build(2 * n + 1, mid + 1, ed);
  tree[n] = Merge(tree[2 * n], tree[2 * n + 1]);
void update(int n, int st, int ed, int i, int j, int
  lazyUpdate(n, st, ed);
  if(st > j or ed < i) return;</pre>
  if(st >= i and ed <= j){
    lazy[n] += v;
    lazyUpdate(n, st, ed);
    return;
  int mid = (st + ed) / 2;
  update(2 * n, st, mid, i, j, v);
  update(2 * n + 1, mid+1, ed, i, i, v);
  tree[n] = Merge(tree[2 * n], tree[2 * n + 1]);
node query(int n, int st, int ed, int i, int j) {
  lazyUpdate(n, st, ed);
  if(st >= i and ed <= j) return tree[n];
int mid = (st + ed) / 2;</pre>
  if (mid < i) return query(2 * n + 1, mid + 1, ed, i,
  else if(mid >= j) return query(2 * n, st, mid, i, j);
  else return Merge(query(2 * n, st, mid, i, j),
     query(2 * n + 1, mid + 1, ed, i, j));
```

### 2.9 sparse\_table

```
int st[K + 1][MAXN];
void build()
   std::copy(array.begin(), array.end(), st[0]);
  for (int i = 1; i <= K; i++)
for (int j = 0; j + (1 << i) <= N; j++)
st[i][j] = f(st[i - 1][j], st[i - 1][j + (1 <<</pre>
          \rightarrow (i - 1))]);
```

#### 3 Geometry 3.1 2D Primitive 3.1.1 Angle

A class for ordering angles (as represented by int points and a number) of rotations around the origin). Useful for rotational sweeping. Sometimes also represents points or vectors.

```
/* Usage:
   vector < Angle > v = \{w[0], w[0].t360() ...\}; //
   int i = 0; rep(i,0,n) { while (v[i] < v[i],t180().
   // sweeps j such that (j-i) represents the number
     of positively oriented triangles with vertices at
     0 and i
struct Angle {
 int x, y;
  int t:
  Angle(int x, int y, int t=0) : x(x), y(y), t(t) {}
  Angle operator-(Angle b) const { return {x-b.x,

    y-b.y, t}; }
int half() const {
    assert(x || y)
    return y < 0 | | (y == 0 \&\& x < 0);
  Angle t90() const { return \{-y, x, t + (half() \&\& x\}\}
 Angle t180() const { return {-x, -y, t + half()}; }
Angle t360() const { return {x, y, t + 1}; }
bool operator<(Angle a, Angle b)
 // add a.dist2() and b.dist2() to also compare
  return make tuple(a.t, a.half(), a.y * (ll)b.x) <</pre>

→ make tuple(b.t, b.half(), a.x * (ll)b.y);

// Given two points, this calculates the smallest
   angle between
// them, i.e., the angle that covers the defined line

→ segment.

pair<Angle, Angle> segmentAngles(Angle a, Angle b) {
  if (b < a) swap(a, b);
  return (b < a.t180() ? make pair(a, b) :

→ make pair(b, a.t360()));
Angle operator+(Angle a, Angle b) { // point a +
   vector b
  Angle r(a.x + b.x, a.y + b.y, a.t);
  if (a.t180() < r) r.t--;
return r.t180() < a ? r.t360() : r;</pre>
int tu = b.t - a.t; a.t = b.t;
  return {a.x*b.x + a.y*b.y, a.x*b.y - a.y*b.x, tu -
  \rightarrow (b < a)}:
```

#### 3.1.2 Line Distance

Returns the signed distance between point p and the line containing points a and b. Positive value on left side and negative on right as seen from a towards b. a==b gives nan. P is supposed to be Point<T> or Point3D<T> where T is e.g. double or long long. It uses products in intermediate steps so watch out for overflow if using int or long long. Using Point3D will always give a non-/s negative distance. For Point3D, call .dist on the result of the cross product.



#include "Point.h"

```
template<class P>
double lineDist(const P& a, const P& b, const P& p) {
 return (double)(b-a).cross(p-a)/(b-a).dist();
```

#### 3.1.3 Line Intersection

If a unique intersection point of the lines going through s1,e1 and s2,e2 exists {1, point} is returned. If no intersection point exists {0, (0,0)} is returned and if infinitely many exists  $\{-1, (0,0)\}$  is returned. The wrong position will be returned if P is Point<ll> and the intersection point does not have integer coordinates. Prod- \( \sigma\_{\text{S1}} \) ucts of three coordinates are used in intermediate steps so watch out for overflow if using int or ll.



```
auto res = lineInter(s1,e1,s2,e2);
   if (res.first == 1)
     cout << "intersection point at " << res.second
    << endl;
#pragma once
#include "Point.h"
template<class P>
pair<int, P> lineInter(P s1, P e1, P s2, P e2) {
 auto d = (e1 - s1).cross(e2 - s2);
 if (d == 0) // if parallel
    return \{-(s1.cross(e1, s2) == 0), P(0, 0)\};
  auto p = \hat{s}2.cross(e1, e2), q = \hat{s}2.cross(e2, \hat{s}1);
  return {1, (s1 * p + e1 * q) / d};
```

#### 3.1.4 Linear Transformation

Apply the linear transformation (translation, rotation po and scaling) which takes line p0-p1 to line q0-q1 to point q0

```
#include "Point.h"
typedef Point<double> P;
P linearTransformation(const P& p0, const P& p1, const
→ P& q0, const P& q1, const P& r) {
 P dp = p1-p0, dq = q1-q0, num(dp.cross(dq),
  → dp.dot(dq));
 return q0 + P((r-p0).cross(num),
```

#### 3.1.5 On Segment

```
/* Description: Returns true iff p lies on the line

→ segment from s to e.

* Use \texttt{(segDist(s,e,p)<=epsilon)} instead when
    using Point<double>.
#include "Point.h"
template<class P> bool onSegment(P s, P e, P p) {
 return p.cross(s, e) == 0 \& (s - p).dot(e - p)
```

#### 3.1.6 Point Sort

res

```
// sort the points in counterclockwise order that
\rightarrow starts from the half line x0, y=0.
using namespace std;
typedef long long ll;
typedef pair <ll, ll> point;
#define x first
#define y second
int main() {
  int n; cin >> n;
  vector <point> p(n);
 for (auto &it : p) scanf("%lld %lld", &it.x, &it.y);
```

```
sort(p.begin(), p.end(), [] (point a, point b) {
   return atan2l(a.y, a.x) < atan2l(b.y, b.x);
});
for (auto it : p) printf("%lld %lld\n", it.x, it.y);
return 0;
}</pre>
```

#### 3.1.7 Point

```
// Class to handle points in the plane. T can be e.g.

→ double or long long. (Avoid int.)

template <class T> int sgn(T x) \{ return (x > 0) - (x < 0) \}
\rightarrow < 0);
template<class T>
struct Point {
 typedef Point P;
 explicit Point(T x=0, T y=0) : x(x), y(y) {}
bool operator<(P p) const { return tie(x,y) <</pre>
 \rightarrow tie(p.x,p.y); }
 bool operator==(P p) const { return
 \rightarrow tie(x,y)==tie(p.x,p.y); }
 P operator+(P p) const { return P(x+p.x, y+p.y);
 P operator-(P p) const { return P(x-p.x, y-p.y); } P operator*(T d) const { return P(x*d, y*d); }
 P operator/(T d) const { return P(x/d, y/d); }
 T dot(P p) const { return x*p.x + y*p.y; }
 T cross(P'p) const { return x*p.y - y*p.x; }
 T cross(P a, P b) const { return

    (a-*this).cross(b-*this); }

 T dist2() const { return x*x + y*y;
 double dist() const { return sqrt((double)dist2()); }
 // angle to x-axis in interval [-pi, pi]
 double angle() const { return atan2(y, x); }
 P unit() const { return *this/dist(); } // makes
  \rightarrow dist()=1
 P perp() const { return P(-y, x); } // rotates +90

→ dearees

 P normal() const { return perp().unit(); }
 // returns point rotated 'a' radians ccw around the
 → origin
 P rotate(double a) const {
  return P(x*cos(a)-y*sin(a),x*sin(a)+y*cos(a)); }
friend ostream& operator<<(ostream& os, P p) {
  return os << "(" << p.x << "," << p.y << ")"; }</pre>
```

#### 3.1.8 Segment Distance

Returns the shortest distance between point p and the line segment from point s to e.

#### 3.1.9 Segment Intersection

If a unique intersection point between the line segments going from s1 to e1 and from s2 to e2 exists then it is returned. If no intersection point exists an empty vector is returned. If infinitely many exist a vector with 2 elements is returned, containing the endpoints of the common line segment. The wrong position will be returned if P is Point<|| > and the intersection point does not have integer coordinates. Products of three coordinates are used in intermediate steps so watch out for overflow if using int or long long.



```
/* Usage:
 * vector<P> inter = segInter(s1,e1,s2,e2);
 * if (sz(inter)==1)
     cout << "segments intersect at " << inter[0] <<
* Status: stress-tested, tested on kattis:intersection
#include "Point.h"
#include "OnSegment.h"
template<class P> vector<P> seqInter(P a, P b, P c, P
 d) {
  auto oa = c.cross(d, a), ob = c.cross(d, b)
ob = a.cross(b, c), od = a.cross(b, d)
  // Checks if intersection is single non-endpoint
  if (sgn(oa) * sgn(ob) < 0 \&\& sgn(oc) * sgn(od) < 0)
    return {(a * ob - b * oa) / (ob - oa)};
  set<P> s;
  if (onSegment(c, d, a)) s.insert(a);
  if (onSegment(c, d, b)) s.insert(b);
  if (onSegment(a, b, c)) s.insert(c);
  if (onSegment(a, b, d)) s.insert(d);
  return {all(s)};
```

#### 3.1.10 Side Of

Returns where p is as seen from s towards e.  $1/0/-1 \Leftrightarrow$  left/on line/right. If the optional argument eps is given 0 is returned if p is within distance eps from the line. P is supposed to be Point<T> where T is e.g. double or long long. It uses products in intermediate steps so watch out for overflow if using int or long long.

#### 3.2 3D 3.2.1 3D Convex Hull

```
#define ll long long
#define sz(x) ((int) (x).size())
#define all(x) (x).begin(), (x).end()
#define vi vector<int>
#define pii pair<int, int>
#define rep(i, a, b) for(int i = (a); i < (b); i++)
using namespace std;
template<typename T>
```

```
using minpq = priority queue<T, vector<T>, greater<T>>;
typedef long double ftype;
struct pt3 {
 ftype x, y, z;
 pt3(ftype x = 0, ftype y = 0, ftype z = 0) : x(x),
  \rightarrow y(y), z(z) {}
  pt3 operator-(const pt3 &o) const {
    return pt3(x - o.x, y - o.y, z - o.z);
  pt3 cross(const pt3 &o) const {
    return pt3(y * o.z - z * o.y, z * o.x - x * o.z, x
     \rightarrow * o.y - y * o.x);
  ftype dot(const pt3 &o) const {
    return x * 0.x + y * 0.y + z * 0.z;
// A face is represented by the indices of its three
   points a. b. c.
// It also stores an outward-facing normal vector q
struct face {
 int a, b, c;
 pt3 a:
// modify this depending on the coordinate sizes in

→ your use case

const ftype EPS = 1e-9;
vector<face> hull3(const vector<pt3> &p) {
 int n = sz(p);
  assert(n >= 3);
  vector<face> f;
  // Consider an edge (a->b) dead if it is not a CCW
     edge of some current face
  // If an edge is alive but not its reverse, this is
     an exposed edge.
  // We should add new faces on the exposed edges.
  vector<vector<bool>>> dead(n, vector<bool>(n, true));
  auto add face = [\&](int a, int b, int c)
    f.push_back({a, b, c, (p[b] - p[a]).cross(p[c] -
     → p[a])});
    dead[a][b] = dead[b][c] = dead[c][a] = false;
  // Initialize the convex hull of the first 3 points
  → as a
// triangular disk with two faces of opposite
      orientation
 add_face(0, 1, 2);
add_face(0, 2, 1);
  rep(i, 3, n) {
    // f2 will be the list of faces invisible to the

→ added point p[i]

    vector<face> f2:
    for(face &F : f)
      if((p[i] - p[F.a]).dot(F.q) > EPS) {
        // this face is visible to the new point, so

→ mark its edges as dead
        dead[F.a][F.b] = dead[F.b][F.c] =
            dead[F.c][F.a] = true;
      }else {
        f2.push back(F);
    // Add a new face for each exposed edge.
    // Only check edges of alive faces for being

→ exposed.

    f.clear();
    for(face &F : f2) {
      int arr[3] = {F.a, F.b, F.c};
      rep(j, 0, 3)
        int a = arr[j], b = arr[(j + 1) % 3];
```

```
if(dead[b][a]) {
        add face(b, a, i);
 f.insert(f.end(), all(f2));
return f;
```

#### 3.2.2 Point3D

Class to handle points in 3D space. T can be e.g. double or long long.

```
template<class T> struct Point3D {
 typedef Point3D P;
  typedef const P& R;
 T x, y, z;
  explicit Point3D(T x=0, T y=0, T z=0) : x(x), y(y),
     z(z) {}
  bool operator<(R p) const {</pre>
   return tie(x, y, z) < tie(p.x, p.y, p.z); }
  bool operator==(R p) const
   return tie(x, y, z) == tie(p.x, p.y, p.z); }
  P operator+(R p) const { return P(x+p.x, y+p.y,
     z+p.z);
 P operator (R p) const { return P(x-p.x, y-p.y,
  \hookrightarrow z-p.z); }
 P operator*(T d) const { return P(x*d, y*d, z*d); } P operator/(T d) const { return P(x/d, y/d, z/d); }
   dot(R p) const \{ return x*p.x + y*p.y + z*p.z; \}
  P cross(R p) const {
   return P(y*p.z - z*p.y, z*p.x - x*p.z, x*p.y -
    \rightarrow y*p.x);
 T dist2() const { return x*x + y*y + z*z; }
  double dist() const { return sqrt((double)dist2()); }|;
 //Azimuthal angle (longitude) to x-axis in interval
 double phi() const { return atan2(y, x); }
 //Zenith angle (latitude) to the z-axis in interval
  double theta() const { return
  \rightarrow atan2(sqrt(x*x+y*y),z); }
  P unit() const { return *this/(T)dist(); } //makes
      dist()=1
  //returns unit vector normal to *this and p
  P normal(P p) const { return cross(p).unit(); }
 //returns point rotated 'angle' radians ccw around
  P rotate(double angle, P axis) const {
    double s = sin(angle), c = cos(angle); P u =
    → axis.unit();
   return u*dot(u)*(1-c) + (*this)*c - cross(u)*s;
```

#### 3.2.3 Polyhedron Volume

Description: Magic formula for the volume of a polyhedron. Faces should point outwards.

```
template<class V, class L>
double signedPolyVolume(const V& p, const L& trilist) {
 double \bar{v} = 0;
  for (auto i : trilist) v +=

→ p[i.a].cross(p[i.b]).dot(p[i.c]);
  return v / 6;
```

#### 3.2.4 Spherical Distance

Returns the shortest distance on the sphere with radius radius be-circles are tangent to each other (in which case .first = .second and the tween the points with azimuthal angles (longitude) f1 ( $\phi_1$ ) and f2 ( $\phi_2$ ) tangent line is perpendicular to the line between the centers). first

(0 = north pole). All angles measured in radians. The algorithm starts find the tangents of a circle with a point set r2 to 0. by converting the spherical coordinates to cartesian coordinates so if that is what you have you can use only the two last rows. dx\*radius is then the difference between the two points in the x direction and d\*radius is the total distance between the points.

```
double sphericalDistance(double f1, double t1,
  double f2, double t2, double radius) {
double dx = sin(t2)*cos(f2) - sin(t1)*cos(f1);
double dy = sin(t2)*sin(f2) - sin(t1)*sin(f1);
  double dz = cos(t2) - cos(t1);
  double d = sqrt(dx*dx + dy*dy' + dz*dz);
  return radius*2*asin(d/2);
```

#### 3.3 Circle

#### 3.3.1 Circle Intersection

Computes the pair of points at which two circles intersect. Returns false in case of no intersection.

```
#include "Point.h"
typedef Point<double> P
bool circleInter(P a,P b,double r1,double r2,pair<P,</pre>
    P>* out) {
  if (a == b) { assert(r1 != r2); return false; }
   P \text{ vec} = b - a;
   double d2 = vec.dist2(), sum = r1+r2, dif = r1-r2,
              p = (d2 + r1*r1 - r2*r2)/(d2*2), h2 = r1*r1

→ p*p*d2;

   if (sum*sum < d2 || dif*dif > d2) return false;
   P \text{ mid} = a + \text{vec*p, per} = \text{vec.perp}() * \text{sqrt}(\text{fmax}(0,
   \stackrel{\hookrightarrow}{*} \stackrel{\text{h2}}{\text{out}} / \stackrel{\text{d2}}{\text{mid}} / \stackrel{\text{per}}{\text{per}}, \stackrel{\text{mid}}{\text{mid}} - \stackrel{\text{per}}{\text{per}};
   return true:
```

#### 3.3.2 Circle Polygon Intersection

Returns the area of the intersection of a circle with a ccw polygon. Time: O(n)

```
#include "Point.h"
typedef Point<double> P;
#define arg(p, q) atan2(p.cross(q), p.dot(q))
double circlePoly(P c, double r, vector<P> ps) {
  auto tri = [&](P p, P q) {
  auto r2 = r * r / 2;
    Pd = q - p;
    auto a = d.dot(p)/d.dist2(), b =
        (p.dist2()-r*r)/d.dist2();
    auto det = a * a - b;
    if (det \le 0) return arg(p, q) * r2;
    auto s = max(0., -a-sqrt(det)), t = min(1.,
     → -a+sqrt(det));
    if (t < 0 \mid | 1 \le s) return arg(p, q) * r2;
P u = p + d * s, v = p + d * t;
    return arg(p,u) * r2 + u.cross(v)/2 + arg(v.q) *
  auto sum = 0.0
  rep(i,0,sz(ps))
    sum += tri(ps[i] - c, ps[(i + 1) % sz(ps)] - c);
  return sum;
```

#### 3.3.3 Circle Tangents

Finds the external tangents of two circles, or internal if r2 is negated. Can return 0, 1, or 2 tangents – 0 if one circle contains the other (or overlaps it, in the internal case, or if the circles are the same); 1 if the

from x axis and zenith angles (latitude) t1 ( $\theta_1$ ) and t2 ( $\theta_2$ ) from z axis and .second give the tangency points at circle 1 and 2 respectively. To

```
#include "Point.h"
template<class P>
vector<pair<P, P>> tangents(P c1, double r1, P c2,

→ double r2) {
 P d = c2 - c1
  double dr = r1 - r2, d2 = d.dist2(), h2 = d2 - dr *
  if (d2' == 0 | | h2 < 0) return {};
  vector<pair<P, P>> out;
 for (double sign : {-1, 1}) {
  P v = (d * dr + d.perp() * sqrt(h2) * sign) / d2;
  out.push_back({c1 + v * r1, c2 + v * r2});
  if (h2 == 0) out.pop back();
  return out;
```

#### 3.3.4 CircumCircle

The circumcirle of a triangle is the circle intersecting all three vertices. ccRadius returns the radius of the circle going through points A, B and C and ccCenter returns the center of the same circle.



```
#include "Point.h"
typedef Point<double> P;
double ccRadius(const P& A, const P& B, const P& C) {
  return (B-A).dist()*(C-B).dist()*(A-C).dist()/
    abs((B-A).cross(C-A))/2;
  ccCenter(const P\& A, const P\& B, const P\& C) { P b = C-A, c = B-A;

→ (b*c.dist2()-c*b.dist2()).perp()/b.cross(c)/2;
```

#### 3.4 Polygon

#### 3.4.1 Hull Diameter

Returns the two points with max distance on a convex hull (ccw, no duplicate/collinear points).

```
#include "Point.h"
typedef Point<ll> P;
array<P, 2> hullDiameter(vector<P> S) {
  int n = sz(S), j = n < 2 ? 0 : 1;
pair<ll, array<P, 2>> res({0, {S[0], S[0]}});
  rep(i,0,j)
     for (;; j = (j + 1) % n) {
       res = \max(\text{res}, \{(S[i] - S[j]).dist2(), \{S[i],
       if ((S[(j+1) % n] - S[j]).cross(S[i+1] - S[j])
        \rightarrow S[i]) >= 0)
         break;
  return res.second;
```

```
3.4.2 Line Hull Intersection
```

Line-convex polygon intersection. The polygon must be ccw and have no collinear points. lineHull(line, poly) returns a pair describing the intersection of a line with the polygon:

(-1,-1) if no collision, (i,-1) if touching the corner i,

(i,i) if along side (i,i+1),

(i,j) if crossing sides (i,i+1) and (j,j+1). In the last case, if a corner i is crossed, this is treated as happening on side (i,i+1). The points are returned in the same order as the line hits the polygon. extrVertex returns the point of a hull with the max projection onto a line.

Time:  $O(\log n)$ 

```
#include "Point.h"
#define cmp(i,j)

    sgn(dir.perp().cross(poly[(i)%n]-poly[(j)%n]))
#define extr(i) cmp(i + 1, i) \Rightarrow 0 && cmp(i, i - 1 +
template <class P> int extrVertex(vector<P>& poly, P

→ dir) {
  int n = sz(poly), lo = 0, hi = n;
  if (extr(0)) return 0;
  while (lo + 1 < hi)
    int \dot{m} = (lo + hi) / 2;
    if (extr(m)) return m;
int ls = cmp(lo + 1, lo), ms = cmp(m + 1, m);
    (ls < ms \mid | (ls == ms \&\& ls == cmp(lo, m)) ? hi :
     \rightarrow lo) = m:
  return lo;
#define cmpL(i) sqn(a.cross(poly[i], b))
template <class P>
array<int, 2> lineHull(P a, P b, vector<P>& poly) {
  int endA = extrVertex(poly, (a - b).perp());
  int endB = extrVertex(poly, (b - a).perp());
  if (cmpL(endA) < 0 \mid | cmpL(endB) > 0)
    return {-1, -1};
  array<int, 2> res;
  rep(1,0,2) {
    int lo = endB, hi = endA, n = sz(poly);
    while ((lo + 1) % n != hi) {
      int \dot{m} = ((lo' + hi + (lo' < hi ? 0 : n)) / 2) % n;
      (cmpL(m) = cmpL(endB) ? lo : hi) = m;
    res[i] = (lo + !cmpL(hi)) % n;
    swap(endA, endB);
  if (res[0] == res[1]) return {res[0], -1};
  if (!cmpL(res[0]) && !cmpL(res[1]))
    switch ((res[0] - res[1] + sz(poly) + 1) %
     \rightarrow sz(poly)) {
      case 0: return {res[0], res[0]};
case 2: return {res[1], res[1]};
  return res;
```

#### 3.4.3 Polygon Center

Returns the center of mass for a polygon. Time: O(n)

#include "Point.h"
typedef Point<double> P;
P polygonCenter(const vector<P>& v) {
 P res(0, 0); double A = 0;
 for (int i = 0, j = sz(v) - 1; i < sz(v); j = i++) {
 res = res + (v[i] + v[j]) \* v[j].cross(v[i]);
}</pre>

```
A += v[j].cross(v[i]);
  return res / A / 3;
3.4.4 Polygon Cut
Returns a vector with the vertices of a polygon with ev-
erything to the left of the line going from s to e cut away.
   vector < P > p = ...;
    p = polygonCut(p, P(0,0), P(1,0));
 * Status: tested but not extensively
#include "Point.h"
#include "lineIntersection.h"
typedef Point<double> P;
vector<P> polygonCut(const vector<P>& poly, P s, P e) {
  vector<P> rés;
  rep(i,0,sz(poly)) {
    P cur = poly[i], prev = i ? poly[i-1] :
     → poly.back();
    bool side = s.cross(e, cur) < 0;</pre>
    if (side != (s.cross(e, prev) < 0))
       res.push back(lineInter(s, e, cur, prev).second);|;
      res.push back(cur);
  return res;
3.5 Closest Pair
Finds the closest pair of points.
Time: O(n \log n)
#include "Point.h"
typedef Point<ll> P;
pair<P, P> closest(vector<P> v) {
  assert(sz(v) > 1);
  set<P>
  sort(all(v), [](P a, P b) { return a.y < b.y; });</pre>
  pair<ll, pair<P, P>> ret{LLONG MAX, {P(), P()}};
  int j = 0;
  for (P p : v)
    P d{1 + (ll)sqrt(ret.first), 0};
    while (v[j].y \le p.y - d.x) S.erase(v[j++]);
    auto lo = S.lower bound(p - d), hi =

→ S.upper bound(p + d);

    for (; lo != hi; ++lo)
  ret = min(ret, {(*lo - p).dist2(), {*lo, p}});
    S.insert(p);
  return ret.second;
3.6 Convex Hull
typedef long long ll;
typedef pair <ll, ll> point;
#define x first
#define y second
inline ll area (point a, point b, point c) {
  return (b.x - a.x) * (c.y - a.y) - (b.y - a.y) *
   \leftarrow (c.x - a.x);
vector <point> convexHull (vector <point> p) {
  int n = p.size(), m = 0;
  if (n < 3) return p;
  vector <point> hull(n + n);
```

sort(p.begin(), p.end());

 $\rightarrow$  p[i]) <= 0) --m;

for (int i = 0; i < n; ++i)

while (m > 1 and area(hull[m - 2], hull[m - 1],

#### 3.7 Minimum Enclosing Circle

```
Expected runtime: O(n)
// Solves Gym 102299J
using namespace std;
typedef long double ld;
typedef pair <ld, ld> point;
#define x first
#define y second
point operator + (const point &a, const point &b) {
  return point(a.x + b.x, a.y + b.y);
point operator - (const point &a, const point &b) {
  return point(a.x - b.x, a.y - b.y);
point operator * (const point &a, const ld &b) {
  return point(a.x * b, a.y * b);
point operator / (const point &a, const ld &b) {
  return point(a.x / b, a.y / b);
const ld EPS = 1e-8;
const ld INF = 1e20
const ld PI = acosl(-1);
inline ld dist (point a, point b) {
  return hypotl(a.x - b.x, a.y - b.y);
inline ld sqDist (point a, point b) {
  return (a.x - b.x) * (a.x - b.x) + (a.y - b.y) *
   \rightarrow (a.y - b.y);
inline ld dot (point a, point b) {
  return a.x * b.x + a.y * b.y;
inline ld cross (point a, point b) {
  return a.x * b.y - a.y * b.x;
inline ld cross (point a, point b, point c) {
  return cross(b - a, c - a);
inline point perp (point a) {
  return point(-a.y, a.x);
// circle through 3 points
pair <point, ld> getCircle (point a, point b, point c)
  pair <point. ld> ret:
  ld den = (ld) 2 * cross(a, b, c);
ret.x.x = ((c.y - a.y) * (dot(b, b) - dot(a, a)) -
     (b.y - a.y) * (dot(c, c) - dot(a, a))) / den;
  ret.x.y = ((b.x - a.x) * (dot(c, c) - dot(a, a)) -
      (c.x - a.x) * (dot(b, b) - dot(a, a))) / den;
  ret.\dot{y} = dist(ret.x, a);
  return ret;
pair <point, ld> minCircleAux (vector <point> &s,

→ point a, point b, int n) {
```

```
ld lo = -INF, hi = INF;
  for (int i = 0; i < n; ++i) {
    auto si = cross(b - a, s[i] - a);

    if (fabs(si) < EPS) continue;
point m = getCircle(a, b, s[i]).x;</pre>
    auto cr = cross(b - a, m - a);
    si < 0 ? hi = min(hi, cr) : lo = max(lo, cr);
  1d v = 0 < lo ? lo : hi < 0 ? hi : 0
  point c = (a + b) * 0.5 + perp(b - a) * v /

    sqDist(a, b);

  return {c, sqDist(a, c)};
pair <point, ld> minCircle (vector <point> &s, point
- a, int n) {
  random_shuffle(s.begin(), s.begin() + n);
  point \bar{b} = s[0], c = (a + b) * \bar{0}.5;
  ld r = sqDist(a, c);
  for (int i = 1; i < n; ++i) {
  if (sqDist(s[i], c) > r * (1 + EPS))
       tie(c, r) = n == s.size() ? minCircle(s, s[i],

→ i) : minCircleAux(s, a, s[i], i);
  return {c, r};
pair <point, ld> minCircle (vector <point> s) {
  assert(!s.empty());
  if (s.size() == 1) return {s[0], 0};
  return minCircle(s, s[0], s.size());
int n; vector <point> p;
int main() {
  cin >> n;
  while (n--) {
    double x, y;
scanf("%lf %lf", &x, &y);
    p.emplace back(x, y);
  pair <point, ld> circ = minCircle(p);
  printf("%0.12f %0.12f %0.12f\n", (double) circ.x.x,

→ (double) circ.x.y, (double) (0.5 * circ.y));
  return 0:
```

```
3.8 Point In Polygon
// Test if a point is inside a convex polygon in O(lg
    n) time
// Solves SPOJ INOROUT typedef long long ll;
typedef pair <ll, ll> point;
#define x first
#define y second
struct seament {
  point P1, P2;
  segment () {}
  segment (point P1, point P2) : P1(P1), P2(P2) {}
inline ll ccw (point A, point B, point C) {
  return (B.x - A.x) * (C.y - A.y) - (C.x - A.x) *
  \hookrightarrow (B.y - A.y);
inline bool pointOnSegment (segment S, point P) {
  11 x = P.x, y = P.y, x1 = S.P1.x, y1 = S.P1.y, x2 = S.P1.y
   \rightarrow S.P2.x, y2 = S.P2.y;
  ll a = x - x1, b = y - y1, c = x2 - x1, d = y2 - y1,
      dot = a * c + b * d, len = c * c + d * d;
  if (x1 == x2 \text{ and } y1 == y2) return x1 == x and y1 == y2
  if (dot < 0 or dot > len) return 0;
```

```
return x1 * len + dot * c == x * len and y1 * len +

→ dot * d == v * len:

const int M = 17
const int N = 10010;
struct polygon {
  int n; //_n > 1
  point p[N]; // clockwise order
  polygon () {}
  polygon (int n, point *T) {
    for (int i = 0; i < n; ++i) p[i] = T[i];
  bool contains (point P, bool strictlyInside) {
    int lo = 1, hi = n - 1;
    while (lo < hi){
      int mid = lo + hi >> 1
      if (ccw(p[0], P, p[mid]) > 0) lo = mid + 1;
      else hi = mid:
    if (ccw(p[0], P, p[lo]) > 0) lo = 1;
    if (!strictlyInside and
        pointOnSegment(segment(p[0], p[n - 1]), P))
       return 1;
    if (!strictlyInside and
        pointOnSegment(segment(p[lo], p[lo - 1]), P))
        return 1:
    if (lo == 1 or ccw(p[0], P, p[n - 1]) == 0) return
    return ccw(p[lo], P, p[lo - 1]) < 0;
```

```
3.9 near_pair
struct pt {
 int x, y, id;
struct cmp x {
  bool operator()(const pt \& a, const pt \& b) const {
    return a.x < b.x \mid | (a.x == b.x \&\& a.y < b.y);
struct cmp_y {
  bool operator()(const pt & a, const pt & b) const {
    return a.y < b.y;
|vector<pt> a;
double mindist;
|pair<int, int> best pair;
void upd ans(const pt & a, const pt & b) {
  double dist = sqrt((a.x - b.x)*(a.x - b.x) + (a.y - b.x)
      b.y)*(a.y - b.y));
  if (dist < mindist) {</pre>
    mindist = dist;
    best pair = {a.id, b.id};
vector<pt> t;
|void rec(int l, int r) {
  if (r - l <= 3) {
    for (int i = l; i < r; ++i) {
      for (int j = i + 1; j < r; ++j) {
        upd ans(a[i], a[j]);
    sort(a.begin() + l, a.begin() + r, cmp y());
    return;
```

```
int m = (l + r) >> 1;
 int midx = a[m].x;
 rec(l, m);
 rec(m, r);
 merge(a.begin() + l, a.begin() + m, a.begin() + m,
  \rightarrow a.begin() + r, t.begin(), cmp y());
 copy(t.begin(), t.begin() + r - l, a.begin() + l);
 for (int i = l; i < r; ++i)
   if (abs(a[i].x - midx) < mindist)</pre>
     for (int j = tsz - 1; j >= 0 \&\& a[i].y - t[j].y
      t[tsz++] = a[i];
void solve(int n)
 t.resize(n)
 sort(a.begin(), a.end(), cmp x());
 mindist = 1E20:
 rec(0, n);
```

```
3.10 sweep
const double EPS = 1E-9;
struct pt {
  double x, y;
struct seg {
  pt p, q;
  int id;
  double get y(double x) const {
    if (abs(\overline{p}.x - q.x) < EPS)
      return p.y;
    return p.y + (q.y - p.y) * (x - p.x) / (q.x - p.x);
bool intersect1d(double l1, double r1, double l2,
   double r2) {
 if (l1 > r1)
    swap(l1, r1);
  if (l2 > r2)
    swap(l2, r2);
  return \max(l1, l2) \ll \min(r1, r2) + EPS;
int vec(const pt& a, const pt& b, const pt& c) {
  double s = (b.x - a.x) * (c.y - a.y) - (b.y - a.y) *
   \leftarrow (c.x - a.x);
  return abs(s) < EPS ? 0 : s > 0 ? +1 : -1;
|bool intersect(const seg& a, const seg& b)
  return intersect1d(a.p.x, a.q.x, b.p.x, b.q.x) &&
    intersect1d(a.p.y, a.q.y, b.p.y, b.q.y) &&
    vec(a.p, a.q, b.p) * vec(a.p, a.q, b.q) <= 0 &&
    vec(b.p, b.q, a.p) * vec(b.p, b.q, a.q) <= 0;
bool operator<(const seg& a, const seg& b)
  double x = max(min(a.p.x, a.q.x), min(b.p.x, b.q.x));
  return a.get y(x) < b.get y(x) - EPS;
struct event {
  double x;
  int tp, id;
```

event() {}

```
event(double x, int tp, int id) : x(x), tp(tp),
  → id(id) {}
  bool operator<(const event& e) const {
    if (abs(x - e.x) > EPS)
      return x < e.x;
    return tp > e.tp;
set<seg> s;
vector<set<seg>::iterator> where;
set<seg>::iterator prev(set<seg>::iterator it) {
 return it == s.begin() ? s.end() : --it;
set<seg>::iterator next(set<seg>::iterator it) {
pair<int, int> solve(const vector<seg>& a) {
  int n = (int)a.size();
  vector<event> e:
  for (int i = 0; i < n; ++i) {
    e.push back(event(min(a[i].p.x, a[i].q.x), +1, i));
    e.push_back(event(max(a[i].p.x, a[i].q.x), -1, i));
  sort(e.begin(), e.end());
  s.clear();
  where resize(a.size());
  for (size t i = 0; i < e.size(); ++i) {
    int id = e[i].id;
    if (e[i].tp = +1) {
      set<seg>::iterator nxt = s.lower bound(a[id]),
      → prv = prev(nxt);
     if (nxt != s.end() && intersect(*nxt, a[id]))
return make_pair(nxt->id, id);
      if (prv != s.end() && intersect(*prv, a[id]))
        return make pair(prv->id, id);
      where[id] = s_insert(nxt, a[id]);
      set<seg>::iterator nxt = next(where[id]), prv =
      → prev(where[id]);
      if (nxt != s.end() && prv != s.end() &&

    intersect(*nxt, *prv))

        return make pair(prv->id, nxt->id);
      s.erase(where[id]);
  return make pair(-1, -1);
```

## 4 Graph

# 4.1 2Sat

```
namespace sat{
 const int MAX = 200010;
 bool vis[MAX];
 vector <int> ed[MAX], rev[MAX]
 int n, m, ptr, dfs_t[MAX], ord[MAX], par[MAX];
 inline int inv(int x)
   return ((x) \le n ? (x + n) : (x - n));
 void init(int vars){
   n = vars, m = vars << 1;
   for (int i = 1; i \le m; i++){
     ed[i].clear();
rev[i].clear();
 inline void add(int a, int b){
   ed[a].push back(b);
   rev[b].push back(a);
```

```
inline void OR(int a, int b){
  add(inv(a), b);
add(inv(b), a);
inline void AND(int a, int b){
  add(a, b);
  add(b, a);
void XOR(int a,int b){
  add(inv(b), a);
  add(a, inv(b));
  add(inv(a), b);
  add(b, inv(a));
inline void XNOR(int a. int b){
  add(a,b);
  add(b,a)
  add(inv(a), inv(b));
add(inv(b), inv(a));
inline void force true(int x){
  add(inv(x), x);
inline void topsort(int s){
  vis[s] = true;
  for(int x : rev[s]) if(!vis[x]) topsort(x);
  dfs t[s] = ++ptr;
inline void dfs(int s, int p){
  par[s] = p;
  vis[s] = true;
  for(int x : ed[s]) if (!vis[x]) dfs(x, p);
void build(){
  CLR(vis);
ptr = 0;
  for(int i=m;i>=1;i--)
    if (!vis[i]) topsort(i);
    ord[dfs t[i]] = i;
  ĆLR(vis);
  for (int i = m; i >= 1; i --){
    int x = ord[i];
    if (!vis[x]) dfs(x, x);
bool satisfy(vector < int>& res){
  build();
  CLR(vis);
  for (int i = 1; i \le m; i++){
    int x = ord[i];
    if (par[x] == par[inv(x)]) return false;
    if (!vis[par[x]]){
      vis[par[x]] = true;
      vis[par[inv(x)]] = false;
  res.clear();
  for (int i = 1; i \le n; i++){
    if (vis[par[i]]) res.push back(i);
  return true:
```

#### 4.2 Centroid decomp

```
vector <int> ed[MAX]
bool isCentroid[MAX]
int sub[MAX], cpar[MAX], clevel[MAX];
int dis[20][MAX];
lvoid calcSubTree(int s,int p) {
```

```
sub[s] = 1;
  for(int x : ed[s]) {
    if(x == p or isCentroid[x]) continue;
    calcSubTree(x,s);
    sub[s] += sub[x];
int nn:
int getCentroid(int s,int p) {
  for(int x : ed[s]) {
    if(!isCentroid[x] && x!=p && sub[x]>(nn/2)) return

→ getCentroid(x,s);

  return s;
void setDis(int s, int from, int p, int lev) {
  dis[from][s] = lev;
  for(int x : ed[s]) -
    if(x == p or isCentroid[x] ) continue;
    setDis(x, from, s, lev+1);
void decompose(int s,int p,int lev) {
  calcSubTree(s,p);
  nn = sub[s];
  int c = getCentroid(s,p);
  setDis(c,lev,p,0);
  isCentroid[c] = true;
  cpar[c] = p;
  clevel[c] = lev;
  for(int x : ed[c])
    if(!isCentroid[x]) decompose(x,c,lev+1);
int ans[MAX];
inline void update(int v) {
  int u = v;
  while(u!=-1) {
    ans[u] = min(ans[u], dis[clevel[u]][v]);
    u = cpar[u];
inline int query(int v) {
  int ret = INF;
  int u = v;
  while(u != -1) {
    ret = min(ret, dis[clevel[u]][v]+ans[u]);
    u = cpar[u];
  return ret;
int main()
  decompose(1,-1,0);
  for(int i=1; i<=n; i++) ans[i] = INF;</pre>
  update(v);
  query(v));
  return 0;
```

#### 4.3 articulation\_point

```
using namespace std;
const int N = 1e5 + 10;
vector<int> g[N];
void dfs(int u, int p) {
 low[u] = vis[u] = ++now; int ch = 0;
 for(int v : q[u]){
   if(v ^ p)
    if(vis[v]) low[u] = min(low[u], vis[v]);
```

```
else {
         ch++; dfs(v, u);
low[u] = min(low[u], low[v]);
if(p + 1 && low[v] >= vis[u]) cut[u] = 1;
         if(low[v] > vis[u]) {
  printf("Bridge %d -- %d\n", u, v);
  } if(p == -1 && ch > 1) cut[u] = 1;
void ArticulationPointAndBridge() {
  now = 0;
  for(int i = 0; i < n; i++) {
    if(!vis[i]) dfs(i, -1);
4.4 bcc
// clear ed[] every test case
// tot -> total number of components
// bcc[i] contains the nodes of the i'th component
// any self loop or multiple edge?
const int MAX = ?:
vector <int> ed[MÁX];
bool cut[MAX];
int tot, Time, low[MAX], st[MAX];
vector <int> bcc[MAX];
stack <int> S;
void popBCC(int s,int x) {
  cut[s] = 1;
  bcc[++tot].pb(s);
  while(bcc[tot]_back() ^ x) {
    bcc[tot].pb(S.top());
    S.pop();
void dfs(int s, int p = -1) {
  S.push(s);
  int ch = 0;
  st[s] = low[s] = ++Time;
  for(int x : ed[s]) {
    if(!st[x]) {
       dfs(x,s);
low[s] = min(low[s],low[x]);
       if(p != -1 \text{ and } low[x] >= st[s]) popBCC(s,x);
       else if(p == -1) if(ch > 1) popBCC(s,x);
    else if(p != x) low[s] = min(low[s],st[x]);
  if(p == -1 \&\& ch > 1) cut[s] = 1;
void processBCC(int n) {
  for(int i=1;i<=n;i++) bcc[i].clear();
CLR(st); CLR(cut);</pre>
  Time = tot = 0;
  for(int i=1; i<=n; i++) {
   if(!st[i]) {</pre>
       dfs(i,-1);
       if(!S.empty()) ++tot;
       while(!S.empty()) {
         bcc[tot].push back(S.top());
         S.pop();
4.5 bridge_tree
const int MAXN = ?;
const int MAXE = ?;
```

```
struct edges {
int u,v;
ara[MAXE];
vector <int> ed[MAXN];
vector <int> isBridge[MAXN];
|vector <int> brTree[MAXN];
bool vis[MAXN];
int st[MAXN], low[MAXN], Time = 0;
int cnum;
|int comp[MAXN];
void findBridge(int s,int par) {
  int i,x,child = 0,j;
   \underline{v}is[s] = 1;
  Time++;

st[s] = low[s] = Time;

for(i=0; i<ed[s].size(); i++) {
     x = ed[s][i]
     if(!vis[x]) {
        child++;
        findBridge(x,s);
        low[s] = min(low[s],low[x]);
        if(low[x] > st[s])
          isBridge[s][i] = 1
          j = lower bound(ed[x].begin(),ed[x].end(),s)-e
               d[x].begin();
          isBridge[x][j] = 1;
     else if(par!=x)
       low[s] = min(low[s], st[x]);
void dfs(int s) {
  int i,x;
vis[s] = 1;
   comp[s] = cnum;
   for(i=0; i<ed[s].size(); i++) {</pre>
     if(!isBridge[s][i]) {
       x = ed[s][i];
if(!vis[x]) dfs(x);
void processBridge(int n,int m) {
  CLR(vis);
   Time = 0
   for(int i=1; i<=n; i++) if(!vis[i]) findBridge(i,-1);</pre>
   CLR(vis);
   for(int i=1; i<=n; i++) {
    if(!vis[i]) { cnum++;
       dfs(i);
  n = cnum;
   for(int i=1; i<=m; i++) {
     if(comp[ara[i].u] != comp[ara[i].v]) {
    br_ree[comp[ara[i].u]].pb(comp[ara[i].v]);
        brTree[comp[ara[i].v]].pb(comp[ara[i].u]);
  }
|int main() {
  int n,m,u,v;
scanf("%d %d",&n,&m);
for(int i=1; i<=m; i++) {</pre>
     sii(u,v);
     ed[u].pb(v);
     ed[v].pb(u);
```

```
isBridge[u].pb(0);
   isBridge[v].pb(0);
   ara[i].u = u;
ara[i].v = v;
 for(int i=1; i<=n; i++) sort(all(ed[i]));</pre>
 processBridge(n,m);
 return 0;
4.6 dinic
namespace dinic {
 using T = int;
 const T INF = 0x3f3f3f3f;
 const int MAXN = 5010;
 int n, src, snk, work[MAXN];
T dist[MAXN];
 struct Edge{
    int to, rev pos;
   T c, f;
 vector <Edge> ed[MAXN];
 void init(int n, int src, int snk) {
   n = n, src = src, snk = snk;
    for(int i=1;i<=n;i++) ed[i].clear();</pre>
 inline void addEdge(int u, int v, T c, T rc = 0) {
    Edge a = \{v, (int)ed[v].size(), c, 0\};
    Edge b = \{u, (int)ed[u].size(), rc, 0\};
    ed[u].push back(a);
    ed[v].push back(b);
 bool dinic bfs() {
    SET(dist);
dist[src] = 0;
    queue <int> q;
    q.push(src);
    while(!q.empty()){
      int u = q.front(); q.pop();
      for(Edge &e : ed[u]){
        if(dist[e.to] == -1 \text{ and } e.f < e.c) 
          dist[e.to] = dist[u] + 1;
          q.push(e.to);
    return (dist[snk]>=0);
 T dinic dfs(int u, T fl){
    if (u == snk) return fl;
    for (; work[u] < (int)ed[u].size(); work[u]++) { Edge &e = ed[u][work[u]];
      if (e.c <= e.f) continue;</pre>
      int v = e.to:
      if (dist[v] == dist[u] + 1)
        T df = dinic dfs(v, min(fl, e.c - e.f));
        if (df > 0){
  e.f += df;
          ed[v][e.rev pos].f -= df;
          return df;
    return 0;
 T solve() {
    T ret = 0:
```

while (dinic bfs()) {

CLR(work);

```
while (T delta = dinic dfs(src, INF)) ret +=

    delta;

    return ret;
int main() {
  int n, m, u, v, c;
  cin >> n >> m:
  dinic::init(n, 1, n);
  while(m--) {
    cin >> u >> v >> c:
    dinic::addEdge(u, v, c, c);
  cout << dinic::solve() << '\n';
  return 0;
4.7 dsu on tree
void calcSubSize(int s,int p) {
  sub[s] = 1;
for(int x : G[s]) {
    if(x==p) continue;
    calcSubSize(x,s);
    sub[s] += sub[x];
void add(int s,int p,int v,int bigchild = -1) {
  freq[color[s]] += v;
  for(int x : G[s]) {
  if(x==p | | x==bigchild) continue;
    add(x,s,v);
void dfs(int s,int p,bool keep) {
  int bigChild = -1;
  for(int x : G[s]) {
    if(x==p) continue;
    if(bigChild==-1 || sub[bigChild] < sub[x] )</pre>

→ bigChild = x;

  for(int x : G[s]) {
    if(x==p || x==bigChild) continue;
    dfs(x,s,0);
  if(bigChild!=-1) dfs(bigChild,s,1);
  add(s,p,1,bigChild);
  if(keep==0)
    add(s,p,-1);
4.8 euler_path
vector <int> ed[MAX+5], sltn;
int inDeg[MAX+5], outDeg[MAX+5];
bool vis[MAX+5];
void dfs(int nd) {
  vis[nd] = true;
  while(ed[nd].size())
    int v = ed[nd].back();
    ed[nd].pop_back();
    dfs(v);
  sltn.pb(nd);
int findEuler (int n) {
```

int src , snk , ret = 1;
bool found\_src = false, found\_snk = false;

for(int u = 1; u <= n; u++) {
 for(int i = 0; i<ed[u].size(); i++) {</pre>

CLR(inDeg); CLR(outDeg);

**int** v = ed[u][i];

```
outDeg[u]++;
       inDeg[v]++;
  int diff;
  for(int i = 1; i<=n; i++) {
  diff = outDeg[i] - inDeg[i];</pre>
    if(diff == 1)
       if(found_src) return 0;
       found sr\overline{c} = true;
       src = i;
    else if (diff == -1) {
       if(found snk) return 0;
       found sn\overline{k} = true;
       snk = i;
    else if(diff != 0) return 0;
  if(!found src) {
     ret = 2;
    for(int i = 1; i <= n; i++) {
  if( outDeg[i] ) {</pre>
         found src = true;
         src = i;
         break;
  if(!found src) return ret:
  CLR(vis);
  sltn.cléar();
  dfs(src);
  for(int i = 1; i<=n; i++) {
    if(outDeg[i] && !vis[i]) return 0;
  for(int i = (int)sltn.size()-1; i>=0; i--)

¬ printf("%d ",sltn[i]);

  puts'("");
  return ret:
4.9 hld
const int N = 3e4 + 5;
vector<int> G[N];
int sz[N], H[N], P[N];
void dfs(int cur, int h)
     sz[cur] = 1;
    H[cur] = h;
     for(int& to : G[cur])
         G[to].erase(find(G[to].begin(), G[to].end(),
            cur));
         P[to] = cur;
         df_s(to, h + 1)
         sz[cur] += sz[to];
         if(sz[to] > sz[G[cur][0]]) swap(G[cur][0], to);
int base[N], pos[N], head[N];
int ptr = 0;
|void hld(int cur)
     pos[cur] = ++ptr;
     base[ptr] = cur;
    for(int to : G[cur]) {
    head[to] = (to == G[cur][0] ? head[cur] : to);
         hld(to):
|segtree ST;
|int query(int u, int v)
```

```
int ret = 0;
    while(head[u] != head[v])
         if(H[head[u]] > H[head[v]]) swap(u, v);
         ret += ST.query(pos[head[v]], pos[v]);
         v = P[head[v]];
    if(H[u] > H[v]) swap(u, v);
ret += ST.query(pos[u], pos[v]);
     return ret;
void update(int u, int val) {
    ST.update(pos[u], val);
void build(int n, int root)
     ptr = 0;
    dfs(root, 0);
head[root] = root;
     hld(root);
    ST = segtree(n);
* clear graph
   call build
 * Prob : sum of values from u to v
4.10 hopcroft_karp
struct HopcroftKarp {
  const int N, M;
  std::vector<std::vector<int>> adj left;
  std::vector<int> matchL, matchR;
  HopcroftKarp(int N, int M, const

    std::vector<std::pair<int, int>>& edge)
       : N(N), M(M), matchL(N, -1), matchR(M, -1),
        → adj left(N) {
    for (auto [l, r] : edge)
  adj_left[l].push_back(r);
  int maxmatching() {
    int sz = 0;
for (bool updated = true; updated;) {
       updated = false;
       static std::vector<int> root(N), prev(N), qq(N);
       static int qi, qj;
       // std::queue<int> q;
       qi = qj = 0;
       std::fill(root.begin(), root.end(), -1),
std::fill(prev.begin(), prev.end(), -1);
       for (int i = 0; i < N; i++)
         if (matchL[i] == -1)
          qq[q]++] = i, root[i] = i, prev[i] = i; g.push(i), root[i] = i;
       while (qi < qj) {
         int u = qq[qi++];
         // int u = q.front(); q.pop();
if (matchL[root[u]] != -1) continue;
for (int v : adj_left[u]) {
           if (matchR[v] = -1) {
              while (v != -1)
                matchR[v] = u, std::swap(matchL[u], v),
                  \rightarrow u = prev[u]:
              updated = true, sz++;
              break;
```

if (prev[matchR[v]] == -1)

 $\rightarrow$  root[u], qq[qj++] = v;

v = matchR[v], prev[v] = u, root[v] =

#### 4.11 hungarian

```
// Given NN matrix A[i][j]. Calculate a permutation
\rightarrow p[i] that minimize A[i][p[i]].
template <typename T>
pair <T, vector <int>> Hungarian (int n, int m, T
\leftarrow c[N][N]) {
  vector <T> v(m), dist(m);
  vector <int> L(n, -1), R(m, -1);
vector <int> index(m), prev(m);
  auto residue = [&] (int i, int j) {return c[i][j] -
   → v[j];};
  iota(index.begin(), index.end(), 0);
  for (int f = \bar{0}; f < n; ++f)
    for (int j = 0; j < m; ++j) {
       dist[j] = residue(f, j), prev[j] = f;
    T w; int i, j, l, s = 0, t = 0;
    while (true) {
       if (s == t) {
         l = s, w = dist[index[t++]];
         for (int k = t; k < m; ++k) {
  j = index[k]; T h = dist[j];</pre>
           if (h <= w) {
   if (h < w) t = s, w = h;
   index[k] = index[t], index[t++] = j;</pre>
         for (int k = s; k < t; ++k) {
            j = index[k];
            if (R[j] < 0) goto augment;
       int q = index[s++], i = R[q];
       for (int k = t; k < m; ++k) {
           = index[k];
         \bar{T} h = residue(i, j) - residue(i, q) + w;
         if (h < dist[i]) {
            dist[j] = h, prev[j] = i;
            if (h == w) {
              if (R[j]' < 0) goto augment;
index[k] = index[t], index[t++] = j;
  augment:
    for (int k = 0; k < 1; ++k) v[index[k]] +=
        dist[index[k]] - w;
      R[j] = i = prev[j], swap(j, L[i]);
    } while (i ^ f);
  for (int i = 0; i < n; ++i) ret += c[i][L[i]];
return {ret, L};</pre>
```

#### 4.12 kuhn

```
namespace bpm{
  const int L = 105;
  const int R = 105;
  vector <int> G[L];
```

```
int matchR[R], matchL[L], vis[L], it;
void init(int n)
  SET(matchL), SET(matchR), CLR(vis);
  for(int i=1;i<=n;i++) G[i].clear();</pre>
inline void addEdge(int u,int v) { G[u].pb(v); }
bool dfs(int s) {
 vis[s] = it;
for(auto x : G[s]) {
   if( matchR[x] == -1 or (vis[matchR[x]] != it and

    dfs(matchR[x])))

      matchL[s] = x; matchR[x] = s;
      return true;
  return false;
int solve() {
  int cnt = 0:
  for(int i=1;i<=n;i++) {</pre>
    if(dfs(i)) cnt++, it++;
  return cnt;
```

// Don't forget to clear ed after test case ends(vt,

#### 4.13 lca

```
→ cost are cleared inside)

using namespace std:
const int MAX = 100010;
int dep[MAX], par[MAX][21];
vector <int> ed[MAX];
void dfs(int s, int p, int d) {
  dep[s] = d, par[s][0] = p;
  for(int x : ed[s]) {
    if(x == p) continue;
    dfs(x, s, d+1);
void preprocess(int root, int n) {
 LG = \lg(n);
  memset(par, -1, sizeof(par));
  dfs(root, -1, 0);
  for(int j=1; j<=LG; j++) {</pre>
    for(int i=1;i<=n;i++) +</pre>
      if(par[i][j-1] != -1) par[i][j] =
       → par[par[i][j-1]][j-1];
  }
|int getLCA(int u, int v) {
  if(dep[u] < dep[v]) swap(u, v);</pre>
  for(int i=LG;i>=0;i--) {
    if(dep[u] - (1 << i) >= dep[v]) u = par[u][i];
  if(u == v) return u;
  for(int i=LG;i>=0;i--) {
    if (par[u][i] != -1 and par[u][i] - par[v][i]) {
      u = par[u][i], v = par[v][i];
  return par[u][0];
```

#### 4.14 manhattan MST

```
using namespace std;
using ll = long long;
```

```
struct UnionFind {
    vector<int> UF;
    int cnt;
    , description find(UF[v]); }
    bool join(int v, int w) {
        if ((v = find(v)) == (w = find(w))) return
        if (UF[v] > UF[w]) swap(v, w);
        UF[v] += UF[w];
UF[w] = v;
        return true;
    bool connected(int v, int w) {
        return find(v) == find(w);
    int getSize(int v) { return -UF[find(v)]; }
template <class T>
struct KruskalMST {
    using Edge = tuple<int, int, T>;
    T mstWeight;
    vector<Edge> mstEdges;
    UnionFind uf;
    KruskalMST(int V, vector<Edge> edges) :

→ mstWeight(),

    uf(V) {
        sort(edges.begin(), edges.end(), [&](const

→ Edge &a, , const Edge &b)

            return qet<2>(a) < qet<2>(b);
        });
for (auto &&e : edges) {
            if (int(mstEdges.size()) >= V - 1) break;
            if (uf.join(get<0>(e), get<1>(e))) {
                mstEdges.push back(e);
                mstWeight += \overline{g}et < 2 > (e);
        }
    }
template <class T>
struct ManhattanMST : public,⊟ KruskalMST<T> {
    using Edge = typename KruskalMST<T>::Edge;
    static vector<Edge>

¬ generateCandidates(vector<pair<T, T>>, P) {
        vector<int> id(P.size());
iota(id.begin(), id.end(), 0);
        vector<Edge> ret;
        ret.reserve(P.size() * 4);
        for (int h = 0; h < 4; h++)
            sort(id.begin(), id.end(), [&](int i, int

→ j) {
                return P[i].first - P[j].first <</pre>
                 → P[j].second - P[i].second;
            map<T, int> M;
            for (int i : id) {
                 auto it = M.lower bound(-P[i].second);
                for (; it != M.en\overline{d}(); it =

→ M.erase(it)) {
                     int j = it->second;
                    T dx = P[i].first - P[j].first, dy
                     → = P[i].second, - P[j].second;
                    if (dy > dx) break;
                     ret emplace back(i, j, dx + dy);
                M[-P[i].second] = i;
```

```
for (auto \&\&p : P) {
                  if (h % 2)
                       p.first = -p.first;
                       swap(p.first, p.second);
         return ret;
    ManhattanMST(const vector<pair<T, T>> &P)
         : KruskalMST<T>(P.size(),

    generateCandidates(P)) {}
};
int main() {
    int N;
     cin >> N:
    vector<pair<ll, ll>> P(N);
for (auto &&p : P) cin >> p.first >> p.second;
    ManhattanMST mst(P);
     cout << mst.mstWeight << '\n';</pre>
    for (auto &&[v, w, weight] : mst.mstEdges) cout <<</pre>

    v << □ , □ << w << '\n'; return 0;
</pre>
```

#### 4.15 mcmf

```
namespace mcmf {
 using T = int;
  const T INF = ?; // 0x3f3f3f3f or

    0x3f3f3f3f3f3f3f3fLL

  const int MAX = ?; // maximum number of nodes
 int n, src, snk;
T dis[MAX], mCap[MAX];
  int par[MAX], pos[MAX];
  bool vis[MAX];
  struct Edge{
   int to, rev pos;
   T cap, cost, flow;
 vector <Edge> ed[MAX];
  void init(int n, int src, int snk) {
   n = n, src = src, snk = snk;
    for(int i=1;i<=n;i++) ed[i].clear();</pre>
  void addEdge(int u, int v, T cap, T cost) {
   Edge a = {v, (int)ed[v].size(), cap, cost, 0};
   Edge b = \{u, (int)ed[u].size(), 0, -cost, 0\};
   ed[u].pb(a);
   ed[v].pb(b);
  inline bool SPFA(){
    CLR(vis);
    for(int i=1; i<=n; i++) mCap[i] = dis[i] = INF;</pre>
    queue <int> q;
   dis[src] = 0;
vis[src] = true;
    q.push(src);
    while(!q.empty()){
      int u = q.front();
      q.pop();
      vis[u] = false;
      for(int i=0; i<(int)ed[u].size(); i++) {</pre>
        Edge &e = ed[u][i];
        int v = e.to;
        if(e.cap > e.flow && dis[v] > dis[u] + e.cost){
          dis[v] = dis[u] + e.cost;
          par[v] = u;
          pos[v] = i;
          mCap[v] = min(mCap[u],e.cap - e.flow);
          if(!vis[v]) {
            vis[v] = true;
```

```
q.push(v);
        }
  return (dis[snk] != INF);
inline pair <T, T> solve() {
  T F = 0, C = 0, f;
  int u, v;
  while(SPFA()){
    u = snk:
    \tilde{f} = mCap[u];
    F += f;
    while(u!=src){
      v = par[u];
      ed[v][pos[u]].flow += f; // edge of v-->u
      ed[u][ed[v][pos[u]].rev_pos].flow -= f;
    \dot{C} += dis[snk] * f;
  return make pair(F,C);
```

#### 4.16 tree hash all root

```
/*Description : Find hashes of a tree when rooted at

→ each possible

node(unrooted tree isomorphism test).Time : O(n)*/
const int sz = 2e5 + 5,mod = 1e9 + 7;
ll hval[sz], h[sz], dp[sz], rans[sz];
void dfs(vector<vector<int>>> &g, int u = 0, int p =
   -1, int val = 0, int up = 0) {
    vector<int> cv, cht; // current child values &
        heights
    if (u > 0) {
         cv.push back(val);
         cht.push back(up);
    for (int v : g[u])
         if (v - p) {
              cv.push back(dp[v]);
              cht.pus\overline{h} back(1 + h[v]);
    sort(cht.begin(), cht.end(), greater<int>());
    if (cv.size() > 1) {
    ll ret[] = {1, 1}; // for biggest &
         -- 2nd-biggest heights
for (int i = 0; i < 2; i++)
    for (int value : cv)
        ret[i] = ret[i] * (hval[cht[i]] +</pre>
         ¬ value) % mod;
rans[u] = ret[0]; // biggest is hash for this
         for (int v : q[u])
              if (v - p) {
                   int id = 1;
                   if (cht[0] - 1 - h[v]) id = 0; // v
                       is not on the biggest height path
                   val = ret[id] * invmod((hval[cht[id])

→ + dp[v]) % mod) % mod;
                   /* division of v subtree hash value */
dfs(g, v, u, val, cht[id] + 1);
    } else if (cv.size()) { // Leaf node u OR vertex

→ - 1 has only one child
         if (!up)
              val = 1:
         else
              val = (val + hval[up]) % mod;
```

```
rans[u] = val;
        for (int v : g[u])
            if (v - p) dfs(g, v, u, val, up + 1);
ll get(vector<vector<int>> &g, int u = 0, int p = -1) {
   h[u] = 0;
vector<ll> childs;
    for (int v : g[u])
        if (v - p) {
            childs.push back(get(g, v, u));
            h[u] = \max(\overline{h}[u], 1 + \overline{h}[v]);
    ll ret = 1;
    for (int value : childs) ret = ret * (hval[h[u]] +

→ value) % mod;

    return dp[u] = ret;
int main() { // can remove the g as param, can change
   to 1-index, multi-test works, no further change
    get(g);
    dfs(g)
    tree[k] = rans[0] = dp[0];
    // rans[i] = tree hash with i as root
```

#### 5 Math 5.1 FWHT

```
const int N = 1 \ll 20;
// apply modulo if necessary
void fwht xor(int *a, int n, int dir = 0) {
    for (int h = 1; h < n; h <<= 1) {
         for (int i = 0; i < n; i + h << 1)
              for (int j = i; j < i + h; ++j) {
                  int x = a[j], y = a[j + h];
a[j] = x + y, a[j + h] = x - y;
if (dir) a[j] >>= 1, a[j + h] >>= 1;
         }
void fwht or(int *a, int n, int dir = 0) {
    for (int h = 1; h < n; h <<= 1) {
         for (int i = 0; i < n; i += h << 1)
              for (int j = i; j < i + h; ++j) {
                  int x = a[j], y = a[j + h];
a[j] = x, a[j + h] = dir ? y - x : x +

→ y;

         }
void fwht and(int *a, int n, int dir = 0) {
    for (int h = 1; h < n; h <<= 1) {
         for (int i = 0; i < n; i += h << 1)
              for (int j = i; j < i + h; ++j) {
                  int x = a[j], y = a[j + h];
                  a[j] = dir^{2} x - y : x + y, a[j + h] =
                      у;
         }
```

#### 5.2 FloorSum

```
long long FloorSumAP(long long a, long long b, long
→ long c, long long n){
if(!a) return (b / c) * (n + 1);
```

```
if(a >= c \text{ or } b >= c) \text{ return } ( (n * (n + 1)) / 2)
      (a / c) + (n + 1) * (b / c) + FloorSumAP(a % c,
 return m * n - FloorSumAP(c, c - b - 1, a, m - 1);
1/\sqrt{0(\log a)} sum^n x * floor(ax + b / c) = g, sum^n
\rightarrow floro(ax + b / c)^2 = h
struct dat {
 long long f, g, h; dat(long long f = 0, long long g = 0, long long h = 0
  \rightarrow 0): f(f), g(g), h(h) {};
long long mul(long long a, long long b){
 return (a * b) % MOD;
dat query(long long a, long long b, long long c, long
\rightarrow long n){
 if(!a) return {mul(n + 1, b / c), mul(mul(mul(b / c,
   \rightarrow n), n + 1), inv2), mul(mul(n + 1, b / c), b /c)};
 long long f, g, h;
  dat nxt;
  if(a >= c or b >= c){
   nxt = query(a % c, b % c, c, n);
   f = (nxt.f + mul(mul(mul(n, n + 1), inv2), a / c)
        + mul(n + 1, b / c)) % MOD;
   g = (nxt.\dot{g} + mul(a / c, mul(mul(n, n + 1), mul(2 *
        n + 1, inv6))) + mul(mul(b / c, mul(n, n +
       1)), inv2)) % MOD;
   h = (nxt.h + 2* mul(b / c, nxt.f) + 2* mul(a)
        c, nxt.q) + mul(mul(a / c, a / c), mul(mul(n,
        n + 1), mul(2 * n + 1, inv6))) + <math>mul(mul(b / 1))
       c, b / c), n + 1) + mul(mul(a / c, b / c),
    \stackrel{\sim}{\sim} mul(n, n + 1)) ) % MOD;
    return {f, g, h};
  long long m = (a * n + b ) / c;
  nxt = query(c, c - b - 1, a, m - 1);
  f = (mul(m, n) - nxt.f) % MOD;
 q = mul(mul(m, mul(n, n + 1)) - nxt.h - nxt.f,
 h = (mul(n, mul(m, m + 1)) - 2 * nxt.q - 2 * nxt.f -

→ f) % MOD;

  return {f, g, h};
```

#### 5.3 NOD

```
N = input()
primes = array containing primes till 10^6
ans = 1
for all p in primes :
    if p*p*p > N:
        break
    count = 1
    while N divisible by p:
        N = N/p
        count = count + 1
ans = ans * count
if N is prime:
    ans = ans * 2
else if N is square of a prime:
    ans = ans * 3
else if N != 1:
    ans = ans * 4
```

#### 5.4 Pollard Rho

```
we#include <bits/stdc++.h>
using namespace std;
typedef long long ll;
typedef unsigned long long ull;
namespace Rho {
```

```
ull mul (ull a, ull b, ull mod)
    ll ret = a * b - mod * (ull) (1.L / mod * a * b);
return ret + mod * (ret < 0) - mod * (ret >= (ll)
  ull bigMod (ull a, ull e, ull mod) {
    ull ret = 1:
       if (e & \hat{1}) ret = mul(ret, a, mod);
       a = mul(a, a, mod), e >>= 1;
    return ret;
  bool isPrime (ull n) {
  if (n < 2 or n % 6 % 4 != 1) return (n | 1) == 3;</pre>
    ull a[] = \{2, 325, 9375, 28178, 450775, 9780504,
         1795265022
    ull s = builtin ctzll(n - 1), d = n >> s;
    for (ull x : a)
       ull p = bigMod(x % n, d, n), i = s;
       while (p != 1 \text{ and } p != n - 1 \text{ and } x \% n \text{ and } i--)
       \rightarrow p = mul(p, p, n);
       if (p != n - 1 \text{ and } i != s) return 0;
    return 1;
  ull pollard (ull n) {
    auto f = [\&] (ull x) {return mul(x, x, n) + 1;};
    ull x = 0, y = 0, t = 0, prod = 2, i = 1, q;
    while (t++ \% 40 \text{ or } \gcd(\text{prod}, n) == 1) {
       if (x == y) x = ++i, y = f(x);
       if ((q = mul(prod, max(x, y) - min(x, y), n)))
          prod = q;
      x = f(x), y = f(f(y));
    return gcd(prod, n);
  vector <ull> factor (ull n) {
    if (n == 1) return {};
    if (isPrime(n)) return {n};
    ull x = pollard(n);
    auto l = factor(x), r = factor(n / x);
    l.insert(l.end(), r.begin(), r.end());
    return l;
int t; ll n;
int main() {
  cin >> t;
  while (t--)
    scanf("%ld", &n);
vector <ull> facs = Rho::factor(n);
sort(facs.begin(), facs.end());
    printf("%d", (int) facs.size())
    for (auto it : facs) printf(" %llu", it);
    puts("");
  return 0;
```

#### . .

#### 5.6 crt

```
//r[i][j]= inverse of p[i] modulo p[j]
//ans= x[0]+x[1]*p[0]+x[2]*(p[0]*p[1])+...+x[k-1]*(p[0
\rightarrow ]*p[1]*p[2]*...*p[k-2])
//ans = ((p[0]*p[1]'*p[2]*...*p[k-1])
for (int i = 0; i < k; ++i) {
    x[i] = a[i];
    for (int j = 0;
                     j < i; ++j) {
        x[i] = r[j][i] * (x[i] - x[j]);
        x[i] = x[i] % p[i];
        if (x[i] < 0)
            x[i] += p[i];
ll mul= p[0],res=x[0],tot=1;
F(i,0,k) tot *= p[i];
F(i,1,k)
res+= x[i]*mul;
res %= tot;
 mul *= p[i];
res %= mul:
return res;
```

#### 5.7 derangement

```
int derangement(int n)
{
    if(!n) return n;
    if(n <= 2) return n-1;
    return (n-1)*(derangement(n-1) + derangement(n-2));
}</pre>
```

#### 5.8 diophantine

```
5.9 discrete log
// Returns minimum x for which a ^x % m = b % m, a
→ and m are coprime.
int solve(int a, int b, int m) {
    a %= m, b %= m;
    int n = sqrt(m) + 1;
    int an = 1;
    for (int i = 0; i < n; ++i)
an = (an * 1ll * a) % m;</pre>
    unordered map<int, int> vals;
    for (int \overline{q} = 0, cur = b; q <= n; ++q) {
        vals[cur] = q;
cur = (cur * 1ll * a) % m;
    for (int p = 1, cur = 1; p <= n; ++p) {
        cur = (cur * 111 * an) % m;
        if (vals.count(cur)) {
             int ans = n * p - vals[cur];
             return ans;
    return -1;
5.10 factorial mod p
// O(log p(n)) gives me n! % p for large n, p
int factmod(int n, int p) {
```

# vector<int> f(p); f[0] = 1; for (int i = 1; i < p; i++) $\dot{f}[i] = f[i-1] * i % p;$ int res = 1: while (n > 1) { if ((n/p) % 2) res = p - res; res = res \* f[n%p] % p; n /= p;**return** res;

#### 5.11 fft

```
typedef complex<double> base:
#define PI acos(-1)
void fft(vector<base> &a, bool invert){
    int n = (int)a.size();
    for (int i = 1, j = 0; i < n; ++i){
         int bit = n >> 1;
         for (; j >= bit; bit >>= 1) j -= bit;
         j += bit;
        if (i < j)swap(a[i], a[j]);</pre>
    for (int len = 2; len <= n; len <<= 1) {
    double ang = 2 * PI / len * (invert ? -1 : 1);</pre>
        base wlen(cos(ang), sin(ang));
        for (int i = 0; i < n; i + len){
             base w(1);
             for (int j = 0; j < len / 2; ++j){
  base u = a[i + j], v = a[i + j + len /</pre>
                    . 2] * w;
                  a[i + j] = u + v;
                  a[i + j + len / 2] = u - v;
                  w^* = wlen;
    if (invert) for (int i = 0; i < n; ++i) a[i] /= n;
void multiply(const vector<int> &a, const vector<int>
    &b, vector<int> &res)
    vector<base> fa(a.begin(), a.end()), fb(b.begin(),
     → b.end());
```

```
size t n = 1;
    while (n < max(a.size(), b.size())) n <<= 1;
    n <<= 1;
    fa.resize(n), fb.resize(n);
fft(fa, false), fft(fb, false);
for (size_t i = 0; i < n; ++i) fa[i] *= fb[i];</pre>
    fft(fa, true); res.resize(n);
    for (size t i = 0; i < n; ++i) res[i] =

    int(fa[i].real() + 0.5);

5.12 find_roots
/*Description: Finds the real roots to a polynomial.
Usage: polyRoots({{2,-3,1}},-1e9,1e9) // solve x2-3x+2
Time: O(n^2\log(1/))*/
struct Poly {
    vector<double> a;
    double operator()(double x) const {
         double val = 0:
        for (int i = a.size(); i--;) (val *= x) +=
         \rightarrow a[i];
        return vál;
    void diff() {
        for (int i = 1; i < a.size(); ++i) a[i - 1] =
         \rightarrow i * a[i];
        a.pop back();
    void divroot(double x0) {
        double b = a.back(), c;
        a.back() = 0;
        for (int i = a.size() - 1; i--;) c = a[i],
         \rightarrow a[i] = a[i + 1] * x0 + b, b = c;
        a.pop back();
vector<double> polyRoots(Poly p, double xmin = -1e9,
    double xmax = 1e9) {
    if (p.a.size() == 2) {
        return {-p.a[0] / p.a[1]};
    vector<double> ret;
    Poly der = p;
    der.diff();
    auto dr = polyRoots(der, xmin, xmax);
    dr.push back(xmin - 1);
    dr.push back(xmax + 1);
    sort(dr.begin(), dr.end());
    for (int i = 0; i < dr.size() - 1; ++i) {
         double l = dr[i], h = dr[i + 1];
         bool sign = p(l) > 0;
         if (sign ^ (p(h) > 0)) {
             for (int it = 0; it < 60; ++it) { //
                 while (h-l > 1e-8)
                 double m = (l + h) / 2, f = p(m);
                 if ((f <= 0) ^ sign)
                      l = m:
                      h = m;
             ret.push back((l + h) / 2);
    return ret;
5.13 gauss_eliminition
const double EPS = 1e-9;
const int INF = 2; // it doesn't actually have to be
```

→ infinity or a big number

```
int gauss (vector < vector<double> > a, vector<double>
    & ans) {
     int n = (int) a.size();
     int m = (int) a[0].size() - 1;
     vector<int> where (m, -1);
     for (int col=0, row=0; col<m && row<n; ++col) {
         int sel = row;
         for (int i=row; i<n; ++i)
    if (abs (a[i][col]) > abs (a[sel][col]))
          if (abs (a[sel][col]) < EPS)</pre>
              continue:
         for (int i=col; i<=m; ++i)
    swap (a[sel][i], a[row][i]);</pre>
         where[col] = row;
          for (int i=0; i<n; ++i)
              if (i != row) {
                   double c = a[i][col] / a[row][col];
                   for (int j=col; j<=m; ++j)</pre>
                        a[i][j] -= a[row][j]^* c;
         ++row;
    ans.assign (m, 0);
    for (int i=0; i<m; ++i)</pre>
         if (where[i] "!= -1)
    ans[i] = a[where[i]][m] / a[where[i]][i];
     for (int i=0; i<n; ++i) {
          double sum = 0;
         for (int j=0; j<m; ++j)
    sum += ans[j] * a[i][j];</pre>
         if (abs (sum - a[i][m]) > EPS)
return 0;
     for (int i=0; i<m; ++i)</pre>
         if (where[i] == -1)
              return INF:
     return 1;
//modular
int gauss (vector < bitset<N> > a, int n, int m,
    bitset<N> & ans) {
    vector<int> where (m, -1);
for (int col=0, row=0; col<m \&\& row<n; ++col) {
         for (int i=row; i<n; ++i)
    if (a[i][col]) {
        swap (a[i], a[row]);</pre>
                   break;
         if (! a[row][col])
              continue;
         where[col] = row;
         for (int i=0; i<n; ++i)
if (i != row && a[i][col])
                   a[i] ^= a[row];
         // The rest of implementation is the same as

→ above

const double EPS = 1E-9;
int compute rank(vector<vector<double>>> A) {
     int n = A.size();
     int m = A[0].size();
     int rank = 0;
     vector<bool> row selected(n, false);
     for (int i = 0; i < m; ++i) {
         int j;
for (j = 0; j < n; ++j) {
```

```
break;
    if (j != n) {
        ++rank;
        row_selected[j] = true;
        for (int p = i + 1; p < m; ++p)
            A[j][p] /= A[j][i];
        for (int k = 0; k < n; ++k)
            if (k != j \&\& abs(A[k][i]) > EPS) {
                for (int p = i + 1; p < m; ++p)
                    A[k][p] -= A[j][p] * A[k][i];
return rank;
```

#### 5.14 gen\_all\_k\_combs

```
vector<int> ans:
void gen(int n, int k, int idx, bool rev) {
    if (k > n | | k < 0)
        return:
    if (!n) ·
        for (int i = 0; i < idx; ++i) {
             if (ans[i])
                 cout << i + 1;
        cout << "\n";
        return:
    ans[idx] = rev;
gen(n - 1, k - rev, idx + 1, false);
    ans[idx] = !rev;
    gen(n-1, k-!rev, idx + 1, true);
void all combinations(int n, int k) {
    ans.resize(n);
gen(n, k, 0, false);
```

#### 5.15 integrate\_adaptive

```
/*Description: Fast integration using an adaptive

→ Simpsons rule.

Usage: double sphereVolume = quad(-1, 1, [](double x)  {
return quad(-1, 1, [&](double y) return quad(-1, 1, [&](double z)
return x*x + y*y + z*z < 1; \}); \}); */
typedef double d;
#define S(a, b) (f(a) + 4 * f((a + b) / 2) + f(b)) *
\rightarrow (b - a) / 6
template <class F>
d \operatorname{rec}(F\& f, d a, d b, d \operatorname{eps}, d S)  {
    d c = (a + b) / 2;

d S1 = S(a, c), S2 = S(c, b), T = S1 + S2;

if (abs(T - S) \le 15 * eps | | b - a < 1e-10)
     return T + (T - S) / 15;
return rec(f, a, c, eps / 2, S1) + rec(f, c, b,
      \rightarrow eps / 2, S2);
template <class F>
d quad(d a, d b, F f, d eps = 1e-8)
     return rec(f, a, b, eps, S(a, b));
```

#### 5.16 lagrange\_interpolation

```
/*Description: A polynomial of degree d can be
   uniquely identified
given its values on d + 1 unique points. O(n) to
→ pre-calculate given
```

```
if (!row selected[j] && abs(A[j][i]) > EPS) the first n points (x=0 to n-1) of the polynomial.
                                                     \rightarrow Then answer each query to interpolate the xth term in O(n). All values
                                                     are done modulo mod, which needs to be a prime as we need its inverse
                                                     modulo. Also
includes an additional helper function called

→ find degree(terms, mod).
                                                    Given at least the first d+2 points of a polynomial of

→ degree d, it finds

                                                    d in roughly O(n log d). Note, n should not exceed mod
                                                     way modular inverse is used. In such cases, we can use
                                                        interpolation
                                                     without modulo in big integers and take the remainder

    later
Time: 0 (n) */
                                                    using namespace std;
                                                    struct Lagrange {
                                                         vector<int> terms, dp;
                                                         int mod, n;
                                                         Lagrange() {}
                                                         Lagrange(const vector<int>& terms, int mod) :
                                                             terms(terms), mod(mod) {
                                                              n = terms.size();
                                                              assert(n <= mod);
                                                              int i, v, f;
for (f = 1, i = 1; i < n; i++) f = (long</pre>
                                                               → long)f * i % mod;
                                                              v = expo(f, mod - 2);
                                                              vector<int> inv(n, v);
for (i = n - 1; i > 0; i--) {
   inv[i - 1] = (long long)inv[i] * i % mod;
                                                              dp.resize(n, 1);
for (i = 0; i < n; i++) {
    dp[i] = (long long)inv[i] * inv[n - i - 1]</pre>
                                                                   dp[i] = (long long)dp[i] * terms[i] % mod;
                                                          int expo(int a, int b) {
                                                              int res = 1;
                                                                   if (b \& 1) res = (long long) res * a % mod;
                                                                   a = (long long)a * a % mod;
                                                                   b >>= 1:
                                                              return res:
                                                         int interpolate(long long x) -
                                                              if (x < n) return terms[x] % mod;</pre>
                                                              x %= mod:
                                                              int i, w;
                                                              vector<int> X(n, 1), Y(n, 1);
for (i = 1; i < n; i++) {
    X[i] = (long long)X[i - 1] * (x - i + 1) %</pre>
                                                                   if (X[i] < 0) X[i] += mod;
                                                              for (i = n - 2; i >= 0; i--)
                                                                   \dot{Y}[i] = (long long)\dot{Y}[i + 1] * (x - i - 1) %
                                                                   if (Y[i] < 0) Y[i] += mod;
                                                               long long res = 0;
                                                              for (i = 0; i < n; i++) {
    w = ((long long)X[i] * Y[i] % mod) * dp[i]</pre>
                                                                        % mod;
                                                                   return res % mod;
```

```
vector<int> get terms(const vector<int>& terms, int
   mod, int l, int r) {
    auto lagrange = Lagrange(terms, mod);
    vector<int> res;
    for (int i = l; i <= r; i++) {
        res.push back(lagrange.interpolate(i));
    return res;
int find degree(const vector<int>& terms. int mod) {
    long long v = mod;
    int k = 1, n = terms.size();
while (v < INT_MAX) {</pre>
        v *= mod:
        k++;
    int l = 1 << 30, r = l + k - 1;
    auto expected = get terms(terms, mod, l, r);
    int low = 1, high = n - 1;
    while ((low + 1) < high)
        int mid = (low + high) >> 1;
        vector<int> v(terms.begin(), terms.begin() +
        if (get terms(v, mod, l, r) == expected)
            high = mid;
       else
low = mid;
    for (int d = low; d <= high; d++) {
        vector<int> v(terms.begin(), terms.begin() +
        if (get terms(v, mod, l, r) == expected)
        → return d - 1:
    return -1;
int main() {
    const int mod = 10000000007;
    vector<int> terms = vector<int>{0, 1, 5, 14, 30};
    auto lagrange = Lagrange(terms, mod);
    assert(lagrange.interpolate(5) == 55);
    assert(lagrange.interpolate(6) == 91);
    assert(lagrange.interpolate(1 << 30) == 300663155);</pre>
    assert(lagrange.interpolate(1LL << 60) ==
    \sim 717860166);
    assert(find degree(terms, mod) == 3);
    terms.pop back();
    assert(find degree(terms, mod) == -1);
    return 0;
```

#### 5.17 matrix expo

```
11 \mod = (1e9) + 7;
struct Matrix{
    int row, col;
    vector<vector<ll>>> mat;
    Matrix(int x, int y){
        row=x;
        mat.assign(row, vector<ll>(col,0));
    Matrix operator *(Matrix &other){
        assert(col==other.row);
        Matrix product(row,other.col);
        for(int i=0;i<row;i++){</pre>
             for(int j=0; j<col; j++){
                 for(int k=0; k<other.col; k++){</pre>
```

k]+(mat[i][j]\*other.mat[j][k])

product.mat[i][k]=(product.mat[i][

%mod)%mod;

```
}
         return product;
};
Matrix expo(Matrix &m, ll n){
    assert(m.row==m.col);
    Matrix ret(m.row,m.col);
    for(int i=0;i<m.row;i++) ret.mat[i][i]=1;</pre>
    while(n){
         if(n\&1) ret=ret*m;
         n/=2:
        m=m*m;
    return ret;
5.18 next_lexicographical_k_comb
bool next combination(vector<int>& a, int n) {
    int k = (int)a.size();
    for (int i = k - 1; i >= 0; i --) {
         if (a[i] < n - k + i + 1) {
             `a[i]++;
             for (int j = i + 1; j < k; j++)

a[j] = a[j - 1] + 1;
             return true:
    return false;
5.19 ntt
const int mod = 998244353;
const int root = 15311432;
const int k = 1 \ll 23;
int root 1;
vector<int> rev;
void pre(int sz){
    root 1 = bigmod(root, mod - 2, mod);
    if (rev.size() == sz) return;
    rev.resize(sz);
    rev[0] = 0;
    int lg n = builtin ctz(sz);
    for (int i = 1; i < \overline{sz}; ++i)
     \rightarrow rev[i]=(rev[i>>1]>>1)|((i&1)<<(lg n-1));
void fft(vector<int> &a, bool inv){
    int n = a.size();
for (int i = 1; i < n - 1; ++i) if (i < rev[i])</pre>

    swap(a[i], a[rev[i]]);

    for (int len = 2; len <= n; len <<= 1) {
         int wlen = inv ? root 1 : root;
         for (int i = len; i < k; i <<= 1) wlen = 1ll *</pre>

→ wlen * wlen % mod;

         for (int st = 0; st < n; st += len){
             int w = 1;
             for (int j = 0; j < len / 2; j++){
                  int ev = a[st + j];
                  int od = 111 * a[st + j + len / 2] * w
                  a[st + j] = ev + od < mod ? ev + od :
                  \rightarrow ev + od - mod;
                  a[st + j + len / 2] = ev - od >= 0?

    ev - od : ev - od + mod;
w = 1ll * w * wlen % mod;

        }
```

```
if (inv){
         int n 1 = bigmod(n, mod - 2, mod);
         for (int \&x : a) x = 111 * x * n 1 % mod;
vector<int> mul(vector<int> &a, vector<int> &b){
    int n = a.size(), m = b.size(), sz = 1;
while (sz < n + m - 1) sz <<= 1;</pre>
    vector<int> x(sz), y(sz), z(sz);
    for (int i = 0; i < sz; ++i){
    x[i] = i < n ? a[i] : 0;
    y[i] = i < m ? b[i] : 0;</pre>
    pre(sz);fft(x, 0);fft(y, 0);
     for (int i = 0: i < sz: ++i) z[i] = 1ll * x[i] *

→ y[i] % mod;

    fft(z, 1);z.resize(n + m - 1);
return z;
5.20 seg_sieve
Segmented Sieve
This code was for 1 \le a \le b \le 2^31-1
Change variable types appropriately.
bool notPrime[ ? ];
void segmented sieve(int a, int b)
    int p, f;
    mem(notPrime, 0);
for (int i = 0; i < tot_prime; i++)</pre>
         p = prime[i];
         if (a \% p == 0)
              f = a:
              f = (a - (a % p) + p);
         f = max(p * p, f);
         for (unsigned j = f; j \le b; j += p)
             notPrime[j - a] = true;
    if (a == 1)
         notPrime[0] = 1;
5.21 stirling
that no box is empty.
int stirling2(int n, int k)
 if(n < k)
 return 0;
 if(k == 1)
 return 1;
 if(dp[n][k] == dp[n][k])
 return dp[n][k];
 return dp[n][k] = stirling2(n-1,k-1) +

    stirling2(n-1,k)*k;

|int stirling1(int n, int k)
 dp[n][k] = stirling1(n-1,k-1) + stirling(n-1,k)*n-1;
5.22 stirling_number_of_the_second_kind
     / k! * sum (-1)^i nCr(k, i) * (k - i) ^ n
ĺĺ f(int n, int k) {
```

ll res = 0;

for (int i = 0; i < k; ++i) {

```
else res = (res + nCr(k, i) * bp(k - i, n,

→ mod) % mod) % mod;

    if (res < 0) res += mod;
    return res * ifac[k] % mod;
5.23 sum_of_totient
using namespace gnu pbds;
const int N = 3e5 + 9, mod = 998244353;
template <const int32 t MOD>
struct modint {
    int32 t value;
    modin\overline{t}() = default;
    modint(int32 t value ) : value(value ) {}
    inline modint<MOD> operator+(modint<MOD> other)
        int32 t c = this->value + other.value;
        return modint<MOD>(c >= MOD ? c - MOD : c);
    inline modint<MOD> operator-(modint<MOD> other)
       const {
        int32 t c = this->value - other.value;
        return modint<MOD>(c < 0 ? c + MOD : c);
    inline modint<MOD> operator*(modint<MOD> other)
        int32 t c = (int64 t)this->value * other.value
        return modint<MOD>(c < 0 ? c + MOD : c);
    inline modint<MOD> &operator+=(modint<MOD> other) {
        this->value += other.value;
        if (this->value >= MOD)
            this->value -= MOD:
        return *this;
    inline modint<MOD> &operator-=(modint<MOD> other) {
        this->value -= other.value;
        if (this->value < 0)</pre>
            this->value += MOD;
        return *this:
   inline modint<MOD> &operator*=(modint<MOD> other) {
        this->value = (int64 t)this->value *

    other.value %
    MOD;

        if (this->value < 0) this->value += MOD;
        return *this;
    inline modint<MOD> operator-() const { return
        modint<MOD>(this->value ? MOD - this->value :
       0); }
    modint<MOD> pow(uint64_t k) const {
        modint<MOD> x =
                    y = 1:
        for (; k; k >>= 1) {
            if (k \& 1) y *= x;
            x *= x;
        return y;
    modint<MOD> inv() const { return pow(MOD - 2); }
       // MOD must be a prime
    inline modint<MOD> operator/(modint<MOD> other)

    const { return *this * other.inv(); }
```

**if** (i & 1) res = (res - nCr(k, i) \* bp(k

→ n, mod) % mod + mod) % mod;

```
inline modint<MOD> operator/=(modint<MOD> other) {
     → return *this *= other.inv(); }
    inline bool operator==(modint<MOD> other) const {
        return value == other.value; }
    inline bool operator!=(modint<MOD> other) const {
       return value != other.value; }
    inline bool operator<(modint<MOD> other) const {
     → return value < other.value; }</pre>
    inline bool operator>(modint<MOD> other) const {

→ return value > other.value: }

template <int32 t MOD>
modint<MOD> operator*(int64 t value, modint<MOD> n) {

¬ return modint<MOD>(value) * n; }

template <int32 t MOD>
modint<MOD> operator*(int32_t value, modint<MOD> n) {

    return modint<MOD>(value % MOD) * n; }

template <int32 t MOD>
istream &operator>>(istream &in, modint<MOD> &n) {

→ return in >> n.value; }

template <int32 t MOD>
ostream & operator << (ostream & out, modint < MOD> n) {

    return out << n.value; }
</pre>
using mint = modint<mod>;
namespace Dirichlet {
// solution for f(x) = phi(x)
const int T = 1e7 + 9;
long long phi[T];
gp hash table<long long, mint> mp;
mint dp[T], inv;
int sz, spf[T], prime[T];
void init() {
    memset(spf, 0, sizeof spf);
    phi[1] = 1;
    sz = 0;
    for (int i = 2; i < T; i++) {
        if (spf[i] == 0) phi[i] = i - 1, spf[i] = i,
         \rightarrow prime[sz++] = i;
        for (int j = 0; j < sz && i * prime[j] < T &&
            prime[j] <= spf[i]; j++)</pre>
            spf[i * prime[j]] = prime[j];
            if (i % prime[j] == 0)
                 phi[i * prime[j]] = phi[i] * prime[j];
            else
                 phi[i * prime[j]] = phi[i] * (prime[j]
                 \rightarrow - 1);
    dp[0] = 0;
    for (int i = 1; i < T; i++) dp[i] = dp[i - 1] +
     → phi[i] % mod;
    inv = 1; // q(1)
mint p c(long long n) {
    if (n \% 2 == 0) return n / 2 \% mod * ((n + 1) \%

→ mod) % mod;

    return (n + 1) / 2 % mod * (n % mod) % mod;
mint p q(long long n) {
    return n % mod:
mint solve(long long x)
    if (x < T) return dp[x];</pre>
    if (mp.find(x) != mp.end()) return mp[x];
    mint ans = 0:
    for (long long i = 2, last; i <= x; i = last + 1) {
    last = x / (x / i);</pre>
        ans += solve(x'/i) * (p_g(last) - p_g(i - 1));
    ans = p c(x) - ans;
```

```
ans /= inv;
     return mp[x] = ans;
    // namespace Dirichlet
 5.24 totient
 int phi(int n) {
     int result = n;
     for (int i = 2; i * i <= n; i++) {
         if (n % i == 0) {
             while (n % i == 0)
                  n /= i;
             result -= result / i;
     if (n > 1)
         result _= result / n;
     return result;
|void phi 1 to n(int n) {
     vector < int > phi(n + 1);
     phi[0] = 0;
phi[1] = 1;
     for (int i = 2; i \le n; i++)
         phi[i] = i;
     for (int i = 2; i \le n; i++) {
         if (phi[i] == i) {
             for (int j = i; j <= n; j += i)
                  phi[j] -= phi[j] / i;
5.25 xor_basis
int basis[d]; // basis[i] keeps the mask of the vector
 → whose f value is i
int sz;
void insertVector(int mask) {
 for (int i = 0; i < d; i++) {
  if ((mask \& 1 << i)) == 0) continue;
   if (!basis[i]) { // If there is no basis vector with
       the i'th bit set, then insert this vector into

    the basis
basis[i] = mask;

    return:
  mask ^= basis[i]; // Otherwise subtract the basis

    vector from this vector

6 Misc
6.1 check
q++ a.cpp -o a.out && q++ ac.cpp -o ac.out && q++
 \rightarrow gen.cpp -o gen && for ((i=0;i<1000;i++))
do
echo $i
 ./gen > in
 ./a.out < inp > out
 ./ac.out < inp > out1
diff out1 out2
if [ $? -ne 0 ]
Ithen
echo "----Input----"
cat in echo "-----"
cat out echo "----Accepted----"
cat out1
break
done
```

```
6.2 debug
```

#### 6.3 interval\_container

```
/* Description: Add and remove intervals from a set of
    disjoint intervals. Will merge the added interval
   with any overlapping intervals in
the set when adding. Intervals are [inclusive,

    exclusive).
Time: 0 (log N)
set<pii>::iterator addInterval(set<pii>& is, int L,
    int R) {
    if (L == R) return is.end();
auto it = is.lower_bound({L, R}), before = it;
    while (it != is.en\overline{d}() && it->first <= R) {
         R = max(R, it->second);
         before = it = is.erase(it);
    if (it != is.begin() && (--it)->second >= L) {
        L = min(L, it->first);
R = max(R, it->second); is.erase(it);
    return is.insert(before, {L, R});
void removeInterval(set<pii>& is, int L, int R) {
    if (L == R) return;
    autò it = áddInterval(is, L, R);
    auto r2 = it->second;
    if (it->first == L)
         is.erase(it);
         (int&)it->second = L;
    if (R != r2) is.emplace(R, r2);
```

#### 6.4 interval\_cover

```
at++:
     if (mx.second == -1) return {};
cur = mx.first;
     R.push back(mx.second);
return R;
```

#### 6.5 pragma

```
// Praamas
#pragma comment(linker, "/stack:200000000")
#pragma GCC optimize("03,unroll-loops")
#pragma GCC target("avx,avx2,fma")
```

#### 6.6 random

```
// shuffle(v.begin(), v.end(),
    default random engine(rnd(1, 1000)));
mt19937 rng(chrono::steady clock::now().time since epo
ch().count());
ll rnd(ll l, ll r)
    return uniform int distribution<ll>(l, r) (rng);
```

#### 6.7 vimre

```
imap jk <Esc>
set mouse=a
set autoindent
set tabstop=4
set shiftwidth=4
    smartindent
    relativenumber
     .aststatus=2
    hlsearch
let mapleader =
nnoremap <leader>s :w<Enter>
nnoremap <leader>y ggVG"+y<CR>
syntax on
vnoremap <leader>/ :s!^!//!<CR> :noh <CR>
vnoremap <leader>u :s!^//!!<CR>
nnoremap <leader>/ :s!^!//!<CR> :noh <CR>
nnoremap <leader>u :s!^//!!<CR>
```

#### 7 Notes

#### 7.1 Counting

#### 1. Lucas Theorem

For non-negative integers m and n and a prime p, the following congruence relation holds: :

$$\begin{pmatrix} m \\ n \end{pmatrix} \equiv \prod_{i=0}^{k} \begin{pmatrix} m_i \\ n_i^i \end{pmatrix} \pmod{p}, \tag{1}$$

where:

$$m = m_k p^k + m_{k-1} p^{k-1} + \dots + m_1 p + m_0,$$

and:

$$n = n_k p^k + n_{k-1} p^{k-1} + \dots + n_1 p + n_0$$

are the base p expansions of m and n respectively. This uses the convention that  $\binom{m}{n} = 0$  if  $m \le n$ .

#### 2. Stirling Numbers of the first kind

S(n,k) counts the number of permutations of n elements with k disjoint cycles.

$$S(n,k) = (n-1) \cdot S(n-1,k) + S(n-1,k-1)$$
 (2)

where, S(0,0) = 1, S(n,0) = S(0,n) = 0

$$\sum_{k=0}^{n} S(n,k) = n! \tag{3}$$

## 3. Stirling Numbers of the second kind

 $S(n,k) \cdot k! = \text{number of ways to color } n \text{ nodes using colors from } 1 \text{ to } k$ such that each color is used at least once.

An r-associated Stirling number of the second kind is the number of ways to partition a set of n objects into k subsets, with each subset containing at least r elements. It is denoted by  $S_r(n,k)$  and obeys the recurrence relation.

$$S_r(n+1,k) = kS_r(n,k) + \binom{n}{r-1}S_r(n-r+1,k-1)$$
 (4)

#### 4. Bell Numbers

Counts the number of partitions of a set.

$$B_{n+1} = \sum_{k=0}^{n} \left(\frac{n}{k}\right) \cdot B_k \tag{5}$$

 $B_n = \sum_{k=0}^n S(n,k)$ , where S(n,k) is stirling number of second kind. 5. Some identities

**Vandermonde's Identify:**  $\sum_{k=0}^{r} {m \choose k} {n \choose r-k} = {m+n \choose r}$ 

**Hockey-Stick Identify**: 
$$n, r \in N, n > r, \sum_{i=r}^{n} \begin{pmatrix} i \\ r \end{pmatrix} = \begin{pmatrix} n+1 \\ r+1 \end{pmatrix}$$

**Involutions**: permutations such that  $p^2$  = identity permutation.  $a_0 = a_1 = 1$  and  $a_n = a_{n-1} + (n-1)a_{n-2}$  for n > 1.

#### 7.2 Fibonacci

Let A.B and n be integer numbers.

$$k = A - B$$

$$F_A F_B = F_{k+1} F_A^2 + F_k F_A F_{A-1} \tag{6}$$

$$\sum_{i=0}^{n} F_i^2 = F_{n+1} F_n \tag{7}$$

$$\sum_{i=0}^{n} F_i F_{i+1} = F_{n+1}^2 - (-1)^n \tag{8}$$

$$\sum_{i=0}^{n} F_i F_{i-1} = \sum_{i=0}^{n-1} F_i F_{i+1}$$
 (9)

$$GCD(F_m, F_n) = F_{GCD(m,n)}$$
(10)

$$\sum_{0 \le k \le n} \binom{n-k}{k} = Fib_{n+1} \tag{11}$$

$$\gcd(F_n, F_{n+1}) = \gcd(F_n, F_{n+2}) = \gcd(F_{n+1}, F_{n+2}) = 1 \tag{12}$$

#### 7.3 Notes

#### 7.4 Geometry

#### 7.4.1 Triangles

Circumradius:  $R = \frac{abc}{4A}$ , Inradius:  $r = \frac{A}{s}$ 

Length of median (divides triangle into two equal-area triangles):  $m_a = \frac{1}{2}\sqrt{2b^2 + 2c^2 - a^2}$ 

Length of bisector (divides angles in two):  $s_a = \sqrt{bc \left[1 - \left(\frac{a}{b+c}\right)^2\right]}$ 

Law of tangents: 
$$\frac{a+b}{a-b} = \frac{\tan \frac{\alpha+\beta}{2}}{\tan \frac{\alpha-\beta}{2}}$$

#### 7.4.2 Quadrilaterals

With side lengths a,b,c,d, diagonals e,f, diagonals angle  $\theta$ , area A and magic flux  $F = b^2 + d^2 - a^2 - c^2$ .

$$4A = 2ef \cdot \sin\theta = F \tan\theta = \sqrt{4e^2f^2 - F^2}$$

For cyclic quadrilaterals the sum of opposite angles is  $180^{\circ}$ , ef =(4) ac + bd, and  $A = \sqrt{(p-a)(p-b)(p-c)(p-d)}$ .

#### 7.4.3 Spherical coordinates

$$x = r \sin \theta \cos \phi \qquad r = \sqrt{x^2 + y^2 + z^2}$$

$$y = r \sin \theta \sin \phi \qquad \theta = a\cos(z/\sqrt{x^2 + y^2 + z^2})$$

$$z = r \cos \theta \qquad \phi = a\tan(2y, x)$$

#### **7.5** Sums

$$1^{2} + 2^{2} + 3^{2} + \dots + n^{2} = \frac{n(2n+1)(n+1)}{6}$$

$$1^{3} + 2^{3} + 3^{3} + \dots + n^{3} = \frac{n^{2}(n+1)^{2}}{4}$$

$$1^{4} + 2^{4} + 3^{4} + \dots + n^{4} = \frac{n(n+1)(2n+1)(3n^{2} + 3n - 1)}{30}$$

$$\sum_{i=1}^{n} i^{m} = \frac{1}{m+1} \left[ (n+1)^{m+1} - 1 - \sum_{i=1}^{n} \left( (i+1)^{m+1} - i^{m+1} - (m+1)i^{m} \right) \right]$$

$$\sum_{i=1}^{n-1} i^{m} = \frac{1}{m+1} \sum_{k=0}^{m} {m+1 \choose k} B_{k} n^{m+1-k}$$

$$\sum_{k=0}^{n} kx^{k} = (x - (n+1)x^{n+1} + nx^{n+2})/(x-1)^{2}$$

#### 7.6 Series

$$e^{x} = 1 + x + \frac{x^{2}}{2!} + \frac{x^{3}}{3!} + \dots, (-\infty < x < \infty)$$
$$\ln(1+x) = x - \frac{x^{2}}{2!} + \frac{x^{3}}{2!} - \frac{x^{4}}{4!} + \dots, (-1 < x \le 1)$$

$$\sqrt{1+x} = 1 + \frac{x}{2} - \frac{x^2}{8} + \frac{2x^3}{32} - \frac{5x^4}{128} + \dots, (-1 \le x \le 1)$$

$$\sin x = x - \frac{x^3}{3!} + \frac{x^5}{5!} - \frac{x^7}{7!} + \dots, (-\infty < x < \infty)$$

$$\cos x = 1 - \frac{x^2}{2!} + \frac{x^4}{4!} - \frac{x^6}{6!} + \dots, (-\infty < x < \infty)$$

$$(x+a)^{-n} = \sum_{k=0}^{\infty} (-1)^k \binom{n+k-1}{k} x^k a^{-n-k}$$

#### 7.7 Pythagorean Triples

The Pythagorean triples are uniquely generated by

$$a = k \cdot (m^2 - n^2), b = k \cdot (2mn), c = k \cdot (m^2 + n^2),$$

with m > n > 0, k > 0,  $m \perp n$ , and either m or n even.

#### 7.8 Number Theory

#### 7.8.1 Primes

p = 962592769 is such that  $2^{21} \mid p - 1$ , which may be useful. For hashing use 970592641 (31-bit number), 31443539979727 (45-bit), 3006703054056749 (52-bit). There are 78498 primes less than 1000000.

Primitive roots exist modulo any prime power  $p^a$ , except for p = 2, a > a2, and there are  $\phi(\phi(p^a))$  many. For p=2, a>2, the group  $\mathbb{Z}_{2a}^{\times}$  is instead isomorphic to  $\mathbb{Z}_2 \times \mathbb{Z}_{2^{a-2}}$ .

#### 7.8.2 Estimates

 $\sum_{d|n} d = O(n \log \log n).$ 

The number of divisors of n is at most around 100 for n < 5e4, 500 for n < 1e7, 2000 for n < 1e10, 200 000 for n < 1e19.

#### 7.8.3 Perfect numbers

n > 1 is called perfect if it equals sum of its proper divisors and 1. Even n is perfect iff  $n = 2^{p-1}(2^p - 1)$  and  $2^p - 1$  is prime (Mersenne's). No odd perfect numbers are vet found.

#### 7.8.4 Carmichael numbers

A positive composite *n* is a Carmichael number  $(a^{n-1} \equiv 1 \pmod{n})$  for all  $a: \gcd(a,n)=1$ ), iff n is square-free, and for all prime divisors p of n, p-1 divides n-1.

#### 7.8.5 Mobius function

 $\mu(1) = 1$ .  $\mu(n) = 0$ , if n is not squarefree.  $\mu(n) = (-1)^s$ , if n is the product of s distinct primes. Let f, F be functions on positive integers. If for all  $n \in N$ ,  $F(n) = \sum_{d|n} f(d)$ , then  $f(n) = \sum_{d|n} \mu(d) F(\frac{n}{d})$ , and vice versa  $\phi(n) = \sum_{d|n} \mu(d) \frac{n}{d}. \quad \sum_{d|n} \mu(d) = 1.$ 

If f is multiplicative, then  $\sum_{d|n} \mu(d) f(d) = \prod_{p|n} (1 - f(p))$  $\sum_{d|n} \mu(d)^2 f(d) = \prod_{p|n} (1 + f(p)).$ 

#### 7.8.6 Legendre symbol

If p is an odd prime,  $a \in \mathbb{Z}$ , then  $\left(\frac{a}{p}\right)$  equals 0, if p|a; 1 if a is a quadratic residue modulo p; and -1 otherwise. Euler's criterion  $\left(\frac{a}{p}\right) = a^{\left(\frac{p-1}{2}\right)} \pmod{p}.$ 

#### 7.8.7 Jacobi symbol

If 
$$n = p_1^{a_1} \cdots p_k^{a_k}$$
 is odd, then  $\left(\frac{a}{n}\right) = \prod_{i=1}^k \left(\frac{a}{p_i}\right)^{k_i}$ .

#### 7.8.8 Primitive roots

If the order of g modulo m (min n > 0:  $g^n \equiv 1 \pmod{m}$ ) is  $\phi(m)$ , then g is called a primitive root. If  $Z_m$  has a primitive root, then it has  $\phi(\phi(m))$  distinct primitive roots.  $Z_m$  has a primitive root iff m is one where  $X^g$  are the elements fixed by g(g.x = x). of 2, 4,  $p^k$ ,  $2p^k$ , where p is an odd prime. If  $Z_m$  has a primitive root g, then for all a coprime to m, there exists unique integer  $i = \operatorname{ind}_g(a)$ modulo  $\phi(m)$ , such that  $g^i \equiv a \pmod{m}$ . ind  $\sigma(a)$  has logarithm-like properties: ind(1) = 0, ind(ab) = ind(a) + ind(b).

If p is prime and a is not divisible by p, then congruence  $x^n \equiv a$ (mod p) has  $\gcd(n, p-1)$  solutions if  $a^{(p-1)/\gcd(n, p-1)} \equiv 1 \pmod{p}$ , and no solutions otherwise. (Proof sketch: let g be a primitive root, and 7.10 Partitions and subsets  $g^i \equiv a \pmod{p}, \ g^u \equiv x \pmod{p}. \ x^n \equiv a \pmod{p} \text{ iff } g^{nu} \equiv g^i \pmod{p}$  7.10.1 Partition function iff  $nu \equiv i \pmod{p}$ .

#### 7.8.9 Discrete logarithm problem

Find x from  $a^x \equiv b \pmod{m}$ . Can be solved in  $O(\sqrt{m})$  time and space with a meet-in-the-middle trick. Let  $n = \lceil \sqrt{m} \rceil$ , and x = ny - z. Equation becomes  $a^{ny} \equiv ba^z \pmod{m}$ . Precompute all values that the RHS can take for z = 0, 1, ..., n - 1, and brute force y on the LHS, each time checking whether there's a corresponding value for RHS.

#### 7.8.10 Pythagorean triples

Integer solutions of  $x^2 + y^2 = z^2$  All relatively prime triples are given by:  $x = 2mn, y = m^2 - n^2, z = m^2 + n^2$  where  $m > n, \gcd(m, n) = 1$  and  $m \not\equiv n \pmod{2}$ . All other triples are multiples of these. Equation  $x^2 + y^2 = 2z^2$  is equivalent to  $(\frac{x+y}{2})^2 + (\frac{x-y}{2})^2 = z^2$ .

#### 7.8.11 Postage stamps/McNuggets problem

Let a, b be relatively-prime integers. There are exactly  $\frac{1}{2}(a-1)(b-1)$ numbers not of form ax + by  $(x, y \ge 0)$ , and the largest is (a-1)(b-1)-1 = ab - a - b.

#### 7.8.12 Fermat's two-squares theorem

Odd prime p can be represented as a sum of two squares iff  $p \equiv$ 1 (mod 4). A product of two sums of two squares is a sum of two squares. Thus, n is a sum of two squares iff every prime of form c(8,k) = 8,0,5040,13068,13132,6769,1960,322,28,1p = 4k + 3 occurs an even number of times in *n*'s factorization.

#### 7.9 Permutations

#### 7.9.1 Factorial

| 1 40001141 |                                    |                   |        |             |              |             |
|------------|------------------------------------|-------------------|--------|-------------|--------------|-------------|
| n          | 1234                               | 5 6               | 7      | 8           | 9            | 10          |
| n!         | 1 2 6 24                           | 120 720           | 5040   | $40320_{-}$ | 362880       | 3628800     |
| $_{-}n$    | 11 1                               | <u>.2 13 </u>     | 14     | 15          | 16           | 17          |
| n!         | 4.0e7 4.8                          | 3e8 6.2e9         | 9 8.7e | 10 1.3e     | 12 2.1e      | 13 3.6e14   |
| n          | 20 25                              | 5 30              | 40     | 50 1        | 00 15        | 0 171       |
| n!         | $\mid 2\mathrm{e}18\ 2\mathrm{e}2$ | $25 \; 3e32 \; 8$ | 3e47.3 | 664 9e      | $157 \; 6e2$ | 62 >DBL_MAX |

#### 7.9.2 Cycles

Let  $g_S(n)$  be the number of n-permutations whose cycle lengths all belong to the set S. Then

$$\sum_{n=0}^{\infty} g_S(n) \frac{x^n}{n!} = \exp\left(\sum_{n \in S} \frac{x^n}{n}\right)$$

#### 7.9.3 Derangements

Permutations of a set such that none of the elements appear in their loriginal position.

$$D(n) = (n-1)(D(n-1) + D(n-2)) = nD(n-1) + (-1)^n = \left\lfloor \frac{n!}{e} \right\rfloor$$

#### 7.9.4 Burnside's lemma

Given a group G of symmetries and a set X, the number of elements of *X up to symmetry* equals

$$\frac{1}{|G|} \sum_{g \in G} |X^g|,$$

If f(n) counts "configurations" (of some sort) of length n, we can ignore rotational symmetry using  $G = \mathbb{Z}_n$  to get

$$g(n) = \frac{1}{n} \sum_{k=0}^{n-1} f(\gcd(n,k)) = \frac{1}{n} \sum_{k|n} f(k)\phi(n/k)$$

Number of ways of writing n as a sum of positive integers, disregarding the order of the summands.

$$p(0) = 1, \ p(n) = \sum_{k \in \mathbb{Z} \setminus \{0\}} (-1)^{k+1} p(n - k(3k - 1)/2)$$

$$p(n) \sim 0.145/n \cdot \exp(2.56\sqrt{n})$$

#### 7.11 General purpose numbers

#### 7.11.1 Stirling numbers of the first kind

Number of permutations on n items with k cycles.

$$c(n,k) = c(n-1,k-1) + (n-1)c(n-1,k), \ c(0,0) = 1$$
  
$$\sum_{k=0}^{n} c(n,k)x^{k} = x(x+1)...(x+n-1)$$

c(n,2) = 0,0,1,3,11,50,274,1764,13068,109584,...

#### 7.11.2 Eulerian numbers

Number of permutations  $\pi \in S_n$  in which exactly k elements are greater than the previous element. k j:s s.t.  $\pi(j) > \pi(j+1)$ , k+1 j:s s.t.  $\pi(j) \ge j$ , k j:s s.t.  $\pi(j) > j$ .

$$E(n,k) = (n-k)E(n-1,k-1) + (k+1)E(n-1,k)$$

$$E(n,0) = E(n,n-1) = 1$$

$$E(n,k) = \sum_{j=0}^{k} (-1)^{j} \binom{n+1}{j} (k+1-j)^{n}$$

#### 7.11.3 Stirling numbers of the second kind

Partitions of n distinct elements into exactly k groups.

$$S(n,k) = S(n-1,k-1) + kS(n-1,k)$$

$$S(n,1) = S(n,n) = 1$$

$$S(n,k) = \frac{1}{k!} \sum_{j=0}^{k} (-1)^{k-j} \binom{k}{j} j^n$$

#### 7.11.4 Bell numbers

Total number of partitions of n distinct elements. B(n) = $1, 1, 2, 5, 15, 52, 203, 877, 4140, 21147, \dots$  For p prime,

$$B(p^m + n) \equiv mB(n) + B(n+1) \pmod{p}$$

#### 7.11.5 Bernoulli numbers

$$\sum_{j=0}^{m} {m+1 \choose j} B_j = 0. \quad B_0 = 1, B_1 = -\frac{1}{2}. \ B_n = 0, \text{ for all odd } n \neq 1.$$

#### 7.11.6 Catalan numbers

$$C_n = \frac{1}{n+1} {2n \choose n} = {2n \choose n} - {2n \choose n+1} = \frac{(2n)!}{(n+1)!n!}$$

$$C_0 = 1$$
,  $C_{n+1} = \frac{2(2n+1)}{n+2}C_n$ ,  $C_{n+1} = \sum C_i C_{n-i}$ 

 $C_n = 1, 1, 2, 5, 14, 42, 132, 429, 1430, 4862, 16796, 58786, \dots$ 

- sub-diagonal monotone paths in an  $n \times n$  grid.
- strings with *n* pairs of parenthesis, correctly nested.
- binary trees with with n + 1 leaves (0 or 2 children).
- ordered trees with n + 1 vertices.
- ways a convex polygon with n+2 sides can be cut into triangles by connecting vertices with straight lines.
- permutations of [n] with no 3-term increasing subseq.

#### 7.12 Inequalities

#### 7.12.1 Titu's Lemma

For positive reals  $a_1, a_2, \dots, a_n$  and  $b_1, b_2, \dots, b_n$ ,

$$\frac{a_1^2}{b_1} + \frac{a_2^2}{b_2} + \ldots + \frac{a_n^2}{b_n} \ge \frac{a_1 + a_2 + \ldots + a_n^2}{b_1 + b_2 + \ldots + b_n}$$

Equality holds if and only if  $a_i = kb_i$  for a non-zero real constant k.

#### **7.13 Games** 7.13.1 Grundy numbers

For a two-player, normal-play (last to move wins) game on a graph  $(V,E): G(x) = \max(\{G(y): (x,y) \in E\}), \text{ where } \max(S) = \min\{n \ge 0: n \not\in E\}$ S. x is losing iff G(x) = 0.

- **7.13.2** Sums of games Player chooses a game and makes a move in it Grundy number of a position is xor of grundy numbers of positions in summed games.
  - Player chooses a non-empty subset of games (possibly, all) and makes moves in all of them A position is losing iff each game is in a losing position.
  - Player chooses a proper subset of games (not empty and not all), and makes moves in all chosen ones. A position is losing iff grundy numbers of all games are equal.
  - Player must move in all games, and loses if can't move in some game A position is losing if any of the games is in a losing position.

#### 7.13.3 Misère Nim

A position with pile sizes  $a_1, a_2, ..., a_n \ge 1$ , not all equal to 1, is losing iff  $a_1 \oplus a_2 \oplus \cdots \oplus a_n = 0$  (like in normal nim.) A position with n piles of size 1 is losing iff n is odd.

#### 7.14 NumberTheory

$$\sum_{k=1}^{n} \frac{1}{\gcd(k,n)} = \sum_{d|n} \frac{1}{d} \cdot \phi\left(\frac{n}{d}\right) = \frac{1}{n} \sum_{d|n} d \cdot \phi(d)$$
 (13)

$$\sum_{k=1}^{n} \frac{k}{\gcd(k,n)} = \frac{n}{2} \cdot \frac{1}{n} \cdot \sum_{d \mid n} d \cdot \phi(d)$$
 (14)

$$\sum_{k=1}^{n} \frac{n}{\gcd(k,n)} = 2 * \sum_{k=1}^{n} \frac{k}{\gcd(k,n)} - 1, \text{ for, } n > 1$$
 (15)

$$\sum_{i=1}^{n} \sum_{j=1}^{n} [\gcd(i,j) = 1] = \sum_{d=1}^{n} \mu(d) \left\lfloor \frac{n}{d} \right\rfloor^{2}$$
 (16)

$$\sum_{i=1}^{n} \sum_{j=1}^{n} \gcd(i,j) = \sum_{d=1}^{n} \phi(d) \left\lfloor \frac{n}{d} \right\rfloor^{2}$$
(17)

$$\sum_{i=1}^{n} \sum_{j=1}^{n} i \cdot j[\gcd(i,j) = 1] = \sum_{i=1}^{n} \phi(i)i^{2}$$
(18)

$$\sum_{i=1}^{n} \sum_{j=1}^{n} \operatorname{lcm}(i,j) = \sum_{l=1}^{n} \left( \frac{\left(1 + \left\lfloor \frac{n}{l} \right\rfloor\right) \left(\left\lfloor \frac{n}{l} \right\rfloor\right)}{2} \right)^{2} = \sum_{d|l} \mu(d) l d$$
 (19)

#### 8 String

#### 8.1 aho corasick

```
using namespace std;
const int N = ?;
const int A = ?;
struct AC {
  int nd, pt;
  int next[N][A], link[N], out_link[N], cnt[N], ans[N];
  vector <int> ed[N], out[N];
  AC(): nd(0), pt(0) { node(); }
  int node() {
    memset(next[nd], 0, sizeof next[nd]);
    link[nd] = out_link[nd] = cnt[nd] = 0;
    ed[nd].clear(), out[nd].clear();
    return nd++;
}
void clear() {
    nd = pt = 0;
    node();
}
```

```
inline int get(char c) { return c - 'a'; }
       void insert(const string &T) {
          int u = 0:
          for (char'c : T) {
            if (!next[u][get(c)]) next[u][get(c)] = node();
            u = next[u][qet(c)];
          ans[pt] = 0;
          out[u].push back(pt++);
       void build() {
          queue <int> q;
for (q.push(0); !q.empty(); ) {
            int u = q.front();
            q.pop();
            int c = 0; c < A; ++c) {
  int v = next[u][c];
  if (!y) next[u][c] = next[link[u]][c];</pre>
                 link[v] = u ? next[link[u]][c] : 0;
out_link[v] = out[link[v]].empty() ?
                  → out link[link[v]] : link[v];
                 ed[link[v]].push back(v);
                 q.push(v);
       void dfs(int s)
          for(int x : ed[s]) dfs(x), cnt[s] += cnt[x];
          for(int e : out[s]) ans[e] = cnt[s];
       void traverse(const string &S) {
          int u = 0;
          for (char'c : S) {
            u = next[u][get(c)];
            cnt[u]++;
          dfs(0);
(18) char str[1000010], pat[505];
     int main() {
              freopen("in.txt","r",stdin);
       AC aho;
       int t,T;
       scanf("%d",&T);
for(int t=1;t<=T;t++) {
          int n;
         scanf("%d",&n);
scanf("%s",str);
for(int i=1;i<=n;i++) {
    scanf("%s",pat);
    aho.insert(pat);</pre>
          aho.build();
          aho.traverse(str);
          printf("Case'%d:\n",t);
          for(int i=0;i<n;i++) {
            printf("%d\n",aho.ans[i]);
          aho.clear();
       return 0;
     8.2 fft_match
     using ld = double:
```

using cd = complex<ld>:

**const** ld eps = 1e-6;

const ld PI = acos(-1.0);

```
void fft(vector<cd>& a, bool invert)
      int n = (int)a.size();
      for(int i = 1, j = 0; i < n; i++) {
           int bit = n >> 1;
for(; j & bit; bit >>= 1) j ^= bit;
           i ^= bit;
           if(i < j) swap(a[i], a[j]);
     for(int len = 2; len <= n; len <= 1) {
    ld ang = 2 * PI / len * (invert ? -1 : 1);</pre>
           cd wlen(cosl(ang), sinl(ang));
           for(int i = 0; i < n; i += len) {</pre>
                cd w(1);

for(int j = 0; j < len / 2; j++) {

   cd u = a[i + j], v = a[i + j + len/2]
                      a[i + j] = u + v;
                      a[i + j + len/2] = u - v;
                      w^* = wlen:
      if(invert) {
           for(cd\&x:a) {
                x /= n;
vector<i64> multiply(vector<cd> const& a, vector<cd>
     vector<cd> fa(a.begin(), a.end()), fb(b.begin(),
      while(n < (int)a.size() + b.size()) {</pre>
           n <<= 1;
      fa.resize(n);
     fb.resize(n)
     fft(fa, false);
     fft(fb, false);
for(int i = 0; i < n; i++) fa[i] *= fb[i];</pre>
     fft(fa, true);
      vector<i64> res(n);
     for(int i = 0; i < n; i++) {
   if(abs(fa[i].imag()) < eps) {
      res[i] = round(fa[i].real());
      if(abs(res[i] - fa[i].real()) > eps) {
       res[i] = -1;
   }
           } else
                res[i] = -1:
      return res:
int main()
            |t| \ll |s|
           Want to check if t occurs as a substring of s
     string s, t; cin >> s >> t;
     vector <cd> sa(s.size()), ta(t.size());
     for (int i = 0; i < s.size(); ++i) {
  ld ang = 2 * PI * (s[i] - 'a' + 1) / 27;</pre>
           sa[i] = cd(cosl(ang), sin(ang));
     for (int i = 0; i < t.size(); ++i) {
    d ang = 2 * PI * (t[i] - 'a' + 1) / 27;
    ta[t.size() - 1 - i] = cd(cosl(ang),</pre>
            → -sin(ang));
```

```
}
vector<i64> mul = multiply(sa, ta);
for (int i = 0; i < sa.size(); ++i) {
    if (mul[i] == ta.size()) {
        cout << i - ta.size() + 1 << ' ';
    }
}
cout << '\n';
}
8.3 hash
struct Hash {</pre>
```

## struct base { string s; **int** b, mod; vector<int> hash, p; **void** init(string & s, int b, int mod) $\{ // b >$ → 26, prime. s = s; b = b, mod = mod; hash\_resize(s.size()); p.resize(s.size()); hash[0] = s[0] - A' + 1; p[0] = 1;for(int i = 1; i < s.size(); ++i) { hash[i] = (ll) hash[i - 1] \* b % mod; hash[i] += s[i] - 'A' + 1; if(hash[i] >= mod) hash[i] -= mod; p[i] = (ll) p[i - 1] \* b % mod; int get(int l, int r) { int ret = hash[r]; if(l) ret -= (ll) hash[l - 1] \* p[r - l + 1] % if(ret < 0) ret += mod; return ret;</pre> void init(string &s) { h[0].init(s, 29, 1e9+7); h[1].init(s, 31, 1e9+9);pair<int, int> get(int l, int r) { **return** { h[0].get(l, r), h[1].get(l, r) }; } H;

#### 8.4 hash\_segtree

```
#define INVALID CHAR
                            - 1
namespace strhash {
  int n;
  const int MAX = 100010;
  int ara[MAX];
  const int MOD[] = {1067737007, 1069815139};
  const int BASE[] = {982451653, 984516781};
  int BP[2][MAX], CUM[2][MAX];
  void init(char *str) {
   n = strlen(str);
   for(int i=0; i< n; i++) ara[i] = str[i]-'0'+1;
 void precal() {
  BP[0][0] = BP[1][0] = 1;
  CUM[0][0] = CUM[1][0] = 1;
    for(int i=1;i<MAX;i++)</pre>
     BP[0][i] = (BP[0][i-1] * (long long) BASE[0])
         % MOD[0]:
      BP[1][i] = (BP[1][i-1] * (long long) BASE[1])

→ % MOD[1];

      CUM[0][i] = (CUM[0][i-1] + (long long) BP[0][i]
     → ) % MOD[1];
```

```
struct node {
  int sz
  int h[2];
node() {}
} tree[4*MAX];
int lazy[4*MAX];
inline void lazyUpdate(int n,int st,int ed) {
  if(lazy[n]!=INVALID CHAR){
    tree[n].h[0] = (lazy[n] * (long long)
    - CUM[0][ed-st]) % MOD[0];
tree[n].h[1] = (lazy[n] * (long long)

→ CUM[1][ed-st]) % MOD[1];

    if(st!=ed){
      lazy[2*n] = lazy[n];
      lazy[2*n+1] = lazy[n];
    lazy[n] = INVALID CHAR;
inline node Merge(node a, node b) {
  node ret;
  ret.h[0] = ( (a.h[0] * (long long) BP[0][b.sz] )
  + b.h[0] ) % MOD[0];
ret.h[1] = ( ( a.h[1] * (long long) BP[1][b.sz] )
  \rightarrow + b.h[1] ) % MOD[1];
  ret.sz = a.sz + b.sz:
  return ret;
inline void build(int n,int st,int ed) {
  lazy[n] = INVALID CHAR;
  if(st==ed) ·
    tree[n].h[0] = tree[n].h[1] = ara[st];
    tree[n].sz = 1;
    return;
  int mid = (st+ed)>>1;
  build(n+n,st,mid);
  build(n+n+1,mid+1,ed);
  tree[n] = Merge(tree[n+n], tree[n+n+1]);
inline void update(int n,int st,int ed,int i,int
lazyUpdate(n,st,ed);
  if(st>j or ed<i) return;</pre>
  if(st>=i and ed<=j) {
    lazy[n] = v;
    lazyUpdate(n,st,ed);
    return;
  int mid = (st+ed)>>1;
  update(n+n,st,mid,i,j,v);
  update(n+n+1,mid+1,ed,i,j,v);
  tree[n] = Merge(tree[n+n], tree[n+n+1]);
inline node query(int n,int st,int ed,int i,int j){
  lazyUpdate(n,st,ed);
  if(st>=i and ed<=j) return tree[n];</pre>
  int mid = (st+ed)/2;
  if(mid<i) return query(n+n+1,mid+1,ed,i,j);</pre>
  else if(mid>=j) return query(n+n,st,mid,i,j);
  else return Merge(query(n+n,st,mid,i,j),query(n+n+
   → 1,mid+1,ed,i,j));
```

```
// returns the longest proper prefix array of pattern p
// where lps[i]=longest proper prefix which is also
→ suffix of p[0...i]
vector<int> build lps(string p) {
 int sz = p.size();
  vector<int> lps;
  lps.assign(sz + 1, 0);
  int j = \bar{0};
  lps[0] = 0;
 for(int i = 1; i < sz; i++) {
  while(j >= 0 && p[i] != p[j]) {
      if(j >= 1) j = lps[j - 1];
      else j = -1;
    lps[i] = j;
  return lps;
vector<int>ans;
// returns matches in vector ans in 0-indexed
void kmp(vector<int> lps, string s, string p) {
 int psz = p.size(), sz = s.size();
  int i = 0;
  for(int i = 0; i < sz; i++)
    while(j >= 0 \&\& p[j] != s[i])
      if(j >= 1) j = lps[j - 1];
      else i = -1;
    i++:
    if(j == psz) {
      j = lps[j - 1];
      // pattern found in string s at position i-psz+1
      ans push back(i - psz + 1);
    // after each loop we have j=longest common suffix
     \rightarrow of s[0..i] which is also prefix of p
8.6 manachar
vector<int> d1(n); // maximum odd length palindrome
// here d1[i]=the palindrome has

→ d1[i]-1 right characters from i

                     // e.g. for aba, d1[1]=2;
for (int i = 0, l = 0, r = -1; i < n; i++) {
 int k = (i > r) ? 1 : min(d1[l + r - i], r - i)
 while (0) <= i - k \&\& i + k < n \&\& s[i - k] == s[i + k]
  k++;
 d1[i] = k--;
 if (i + k > r) {
    l = i - k;
    r = i + k;
vector<int> d2(n); // maximum even length palindrome
// here d2[i]=the palindrome has

    d2[i]-1 right characters from i

                    // e.g. for abba, d2[2]=2;
for (int i = 0, l = 0, r = -1; i < n; i++) {
 int k = (i > r) ? 0 : min(d2[l + r - i + 1], r - i +
  while (0 \le i - k - 1) \& i + k \le n \& s[i - k - 1]
     == s[i + k]) {
    k++;
```

8.5 kmp

d2[i] = k--;

```
if (i + k > r) {
    l = i - k - 1;
r = i + k;
8.7 palindromic_tree
const int A = 26;
const int N = 300010;
char s[N]; long long ans;
int last, ptr, nxt[N][A], link[N], len[N], occ[N];
void feed (int at) -
  while (s[at - len[last] - 1] != s[at]) last =
      link[last];
  int ch = s[at] - 'a', temp = link[last];
  while (s[at - len[temp] - 1] != s[at]) temp =
  if (!nxt[last][ch]) {
   nxt[last][ch] = ++ptr, len[ptr] = len[last] + 2;
                                                                  int SA[N], iSA[N]; // SA is the suffix array, iSA[i]
                                                                 int cnt[N], nxt[N]; // Internal stuff
bool bh[N], b2h[N]; // Internal stuff
int lcn[N]: // Starce/larger
     link[ptr] = len[ptr] = 1 ? 2 : nxt[temp][ch];
   last = nxt[last][ch], ++occ[last];
int main() {
  len[1] = -1, len[2] = 0, link[1] = link[2] = 1, last
   \rightarrow = ptr = 2;
  scanf("%s", s + 1);
  for (int i = 1, n = strlen(s + 1); i <= n; ++i)</pre>
   \rightarrow feed(i);
  for (int i = ptr; i > 2; --i) ans = max(ans, len[i]
   * 1LL * occ[i]), occ[link[i]] += occ[i];
  printf("%lld\n", ans);
  return 0;
8.8 persistant trie
const int MAX = 200010;
const int B = 19;
int root[MAX], ptr = 0;
struct node {
```

```
int ara[2], sum;
node() {}
\ tree \( MAX \( * (B+1) \);
void insert(int prevnode, int &curRoot, int val) {
  curRoot = ++ptr;
  int curnode = curRoot;
  for(int i = B; i >= 0; i--) {
  bool bit = val & (1 << i);
  tree[curnode] = tree[prevnode];</pre>
    tree[curnode].ara[bit] = ++ptr;
    tree[curnode].sum += 1
    prevnode = tree[prevnode].ara[bit];
    curnode = tree[curnode].ara[bit];
  tree[curnode] = tree[prevnode];
  tree[curnode].sum += 1;
int find xor max(int prevnode, int curnode, int x) {
  int ans = 0;
  for(int i = B; i >= 0; i--) {
    bool bit = x \& (1 << \bar{i});
    if(tree[tree[curnode].ara[bit ^ 1]].sum >
     - tree[tree[prevnode].ara[bit ^ 1]].sum) {
curnode = tree[curnode].ara[bit ^ 1];
      prevnode = tree[prevnode].ara[bit ^ 1];
       ans = ans | (1 << i);
    else {
       curnode = tree[curnode].ara[bit]
       prevnode = tree[prevnode].ara[bit];
```

```
return ans;
void solve() {
  int n, q, L, R, K;
  cin >> n;
  for(int i=1;i<=n;i++) cin >> ara[i];
  for(int i=1;i<=q;i++) {</pre>
    cin >> L >> R >> K;
    cout << find xor max(root[L-1],root[R],K) << endl;</pre>
```

char s[N]; // Suffix array will be built for this

stores the rank of the i'th suffix

```
8.9 suffix_array
```

// Everything is 0-indexed

```
int lcp[N]; // Stores lcp of SA[i] and SA[i + 1];
\rightarrow lcp[n - 1] = 0
int lcpSparse[LOGN][N]; // lcpSparse[i][j] =
 \rightarrow min(lcp[j], ..., lcp[j - 1 + (1 << i)])
|void buildSA(int n) {
  for (int i = 0; i < n; i++) SA[i] = i;
  sort(SA, SA + n, [](int i, int j) { return s[i] <
   \hookrightarrow s[j]; \});
  for (int i = 0; i < n; i++)  { bh[i] = i == 0 || s[SA[i]] != s[SA[i - 1]];
    b2h[i] = 0;
  for (int h = 1; h < n; h <<= 1) {
    int tot = 0;
    for (int i = 0, j; i < n; i = j) {
       i = i + 1;
       while (j < n && !bh[j]) j++;
       nxt[i] = j; tot++;
    } if (tot == n) break;
    for (int i = 0; i < n; i = nxt[i])
       for (int j = i; j < nxt[i]; j++) iSA[SA[j]] = i;
       cnt[i] = 0;
    b2h[iSA[n - h]] = 1;
    for (int i = 0; i < n; i = nxt[i]) {
  for (int j = i; j < nxt[i]; j++) {</pre>
         int s = SA[j] - h;
         if (s < 0) continue;
         int head = iSA[s];
         iSA[s] = head + cnt[head]++;
         b2h[iSA[s]] = 1;
       for (int j = i; j < nxt[i]; j++) {</pre>
         int s = SA[j] - h;
         if (s < 0 | | !b2h[iSA[s]]) continue;</pre>
         for (int k = iSA[s] + 1; !bh[k] \&\& b2h[k];
             k++) b2h[k] = 0;
    for (int i = 0; i < n; i++) {
       SA[iSA[i]] = i
       bh[i] [= b2h[i];
  for (int i = 0; i < n; i++) iSA[SA[i]] = i;
|void buildLCP(int n) {
| for (int i = 0, k = 0; i < n; i++) {
```

```
if (iSA[i] == n - 1) {
      k = 0;
      lcp[n' - 1] = 0;
     continue:
    int j = SA[iSA[i] + 1];
   while (i + k < n \&\& j + k < n \&\& s[i + k] == s[j + k]
   if (k) k--;
void buildLCPSparse(int n) {
 for (int i = 0; i < n; i++) lcpSparse[0][i] = lcp[i];
 for (int i = 1; i < LOGN; i++) {
    for (int j = 0; j < n; j++) {
      lcpSparse[i][j] = min(lcpSparse[i - 1][j],
          lcpSparse[i - 1][min(n - 1, j + (1 << (i -
      □ 1)))]);
pair<int, int> minLCPRange(int n, int from, int

→ minLCP) {
 int r = from;
 for (int i = LOGN - 1; i >= 0; i--) {
   int jump = 1 << i;
   if (r + jump < n and lcpSparse[i][r] >= minLCP) r
 int l = from;
 for (int i = LOGN - 1; i >= 0; i--) {
    int jump = 1 << i;
    if (\(\bar{l}\) - jump >= 0 and lcpSparse[i][l - jump] >=

→ minLCP) l -= jump;

 return make pair(l, r);
```

#### 8.10 suffix\_automata

```
namespace sa{
 const int MAXN = 100005 << 1;
 const int MAXC = 26;
  char str[MAXN];
 int n, sz, last;
 int len[MAXN], link[MAXN], ed[MAXN][MAXC], cnt[MAXN];
bool terminal[MAXN];
 vector <int> G[MAXN];
 void init()

\begin{array}{lll}
SET(ed[0]); \\
len[0] = 0, \\
link[0] = -1, \\
sz = 1, \\
last = 0,
\end{array}

    terminal[0] = false;

 inline int scale(char c) { return c-'a'; }
  void extend(char c) {
    int cur = sz++;
    terminal[cur] = false;
    cnt[cur] = 1;
    SET(ed[cur]);
len[cur] = len[last] + 1;
    int p = last;
    while (p != -1 \&\& ed[p][c]==-1) {
      ed[p][c] = cur;
      p = link[p];
    if (p == -1) link[cur] = 0;
    else {
```

```
int q = ed[p][c];
    if (len[p] + 1 == len[q]) link[cur] = q;
    else {
      int clone = sz++
      len[clone] = len[p] + 1;
      memcpy(ed[clone],ed[q],sizeof(ed[q]));
      link[clone] = link[q];
      while (p != -1 \&\& ed[p][c] == q) {
        ed[p][c] = clone;
        p = link[p];
      link[q] = link[cur] = clone;
      cnt[clone] = 0;
      terminal[clone] = false;
  last = cur;
void dfs(int s) {
 for(auto x : G[s]) dfs(x), cnt[s] += cnt[x];
void build() {
 init();
 int n = strlen(str);
 for(int i=0;i<n;i++) extend(scale(str[i]));</pre>
 for(int i=1;i<sz;i++) G[link[i]].pb(i);</pre>
 dfs(0);
 for(int i=0;i<sz;i++) G[i].clear();</pre>
 for(int i=last;i!=-1;i=link[i]) terminal[i] = true;
```

```
8.11 trie
#define N
                  200000
#define S
int root.now:
int nxt[N][S], cnt[N];
void init(){
  root = now = 1;
  CLR(nxt), CLR(cnt);
inline int scale(char ch) { return (ch - 'a'); }
inline void Insert(char s[],int sz){
 int cur = root, to;
 for(int i=0; i < sz; i++){
  to = scale(s[i]);
  if( !nxt[cur][to] ) nxt[cur][to] = ++now;</pre>
    cur = nxt[cur][to];
  cnt[cur]++;
inline bool Find(char s[],int sz){
  int cur = root, to;
for(int i=0; i<sz; i++){</pre>
    to = scale(s[i])
    if( !nxt[cur][to] ) return false;
    cur = nxt[cur][to];
  return (cnt[cur]!=0);
inline void Delete(char s[],int sz){
  int cur = root, to;
  for(int i=0 ; i<sz ; i++){</pre>
    to = scale(s[i])
    cur = nxt[cur][to];
  cnt[cur]--;
```

8.12 z\_algo

const int N = 100010;

```
char s[N];
int t, n, z[N];
int main() {
    scanf("%s", s);
    n = strlen(s), z[0] = n;
    int L = 0, R = 0;
    for (int i = 1; i < n; ++i) {
        if (i > R) {
            L = R = i;
            while (R < n && s[R - L] == s[R]) ++R;
            z[i] = R - L; --R;
        } else {
        int k = i - L;
        if (z[k] < R - i + 1) z[i] = z[k];
        else {
            L = i;
            while (R < n && s[R - L] == s[R]) ++R;
            z[i] = R - L; --R;
        }
    }
}</pre>
```