

DU_Kronos

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```
if(cur >= len ) cur = len-1;
   while( cur < len-1 && getY(cur+1,x) >= getY(cur,x)
    → ) cur++; // <= to minimize, >= to maximize
    return getY(cur,x);
  inline ll TS( ll x ) {
   int low = 0, high = len-1 , mid ;
   while ( high - low > 1 ) {
      mid = low + high >> 1;
      if(qetY(mid,x) < qetY(mid+1,x)) low = mid + 1;
       → // > to minimize , < to maximize
     else high = mid;
   return max(getY(low,x),getY(high,x)); // adjust

→ min/max

};'
CHT cht;
cht.init();
```

1.2 DC Optimization

```
void compute(int L, int R, int optL, int optR){
 if(L > R) return;
 int M = L + R \gg 1
 pair<ll, int> best(1LL << 60, -1);</pre>
 for(int k = optL; k <= min(M, optR); k++){</pre>
   best = min(best, \{dp[prv][k] + C[k + 1][M], k\});
 dp[now][M] = best.ff;
 compute(L, M - 1, optL, best.ss);
 compute(M + 1, R, best.ss, optR);
```

1.3 SOS DP

```
S(mask.i) denotes those submasks of mask which
from mask only in the least significant i+1 bits (
- 0, 1, 2, ...., i --> 0 based indexing)
Example : S(1011010,3) =
\rightarrow {1011010, 1010010, 1011000, 1010000}
if(bit i is on)
    S(mask, i)' = S(mask, i - 1) U S(mask ^ (1 << i),
     \rightarrow i - 1)
    S(mask, i) = S(mask, i - 1)
Let Sum(mask,i) denote the sum of the all A[x]
→ where x element of S(mask,i)
So, Sum(mask,N-1) will contain SOS DP result for a

→ particular mask

Recurrence : if(i'th bit is on)
    Sum(mask, i) = Sum(mask, i-1) + Sum(mask ^
      \hookrightarrow (1<<i), i-1)
     Sum(mask, i) = Sum(mask, i-1)
for(int mask = 0; mask < (1<<N); ++mask){
     dp[mask][-1] = A[mask];
     for(int i = 0; i < N; ++i){
         if(mask & (1 << i)) dp[mask][i] =
              dp[mask][i-1] + dp[mask ^ (1<<i)][i-1];
         else dp[mask][i] = dp[mask][i-1];
     F[mask] = dp[mask][N-1];
for(int i = 0; i < (1 << N); ++i) F[i] = A[i];
for(int i = 0;i < N; ++i)
    for(int mask = 0; mask < (1<<N); ++mask){</pre>
         if(mask \& (1<<i))
              F[mask] += F[mask^(1<<i)];
```

```
* c = x\&-x, r = x+c; (((r^x) >> 2)/c) \mid r is the next
    number after x with the same number of bits set.
```

```
1.4 dynamic cht
//add lines with -m and -b and return -ans to
//make this code work for minimums.(not -x)
const ll_inf = -(1LL << 62);</pre>
|struct line {
  ll m, b;
  mutable function<const line*() > succ;
  bool operator < (const line& rhs) const {</pre>
     if (rhs.b != inf) return m < rhs.m;</pre>
     const line* s = succ();
    if (!s) return 0;
ll x = rhs.m;
     return b - s->b < (s->m - m) * x;
struct CHT : public multiset<line> {
  bool bad(iterator y) {
     auto z = next(y);
     if (y == begin()) {
       if (z == end()) return 0;
       return y \rightarrow m == z \rightarrow m \&\& y \rightarrow b <= z \rightarrow b;
     auto x = prev(y);
     if (z == end()) return y -> m == x -> m &  y -> b
     return 1.0 * (x -> b - y -> b) * (z -> m - y -> m)
     \Rightarrow >= 1.0 * (y -> b - z -> b) * (y -> m - x -> m);
  void add(ll m, ll b) {
     auto y = insert({ m, b });
     y->succ = [ = ] { return next(y) == end() ? 0 :
      \hookrightarrow &*next(y); };
     if (bad(y)) {
       erase(y);
       return;
     while (next(y) != end() \&\& bad(next(y)))
         erase(next(y));
     while (y != begin() && bad(prev(y)))

    erase(prev(y));

  11 query(ll x) {
     auto l = *lower bound((line) {
       x, inf
     return l.m * x + l.b;
ľĆĤT* cht;
ĺl a[N], b[N];
ios_base::sync_with_stdio(0);
  cin.tie(0);
  int n;
  cin >> n:
   for(int i = 0; i < n; i++) cin >> a[i];
  for(int i = 0; i < n; i++) cin >> b[i];
   cht = new CHT(
  cht \rightarrow add(-b[0], 0);
  ll ans = 0;
  for(int i = 1; i < n; i++)
    ans = -cht -> query(a[i]);
     cht -> add(-b[i], -ans);
```

cout << ans << nl;

return 0;

```
2 Data Structure
2.1 2D_segtree
struct Point {
    int x, y, mx;
Point() {}
    Point(int x, int y, int mx) : x(x), y(y), mx(mx) {}
    bool operator < (const Point& other) const {</pre>
         return mx < other.mx;</pre>
struct Segtree2d {
    // I didn't calculate the exact size needed in
        terms of 2D_container size.
       If anyone, please edit the answer.
    // It's just a safe size to store nodes for MAX *
        MAX 2D grids which won't cause stack overflow
    Point T[500000]; // TODO: calculate the accurate

→ space needed
    int n, m;
    // initialize and construct segment tree
    void init(int n, int m) {
         this -> n = n;
         this -> m = m;
         build(1, 1, 1, n, m);
    // build a 2D segment tree from data [ (a1, b1),
     // Tìme: O(n logn)
    Point build(int node, int a1, int b1, int a2, int
     → b2) {
         // out of range
        if (a1 > a2 \text{ or } b1 > b2)
             return def();
         // if it is only a single index, assign value
             to node
         if (a1 == a2 \text{ and } b1 == b2)
             return T[node] = Point(a1, b1, P[a1][b1]);
         // split the tree into four segments
        T[node] = def();
        T[node] = maxNode(T[node], build(4 * node - 2,
        al, bl, (a1 + a2) / 2, (b1 + b2) / 2));
T[node] = maxNode(T[node], build(4 * node - 1,
             (a1 + a2) / 2 + 1, b1, a2, (b1 + b2) / 2
        T[node] = maxNode(T[node], build(4 * node + 0,
             a1, (b1 + b2) / 2 + 1, (a1 + a2) / 2, b2)
        T[node] = maxNode(T[node], build(4 * node + 1,
             (a1 + a2) / 2 + 1, (b1 + b2) / 2 + 1, a2,

    b2) );

         return Tinodel;
    // helper function for query(int, int, int, int);
    Point query(int node, int al, int bl, int a2, int

→ b2, int x1, int y1, int x2, int y2) {
         // if we out of range, return dummy
        if (x1 > a2 or y1 > b2 or x2 < a1 or y2 < b1
         \rightarrow or a1 > a2 or b1 > b2)
             return def();
         // if it is within range, return the node
        if (x1 \le a1 \text{ and } y1 \le b1 \text{ and } a2 \le x2 \text{ and } b2
         return T[node];
        // split into four segments
```

```
Point mx = def();
        mx = maxNode(mx, query(4 * node - 2, a1, b1,
             (a1 + a2) / 2, (b1 + b2) / 2, x1, y1, x2,
        mx = maxNode(mx, query(4 * node - 1, (a1 + a2))
            /2 + 1, b1, a2, (b1 + b2) /2, x1, y1,
            x2, y2));
        mx = maxNode(mx, query(4 * node + 0, a1, (b1 +
            b2) / 2 + 1, (a1 + a2) / 2, b2, x1, y1,
         \stackrel{\sim}{\sim} x2, y2));
        mx = maxNode(mx, query(4 * node + 1, (a1 + a2))
            /2 + 1, (b1 + b2) /2 + 1, a2, b2, x1,

    y1, x2, y2));

         // return the maximum value
        return mx;
    // query from range [ (x1, y1), (x2, y2) ]
    // Time: 0(logn)
    Point query(int x1, int y1, int x2, int y2)
        return query(1, 1, 1, n, m, x1, y1, x2, y2);
    // helper function for update(int, int, int);
    Point update(int node, int al, int bl, int a2, int
        b2, int x, int y, int value) {
if (a1 > a2 or b1 > b2)
            return def();
        if (x > a2 \text{ or } y > b2 \text{ or } x < a1 \text{ or } y < b1)
            return T[node];
        if (x == a1 \text{ and } y == b1 \text{ and } x == a2 \text{ and } y ==
            return T[node] = Point(x, y, value);
        Point mx = def();
        mx = maxNode(mx, update(4 * node - 2, a1, b1,
             (a1 + a2) / 2, (b1 + b2) / 2, x, y, value)
         □);
        mx = maxNode(mx, update(4 * node - 1, (a1 +
            a2) / 2 + 1, b1, a2, (b1 + b2) / 2, x, y,
        mx = maxNode(mx, update(4 * node + 0, a1, (b1)))
            + b2) / 2 + 1, (a1 + a2) / 2, b2, x, y,

    value));

        mx = maxNode(mx, update(4 * node + 1, (a1 +
            a2) / 2 + 1, (b1 + b2) / 2 + 1, a2, b2, x,
         return T[node] = mx;
    // update the value of (x, y) index to 'value'
    // Time: 0(logn)
    Point update(int x, int y, int value) {
        return update(1, 1, 1, n, m, x, y, value);
    // utility functions; these functions are virtual
        because they will be overridden in child class
    virtual Point maxNode(Point a, Point b) {
        return max(a, b);
    // dummy node
   virtual Point def() {
    return Point(0, 0, -INF);
/* 2D Segment Tree for range minimum query; a override

→ of Seatree2d class */
struct Segtree2dMin : Segtree2d {
    // overload maxNode() function to return minimum
    Point maxNode(Point a, Point b) {
        return min(a, b);
```

```
Point def()_{
        return Point(0, 0, INF);
2.2 Segtree beats desc
Description: For update Ai Ai mod x and similar, keep
 max in node and lazily update whenever min = max. For
Ai min(Ai, x) and similar, keep range max, second max
    in node and
lazily update whenever x > second max.
Time: O(\log^2 N), (\log N)
2.3 gp_hash_table
using namespace gnu pbds;
const int RANDOM = chrono::high resolution clock::now(;
→ ).time since epoch().count();
using namespace gnu pbds;
struct chash -
  const int RANDOM = (long
      long)(make unique<char>().get()) ^
      chrono::high resolution clock::now().time since
      epoch().count();
  static unsigned long long hash f(unsigned long long
    x += 0x9e3779b97f4a7c15
    x = (x ^ (x >> 30)) * 0xbf58476d1ce4e5b9;
x = (x ^ (x >> 27)) * 0x94d049bb133111eb;
return x ^ (x >> 31);
  static unsigned hash combine(unsigned a, unsigned b)
      { return a * 31 + b; }
  ll operator()(ll x) const { return hash f(x)^RANDOM;
gp hash table<key, long long, chash> table;
2.4 iterative_segtree
const int N = 500010;
int n, a[N], tree[N << 1];</pre>
|void init() {
    for (int i = 0; i < n; ++i) tree[n + i] = a[i];
    for (int i = n' - 1; i >= 0; --i)
        tree[i] = min(tree[i << 1], tree[i << 1 | 1]);</pre>
void update(int p, int v) {
    for (tree[p += n] = v; p > 1; p >>= 1)
        tree[p] >> 1] = min(tree[p], tree[p ^ 1]);
|int query(int l, int r) {
    int ret = INT MAX;
    for (l += n, \bar{r} += n; l < r; l >>= 1, r >>= 1) {
          f (l & 1) ret = min(ret, tree[l++]);
        if (r \& 1) ret = min(ret, tree[--r]);
    return ret;
2.5 mos_algo
struct data{
  int l, r, id, bn;
data() {}
```

data(int l, int r, int id){

```
l = l, r = r, id = id;
    bn = l' / block'sz;
 bool operator < (const data& other) const{</pre>
    if (bn != other.bn) return (bn < other.bn);</pre>
    return ((bn \& 1) ? (r < other.r) : (r > other.r));
int curL = 0, curR = -1;
for(int i = 0; i < 0.sz; i++){
 while(curL > Q[i].L){
    curL--; add(curL);
 while(curR < 0[i].R)
    curR++; add(curR);
 while(curL < Q[i].L){
    remove(curL); curL++;
 while(curR > 0[i].R){
    remove(curR); curR--;
2.6 ordered set
using namespace std;
using namespace gnu pbds;
typedef tree<
   int
    null'type
    less < int > , // "less equal<int>," for multiset
    rb tree tag,
    tree order statistics node update > ordered set;
ordered set OS;
2.7 persistant_segtree
/* Persistent Segment Tree using static Array
 Point Update, Range Sum
 Initialize ncnt to 0 in every test case */
const int MAX = 100010;
```

```
int ncnt = 0;
struct node {
 int sum:
  int left,right;
 node() {}
  node(int val) {
    sum = val;
    left = riaht = -1:
} free[ ? ];
int ara[MAX];
int version[MAX];
void build(int n,int st,int ed) {
  if (st==ed) {
    tree[n] = node(ara[st]);
    return;
  int mid = (st+ed) / 2;
  tree[n].left = ++ncnt;
  tree[n].right = ++ncnt;
  build(tree[n].left, st, mid);
  build(tree[n].right, mid+1, ed);
  tree[n].sum = tree[tree[n].left].sum +

    tree[tree[n].right].sum;

void update(int prev,int cur,int st,int ed,int id, int
```

```
(id > ed or id < st) return;
  if (st == ed) {
    tree[cur] = node(val);
    return;
  int mid = (st+ed) / 2;
  if (id <= mid)
    tree[curl.right = tree[prevl.right:
    tree[cur].left = ++ncnt;
    update(tree[prev].left,tree[cur].left, st, mid,
    → id. val):
  else {
   tree[cur].left = tree[prev].left;
tree[cur].right = ++ncnt;
    update(tree[prev].right, tree[cur].right, mid+1,

    ed. id. val):

  tree[cur].sum = tree[tree[cur].left].sum +

    tree[tree[cur].right].sum;

int query(int n,int st,int ed,int i,int j){
  if(st>=i && ed<=j) return tree[n].sum;</pre>
  int mid = (st+ed)/2;
  if(mid<i) return query(tree[n].right,mid+1,ed,i,j);</pre>
  else if(mid>=j) return

¬ query(tree[n].left,st,mid,i,j);

  else return query(tree[n].left,st,mid,i,j) +
     query(tree[n].right,mid+1,ed,i,j);
int main() {
  int n,q,l,r,k;
  sii(n,q);
  version[0] = ++ncnt;
  build(version[0],1,n);
  version[1] = ++ncnt;
  update(version[0], version[1], 1, n, id, val);
  query(version[0],1,n,id,id);
  query(version[1],1,n,id,id);
  return 0;
```

2.8 segment_tree

```
int ara[MAX];
struct node {
int sum;
} tree[4 * MAX];
int lazy[4 * MAX];
node Merge(node a, node b) {
  node ret;
  ret.sum = a.sum + b.sum;
  return ret;
void lazyUpdate(int n, int st, int ed) {
  if(lazy[n] != 0){
    tree[n].sum += ((ed - st + 1) * lazy[n]);
    if(st != ed){
      lazy[2 * n] += lazy[n];
lazy[2 * n + 1] += lazy[n];
    lazy[n] = 0;
void build(int n, int st, int ed) {
  lazy[n] = 0;
  if(st == ed){
    tree[n].sum = ara[st];
    return;
  int mid = (st + ed) / 2;
```

```
build(2 * n, st, mid);
build(2 * n + 1, mid + 1, ed);
  tree[n] = Merge(tree[2 * n], tree[2 * n + 1]);
void update(int n, int st, int ed, int i, int j, int

    ∨) {
 lazyUpdate(n, st, ed);
  if(st > j or ed < i) return;</pre>
  if(st >= i and ed <= i){
    lazy[n] += v;
    lazyUpdate(n, st, ed);
    return;
  int mid = (st + ed) / 2;
update(2 * n, st, mid, i, j, v);
  update(2 * n + 1, mid+1, ed, i, j, v);
  tree[n] = Merge(tree[2 * n], tree[2 * n + 1]);
|node query(int n, int st, int ed, int i, int j) {
  lazyUpdate(n, st, ed);
  if(st >= i and ed <= j) return tree[n];</pre>
  int mid = (st + ed) / 2;
  if (mid < i) return query (2 * n + 1, mid + 1, ed, i,

→ j);

  else if(mid >= j) return query(2 * n, st, mid, i, j);
  else return Merge(query(2 * n, st, mid, i, j),
   \rightarrow query(2 * n + 1, mid + 1, ed, i, j));
```

2.9 sparse_table

```
int st[K + 1][MAXN];
|void build() {
  std::copy(array.begin(), array.end(), st[0]);
  for (int i = 1; i <= K; i++)
    for (int j = 0; j + (1 << i) <= N; j++)
      st[i][j] = f(st[i - 1][j], st[i - 1][j + (1 <<
       \rightarrow (i - 1))]);
```

3 Geometry 3.1 2D Primitive 3.1.1 Angle

of rotations around the origin). Useful for rotational sweeping. Sometimes also represents points or vectors.

```
/* Usage:
* vector<Angle> v = \{w[0], w[0].t360() ...\}; //
    sorted
* int j = 0; rep(i,0,n) { while (v[j] < v[i].t180())
 \hookrightarrow ++j; }
   // sweeps j such that (j-i) represents the number
     of positively oriented triangles with vertices at
    0 and i
|struct Angle {
  int x, y;
  int t;
  Angle(int x, int y, int t=0) : x(x), y(y), t(t) {}
  Angle operator-(Angle b) const { return {x-b.x,
  y-b.y, t}; }
  int half() const {
    assert(x || y);
    return y < 0 | | (y == 0 \&\& x < 0);
  Angle t90() const { return \{-y, x, t + (half() \&\& x\}\}
 -->= 0)); }
Angle t180() const { return {-x, -y, t + half()}; }
  Angle t360() const { return \{x, y, t + 1\}; \}
```

```
bool operator<(Angle a, Angle b) {</pre>
      // add a.dist2() and b.dist2() to also compare
                  distances
      return make tuple(a.t, a.half(), a.y * (ll)b.x) <</pre>

→ make tuple(b.t, b.half(), a.x * (ll)b.y);

// Given two points, this calculates the smallest
           angle between
// them, i.e., the angle that covers the defined line

→ seament.

pair<Angle, Angle> segmentAngles(Angle a, Angle b) {
      if (b < a) swap(a, b);
      return (b < a.t180() ? make pair(a, b) :

→ make pair(b, a.t360()));

Angle operator+(Angle a, Angle b) { // point a +
             vector b
      Angle r(a.x + b.x, a.y + b.y, a.t);
      if (a.t180() < r) r.t--:
      return r.t180() < a ? r.t360() : r;</pre>
Angle angleDiff(Angle a, Angle b) { // angle b - angle
      int tu = b.t - a.t; a.t = b.t;
      return \{a.x*b.x + \dot{a}.y*b.y, a.\dot{x}*b.y - a.y*b.x, tu -
         \rightarrow (b < a)};
```

3.1.2 Line Distance

Returns the signed distance between point p and the line containing points a and b. Positive value on left side and negative on right as seen from a towards b. a==b gives nan. P is supposed to be Point<T> or Point3D<T> where T is e.g. double or long long. It uses products in intermediate steps so watch out for overflow if using int or long long. Using Point3D will always give a non-/s negative distance. For Point3D, call .dist on the result of the cross product.



```
#include "Point.h"
                                                             template<class P>
A class for ordering angles (as represented by int points and a number double lineDist(const P& a, const P& b, const P& p) {
                                                               return (double)(b-a).cross(p-a)/(b-a).dist();
```

3.1.3 Line Intersection

If a unique intersection point of the lines going through s1,e1 and s2,e2 exists {1, point} is returned. If no intersection point exists {0, (0,0)} is returned and if infinitely many exists {-1, (0,0)} is returned. The wrong position will be returned if P is Point<ll> and the intersection point does not have integer coordinates. Prod- < ucts of three coordinates are used in intermediate steps so watch out for overflow if using int or ll.



```
auto res = lineInter(s1,e1,s2,e2);
   if (res.first == 1)
  cout << "intersection point at " << res.second</pre>
     << endl;
#pragma once
#include "Point.h"
template<class P>
pair<int, P> lineInter(P s1, P e1, P s2, P e2) {
 auto d = (e1 - s1).cross(e2 - s2);
 if (d == 0) // if parallel
    return \{-(s1.cross(e1, s2) == 0), P(0, 0)\};
```

```
auto p = s2.cross(e1, e2), q = s2.cross(e2, s1);
  return {1, (s1 * p + e1 * q) / d};
3.1.4 Linear Transformation
Apply the linear transformation (translation, rotation po
and scaling) which takes line p0-p1 to line q0-q1 to point q0
#include "Point.h"
```

```
typedef Point<double> P;
P linearTransformation(const P& p0, const P& p1, const
   P\& q0, const P\& q1, const P\& r) {
 P dp = p1-p0, dq = q1-q0, num(dp.cross(dq),
    dp.dot(dq));
 return q0 + P((r-p0).cross(num),
```

3.1.5 On Segment

```
/* Description: Returns true iff p lies on the line

→ segment from s to e.

* Use \texttt{(seqDist(s,e,p)<=epsilon)} instead when
    using Point<double>.
#include "Point.h"
template < class P > bool on Segment (P s, P e, P p) {
 return p.cross(s, e) == 0 \& \& (s - p).dot(e - p) <= 0;
```

3.1.6 Point Sort

```
// sort the points in counterclockwise order that
\rightarrow starts from the half line x0, y=0.
using namespace std;
typedef long long ll;
typedef pair <ll, ll> point;
#define x first
#define y second
int main() {
  int n; cin >> n;
  vector <point> p(n);
  for (auto &it : p) scanf("%lld %lld", &it.x, &it.y);
  sort(p.begin(), p.end(), [] (point a, point b) {
    return atan2l(a.y, a.x) < atan2l(b.y, b.x);</pre>
  for (auto it : p) printf("%lld %lld\n", it.x, it.y);
  return 0;
```

3.1.7 Point

```
// Class to handle points in the plane. T can be e.g.

→ double or long long. (Avoid int.)
template <class T> int sgn(T x) \{ return (x > 0) - (x < 0) \}
\rightarrow < 0);
templaté<class T>
struct Point {
 typedef Point P;
Tx, y
 explicit Point(T x=0, T y=0) : x(x), y(y) {}
 bool operator<(P p) const { return tie(x,y) <</pre>
\begin{tabular}{lll} &\rightarrow & \text{tie}(p.x,p.y)\,; & \\ & \text{bool} & \text{operator} {==}(P\ p) & \text{const} \ \{\ \text{return} \end{tabular}
    tie(x,y)==tie(p.x,p.y); 
 P operator+(P p) const { return P(x+p.x, y+p.y); ]
 P operator-(P p) const { return P(x-p.x, y-p.y); } P operator*(T d) const { return P(x*d, y*d); }
 P operator/(T d) const { return P(x/d, y/d); }
 T dot(P p) const { return x*p.x + y*p.y; }
 T cross(P p) const { return x*p.y - y*p.x; }
```

```
T cross(P a, P b) const { return
    (a-*this).cross(b-*this); }
T dist2() const { return x*x + y*y;
double dist() const { return sqrt((double)dist2()); }
// angle to x-axis in interval [-pi, pi]
double angle() const { return atan2(y, x); }
P unit() const { return *this/dist(); } // makes
  dist()=1
P perp() const { return P(-y, x); } // rotates +90
→ degrees
P normal() const { return perp().unit(); }
// returns point rotated 'a' radians ccw around the
   origin
P rotate(double a) const {
 return P(x*cos(a)-y*sin(a),x*sin(a)+y*cos(a)); }
friend ostream& operator<<(ostream& os, P p) {</pre>
return os << "("' << p.x << "," << p.y << '')"; }
```

3.1.8 Segment Distance

Returns the shortest distance between point p and the line segment from point s to e.

```
/* Usage:
   Point<double> a, b(2,2), p(1,1);
bool onSegment = segDist(a,b,p) < 1e-10;
  Status: tested
#pragma once
#include "Point.h"
typedef Point<double> P;
double segDist(P& s, P& e, P& p) {
  if (s==e) return (p-s).dist();
  auto d = (e-s).dist2(), t =
  \rightarrow min(d,max(.0,(p-s).dot(e-s)));
  return ((p-s)*d-(e-s)*t).dist()/d;
```

3.1.9 Segment Intersection

If a unique intersection point between the line segments going from s1 to e1 and from s2 to e2 exists then it is returned. If no intersection point exists an empty vector is returned. If infinitely many exist a vector with 2 elements is returned, containing the endpoints of the common line segment. The wrong position will be returned if P is Point<ll> and the intersection point does \$1 not have integer coordinates. Products of three coordinates are used in intermediate steps so watch out for overflow if using int or long long.

```
/* Usage:
 * vector<P> inter = segInter(s1,e1,s2,e2);
 * if (sz(inter)==1)
     cout << "segments intersect at " << inter[0] <<
     endl;
 * Status: stress-tested, tested on kattis:intersection
#include "Point.h"
#include "OnSegment.h"
template<class P> vector<P> segInter(P a, P b, P c, P
 → d) {
  auto oa = c.cross(d, a), ob = c.cross(d, b);
       oc = a.cross(b, c), od = a.cross(b, d);
  // Checks if intersection is single non-endpoint
  if (sgn(oa) * sgn(ob) < 0 \&\& sgn(oc) * sgn(od) < 0)
    return {(a * ob - b * oa) / (ob - oa)};
  set<P> s;
  if (onSegment(c, d, a)) s.insert(a);
  if (onSegment(c, d, b)) s.insert(b);
```

```
if (onSegment(a, b, c)) s.insert(c);
if (onSegment(a, b, d)) s.insert(d);
return {all(s)};
```

3.1.10 Side Of

Returns where *p* is as seen from *s* towards *e*. $1/0/-1 \Leftrightarrow \text{left/on}$ line/right. If the optional argument *eps* is given 0 is returned if *p* is within distance *eps* from the line. P is supposed to be Point<T> where T is e.g. double or long long. It uses products in intermediate steps so watch out for overflow if using int or long long.

```
/* Usage:
* bool left = sideOf(p1,p2,q)==1;
* Status: tested
#include "Point.h"
template<class P>
int sideOf(P s, P e, P p) { return sgn(s.cross(e, p));
template<class P>
int sideOf(const P& s, const P& e, const P& p, double
→ eps) {
 auto a = (e-s).cross(p-s);
 double l = (e-s).dist()*eps;
 return (a > l) - (a < -l);
```

3.2 3D 3.2.1 3D Convex Hull

```
#define ll long long
#define sz(x) ((int) (x).size())
#define all(x) (x).begin(), (x).end()
#define vi vector<int>
#define pii pair<int, int>
#define rep(i, a, b) for(int i = (a); i < (b); i++)
using namespace std;
template<typename T>
using minpq = priority queue<T, vector<T>, greater<T>>;
typedef long double ftype;
struct pt3 {
  ftype x, y, z;
  pt3(ftype x = 0, ftype y = 0, ftype z = 0) : x(x),
   \rightarrow y(y), z(z) {}
  pt3 operator-(const pt3 &o) const {
    return pt3(x - o.x, y - o.y, z - o.z);
  pt3 cross(const pt3 &o) const {
    return pt3(y * o.z - z * o.y, z * o.x - x * o.z, x
     \rightarrow * o.y - y * o.x);
  ftype dot(const pt3 &o) const {
    return x * 0.x + y * 0.y + z * 0.z;
// A face is represented by the indices of its three
   points a, b, c.
// It also stores an outward-facing normal vector q
struct face {
  int a, b, c;
  pt3 q;
// modify this depending on the coordinate sizes in

→ your use case

const ftype EPS = 1e-9;
vector<face> hull3(const vector<pt3> &p) {
  int n = sz(p);
  assert(n >= 3);
  vector<face> f;
```

```
// Consider an edge (a->b) dead if it is not a CCW
   edge of some current face
// If an edge is alive but not its reverse, this is
   an exposed edge.
// We should add new faces on the exposed edges.
vector<vector<bool>> dead(n, vector<bool>(n, true));
auto add_face = [&](int a, int b, int c) +
 f.push_back({a, b, c, (p[b] - p[a]).cross(p[c] -
     p[a])});
  dead[a][b] = dead[b][c] = dead[c][a] = false;
// Initialize the convex hull of the first 3 points
// triangular disk with two faces of opposite
add_face(0, 1, 2);
add face(0, 2, 1);
rep(i, 3, n) {
 // f2 will be the list of faces invisible to the

→ added point p[i]

  vector<face> f2;
  for(face &F : f)
    if((p[i] - p[f.a]).dot(F.q) > EPS) {
      // this face is visible to the new point, so

→ mark its edges as dead
      dead[F.a][F.b] = dead[F.b][F.c] =

    dead[F.c][F.a] = true;

    }else {
      f2.push back(F);
  // Add a new face for each exposed edge.
 // Only check edges of alive faces for being
     exposed.
 f.clear();
for(face &F : f2) {
    int arr[3] = {F.a, F.b, F.c};
    rep(j, 0, 3)
      int a = arr[j], b = arr[(j + 1) % 3];
      if(dead[b][a]) {
        add face(b, a, i);
  f.insert(f.end(), all(f2));
return f;
```

3.2.2 Point3D

Class to handle points in 3D space.T can be e.g. double or long long.

```
template<class T> struct Point3D {
 typedef Point3D P;
 typedef const P& R;
 explicit Point3D(T x=0, T y=0, T z=0) : x(x), y(y),
    z(z) {}
 bool operator<(R p) const {</pre>
   return tie(x, y, z) < tie(p.x, p.y, p.z); }
 bool operator==(R p) const {
   return tie(x, y, z) == tie(p.x, p.y, p.z); }
 P operator+(R p) const { return P(x+p.x, y+p.y,
     z+p.z);
 P operator-(R p) const { return P(x-p.x, y-p.y,
     z-p.z);
 P operator*(T d) const { return P(x*d, y*d, z*d);
 P operator/(T d) const { return P(x/d, y/d, z/d); }
 T dot(R p) const { return x*p.x + y*p.y + z*p.z; }
 P cross(R p) const {
   return P(y*p.z - z*p.y, z*p.x - x*p.z, x*p.y -
    \rightarrow y*p.x);
```

```
dist2() const { return x*x + y*y + z*z; }
double dist() const { return sqrt((double)dist2()); }
//Azimuthal angle (longitude) to x-axis in interval
double phi() const { return atan2(y, x); }
//Zenith angle (latitude) to the z-axis in interval
double theta() const { return
\rightarrow atan2(sqrt(x*x+y*y),z); }
P unit() const { return *this/(T)dist(); } //makes
   dist()=1
//returns unit vector normal to *this and p
P normal(P p) const { return cross(p).unit(); }
//returns point rotated 'angle' radians ccw around
P rotate(double angle, P axis) const {
  double s = sin(angle), c = cos(angle); P u =
     axis.unit();
  return u*dot(u)*(1-c) + (*this)*c - cross(u)*s;
```

3.2.3 Polyhedron Volume

Description: Magic formula for the volume of a polyhedron. Faces should point outwards.

3.2.4 Spherical Distance

Returns the shortest distance on the sphere with radius radius between the points with azimuthal angles (longitude) f1 (ϕ_1) and f2 (ϕ_2) from x axis and zenith angles (latitude) f1 (f1) and f2 (f2) from z axis (f1) and f2 (f2) from z axis (f1) and f2 (f2) from z axis by converting the spherical coordinates to cartesian coordinates so if that is what you have you can use only the two last rows. dx*radius is then the difference between the two points in the x direction and d*radius is the total distance between the points.

3.3 Circle

3.3.1 Circle Intersection

Computes the pair of points at which two circles intersect. Returns false in case of no intersection.

```
*out = {mid + per, mid - per};
return true;
```

3.3.2 Circle Polygon Intersection

Returns the area of the intersection of a circle with a ccw polygon. Time: O(n)

```
#include "Point.h"
typedef Point<double> P;
#define arg(p, q) atan2(p.cross(q), p.dot(q))
double circlePoly(P c, double r, vector<P> ps) {
 auto tri = [\&](Pp, Pq) {
    auto r2 = r * r' / 2;
    Pd = q - p;
    auto a = d.dot(p)/d.dist2(), b =
       (p.dist2()-r*r)/d.dist2();
    auto det = a * a - b;
    if (det <= 0) return arg(p, q) * r2;</pre>
    auto s = max(0., -a-sqrt(det)), t = min(1.,
       -a+sqrt(det));
   if (t < 0 \mid | 1 \le s) return arg(p, q) * r2;
P u = p + d * s, v = p + d * t;
    return arg(p,u) * r2 + u.cross(v)/2 + arg(v,q) *

    r2;

 auto sum = 0.0
  rep(i,0,sz(ps))
    sum += tri(ps[i] - c, ps[(i + 1) % sz(ps)] - c);
  return sum;
```

3.3.3 Circle Tangents

Finds the external tangents of two circles, or internal if r2 is negated. Can return 0, 1, or 2 tangents -0 if one circle contains the other (or overlaps it, in the internal case, or if the circles are the same); 1 if the circles are tangent to each other (in which case .first = .second and the tangent line is perpendicular to the line between the centers). .first and .second give the tangency points at circle 1 and 2 respectively. To find the tangents of a circle with a point set r2 to 0.

3.3.4 CircumCircle

The circumcirle of a triangle is the circle intersecting all three vertices. ccRadius returns the radius of the circle going through points A, B and C and ccCenter returns the center of the same circle.



```
#include "Point.h"
typedef Point<double> P;
double ccRadius(const P& A, const P& B, const P& C) {
  return (B-A).dist()*(C-B).dist()*(A-C).dist()/
    abs((B-A).cross(C-A))/2;
```

```
ccCenter(const P& A, const P& B, const P& C) {
P b = C-A, c = B-A;
return A +
- (b*c.dist2()-c*b.dist2()).perp()/b.cross(c)/2;
```

3.4 Polygon 3.4.1 Hull Diameter

Returns the two points with max distance on a convex hull (ccw, no duplicate/collinear points).

3.4.2 Line Hull Intersection

Line-convex polygon intersection. The polygon must be ccw and have no collinear points. lineHull(line, poly) returns a pair describing the intersection of a line with the polygon:

```
(-1,-1) if no collision,
```

(i,-1) if touching the corner i,

(i,i) if along side (i,i+1),

(i,j) if crossing sides (i,i+1) and (j,j+1). In the last case, if a corner i is crossed, this is treated as happening on side (i,i+1). The points are returned in the same order as the line hits the polygon. extrVertex returns the point of a hull with the max projection onto a line. Time: $O(\log n)$

```
#include "Point.h"
#define cmp(i,j)

    sgn(dir.perp().cross(poly[(i)%n]-poly[(j)%n]))
#define extr(i) cmp(i + 1, i) \Rightarrow 0 && cmp(i, i - 1 +
\rightarrow n) < 0
template <class P> int extrVertex(vector<P>& poly, P

    dir) {

  int n = sz(poly), lo = 0, hi = n;
  if (extr(0)) return 0;
  while (lo + 1 < hi) {
   int m = (lo + hi) / 2;
    if (extr(m)) return m;
    int ls = cmp(lo + 1, lo), ms = cmp(m + 1, m);
    (ls < ms \mid | (ls == ms \&\& ls == cmp(lo, m)) ? hi :
     \rightarrow lo) = m;
  return lo;
#define cmpL(i) sgn(a.cross(poly[i], b))
template <class P>
array<int, 2> lineHull(P a, P b, vector<P>& poly) {
  int endA = extrVertex(poly, (a - b).perp());
  int endB = extrVertex(poly, (b - a).perp());
  if (cmpL(endA) < 0 \mid | cmpL(endB) > 0)
    return {-1, -1};
  array<int, 2> res;
  rep(i,0,2) {
```

3.4.3 Polygon Center

Returns the center of mass for a polygon. Time: O(n)

```
#include "Point.h"
typedef Point<double> P;
P polygonCenter(const vector<P>& v) {
   P res(0, 0); double A = 0;
   for (int i = 0, j = sz(v) - 1; i < sz(v); j = i++) {
      res = res + (v[i] + v[j]) * v[j].cross(v[i]);
      A += v[j].cross(v[i]);
   }
   return res / A / 3;</pre>
```

3.4.4 Polygon Cut

Returns a vector with the vertices of a polygon with everything to the left of the line going from s to e cut away.

```
vector<P> p = ...;
    p = polygonCut(p, P(0,0), P(1,0));
   Status: tested but not extensively
#include "Point.h"
#include "lineIntersection.h"
typedef Point<double> P;
vector<P> polygonCut(const vector<P>& poly, P s, P e)
  vector<P'> réš;
  rep(i,0,sz(poly)) {
    P`cur' = poly[i], prev = i ? poly[i-1] :
        poly.back();
    bool side = s.cross(e, cur) < 0;</pre>
    if (side != (s.cross(e, prev) < 0))</pre>
      res.push back(lineInter(s, e, cur, prev).second);
    if (side)
      rès.push back(cur);
  return res;
```

3.5 Closest Pair

Finds the closest pair of points. Time: $O(n \log n)$

```
#include "Point.h"
typedef Point<ll> P;
pair<P, P> closest(vector<P> v) {
   assert(sz(v) > 1);
   set<P> S;
   sort(all(v), [](P a, P b) { return a.y < b.y; });
   pair<ll, pair<P, P>> ret{LLONG_MAX, {P(), P()}};
   int j = 0;
   for (P p : v) {
```

3.6 Convex Hull

```
typedef long long ll;
typedef pair <ll, ll> point;
#define x first
#define y second
inline | l area (point a, point b, point c) {
  return (b.x - a.x) * (c.y - a.y) - (b.y - a.y) *
  vector <point> convexHull (vector <point> p) {
  int n = p.size(), m = 0;
  if (n < 3) return p;
  vector <point> hull(n + n);
  sort(p.begin(), p.end());
  for (int i = 0; i < n; ++i)
    while (m > 1) and area (hull[m - 2], hull[m - 1],
     \rightarrow p[i]) \leq 0) --m;
    hull[m++] = p[i];
  for (int i = n - 2, j = m + 1; i >= 0; --i) {
    while (m >= j and area(hull[m - 2], hull[m - 1],
     \rightarrow p[i]) \leq 0) --m;
    hull[m++1] = p[i]:
  hull.resize(m - 1); return hull;
```

3.7 Minimum Enclosing Circle

```
// Expected runtime: O(n)
// Solves Gvm 102299J
using namespace std;
typedef long double ld;
typedef pair <ld, ld> point;
#define x first
#define y second
point operator + (const point &a, const point &b) {
  return point(a.x + b.x, a.y + b.y);
point operator - (const point &a, const point &b) {
  return point(a.x - b.x, a.y - b.y);
point operator * (const point &a, const ld &b) {
  return point(a.x * b, a.y * b);
point operator / (const point &a, const ld &b) {
  return point(a.x / b, a.y / b);
const ld EPS = 1e-8;
const ld INF = 1e20
const ld PI = acosl(-1);
inline ld dist (point a, point b) {
  return hypotl(a.x - b.x, a.y - b.y);
|inline ld sqDist (point a, point b) {
  return (a.x - b.x) * (a.x - b.x) + (a.y - b.y) *
   \rightarrow (a.y - b.y);
```

```
inline ld dot (point a, point b) {
  return a.x * b.x + a.y * b.y;
inline ld cross (point a, point b) {
  return a.x * b.y - a.y * b.x;
inline ld cross (point a, point b, point c) {
  return cross(b - a, c - a);
inline point perp (point a) {
  return point(-a.y, a.x);
// circle through 3 points
pair <point, ld> getCircle (point a, point b, point c)
  pair <point, ld> ret;
  ld den = (ld) 2 * cross(a, b, c);
  ret.x.x = ((c.y - a.y) * (dot(b, b) - dot(a, a)) -
   \rightarrow (b.y - a.y) * (dot(c, c) - dot(a, a))) / den;
  ret.x.y = ((b.x - a.x) * (dot(c, c) - dot(a, a)) -
  \leftarrow (c.x - a.x) * (dot(b, b) - dot(a, a))) / den;
  ret.y = dist(ret.x, a);
  return ret;
pair <point, ld> minCircleAux (vector <point> &s,
  point a, point b, int n) {
  ld'lo = -INF, hi = INF;
  for (int i = 0; i < n; ++i) {
    auto si = cross(b - a, s[i] - a);</pre>
    if (fabs(si) < EPS) continue
    point m = getCircle(a, b, s[i]).x;
    auto cr = cross(b - a, m - a);
    si < 0 ? hi = min(hi, cr) : lo = max(lo, cr);
  1d \ v = 0 < lo ? lo : hi < 0 ? hi : 0;
  point c = (a + b) * 0.5 + perp(b - a) * v /

    sqDist(a, b);

  return {c, sqDist(a, c)};
pair <point, ld> minCircle (vector <point> &s, point
\rightarrow a. int n)
  random shuffle(s.begin(), s.begin() + n);
  point \bar{b} = s[0], c = (a + b) * 0.5;
  ld r = sqDist(a, c);
  for (int i = 1; i < n; ++i) {
  if (sqDist(s[i], c) > r * (1 + EPS))
      tie(c, r) = n == s.size() ? minCircle(s, s[i],

→ i) : minCircleAux(s, a, s[i], i);

  return {c, r};
pair <point, ld> minCircle (vector <point> s) {
  assert(!s.empty());
  if (s.size() == 1) return {s[0],
  return minCircle(s, s[0], s.size());
int n; vector <point> p;
int main() {
  cin >> n;
  while (n--) {
    double x, y;
scanf("%lf %lf", &x, &y);
    p.emplace back(x, y);
  pair <point, ld> circ = minCircle(p);
  printf("%0.12f %0.12f %0.12f\n", (double) circ.x.x,
     (double) circ.x.y, (double) (0.5 * circ.y));
  return 0;
```

```
3.8 Point In Polygon
// Test if a point is inside a convex polygon in O(lg
// Solves SPOJ INOROUT typedef long long ll;
typedef pair <ll, ll> point;
#define x first
#define y second
struct segment {
  point P1, P2;
  segment () {}
  segment (point P1, point P2) : P1(P1), P2(P2) {}
|inline ll ccw (point A, point B, point C) {
  return (B.x - A.x) * (C.y - A.y) - (C.x - A.x) *
   \hookrightarrow (B.v - A.v);
inline bool pointOnSegment (segment S, point P) {
  ll x = P.x, y = P.y, x1 = S.P1.x, y1 = S.P1.y, x2 =
   \rightarrow S.P2.x, y2 = S.P2.y;
  ll a = x - x1, b = y - y1, c = x2 - x1, d = y2 - y1,
      dot = a * c + b * d, len = c * c + d * d;
  if (x1 == x2 \text{ and } y1 == y2) return x1 == x and y1 == y2
  if (dot < 0 or dot > len) return 0;
  return x1 * len + dot * c == x * len and y1 * len +
   \rightarrow dot * d == y * len;
const int M = 17
const int N = 10010;
|struct polygon {
  int n; // n > 1
  point p[N]; // clockwise order
  polvaon () {}
  polygon (int n, point *T) {
    for (int i = 0; i < n; ++i) p[i] = T[i];
  bool contains (point P, bool strictlyInside) {
    int lo = 1, hi = n - 1;
    while (lo < hi){
      int mid = lo + hi >> 1
      if (ccw(p[0], P, p[mid]) > 0) lo = mid + 1;
      else hi = mid;
    if (ccw(p[0], P, p[lo]) > 0) lo = 1;
    if (!strictlyInside and
        pointOnSegment(segment(p[0], p[n - 1]), P))
        return 1
    if (!strictlyInside and
        pointOnSegment(segment(p[lo], p[lo - 1]), P))
        return 1:
    if (lo == 1 or ccw(p[0], P, p[n - 1]) == 0) return
    return ccw(p[lo], P, p[lo - 1]) < 0;
3.9 near_pair
struct pt {
 int x, y, id;
```

bool operator()(const pt & a, const pt & b) const {

return a.x < b.x | | (a.x == b.x && a.y < b.y);

struct cmp x {

};

```
struct cmp_y {
  bool operator()(const pt & a, const pt & b) const {
    return a.y < b.y;
int n;
vector<pt> a;
double mindist:
pair<int, int> best pair;
void upd ans(const pt & a, const pt & b) {
  double dist = sqrt((a.x - b.x)*(a.x - b.x) + (a.y - b.x)
   \rightarrow b.y)*(a.y - b.y));
  if (dist < mindist) {</pre>
    mindist = dist;
    best pair = {a.id, b.id};
vector<pt> t;
void rec(int l, int r) {
  if (r - l <= 3) {
    for (int i = l; i < r; ++i) {
  for (int j = i + 1; j < r; ++j) {</pre>
        upd ans(a[i], a[j]);
    sort(a.begin() + l, a.begin() + r, cmp y());
    return;
  int m = (l + r) >> 1;
  int midx = a[m].x;
  rec(l, m);
  rec(m, r);
  merge(a.begin() + l, a.begin() + m, a.begin() + m,
  \rightarrow a.begin() + r, t.begin(), cmp y());
  copy(t.begin(), t.begin() + r - l, a.begin() + l);
  for (int i = l; i < r; ++i) {
    if (abs(a[i].x - midx) < mindist) {</pre>
      for (int j = tsz - 1; j >= 0 \&\& a[i].y - t[j].y
       upd ans(a[i], t[j]);
      t[tsz++] = a[i];
void solve(int n)
  t.resize(n);
  sort(a.begin(), a.end(), cmp x());
  mindist = 1E20;
  rec(0, n);
3.10 sweep
const double EPS = 1E-9;
struct pt {
  double x, y;
struct seg {
  pt p, q;
  int id;
```

double get y(double x) const {

return p.y + (q.y - p.y) * (x - p.x) / (q.x - p.x);

bool intersect1d(double l1, double r1, double l2,

if (abs(p.x - q.x) < EPS)

return p.y;

double r2) {

```
if (l1 > r1)
    swap(l1, r1);
  if (l2 > r2)
  swap(l2, r2);
  return max(l1, l2) \le min(r1, r2) + EPS;
int vec(const pt& a, const pt& b, const pt& c) {
  double s = (b.x - a.x) * (c.y - a.y) - (b.y - a.y) *
     (c.x - a.x);
  return abs(s) < EPS ? 0 : s > 0 ? +1 : -1;
bool intersect(const seg& a, const seg& b)
  return intersect1d(a.p.x, a.q.x, b.p.x, b.q.x) &&
    intersect1d(a.p.y, a.q.y, b.p.y, b.q.y) &&
    vec(a.p, a.q, b.p) * vec(a.p, a.q, b.q) <= 0 &&
    vec(b.p, b.q, a.p) * vec(b.p, b.q, a.q) <= 0;
bool operator<(const seg& a, const seg& b)
  double x = max(min(a.p.x, a.q.x), min(b.p.x, b.q.x));
  return a.get y(x) < b.get y(x) - EPS;
struct event {
  double x;
  int tp, id;
  event() {}
  event(double x, int tp, int id) : x(x), tp(tp),

    id(id) {}
  bool operator<(const event& e) const {</pre>
    if (abs(x - e.x) > EPS)
      return x < e.x;
    return tp > e.tp;
set<seg> s;
vector<set<seg>::iterator> where;
set<seg>::iterator prev(set<seg>::iterator it) {
 return it == s.begin() ? s.end() : --it;
set<seg>::iterator next(set<seg>::iterator it) {
 return ++it;
pair<int, int> solve(const vector<seg>& a) {
  int n = (int)a.size();
  vector<event> e;
  for (int i = 0; i < n; ++i) {
    e.push back(event(min(a[i].p.x, a[i].q.x), +1, i));
    e.push_back(event(max(a[i].p.x, a[i].q.x), -1, i));
  sort(e.begin(), e.end());
  s.clear();
  where.resize(a.size());
  for (size_t i = 0; i < e.size(); ++i) {</pre>
    int id = e[i].id;
    if (e[i].tp == +1) {
      set<seg>::iterator nxt = s.lower bound(a[id]),
         prv = prev(nxt);
      if (nxt != s.end() && intersect(*nxt, a[id]))
        return make pair(nxt->id, id);
      if (prv != s.end() && intersect(*prv, a[id]))
        return make pair(prv->id, id);
      where[id] = s_insert(nxt, a[id]);
    } else ·
      set<seg>::iterator nxt = next(where[id]), prv =

    prev(where[id]);

      if (nxt != s.end() && prv != s.end() &&

→ intersect(*nxt, *prv))
        return make pair(prv->id, nxt->id);
      s.erase(where[id]);
```

```
return make pair(-1, -1);
4 Graph
4.1 2Sat
namespace sat{
  const int MAX = 200010;
  bool vis[MAX];
  vector <int> ed[MAX], rev[MAX]
  int n, m, ptr, dfs_t[MAX], ord[MAX], par[MAX];
  inline int inv(int x){
    return ((x) \le n ? (x + n) : (x - n));
  void init(int vars){
    n = vars, m = vars << 1;
    for (int i = 1; i <= m; i++){
  ed[i].clear();
  rev[i].clear();</pre>
  inline void add(int a, int b){
    ed[a].push back(b);
    rev[b].push back(a);
  inline void OR(int a, int b){
    add(inv(a), b);
    add(inv(b), a);
  inline void AND(int a, int b){
    add(a, b);
    add(b, a);
  void XOR(int a,int b){
    add(inv(b), a);
add(a, inv(b));
    add(inv(a), b);
    add(b, inv(a));
  inline void XNOR(int a, int b){
    add(a,b);
    add(b,a)
    add(inv(a), inv(b));
add(inv(b), inv(a));
  inline void force true(int x){
    add(inv(x), x);
  inline void topsort(int s){
    vis[s] = true;
    for(int x : rev[s]) if(!vis[x]) topsort(x);
    dfs t[s] = ++ptr;
  inline void dfs(int s. int p){
    par[s] = p;
    vis[s] = true;
    for(int x : ed[s]) if (!vis[x]) dfs(x, p);
  void build(){
    CLR(vis);
    for(int i=m;i>=1;i--)
      if (!vis[i]) topsort(i);
      ord[dfs t[i]] = i;
    for (int i = m; i >= 1; i--){
      int x = ord[i];
      if (!vis[x]) dfs(x, x);
```

```
}
bool satisfy(vector <int>& res){
  build();
  CLR(vis);
  for (int i = 1; i <= m; i++){
    int x = ord[i];
    if (par[x] == par[inv(x)]) return false;
    if (!vis[par[x]]){
      vis[par[x]] = true;
      vis[par[inv(x)]] = false;
    }
}
res.clear();
for (int i = 1; i <= n; i++){
    if (vis[par[i]]) res.push_back(i);
}
return true;
}</pre>
```

4.2 Centroid decomp

```
vector <int> ed[MAX]:
bool isCentroid MAX
int sub[MAX], cpar[MAX], clevel[MAX];
int dis[20][MAX];
void calcSubTree(int s,int p) {
  sub[s] = 1;
  for(int x : ed[s]) {
   if(x == p or isCentroid[x]) continue;
    calcSubTree(x,s);
    sub[s] += sub[x1;
int nn;
int getCentroid(int s,int p) {
 for(int x : ed[s]) {
    if(!isCentroid[x] && x!=p && sub[x]>(nn/2)) return

→ getCentroid(x,s);

  return s;
void setDis(int s, int from, int p, int lev) {
  dis[from][s] = lev;
  for(int x : ed[s]) {
    if(x == p or isCentroid[x] ) continue;
    setDis(x, from, s, lev+1);
void decompose(int s,int p,int lev) {
  calcSubTree(s,p);
  nn = sub[s];
  int c = getCentroid(s,p);
  setDis(c,lev,p,0);
  isCentroid[c] = true;
  cpar[c] = p;
  clevel[c] = lev;
  for(int x : ed[c])
    if(!isCentroid[x]) decompose(x,c,lev+1);
int ans[MAX];
inline void update(int v) {
 int u = v;
  while(u!=-1)
    ans[u] = min(ans[u], dis[clevel[u]][v]);
    u = cpar[u];
```

```
inline int query(int v) {
  int ret = INF;
  int u = v;
 while(u != -1) {
   ret = min(ret, dis[clevel[u]][v]+ans[u]);
   u = cpar[u];
  return ret;
int main()
  decompose(1,-1,0);
  for(int i=1; i<=n; i++) ans[i] = INF;
  update(v);
 query(v));
  return 0;
```

4.3 articulation point

```
using namespace std;
const int N = 1e5 + 10:
vector<int> g[N];
int vis[N], \tilde{low}[N], cut[N], now = 0, n, m;
void dfs(int u, int p) {
  low[u] = vis[u] = ++now; int ch = 0;
  for(int v : g[u]){
    if(v ^ p)
      if(vis[v]) low[u] = min(low[u], vis[v]);
        ch++; dfs(v, u);
low[u] = min(low[u], low[v]);
        if(p + 1 \&\& low[v] >= vis[u]) cut[u] = 1;
        if(low[v] > vis[u]) {
  printf("Bridge %d -- %d\n", u, v);
  } if(p == -1 \&\& ch > 1) cut[u] = 1;
void ArticulationPointAndBridge() {
  now = 0;
  for(int i = 0; i < n; i++) {
    if(!vis[i]) dfs(i, -1);
```

```
4.4 bcc
// clear ed[] every test case
// tot -> total number of components
// bcc[i] contains the nodes of the i'th component
// any self loop or multiple edge?
const int MAX = ?;
vector <int> ed[MAX];
bool cut[MAX];
int tot, Time, low[MAX], st[MAX];
vector <int>_bcc[MAX];
stack <int> S;
void popBCC(int s,int x) {
  cut[s] = 1;
bcc[++tot].pb(s);
  while(bcc[tot].back() ^ x) {
  bcc[tot].pb(S.top());
    S.pop();
void dfs(int s, int p = -1) {
  S.push(s);
  int ch = 0;
  st[s] = low[s] = ++Time;
  for(int x : ed[s]) {
    if(!st[x]) {
       ch++;
```

```
dfs(x,s);
low[s] = min(low[s],low[x]);
if(p != -1 and low[x] >= st[s]) popBCC(s,x);
       else if(p == -1) if(ch > 1) popBCC(s,x);
    else if(p != x) low[s] = min(low[s],st[x]);
  if(p == -1 \&\& ch > 1) cut[s] = 1;
void processBCC(int n) {
  for(int i=1;i<=n;i++) bcc[i].clear();</pre>
  CLR(st); CLR(cut);
  Time = tot = 0;
for(int i=1; i<=n; i++) {
    if(!st[i]) {
       dfs(i,-1);
if(!S.empty()) ++tot;
       while(!S.empty()) {
         bcc[tot].push back(S.top());
         S.pop();
```

4.5 bridge tree const int MAXN = ?:

const int MAXE = ?;

```
|struct edges {
 int u,v;
|} ara[MÁXÉ];
vector <int> ed[MAXN];
vector <int> isBridge[MAXN];
|vector <int> brTree[MAXN];
bool vis[MAXN]:
int st[MAXN], low[MAXN], Time = 0;
int cnum;
int comp[MAXN];
|void findBridge(int s,int par) {
  int i,x,child = 0,j;
  vis[s] = 1;
  Time++;

st[s] = low[s] = Time;

for(i=0; i<ed[s].size(); i++) {
     x = ed[s][i];
if(!vis[x]) {
        child++;
        findBridge(x,s);
        low[s] = min(low[s], low[x]);
        if(low[x] > st[s]) {
  isBridge[s][i] = 1
          j = lower bound(ed[x].begin(),ed[x].end(),s)-e
           \rightarrow d[x].begin();
          isBridge[x][j] = 1;
     else if(par!=x)
        low[s] = min(low[s], st[x]);
void dfs(int s) {
  int i,x;
vis[s] = 1;
comp[s] = cnum;
  for(i=0; i<ed[s].size(); i++) {</pre>
    if(!isBridge[s][i]) {
    x = ed[s][i];
    if(!vis[x]) dfs(x);
```

```
void processBridge(int n,int m) {
  CLR(vis);
  Time = 0
  for(int i=1; i<=n; i++) if(!vis[i]) findBridge(i,-1);</pre>
  CLR(vis);
  for(int i=1; i<=n; i++) {
    if(!vis[i]) {
    cnum++;
    dfs(i);
  n = cnum;
  for(int i=1; i<=m; i++) {</pre>
     if(comp[ara[i].u] != comp[ara[i].v]) {
   brTree[comp[ara[i].u]].pb(comp[ara[i].v]);
   brTree[comp[ara[i].v]].pb(comp[ara[i].u]);
int main() {
  int n,m,u,v;
scanf("%d %d",&n,&m);
for(int i=1; i<=m; i++) {</pre>
     sii(u,v);
     ed[u].pb(v);
     ed[v].pb(u);
     isBridge[u].pb(0);
     isBridge[v].pb(0);
     ara[i].u = u;
ara[i].v = v;
  for(int i=1; i<=n; i++) sort(all(ed[i]));</pre>
  processBridge(n,m);
  return 0;
```

4.6 dinic

```
namespace dinic {
 using T = int;
 const T INF = 0x3f3f3f3f;
 const int MAXN = 5010;
 int n, src, snk, work[MAXN];
 T dist[MAXN];
 struct Edge{
   int to, rev pos;
   Tc, f;
 vector <Edge> ed[MAXN];
 void init(int n, int src, int snk) {
   n = n, src = src, snk = snk;
   for(int i=1;i<=n;i++) ed[i].clear();
 inline void addEdge(int u, int v, T c, T rc = 0) {
   Edge a = \{v, (int)ed[v].size(), c, 0\};
   Edge b = \{u, (int)ed[u].size(), rc, 0\};
   ed[u].push back(a);
   ed[v].push_back(b);
 bool dinic bfs() {
   SET(dist);
   dist[src] = 0;
   queue <int> q;
   q.push(src);
   while(!q.empty()){
     int u = q.front(); q.pop();
     for(Edge &e : ed[u]){
       if(dist[e.to] == -1 \text{ and } e.f < e.c) 
          dist[e.to] = dist[u] + 1;
```

```
q.push(e.to);
    return (dist[snk]>=0);
 T dinic_dfs(int u, T fl){
    if (u == snk) return fl;
    for (; work[u] < (int)ed[u].size(); work[u]++) {</pre>
      Edge &e = ed[u][work[u]];
      if (e.c <= e.f) continue;</pre>
      int v = e.to;
      if (dist[v] == dist[u] + 1)
        T df = dinic dfs(v, min(fl, e.c - e.f));
        if (df > 0){
          e.f += df:
          ed[v][e.rev_pos].f -= df;
          return df;
    return 0;
    <u>s</u>olve()
    T ret = 0:
    while (dinic bfs()) {
      CLR(work);
      while (T delta = dinic dfs(src, INF)) ret +=

→ delta:

    return ret;
int main() {
  int n, m, u, v, c;
  cin >> n >> m;
  dinic::init(n, 1, n);
  while(m--) {
    cin >> u >> v >> c;
    dinic::addEdge(u, v, c, c);
  cout << dinic::solve() << '\n';
  return 0;
4.7 dsu_on_tree
```

```
void calcSubSize(int s,int p) {
 sub[s] = 1;
 for(int x : G[s]) {
   if(x==p) continue;
   calcSubSize(x,s);
   sub[s] += sub[x];
void add(int s,int p,int v,int bigchild = -1) {
 freq[color[s]] += v;
 for(int x : G[s]) {
   if(x==p || x==bigchild) continue;
   add(x,s,v);
void dfs(int s,int p,bool keep) {
 int bigChild = -1;
 for(int x : G[s]) {
   if(x==p) continue;
   if(bigChild==-1 || sub[bigChild] < sub[x] )</pre>
    → bigChild = x;
 for(int x : G[s]) ∃
   if(x==p || x==bigChild) continue;
   dfs(x,s,0);
 if(bigChild!=-1) dfs(bigChild,s,1);
```

```
add(s,p,1,bigChild);
  if(keep==0)
    add(s,p,-1);
4.8 euler_path
vector <int> ed[MAX+5], sltn;
int inDeg[MAX+5], outDeg[MAX+5];
bool vis[MAX+5];
void dfs(int nd) {
  vis[nd] = true;
  while(ed[nd].size())
    int v = ed[nd].back();
    ed[nd].pop_back();
    dfs(v);
  sltn.pb(nd);
|int findEuler (int n) {
  int src , snk , ret = 1;
bool found_src = false, found_snk = false;
  CLR(inDeg); CLR(outDeg);
  for(int u = 1; u <= n; u++) {
    for(int i = 0; i<ed[u].size(); i++) {</pre>
       int v = ed[u][i];
       outDeg[u]++;
       inDeg[v]++;
  int diff;
  for(int i = 1; i<=n; i++) {
   diff = outDeg[i] - inDeg[i];</pre>
    if(diff == 1)
       if(found src) return 0;
       found sr\overline{c} = true;
       src = i;
    else if (diff == -1) {
       if(found snk) return 0;
       found sn\overline{k} = true;
       snk = i:
    else if(diff != 0) return 0;
  if(!found src) {
     ret = 2;
     for(int i = 1 ; i <= n ; i++) {
       if( outDeg[i] ) {
         found src = true;
         src = i;
         break;
  if(!found src) return ret;
  CLR(vis);
  sltn.cléar();
  dfs(src);
  for(int i = 1; i<=n; i++) {
    if(outDeg[i] && !vis[i]) return 0;
  for(int i = (int)sltn.size()-1; i>=0; i--)
   → printf("%d ",sltn[i]);
  puts'("");
  return ret;
4.9 hld
const int N = 3e4 + 5;
```

vector<int> G[N];

int sz[N], H[N], P[N];

```
void dfs(int cur, int h)
    sz[cur] = 1;
H[cur] = h;
    for(int& to : G[cur])
         G[to].erase(find(G[to].begin(), G[to].end(),
        P[to] = cur;
dfs(to, h + 1);
sz[cur] += sz[to];
         if(sz[to] > sz[G[cur][0]]) swap(G[cur][0], to);
int base[N], pos[N], head[N];
int ptr = 0;
void hld(int cur)
     pos[cur] = ++ptr;
    base[ptr] = cur;
    for(int to : G[cur]) {
    head[to] = (to == G[cur][0] ? head[cur] : to);
         hld(to);
segtree ST;
int query(int u, int v)
     int ret = 0:
    while(head[u] != head[v])
         if(H[head[u]] > H[head[v]]) swap(u, v);
         ret += ST.query(pos[head[v]], pos[v]);
         v = P[head[v]];
    if(H[u] > H[v]) swap(u, v);
    ret += ST.query(pos[u], pos[v]);
    return ret;
lvoid update(int u, int val) {
    ST.update(pos[u], val);
void build(int n, int root)
    ptr = 0:
    dfs(root, 0);
    head[root] = root;
    hld(root);
    ST = seqtree(n);
   clear graph
 * call build
* Prob : sum of values from u to v
```

4.10 hopcroft_karp

```
struct HopcroftKarp {
 const int N, M;
 std::vector<std::vector<int>> adj left;
 std::vector<int> matchL, matchR;
 HopcroftKarp(int N, int M, const

    std::vector<std::pair<int, int>>& edge)
      : N(N), M(M), matchL(N, -1), matchR(M, -1),
      → adj left(N) {
    for (auto [l, r] : edge)
  adj_left[l].push_back(r);
 int maxmatching() {
    int sz = 0;
    for (bool updated = true; updated;) {
```

```
updated = false;
  static std::vector<int> root(N), prev(N), gq(N);
  static int qi, qj;
// std::queue<int> q;
  qi = qj = 0;
std::fill(root.begin(), root.end(), -1),
  std::fill(prev.begin(), prev.end(), -1);
  for (int i = 0; i < N; i++)
    if (matchL[i] == -1
     qq[qj++] = i, root[i] = i, prev[i] = i;
q.push(i), root[i] = i;
  while (qi < qj) {
    int u = qq[qi++];
    // int u = q.front(); q.pop();
    if (matchL[root[u]] != -1) continue;
    for (int v : adj left[u]) {
      if (matchR[v] = -1) {
        while (v != -1)
           matchR[v] = u, std::swap(matchL[u], v),
           \rightarrow u = prev[u];
         updated = true, sz++;
        break;
      if (prev[matchR[v]] == -1)
        v = matchR[v], prev[v] = u, root[v] =
         → root[u], qq[qj++] = v;
      // v = matchR[v], prev[v] = u, root[v] =
       \rightarrow root[u], q.push(v);
 }
return sz;
```

4.11 hungarian

```
// Given NN matrix A[i][j]. Calculate a permutation
\neg p[i] that minimize A[i][p[i]].
template <typename T>
pair <T, vector <int>> Hungarian (int n, int m, T
c[N][N]) {
 vector <T> v(m), dist(m);
  vector \langle int \rangle L(n, -1), R(m, -1);
  vector <int> index(m), prev(m);
  auto residue = [\&] (int i, int j) {return c[i][j] -

    ∨[j];};

  iota(index.begin(), index.end(), 0);
  for (int f = 0; f < n; ++f) {
  for (int j = 0; j < m; ++j) {
      dist[j] = residue(f, j), prev[j] = f;
    T w; int i, j, l, s = 0, t = 0;
    while (true) {
      if (s == t) {
        l = s, w = dist[index[t++]];
        for (int k = t; k < m; ++k) {
   j = index[k]; T h = dist[j];</pre>
           if (h <= w) {
             if (h < w) t = s, w = h
             index[k] = index[t], index[t++] = j;
        for (int k = s; k < t; ++k) {
            = index[k];
           if (R[j] < 0) goto augment;
      int q = index[s++], i = R[q];
      for (int k = t; k < m; ++k) {
          index[k];
        T h = residue(i, j) - residue(i, q) + w;
```

```
if (h < dist[j]) {
        dist[j] = h, prev[j] = i;
        if (h == w) {
          if (R[i] < 0) goto augment;</pre>
          index[k] = index[t], index[t++] = j;
augment:
  for (int k = 0; k < l; ++k) v[index[k]] +=
     dist[index[k]] - w;
  do
    R[i] = i = prev[j], swap(j, L[i]);
  } while (i ^ f);
1 \text{ ret} = 0;
for (int i = 0; i < n; ++i) ret += c[i][L[i]];
return {ret, L};
```

4.12 kuhn

```
namespace bpm{
 const int L = 105;
 const int R = 105;
 vector <int> G[L];
 int matchR[R], matchL[L], vis[L], it;
 void init(int n) +
    SET(matchL), SET(matchR), CLR(vis);
    for(int i=1;i<=n;i++) G[i].clear();</pre>
 inline void addEdge(int u,int v) { G[u].pb(v); }
 bool dfs(int s) {
  vis[s] = it;
  for(auto x : G[s]) {
    if( matchR[x] == -1 or (vis[matchR[x]] != it and

    dfs(matchR[x])) )

        matchL[s] = x; matchR[x] = s;
        return true:
    return false;
 int solve() {
   int cnt = 0;
    for(int i=1;i<=n;i++) {
      if(dfs(i)) cnt++, it++;
   return cnt;
```

```
// Don't forget to clear ed after test case ends(vt,
using namespace std;
const int MAX = 100010;
int dep[MAX], par[MAX][21];
vector <int> ed[MAX];
void dfs(int s, int p, int d) {
  dep[s] = d, par[s][0] = p;
  for(int x : ed[s]) {
   if(x == p) continue;
    dfs(x, s, d+1);
|void preprocess(int root, int n) {
        lq(n);
 memset(par, -1, sizeof(par));
```

```
dfs(root, -1, 0);
 for(int j=1; j<=LG; j++) {</pre>
    for(int i=1;i<=n;i++) {
      if(par[i][j-1]] = -1) par[i][j] =
      → par[par[i][j-1]][j-1];
int getLCA(int u, int v) {
 if(dep[u] < dep[v]) swap(u, v);
 for(int i=LG;i>=0;i--) {
    if(dep[u] - (1 << i) >= dep[v]) u = par[u][i];
 if(u == v) return u;
 for(int i=LG;i>=0;i--) {
    if (par[u][i] != -1 and par[u][i] - par[v][i]) {
      u = par[u][i], v = par[v][i];
  return par[u][0];
```

4.14 manhattan_MST

```
using namespace std;
using ll = long long;
struct UnionFind {
    vector<int> UF;
    int cnt;
    UnionFind(int N) : UF(N, -1),
,⊕ cnt(N) {}
    int find(int v) { return UF[v] < 0 ? v : UF[v] =</pre>
                                  ,
 find(UF[v]); }
    bool join(int v, int w) {
        if ((v = find(v)) == (w = find(w))) return

→ false;

        if (UF[v] > UF[w]) swap(v, w);
        UF[v] += UF[w];
        UF[w] = v;
        cnt-
        return true;
    bool connected(int v, int w) {
        return find(v) = find(w);
    int getSize(int v) { return -UF[find(v)]; }
template <class T>
struct KruskalMST {
    using Edge = tuple<int, int, T>;
    T mstWeight;
    vector<Edge> mstEdges;
    UnionFind uf;
    KruskalMST(int V, vector<Edge> edges) :

→ mstWeight(),
                                              → uf(V) {
        sort(edges.begin(), edges.end(), [&](const

→ Edge &a, , const Edge &b) 

→
            return qet<2>(a) < qet<2>(b);
        for (auto &&e : edges) {
            if (int(mstEdges.size()) >= V - 1) break;
            if (uf.join(get<0>(e), get<1>(e))) {
                mstEdges.push back(e);
                mstWeight += \overline{q}et < 2 > (e);
        }
template <class T>
```

```
using Edge = typename KruskalMST<T>::Edge;
    static vector<Edge>
        generateCandidates(vector<pair<T, T>>, P) {
        vector<int> id(P.size());
iota(id.begin(), id.end(), 0);
        vector<Edge> ret;
        ret.reserve(P.size() * 4);
        for (int h = 0; h < 4; h++)
             sort(id.begin(), id.end(), [&](int i, int
                 return P[i].first - P[j].first <</pre>
                  → P[j].second - P[i].second;
             map<T, int> M;
            for (int i : id) {
   auto it = M.lower_bound(-P[i].second);
                 for (; it != M.en\overline{d}(); it =
                  → M.erase(it)) {
                      int j = it->second;
                      T dx = P[i].first - P[j].first, dy
                      → = P[i].second, - P[j].second;
                      if (dy > dx) break;
                      ret.emplace back(i, j, dx + dy);
                 M[-P[i].second] = i;
             for (auto &&p : P) {
                 if (h % 2)
                     p.first = -p.first;
                 else
                      swap(p.first, p.second);
        return ret;
    ManhattanMST(const vector<pair<T, T>> &P)
        : KruskalMST<T>(P.size(),

→ generateCandidates(P)) {}
};
int main() {
    int N:
    cin >> N;
    vector<pair<ll, ll>> P(N);
    for (auto &&p : P) cin >> p.first >> p.second;
    ManhattanMST mst(P);
    cout << mst.mstWeight << '\n';</pre>
    for (auto &&[v, w, weight] : mst.mstEdges) cout <<

    v << □ ,. □ << w << '\n'; return 0;
</pre>
```

```
Edge a = \{v, (int)ed[v].size(), cap, cost, 0\};
  Edge b = \{u, (int)ed[u].size(), 0, -cost, 0\};
  ed[u].pb(a);
  ed[v].pb(b);
inline bool SPFA(){
 CLR(vis);
  for(int i=1; i<=n; i++) mCap[i] = dis[i] = INF;</pre>
  queue <int> q;
  dis[src] = 0;
 vis[src] = true;
 q.push(src);
  while(!q.empty()){
    int u = q.front();
    q.pop();
    vis[u] = false;
    for(int i=0; i<(int)ed[u].size(); i++) {</pre>
      Edge &e = ed[\hat{u}][i];
      int v = e.to;
      if(e.cap > e.flow && dis[v] > dis[u] + e.cost){
        dis[v] = dis[u] + e.cost;
        par[v] = u;
        pos[v] = i;
        mCap[v] = min(mCap[u],e.cap - e.flow);
        if(!vis[v]) {
          vis[v] = true;
          q.pūsh(v);
  return (dis[snk] != INF);
inline pair <T, T> solve() {
 T F = 0, C = 0, f;
 int u, v;
 while(SPFA()){
    u = snk:
    f = mCap[u];
    while(u!=src){
      v = par[u]
      ed[v][pos[u]].flow += f; // edge of v-->u
      ed[u][ed[v][pos[u]].rev_pos].flow -= f;
    \dot{C} += dis[snk] * f;
  return make pair(F,C);
```

4.15 mcmf

```
namespace mcmf {
 using T = int;
 const T INF = ?; // 0x3f3f3f3f or

→ 0x3f3f3f3f3f3f3f3fLL
  const int MAX = ?; // maximum number of nodes
  int n, src, snk;
 T dis[MAX], mCap[MAX];
  int par[MAX], pos[MAX];
  bool vis[MAX];
  struct Edge{
   int to, rev pos;
   T cap, cost, flow;
 vector <Edge> ed[MAX];
  void init(int n, int src, int snk) {
   n = n, src = src, snk = snk;
   for(int i=1;i<=n;i++) ed[i].clear();</pre>
 void addEdge(int u, int v, T cap, T cost) {
```

```
4.16 tree_hash_all_root
/*Description : Find hashes of a tree when rooted at

→ each possible

node(unrooted tree isomorphism test).Time : O(n)*/
const int sz = 2e5 + 5, mod = 1e9 + 7;
ll hval[sz], h[sz], dp[sz], rans[sz];
void dfs(vector<vector<int>>> &q, int u = 0, int p =
\rightarrow -1, int val = 0, int up = 0) {
    vector<int> cv, cht; // current child values &

→ heights

    if (u > 0) {
        cv.push back(val);
        cht.push back(up);
    for (int v : g[u])
        if (v - p) {
            cv.push back(dp[v]);
            cht.push back(1 + h[v]);
```

```
sort(cht.begin(), cht.end(), greater<int>());
    if (cv.size() > 1) {
         ll ret[] = \{1, 1\}; // for biggest &

→ 2nd-biggest heights

        for (int i = 0; i < 2; i++)
             for (int value : cv)
    ret[i] = ret[i] * (hval[cht[i]] +

→ value) % mod;

        rans[u] = ret[0]; // biggest is hash for this
        for (int v : g[u])
             if (v - p) {
                 int id = 1;
                 if (cht[0] - 1 - h[v]) id = 0; // v

    is not on the biggest height path

                 val = ret[id] * invmod((hval[cht[id]))
                    + dp[v]) % mod) % mod;
                 /* division of v subtree hash value */
dfs(g, v, u, val, cht[id] + 1);
    } else if (cv.size()) { // Leaf node u OR vertex

→ - 1 has only one child
        if (!up)
            val = 1;
        else
             val = (val + hval[up]) % mod;
         rans[u] = val;
        for (int v : g[u])
             if (v - p) dfs(g, v, u, val, up + 1);
il get(vector<vector<int>> \delta g, int u = 0, int p = -1) {
    h[ù] = 0:
    vector childs:
    for (int v : q[u])
        if (v - p) {
             childs.push_back(get(g, v, u));
            h[u] = max(\overline{h}[u], 1 + h[v]);
    ll ret = 1;
    for (int value : childs) ret = ret * (hval[h[u]] +
     → value) % mod;
    return dp[u] = ret;
int main() { // can remove the g as param, can change

→ to 1-index, multi-test works, no further change
    qet(q);
    dfs(q);
    tree[k] = rans[0] = dp[0];
    // rans[i] = tree hash with i as root
```

5 Math

```
5.1 FWHT
const int N = 1 \ll 20;
// apply modulo if necessary
void fwht xor(int *a, int n, int dir = 0) {
    for (int h = 1; h < n; h <<= 1) {
        for (int i = 0; i < n; i += h << 1)
            for (int j' = i; j' < i + h; ++j) {
                int x = a[j], y = a[j + h];
                a[j] = x + y, a[j + h] = x - y;
                if (dir) a[j] >>= 1, a[j + h] >>= 1;
       }
void fwht or(int *a, int n, int dir = 0) {
    for (int h = 1; h < n; h <<= 1) {
```

for (int i = 0; i < n; i += h << 1) {

```
for (int j = i; j < i + h; ++j)
                                                              5.3 NOD
                 int x = a[j], y = a[j + h];
                                                              |N = input()|
                 a[j] = x, \bar{a}[j + h] = dir ? y - x : x +
                                                              primes = array containing primes till 10^6
void fwht and(int *a, int n, int dir = 0) {
    for (int h = 1; h < n; h <<= 1) {
        for (int i = 0; i < n; i += h << 1)
            for (int j = i; j < i + h; ++j)
int x = a[j], y = a[j + h];
                                                              if N is prime:
                 a[j] = dir ? x - y : x + y, a[j + h] =

    y;

        }
   }
```

```
5.2 FloorSum
long long FloorSumAP(long long a, long long b, long
   long c, long long n){
  if(!a) return (b / c) * (n + 1);
  if(a >= c or b >= c) return ( (n * (n + 1) ) / 2) *
      (a / c) + (n + 1) * (b / c) + FloorSumAP(a % c,
 \Rightarrow b % c, c, n);
long long m = (a * n + b) / c;
  return m * n - FloorSumAP(c, c - b - 1, a, m - 1);
1/\sqrt{0(\log a)} sum^n x * floor(ax + b / c) = g, sum^n
\rightarrow floro(ax + b / c)^2 = h
struct dat {
  long long f, g, h;
  dat(long long f = 0, long long q = 0, long long h =
  \rightarrow 0) : f(f), g(g), h(h) {};
long long mul(long long a, long long b){
 return (a * b) % MOD;
dat query(long long a, long long b, long long c, long
 if(!a) return {mul(n + 1, b / c), mul(mul(mul(b / c,
  \rightarrow n), n + 1), inv2), mul(mul(n + 1, b / c), b /c)};
  long long f, g, h;
  dat nxt;
  if(a >= c or b >= c){
    nxt = query(a % c, b % c, c, n);
    f = (nxt.f + mul(mul(mul(n, n + 1), inv2), a / c)
       + mul(n + 1, b / c)) % MOD;
    g = (nxt.g + mul(a / c, mul(mul(n, n + 1), mul(2 *
        n + 1, inv6))) + mul(mul(b / c, mul(n, n +

☐ 1)), inv2)) % MOD;
    h = (nxt.h + 2) * mul(b / c, nxt.f) + 2 * mul(a / c)
        c, nxt.q) + mul(mul(a / c, a / c), mul(mul(n,
        n + 1, mul(2 * n + 1, inv6))) + <math>mul(mul(b / mul))
       (c, b / c), n + 1) + mul(mul(a / c, b / c),
        mul(n, n + 1)) % MOD;
    return {f, g, h};
  long long m = (a * n + b ) / c;
  nxt = query(c, c - b - 1, a, m - 1);
  f = (mul(m, n) - nxt.f) \% MOD;
  g = mul(mul(m, mul(n, n + 1)) - nxt.h - nxt.f,
  h = (mul(n, mul(m, m + 1)) - 2 * nxt.q - 2 * nxt.f -

→ f) % MOD;

  return {f, q, h};
```

```
ans = ans * 2
else if N is square of a prime:
    ans = ans * 3
else if N != 1;
    ans = ans * 4
5.4 Pollard Rho
we#include <bits/stdc++.h>
using namespace std;
typedef long long ll;
typedef unsigned long long ull;
namespace Rho {
  ull mul (ull a, ull b, ull mod) {
    ll ret = a * b - mod * (ull) (1.L / mod * a * b);
    return ret + mod * (ret < 0) - mod * (ret >= (11)
  ull bigMod (ull a, ull e, ull mod) {
    ull ret = 1;
    while (e) {
      if (e \& 1) ret = mul(ret, a, mod);
      a = mul(a, a, mod), e >>= 1;
    return ret;
  bool isPrime (ull n) {
    if (n < 2 \text{ or } n \% 6 \% 4 != 1) \text{ return } (n | 1) == 3;
    ull a[] = \{2, 325, 9375, 28178, 450775, 9780504,

→ 1795265022

    ull s = builtin ctzll(n - 1), d = n >> s;
    for (ull x : a)
      ull p = bigMod(x % n, d, n), i = s;
      while (p != 1 \text{ and } p != n - 1 \text{ and } x % n \text{ and } i--)
       \rightarrow p = mul(p, p, n);
      if (p != n - 1 and i != s) return 0;
    return 1;
  ull pollard (ull n) {
    auto f = [\&] (ull x) {return mul(x, x, n) + 1;};
    ull x = 0, y = 0, t = 0, prod = 2, i = 1, q;
    while (t++ \% 40 \text{ or } \gcd(\text{prod}, n) == 1) {
      if (x == y) x = ++i, y = f(x);
      if ((q = mul(prod, max(x, y) - min(x, y), n)))
       \rightarrow prod = q;
      x = f(x), y = f(f(y));
    return gcd(prod, n);
  vector <ull> factor (ull n) {
    if (n == 1) return {};
    if (isPrime(n)) return {n};
    ull x = pollard(n);
    auto l = factor(x), r = factor(n / x);
    l.insert(l.end(), r.begin(), r.end());
```

for all p in primes : if p*p*p > N:

break

return l;

N = N/p

while N divisible by p:

count = count + 1 ans = ans * count

```
int t; ll n;
int main() {
 cin >> t:
 while (t--)
   scanf("%ild", &n);
vector <ull> facs = Rho::factor(n);
    sort(facs.begin(), facs.end());
    printf("%d", (int) facs.size());
    for (auto it : facs) printf(" %llu", it);
    puts("");
 return 0:
```

5.5 catalan

```
//Recursive
const int MOD = ....
const int MAX = ....
int catalan[MAX];
void init() {
   catalan[0] = catalan[1] = 1;
    for (int i=2; i<=n; i++) {
    catalan[i] = 0;</pre>
          for (int j=0; j < i; j++)
               catalan[i] += (catalan[j] *

    catalan[i-j-1]) % MOD;

              if (catalan[i] >= MOD) {
   catalan[i] -= MOD;
//Analytical formula:
ans = ncr(2*n,n) - ncr(2*n,n-1) = ncr(2*n,n)/(n+1)
```

```
//r[i][j]= inverse of p[i] modulo p[j] //ans= x[0]+x[1]*p[0]+x[2]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1]*(p
    \rightarrow ]*p[1]*p[2]*...*p[k-2])
  //ans %= ((p[0]*p[1]*p[2]*...*p[k-1])
 for (int i = 0; i < k; ++i) {
                            x[i] = a[i];
                             for (int j = 0; j < i; ++j) {
    x[i] = r[j][i] * (x[i] - x[j]);
                                                         x[i] = x[i] % p[i];
                                                        if (x[i] < 0)
x[i] += p[i];
  ll mul= p[0], res=x[0], tot=1;
F(i,0,k) tot *= p[i];
 F(i,1,k)
     res+= x[i]*mul;
res %= tot;
mul *= p[i];
  res %= mul;
  return res:
```

5.7 derangement

```
int derangement(int n)
if(!n) return n;
if(n <= 2) return n-1;
return (n-1)*(derangement(n-1) + derangement(n-2));
```

```
5.8 diophantine

void print_solution(int a, int b, int c)
{
    int x, y;
    if (a == 0 && b == 0) {
        if (c == 0) {
            cout << "Infinite Solutions Exist" << endl;
        }
        else {
            cout << "No Solution exists" << endl;
        }
    int gcd = gcd_extend(a, b, x, y);
    if (c % gcd != 0) {
            cout << "No Solution exists" << endl;
    }
    else {
        cout << "x = " << x * (c / gcd) << ", y = " << y * (c / gcd) << endl;
    }
}</pre>
```

5.9 discrete_log

```
// Returns minimum x for which a ^x \% m = b ^x \% m, a

→ and m are coprime.

int solve(int a, int b, int m) {
    a \% m, b \% m;
    int n = sqrt(m) + 1;
    int an = 1;
    for (int i = 0; i < n; ++i)
an = (an * 1ll * a) % m;
    unordered map<int, int> vals;
    for (int \overline{q} = 0, cur = b; q \le n; ++q) {
        vals[cur] = q;
        cur = (cur * 111 * a) % m;
    for (int p = 1, cur = 1; p <= n; ++p) {
        cur = (cur * 1ll * an) % m;
        if (vals.count(cur)) {
             int ans = n * p - vals[cur];
             return ans;
    return -1;
```

5.10 factorial_mod_p

```
// O(log_p(n)) gives me n! % p for large n, p
int factmod(int n, int p) {
   vector<int> f(p);
   f[0] = 1;
   for (int i = 1; i < p; i++)
        f[i] = f[i-1] * i % p;

   int res = 1;
   while (n > 1) {
        if ((n/p) % 2)
            res = p - res;
        res = res * f[n%p] % p;
   }
   return res;
}
```

5.11 fft.

```
typedef complex<double> base;
#define PI acos(-1)
void fft(vector<base> &a, bool invert){
   int n = (int)a.size();
   for (int i = 1, j = 0; i < n; ++i){
      int bit = n >> 1;
```

```
for (; j >= bit; bit >>= 1) j -= bit;
          j += bit;
         if (i < j)swap(a[i], a[j]);
     for (int len = 2; len <= n; len <<= 1){
    double ang = 2 * PI / len * (invert ? -1 : 1);</pre>
          base wlen(cos(ang), sin(ang));
         for (int i = 0; i < n; i + len){
              base w(1);
for (int j = 0; j < len / 2; ++j){
  base u = a[i + j], v = a[i + j + len /</pre>
                   a[i + j] = u + v;
                   a[i + j + len / 2] = u - v;
                   w *= wlen:
    if (invert) for (int i = 0; i < n; ++i) a[i] /= n;
void multiply(const vector<int> &a, const vector<int>
    &b, vector<int> &res){
    vector<base> fa(a.begin(), a.end()), fb(b.begin(),
     → b.end());
     size t n = 1;
     while (n < max(a.size(), b.size())) n <<= 1;</pre>
    fa.resize(n), fb.resize(n);
fft(fa, false), fft(fb, false);
     for (size t i = 0; i < n; ++i) fa[i] *= fb[i];</pre>
     fft(fa, true); res.resize(n);
     for (size t i = 0; i < n; ++i) res[i] =

    int(fa[i].real() + 0.5);
```

5.12 find_roots

```
/*Description: Finds the real roots to a polynomial.
Usage: polyRoots({{2,-3,1}},-1e9,1e9) // solve x2-3x+2
\overrightarrow{\text{Time}} = 0
(n^2\log(1/))^*/
struct Poly {
    vector<double> a;
    double operator()(double x) const {
        double val = 0;
        for (int i = a.size(); i--;) (val *= x) +=
         → a[i];
        return val;
    void diff() {
        for (int i = 1; i < a.size(); ++i) a[i - 1] =
         \rightarrow i * a[i];
        a.pop back();
    void divroot(double x0) {
        double b = a.back(), c;
a.back() = 0;
        for (int i = a.size() - 1; i--;) c = a[i],
         \rightarrow a[i] = a[i + 1] * x0 + b, b = c;
        a.pop back();
vector<double> polyRoots(Poly p, double xmin = -1e9,
    double xmax = 1e9) {
    if (p.a.size() == 2) {
        return {-p.a[0] / p.a[1]};
    vector<double> ret;
    Poly der = p;
    der.diff();
    auto dr = polyRoots(der, xmin, xmax);
    dr.push back(xmin - 1);
    dr.push back(xmax + 1);
```

5.13 gauss_eliminition

```
const double EPS = 1e-9:
const int INF = 2; // it doesn't actually have to be

→ infinity or a big number

int gauss (vector < vector<double> > a, vector<double>
    & ans) {
     int n = (int) a.size();
     int m = (int) a[0].size() - 1;
     vector<int> where (m, -1);
     for (int col=0, row=0; col<m && row<n; ++col) {</pre>
          int sel = row;
          for (int i=row; i<n; ++i)
               if (abs (a[i][col]) > abs (a[sel][col]))
          if (abs (a[sel][col]) < EPS)
               continue;
          for (int i=col; i<=m; ++i)
    swap (a[sel][i], a[row][i]);</pre>
          where[col] = row;
          for (int i=0; i<n; ++i)</pre>
               if (i != row) {
                   double c = a[i][col] / a[row][col];
for (int j=col; j<=m; ++j)
    a[i][j] -= a[row][j] * c;</pre>
          ++row:
     ans.assign (m, 0);
     for (int i=0; i<m; ++i)</pre>
          if (where[i] != -1)
               ans[i] = a[where[i]][m] / a[where[i]][i];
     for (int i=0; i<n; ++i) {
          double sum = 0;
          for (int j=0; j<m; ++j)
    sum += ans[j] * a[i][j];</pre>
         if (abs (sum - a[i][m]) > EPS)
               return 0;
     for (int i=0; i<m; ++i)
    if (where[i] == -1)</pre>
               return INF;
     return 1;
|int gauss (vector < bitset<N> > a, int n, int m,

→ bitset<N> & ans) {
     vector<int> where (m, -1); for (int col=0, row=0; col<m, && row<n; ++col) {
          for (int i=row; i<n; ++i)
               if (a[i][col]) {
                    swap (a[i], a[row]);
                   break;
```

```
if (! a[row][col])
            continue;
        where[col] = row;
        for (int i=0; i<n; ++i)
            if (i != row && a[i][col])
                `a[i] ^= a[row];
        // The rest of implementation is the same as

→ above

//rank
const double EPS = 1E-9;
int compute rank(vector<vector<double>> A) {
    int n = A.size();
    int m = A[0].size();
    int rank = 0;
    vector<bool> row selected(n, false);
    for (int i = 0; \overline{i} < m; ++i) {
        int j;
        for (j = 0; j < n; ++j) {
            if (!row selected[j] \&\& abs(A[j][i]) > EPS)
                break:
        if (j != n) {
            ++rank;
            row selected[j] = true;
            for (int p = i + 1; p < m; ++p)
                A[j][p] /= A[j][i];
            for (int k = 0; k < n; ++k)
                if (k != j \& abs(A[k][i]) > EPS) {
                    for (int p = i + 1; p < m; ++p)
                        A[k][p] -= A[j][p] * A[k][i];
    return rank;
```

5.14 gen_all_k_combs

```
vector<int> ans:
void gen(int n, int k, int idx, bool rev) {
   if (k > n \mid | k < 0)
        return;
    if (!n)
        for (int i = 0; i < idx; ++i) {
            if (ans[i])
                 cout << i + 1;
        cout << "\n";
        return;
    ans[idx] = rev;
    gen(n-1, k-rev, idx + 1, false);
    ans[idx] = !rev;
    gen(n - 1, k - !rev, idx + 1, true);
void all combinations(int n, int k) {
    ans.resize(n);
    gen(n, k, 0, false);
```

5.15 integrate_adaptive

```
/*Description: Fast integration using an adaptive
— Simpsons rule.
Usage: double sphereVolume = quad(-1, 1, [](double x) {
return quad(-1, 1, [&](double y) {
return quad(-1, 1, [&](double z) {
```

```
5.16 lagrange interpolation
/*Description: A polynomial of degree d can be

→ uniquely identified

given its values on d + 1 unique points. O(n) to

    pre-calculate given

the first n points (x=0 \text{ to } n-1) of the polynomial.
→ Then answer each
query to interpolate the xth term in O(n). All values

→ are done modulo

mod, which needs to be a prime as we need its inverse
   modulo, Also
includes an additional helper function called

→ find degree(terms, mod).

Given at least the first d+2 points of a polynomial of

→ degree d, it finds

d in roughly O(n log d). Note, n should not exceed mod

    → due to the way modular inverse is used. In such cases, we can use

→ interpolation

without modulo in big integers and take the remainder

→ later

Time: 0 (n) */
using namespace std;
struct Lagrange {
    vector<int> terms, dp;
    int mod, n;
    Lagrange() {}
    Lagrange(const vector<int>& terms, int mod) :

    terms(terms), mod(mod) {

        n = terms.size();
        assert(n <= mod);
        int i, v, f;
for (f = 1, i = 1; i < n; i++) f = (long</pre>
        → long)f * i % mod;
        v = expo(f, mod - 2);
        vector int inv(n, v);
        for (i = n - 1; i > 0; i - -)
            inv[i - 1] = (long long)inv[i] * i % mod;
       dp[i] = (long long)dp[i] * terms[i] % mod;
    int expo(int a, int b) {
        int res = 1;
        while (b) {
            if (b \& 1) res = (long long) res * a % mod;
            a = (long long)a * a % mod;
            b >>= 1:
        return res;
```

```
int interpolate(long long x) {
        if (x < n) return terms[x] % mod;</pre>
        x %= mod;
        int i, w;
         vector<int> X(n, 1), Y(n, 1);
        for (i = 1; i < n; i++) {
    X[i] = (long long)X[i - 1] * (x - i + 1) %
             if (X[i] < 0) X[i] += mod:
        for (i = n - 2; i >= 0; i--) {
    Y[i] = (long long)Y[i + 1] * (x - i - 1) %
             if (Y[i] < 0) Y[i] += mod;
        long long res = 0;
        for (i = 0; i < n; i++) {
    w = ((long long)X[i] * Y[i] % mod) * dp[i]
                % mod:
             \inf_{\text{res}} ((n - i + 1) \& 1) w = \text{mod} - w;
        return res % mod:
vector<int> get terms(const vector<int>& terms, int
   mod, int l, int r) {
    auto lagrange = Lagrange(terms, mod);
    vector<int> res;
    for (int i = l; i <= r; i++) {
        res.push back(lagrange.interpolate(i));
    return res:
int find degree(const vector<int>& terms, int mod) {
    long long v = mod;
    int^{k} = 1, n = terms.size();
    while (v < INT MAX) {
        v *= mod:
        k++;
    int l = 1 << 30, r = l + k - 1;
    auto expected = get terms(terms, mod, l, r);
    int low = 1, high = n - 1;
    while ((low + 1) < high)
        int mid = (low + high) >> 1;
        vector<int> v(terms.begin(), terms.begin() +
        if (get terms(v, mod, l, r) == expected)
             high = mid;
        else_
             low = mid;
    for (int d = low; d <= high; d++) {</pre>
        vector<int> v(terms.begin(), terms.begin() +
        if (get terms(v, mod, l, r) == expected)
         \rightarrow return d - 1;
    return -1;
int main() {
    const int mod = 10000000007;
    vector<int> terms = vector<int>{0, 1, 5, 14, 30};
    auto lagrange = Lagrange(terms, mod);
    assert(lagrange.interpolate(5) == 55);
    assert(lagrange.interpolate(6) == 91);
    assert(lagrange.interpolate(1 << 30) == 300663155);
    assert(lagrange.interpolate(1LL << 60) ==
     → 717860166):
    assert(find degree(terms, mod) == 3);
```

```
terms.pop back();
    assert(find degree(terms, mod) == -1);
                                                               void fft(vector<int> &a, bool inv){
                                                                   int n = a.size();
for (int i = 1; i < n - 1; ++i) if (i < rev[i])</pre>
    return 0;
                                                                    → swap(a[i], a[rev[i]]);
                                                                   for (int len = 2; len <= n; len <<= 1) {
5.17 matrix_expo
                                                                        int wlen = inv ? root 1 : root;
11 \mod = (1e9) + 7;
                                                                        for (int i = len; i < k; i <<= 1) wlen = 1ll *
struct Matrix{

→ wlen * wlen % mod;

    int row, col;
                                                                        for (int st = 0; st < n; st += len){
    vector<vector<ll>> mat;
                                                                            int w = 1;
for (int j = 0; j < len / 2; j++){</pre>
    Matrix(int x, int y){
         row=x;
                                                                                 int ev = a[st + j];
         col=y;
                                                                                 int od = 111 * a[st + j + len / 2] * w
         mat.assign(row, vector<ll>(col,0));
                                                                                 a[st + j]' = ev + od < mod ? ev + od :
    Matrix operator *(Matrix &other){
                                                                                     ev + od - mod;
         assert(col==other.row);
                                                                                 a[st + j + len / 2] = ev - od >= 0?
         Matrix product(row,other.col);
                                                                                 \rightarrow ev - od : ev - od + mod;
         for(int i=0;i<row;i++){</pre>
                                                                                 w = 111 * w * wlen % mod;
             for(int j=0; j<col; j++){
                  for(int k=0;k<other.col;k++){</pre>
                      product.mat[i][k]=(product.mat[i][
|
                           k]+(mat[i][j]*other.mat[i][k])
                                                                   if (inv){
                       ⇒ %mod)%mod:
                                                                        int n 1 = bigmod(n, mod - 2, mod);
                                                                        for (int \&x : a) x = 111 * x * n 1 % mod;
         return product;
                                                               vector<int> mul(vector<int> &a, vector<int> &b){
                                                                   int n = a.size(), m = b.size(), sz = 1;
while (sz < n + m - 1) sz <<= 1;</pre>
};
Matrix expo(Matrix &m, ll n){
                                                                   vector<int> x(sz), y(sz), z(sz);
                                                                   for (int i = 0; i < sz; ++i){
    x[i] = i < n ? a[i] : 0;
    y[i] = i < m ? b[i] : 0;</pre>
    assert(m.row==m.col):
    Matrix ret(m.row,m.col);
    for(int i=0;i<m.row;i++) ret.mat[i][i]=1;</pre>
    while(n){
         if(n\&1) ret=ret*m;
                                                                   pre(sz);fft(x, 0);fft(y, 0);
         n/=2;
                                                                   for (int i = 0; i < sz; ++i) z[i] = 1ll * x[i] *
         m=m*m;
                                                                    \rightarrow y[i] % mod;
                                                                   fft(z, 1); z.resize(n + m - 1);
    return ret;
                                                                   return z;
5.18 next_lexicographical_k_comb
                                                               5.20 seg_sieve
bool next combination(vector<int>& a, int n) {
    int k = (int)a.size();
    for (int i = k - 1; i \ge 0; i - -) {
                                                               This code was for 1 <= a <= b <= 2^31-1
                                                               Change variable types appropriately.
         if (a[i] < n - k + i + 1) {
             a[i]++;
                                                               bool notPrime[ ? ];
             for (int j = i + 1; j < k; j++)
a[j] = a[j - 1] + 1;
                                                               void segmented sieve(int a, int b)
             return true;
                                                                   int p, <u>f</u>;
                                                                   mem(notPrime, 0);
                                                                   for (int i = 0; i < tot prime; i++)
    return false:
                                                                        p = prime[i];
5.19 ntt
                                                                        if (a \% p == 0)
                                                                            f = a;
const int mod = 998244353;
const int root = 15311432;
                                                                            f = (a - (a \% p) + p);
const int k = 1 << 23;
                                                                        f = max(p * p, f);
int root 1;
                                                                        for (unsigned j = f; j \le b; j += p)
vector<int> rev;
                                                                            notPrime[j - a] = true;
void pre(int sz){
    root 1 = bigmod(root, mod - 2, mod);
                                                                   if (a == 1)
    if (rev.size() == sz) return;
                                                                        notPrime[0] = 1;
    rev.resize(sz);
```

5.21 stirling

that no box is empty.

|int stirling2(int n, int k)

rev[0] = 0;

int lg_n = __builtin_ctz(sz);

 \rightarrow rev[i]=(rev[i>>1]>>1)|((i&1)<<(lq n-1));

for (int i = 1; i < sz; ++i)

```
if(n < k)
return 0:
if(k == 1)
return 1
if(dp[n][k] == dp[n][k])
return dp[n][k];
return dp[n][k] = stirling2(n-1,k-1) +

    stirling2(n-1,k)*k;

is empty
int stirling1(int n, int k)
dp[n][k] = stirling1(n-1,k-1) + stirling(n-1,k)*n-1;
5.22 stirling_number_of_the_second_kind
    / k! * sum (-1)^i nCr(k, i) * (k - i) ^ n
11 f(int n, int k) {
    ll res = 0;
    for (int i = 0; i < k; ++i) {
        if (i \& 1) res = (res - nCr(k, i) * bp(k - i,
         \rightarrow n, mod) % mod + mod) % mod;
        else res = (res + nCr(k, i) * bp(k - i, n,

→ mod) % mod) % mod;

   if (res < 0) res += mod;
    return res * ifac[k] % mod;
5.23 sum_of_totient
using namespace __gnu_pbds;
const int N = 3e5 + 9, mod = 998244353;
template <const int32 t MOD>
struct modint -
    int32 t value:
    modint() = default;
    modint(int32_t value) : value(value) {}
    inline modint<MOD> operator+(modint<MOD> other)
        int32 t c = this->value + other.value;
        return modint<MOD>(c >= MOD ? c - MOD : c);
    inline modint<MOD> operator-(modint<MOD> other)
        int32 t c = this->value - other.value;
        return modint<MOD>(c < 0 ? c + MOD : c):
    inline modint<MOD> operator*(modint<MOD> other)
       const {
        int32 t c = (int64 t)this->value * other.value
        return modint<MOD>(c < 0 ? c + MOD : c);
    inline modint<MOD> &operator+=(modint<MOD> other) {
        this->value += other.value:
        if (this->value >= MOD)
            this->value -= MOD;
        return *this;
    inline modint<MOD> &operator-=(modint<MOD> other) {
        this->value -= other.value;
        if (this->value < 0)</pre>
            this->value += MOD;
        return *this;
   inline modint<MOD> &operator*=(modint<MOD> other) {
        this->value = (int64 t)this->value *

→ other.value %
```

```
if (this->value < 0) this->value += MOD;
        return *this;
    inline modint<MOD> operator-() const { return
        modint<MOD>(this->value ? MOD - this->value :
       0); }
    modint<MOD> pow(uint64_t k) const {
        modint<MOD> x =
                         *this.
                     v = 1;
        for (; k; k >>= 1) {
    if (k & 1) y *= x;
            x *= x;
        return v;
    modint<MOD> inv() const { return pow(MOD - 2); }
    → // MOD must be a prime
    inline modint<MOD> operator/(modint<MOD> other)
        const { return *this * other.inv(): }
    inline modint<MOD> operator/=(modint<MOD> other) {
        return *this *= other.inv(); }
    inline bool operator==(modint<MOD> other) const {
       return value == other.value; }
    inline bool operator!=(modint<MOD> other) const {

    return value != other.value; }

    inline bool operator<(modint<MOD> other) const {

→ return value < other.value; }
</p>
    inline bool operator>(modint<MOD> other) const {

→ return value > other.value: }

template <int32 t MOD>
modint<MOD> operator*(int64 t value, modint<MOD> n) {

→ return modint<MOD>(value) * n: }

template <int32 t MOD>
modint<MOD> operator*(int32 t value, modint<MOD> n) {

→ return modint<MOD>(value % MOD) * n; }

template <int32 t MOD>
istream & operator>>(istream &in, modint<MOD> &n) {

    return in >> n.value; }

template <int32 t MOD>
ostream & operator << (ostream & out, modint < MOD> n) {

    return out << n.value; }
</pre>
using mint = modint<mod>;
namespace Dirichlet {
// solution for f(x) = phi(x)
const int T = 1e7 + 9;
long long phi[T];
gp hash table<long long, mint> mp;
mint dp[T], inv;
int sz, spf[T], prime[T];
void init() {
    memset(spf, 0, sizeof spf);
    phi[1] = 1;
    SZ = 0;
    for (int i = 2; i < T; i++) {
        if (spf[i] == 0) phi[i] = i - 1, spf[i] = i,
         \rightarrow prime[sz++] = i:
        for (int j = 0; j < sz && i * prime[j] < T &&
            prime[j] <= spf[i]; j++) {
spf[i * prime[j]] = prime[j];</pre>
            if (i % prime[i] == 0)
                 phi[i * prime[j]] = phi[i] * prime[j];
            else
                 phi[i * prime[j]] = phi[i] * (prime[j]

→ - 1);
    dp[0] = 0;
```

```
for (int i = 1; i < T; i++) dp[i] = dp[i - 1] +
     → phi[i] % mod;
    inv = 1; // q(1)
|mint p c(long long n) {
    if (n \% 2 == 0) return n / 2 \% mod * ((n + 1) \%
     → mod) % mod;
    return (n + 1) / 2 % mod * (n % mod) % mod;
mint p q(long long n) {
    return n % mod;
mint solve(long long x)
    if (x < T) return dp[x];</pre>
    if (mp.find(x) != mp.end()) return mp[x];
    mint ans = 0;
    for (long long i = 2, last; i <= x; i = last + 1)
    last = x / (x / i);</pre>
         ans += solve(x / i) * (p g(last) - p_g(i - 1));
    ans = p c(x) - ans;
    ans /=\overline{i}nv;
    return mp[x] = ans;
   // namespace Dirichlet
5.24 totient
int phi(int n) {
    int result = n;
for (int i = 2; i * i <= n; i++) {</pre>
         if (n % i == 0) {
             while (n % i == 0)
                  n \neq i;
             result -= result / i;
    if (n > 1)
         `result -= result / n;
    return result;
vector<int> phi(n + 1):
     phi[0] = 0;
     phi[1] = 1;
     for (int i = 2; i <= n; i++)
         phi[i] = i;
    for (int i = 2; i <= n; i++) {
   if (phi[i] == i) {</pre>
             for (int j = i; j <= n; j += i)
    phi[j] -= phi[j] / i;</pre>
    }
5.25 xor basis
int basis[d]; // basis[i] keeps the mask of the vector
→ whose f value is i
int sz;
void insertVector(int mask) {
 for (int i = 0; i < d; i++) {
  if ((mask \& 1 << i) == 0) continue;
  if (!basis[i]) { // If there is no basis vector with
       the i'th bit set, then insert this vector into

    the basis
basis[i] = mask;

   return:
  mask ^= basis[i]; // Otherwise subtract the basis

    vector from this vector
```

```
6 Misc
6.1 check
```

6.2 debug

6.3 flags

```
g++ -std=c++17 -Wall -Wextra -pedantic -Wshadow
-Wformat=2 -Wfloat-equal -Wconversion -Wlogical-op
-Wshift-overflow=2 -Wduplicated-cond -Wcast-qual
-Wcast-align -D_GLIBCXX_DEBUG
-D_GLIBCXX_DEBUG_PEDANTIC -D_FORTIFY_SOURCE=2
-fsanitize=address -fsanitize=undefined
-fno-sanitize-recover -fstack-protector -o $1
-Wno-unused-result $1.cpp
```

6.4 interval container

```
/* Description: Add and remove intervals from a set of
    disjoint intervals. Will merge the added interval
    with any overlapping intervals in
the set when adding. Intervals are [inclusive,
        exclusive).
Time: 0 (log N)
*/
setpii>::iterator addInterval(setpii>& is, int L,
    int R) {
    if (L == R) return is.end();
    auto it = is.lower bound({L, R}), before = it;
    while (it != is.end() & it->first <= R) {
        R = max(R, it->second);
        before = it = is.erase(it);
```

```
if (it != is.begin() && (--it)->second >= L) {
        L = min(L, it->first);
R = max(R, it->second); is.erase(it);
   return is.insert(before, {L, R});
void removeInterval(set<pii>& is, int L, int R) {
    if (L == R) return;
   auto it = addInterval(is, L, R);
    auto r2 = it->second;
    if (it->first == L)
        is.erase(it);
         (int&)it->second = L;
   if (R != r2) is.emplace(R, r2);
```

6.5 interval cover

```
template <class T>
vi cover(pair<T, T> G, vector<pair<T, T>> I) {
    vi S(sz(I)), R;
iota(all(S), 0);
sort(all(S), [&](int a, int b) { return I[a] <
    int at = 0;
    while (cur < G.second) { // (A)
   pair<T, int> mx = make_pair(cur, -1);
         while (at < sz(I) && I[S[at]].first <= cur) {
    mx = max(mx, make_pair(I[S[at]].second,</pre>
              if (mx.second == -1) return {};
          cur = mx.first;
         R.push back(mx.second);
    return R;
```

6.6 pragma

```
// Pragmas
#pragma comment(linker, "/stack:200000000")
#pragma GCC optimize("03,unroll-loops")
#pragma GCC target("avx,avx2,fma")
```

6.7 random

```
// shuffle(v.begin(), v.end(),
→ default random engine(rnd(1, 1000)));
mt19937 rng(chrono::steady clock::now().time since epo
ch().count());
ll rnd(ll l, ll r)
    return uniform int distribution<ll>(l, r) (rng);
```

6.8 vimre

```
imap jk <Esc>
set nú
set mouse=a
set autoindent
set tabstop=4
   shiftwidth=4
   smartindent
   relativenumber
    laststatus=2
   hlsearch
let mapleader =
nnoremap <leader>s :w<Enter>
nnoremap <leader>y ggVG"+y<CR>
syntax on
vnoremap <leader>/ :s!^!//!<CR> :noh <CR>
```

```
vnoremap <leader>u :s!^//!!<CR>
nnoremap <leader>/ :s!^!//!<CR> :noh <CR>
nnoremap <leader>u :s!^//!!<CR>
```

7 Notes

7.1 Counting

1. Lucas Theorem

For non-negative integers m and n and a prime p, the following congruence relation holds: :

$$\begin{pmatrix} m \\ n \end{pmatrix} \equiv \prod_{i=0}^{k} \begin{pmatrix} m_i \\ n_i^i \end{pmatrix} \pmod{p}, \tag{1}$$

where:

$$m = m_k p^k + m_{k-1} p^{k-1} + \dots + m_1 p + m_0$$

and:

$$n = n_k p^k + n_{k-1} p^{k-1} + \dots + n_1 p + n_0$$

are the base p expansions of m and n respectively. This uses the convention that $\binom{m}{n} = 0$ if $m \le n$.

2. Stirling Numbers of the first kind

S(n,k) counts the number of permutations of n elements with k disioint cycles.

$$S(n,k) = (n-1) \cdot S(n-1,k) + S(n-1,k-1)$$
 (2)

where, S(0,0) = 1, S(n,0) = S(0,n) = 0

$$\sum_{k=0}^{n} S(n,k) = n! \tag{3}$$

3. Stirling Numbers of the second kind

 $S(n,k) \cdot k! = \text{number of ways to color } n \text{ nodes using colors from 1 to } k \mid \text{With side lengths } a,b,c,d, \text{ diagonals } e,f, \text{ diagonals angle } \theta, \text{ area } A$ such that each color is used at least once.

An r-associated Stirling number of the second kind is the number of ways to partition a set of n objects into k subsets, with each subset containing at least r elements. It is denoted by $S_r(n,k)$ and obeys the recurrence relation.

$$S_r(n+1,k) = kS_r(n,k) + \binom{n}{r-1}S_r(n-r+1,k-1)$$
 (4)

4. Bell Numbers

Counts the number of partitions of a set.

$$B_{n+1} = \sum_{k=0}^{n} \left(\frac{n}{k}\right) \cdot B_k \tag{5}$$

 $B_n = \sum_{k=0}^n S(n,k)$, where S(n,k) is stirling number of second kind.

5. Some identities

Vandermonde's Identify:
$$\sum_{k=0}^{r} {m \choose k} {n \choose r-k} = {m+n \choose r}$$

Hockey-Stick Identify:
$$n, r \in N, n > r, \sum_{i=r}^{n} \begin{pmatrix} i \\ r \end{pmatrix} = \begin{pmatrix} n+1 \\ r+1 \end{pmatrix}$$

Involutions: permutations such that p^2 = identity permutation. $a_0 = a_1 = 1$ and $a_n = a_{n-1} + (n-1)a_{n-2}$ for n > 1.

7.2 Fibonacci

Let A, B and n be integer numbers.

$$k = A - B$$

$$F_A F_B = F_{k+1} F_A^2 + F_k F_A F_{A-1} \tag{6}$$

$$\sum_{i=0}^{n} F_i^2 = F_{n+1} F_n \tag{7}$$

$$\sum_{i=0}^{n} F_i F_{i+1} = F_{n+1}^2 - (-1)^n \tag{8}$$

$$\sum_{i=0}^{n} F_i F_{i-1} = \sum_{i=0}^{n-1} F_i F_{i+1}$$
 (9)

$$GCD(F_m, F_n) = F_{GCD(m,n)}$$
(10)

$$\sum_{0 \le k \le n} \binom{n-k}{k} = Fib_{n+1} \tag{11}$$

$$\gcd(F_n, F_{n+1}) = \gcd(F_n, F_{n+2}) = \gcd(F_{n+1}, F_{n+2}) = 1$$
 (12)

7.3 Notes

7.4 Geometry

7.4.1 Triangles

Circumradius: $R=\frac{abc}{4A}$, Inradius: $r=\frac{A}{s}$ Length of median (divides triangle into two equal-area triangles): (2) $m_a = \frac{1}{2}\sqrt{2b^2 + 2c^2 - a^2}$

Length of bisector (divides angles in two):
$$s_a = \sqrt{bc \left[1 - \left(\frac{a}{b+c}\right)^2\right]}$$

Law of tangents:
$$\frac{a+b}{a-b} = \frac{\tan \frac{\alpha+\beta}{2}}{\tan \frac{\alpha-\beta}{2}}$$

7.4.2 Quadrilaterals

and magic flux $F = b^2 + d^2 - a^2 - c^2$:

$$4A = 2ef \cdot \sin \theta = F \tan \theta = \sqrt{4e^2 f^2 - F^2}$$

For cyclic quadrilaterals the sum of opposite angles is 180° , ef =ac+bd, and $A=\sqrt{(p-a)(p-b)(p-c)(p-d)}$.

(4) 7.4.3 Spherical coordinates

$$x = r \sin \theta \cos \phi \qquad r = \sqrt{x^2 + y^2 + z^2}$$

$$y = r \sin \theta \sin \phi \qquad \theta = a\cos(z/\sqrt{x^2 + y^2 + z^2})$$

$$z = r \cos \theta \qquad \phi = a\tan(2(y, x))$$

(5)
$$1^{2} + 2^{2} + 3^{2} + \dots + n^{2} = \frac{n(2n+1)(n+1)}{6}$$

$$1^{3} + 2^{3} + 3^{3} + \dots + n^{3} = \frac{n^{2}(n+1)^{2}}{4}$$

$$1^{4} + 2^{4} + 3^{4} + \dots + n^{4} = \frac{n(n+1)(2n+1)(3n^{2} + 3n - 1)}{30}$$

$$\sum_{i=1}^{n} i^{m} = \frac{1}{m+1} \left[(n+1)^{m+1} - 1 - \sum_{i=1}^{n} \left((i+1)^{m+1} - i^{m+1} - (m+1)i^{m} \right) \right]$$

$$\sum_{i=1}^{n-1} i^{m} = \frac{1}{m+1} \sum_{k=0}^{m} {m+1 \choose k} B_{k} n^{m+1-k}$$
ion.
$$\sum_{k=0}^{n} k x^{k} = (x - (n+1)x^{n+1} + nx^{n+2})/(x-1)^{2}$$

7.6 Series

$$e^{x} = 1 + x + \frac{x^{2}}{2!} + \frac{x^{3}}{3!} + \dots, (-\infty < x < \infty)$$

$$\ln(1+x) = x - \frac{x^{2}}{2} + \frac{x^{3}}{3} - \frac{x^{4}}{4} + \dots, (-1 < x \le 1)$$

$$\sqrt{1+x} = 1 + \frac{x}{2} - \frac{x^{2}}{8} + \frac{2x^{3}}{32} - \frac{5x^{4}}{128} + \dots, (-1 \le x \le 1)$$

$$\sin x = x - \frac{x^{3}}{3!} + \frac{x^{5}}{5!} - \frac{x^{7}}{7!} + \dots, (-\infty < x < \infty)$$

$$\cos x = 1 - \frac{x^{2}}{2!} + \frac{x^{4}}{4!} - \frac{x^{6}}{6!} + \dots, (-\infty < x < \infty)$$

$$(x+a)^{-n} = \sum_{k=0}^{\infty} (-1)^{k} \binom{n+k-1}{k} x^{k} a^{-n-k}$$

7.7 Pythagorean Triples

The Pythagorean triples are uniquely generated by

$$a = k \cdot (m^2 - n^2), b = k \cdot (2mn), c = k \cdot (m^2 + n^2),$$

with m > n > 0, k > 0, $m \perp n$, and either m or n even.

7.8 Number Theory

7.8.1 Primes

hashing use 970592641 (31-bit number), 31443539979727 (45-bit), 1 = ab - a - b. 3006703054056749 (52-bit). There are 78498 primes less than 7.8.12 Fermat's two-squares theorem

Primitive roots exist modulo any prime power p^a , except for $p = 2, a > 1 \pmod{4}$. A product of two sums of two squares is a sum of two 2, and there are $\phi(\phi(p^a))$ many. For p=2,a>2, the group \mathbb{Z}_{2a}^{\times} is squares. Thus, n is a sum of two squares iff every prime of form greater than the previous element. k j:s s.t. $\pi(j) > \pi(j+1)$, k+1 j:s s.t. instead isomorphic to $\mathbb{Z}_2 \times \mathbb{Z}_{2a-2}$.

7.8.2 Estimates

 $\sum_{d|n} d = O(n \log \log n).$

The number of divisors of n is at most around 100 for n < 5e4, 500 for n < 1e7, 2000 for n < 1e10, 200 000 for n < 1e19.

7.8.3 Perfect numbers

n > 1 is called perfect if it equals sum of its proper divisors and 1. Even *n* is perfect iff $n = 2^{p-1}(2^p - 1)$ and $2^p - 1$ is prime (Mersenne's). No **7.9.2** Cycles odd perfect numbers are yet found.

7.8.4 Carmichael numbers

A positive composite *n* is a Carmichael number $(a^{n-1} \equiv 1 \pmod{n})$ for all a: gcd(a, n) = 1), iff n is square-free, and for all prime divisors p of n. p-1 divides n-1.

7.8.5 Mobius function

 $\mu(1) = 1$. $\mu(n) = 0$, if *n* is not squarefree. $\mu(n) = (-1)^s$, if *n* is the product of s distinct primes. Let f, \overline{F} be functions on positive integers. If for 7.9.3 Derangements all $n \in N$, $F(n) = \sum_{d|n} f(d)$, then $f(n) = \sum_{d|n} \mu(d) F(\frac{n}{d})$, and vice versa. Permutations of a set such that none of the elements appear in their $\phi(n) = \sum_{d \mid n} \mu(d) \frac{n}{d}$. $\sum_{d \mid n} \mu(d) = 1$.

If f is multiplicative, then $\sum_{d|n} \mu(d) f(d) = \prod_{p|n} (1 - f(p)),$ $\begin{array}{ll} \sum_{d\mid n} \mu(d)^2 f(d) = \prod_{p\mid n} (1+f(p)). \\ \textbf{7.8.6} \quad \textbf{Legendre symbol} \end{array}$

If p is an odd prime, $a \in \mathbb{Z}$, then $\left(\frac{a}{p}\right)$ equals 0, if p|a; 1 if a is a quadratic residue modulo p; and -1 otherwise. Euler's criterion: of X up to symmetry equals $\left(\frac{a}{p}\right) = a^{\left(\frac{p-1}{2}\right)} \pmod{p}.$

7.8.7 Jacobi symbol

If $n = p_1^{a_1} \cdots p_k^{a_k}$ is odd, then $\left(\frac{a}{n}\right) = \prod_{i=1}^k \left(\frac{a}{n_i}\right)^{k_i}$.

7.8.8 Primitive roots

If the order of g modulo m (min n > 0: $g^n \equiv 1 \pmod{m}$) is $\phi(m)$, then rotational symmetry using $G = \mathbb{Z}_n$ to get g is called a primitive root. If Z_m has a primitive root, then it has $\phi(\phi(m))$ distinct primitive roots. Z_m has a primitive root iff m is one of 2, 4, p^k , $2p^k$, where p is an odd prime. If Z_m has a primitive root g, then for all a coprime to m, there exists unique integer $i = \text{ind}_g(a)$ modulo $\phi(m)$, such that $g^i \equiv a \pmod{m}$. ind_g(a) has logarithm-like properties: ind(1) = 0, ind(ab) = ind(a) + ind(b).

If p is prime and a is not divisible by p, then congruence $x^n \equiv a \pmod{p}$ has $\gcd(n, p-1)$ solutions if $a^{(p-1)/\gcd(n, p-1)} \equiv 1 \pmod{p}$, and no solutions otherwise. (Proof sketch: let g be a primitive root, and $g^i \equiv a \pmod{p}, g^u \equiv x \pmod{p}. x^n \equiv a \pmod{p} \text{ iff } g^{nu} \equiv g^i \pmod{p}$ iff $nu \equiv i \pmod{p}$.)

7.8.9 Discrete logarithm problem

Find x from $a^x \equiv b \pmod{m}$. Can be solved in $O(\sqrt{m})$ time and space with a meet-in-the-middle trick. Let $n = \lceil \sqrt{m} \rceil$, and x = ny - z. Equation becomes $a^{ny} \equiv ba^z \pmod{m}$. Precompute all values that the RHS can take for $z = 0, 1, \dots, n-1$, and brute force y on the LHS, each time checking whether there's a corresponding value for RHS.

7.8.10 Pythagorean triples

Integer solutions of $x^2 + y^2 = z^2$ All relatively prime triples are given by: $x = 2mn, y = m^2 - n^2, z = m^2 + n^2$ where $m > n, \gcd(m, n) = 1$ and $m \not\equiv n \pmod{2}$. All other triples are multiples of these. Equation $x^2 + y^2 = 2z^2$ is equivalent to $(\frac{x+y}{2})^2 + (\frac{x-y}{2})^2 = z^2$.

7.8.11 Postage stamps/McNuggets problem

Let a, b be relatively-prime integers. There are exactly $\frac{1}{2}(a-1)(b-1)$ p = 962592769 is such that $2^{21} \mid p-1$, which may be useful. For numbers not of form ax + by $(x, y \ge 0)$, and the largest is (a-1)(b-1)

Odd prime p can be represented as a sum of two squares iff $p \equiv |7.11.2|$ Eulerian numbers p = 4k + 3 occurs an even number of times in *n*'s factorization.

7.9 Permutations

).1	Factorial	
	n	$1\ 2\ 3\ 4\ 5\ 6\ 7\ 8\ 9\ 10$
	n!	$1\ 2\ 6\ 24\ 120\ 720\ 5040\ 40320\ 362880\ 3628800$
	n	11 12 13 14 15 16 17
	n!	4.0e7 4.8e8 6.2e9 8.7e10 1.3e12 2.1e13 3.6e14
	n	20 25 30 40 50 100 150 171
	n	2018 2025 3032 8047 3064 90157 60262 SDRI MAY

Let $g_S(n)$ be the number of *n*-permutations whose cycle lengths all belong to the set S. Then

$$\sum_{n=0}^{\infty} g_S(n) \frac{x^n}{n!} = \exp\left(\sum_{n \in S} \frac{x^n}{n}\right)$$

original position.

$$D(n) = (n-1)(D(n-1) + D(n-2)) = nD(n-1) + (-1)^n = \left| \frac{n!}{e} \right|$$

$$\frac{1}{|G|} \sum_{g \in G} |X^g|,$$

where X^g are the elements fixed by g(g.x = x).

If f(n) counts "configurations" (of some sort) of length n, we can ignore

$$g(n) = \frac{1}{n} \sum_{k=0}^{n-1} f(\gcd(n,k)) = \frac{1}{n} \sum_{k|n} f(k)\phi(n/k)$$

7.10 Partitions and subsets

7.10.1 Partition function

Number of ways of writing n as a sum of positive integers, disregarding the order of the summands.

$$p(0) = 1, \ p(n) = \sum_{k \in \mathbb{Z} \setminus \{0\}} (-1)^{k+1} p(n - k(3k - 1)/2)$$

$$p(n) \sim 0.145/n \cdot \exp(2.56\sqrt{n})$$

7.11 General purpose numbers

7.11.1 Stirling numbers of the first kind

Number of permutations on n items with k cycles.

$$c(n,k) = c(n-1,k-1) + (n-1)c(n-1,k), \ c(0,0) = 1$$

$$\sum_{k=0}^{n} c(n,k)x^{k} = x(x+1)...(x+n-1)$$

c(8,k) = 8,0,5040,13068,13132,6769,1960,322,28,1c(n,2) = 0,0,1,3,11,50,274,1764,13068,109584,...

Number of permutations $\pi \in S_n$ in which exactly k elements are $\pi(i) \ge i, k \text{ i:s s.t. } \pi(i) > i.$

$$E(n,k) = (n-k)E(n-1,k-1) + (k+1)E(n-1,k)$$

$$E(n,0) = E(n,n-1) = 1$$

$$E(n,k) = \sum_{j=0}^{k} (-1)^{j} \binom{n+1}{j} (k+1-j)^{n}$$

7.11.3 Stirling numbers of the second kind

Partitions of *n* distinct elements into exactly *k* groups.

$$S(n,k) = S(n-1,k-1) + kS(n-1,k)$$

$$S(n,1) = S(n,n) = 1$$

$$S(n,k) = \frac{1}{k!} \sum_{j=0}^{k} (-1)^{k-j} \binom{k}{j} j^{n}$$

7.11.4 Bell numbers

Total number of partitions of n distinct elements. B(n) =

$$B(p^m + n) \equiv mB(n) + B(n+1) \pmod{p}$$

7.11.5 Bernoulli numbers

$$\sum_{j=0}^{m} \binom{m+1}{j} B_j = 0. \quad B_0 = 1, B_1 = -\frac{1}{2}. \ B_n = 0, \text{ for all odd } n \neq 1.$$

7.11.6 Catalan numbers

$$C_n = \frac{1}{n+1} {2n \choose n} = {2n \choose n} - {2n \choose n+1} = \frac{(2n)!}{(n+1)!n!}$$

$$C_0 = 1, \ C_{n+1} = \frac{2(2n+1)}{n+2}C_n, \ C_{n+1} = \sum C_i C_{n-i}$$

 $C_n = 1, 1, 2, 5, 14, 42, 132, 429, 1430, 4862, 16796, 58786, \dots$

- sub-diagonal monotone paths in an $n \times n$ grid.
- strings with *n* pairs of parenthesis, correctly nested.
- binary trees with with n + 1 leaves (0 or 2 children).
- ordered trees with n+1 vertices.
- ways a convex polygon with n+2 sides can be cut into triangle by connecting vertices with straight lines.
- permutations of [n] with no 3-term increasing subseq.

7.12 Inequalities

7.12.1 Titu's Lemma

For positive reals $a_1, a_2, ..., a_n$ and $b_1, b_2, ..., b_n$,

$$\frac{{a_1}^2}{b_1} + \frac{{a_2}^2}{b_2} + \ldots + \frac{{a_n}^2}{b_n} \ge \frac{a_1 + a_2 + \ldots + a_n^2}{b_1 + b_2 + \ldots + b_n}$$

Equality holds if and only if $a_i = kb_i$ for a non-zero real constant k. **7.13** Games

7.13.1 Grundy numbers

For a two-player, normal-play (last to move wins) game on a graph Are you sure your algorithm works? (V,E): $G(x) = \max(\{G(y):(x,y)\in E\})$, where $\max(S) = \min\{n\geq 0:n\not\in S\}$. x is losing iff G(x)=0. S. x is losing iff G(x) = 0.

- of a position is xor of grundy numbers of positions in summed Are you clearing all data structures between test
- games.

 Player chooses a non-empty subset of games (possibly, all) and Any undefined behavior (array out of bounds)?
- Player chooses a proper subset of games (not empty and not all), Return vs continue vs break?

 and makes moves in all chosen ones. A position is losing iff Are you sure the STL functions you use work as you grundy numbers of all games are equal.
- Player must move in all games, and loses if can't move in some Create some test cases to run your algorithm on. game A position is losing if any of the games is in a losing posi- Go through the algorithm for a simple case. tion.

7.13.3 Misère Nim

A position with pile sizes $a_1, a_2, ..., a_n \ge 1$, not all equal to 1, is losing Go for a small walk, e.g. to the toilet. iff $a_1 \oplus a_2 \oplus \cdots \oplus a_n = 0$ (like in normal nim.) A position with n piles of Rewrite your solution from the start or let a teammate size 1 is losing iff n is odd.

7.14 NumberTheory

$$\sum_{k=1}^{n} \frac{1}{\gcd(k,n)} = \sum_{d|n} \frac{1}{d} \cdot \phi\left(\frac{n}{d}\right) = \frac{1}{n} \sum_{d|n} d \cdot \phi(d)$$

$$\sum_{k=1}^{n} \frac{k}{\gcd(k,n)} = \frac{n}{2} \cdot \frac{1}{n} \cdot \sum_{d \mid n} d \cdot \phi(d)$$

$$\sum_{k=1}^{n} \frac{n}{\gcd(k,n)} = 2 * \sum_{k=1}^{n} \frac{k}{\gcd(k,n)} - 1, \text{ for, } n > 1$$

$$\sum_{i=1}^{n} \sum_{j=1}^{n} [\gcd(i,j) = 1] = \sum_{d=1}^{n} \mu(d) \left\lfloor \frac{n}{d} \right\rfloor^{2}$$

$$\sum_{i=1}^{n} \sum_{j=1}^{n} \gcd(i,j) = \sum_{d=1}^{n} \phi(d) \left\lfloor \frac{n}{d} \right\rfloor^{2}$$

$$\sum_{i=1}^{n} \sum_{j=1}^{n} i \cdot j[\gcd(i,j) = 1] = \sum_{i=1}^{n} \phi(i)i^{2}$$

$$\sum_{i=1}^{n} \sum_{j=1}^{n} \operatorname{lcm}(i,j) = \sum_{l=1}^{n} \left(\frac{\left(1 + \left\lfloor \frac{n}{l} \right\rfloor\right) \left(\left\lfloor \frac{n}{l} \right\rfloor\right)}{2} \right)^{2} = \sum_{d|l} \mu(d) l d$$
 (19)

7.15 troubleshoot

Write down most of your thoughts, even if you re not whether they re useful. Give your variables (and files) meaningful names. Stay organized and don<u>⊓</u>t leave papers all over the You should know what your code is doing ...

Write a few simple test cases **if** sample is not enough.

Are time limits close? If so, generate max cases. Is the memory usage fine? Could anything overflow?

Remove débug output. Make sure to submit the right file.

Print your solution! Print debug output as well. Read the full problem statement again.

Have you understood the problem correctly?

Did you consider corner cases (ex. n=1)? 7.13.2 Sums of games
Player chooses a game and makes a move in it Grundy number Is your output format correct? (including whitespace)

makes moves in all of them A position is losing iff each game is in a losing position.

Any overflows or NaNs (or shifting ll by >=64 bits)? Confusing N and M, i and j, etc.?

Confusing ++i and i++?

Add some assertions, maybe resubmit.

Go through this list again.

Explain your algorithm to a teammate. Ask the teammate to look at your code.

→ do it.

Geometry:

Work with ints if possible.
(13) Correctly account for numbers close to (but not) zero.

for functions like acos make sure absolute val of input is not

(14) (slightly) greater than one. Correctly deal with vertices that are collinear,

coplanar (in 3D), etc. Subtracting a point from every other (but not itself)?

Have you tested all corner cases locally? (16) Any uninitialized variables? Aré you reading or writing outside the range of any

Any assertions that might fail? (17) Any possible division by 0? (mod 0 for example) Any possible infinite récursion?

(18) Invalidated pointers or iterators? Are you using too much memory? Debug with resubmits (e.g. remapped signals, see → Various). Time limit exceeded: (19) Do you have any possible infinite loops?
What's your complexity? Large TL does not mean that → something simple (like NlogN) isnut intended. Are you copying a lot of unnecessary data? (References) Avoid vector, map. (use arrays/unordered map) How big is the input and output? (consider FastIO) What do your teammates think about your algorithm? Calling count() on multiset? Memory limit exceeded: What is the max amount of memory your algorithm should Are you clearing all data structures between test → cases? If using pointers try BumpAllocator.

8 String

8.1 aho_corasick

using namespace std;

```
const int N = ?;
const int A = ?;
struct AC {
  int nd, pt;
  int next[N][A], link[N], out link[N], cnt[N], ans[N];
  vector <int> ed[N], out[N];
  AC(): nd(0), pt(0) { node(); }
  int node() {
    memset(next[nd], 0, sizeof next[nd]);
link[nd] = out link[nd] = cnt[nd] = 0;
ed[nd].clear(), out[nd].clear();
     return nd++:
  void clear() {
    nd = pt = 0;
    node();
  inline int get(char c) { return c - 'a'; }
  void insert(const string &T) {
     int u = 0;
     for (char c : T) {
   if (!next[u][get(c)]) next[u][get(c)] = node();
   u = next[u][get(c)];
     ans[pt] = 0;
     out[u].push back(pt++);
  void build() {
     aueue <int> a:
     for (q.push(0); !q.empty(); ) {
       int u = q.front();
       q.pop();
       int c = 0; c < A; ++c) {
   int v = next[u][c];
   if (!v) next[u][c] = next[link[u]][c];</pre>
            link[v] = u ? next[link[u]][c] : 0;
out_link[v] = out[link[v]].empty() ?
              → out link[link[v]] : link[v];
             ed[link[v]].push back(v);
             q.push(v);
```

```
void dfs(int s) {
    for(int x : ed[s]) dfs(x), cnt[s] += cnt[x];
    for(int e : out[s]) ans[e] = cnt[s];
  void traverse(const string &S) {
    int u = 0:
    for (char c : S) {
  u = next[u][get(c)];
      cnt[u]++;
    dfs(0);
char str[1000010], pat[505];
int main() {
        freopen("in.txt","r",stdin);
  AC aho;
 int t,T;
scanf("%d",&T);
  for(int t=1;t<=T;t++) {</pre>
    int n;
    scanf("%d",&n);
scanf("%s",str);
    for(int i=1;i<=n;i++) {
      scanf("%s",pat);
      aho.insert(pat);
    aho.build();
    aho.traverse(str);
    printf("Case `%d:\n",t);
    for(int i=0;i<n;i++) {
      printf("%d\n",aho.ans[i]);
    aho.clear();
  return 0:
```

8.2 fft_match

```
using ld = double;
using cd = complex<ld>;
const ld PI = acos(-1.0);
const ld eps = 1e-6;
void fft(vector<cd>& a, bool invert)
    int n = (int)a.size();
    for(int i = 1, j = 0; i < n; i++) {
          int bit = n >> 1
         for(; j & bit; bit >>= 1) j ^= bit;
          i ^= bit;
         if(i < j) swap(a[i], a[j]);
    for(int len = 2; len <= n; len <<= 1) {
    ld ang = 2 * PI / len * (invert ? -1 : 1);</pre>
         cd wlen(cosl(ang), sinl(ang));
         for(int i = 0; i < n; i += len) {
              cd w(1);
              for(int j = 0; j < len / 2; j++) {
  cd u = a[i + j], v = a[i + j + len/2]</pre>
                   a[i + j] = u + v;
                   a[i + j + len/2] = u - v;
                   w^* = wlen;
         }
    if(invert) {
         for(cd\&x:a) {
              x /= n;
```

```
vector<i64> multiply(vector<cd> const& a. vector<cd>
    vector<cd> fa(a.begin(), a.end()), fb(b.begin(),

→ b.end());

     int n = 1;
    while(n < (int)a.size() + b.size()) {</pre>
         n <<= 1;
     fa.resize(n);
    fb.resize(n);
     fft(fa, false);
    fft(fb, false);
     for(int i = 0; i < n; i++) fa[i] *= fb[i];
    fft(fa, true);
    vector<i64> res(n);
    for(int i = 0; i < n; i++) {
         if(abs(fa[i].imag()) < eps) </pre>
              res[i] = round(fa[i].real());
             if(abs(res[i] - fa[i].real()) > eps) {
                  res[i] = -1;
         } else {
             res[i] = -1;
    return res:
|int main()
         |t| \ll |s|
         Want to check if t occurs as a substring of s
    string s, t; cin >> s >> t;
    vector <cd> sa(s.size()), ta(t.size());
    for (int i = 0; i < s.size(); ++i) {
  ld ang = 2 * PI * (s[i] - 'a' + 1) / 27;
         sa[i] = cd(cosl(ang), sin(ang));
    for (int i = 0; i < t.size(); ++i) {
    ld_ang = 2 * PI * (t[i] - 'a' + 1) / 27;</pre>
         ta[t.size() - 1 - i] = cd(cosl(ang),
          → -sin(ang));
    vector<i64> mul = multiply(sa, ta);
    for (int i = 0; i < sa.size(); ++i) {</pre>
         if (mul[i] == ta.size()) {
             cout << i - ta.size() + 1 << ' ';
     cout << '\n';
8.3 hash
```

```
struct Hash {
   struct base {
       string s; int b, mod;
       vector<int> hash, p;
       void init(string & s, int b, int mod) \{ // b >
        → 26. prime.
          s = \underline{\hspace{0.1cm}} s; b = \underline{\hspace{0.1cm}} b, mod = \underline{\hspace{0.1cm}} mod;
           hash_resize(s.size());
           p.resize(s.size());
           hash[0] = s[0] - A' + 1; p[0] = 1;
          for(int i = 1; i < s.size(); ++i) {
    hash[i] = (ll) hash[i - 1] * b % mod;
    hash[i] += s[i] - 'A' + 1;
    if(hash[i] >= mod) hash[i] -= mod;
    p[i] = (ll) p[i - 1] * b % mod;
```

```
int get(int l, int r) {
     int ret = hash[r];
if(l) ret -= (ll) hash[l - 1] * p[r - l + 1] %
     if(ret < 0) ret += mod;
     return ret;
void init(string &s) {
h[0].init(s, 29, 1e9+7);
h[1].init(s, 31, 1e9+9);
pair<int, int> get(int l, int r) {
   return { h[0].get(l, r), h[1].get(l, r) };
```

```
8.4 hash_segtree
                              - 1
#define INVALID CHAR
namespace strhash {
 int n;
  const int MAX = 100010;
  int ara[MAX];
 const int MOD[] = {1067737007, 1069815139};
const int BASE[] = {982451653, 984516781};
  int BP[2][MAX], CUM[2][MAX];
  void init(char *str) {
    n = strlen(str);
    for(int i=0;i<n;i++) ara[i] = str[i]-'0'+1;
 void precal() {
  BP[0][0] = BP[1][0] = 1;
  CUM[0][0] = CUM[1][0] = 1;
    for(int i=1;i<MAX;i++)
      BP[0][i] = (BP[0][i-1] * (long long) BASE[0])
         % MOD[0];
      BP[1][i] = (BP[1][i-1] * (long long) BASE[1])
       → % MOD[1]:
      CUM[0][i] = (CUM[0][i-1] + (long long) BP[0][i]
      → ) % MOD[1];
  struct node {
    int sz;
    int h[2]
    node()
  } tree[4*MAX];
  int lazy[4*MAX];
  inline void lazyUpdate(int n,int st,int ed) {
    if(lazy[n]!=INVALID CHAR){
      tree[n].h[0] = (lazy[n] * (long long)

    CUM[0][ed-st]) % MOD[0];
tree[n].h[1] = (lazy[n] * (long long))

→ CUM[1][ed-st]) % MOD[1];
      if(st!=ed){
        lazy[2*n] = lazy[n];
        lazy[2*n+1] = lazy[n];
      lazy[n] = INVALID CHAR;
  inline node Merge(node a, node b) {
    node ret;
    ret.h[0] = ((a.h[0] * (long long) BP[0][b.sz])
    \rightarrow + b.h[0] ) % MOD[0];
```

```
ret.h[1] = ( (a.h[1] * (long long) BP[1][b.sz] )
    \rightarrow + b.h[1] ) % MOD[1];
    ret.sz = a.sz + b.sz;
    return ret;
  inline void build(int n,int st,int ed) {
    lazy[n] = INVALÌD CHAŔ;
    if(st==ed)
      tree[n].h[0] = tree[n].h[1] = ara[st];
      tree[n].sz = 1;
      return;
    int_mid = (st+ed) >> 1;
    build(n+n,st,mid);
    build(n+n+1,mid+1,ed);
    tree[n] = Merge(tree[n+n], tree[n+n+1]);
  inline void update(int n,int st,int ed,int i,int
  lazyUpdate(n,st,ed);
    if(st>j or ed<i) return;</pre>
    if(st>=i and ed<=j) {
      lazy[n] = v;
      lazyUpdate(n,st,ed);
      return;
    int mid = (st+ed) >> 1;
    update(n+n,st,mid,i,j,v);
    update(n+n+1,mid+1,ed,i,j,v);
    tree[n] = Merge(tree[n+n], tree[n+n+1]);
  inline node guery(int n,int st,int ed,int i,int j){
    lazyUpdate(n,st,ed);
    if(st>=i and ed<=j) return tree[n];</pre>
    int mid = (st+ed)/2;
    if(mid<i) return query(n+n+1,mid+1,ed,i,j);</pre>
    else if(mid>=j) return query(n+n,st,mid,i,j);
    else return Merge(query(n+n,st,mid,i,j),query(n+n+)

→ 1.mid+1.ed.i.i)):
8.5 kmp
// returns the longest proper prefix array of pattern p
// where lps[i]=longest proper prefix which is also
\rightarrow suffix of p[0...i]
vector<int> build lps(string p) {
  int sz = p.size();
  vector<int> lps;
  lps.assign(sz + 1, 0);
  int i = 0:
  lps[0] = 0;
  for(int i = 1; i < sz; i++)
    while(j >= 0 \& p[i] != p[j]) {
      if(j >= 1) j = lps[j - 1];
      else i = -1.
    [[i] = j;
  return lps;
vector<int>ans:
// returns matches in vector ans in 0-indexed
```

void kmp(vector<int> lps, string s, string p) {

int psz = p.size(), sz = s.size();

while(j >= 0 && p[j] != s[i])

if(j >= 1) j = lps[j - 1];

for(int i = 0; i < sz; i++)

int j = 0;

```
else j = -1;
    if(j == psz) {
      i = lps[j - 1];
      // pattern found in string s at position i-psz+1
      ans.push back(i - psz + 1);
    // after each loop we have j=longest common suffix
       of s[0..i] which is also prefix of p
8.6 manachar
vector<<mark>int</mark>> d1(n); // maximum odd length palindrome
   centered at i
                     // here d1[i]=the palindrome has
                        d1[i]-1 right characters from i
                     // e.g. for aba, d1[1]=2;
for (int i = 0, l = 0, r = -1; i < n; i++) {
  int k = (i > r) ? 1 : min(d1[l + r - i], r - i)
  while (0 \le i - k \&\& i + k < n \&\& s[i - k] == s[i + k]
   k++;
  d1[i] = k--;
  if_{i}(i + k > r) {
    l = i - k;
    r = i + k:
vector<int> d2(n); // maximum even length palindrome
// here d2[i]=the palindrome has
                        d2[i]-1 right characters from i
                     // e.g. for abba, d2[2]=2;
for (int i = 0, l = 0, r = -1; i < n; i++) {
 int k = (i > r) ? 0 : min(d2[l + r - i + 1], r - i +
  while (0 \le i - k - 1) \& i + k < n \& s[i - k - 1]
   \Rightarrow == s[i + k]) \{
    k++;
  d2[i] = k--;
  if^{-}(i + k > r)  {
    l`= i - k - 1;
    r = i + k;
8.7 palindromic_tree
const int A = 26;
const int N = 300010;
char s[N]; long long ans;
int last, ptr, nxt[N][A], link[N], len[N], occ[N];
void feed (int at)
  while (s[at - len[last] - 1] != s[at]) last =
      link[last];
  int ch = s[at] - 'a', temp = link[last];
  while (s[at - len[temp] - 1] != s[at]) temp =
   → link[temp];
  if (!nxt[last][ch])
    nxt[last][ch] = ++ptr, len[ptr] = len[last] + 2;
    link[ptr] = len[ptr] == 1 ? 2 : nxt[temp][ch];
  last = nxt[last][ch], ++occ[last];
|int main() {
  len[1] = -1, len[2] = 0, link[1] = link[2] = 1, last
  \rightarrow = ptr = 2;
  scanf("%s", s + 1);
  for (int i = 1, n = strlen(s + 1); i <= n; ++i)</pre>
  → feed(i);
```

```
for (int i = ptr; i > 2; --i) ans = max(ans, len[i]
     * 1LL * occ[i]), occ[link[i]] += occ[i];
  printf("%lld\n", ans);
  return 0;
8.8 persistant_trie
const int MAX = 200010;
const int B = 19;
int root[MAX], ptr = 0;
struct node {
  int ara[2], sum;
  node()
}    tree[ MAX * (B+1) ];
void insert(int prevnode, int &curRoot, int val) {
  curRoot = ++ptr;
  int curnode = curRoot;
  for(int i = B; i \ge 0; i - -) {
    bool bit = val & (1 \ll i);
    tree[curnode] = tree[prevnode];
    tree[curnode].ara[bit] = ++ptr;
    tree[curnode].sum += 1
    prevnode = tree[prevnode].ara[bit];
    curnode = tree[curnode].ara[bit];
  tree[curnode] = tree[prevnode];
  tree[curnode].sum += 1;
int find xor max(int prevnode, int curnode, int x) {
  int ans = \overline{0};
  for(int i = B; i >= 0; i--) {
    pool bit = x & (1 << i);</pre>
    if(tree[tree[curnode].ara[bit ^ 1]].sum >
     - tree[tree[prevnode].ara[bit ^ 1]].sum) {
  curnode = tree[curnode].ara[bit ^ 1];
      prevnode = tree[prevnode].ara[bit ^1];
      ans = ans | (1 << i);
    else {
      curnode = tree[curnode].ara[bit]
      prevnode = tree[prevnode].ara[bit];
  return ans:
void solve() {
  int n, q, L, R, K;
  cin >> n;
  for(int i=1;i<=n;i++) cin >> ara[i];
  for(int i=1;i<=q;i++) {
    cin >> L >> R >> K:
    cout << find xor max(root[L-1],root[R],K) << endl;</pre>
8.9 suffix array
// Everything is 0-indexed
char s[N]; // Suffix array will be built for this

→ strina

int SA[N], iSA[N]; // SA is the suffix array, iSA[i]
    stores the rank of the i'th suffix
int cnt[N], nxt[N]; // Internal stuff
bool bh[N], b2h[N]; // Internal stuff
int lcp[N]; // Stores lcp of SA[i] and SA[i + 1];
```

lcp[n - 1] = 0

int lcpSparse[LOGN][N]; // lcpSparse[i][j] =

min(lcp[i], ..., lcp[i - 1 + (1 << i)])

```
void buildSA(int n) {
  for (int i = 0; i < n; i++) SA[i] = i;
  sort(SA, SA + n, [](int i, int j) { return s[i] <</pre>
  for (int i = 0; i < n; i++) {
  bh[i] = i == 0 || s[SA[i]] != s[SA[i - 1]];</pre>
    b2h[i] = 0;
  for (int h = 1; h < n; h <<= 1) {
    int tot = 0;
    for (int i = 0, j; i < n; i = j) {
       j = i + 1;
      while (j < n && !bh[j]) j++;
      nxt[i] = j; tot++;
    } if (tot == n) break;
    for (int i = 0; i < n; i = nxt[i])</pre>
       for (int j = i; j < nxt[i]; j++) iSA[SA[j]] = i;</pre>
       cnt[i] = 0;
    cnt[iSA[n - h]]++;
b2h[iSA[n - h]] = 1;
    for (int i = 0; i < n; i = nxt[i])
      for (int j = i; j < nxt[i]; j++) {
  int s = SA[j] - h;</pre>
         if (s < 0) continue;
         int head = iSA[s];
         iSA[s] = head + cnt[head]++;
b2h[iSA[s]] = 1;
       for (int j = i; j < nxt[i]; j++) {
         int s = SA[i] - h;
         if (s < 0 | | !b2h[iSA[s]]) continue;</pre>
         for (int k = iSA[s] + 1; !bh[k] \&\& b2h[k];
          \rightarrow k++) b2h[k] = 0;
    for (int i = 0; i < n; i++) {
    SA[iSA[i]] = i;</pre>
      bh[i] [= b2h[i];
  for (int i = 0; i < n; i++) iSA[SA[i]] = i;
void buildLCP(int n) {
  for (int i = 0, k = 0; i < n; i++) {
    if (iSA[i] == n - 1) {
       k = 0;
       lcp[n - 1] = 0;
       continue;
    int j = SA[iSA[i] + 1];
    while (i + k < n \&\& j + k < n \&\& s[i + k] == s[j + k]
    lcp[i$A[i]]' = k;
    if (k) k--;
void buildLCPSparse(int n) {
  for (int i = 0; i < n; i++) lcpSparse[0][i] = lcp[i];</pre>
  for (int i = 1; i < LOGN; i++) {
    for (int j = 0; j < n; j++) {
   lcpSparse[i][j] = min(lcpSparse[i - 1][j],</pre>
           lcpSparse[i - 1][min(n - 1, j + (1 << (i -
        □ 1)))]);
pair<int, int> minLCPRange(int n, int from, int

→ minLCP) {
  int r = from;
```

```
for (int i = LOGN - 1; i >= 0; i--) {
   int jump = 1 << i;
    if (r + jump < n and lcpSparse[i][r] >= minLCP) r

→ += jump;

 int l = from;
 for (int i = LOGN - 1; i >= 0; i --) {
   int jump = 1 << i;
   if (l - jump >= 0 and lcpSparse[i][l - jump] >=

→ minLCP) l -= iump:

 return make pair(l, r);
8.10 suffix_automata
namespace <mark>sa</mark>{
 const int MAXN = 100005 << 1;
 const int MAXC = 26;
 char str[MAXN];
 int n, sz, last;
 int len[MAXN], link[MAXN], ed[MAXN][MAXC], cnt[MAXN];
 bool terminal[MAXN]
 vector <int> G[MAXN];
 void init() {
    \begin{array}{l} \overline{SET(ed[0])}; \\ len[0] = 0, \ link[0] = -1, \ sz = 1, \ last = 0, \end{array}

    terminal[0] = false;

 inline int scale(char c) { return c-'a'; }
 void extend(char c) {
   int cur = sz++;
   terminal[cur] = false;
    cnt[cur] = 1;
    SET(ed[cur]);
len[cur] = len[last] + 1;
    int p = last;
    while (p != -1 \&\& ed[p][c]==-1) {
      ed[p][c] = cur;
      p = link[p];
    if (p == -1) link[cur] = 0;
    else {
      int q = ed[p][c];
      if (len[p] + 1 == len[q]) link[cur] = q;
      else {
        int clone = sz++;
len[clone] = len[p] + 1;
        memcpy(ed[clone],ed[q],sizeof(ed[q]));
        link[clone] = link[q];
        while (p != -1 \&\& ed[p][c] == q) {
          ed[p][c] = clone;
          p = link[p];
        link[q] = link[cur] = clone;
        cnt[clone] = 0;
        terminal[clone] = false;
    last = cur;
 void dfs(int s) {
  for(auto x : G[s]) dfs(x), cnt[s] += cnt[x];
 void build() {
   init();
    int n = strlen(str);
    for(int i=0;i<n;i++) extend(scale(str[i]));</pre>
    for(int i=1;i<sz;i++) G[link[i]].pb(i);</pre>
    for(int i=0;i<sz;i++) G[i].clear();</pre>
```

```
for(int i=last;i!=-1;i=link[i]) terminal[i] = true;
8.11 trie
#define N
#define S
                  200000
26
int root.now:
int nxt[N][S], cnt[N];
void init(){
  root = now = 1;
  CLR(nxt), CLR(cnt);
inline int scale(char ch) { return (ch - 'a'); }
inline void Insert(char s[],int sz){
  int cur = root, to;
  for(int i=0 ; i < sz ; i++){</pre>
    to = scale(s[i]);
    if( !nxt[cur][to] ) nxt[cur][to] = ++now;
cur = nxt[cur][to];
  cnt[cur]++;
inline bool Find(char s[],int sz){
  int cur = root, to;
for(int i=0; i<sz; i++){</pre>
    to' = scale(s[i])
    if( !nxt[cur][to]) return false;
    cur = nxt[cur][to];
  return (cnt[cur]!=0);
inline void Delete(char s[],int sz){
  int cur = root, to;
  for(int i=0 ; i<sz ; i++){
    to = scale(s[i])
    cur = nxt[cur][to];
  cnt[cur]--;
8.12 z_algo
const int N = 100010;
char s[N];
int t, n, z[N];
int main() -
  scanf("%s
  n = strlen(s), z[0] = n;
  int L = 0, R = 0
  for (int i = 1; i < n; ++i) {
    if (i > R) {
      L = R = i:
      while (\underline{R} < \underline{n} \&\& \underline{s}[R - L] == s[R]) ++R;
      z[i] = R - L; --R;
    } else {
```

int k = i - L;

z[i] = R - L; --R;

else {

if (z[k] < R - i + 1) z[i] = z[k];

while (R < n && s[R - L] == s[R]) ++R;