

DU_Scolps

Asif Jawad, Sakib Hassan, Bholanath Das Niloy

Team Reference Document for NCPC 2024

Updated: February 2024

Contents			5 Math			7.12.1 Titu's Lemma
U	Contents	l,	5.1	FWHT	13 13	7.13 Games
1	1 DP	1	5.2	FloorSum	13	7.13.1 Grundy numbers
_	1 DI 11 CHT	11	5.3	NOD	$\bar{13}$	7.13.2 Sums of games
	1.1 CHT 1.2 DC_Optimization	히	5.4	Pollard Rho	$\bar{1}\bar{3}$	
	1.3 SOS_DP	5	5.5	catalan	14	7.13.3 Misère Nim
	1.4 dynamic cht	4	5.6	crt	14	7.14 NumberTheory
	1.4 dynamic_cnt	Z	5.7	derangement	14	
_			5.8	diophantine	14	8 String 19
2	2 Data _S tructure	2	5.9	discrete_log	14	8.1 aho_corasick
	2.1 2D_segtree	2) factorial_mod_p	14	$8.2 \text{ hash} \dots 20$
	2.2 Segtree_beats_desc	3		fft	14	8.3 hash_segtree
	2.3 gp_hash_table	3		gauss_eliminition		8.4 kmp
			5.1	gauss_emminuon		
	2.4 iterative_segtree	3	5.1	gen_all_k_combs	15	
	2.5 mos_algo	3		integrate_adaptive		8.6 palindromic_tree
	2.6 ordered_set	3	5.1	ó matrix_expo	15	8.7 persistant_trie
	2.7 persistant_segtree	3	5.1	next_lexicographical_k_comb	15	
	2.8 segment_tree	3		' ntt	15	8.9 suffix_automata
	2.9 sparse_table	4		8 seg_sieve	16	8.10 trie
	······································		5.1°	stirling	16	8.11 z_algo
3	3 Geometry	4	5.2	stirling_number_of_the_second_kind	16	
U	3.1 2D Primitive	7	5.2	sum_of_totient	16	
	3.1.1 Angle	4	5.4	totient		
		4		a totient	17	1 DP 1.1 CHT
	3.1.2 Line Distance	4	0.2	O AUI_Dasis	11	
	3.1.3 Line Intersection	4 6	6 Mi		17	ll M[MAX] , C[MAX];
	3.1.4 Linear Transformation	41		interval container	17 17	struct CHT {
		4	6.2		$\frac{17}{17}$	<pre>int len , cur; void init() {</pre>
	3.1.6 Point Sort	4	6.3	pragma	$\frac{17}{17}$	void init() {
	3.1.8 Segment Distance	5				len = 0 ,cur = 0;
		5	$\frac{6.4}{6.5}$	random	$\begin{array}{c} 17 \\ 17 \end{array}$	}
	3.1.9 Segment Intersection	9	0.5	VIIIIrC	11	<pre>inline bool isBad(ll nm,ll nc) {</pre>
	3.1.10 Side Of	5/2	7 No	.00	17	<pre>inline bool isBad(ll nm,ll nc) { return ((C[len-1]-C[len-2])/(double)(M[len-2]-M[l])</pre>
	3.2 3D	51,	7.1		$\frac{1}{17}$	on $ $ $ $ $ $ $ $ $ $ $ $ $ $ $ $
	3.2.1 3D Convex Hull	2	7.1	0] =): ', (
	3.2.3 Polyhedron Volume	6	$7.3^{1.2}$	Fibonacci	18 18	//return ((C[len-1]-C[len-2])*(M[len-2]-nm) >=
		0	7.3	Geometry	18	
	3.2.4 Spherical Distance	0	1.1			}
	3.3 Circle	6				<pre>inline void addLine(ll nm,ll nc) {</pre>
	3.3.1 Circle Intersection	9		7.4.2 Quadrilaterals		<pre>if(len == 0) M[len] = nm , C[len] = nc , ++len; else if(M[len-1] == nm) {</pre>
	3.3.2 Circle Polygon Intersection	9		7.4.3 Spherical coordinates	18	else if(M[len-1] == nm) {
	3.3.3 Circle Tangents	6	$\frac{7.5}{2}$	Sums	18	<pre>if(C[len-1] <= nc) return:</pre>
	3.3.4 CircumCircle	6	7.6	Series	18	<pre>if(C[len-1] <= nc) return; else C[len-1] = nc;</pre>
	3.4 Polygon	6	7.7	Pythagorean Triples	18	}
	3.4.1 Hull Diameter	6	7.8	Number Theory	18	élse {
	3.4.2 Line Hull Intersection	7		7.8.1 Primes	18	<pre>while(len >= 2 &\& isBad(nm,nc))len; M[len] = nm , C[len] = nc , ++len;</pre>
	3.4.3 Polygon Center	7		7.8.2 Estimates	18	M[len] = nm , C[len] = nc , ++len;
	3.4.4 Polygon Cut	7		7.8.3 Perfect numbers	$\frac{18}{18}$	
	3.5 Closest Pair	7		7.8.4 Carmichael numbers		
	3.5 Closest Pair 3.6 Convex Hull 3.7 Minimum Enclosing Circle	$\overline{2}$		7.8.5 Mobius function	18	<pre>inline ll getY(int id , ll x) { return (M[id]*x + C[id]);</pre>
	3.7 Minimum Enclosing Circle	7				return (M[id]*x + C[id]);
	3.8 Point In Polygon	8		7.8.7 Jacobi symbol	18	}
	3.9 near_pair	8		7.8.8 Primitive roots	18	<pre>inline ll sortedQuery(ll x) {</pre>
	3.10 sweep	8		7.8.9 Discrete logarithm problem	18	if(cur >= len) cur = len-1:
		Ĭ		7.8.10 Pythagorean triples		l while(cur < len-1 && getY(cur+1 x) >= getY(cur x)
4	4 Graph	9		7.8.11 Postage stamps/McNuggets problem	18	
-	4.1 2Sat	ด		7.8.12 Fermat's two-squares theorem	18	→) cur++; return getY(cur,x);
	4.1 25at	å	7.9	Permutations	18	}
	4.2 Centroid_decomp	9		7.9.1 Factorial	18	inlino 11 TC/ 11 v) (
	4.3 articulation_point	19			18	inline ll TS(ll x) { int low = 0, high = len-1 , mid ;
		10		7.9.3 Derangements	19	Int tow - 0, high - ten-1 , mid ,
	4.5 bridge_tree	10		7.9.4 Burnside's lemma	19	mid - low + high >> 1:
	4.6 dinic	10	7.1	Partitions and subsets	19	
	4.7 dsu_on_tree			7.10.1 Partition function	19	\mathbf{I} (get \mathbf{I} (mid, \mathbf{X}) < get \mathbf{I} (mid+1, \mathbf{X}) tow = mid + 1;
	4.8 euler_path	11	7.1	General purpose numbers	19	else high = mid;
		11		7.11.1 Stirling numbers of the first kind	19	}
	4.10 hopcroft_karp	11		7.11.2 Eulerian numbers	19	return max(getY(low,x),getY(high,x));
	4.11 hungarian	12		7.11.3 Stirling numbers of the second kind	19	\1. ^J
	4.12 kuhn	12				
	4.13 lca	12		7.11.4 Bell numbers	19	CHT cht;
	4.14 manhattan_MST	12		7.11.6 Catalan numbers	19	Circ. Init(),
	4.15 mcmf		7.13	Inequalities	19	

void compute(int L, int R, int optL, int optR){

1.2 DC_Optimization

if(L > R) return;

int $M = L + R \gg 1$

```
9
```

```
pair<ll, int> best(1LL << 60, -1);
  for(int k = optL; k <= min(M, optR); k++){</pre>
    best = min(best, \{dp[prv][k] + C[k + 1][M], k\});
  dp[now][M] = best.ff;
  compute(L, M - 1, optL, best.ss);
  compute(M + 1, R, best.ss, optR);
1.3 SOS DP
for(int mask = 0; mask < (1<<N); ++mask){
         dp[mask][-1] = A[mask];
         for(int i = 0; i < N; ++i){
             if(mask \& (1 << i)) dp[mask][i] =
              \rightarrow dp[mask][i-1] + dp[mask \land (1<<i)][i-1];
             else dp[mask][i] = dp[mask][i-1];
         f[mask] = dp[mask][N-1];
    for(int i = 0; i < (1 << N); ++i) F[i] = A[i];
    for(int i = 0; i < N; ++i)
         for(int mask = 0; mask < (1 << N); ++mask){
             if(mask & (1<<i))
                  F[mask] += F[mask^(1<<i)];
1.4 dynamic_cht
//add lines with -m and -b and return -ans to
//make this code work for minimums.(not -x)
const ll inf = -(1LL << 62);</pre>
struct line {
  ll m, b;
  mutable function<const line*() > succ;
  bool operator < (const line& rhs) const {</pre>
    if (rhs.b != inf) return m < rhs.m:</pre>
    const line* s = succ();
    if (!s) return 0;
    ll x = rhs.m;
    return b - s->b < (s->m - m) * x;
struct CHT : public multiset<line> {
  bool bad(iterator y) {
    auto z = next(y)
    if (y == begin()) {
       if(z == end()) return 0;
       return y \rightarrow m == z \rightarrow m \&\& y \rightarrow b <= z \rightarrow b;
    auto x = prev(y);
    if (z == end()) return y -> m == x -> m &  y -> b
     \rightarrow <= x -> b;
    return 1.0 * (x -> b - y -> b) * (z -> m - y -> m)
     \Rightarrow >= 1.0 * (v -> b - z -> b) * (v -> m - x -> m);
  void add(ll m, ll b) {
    auto y = insert({'m, b });
y->succ = [ = ] { return next(y) == end() ? 0 :
     if (bad(y)) {
       erase(y);
      return;
    while (next(y) != end() \&\& bad(next(y)))
        erase(next(v));
    while (y != begin() \&\& bad(prev(y)))
     → erase(prev(y));
```

```
ll query(ll x) {
     auto 1 = *lower bound((line) {
      x, inf
    return l.m * x + l.b;
ľCHT* cht;
ll a[N], b[N]
|int32 t main() {
  ios base::sync with stdio(0);
  cin_tie(0):
  int n;
  cin >> n;
  for(int i = 0; i < n; i++) cin >> a[i];
  for(int i = 0; i < n; i++) cin >> b[i];
cht = new CHT();
  cht -> add(-b[0], 0);
  ll ans = 0;
  for(int i = 1; i < n; i++)
    ans = -cht \rightarrow query(a[i]);
     cht -> add(-b[i], -ans);
  cout << ans << nl;
  return 0;
|2 Data<sub>S</sub>tructure
2.1 2D_segtree
```

```
struct Point {
    int x, y, mx;
Point() {}
    Point(int x, int y, int mx) : x(x), y(y), mx(mx) {}
    bool operator < (const Point& other) const {</pre>
        return mx < other.mx;</pre>
struct Segtree2d {
    // I didn't calculate the exact size needed in
    terms of 2D container size.
// If anyone, please edit the answer.
    // It's just a safe size to store nodes for MAX *
        MAX 2D grids which won't cause stack overflow
    Point T[500000]; // TODO: calculate the accurate

→ space needed
    int n, m;
    // initialize and construct segment tree
    void init(int n, int m) {
         this \rightarrow n = n;
         this -> m = m:
         build(1, 1, 1, n, m);
    // build a 2D segment tree from data [ (a1, b1),
     // Tìme: O(n logn)
    Point build(int node, int a1, int b1, int a2, int
     → b2) {
         // out of range
        if (a1 > a2 or b1 > b2)
             return def();
         // if it is only a single index, assign value
             to node
        if (a1 == a2 and b1 == b2)
    return T[node] = Point(a1, b1, P[a1][b1]);
            split the tree into four segments
         T[node] = def();
         T[node] = maxNode(T[node], build(4 * node - 2,
         \rightarrow a1, b1, (a1 + a2) / 2, (b1 + b2) / 2 ));
```

```
T[node] = maxNode(T[node], build(4 * node - 1,
         (a1 + a2) / 2 + 1, b1, a2, (b1 + b2) / 2
    T[node] = maxNode(T[node], build(4 * node + 0,
        a1, (b1 + b2) / 2 + 1, (a1 + a2) / 2, b2)
    T[node] = maxNode(T[node], build(4 * node + 1,
         (a1 + a2) / 2 + 1, (b1 + b2) / 2 + 1, a2,
        b2)):
    return T[node];
// helper function for query(int, int, int, int);
Point query(int node, int al, int bl, int a2, int
\rightarrow b2, int x1, int y1, int x2, int y2) {
    // if we out of range, return dummy
    if (x1 > a2 \text{ or } y1 > b2 \text{ or } x2 < a1 \text{ or } y2 < b1
    \rightarrow or a1 > a2 or b1 > b2)
        return def();
    // if it is within range, return the node
    if (x1 \le a1 \text{ and } y1 \le b1 \text{ and } a2 \le x2 \text{ and } b2
    <= y2)
         return T[node];
    // split into four segments
    Point mx = def();
    mx = maxNode(mx, query(4 * node - 2, a1, b1,
         (a1 + a2) / 2, (b1 + b2) / 2, x1, y1, x2,
    mx = maxNode(mx, query(4 * node - 1, (a1 + a2))
        /2 + 1, b1, a2, (b1 + b2) /2, x1, y1,
     \equiv x2. v2) ):
    mx = maxNode(mx, query(4 * node + 0, a1, (b1 +
        b2) / 2 + 1, (a1 + a2) / 2, b2, x1, y1,
     \stackrel{\sim}{\rightarrow} x2, y2));
    mx = maxNode(mx, query(4 * node + 1, (a1 + a2))
        /2 + 1, (b1 + b2) /2 + 1, a2, b2, x1,

    y1, x2, y2));

    // return the maximum value
    return mx;
// query from range [ (x1, y1), (x2, y2) ]
// ˈtimeː O(logn)
Point query(int x1, int y1, int x2, int y2)
    return query(1, 1, 1, n, m, x1, y1, x2, y2);
// helper function for update(int, int, int);
Point update(int node, int al, int bl, int a2, int
→ b2, int x, int y, int value) {
    if (a1 > a2 \text{ or } b1 > b2)
        return def();
    if (x > a2 \text{ or } y > b2 \text{ or } x < a1 \text{ or } y < b1)
        return T[node];
    if (x == a1 \text{ and } y == b1 \text{ and } x == a2 \text{ and } y ==

→ b2)

         return T[node] = Point(x, y, value);
    Point mx = def();
    mx = maxNode(mx, update(4 * node - 2, a1, b1,
         (a1 + a2) / 2, (b1 + b2) / 2, x, y, value)
    ≒ );
    mx = maxNode(mx, update(4 * node - 1, (a1 +
         a2) / 2 + 1, b1, a2, (b1 + b2) / 2, x, y,

    value))

    mx = maxNode(mx, update(4 * node + 0, a1, (b1))
        + b2) / 2 + 1, (a1 + a2) / 2, b2, x, y,

    value))

    mx = maxNode(mx, update(4 * node + 1, (a1 +
        a2) / 2 + 1, (b1 + b2) / 2 + 1, a2, b2, x,
        y, value) );
    return T[node] = mx;
```

```
// update the value of (x, y) index to 'value'
    // Time: 0(logn)
    Point update(int x, int y, int value) {
        return update(1, 1, 1, n, m, x, y, value);
    // utility functions; these functions are virtual
    → because they will be overridden in child class
    virtual Point maxNode(Point a, Point b) {
        return max(a, b);
    // dummy node
   virtual Point def() {
    return Point(0, 0, -INF);
/* 2D Segment Tree for range minimum query; a override

→ of Segtree2d class */
struct Segtree2dMin : Segtree2d {
    // overload maxNode() function to return minimum
    Point maxNode(Point a, Point b) {
        return min(a, b);
    Point def() {
        return Point(0, 0, INF);
```

2.2 Segtree_beats_desc

```
Description: For update Ai Ai mod x and similar, keep
max in node and lazily update whenever min = max. For
→ update
Ai min(Ai, x) and similar, keep range max, second max
    in node and
lazily update whenever x > second max.
Time: O(\log^2 N), (\log N)
```

2.3 gp hash table

```
using namespace gnu pbds;
const int RANDOM = chrono::high resolution clock::now(
→ ).time since epoch().count();
using namespace __gnu_pbds;
struct chash
  const int RANDOM = (long
      long)(make unique<char>().get()) ^
      chrono::high resolution clock::now().time since
      epoch().count();
  static unsigned long long hash f(unsigned long long
    x += 0x9e3779b97f4a7c15;
   x = (x ^ (x >> 30)) * 0xbf58476d1ce4e5b9;
x = (x ^ (x >> 27)) * 0x94d049bb133111eb;
return x ^ (x >> 31);
  static unsigned hash combine(unsigned a, unsigned b)
      { return a * 31 + b; }
  ll operator()(ll x) const { return hash f(x)^RANDOM;
  → }
qp hash table<key, long long, chash> table;
```

2.4 iterative_segtree

```
const int N = 500010;
int n, a[N], tree[N << 1];</pre>
void init() {
```

```
for (int i = 0; i < n; ++i) tree[n + i] = a[i];</pre>
    for (int i = n - 1; i >= 0; --i)
        tree[i] = min(tree[i << 1], tree[i << 1 | 1]);
void update(int p, int v) {
    for (tree[p += n] = v; p > 1; p >>= 1)
        tree[p >> 1] = min(tree[p], tree[p ^{\uparrow} 1]);
|int query(int l, int r) {
    int ret = INT MAX:
    for (l += n, r += n; l < r; l >>= 1, r >>= 1) {
         if (l & 1) ret = min(ret, tree[l++]);
        if (r \& 1) ret = min(ret, tree[--r]);
    return ret;
2.5 mos algo
```

```
struct data{
  int l, r, id, bn;
data() {}
  data(int l, int r, int id){
    l = l, r = r, id = i\overline{d};
    bn = l / block sz;
  bool operator < (const data& other) const{</pre>
    if (bn != other.bn) return (bn < other.bn);</pre>
    return ((bn \& 1) ? (r < other.r) : (r > other.r));
int curL = 0, curR = -1;
for(int i = 0; i < 0.sz; i++){
  while(curL > Q[i].L){
    curl--; add(curl);
  while(curR < 0[i].R){</pre>
    curR++; add(curR);
  while(curL < Q[i].L){
    remove(curL); curL++;
  while(curR > Q[i].R){
    remove(curR); curŔ--;
```

2.6 ordered set

```
using namespace std:
using namespace gnu pbds;
typedef tree<
    int
    less < int > , // "less equal<int>," for multiset
    rb tree tag,
    tree order statistics node update > ordered_set;
ordered set OS;
```

2.7 persistant_segtree

```
const int MAX = 100010;
lint ncnt = 0;
|struct node {
  int sum:
  int left, right;
  node() {}
  node(int val) {
    sum = val:
    left = right = -1;
} tree[ ? ];
```

```
int ara[MAX];
int version[MAX];
void build(int n.int st.int ed) {
  if (st==ed) {
    tree[n] = node(ara[st]);
    return;
  int mid = (st+ed) / 2;
  tree[n].left = ++ncnt;
  tree[n].right = ++ncnt;
  build(tree[n].left, st, mid);
  build(tree[n].right, mid+1, ed);
  tree[n].sum = tree[tree[n].left].sum +

    tree[tree[n].right].sum;

void update(int prev,int cur,int st,int ed,int id, int
   val)
  if (id > ed or id < st) return;</pre>
  if (st == ed) {
    tree[cur] = node(val);
    return;
  int mid = (st+ed) / 2;
  if (id <= mid) ·
    tree[cur].right = tree[prev].right;
    tree[cur].left = ++ncnt;
    update(tree[prev].left,tree[cur].left, st, mid,
     \rightarrow id, val);
  else {
    tree[cur].left = tree[prev].left;
    tree[cur].right = ++ncnt;
    update(tree[prev].right, tree[cur].right, mid+1,
     → ed. id. val):
  tree[cur].sum = tree[tree[cur].left].sum +

    tree[tree[cur].right].sum;

int query(int n,int st,int ed,int i,int j){
 if(st>=i && ed<=j) return tree[n].sum;</pre>
  int mid = (st+ed)/2;
  if(mid<i) return query(tree[n].right,mid+1,ed,i,j);</pre>
  else if(mid>=j) return

¬ query(tree[n].left,st,mid,i,j);
  else return query(tree[n].left,st,mid,i,j) +

¬ query(tree[n].right,mid+1,ed,i,j);

int main() {
 int n,q,l,r,k;
  sii(n,q);
  version[0] = ++ncnt;
  build(version[0],1,n);
  version[1] = ++ncnt;
  update(version[0], version[1], 1, n, id, val);
 query(version[0],1,n,id,id);
query(version[1],1,n,id,id);
  return 0;
```

2.8 segment_tree

```
int ara[MAX];
struct node {
 int sum:
} tree[4 * MAX];
int lazy[4 * MAX];
node Merge(node a, node b) {
```

```
node ret;
  ret.sum = a.sum + b.sum;
  return ret;
void lazyUpdate(int n, int st, int ed) {
  if(lazy[n] != 0){
    tree[n].sum += ((ed - st + 1) * lazv[n]);
    if(st != ed){
  lazy[2 * n] += lazy[n];
       lazy[2 * n + 1] += lazy[n];
     lazy[n] = 0;
void build(int n, int st, int ed) {
  lazy[n] = 0;
  if(st == ed){
    tree[n].sum = ara[st];
    return:
 int mid = (st + ed) / 2;
build(2 * n, st, mid);
build(2 * n + 1, mid + 1, ed);
  \tilde{\text{tree}}[\tilde{n}] = \text{Merge}(\tilde{\text{tree}}[2 * \tilde{n}], \text{tree}[2 * n + 1]);
void update(int n, int st, int ed, int i, int j, int

    ∨ ) {
  lazyUpdate(n, st, ed);
  if(st > j or ed < i) return;</pre>
  if(st >= i and ed <= j){
     lazy[n] += v;
    lazyUpdate(n, st, ed);
    return;
  int mid = (st + ed) / 2;
  update(2 * n, st, mid, i, j, v);
  update(2 * n + 1, mid+1, ed, i, j, v);
  tree[n] = Merge(tree[2 * n], tree[2 * n + 1]);
node query(int n, int st, int ed, int i, int j) {
  lazyUpdate(n, st, ed);
  if(st >= i and ed <= j) return tree[n];</pre>
  int mid = (st + ed) / 2;
  if (mid < i) return query (2 * n + 1, mid + 1, ed, i,
  else if(mid >= j) return query(2 * n, st, mid, i, j);
  else return Merge(query(2 * n, st, mid, i, j),
   \rightarrow query(2 * n + 1, mid + 1, ed, i, j));
```

2.9 sparse_table

```
int st[K + 1][MAXN];
void build() {
  std::copy(array.begin(), array.end(), st[0]);
  for (int i = 1; i \le K; i++)
for (int j = 0; j + (1 << i) <= N; j++)
       st[i][j] = f(st[i-1][j], st[i-1][j+(1 <<
        \hookrightarrow (i - 1))]);
```

3 Geometry

3.1 2D Primitive

3.1.1 Angle

A class for ordering angles (as represented by int points and a number #include "Point.h" of rotations around the origin). Useful for rotational sweeping. Some-template<class P> times also represents points or vectors.

```
/* Usage:
 * vector < Angle > v = \{w[0], w[0].t360() ...\}; //

→ sorted
```

```
int j = 0; rep(i,0,n) { while (v[j] < v[i].t180()
            ++j; }
          // sweeps i such that (i-i) represents the number
             of positively oriented triangles with vertices at
             0 and i
|struct Angle {
     int x, y;
     int t:
     Angle(int x, int y, int t=0) : x(x), y(y), t(t) {}
     Angle operator-(Angle b) const { return {x-b.x,
                y-b.y, t}; }
     int half() const {
          assert(x || y);
           return y < 0 | | (y == 0 \&\& x < 0);
     Angle t90() const { return \{-y, x, t + (half() \&\& x\}\}
     Angle t180() const { return {-x, -y, t + half()}; }
Angle t360() const { return {x, y, t + 1}; }
bool operator<(Angle a, Angle b) {
    // add a.dist2() and b.dist2() to also compare
     return make tuple(a.t, a.half(), a.y * (ll)b.x) <</pre>

→ make tuple(b.t, b.half(), a.x * (ll)b.y);

 // Given two points, this calculates the smallest

→ angle between

 // them, i.e., the angle that covers the defined line
pair<Angle, Angle> segmentAngles(Angle a, Angle b) {
     if (b < a) swap(a, b);
     return (b < a.t180() ? make pair(a, b) :
       \rightarrow make pair(b, a.t360()));
Angle operator+(Angle a, Angle b) { // point a +
     Angle r(a.x + b.x, a.y + b.y, a.t);
     if (a.t180() < r) r.t--
     return r.t180() < a ? r.t360() : r;
Angle angleDiff(Angle a, Angle b) { // angle b - angle
     int tu = b.t - a.t; a.t = b.t;
     return \{a.x*b.x + a.y*b.y, a.x*b.y - a.y*b.x, tu - a.y*b
       \rightarrow (b < a)};
```

3.1.2 Line Distance

Returns the signed distance between point p and the line containing points a and b. Positive value on left side and negative on right as seen from a towards b. a==b gives nan. P is supposed to be Point<T> or Point3D<T> where T is e.g. double or long long. It uses products in intermediate steps so watch out for overflow if using int or long long. Using Point3D will always give a non-/s negative distance. For Point3D, call .dist on the result of the cross product.

```
double lineDist(const P& a, const P& b, const P& p) {
 return (double)(b-a).cross(p-a)/(b-a).dist();
```

3.1.3 Line Intersection

If a unique intersection point of the lines going through s1,e1 and s2,e2 exists {1, point} is returned. If no intersection point exists {0, (0,0)} is returned and if infinitely many exists $\{-1, (0,0)\}$ is returned. The wrong position will be returned if P is Point<ll> and the intersection point does not have integer coordinates. Prod- <1 ucts of three coordinates are used in intermediate steps so watch out for overflow if using int or ll.

```
/* Usage:
   auto res = lineInter(s1,e1,s2,e2);
   if (res.first == 1)
  cout << "intersection point at " << res.second</pre>
     << endl:
#pragma once
#include "Point.h"
template<class P>
pair<int, P> lineInter(P s1, P e1, P s2, P e2) {
 auto d = (e1 - s1).cross(e2 - s2);
  if (d == 0) // if parallel
    return \{-(s1.cross(e1, s2) == 0), P(0, 0)\};
  auto p = s2.cross(e1, e2), q = s2.cross(e2, s1);
  return {1, (s1 * p + e1 * q) / d};
```

3.1.4 Linear Transformation

Apply the linear transformation (translation, rotation p0 and scaling) which takes line p0-p1 to line q0-q1 to point q0

```
#include "Point.h"
typedef Point<double> P;
P linearTransformation(const P& p0, const P& p1, const
   P& q0, const P& q1, const P& r) {
 P dp = p1-p0, dq = q1-q0, num(dp.cross(dq),
     dp.dot(dq));
 return q0 + P((r-p0), cross(num),
    (r-p0).dot(num))/dp.dist2();
```

3.1.5 On Segment

```
/* Description: Returns true iff p lies on the line
   segment from s to e.
* Use \texttt{(segDist(s,e,p)<=epsilon)} instead when
    using Point<double>.
#include "Point.h"
template<class P> bool onSegment(P s, P e, P p) {
 return p.cross(s, e) == 0 \& (s - p).dot(e - p) <= 0;
```

```
3.1.6 Point Sort
// sort the points in counterclockwise order that
\rightarrow starts from the half line x0, y=0.
using namespace std;
typedef long long ll;
typedef pair <ll, ll> point;
#define x first
#define y second
int main() {
 int n: cin >> n:
 vector <point> p(n);
 for (auto &it : p) scanf("%lld %lld", &it.x, &it.y);
 sort(p.begin(), p.end(), [] (point a, point b) {
    return atan2l(a.y, a.x) < atan2l(b.y, b.x);</pre>
 for (auto it : p) printf("%lld %lld\n", it.x, it.y);
```

```
return 0;
}
```

3.1.7 Point

```
// Class to handle points in the plane. T can be e.g.

→ double or long long. (Avoid int.)

template <class T> int sgn(T x) \{ return (x > 0) - (x < 0) \}
templaté<class T>
struct Point {
 typedef Point P;
 T x, y;
 explicit Point(T x=0, T y=0) : x(x), y(y) {}
 bool operator<(P p) const { return tie(x,y) <</pre>
 \begin{tabular}{lll} &\rightarrow & \text{tie}(p.x,p.y)\,; & \\ & \text{bool} & \text{operator} {==}(P\ p) & \text{const} \ \{\ \text{return} \end{tabular}
  \rightarrow tie(x,y)==tie(p.x,p.y); }
 P operator+(P p) const { return P(x+p.x, y+p.y); }
P operator-(P p) const { return P(x-p.x, y-p.y); }
P operator*(T d) const { return P(x*d, y*d); }
P operator/(T d) const { return P(x*d, y*d); }
P operator/(T d) const { return P(x*d, y*d); }
 T dot(P p) const { return x*p.x + y*p.y; }
 T cross(P'p) const { return x*p.y - y*p.x; }
 T cross(P a, P b) const { return
      (a-*this).cross(b-*this); }
 T dist2() const { return x*x + y*y;
 double dist() const { return sqrt((double)dist2()); }
 // angle to x-axis in interval [-pi, pi]
 double angle() const { return atan2(y, x); }
 P unit() const { return *this/dist(); } // makes
 P perp() const { return P(-y, x); } // rotates +90
  → degrees
 P normal() const { return perp().unit(); }
 // returns point rotated 'a' radians ccw around the
  → origin
 P rotate(double a) const {
  return P(x*cos(a)-y*sin(a),x*sin(a)+y*cos(a)); }
 friend ostream& operator<<(ostream& os, P p) {
  return os << "(" << p.x << "," << p.y << ")"; }</pre>
```

3.1.8 Segment Distance

Returns the shortest distance between point p and the line segment from point s to e.

3.1.9 Segment Intersection



```
Usage:
 * vector<P> inter = segInter(s1,e1,s2,e2);
 * if (sz(inter)==1)
     cout << "segments intersect at " << inter[0] <<
 * Status: stress-tested, tested on kattis:intersection
#include "Point.h"
#include "OnSegment.h"
template<class P> vector<P> seqInter(P a, P b, P c, P
 d) {
  auto oa = c.cross(d, a), ob = c.cross(d, b)
oc = a.cross(b, c), od = a.cross(b, d)
  // Checks if intersection is single non-endpoint
  if (sgn(oa) * sgn(ob) < 0 \&\& sgn(oc) * sgn(od) < 0)
    return {(a * ob - b * oa) / (ob - oa)};
  set<P> s;
  if (onSegment(c, d, a)) s.insert(a);
  if (onSegment(c, d, b)) s.insert(b);
  if (onSegment(a, b, c)) s.insert(c);
  if (onSegment(a, b, d)) s.insert(d);
  return {all(s)};
```

3.1.10 Side Of

res

Returns where p is as seen from s towards e. $1/0/-1 \Leftrightarrow$ left/on line/right. If the optional argument eps is given 0 is returned if p is within distance eps from the line. P is supposed to be Point<T> where T is e.g. double or long long. It uses products in intermediate steps so watch out for overflow if using int or long long.

3.2 3D 3.2.1 3D Convex Hull

```
#define ll long long
#define sz(x) ((int) (x).size())
#define all(x) (x).begin(), (x).end()
#define vi vector<int>
#define pii pair<int, int>
#define rep(i, a, b) for(int i = (a); i < (b); i++)
using namespace std;
template<typename T>
```

```
using minpq = priority queue<T, vector<T>, greater<T>>;
typedef long double ftype;
struct pt3 {
 ftype x, y, z;
 pt3(ftype x = 0, ftype y = 0, ftype z = 0) : x(x),
  \rightarrow y(y), z(z) {}
  pt3 operator-(const pt3 &o) const {
    return pt3(x - o.x, y - o.y, z - o.z);
  pt3 cross(const pt3 &o) const {
    return pt3(y * o.z - z * o.y, z * o.x - x * o.z, x
    \rightarrow * o.y - y * o.x);
  ftvpe dot(const pt3 &o) const {
    return x * 0.x + y * 0.y + z * 0.z;
// A face is represented by the indices of its three
   points a. b. c.
// It also stores an outward-facing normal vector q
struct face {
 int a, b, c;
 pt3 a:
// modify this depending on the coordinate sizes in

→ your use case

const ftype EPS = 1e-9;
vector<face> hull3(const vector<pt3> &p) {
 int n = sz(p);
  assert(n >= 3);
  vector<face> f;
  // Consider an edge (a->b) dead if it is not a CCW
     edge of some current face
  // If an edge is alive but not its reverse, this is
     an exposed edge.
  // We should add new faces on the exposed edges.
  vector<vector<bool>>> dead(n, vector<bool>(n, true));
  auto add face = [\&](int a, int b, int c)
   f.push_back({a, b, c, (p[b] - p[a]).cross(p[c] -
     → p[a])});
    dead[a][b] = dead[b][c] = dead[c][a] = false;
  // Initialize the convex hull of the first 3 points
  → as a
// triangular disk with two faces of opposite
      orientation
  add face(0, 1, 2);
  add face (0, 2, 1);
  rep(i, 3, n) {
    // f2 will be the list of faces invisible to the

→ added point p[i]

    vector<face> f2:
    for(face &F : f)
      if((p[i] - p[F.a]).dot(F.q) > EPS) {
        // this face is visible to the new point, so

→ mark its edges as dead

        dead[F.a][F.b] = dead[F.b][F.c] =
            dead[F.c][F.a] = true;
      }else {
        f2.push back(F);
    // Add a new face for each exposed edge.
    // Only check edges of alive faces for being

→ exposed.

    f.clear();
    for(face &F : f2) {
      int arr[3] = {F.a, F.b, F.c};
      rep(j, 0, 3)
        int a = arr[j], b = arr[(j + 1) % 3];
```

```
if(dead[b][a]) {
        add face(b, a, i);
 f.insert(f.end(), all(f2));
return f;
```

3.2.2 Point3D

Class to handle points in 3D space. T can be e.g. double or long long.

```
template<class T> struct Point3D {
 typedef Point3D P;
 typedef const P& R;
 T x, y, z;
 explicit Point3D(T x=0, T y=0, T z=0) : x(x), y(y),
    z(z) {}
 bool operator<(R p) const {</pre>
   return tie(x, y, z) < tie(p.x, p.y, p.z); }
 bool operator==(R p) const
   return tie(x, y, z) == tie(p.x, p.y, p.z); }
 P operator+(R p) const { return P(x+p.x, y+p.y,
    z+p.z);
 P operator (R p) const { return P(x-p.x, y-p.y,
  P operator*(T d) const { return P(x*d, y*d, z*d); } P operator/(T d) const { return P(x/d, y/d, z/d); }
   dot(R p) const \{ return x*p.x + y*p.y + z*p.z; \}
 P cross(R p) const {
   return P(y*p.z - z*p.y, z*p.x - x*p.z, x*p.y -
    \rightarrow y*p.x);
 T dist2() const { return x*x + y*y + z*z; }
 double dist() const { return sqrt((double)dist2()); } | }
 //Azimuthal angle (longitude) to x-axis in interval
 double phi() const { return atan2(y, x); }
 //Zenith angle (latitude) to the z-axis in interval
 double theta() const { return
  \rightarrow atan2(sqrt(x*x+y*y),z); }
 P unit() const { return *this/(T)dist(); } //makes
     dist()=1
 //returns unit vector normal to *this and p
 P normal(P p) const { return cross(p).unit(); }
 //returns point rotated 'angle' radians ccw around
 P rotate(double angle, P axis) const {
   double s = sin(angle), c = cos(angle); P u =
    → axis.unit();
   return u*dot(u)*(1-c) + (*this)*c - cross(u)*s;
```

3.2.3 Polyhedron Volume

Description: Magic formula for the volume of a polyhedron. Faces should point outwards.

```
template<class V, class L>
double signedPolyVolume(const V& p, const L& trilist) {
 double \bar{v} = 0;
  for (auto i : trilist) v +=

→ p[i.a].cross(p[i.b]).dot(p[i.c]);
  return v / 6;
```

3.2.4 Spherical Distance

(0 = north pole). All angles measured in radians. The algorithm starts find the tangents of a circle with a point set r2 to 0. by converting the spherical coordinates to cartesian coordinates so if that is what you have you can use only the two last rows. dx*radius is then the difference between the two points in the x direction and d*radius is the total distance between the points.

```
double sphericalDistance(double f1, double t1,
  double f2, double t2, double radius) {
double dx = sin(t2)*cos(f2) - sin(t1)*cos(f1);
double dy = sin(t2)*sin(f2) - sin(t1)*sin(f1);
  double dz = cos(t2) - cos(t1);
  double d = sqrt(dx*dx + dy*dy' + dz*dz);
  return radius*2*asin(d/2);
```

3.3 Circle

3.3.1 Circle Intersection

Computes the pair of points at which two circles intersect. Returns false in case of no intersection.

```
#include "Point.h"
typedef Point<double> P
bool circleInter(P a,P b,double r1,double r2,pair<P,</pre>
 → P>* out) {
  if (a == b) { assert(r1 != r2); return false; }
  P \text{ vec} = b - a;
  double d2 = vec.dist2(), sum = r1+r2, dif = r1-r2,
          p = (d2 + r1*r1 - r2*r2)/(d2*2), h2 = r1*r1

→ p*p*d2;

  if (sum*sum < d2 || dif*dif > d2) return false;
  P \text{ mid} = a + \text{vec*p, per} = \text{vec.perp}() * \text{sqrt}(\text{fmax}(0,
  *out = {mid + per, mid - per};
  return true:
```

3.3.2 Circle Polygon Intersection

Returns the area of the intersection of a circle with a ccw polygon. Time: O(n)

```
#include "Point.h"
typedef Point<double> P;
#define arg(p, q) atan2(p.cross(q), p.dot(q))
double circlePoly(P c, double r, vector<P> ps) {
  auto tri = [&](P p, P q) {
  auto r2 = r * r / 2;
    Pd = q - p;
    auto a = d.dot(p)/d.dist2(), b =
        (p.dist2()-r*r)/d.dist2();
    auto det = a * a - b;
    if (det \le 0) return arg(p, q) * r2;
    auto s = max(0., -a-sqrt(det)), t = min(1.,
     → -a+sqrt(det));
    if (t < 0 \mid | 1 \le s) return arg(p, q) * r2;
P u = p + d * s, v = p + d * t;
    return arg(p,u) * r2 + u.cross(v)/2 + arg(v.q) *
  auto sum = 0.0
  rep(i,0,sz(ps))
    sum += tri(ps[i] - c, ps[(i + 1) % sz(ps)] - c);
  return sum;
```

3.3.3 Circle Tangents

Finds the external tangents of two circles, or internal if r2 is negated. Can return 0, 1, or 2 tangents – 0 if one circle contains the other (or overlaps it, in the internal case, or if the circles are the same); 1 if the Returns the shortest distance on the sphere with radius radius be-circles are tangent to each other (in which case .first = .second and the tween the points with azimuthal angles (longitude) f1 (ϕ_1) and f2 (ϕ_2) tangent line is perpendicular to the line between the centers). first

from x axis and zenith angles (latitude) t1 (θ_1) and t2 (θ_2) from z axis and .second give the tangency points at circle 1 and 2 respectively. To

```
#include "Point.h"
template<class P>
vector<pair<P, P>> tangents(P c1, double r1, P c2,

→ double r2) {
 P d = c2 - c1
  double dr = r1 - r2, d2 = d.dist2(), h2 = d2 - dr *
  if (d2' == 0 | | h2 < 0) return {};
  vector<pair<P, P>> out;
 for (double sign : {-1, 1}) {
  P v = (d * dr + d.perp() * sqrt(h2) * sign) / d2;
  out.push_back({c1 + v * r1, c2 + v * r2});
  if (h2 == 0) out.pop back();
  return out;
```

3.3.4 CircumCircle

The circumcirle of a triangle is the circle intersecting all three vertices. ccRadius returns the radius of the circle going through points A, B and C and ccCenter returns the center of the same circle.



```
#include "Point.h"
typedef Point<double> P;
double ccRadius(const P& A, const P& B, const P& C) {
  return (B-A).dist()*(C-B).dist()*(A-C).dist()/
   abs((B-A).cross(C-A))/2;
  ccCenter(const P& A, const P& B, const P& C) {
  P b = C - A. c = B - A:

→ (b*c.dist2()-c*b.dist2()).perp()/b.cross(c)/2;
```

3.4 Polygon

3.4.1 Hull Diameter

Returns the two points with max distance on a convex hull (ccw, no duplicate/collinear points).

```
#include "Point.h"
typedef Point<ll> P;
array<P, 2> hullDiameter(vector<P> S) {
  int n = sz(S), j = n < 2 ? 0 : 1;
  pair<ll, array<P, 2>> res({0, {S[0], S[0]}});
  rep(i,0,j)
    for (;; j = (j + 1) % n) {
      res = \max(\text{res}, \{(S[i] - S[j]).dist2(), \{S[i],
      if ((S[(j+1) % n] - S[j]).cross(S[i+1] - S[j])
       \rightarrow S[i]) >= 0)
        break;
  return res.second;
```

3.4.2 Line Hull Intersection

Line-convex polygon intersection. The polygon must be ccw and have no collinear points. lineHull(line, poly) returns a pair describing the intersection of a line with the polygon:

(-1,-1) if no collision,

(i,-1) if touching the corner i, (i,i) if along side (i,i+1),

(i,j) if crossing sides (i,i+1) and (j,j+1). In the last case, if a corner i is crossed, this is treated as happening on side (i,i+1). The points are returned in the same order as the line hits the polygon. extrVertex returns the point of a hull with the max projection onto a line.

Time: $O(\log n)$

```
#include "Point.h"
#define cmp(i,j)

    sgn(dir.perp().cross(poly[(i)%n]-poly[(j)%n]))
#define extr(i) cmp(i + 1, i) \Rightarrow 0 && cmp(i, i - 1 +
template <class P> int extrVertex(vector<P>& poly, P

→ dir) {
  int n = sz(poly), lo = 0, hi = n;
  if (extr(0)) return 0;
  while (lo + 1 < hi)
    int \dot{m} = (lo + hi) / 2;
    if (extr(m)) return m;
int ls = cmp(lo + 1, lo), ms = cmp(m + 1, m);
    (ls < ms \mid | (ls == ms \&\& ls == cmp(lo, m)) ? hi :
     \rightarrow lo) = m:
  return lo;
#define cmpL(i) sqn(a.cross(poly[i], b))
template <class P>
array<int, 2> lineHull(P a, P b, vector<P>& poly) {
  int endA = extrVertex(poly, (a - b).perp());
  int endB = extrVertex(poly, (b - a).perp());
  if (cmpL(endA) < 0 \mid | cmpL(endB) > 0)
    return {-1, -1};
  array<int, 2> res;
  rep(1,0,2) {
    int lo = endB, hi = endA, n = sz(poly);
    while ((lo + 1) % n != hi) {
      int \dot{m} = ((lo' + hi + (lo' < hi ? 0 : n)) / 2) % n;
      (cmpL(m) = cmpL(endB) ? lo : hi) = m;
    res[i] = (lo + !cmpL(hi)) % n;
    swap(endA, endB);
  if (res[0] == res[1]) return {res[0], -1};
  if (!cmpL(res[0]) && !cmpL(res[1]))
    switch ((res[0] - res[1] + sz(poly) + 1) %
     \rightarrow sz(poly)) {
      case 0: return {res[0], res[0]};
case 2: return {res[1], res[1]};
  return res;
```

3.4.3 Polygon Center

Returns the center of mass for a polygon.

Time: O(n)

```
#include "Point.h"
typedef Point<double> P;
P polygonCenter(const vector<P>& v) {
  P res(0, 0); double A = 0;
  for (int i = 0, j = sz(v) - 1; i < sz(v); j = i++) {
    res = res + (v[i] + v[j]) * v[j].cross(v[i]);
}</pre>
```

```
A += v[j].cross(v[i]);
  return res / A / 3;
3.4.4 Polygon Cut
Returns a vector with the vertices of a polygon with ev-
erything to the left of the line going from s to e cut away.
   vector < P > p = ...;
    p = polygonCut(p, P(0,0), P(1,0));
 * Status: tested but not extensively
#include "Point.h"
#include "lineIntersection.h"
typedef Point<double> P;
vector<P> polygonCut(const vector<P>& poly, P s, P e) {
  vector<P> rés;
  rep(i,0,sz(poly)) {
    P cur = poly[i], prev = i ? poly[i-1] :

→ poly.back();

    bool side = s.cross(e, cur) < 0;</pre>
    if (side != (s.cross(e, prev) < 0))
       res.push back(lineInter(s, e, cur, prev).second);|;
      res.push back(cur);
  return res;
3.5 Closest Pair
Finds the closest pair of points.
Time: O(n \log n)
#include "Point.h"
typedef Point<ll> P;
pair<P, P> closest(vector<P> v) {
  assert(sz(v) > 1);
  set<P>
  sort(all(v), [](P a, P b) { return a.y < b.y; });</pre>
  pair<ll, pair<P, P>> ret{LLONG MAX, {P(), P()}};
  int j = 0;
  for (P p : v)
    P d{1 + (ll)sqrt(ret.first), 0};
    while (v[j].y \le p.y - d.x) S.erase(v[j++]);
    auto lo = S.lower bound(p - d), hi =

→ S.upper bound(p + d);

    for (; lo != hi; ++lo)
  ret = min(ret, {(*lo - p).dist2(), {*lo, p}});
    S.insert(p);
  return ret.second;
3.6 Convex Hull
typedef long long ll;
typedef pair <ll, ll> point;
#define x first
#define y second
inline ll area (point a, point b, point c) {
  return (b.x - a.x) * (c.y - a.y) - (b.y - a.y) *
   \leftarrow (c.x - a.x);
vector <point> convexHull (vector <point> p) {
  int n = p.size(), m = 0;
  if (n < 3) return p;
  vector <point> hull(n + n);
  sort(p.begin(), p.end());
  for (int i = 0; i < n; ++i)
    while (m > 1 \text{ and area}(hull[m - 2], hull[m - 1],
```

 \rightarrow p[i]) <= 0) --m;

3.7 Minimum Enclosing Circle

```
Expected runtime: O(n)
// Solves Gym 102299J
using namespace std;
typedef long double ld;
typedef pair <ld, ld> point;
#define x first
#define y second
point operator + (const point &a, const point &b) {
  return point(a.x + b.x, a.y + b.y);
point operator - (const point &a, const point &b) {
  return point(a.x - b.x, a.y - b.y);
point operator * (const point &a, const ld &b) {
  return point(a.x * b, a.y * b);
point operator / (const point &a, const ld &b) {
  return point(a.x / b, a.y / b);
const ld EPS = 1e-8:
const ld INF = 1e20
const ld PI = acosl(-1);
inline ld dist (point a, point b) {
  return hypotl(a.x - b.x, a.y - b.y);
inline ld sqDist (point a, point b) {
  return (a.x - b.x) * (a.x - b.x) + (a.y - b.y) *
   \rightarrow (a.y - b.y);
inline ld dot (point a, point b) {
  return a.x * b.x + a.y * b.y;
inline ld cross (point a, point b) {
  return a.x * b.y - a.y * b.x;
inline ld cross (point a, point b, point c) {
  return cross(b - a, c - a);
inline point perp (point a) {
  return point(-a.y, a.x);
// circle through 3 points
pair <point, ld> getCircle (point a, point b, point c)
  pair <point. ld> ret:
  ld den = (ld) 2 * cross(a, b, c);
ret.x.x = ((c.y - a.y) * (dot(b, b) - dot(a, a)) -
     (b.y - a.y) * (dot(c, c) - dot(a, a))) / den;
  ret.x.y = ((b.x - a.x) * (dot(c, c) - dot(a, a)) -
      (c.x - a.x) * (dot(b, b) - dot(a, a))) / den;
  ret.\dot{y} = dist(ret.x, a);
  return ret;
pair <point, ld> minCircleAux (vector <point> &s,

→ point a, point b, int n) {
```

```
ld lo = -INF, hi = INF;
  for (int i = 0; i < n; ++i) {
  auto si = cross(b - a, s[i] - a);</pre>
    if (fabs(si) < EPS) continue;</pre>
    point m = getCircle(a, b, s[i]).x;
    auto cr = cross(b - a, m - a);
    si < 0 ? hi = min(hi, cr) : lo = max(lo, cr);
  1d v = 0 < lo ? lo : hi < 0 ? hi : 0
  point c = (a + b) * 0.5 + perp(b - a) * v /

    sqDist(a, b);

  return {c, sqDist(a, c)};
pair <point, ld> minCircle (vector <point> &s, point
- a, int n) {
  random_shuffle(s.begin(), s.begin() + n);
  point \bar{b} = s[0], c = (a + b) * \bar{0}.5;
  ld r = sqDist(a, c);
  for (int i = 1; i < n; ++i) {
  if (sqDist(s[i], c) > r * (1 + EPS))
      tie(c, r) = n == s.size() ? minCircle(s, s[i],

→ i) : minCircleAux(s, a, s[i], i);
  return {c, r};
pair <point, ld> minCircle (vector <point> s) {
  assert(!s.empty());
  if (s.size() == 1) return {s[0], 0};
  return minCircle(s, s[0], s.size());
int n; vector <point> p;
int main() {
  cin >> n;
  while (n--) {
    double x, y;
scanf("%lf %lf", &x, &y);
    p.emplace back(x, y);
  pair <point, ld> circ = minCircle(p);
  printf("%0.12f %0.12f %0.12f\n", (double) circ.x.x,
  return 0:
```

```
3.8 Point In Polygon
// Test if a point is inside a convex polygon in O(lg
    n) time
// Solves SPOJ INOROUT typedef long long ll;
typedef pair <ll, ll> point;
#define x first
#define y second
struct seament {
  point P1, P2;
  segment () {}
  segment (point P1, point P2) : P1(P1), P2(P2) {}
inline ll ccw (point A, point B, point C) {
  return (B.x - A.x) * (C.y - A.y) - (C.x - A.x) *
   \hookrightarrow (B.y - A.y);
inline bool pointOnSegment (segment S, point P) {
  ll x = P.x, y = P.y, x1 = S.P1.x, y1 = S.P1.y, x2 =
   \rightarrow S.P2.x, y2 = S.P2.y;
  ll a = x - x1, b = y - y1, c = x2 - x1, d = y2 - y1,
      dot = a * c + b * d, len = c * c + d * d:
  if (x1 == x2 \text{ and } y1 == y2) return x1 == x and y1 == y2
  if (dot < 0 or dot > len) return 0;
```

```
return x1 * len + dot * c == x * len and y1 * len +

→ dot * d == v * len:

const int M = 17
const int N = 10010;
|struct polygon {
  int n; //_n > 1
  point p[N]; // clockwise order
  polygon () {}
  polygon (int n, point *T) {
    for (int i = 0; i < n; ++i) p[i] = T[i];
  bool contains (point P, bool strictlyInside) {
    int lo = 1, hi = n - 1;
    while (lo < hi){
      int mid = lo + hi >> 1
      if (ccw(p[0], P, p[mid]) > 0) lo = mid + 1;
      else hi = mid:
    if (ccw(p[0], P, p[lo]) > 0) lo = 1;
    if (!strictlyInside and
        pointOnSegment(segment(p[0], p[n - 1]), P))
        return 1;
    if (!strictlyInside and
        pointOnSegment(segment(p[lo], p[lo - 1]), P))
        return 1:
    if (lo == 1 or ccw(p[0], P, p[n - 1]) == 0) return
    return ccw(p[lo], P, p[lo - 1]) < 0;
3.9 near_pair
struct pt {
 int x, y, id;
struct cmp x {
  bool operator()(const pt & a, const pt & b) const {
    return a.x < b.x \mid | (a.x == b.x \&\& a.y < b.y);
struct cmp_y {
  bool operator()(const pt & a, const pt & b) const {
    return a.y < b.y;
|vector<pt> a;
double mindist;
|pair<int, int> best pair;
void upd ans(const pt & a, const pt & b) {
  double dist = sqrt((a.x - b.x)*(a.x - b.x) + (a.y - b.x)
      b.y)*(a.y - b.y));
  if (dist < mindist) {</pre>
    mindist = dist;
    best pair = {a.id, b.id};
vector<pt> t;
|void rec(int l, int r) {
  if (r - l <= 3) {
    for (int i = l; i < r; ++i) {
      for (int j = i + 1; j < r; ++j) {
        upd ans(a[i], a[j]);
    sort(a.begin() + l, a.begin() + r, cmp y());
```

return;

```
int m = (l + r) >> 1;
  int midx = a[m].x;
  rec(l, m);
  rec(m, r);
  merge(a.begin() + l, a.begin() + m, a.begin() + m,
  \rightarrow a.begin() + r, t.begin(), cmp y());
  copy(t.begin(), t.begin() + r - l, a.begin() + l);
  for (int i = l; i < r; ++i)
    if (abs(a[i].x - midx) < mindist)</pre>
      for (int j = tsz - 1; j >= 0 \&\& a[i].y - t[j].y
      t[tsz++] = a[i];
void solve(int n)
  t.resize(n)
  sort(a.begin(), a.end(), cmp x());
  mindist = 1E20:
  rec(0, n);
3.10 sweep
const double EPS = 1E-9;
struct pt {
  double x, y;
struct seg {
  pt p, q;
  int id;
  double get y(double x) const {
    if (abs(\overline{p}.x - q.x) < EPS)
      return p.y;
    return p.y + (q.y - p.y) * (x - p.x) / (q.x - p.x);
bool intersect1d(double l1, double r1, double l2,
   double r2) {
 if (l1 > r1)
    swap(l1, r1);
  if (l2 > r2)
    swap(l2, r2);
  return \max(l1, l2) \ll \min(r1, r2) + EPS;
int vec(const pt& a, const pt& b, const pt& c) {
  double s = (b.x - a.x) * (c.y - a.y) - (b.y - a.y) *
   \leftarrow (c.x - a.x);
  return abs(s) < EPS ? 0 : s > 0 ? +1 : -1;
|bool intersect(const seg& a, const seg& b)
  return intersect1d(a.p.x, a.q.x, b.p.x, b.q.x) &&
    intersect1d(a.p.y, a.q.y, b.p.y, b.q.y) &&
    vec(a.p, a.q, b.p) * vec(a.p, a.q, b.q) <= 0 &&
    vec(b.p, b.q, a.p) * vec(b.p, b.q, a.q) <= 0;
|bool operator<(const seg& a, const seg& b)
  double x = max(min(a.p.x, a.q.x), min(b.p.x, b.q.x));
  return a.get y(x) < b.get y(x) - EPS;
```

struct event {

int tp, id;

event() {}

double x;

```
event(double x, int tp, int id) : x(x), tp(tp),
  → id(id) {}
  bool operator<(const event& e) const {
    if (abs(x - e.x) > EPS)
      return x < e.x;
    return tp > e.tp;
set<seg> s;
vector<set<seg>::iterator> where;
set<seg>::iterator prev(set<seg>::iterator it) {
 return it == s.begin() ? s.end() : --it;
set<seg>::iterator next(set<seg>::iterator it) {
pair<int, int> solve(const vector<seg>& a) {
  int n = (int)a.size();
  vector<event> e:
  for (int i = 0; i < n; ++i) {
    e.push back(event(min(a[i].p.x, a[i].q.x), +1, i));
    e.push_back(event(max(a[i].p.x, a[i].q.x), -1, i));
  sort(e.begin(), e.end());
  s.clear();
  where resize(a.size());
  for (size t i = 0; i < e.size(); ++i) {
    int id = e[i].id;
    if (e[i].tp = +1) {
      set<seg>::iterator nxt = s.lower bound(a[id]),
      → prv = prev(nxt);
     if (nxt != s.end() && intersect(*nxt, a[id]))
return make_pair(nxt->id, id);
      if (prv != s.end() && intersect(*prv, a[id]))
        return make pair(prv->id, id);
      where[id] = s_insert(nxt, a[id]);
      set<seg>::iterator nxt = next(where[id]), prv =
      → prev(where[id]);
      if (nxt != s.end() && prv != s.end() &&

    intersect(*nxt, *prv))

        return make pair(prv->id, nxt->id);
      s.erase(where[id]);
  return make pair(-1, -1);
```

4 Graph

4.1 2Sat

```
namespace sat{
 const int MAX = 200010;
 bool vis[MAX];
 vector <int> ed[MAX], rev[MAX]
 int n, m, ptr, dfs t[MAX], ord[MAX], par[MAX];
 inline int inv(int x)
   return ((x) \le n ? (x + n) : (x - n));
 void init(int vars){
   n = vars, m = vars << 1;
   for (int i = 1; i \le m; i++){
     ed[i].clear();
rev[i].clear();
 inline void add(int a, int b){
   ed[a].push back(b);
   rev[b].push back(a);
```

```
inline void OR(int a, int b){
  add(inv(a), b);
add(inv(b), a);
inline void AND(int a, int b){
  add(a, b);
  add(b, a);
void XOR(int a,int b){
  add(inv(b), a);
  add(a, inv(b));
  add(inv(a), b);
  add(b, inv(a));
inline void XNOR(int a. int b){
  add(a,b);
  add(b,a)
  add(inv(a), inv(b));
add(inv(b), inv(a));
inline void force true(int x){
  add(inv(x), x);
inline void topsort(int s){
  vis[s] = true;
  for(int x : rev[s]) if(!vis[x]) topsort(x);
  dfs t[s] = ++ptr;
inline void dfs(int s, int p){
  par[s] = p;
  vis[s] = true;
  for(int x : ed[s]) if (!vis[x]) dfs(x, p);
void build(){
  CLR(vis);
ptr = 0;
  for(int i=m;i>=1;i--)
    if (!vis[i]) topsort(i);
    ord[dfs t[i]] = i;
  ĆLR(vis);
  for (int i = m; i >= 1; i --) {
    int x = ord[i];
    if (!vis[x]) dfs(x, x);
bool satisfy(vector < int>& res){
  build():
  CLR(vis);
  for (int i = 1; i \le m; i++){
    int x = ord[i];
    if (par[x] == par[inv(x)]) return false;
    if (!vis[par[x]]){
      vis[par[x]] = true;
      vis[par[inv(x)]] = false;
  res.clear();
  for (int i = 1; i \le n; i++){
    if (vis[par[i]]) res.push back(i);
  return true:
```

4.2 Centroid decomp

```
vector <int> ed[MAX]
bool isCentroid[MAX]
int sub[MAX], cpar[MAX], clevel[MAX];
int dis[20][MAX];
lvoid calcSubTree(int s,int p) {
```

```
sub[s] = 1;
  for(int x : ed[s]) {
    if(x == p or isCentroid[x]) continue;
    calcSubTree(x,s);
    sub[s] += sub[x];
int nn:
int getCentroid(int s,int p) {
  for(int x : ed[s]) {
    if(!isCentroid[x] && x!=p && sub[x]>(nn/2)) return

→ getCentroid(x,s);

  return s;
void setDis(int s, int from, int p, int lev) {
  dis[from][s] = lev;
  for(int x : ed[s]) -
    if(x == p or isCentroid[x] ) continue;
    setDis(x, from, s, lev+1);
void decompose(int s,int p,int lev) {
  calcSubTree(s,p);
  nn = sub[s];
  int c = getCentroid(s,p);
  setDis(c,lev,p,0);
  isCentroid[c] = true;
  cpar[c] = p;
  clevel[c] = lev;
  for(int x : ed[c])
    if(!isCentroid[x]) decompose(x,c,lev+1);
int ans[MAX];
inline void update(int v) {
  int u = v;
  while(u!=-1) {
    ans[u] = min(ans[u], dis[clevel[u]][v]);
    u = cpar[u];
inline int query(int v) {
  int ret = INF;
  int u = v;
  while(u != -1) {
    ret = min(ret, dis[clevel[u]][v]+ans[u]);
    u = cpar[u];
  return ret;
int main()
  decompose(1,-1,0);
  for(int i=1; i<=n; i++) ans[i] = INF;</pre>
  update(v);
  query(v));
  return 0;
```

4.3 articulation_point

```
using namespace std;
const int N = 1e5 + 10;
vector<int> g[N];
void dfs(int u, int p) {
 low[u] = vis[u] = ++now; int ch = 0;
 for(int v : q[u]){
   if(v ^ p)
    if(vis[v]) low[u] = min(low[u], vis[v]);
```

```
else {
         ch++; dfs(v, u);
low[u] = min(low[u], low[v]);
if(p + 1 && low[v] >= vis[u]) cut[u] = 1;
                                                                  struct edges {
                                                                 int u,v;
ara[MAXE];
         if(low[v] > vis[u]) {
  printf("Bridge %d -- %d\n", u, v);
                                                                  vector <int> ed[MAXN];
                                                                  vector <int> isBridge[MAXN];
                                                                  |vector <int> brTree[MAXN];
                                                                  bool vis[MAXN];
                                                                  int st[MAXN], low[MAXN], Time = 0;
  } if(p == -1 && ch > 1) cut[u] = 1;
                                                                  int cnum;
                                                                  |int comp[MAXN];
void ArticulationPointAndBridge() {
                                                                  void findBridge(int s,int par) {
  now = 0;
                                                                    int i,x,child = 0,j;
  for(int i = 0; i < n; i++) {
                                                                    \underline{v}is[s] = 1;
    if(!vis[i]) dfs(i, -1);
                                                                    Time++;

st[s] = low[s] = Time;

for(i=0; i<ed[s].size(); i++) {
                                                                      x = ed[s][i]
4.4 bcc
                                                                      if(!vis[x]) {
// clear ed[] every test case
                                                                         child++;
// tot -> total number of components
                                                                         findBridge(x,s);
// bcc[i] contains the nodes of the i'th component
                                                                         low[s] = min(low[s],low[x]);
// any self loop or multiple edge?
                                                                         if(low[x] > st[s])
                                                                           isBridge[s][i] = 1
const int MAX = ?:
vector <int> ed[MÁX];
                                                                           j = lower bound(ed[x].begin(),ed[x].end(),s)-e
bool cut[MAX];
                                                                               d[x].begin();
int tot, Time, low[MAX], st[MAX];
                                                                           isBridge[x][j] = 1;
vector <int> bcc[MAX];
stack <int> S;
                                                                      else if(par!=x)
void popBCC(int s,int x) {
  cut[s] = 1;
                                                                         low[s] = min(low[s], st[x]);
  bcc[++tot].pb(s);
  while(bcc[tot]_back() ^ x) {
    bcc[tot].pb(S.top());
                                                                  |void dfs(int s) {
    S.pop();
                                                                    int i,x;
vis[s] = 1;
                                                                    comp[s] = cnum;
                                                                    for(i=0; i<ed[s].size(); i++) {</pre>
void dfs(int s, int p = -1) {
                                                                      if(!isBridge[s][i]) {
  S.push(s);
                                                                         x = ed[s][i];
if(!vis[x]) dfs(x);
  int ch = 0;
  st[s] = low[s] = ++Time;
  for(int x : ed[s]) {
    if(!st[x]) {
      dfs(x,s);
low[s] = min(low[s],low[x]);
                                                                  void processBridge(int n,int m) {
                                                                    CLR(vis);
       if(p != -1 \text{ and } low[x] >= st[s]) popBCC(s,x);
                                                                    Time = 0
       else if(p == -1) if(ch > 1) popBCC(s,x);
                                                                    for(int i=1; i<=n; i++) if(!vis[i]) findBridge(i,-1);</pre>
    else if(p != x) low[s] = min(low[s],st[x]);
                                                                    CLR(vis);
                                                                    for(int i=1; i<=n; i++) {
                                                                      if(!vis[i]) { cnum++;
  if(p == -1 \&\& ch > 1) cut[s] = 1;
                                                                         dfs(i);
void processBCC(int n) {
  for(int i=1;i<=n;i++) bcc[i].clear();
CLR(st); CLR(cut);</pre>
                                                                    n = cnum;
  Time = tot = 0;
  for(int i=1; i<=n; i++) {
   if(!st[i]) {</pre>
                                                                    for(int i=1; i<=m; i++) {
                                                                      if(comp[ara[i].u] != comp[ara[i].v]) {
    brTree[comp[ara[i].u]].pb(comp[ara[i].v]);
       dfs(i,-1);
                                                                         brTree[comp[ara[i].v]].pb(comp[ara[i].u]);
       if(!S.empty()) ++tot;
       while(!S.empty()) {
                                                                    }
         bcc[tot].push back(S.top());
         S.pop();
                                                                  |int main() {
                                                                    int n,m,u,v;
scanf("%d %d",&n,&m);
for(int i=1; i<=m; i++) {</pre>
                                                                      sii(u,v);
4.5 bridge_tree
                                                                      ed[u].pb(v);
const int MAXN = ?;
```

const int MAXE = ?;

ed[v].pb(u);

```
ara[i].v = v;
 for(int i=1; i<=n; i++) sort(all(ed[i]));</pre>
 processBridge(n,m);
 return 0;
4.6 dinic
namespace dinic {
 using T = int;
 const T INF = 0x3f3f3f3f;
 const int MAXN = 5010;
 int n, src, snk, work[MAXN];
T dist[MAXN];
 struct Edge{
    int to, rev pos;
   T c, f;
 vector <Edge> ed[MAXN];
 void init(int n, int src, int snk) {
   n = n, src = src, snk = snk;
    for(int i=1;i<=n;i++) ed[i].clear();</pre>
 inline void addEdge(int u, int v, T c, T rc = 0) {
    Edge a = \{v, (int)ed[v].size(), c, 0\};
    Edge b = \{u, (int)ed[u].size(), rc, 0\};
    ed[u].push back(a);
    ed[v].push back(b);
 bool dinic bfs() {
    SET(dist);
dist[src] = 0;
    queue <int> q;
    q.push(src);
    while(!q.empty()){
      int u = q.front(); q.pop();
      for(Edge &e : ed[u]){
        if(dist[e.to] == -1 \text{ and } e.f < e.c) 
          dist[e.to] = dist[u] + 1;
          q.push(e.to);
    return (dist[snk]>=0);
 T dinic dfs(int u, T fl){
    if (u == snk) return fl;
    for (; work[u] < (int)ed[u].size(); work[u]++) { Edge &e = ed[u][work[u]];
      if (e.c <= e.f) continue;</pre>
      int v = e.to:
      if (dist[v] == dist[u] + 1)
        T df = dinic dfs(v, min(fl, e.c - e.f));
        if (df > 0){
          e.f += df;
          ed[v][e.rev pos].f -= df;
          return df;
    return 0;
 T solve() {
    T ret = 0:
    while (dinic bfs()) {
      CLR(work);
```

isBridge[u].pb(0);

isBridge[v].pb(0);

ara[i].u = u;

```
int n, m, u, v, c;
  cin >> n >> m:
  dinic::init(n, 1, n);
  while(m--) {
    cin >> u >> v >> c:
    dinic::addEdge(u, v, c, c);
  cout << dinic::solve() << '\n';
  return 0;
4.7 dsu on tree
void calcSubSize(int s,int p) {
  sub[s] = 1;
for(int x : G[s]) {
    if(x==p) continue;
    calcSubSize(x,s);
    sub[s] += sub[x];
void add(int s,int p,int v,int bigchild = -1) {
  freq[color[s]] += v;
  for(int x : G[s]) {
    if(x==p || x==bigchild) continue;
    add(x,s,v);
void dfs(int s,int p,bool keep) {
  int bigChild = -1;
  for(int x : G[s]) {
    if(x==p) continue;
    if(bigChild==-1 || sub[bigChild] < sub[x] )</pre>

→ bigChild = x;

  for(int x : G[s]) {
    if(x==p || x==bigChild) continue;
    dfs(x,s,0);
  if(bigChild!=-1) dfs(bigChild,s,1);
  add(s,p,1,bigChild);
  if(keep==0)
    add(s,p,-1);
4.8 euler_path
vector <int> ed[MAX+5], sltn;
int inDeg[MAX+5], outDeg[MAX+5];
bool vis[MAX+5];
void dfs(int nd) {
  vis[nd] = true;
  while(ed[nd].size())
    int v = ed[nd].back();
    ed[nd].pop_back();
    dfs(v);
  sltn.pb(nd);
int findEuler (int n) {
  int src , snk , ret = 1;
bool found_src = false, found_snk = false;
  CLR(inDeg); CLR(outDeg);
  for(int u = 1; u <= n; u++) {
  for(int i = 0; i<ed[u].size(); i++) {</pre>
      int v = ed[u][i];
```

while (T delta = dinic dfs(src, INF)) ret +=

→ delta;

return ret;

int main() {

```
outDeg[u]++;
       inDeg[v]++;
  int diff;
  for(int i = 1; i<=n; i++) {
  diff = outDeg[i] - inDeg[i];</pre>
    if(diff == 1)
       if(found_src) return 0;
       found sr\overline{c} = true;
       src = i;
    else if (diff == -1) {
       if(found snk) return 0;
       found sn\overline{k} = true;
       snk = i;
    else if(diff != 0) return 0;
  if(!found src) {
     ret = 2;
    for(int i = 1; i <= n; i++) {
  if( outDeg[i] ) {</pre>
         found src = true;
         src = i;
         break;
  if(!found src) return ret;
  CLR(vis);
  sltn.clear();
  dfs(src);
  for(int i = 1; i<=n; i++) {
    if(outDeg[i] && !vis[i]) return 0;
  for(int i = (int)sltn.size()-1; i>=0; i--)

¬ printf("%d ",sltn[i]);

  puts'("");
  return ret:
4.9 hld
const int N = 3e4 + 5;
vector<int> G[N];
int sz[N], H[N], P[N];
void dfs(int cur, int h)
     sz[cur] = 1;
    H[cur] = h;
     for(int& to : G[cur])
         G[to].erase(find(G[to].begin(), G[to].end(),
            cur));
         P[to] = cur;
         dfs(to, h + 1)
         sz[cur] += sz[to];
         if(sz[to] > sz[G[cur][0]]) swap(G[cur][0], to);
int base[N], pos[N], head[N];
int ptr = 0;
|void hld(int cur)
     pos[cur] = ++ptr;
     base[ptr] = cur;
    for(int to : G[cur]) {
    head[to] = (to == G[cur][0] ? head[cur] : to);
         hld(to):
|segtree ST;
|int query(int u, int v)
```

```
11
    int ret = 0;
    while(head[u] != head[v])
        if(H[head[u]] > H[head[v]]) swap(u, v);
        ret += ST.query(pos[head[v]], pos[v]);
        v = P[head[v]];
    if(H[u] > H[v]) swap(u, v);
ret += ST.query(pos[u], pos[v]);
    return ret;
void update(int u, int val) {
    ST.update(pos[u], val);
void build(int n, int root)
    ptr = 0;
    dfs(root, 0);
head[root] = root;
    hld(root);
    ST = segtree(n);
* clear graph
   call build
 * Prob : sum of values from u to v
4.10 hopcroft_karp
struct HopcroftKarp {
  const int N, M;
  std::vector<std::vector<int>> adj left;
  std::vector<int> matchL, matchR;
  HopcroftKarp(int N, int M, const

    std::vector<std::pair<int, int>>& edge)
      : N(N), M(M), matchL(N, -1), matchR(M, -1),
       → adj left(N) {
    for (auto [l, r] : edge)
      adj left[l].push back(r);
  int maxmatching() {
    int sz = 0;
    for (bool updated = true; updated;) {
      updated = false;
```

static std::vector<int> root(N), prev(N), qq(N);

qq[q]++] = i, root[i] = i, prev[i] = i; q.push(i), root[i] = i;

matchR[v] = u, std::swap(matchL[u], v),

v = matchR[v], prev[v] = u, root[v] =

std::fill(root.begin(), root.end(), -1),
std::fill(prev.begin(), prev.end(), -1);

// int u = q.front(); q.pop();
if (matchL[root[u]] != -1) continue;
for (int v : adj_left[u]) {

static int qi, qj;

while (qi < qj) {

break;

int u = qq[qi++];

qi = qj = 0;

// std::queue<int> q;

for (int i = 0; i < N; i++)

if (matchR[v] = -1) {

 \rightarrow u = prev[u]:

updated = true, sz++;

if (prev[matchR[v]] == -1)

 \rightarrow root[u], qq[qj++] = v;

while (v != -1)

if (matchL[i] == -1)

 \rightarrow root[u], q.push(v);

// v = matchR[v], prev[v] = u, root[v] =

```
12
```

```
return sz;
4.11 hungarian
// Given NN matrix A[i][j]. Calculate a permutation
\rightarrow p[i] that minimize A[i][p[i]].
template <typename T>
pair <T, vector <int>> Hungarian (int n, int m, T
vector <T> v(m), dist(m);
  vector <int> L(n, -1), R(m, -1);
vector <int> index(m), prev(m);
  auto residue = [&] (int i, int j) {return c[i][j] -
   → v[j];};
  iota(index.begin(), index.end(), 0);
  for (int f = 0; f < n; ++f) {
    for (int j = 0; j < m; ++j) {
       dist[j] = residue(f, j), prev[j] = f;
    \bar{T} w; int i, j, l, s = 0, t = 0;
    while (true) {
       if (s == t) {
         l = s, w = dist[index[t++]];
         for (int k = t; k < m; ++k) {
   j = index[k]; T h = dist[j];</pre>
           if (h <= w) {
   if (h < w) t = s, w = h;
   index[k] = index[t], index[t++] = j;</pre>
         for (int k = s; k < t; ++k) {
           j = index[k];
           if (R[j] < 0) goto augment;
       int q = index[s++], i = R[q];
       for (int k = t; k < m; ++k) {
           = index[k];
         \hat{T} h = residue(i, j) - residue(i, q) + w;
         if (h < dist[j]) {
           dist[j] = h, prev[j] = i;
           if (h == w) {
             if (R[j] < 0) goto augment;</pre>
             index[k] = index[t], index[t++] = j;
  augment:
    for (int k = 0; k < 1; ++k) v[index[k]] +=
       dist[index[k]] - w;
      R[j] = i = prev[j], swap(j, L[i]);
    } while (i ^ f);
  for (int i = 0; i < n; ++i) ret += c[i][L[i]];
return {ret, L};</pre>
```

4.12 kuhn

```
namespace bpm{
  const int L = 105;
  const int R = 105;
  vector <int> G[L];
```

```
int matchR[R], matchL[L], vis[L], it;
  void init(int n) -
    SET(matchL), SET(matchR), CLR(vis);
    for(int i=1;i<=n;i++) G[i].clear();</pre>
 inline void addEdge(int u,int v) { G[u].pb(v); }
  bool dfs(int s) {
    vis[s] = it;
    for(auto x : G[s]) {
   if( matchR[x] == -1 or (vis[matchR[x]] != it and

    dfs(matchR[x])))

        matchL[s] = x; matchR[x] = s;
        return true;
    return false;
  int solve() {
    int cnt = 0:
    for(int i=1;i<=n;i++) {</pre>
      if(dfs(i)) cnt++, it++;
    return cnt;
4.13 lca
// Don't forget to clear ed after test case ends(vt,

→ cost are cleared inside)
```

```
using namespace std;
const int MAX = 100010:
int dep[MAX], par[MAX][21];
|vector <int> ed[MAX];
void dfs(int s, int p, int d) {
  dep[s] = d, par[s][0] = p;
  for(int x : ed[s]) {
    if(x == p) continue;
    dfs(x, s, d+1);
void preprocess(int root, int n) {
  LG = \lg(n);
  memset(par, -1, sizeof(par));
  dfs(root, -1, 0);
  for(int j=1; j<=LG; j++) {</pre>
    for(int i=1;i<=n;i++)
      if(par[i][j-1] != -1) par[i][j] =
       → par[par[i][j-1]][j-1];
  }
|int getLCA(int u, int v) {
  if(dep[u] < dep[v]) swap(u, v);</pre>
  for(int i=LG;i>=0;i--) {
    if(dep[u] - (1 << i) >= dep[v]) u = par[u][i];
  if(u == v) return u;
  for(int i=LG;i>=0;i--) {
    if (par[u][i] != -1  and par[u][i] - par[v][i]) {
      u = par[u][i], v = par[v][i];
  return par[u][0];
```

```
4.14 manhattan_MST
using namespace std;
using ll = long long;
```

```
struct UnionFind {
    vector<int> UF;
    int cnt;
    , description find(UF[v]); }
    bool join(int v, int w) {
        if ((v = find(v)) == (w = find(w))) return
        if (UF[v] > UF[w]) swap(v, w);
        UF[v] += UF[w];
UF[w] = v;
        return true;
    bool connected(int v, int w) {
        return find(v) == find(w);
    int getSize(int v) { return -UF[find(v)]; }
template <class T>
|struct KruskalMST {
    using Edge = tuple<int, int, T>;
    T mstWeight;
    vector<Edge> mstEdges;
    UnionFind uf;
    KruskalMST(int V, vector<Edge> edges) :

→ mstWeight(),

  uf(V) {
        sort(edges.begin(), edges.end(), [&](const

→ Edge &a, , const Edge &b)

            return qet<2>(a) < qet<2>(b);
        });
for (auto &&e : edges) {
            if (int(mstEdges.size()) >= V - 1) break;
            if (uf.join(get<0>(e), get<1>(e))) {
                mstEdges.push back(e);
                mstWeight += \overline{g}et < 2 > (e);
        }
    }
template <class T>
struct ManhattanMST : public,⊟ KruskalMST<T> {
    using Edge = typename KruskalMST<T>::Edge;
    static vector<Edge>

¬ generateCandidates(vector<pair<T, T>>,, P) {
        vector<int> id(P.size());
iota(id.begin(), id.end(), 0);
        vector<Edge> ret;
        ret.reserve(P.size() * 4);
        for (int h = 0; h < 4; h++)
            sort(id.begin(), id.end(), [&](int i, int

→ j) {
                return P[i].first - P[j].first <</pre>
                 → P[i].second - P[i].second:
            map<T, int> M;
            for (int i : id) {
                auto it = M.lower bound(-P[i].second);
                for (; it != M.en\overline{d}(); it =

→ M.erase(it)) {
                     int j = it->second;
                    T dx = P[i].first - P[j].first, dy
                     → = P[i].second, - P[j].second;
                    if (dy > dx) break;
                     ret.emplace back(i, j, dx + dy);
                M[-P[i].second] = i;
```

```
q.push(v);
                                                                                                                        long long m = (a * n + b) / c;
            for (auto &&p : P) {
                                                                                                                        return m * n - FloorSumAP(c, c - b - 1, a, m - 1);
                 if (h % 2)
                     p.first = -p.first;
                                                                                                                      \rightarrow floro(ax + b / c)^2 = h
                     swap(p.first, p.second);
                                                               return (dis[snk] != INF);
                                                                                                                      struct dat {
                                                                                                                        long long f, g, h;
dat(long long f = 0, long long g = 0, long long h =
                                                             inline pair <T, T> solve() {
        return ret;
                                                               T F = 0, C = 0, f;
                                                                                                                        \rightarrow 0) : f(f), g(g), h(h) {};
                                                               int u, v;
    ManhattanMST(const vector<pair<T, T>> &P)
                                                               while(SPFA()){
                                                                                                                      long long mul(long long a, long long b){
        : KruskalMST<T>(P.size(),
                                                                 u = snk:
                                                                                                                        return (a * b) % MOD;

    generateCandidates(P)) {}

                                                                 f = mCap[u];
                                                                 F += f:
int main() {
                                                                                                                     |dat query(long long a, long long b, long long c, long
                                                                 while(u!=src){
    int N;
                                                                                                                      \rightarrow long n){
                                                                   v = par[u];
    cin >> N:
                                                                                                                        if(!a) return {mul(n + 1, b / c), mul(mul(mul(b / c,
                                                                   ed[v][pos[u]].flow += f; // edge of v-->u
    vector<pair<ll, ll>> P(N);
                                                                                                                         \rightarrow n), n + 1), inv2), mul(mul(n + 1, b / c), b /c)};
    for (auto &&p : P) cin >> p.first >> p.second;
                                                                   ed[u][ed[v][pos[u]].rev_pos].flow -= f;
                                                                                                                        long long f, g, h;
    ManhattanMST mst(P)
                                                                                                                        dat nxt;
    cout << mst.mstWeight << '\n';</pre>
                                                                                                                        if(a >= c \text{ or } b >= c){}
    for (auto &&[v, w, weight] : mst.mstEdges) cout <<</pre>
                                                                 \dot{C} += dis[snk] * f;
                                                                                                                          nxt = query(a % c, b % c, c, n);

    v << □ , □ << w << '\n'; return 0;
</pre>
                                                                                                                          f = (nxt.f + mul(mul(mul(n, n + 1), inv2), a / c)
                                                               return make pair(F,C);
                                                                                                                          g = (nxt.g + mul(n + 1, b / c)) % MOD;
4.15 mcmf
                                                                                                                              n + 1, inv6))) + mul(mul(b / c, mul(n, n +
namespace mcmf {
                                                                                                                              1)), inv2)) % MOD;
                                                           5 Math
  using T = int;
                                                                                                                          h = (nxt.h + 2 * mul(b / c, nxt.f) + 2 * mul(a / c)
                                                           5.1 FWHT
  const T INF = ?; // 0x3f3f3f3f or
                                                                                                                              c, nxt.q) + mul(mul(a / c, a / c), mul(mul(n,
  → 0x3f3f3f3f3f3f3f3fLL
                                                           const int N = 1 \ll 20;
                                                                                                                              n + 1), mul(2 * n + 1, inv6))) + <math>mul(mul(b / a))
  const int MAX = ?; // maximum number of nodes
                                                           // apply modulo if necessary
                                                                                                                             c, b / c), n + 1) + mul(mul(a / c, b / c),
                                                           void fwht xor(int *a, int n, int dir = 0) {
                                                                                                                              mul(n, n + 1)) % MOD;
  int n, src, snk;
                                                               for (int h = 1; h < n; h <<= 1) {
                                                                                                                          return {f, g, h};
  T dis[MAX], mCap[MAX];
                                                                   for (int i = 0; i < n; i += h << 1)
  int par[MAX], pos[MAX];
                                                                       for (int j = i; j < i + h; ++j) {
                                                                                                                        long long m = (a * n + b) / c;
  bool vis[MAX];
                                                                            int x = a[j], y = a[j + h];
                                                                                                                        nxt = query(c, c - b - 1, a, m - 1);
  struct Edge{
                                                                           a[j] = x + y, a[j + h] = x - y;
                                                                                                                        f = (mul(m, n) - nxt.f) % MOD;
    int to, rev pos;
                                                                           if (dir) a[j] >>= 1, a[j + h] >>= 1;
                                                                                                                        g = mul(mul(m, mul(n, n + 1)) - nxt.h - nxt.f,
    T cap, cost, flow;
                                                                                                                        h = (mul(n, mul(m, m + 1)) - 2 * nxt.q - 2 * nxt.f -
  vector <Edge> ed[MAX];

→ f) % MOD:

  void init(int n, int src, int snk) {
                                                                                                                        return {f, g, h};
                                                           void fwht or(int *a, int n, int dir = 0) {
    n = n, src = src, snk = snk;
                                                               for (int h = 1; h < n; h <<= 1) {
    for(int i=1;i<=n;i++) ed[i].clear();</pre>
                                                                   for (int i = 0; i < n; i += h << 1)
                                                                       for (int j = i; j < i + h; ++j) {
  void addEdge(int u, int v, T cap, T cost) {
                                                                                                                      5.3 NOD
                                                                           int x = a[j], y = a[j + h];
a[j] = x, a[j + h] = dir ? y - x : x +
    Edge a = \{v, (int)ed[v].size(), cap, cost, \theta\};
                                                                                                                     N = input()
    Edge b = \{u, (int)ed[u].size(), 0, -cost, 0\};
                                                                             پ َ y;
                                                                                                                     primes = array containing primes till 10^6
    ed[u].pb(a);
    ed[v].pb(b);
                                                                                                                      for all p in primes :
                                                                                                                          if p*p*p > N:
  inline bool SPFA(){
                                                                                                                              break
    CLR(vis);
                                                           void fwht and(int *a, int n, int dir = 0) {
    for(int i=1; i<=n; i++) mCap[i] = dis[i] = INF;</pre>
                                                               for (int h = 1; h < n; h <<= 1) {
                                                                                                                          while N divisible by p:
    queue <int> q;
                                                                   for (int i = 0; i < n; i += h << 1)
                                                                                                                              N = N/p
    dis[src] = 0;
vis[src] = true;
                                                                                                                          count = count + 1
                                                                       for (int j = i; j < i + h; ++j) {
                                                                            int x = a[j], y = a[j + h];
                                                                                                                      if N is prime:
    q.push(src);
                                                                           a[j] = dir^{?} x - y : x + y, a[j + h] =
                                                                                                                          ans = ans * 2
    while(!q.empty()){

    y;

                                                                                                                      else if N is square of a prime:
      int u = q.front();
                                                                                                                          ans = ans * 3
                                                                   }
      q.pop();
                                                                                                                      else if N != 1:
                                                               }
      vis[u] = false;
                                                                                                                          ans = ans *
      for(int i=0; i<(int)ed[u].size(); i++) {</pre>
        Edge \&e = ed[u][i];
                                                           5.2 FloorSum
        int v = e.to;
                                                                                                                      5.4 Pollard Rho
        if(e.cap > e.flow && dis[v] > dis[u] + e.cost){\frac{1}{0}0(log a) sum^n floor(ax + b / c) = f
          dis[v] = dis[u] + e.cost;
                                                           long long FloorSumAP(long long a, long long b, long
                                                                                                                      we#include <bits/stdc++.h>
          par[v] = u;
                                                              long c, long long n){
                                                                                                                      using namespace std;
          pos[v] = i;
```

if(!a) return (b / c) * (n + 1); if(a >= c or b >= c) return ((n * (n + 1)) / 2) *

(a / c) + (n + 1) * (b / c) + FloorSumAP(a % c,

mCap[v] = min(mCap[u],e.cap - e.flow);

if(!vis[v]) {

vis[v] = true;

typedef long long ll;

namespace Rho {

typedef unsigned long long ull;

13

```
ull mul (ull a, ull b, ull mod)
    ll ret = a * b - mod * (ull) (1.L / mod * a * b);
return ret + mod * (ret < 0) - mod * (ret >= (ll)
  ull bigMod (ull a, ull e, ull mod) {
    ull ret = 1:
       if (e \& 1) ret = mul(ret, a, mod);
       a = mul(a, a, mod), e >>= 1;
    return ret;
  bool isPrime (ull n) {
   if (n < 2 or n % 6 % 4 != 1) return (n | 1) == 3;</pre>
    ull a[] = \{2, 325, 9375, 28178, 450775, 9780504,
         1795265022
    ull s = builtin ctzll(n - 1), d = n >> s;
    for (ull x : a) {
   ull p = bigMod(x % n, d, n), i = s;
       while (p != 1 \text{ and } p != n - 1 \text{ and } x \% \text{ n and } i--)
        \rightarrow p = mul(p, p, n);
       if (p != n - 1 \text{ and } i != s) return 0;
    return 1;
  ull pollard (ull n) {
    auto f = [\&] (ull x) {return mul(x, x, n) + 1;};
    ull x = 0, y = 0, t = 0, prod = 2, i = 1, q;
    while (t++ \% 40 \text{ or } \gcd(\text{prod}, n) == 1) {
       if (x == y) x = ++i, y = f(x);
       if ((q = mul(prod, max(x, y) - min(x, y), n)))
           prod = q;
      x = f(x), y = f(f(y));
    return gcd(prod, n);
  vector <ull> factor (ull n) {
    if (n == 1) return {};
    if (isPrime(n)) return {n};
    ull x = pollard(n);
    auto l = factor(x), r = factor(n / x);
    l.insert(l.end(), r.begin(), r.end());
    return l;
int t; ll n;
int main() {
  cin >> t;
  while (t--)
    nile (t--) {
    scanf("%lld", &n);
    vector <ull> facs = Rho::factor(n);
    sort(facs.begin(), facs.end());
    printf("%d", (int) facs.size());
    for (auto it : facs) printf(" %llu", it);
    puts("");
  return 0;
```

5.5 catalan

```
//Recursive
const int MOD = ....
const int MAX = ....
int catalan[MAX];
void init()
     catalan[0] = catalan[1] = 1;
    for (int i=2; i<=n; i++) {
    catalan[i] = 0;
          for (int j=0; j < i; j++) {
    catalan[i] += (catalan[j] *</pre>

    catalan[i-j-1]) % MOD;
```

```
if (catalan[i] >= MOD) {
   catalan[i] -= MOD;
               }
  //Analytical formula:
|ans = ncr(2*n,n)-ncr(2*n,n-1) = ncr(2*n,n)/(n+1)
  //r[i][j]= inverse of p[i] modulo p[j]
 1/ans = x[0] + x[1] * p[0] + x[2] * (p[0] * p[1]) + ... + x[k-1] * (p[0] * p[1]) + ... + x[
 - ]*p[1]*p[2]*...*p[k-2])
//ans %= ((p[0]*p[1]*p[2]*...*p[k-1])
|for (int i = 0; i < k; ++i) {
                \hat{x}[i] = a[i];
                 for (int^{-}j = 0; j < i; ++j) {
                               \dot{x}[i] = r[j][i] * (x[i] - x[j]);
                              x[i] = x[i] % p[i];
                              if (x[i] < 0)
x[i] += p[i];
 ll mul= p[0],res=x[0],tot=1;
F(i,0,k) tot *= p[i];
F(i,1,k)
   res+= x[i]*mul;
res %= tot;
    mul *= p[i];
  res %= mul;
 return res:
 5.7 derangement
  int derangement(int n)
|if(!n) return n;
lif(n <= 2) return n-1;
 return (n-1)*(derangement(n-1) + derangement(n-2));
 5.8 diophantine
  void print solution(int a, int b, int c)
                int x, y;
               if (a == 0 \&\& b == 0) {
                               if (c == 0) {
                                              cout << "Infinite Solutions Exist" << endl;
```

```
else {
         cout << "No Solution exists" << endl;
int gcd = gcd extend(a, b, x, y);
if (c % acd != 0) {
    cout<< "No Solution exists"<< endl;</pre>
    \dot{\text{cout}} << "x = " << x * (c / qcd) << ", y = " <<
     \rightarrow v * (c / qcd) << endl;
```

// Returns minimum x for which a ^ x % m = b % m, a

5.9 discrete_log

→ and m are coprime.

a %= m, b %= m; int n = sqrt(m) + 1;

|int solve(int a, int b, int m) {

```
5.11 fft
typedef complex<double> base;
#define PI acos(-1)
void fft(vector<base> &a, bool invert){
    int n = (int)a.size()
    for (int i = 1, j = 0; i < n; ++i){
         int bit = n >> 1:
         for (; j >= bit; bit >>= 1) j -= bit;
         j += bit;
         if (i < j)swap(a[i], a[j]);</pre>
    for (int len = 2; len <= n; len <<= 1){
    double ang = 2 * PI / len * (invert ? -1 : 1);</pre>
         base wlen(cos(ang), sin(ang));
         for (int i = 0; i < n; i + = len){
             base w(1);
             for (int j = 0; j < len / 2; ++j){
  base u = a[i + j], v = a[i + j + len /</pre>
                  a[i + j] = u + v;
                  a[i + j + len / 2] = u - v;
                  w^* = wlen;
    if (invert) for (int i = 0; i < n; ++i) a[i] /= n;</pre>
void multiply(const vector<int> &a, const vector<int>
   &b, vector<int> &res){
    vector<base> fa(a.begin(), a.end()), fb(b.begin(),

    b.end());
    size t n = 1:
    while (n < max(a.size(), b.size())) n <<= 1;</pre>
    fa.resize(n), fb.resize(n);
```

int an = 1;

return -1;

5.10 factorial_mod_p

int res = 1:

while (n > 1) {

n /= p;

return res;

int factmod(int n, int p) {

if ((n/p) % 2)

vector<int> f(p); f[0] = 1;

for (int i = 0; i < n; ++i) an = (an * 1ll * a) % m;

unordered map<int. int> vals:

return ans;

for (int i = 1; i < p; i++) f[i] = f[i-1] * i % p;

res'é p - res:

res = res * f[n%p] % p;

vals[cur] = q;

for (int q = 0, cur = b; $q \le n$; ++q) {

for (int p = 1, cur = 1; $p \le n$; ++p) {

// O(log p(n)) gives me n! % p for large n, p

int ans = n * p - vals[cur];

cur = (cur * 111 * a) % m;

cur = (cur * 111 * an) % m; if (vals.count(cur)) {

14

```
15
```

```
fft(fa, false), fft(fb, false);
for (size_t i = 0; i < n; ++i) fa[i] *= fb[i];</pre>
    fft(fa, true); res.resize(n);
    for (size t i = 0; i < n; ++i) res[i] =

    int(fa[i].real() + 0.5);

5.12 gauss_eliminition
const double EPS = 1e-9:
const int INF = 2; // it doesn't actually have to be

→ infinity or a big number

int gauss (vector < vector < double> > a, vector < double>

→ & ans) {
    int n = (int) a.size();
    int m = (int) a[0].size() - 1;
    vector<int> where (m, -1);
    for (int col=0, row=0; col<m && row<n; ++col) {
         int sel = row;
         for (int i=row; i<n; ++i)
             if (abs (a[i][col]) > abs (a[sel][col]))
                  sel = i
        if (abs (a[sel][col]) < EPS)</pre>
             continue;
         for (int i=col; i<=m; ++i)</pre>
             swap (a[sel][i], a[row][i]);
        where[col] = row;
        for (int i=0; i<n; ++i)</pre>
             for (int j=col; j<=m; ++j)
a[i][j] -= a[row][j] * c;
        ++row;
    ans.assign (m, \theta);
    for (int i=0; i<m; ++i)</pre>
         if (where[i] != -1)
             `ans[i] = a[where[i]][m] / a[where[i]][i];
    for (int i=0; i<n; ++i) {
        double sum = 0;

for (int j=0; j<m; ++j)

sum += ans[j] * a[i][j];
        if (abs (sum - a[i][m]) > EPS)
             return 0;
    for (int i=0; i<m; ++i)
         if (where[i] \stackrel{.}{=} -1)
             return İNF;
    return 1:
int gauss (vector < bitset<N> > a, int n, int m,

    bitset<N> & ans) {

    vector<int> where (m. -1):
    for (int col=0, row=0; col<m && row<n; ++col) {</pre>
         for (int i=row; i<n; ++i)</pre>
             if (a[i][col]) {
    swap (a[i], a[row]);
                  break:
        if (! a[row][col])
             continue:
        where[col] = row;
         for (int i=0; i<n; ++i)
             if (i != row && a[i][col])
a[i] ^= a[row];
    }
         // The rest of implementation is the same as
         → above
```

```
const double EPS = 1E-9;
int compute rank(vector<vector<double>> A) {
     int n = A.size();
     int m = A[0].size();
     int rank = 0;
     vector<bool> row selected(n, false);
     for (int i = 0: i < m: ++i) {
         if (!row selected[j] && abs(A[j][i]) > EPS)
                   break:
          if (j != n) {
               ++rank:
               row selected[j] = true;
               for (int p = i + 1; p < m; ++p)
                   A[j][p] /= A[j][i];
              for (int k = 0; k < n; ++k) {
    if (k != j \&\& abs(A[k][i]) > EPS) {
                        for (int p = i + 1; p < m; ++p)
                             A[k][p] -= A[j][p] * A[k][i];
         }
     return rank;
5.13 gen_all_k_combs
 vector<int> ans:
void gen(int n, int k, int idx, bool rev) {
   if (k > n || k < 0)</pre>
          return;
     if (!n)
          for (int i = 0; i < idx; ++i) {
               if (ans[i])
                   cout \ll i + 1;
          cout << "\n";
          return;
     ans[idx] = rev;
     gen(n-1, k-rev, idx + 1, false);
     ans[idx] = !rev;
     gen(n - 1, k - !rev, idx + 1, true);
void all combinations(int n, int k) {
     ans.resize(n);
     gen(n, k, 0, false);
 5.14 integrate_adaptive
 /*Description: Fast integration using an adaptive

→ Simpsons rule.

Usage: double sphereVolume = quad(-1, 1, [](double x) {
return quad(-1, 1, [\&](double y) return quad(-1, 1, [\&](double z)
 return \dot{x}^*x + \dot{v}^*y + z^*z < 1; \}); \}); \}); */
typedef double d:
#define S(a, b) (f(a) + 4 * f((a + b) / 2) + f(b)) *
 \rightarrow (b - a) / 6
template <class F>
|d rec(F& f, d a, d b, d eps, d S) {
     \begin{array}{l} d \ c = (a + b) \ / \ 2; \\ d \ S1 = S(a, c), \ S2 = S(c, b), \ T = S1 + S2; \\ if \ (abs(T - S) <= 15 * eps | | b - a < 1e-10) \end{array}
          return T + (T - S) / 15;
     return rec(f, a, c, eps \sqrt{2}, S1) + rec(f, c, b,
      \rightarrow eps / 2, S2);
|}
```

```
template <class F>
d quad(d a, d b, F f, d eps = 1e-8) {
    return rec(f, a, b, eps, S(a, b));
5.15 matrix expo
11 \mod = (1e9) + 7:
struct Matrix{
    int row, col;
    vector<vector<ll>> mat;
    Matrix(int x, int y){
        row=x:
        col=y;
        mat.assign(row, vector<ll>(col,0));
    Matrix operator *(Matrix &other){
        assert(col==other.row);
        Matrix product(row,other.col);
        for(int i=0;i<row;i++){</pre>
             for(int j=0; j < col; j++) {</pre>
                 for(int k=0;k<other.col;k++){</pre>
                     product.mat[i][k]=(product.mat[i][<sub>]</sub>
                          k]+(mat[i][i]*other.mat[i][k])
                      ⇒ %mod)%mod;
        return product;
Matrix expo(Matrix &m, ll n){
    assert(m.row==m.col);
    Matrix ret(m.row,m.col);
    for(int i=0;i<m.row;i++) ret.mat[i][i]=1;</pre>
    while(n){
        if(n\&1) ret=ret*m;
        n/=2;
m=m*m;
    return ret:
5.16 next_lexicographical_k_comb
bool next combination(vector<int>& a, int n) {
    int k = (int)a.size();
    for (int i = k - 1; i >= 0; i - -) {
        if (a[i] < n - k + i + 1)  {
             a[i]++;
             for (int j = i + 1; j < k; j++)
                 a[j] = a[\bar{j} - \bar{1}] + 1;
             return true:
    return false:
5.17 ntt
const int mod = 998244353;
const int root = 15311432;
const int k = 1 << 23;
int root 1:
vector<int> rev:
void pre(int sz){
    root 1 = bigmod(root, mod - 2, mod);
    if (rev.size() == sz) return;
```

rev.resize(sz);

int lq n = builtin ctz(sz);

rev[0] = 0;

```
16
```

```
for (int i = 1; i < sz; ++i)
     \rightarrow rev[i]=(rev[i>>1]>>1)|((i&1)<<(lq n-1));
void fft(vector<int> &a, bool inv){
    int n = a.size();
    for (int i = 1; i < n - 1; ++i) if (i < rev[i])

    swap(a[i], a[rev[i]]);

    for (int len = 2; len <= n; len <<= 1) {
         int wlen = inv ? root 1 : root;
         for (int i = len; i < k; i <<= 1) wlen = 1ll *

→ wlen * wlen % mod;

         for (int st = 0; st < n; st += len){</pre>
             int w = 1;
             for (int j = 0; j < len / 2; j++){
                  int ev = a[st + j];
                  int od = 1ll * a[st + j + len / 2] * w
                  a[st + j]' = ev + od < mod ? ev + od :
                   \rightarrow ev + od - mod:
                  a[st + j + len / 2] = ev - od >= 0?
                  \rightarrow ev - od : ev - od + mod:
                  w = 111 * w * wlen % mod;
        }
    if (inv){
         int n = bigmod(n, mod - 2, mod);
         for (\overline{int} \&x : a) x = 111 * x * n 1 % mod;
vector<int> mul(vector<int> &a, vector<int> &b){
    int n = a.size(), m = b.size(), sz = 1;
    while (sz < n + m - 1) sz <<= 1;
vector<int> x(sz), y(sz), z(sz);
    for (int i = 0; i < sz; ++i){
    x[i] = i < n ? a[i] : 0;</pre>
         v[\bar{i}] = \bar{i} < m ? b[\bar{i}] : 0;
    pre(sz);fft(x, 0);fft(y, 0);
    for (int i = 0; i < sz; ++i) z[i] = 111 * x[i] *

→ y[i] % mod;

    fft(z, 1); z.resize(n + m - 1);
    return z;
```

5.18 seg_sieve

```
Seamented Sieve
This code was for 1 <= a <= b <= 2^31-1
Change variable types appropriately.
bool notPrime[ ? ];
void segmented_sieve(int a, int b)
    int p, f;
    mem(notPrime, 0);
    for (int i = 0; i < tot prime; i++)
         p = prime[i];
        if (a \% p == 0)
             f = a;
             f = (a - (a \% p) + p);
         f = max(p * p, f);
        for (unsigned j = f; j \le b; j += p)
             notPrime[i - a] = true;
    if (a == 1)
        notPrime[0] = 1;
```

```
5.19 stirling
that no box is empty.
int stirling2(int n, int k)
 if(n < k)
 return 0:
 if(k == 1)
 return 1
 if(dp[n][k] == dp[n][k])
 return dp[n][k];
 return dp[n][k] = stirling2(n-1,k-1) +

    stirlina2(n-1.k)*k:
is empty
int stirling1(int n, int k)
 dp[n][k] = stirling1(n-1,k-1) + stirling(n-1,k)*n-1;
5.20 stirling_number_of_the_second_kind
   1 / k! * sum (-1)^i nCr(k, i) * (k - i) ^ n
ll f(int n, int k) {
    ll res = 0:
    for (int i = 0; i < k; ++i) {
        if (i \& 1) res = (res - nCr(k, i) * bp(k - i, i)

_ n, mod) % mod + mod) % mod;

        else res = (res + nCr(k, i) * bp(k - i, n,

→ mod) % mod) % mod;

    if (res < 0) res += mod;
    return res * ifac[k] % mod;
5.21 sum_of_totient
using namespace __gnu_pbds;
const int N = 3e5 + 9, mod = 998244353;
template <const int32 t MOD>
struct modint
    int32 t value;
    modin\overline{t}() = default;
    modint(int32 t value) : value(value) {}
    inline modint<MOD> operator+(modint<MOD> other)
        int32 t c = this->value + other.value;
        return modint<MOD>(c >= MOD ? c - MOD : c);
    inline modint<MOD> operator-(modint<MOD> other)
     int32 t c = this->value - other.value;
        return modint<MOD>(c < \theta ? c + MOD : c);
    inline modint<MOD> operator*(modint<MOD> other)
        int32 t c = (int64_t)this->value * other.value
        return modint<MOD>(c < 0 ? c + MOD : c);
    inline modint<MOD> &operator+=(modint<MOD> other) {
        this->value += other.value;
        if (this->value >= MOD)
            this->value -= MOD;
        return *this:
    inline modint<MOD> &operator-=(modint<MOD> other) {
        this->value -= other.value:
        if (this->value < 0)</pre>
            this->value += MOD;
        return *this;
```

```
inline modint<MOD> &operator*=(modint<MOD> other) {
        this->value = (int64 t)this->value

    other.value %
    MOD;

        if (this->value < 0) this->value += MOD;
        return *this;
    inline modint<MOD> operator-() const { return
        modint<MOD>(this->value ? MOD - this->value :
    modint<MOD> pow(uint64_t k) const {
        modint < MOD > x =
                    y = 1;
        for (; k; k >>= 1) {
   if (k & 1) y *= x;
            x *= x;
        return y;
    modint<MOD> inv() const { return pow(MOD - 2); }
    → // MOD must be a prime
    inline modint<MOD> operator/(modint<MOD> other)

    const { return *this * other.inv(); }

    inline modint<MOD> operator/=(modint<MOD> other) {
        return *this *= other.inv(); }
    inline bool operator==(modint<MOD> other) const {
       return value == other.value; }
    inline bool operator!=(modint<MOD> other) const {

    return value != other.value; }

    inline bool operator<(modint<MOD> other) const {
    → return value < other.value; }</pre>
    inline bool operator>(modint<MOD> other) const {
    → return value > other.value; }
template <int32 t MOD>
modint<MOD> operator*(int64 t value, modint<MOD> n) {
→ return modint<MOD>(value) * n; }
template <int32 t MOD>
modint<MOD> operator*(int32 t value, modint<MOD> n) {
→ return modint<MOD>(value % MOD) * n; }
template <int32 t MOD>
istream & operator >> (istream & in, modint < MOD > & n) {
→ return in >> n.value; }
template <int32 t MOD>
ostream & operator << (ostream & out, modint < MOD> n) {

    return out << n.value; }
</pre>
using mint = modint<mod>;
namespace Dirichlet {
// solution for f(x) = phi(x)
const int T = 1e7 + 9;
long long phi[T];
gp hash table<long long, mint> mp;
mint dp[T], inv;
int sz, spf[T], prime[T];
void init() {
    memset(spf, 0, sizeof spf);
    phi[1] = 1;
    sz = 0;
    for (int i = 2; i < T; i++) {
        if (spf[i] == 0) phi[i] = i - 1, spf[i] = i,

    prime[sz++] = i:

        for (int j = 0; j < sz && i * prime[j] < T &&

    prime[j] <= spf[i]; j++)
</pre>
            spf[i * prime[j]] = prime[j];
            if (i % prime[i] == 0)
                 phi[i * prime[j]] = phi[i] * prime[j];
                phi[i * prime[j]] = phi[i] * (prime[j]

→ - 1);
```

```
dp[0] = 0;
    for (int i = 1; i < T; i++) dp[i] = dp[i - 1] +
     → phi[i] % mod;
    inv = 1; // g(1)
mint p c(long long n) {
    if^{-}(n \% 2 == 0) return n / 2 \% mod * ((n + 1) %)
    return (n + 1) / 2 % mod * (n % mod) % mod;
mint p q(long long n) {
    return n % mod:
mint solve(long long x)
    if (x < T) return dp[x]:
    if (mp.find(x) != mp.end()) return mp[x];
    mint ans = 0:
    for (long long i = 2, last; i <= x; i = last + 1) {
    last = x / (x / i);
    ans += solve(x / i) * (p_g(last) - p_g(i - 1));</pre>
    ans = p c(x) - ans;
    ans /= \overline{i}nv;
    return mp[x] = ans;
}; // namespace Dirichlet
5.22 totient
```

int phi(int n) { int result = n; for (int i = 2; i * i <= n; i++) { result -= result / i: **if** (n > 1)result -= result / n; return result; void phi 1 to n(int n) { vector < int > phi(n + 1);phi[0] = 0: phi[1] = 1;for (int i = 2; i <= n; i++) phi[i] = i;</pre> for (int i = 2; i <= n; i++) { if (phi[i] == i) {</pre> for (int j = i; j <= n; j += i) phi[j] -= phi[j] / i;

5.23 xor_basis

```
int basis[d]; // basis[i] keeps the mask of the vector

→ whose f value is i

int sz;
void insertVector(int mask) {
 for (int i = 0; i < d; i++) {
  if ((mask & 1 << i) == 0) continue;
if (!basis[i]) { // If there is no basis vector with</pre>
       the i'th bit set, then insert this vector into

    the basis
basis[i] = mask;

   return;
  mask ^= basis[i]; // Otherwise subtract the basis

→ vector from this vector
```

```
6 Misc
6.1 interval container
/* Description: Add and remove intervals from a set of
    disjoint intervals. Will merge the added interval
   with any overlapping intervals in
the set when adding. Intervals are [inclusive,

    exclusive).
Time: 0 (log N)
set<pii>::iterator addInterval(set<pii>& is, int L,
    int R) {
    if (L == R) return is.end();
auto it = is.lower_bound({L, R}), before = it;
    while (it = is.en\overline{d}() && it->first <= R) {
         R = max(R, it->second);
         before = it = is.erase(it);
    if (it != is.begin() \&\& (--it)->second >= L) {
        L = min(L, it->first);
R = max(R, it->second); is.erase(it);
    return is.insert(before, {L, R});
void removeInterval(set<pii>& is, int L, int R) {
    if (L == R) return;
    auto it = addInterval(is, L, R);
    auto r2 = it->second;
    if (it->first == L)
         is.erase(it);
         (int&)it->second = L:
    if (R != r2) is.emplace(R, r2);
```

```
6.2 interval cover
```

```
template <class T>
vi cover(pair<T, T> G, vector<pair<T, T>> I) {
    vi S(sz(I)), R;
iota(all(S), 0);
sort(all(S), [&](int a, int b) { return I[a] <</pre>
    → I[b]; });
T cur = G.first;
     int at = 0;
     while (cur < G.second) \{ // (A) \}
          pair<T, int> mx = make pair(cur, -1);
          while (at < sz(I) && I[S[at]].first <= cur) {
    mx = max(mx, make_pair(I[S[at]].second,</pre>
               if (mx.second == -1) return {};
          cur = mx.first;
          R.push back(mx.second);
     return R;
```

6.3 pragma

```
// Praamas
#pragma comment(linker, "/stack:200000000")
#pragma GCC optimize("03,unroll-loops")
#pragma GCC target("avx,avx2,fma")
```

6.4 random

```
// shuffle(v.begin(), v.end(),
→ default random engine(rnd(1, 1000)));
mt19937 rnq(chrono::steady clock::now().time since epo

    ch().count());
```

```
ll rnd(ll l, ll r) {
    return uniform_int_distribution<ll>(l, r) (rng);
```

6.5 vimre

```
imap jk <Esc>
set nu
set mouse=a
set autoindent
set tabstop=4
set shiftwidth=4
set smartindent
set relativenumber
set laststatus=2
set hlsearch
let mapleader = " "
nnoremap <leader>s :w<Enter>
nnoremap <leader>y ggVG"+y<CR>
syntax on
vnoremap <leader>/ :s!^!//!<CR> :noh <CR>
vnoremap <leader>u :s!^//!!<CR>
nnoremap <leader>/ :s!^!//!<CR> :noh <CR>
nnoremap <leader>u :s!^//!!<CR>
```

7 Notes

7.1 Counting

1. Lucas Theorem

For non-negative integers m and n and a prime p, the following congruence relation holds: :

$$\begin{pmatrix} m \\ n \end{pmatrix} \equiv \prod_{i=0}^{k} \begin{pmatrix} m_i \\ n_i \end{pmatrix} \pmod{p},$$
 (1)

where:

$$m = m_k p^k + m_{k-1} p^{k-1} + \dots + m_1 p + m_0$$

and:

$$n = n_k p^k + n_{k-1} p^{k-1} + \dots + n_1 p + n_0$$

are the base p expansions of m and n respectively. This uses the convention that $\binom{m}{n} = 0$ if $m \le n$.

2. Stirling Numbers of the first kind

S(n,k) counts the number of permutations of n elements with k disjoint cycles.

$$S(n,k) = (n-1) \cdot S(n-1,k) + S(n-1,k-1)$$
 (2)

where, S(0,0) = 1, S(n,0) = S(0,n) = 0

$$\sum_{k=0}^{n} S(n,k) = n!$$
 (3)

3. Stirling Numbers of the second kind

 $S(n,k) \cdot k! = \text{number of ways to color } n \text{ nodes using colors from } 1 \text{ to } k$ such that each color is used at least once.

An r-associated Stirling number of the second kind is the number of ways to partition a set of n objects into k subsets, with each subset containing at least r elements. It is denoted by $S_r(n,k)$ and obeys the recurrence relation.

$$S_r(n+1,k) = kS_r(n,k) + \binom{n}{r-1}S_r(n-r+1,k-1)$$
 (4)

4. Bell Numbers

Counts the number of partitions of a set

$$B_{n+1} = \sum_{k=0}^{n} \left(\frac{n}{k}\right) \cdot B_k \tag{5}$$

 $B_n = \sum_{k=0}^n S(n,k)$, where S(n,k) is stirling number of second kind.

5. Some identities

Vandermonde's Identify: $\sum_{k=0}^{r} {m \choose k} {n \choose r-k} = {m+n \choose r}$

Hockey-Stick Identify:
$$n, r \in N, n > r, \sum_{i=r}^{n} \begin{pmatrix} i \\ r \end{pmatrix} = \begin{pmatrix} n+1 \\ r+1 \end{pmatrix}$$

Involutions: permutations such that p^2 = identity permutation. $a_0 = a_1 = 1$ and $a_n = a_{n-1} + (n-1)a_{n-2}$ for n > 1.

7.2 Fibonacci

Let A, B and n be integer numbers

$$k = A - B$$

$$F_A F_B = F_{k+1} F_A^2 + F_k F_A F_{A-1} \tag{6}$$

$$\sum_{i=0}^{n} F_i^2 = F_{n+1} F_n \tag{7}$$

$$\sum_{i=0}^{n} F_i F_{i+1} = F_{n+1}^2 - (-1)^n \tag{8}$$

$$\sum_{i=0}^{n} F_i F_{i-1} = \sum_{i=0}^{n-1} F_i F_{i+1}$$
 (9)

$$GCD(F_m, F_n) = F_{GCD(m,n)}$$
(10)

$$\sum_{0 \le k \le n} \binom{n-k}{k} = Fib_{n+1} \tag{11}$$

$$\gcd(F_n, F_{n+1}) = \gcd(F_n, F_{n+2}) = \gcd(F_{n+1}, F_{n+2}) = 1$$
 (12)

7.3 Notes 7.4 Geometry

7.4.1 Triangles

Circumradius:
$$R=\frac{abc}{4A}$$
, Inradius: $r=\frac{A}{s}$
Length of median (divides triangle into two equal-area triangles)

 $m_a = \frac{1}{2}\sqrt{2b^2 + 2c^2 - a^2}$

Length of bisector (divides angles in two): $s_a = \sqrt{bc \left[1 - \left(\frac{a}{b+c}\right)^2\right]}$

Law of tangents: $\frac{a+b}{a-b} = \frac{\tan \frac{\alpha+\beta}{2}}{\tan \frac{\alpha-\beta}{2}}$

7.4.2 Quadrilaterals

With side lengths a,b,c,d, diagonals e,f, diagonals angle θ , area and magic flux $F = b^2 + d^2 - a^2 - c^2$:

$$4A = 2ef \cdot \sin \theta = F \tan \theta = \sqrt{4e^2f^2 - F^2}$$

For cyclic quadrilaterals the sum of opposite angles is 180° , ef =ac+bd, and $A=\sqrt{(p-a)(p-b)(p-c)(p-d)}$.

7.4.3 Spherical coordinates

$$x = r \sin \theta \cos \phi \qquad r = \sqrt{x^2 + y^2 + z^2}$$

$$y = r \sin \theta \sin \phi \qquad \theta = \frac{1}{2} \cos(z) \sqrt{x^2 + y^2 + z^2}$$

$$z = r \cos \theta \qquad \phi = \frac{1}{2} \sin(z) \sqrt{x^2 + y^2 + z^2}$$

1.3 Sums
$$1^{2} + 2^{2} + 3^{2} + \dots + n^{2} = \frac{n(2n+1)(n+1)}{6}$$
(5)
$$1^{3} + 2^{3} + 3^{3} + \dots + n^{3} = \frac{n^{2}(n+1)^{2}}{4}$$

$$1^{4} + 2^{4} + 3^{4} + \dots + n^{4} = \frac{n(n+1)(2n+1)(3n^{2} + 3n - 1)}{30}$$

$$\sum_{i=1}^{n} i^{m} = \frac{1}{m+1} \left[(n+1)^{m+1} - 1 - \sum_{i=1}^{n} ((i+1)^{m+1} - i^{m+1} - (m+1)i^{m}) \right]$$

$$\sum_{i=1}^{n-1} i^{m} = \frac{1}{m+1} \sum_{k=0}^{m} {m+1 \choose k} B_{k} n^{m+1-k}$$
ion.
$$\sum_{k=0}^{n} kx^{k} = (x - (n+1)x^{n+1} + nx^{n+2})/(x-1)^{2}$$
7.6 Series

$e^x = 1 + x + \frac{x^2}{2!} + \frac{x^3}{3!} + \dots, (-\infty < x < \infty)$ $\ln(1+x) = x - \frac{x^2}{2} + \frac{x^3}{3} - \frac{x^4}{4} + \dots, (-1 < x \le 1)$ $\sqrt{1+x} = 1 + \frac{x}{2} - \frac{x^2}{8} + \frac{2x^3}{32} - \frac{5x^4}{128} + \dots, (-1 \le x \le 1)$ $\sin x = x - \frac{x^3}{3!} + \frac{x^5}{5!} - \frac{x^7}{7!} + \dots, (-\infty < x < \infty)$ $\cos x = 1 - \frac{x^2}{2!} + \frac{x^4}{4!} - \frac{x^6}{6!} + \dots, (-\infty < x < \infty)$ $(x+a)^{-n} = \sum_{k=0}^{\infty} (-1)^k \binom{n+k-1}{k} x^k a^{-n-k}$

(10) 7.7 Pythagorean Triples

The Pythagorean triples are uniquely generated by

$$a = k \cdot (m^2 - n^2), b = k \cdot (2mn), c = k \cdot (m^2 + n^2),$$

(12) with m > n > 0, k > 0, $m \perp n$, and either m or n even.

7.8 Number Theory **7.8.1** Primes

p = 962592769 is such that $2^{21} | p - 1$, which may be useful. For hashing use 970592641 (31-bit number), 31443539979727 (45-bit), 3006703054056749 (52-bit). There are 78498 primes less than

Primitive roots exist modulo any prime power p^a , except for p = 2, a > 1 = ab - a - b. 2, and there are $\phi(\phi(p^a))$ many. For p=2, a>2, the group \mathbb{Z}_{2a}^{\times} is 7.8.12 Fermat's two-squares theorem instead isomorphic to $\mathbb{Z}_2 \times \mathbb{Z}_{2a-2}$.

7.8.2 Estimates

 $\sum_{d|n} d = O(n \log \log n).$

The number of divisors of n is at most around 100 for n < 5e4, 500 for p = 4k + 3 occurs an even number of times in n's factorization. n < 1e7, 2000 for n < 1e10, 200 000 for n < 1e19.

7.8.3 Perfect numbers

n > 1 is called perfect if it equals sum of its proper divisors and 1. Even n is perfect iff $n = 2^{p-1}(2^p - 1)$ and $2^p - 1$ is prime (Mersenne's). No odd perfect numbers are vet found.

7.8.4 Carmichael numbers

A positive composite *n* is a Carmichael number $(a^{n-1} \equiv 1 \pmod{n})$ for all $a: \gcd(a, n) = 1$, iff n is square-free, and for all prime divisors p of n, p-1 divides n-1.

7.8.5 Mobius function

 $\mu(1) = 1$. $\mu(n) = 0$, if n is not squarefree. $\mu(n) = (-1)^s$, if n is the product of s distinct primes. Let f, \overline{F} be functions on positive integers. If for all $n \in N$, $F(n) = \sum_{d|n} f(d)$, then $f(n) = \sum_{d|n} \mu(d) F(\frac{n}{d})$, and vice versa. $\phi(n) = \sum_{d|n} \mu(d) \frac{n}{d}$. $\sum_{d|n} \mu(d) = 1$.

If f is multiplicative, then $\sum_{d|n} \mu(d) f(d) = \prod_{p|n} (1 - f(p)),$ $\sum_{d|n} \mu(d)^2 f(d) = \prod_{p|n} (1 + f(p)).$

7.8.6 Legendre symbol

If p is an odd prime, $a \in \mathbb{Z}$, then $\left(\frac{a}{p}\right)$ equals 0, if p|a; 1 if a is a quadratic residue modulo p; and -1 otherwise. Euler's criterion: $\left(\frac{a}{p}\right) = a^{\left(\frac{p-1}{2}\right)} \pmod{p}.$

7.8.7 Jacobi symbol

If
$$n = p_1^{a_1} \cdots p_k^{a_k}$$
 is odd, then $\left(\frac{a}{n}\right) = \prod_{i=1}^k \left(\frac{a}{p_i}\right)^{k_i}$.

7.8.8 Primitive roots

If the order of g modulo m (min n > 0: $g^n \equiv 1 \pmod{m}$) is $\phi(m)$, then g is called a primitive root. If Z_m has a primitive root, then it has $\phi(\phi(m))$ distinct primitive roots. Z_m has a primitive root iff m is one of 2, 4, p^k , $2p^k$, where p is an odd prime. If Z_m has a primitive root g, then for all a coprime to m, there exists unique integer $i = \operatorname{ind}_g(a)$ modulo $\phi(m)$, such that $g^i \equiv a \pmod{m}$. ind_{σ}(a) has logarithm-like properties: ind(1) = 0, ind(ab) = ind(a) + ind(b).

If p is prime and a is not divisible by p, then congruence $x^n \equiv a$ \pmod{p} has $\gcd(n, p-1)$ solutions if $a^{(p-1)/\gcd(n, p-1)} \equiv 1 \pmod{p}$. and no solutions otherwise. (Proof sketch: let g be a primitive root, and $g^i \equiv a \pmod{p}, g^u \equiv x \pmod{p}. x^n \equiv a \pmod{p} \text{ iff } g^{nu} \equiv g^i \pmod{p}$ iff $nu \equiv i \pmod{p}$.)

7.8.9 Discrete logarithm problem

Find x from $a^x \equiv b \pmod{m}$. Can be solved in $O(\sqrt{m})$ time and space with a meet-in-the-middle trick. Let $n = \lceil \sqrt{m} \rceil$, and x = ny - z. Equation becomes $a^{ny} \equiv ba^z \pmod{m}$. Precompute all values that the RHS can take for z = 0, 1, ..., n - 1, and brute force y on the LHS, each time checking whether there's a corresponding value for RHS.

7.8.10 Pythagorean triples

Integer solutions of $x^2 + y^2 = z^2$ All relatively prime triples are given by: x = 2mn, $y = m^2 - n^2$, $z = m^2 + n^2$ where m > n, gcd(m, n) = 1 and $m \ne n \pmod{2}$. All other triples are multiples of these. Equation $x^2 + y^2 = 2z^2$ is equivalent to $(\frac{x+y}{2})^2 + (\frac{x-y}{2})^2 = z^2$.

7.8.11 Postage stamps/McNuggets problem

Let a, b be relatively-prime integers. There are exactly $\frac{1}{2}(a-1)(b-1)$ numbers not of form ax + by $(x, y \ge 0)$, and the largest is (a-1)(b-1)

Odd prime p can be represented as a sum of two squares iff $p \equiv$ 1 (mod 4). A product of two sums of two squares is a sum of two squares. Thus, n is a sum of two squares iff every prime of form

7.9 Permutations

7.9.1 Factorial

7.9.2 Cycles

Let $g_S(n)$ be the number of *n*-permutations whose cycle lengths all belong to the set S. Then

$$\sum_{n=0}^{\infty} g_S(n) \frac{x^n}{n!} = \exp\left(\sum_{n \in S} \frac{x^n}{n}\right)$$

7.9.3 Derangements

Permutations of a set such that none of the elements appear in their Total number of partitions of n distinct elements. original position.

$$D(n) = (n-1)(D(n-1) + D(n-2)) = nD(n-1) + (-1)^n = \left\lfloor \frac{n!}{e} \right\rfloor$$

7.9.4 Burnside's lemma

Given a group G of symmetries and a set X, the number of elements of X up to symmetry equals

$$\frac{1}{|G|} \sum_{g \in G} |X^g|,$$

where X^g are the elements fixed by g(g.x = x).

If f(n) counts "configurations" (of some sort) of length n, we can ignore rotational symmetry using $G = \mathbb{Z}_n$ to get

$$g(n) = \frac{1}{n} \sum_{k=0}^{n-1} f(\gcd(n,k)) = \frac{1}{n} \sum_{k|n} f(k)\phi(n/k)$$

7.10 Partitions and subsets

7.10.1 Partition function

Number of ways of writing n as a sum of positive integers, disregarding the order of the summands.

$$p(0) = 1, \ p(n) = \sum_{k \in \mathbb{Z} \setminus \{0\}} (-1)^{k+1} p(n - k(3k - 1)/2)$$

$$p(n) \sim 0.145/n \cdot \exp(2.56\sqrt{n})$$

7.11 General purpose numbers

7.11.1 Stirling numbers of the first kind

Number of permutations on n items with k cycles.

$$c(n,k) = c(n-1,k-1) + (n-1)c(n-1,k), \ c(0,0) = 1$$

$$\sum_{k=0}^{n} c(n,k)x^{k} = x(x+1)...(x+n-1)$$

c(8,k) = 8,0,5040,13068,13132,6769,1960,322,28,1c(n,2) = 0,0,1,3,11,50,274,1764,13068,109584,...

7.11.2 Eulerian numbers

Number of permutations $\pi \in S_n$ in which exactly k elements are greater than the previous element. k j:s s.t. $\pi(j) > \pi(j+1)$, k+1 j:s s.t. $\pi(i) \ge i$, k i:s s.t. $\pi(i) > i$.

$$E(n,k) = (n-k)E(n-1,k-1) + (k+1)E(n-1,k)$$

$$E(n,0) = E(n,n-1) = 1$$

$$E(n,k) = \sum_{i=0}^{k} (-1)^{j} \binom{n+1}{j} (k+1-j)^{n}$$

7.11.3 Stirling numbers of the second kind

Partitions of *n* distinct elements into exactly *k* groups.

$$S(n,k) = S(n-1,k-1) + kS(n-1,k)$$

$$S(n,1) = S(n,n) = 1$$

$$S(n,k) = \frac{1}{k!} \sum_{j=0}^{k} (-1)^{k-j} \binom{k}{j} j^n$$

7.11.4 Bell numbers

B(n) = $1, 1, 2, 5, 15, 52, 203, 877, 4140, 21147, \dots$ For p prime,

$$B(p^m + n) \equiv mB(n) + B(n+1) \pmod{p}$$

7.11.5 Bernoulli numbers

$$\sum_{j=0}^{m} \binom{m+1}{j} B_j = 0. \quad B_0 = 1, B_1 = -\frac{1}{2}. \ B_n = 0, \text{ for all odd } n \neq 1.$$

7.11.6 Catalan numbers

$$C_n = \frac{1}{n+1} \binom{2n}{n} = \binom{2n}{n} - \binom{2n}{n+1} = \frac{(2n)!}{(n+1)!n!}$$

$$C_0 = 1, \ C_{n+1} = \frac{2(2n+1)}{n+2}C_n, \ C_{n+1} = \sum C_i C_{n-i}$$

 $C_n = 1, 1, 2, 5, 14, 42, 132, 429, 1430, 4862, 16796, 58786, \dots$

- sub-diagonal monotone paths in an $n \times n$ grid.
- strings with *n* pairs of parenthesis, correctly nested.
- binary trees with with n + 1 leaves (0 or 2 children).
- ordered trees with n+1 vertices.
- ways a convex polygon with n+2 sides can be cut into triangles by connecting vertices with straight lines.
- permutations of [n] with no 3-term increasing subseq.

7.12 Inequalities

7.12.1 Titu's Lemma

For positive reals a_1, a_2, \dots, a_n and b_1, b_2, \dots, b_n ,

$$\frac{{a_1}^2}{b_1} + \frac{{a_2}^2}{b_2} + \ldots + \frac{{a_n}^2}{b_n} \ge \frac{a_1 + a_2 + \ldots + {a_n}^2}{b_1 + b_2 + \ldots + b_n}$$

Equality holds if and only if $a_i = kb_i$ for a non-zero real constant k.

7.13 Games

7.13.1 Grundy numbers

For a two-player, normal-play (last to move wins) game on a graph $(V,E): G(x) = \max(\{G(y): (x,y) \in E\}), \text{ where } \max(S) = \min\{n \ge 0: n \not\in E\}$ S. x is losing iff G(x) = 0.

7.13.2 Sums of games

- Player chooses a game and makes a move in it Grundy number of a position is xor of grundy numbers of positions in summed gamês.
- Player chooses a non-empty subset of games (possibly, all) and makes moves in all of them A position is losing iff each game is in a losing position.
- Player chooses a proper subset of games (not empty and not all), and makes moves in all chosen ones. A position is losing iff grundy numbers of all games are equal.
- Player must move in all games, and loses if can't move in some game A position is losing if any of the games is in a losing position.

7.13.3 Misère Nim

A position with pile sizes $a_1, a_2, ..., a_n \ge 1$, not all equal to 1, is losing iff $a_1 \oplus a_2 \oplus \cdots \oplus a_n = 0$ (like in normal nim.) A position with n piles of size 1 is losing iff n is odd.

7.14 NumberTheory

$$\sum_{k=1}^{n} \frac{1}{\gcd(k,n)} = \sum_{d|n} \frac{1}{d} \cdot \phi\left(\frac{n}{d}\right) = \frac{1}{n} \sum_{d|n} d \cdot \phi(d)$$
 (13)

$$\sum_{k=1}^{n} \frac{k}{\gcd(k,n)} = \frac{n}{2} \cdot \frac{1}{n} \cdot \sum_{d|n} d \cdot \phi(d)$$
 (14)

$$\sum_{k=1}^{n} \frac{n}{\gcd(k,n)} = 2 * \sum_{k=1}^{n} \frac{k}{\gcd(k,n)} - 1, \text{ for, } n > 1$$
 (15)

$$\sum_{i=1}^{n} \sum_{j=1}^{n} [\gcd(i,j) = 1] = \sum_{d=1}^{n} \mu(d) \left\lfloor \frac{n}{d} \right\rfloor^{2}$$
 (16)

$$\sum_{i=1}^{n} \sum_{j=1}^{n} \gcd(i,j) = \sum_{d=1}^{n} \phi(d) \left\lfloor \frac{n}{d} \right\rfloor^{2}$$
(17)

$$\sum_{i=1}^{n} \sum_{j=1}^{n} i \cdot j[\gcd(i,j) = 1] = \sum_{i=1}^{n} \phi(i)i^{2}$$
(18)

$$\sum_{i=1}^{n} \sum_{j=1}^{n} \operatorname{lcm}(i,j) = \sum_{l=1}^{n} \left(\frac{\left(1 + \left\lfloor \frac{n}{l} \right\rfloor\right) \left(\left\lfloor \frac{n}{l} \right\rfloor\right)}{2} \right)^{2} = \sum_{d|l} \mu(d) l d$$
 (19)

8 String

8.1 aho_corasick

```
using namespace std;
const int N = ?;
const int A = ?;
struct AC {
  int nd, pt;
  int next[N][A], link[N], out link[N], cnt[N], ans[N];
  vector <int> ed[N], out[N];
  AC(): nd(0), pt(0) { node(); }
  int node() {
    memset(next[nd], 0, sizeof next[nd]);
link[nd] = out link[nd] = cnt[nd] = 0;
ed[nd].clear(), out[nd].clear();
     return nd++:
  void clear() {
    nd = pt = 0;
     node();
  inline int get(char c) { return c - 'a'; }
  void insert(const string &T) {
    int u = 0;
for (char c : T) {
   if (!next[u][get(c)]) next[u][get(c)] = node();
       u = next[u][qet(c)];
     ans[pt] = 0;
     out[u].push back(pt++);
  void build() {
     queue <int> q;
     for (q.push(0); !q.empty(); ) {
       int u = q.front();
       q.pop();
       q.pop(),
for (int c = 0; c < A; ++c) {
   int v = next[u][c];
   if (!v) next[u][c] = next[link[u]][c];</pre>
             link[v] = u ? next[link[u]][c] : 0;
```

```
out link[v] = out[link[v]].empty() ?
               out link[link[v]] : link[v];
          ed[link[v]].push back(v);
          q.push(v);
  void dfs(int s)
    for(int x : ed[s]) dfs(x), cnt[s] += cnt[x];
    for(int e : out[s]) ans[e] = cnt[s];
  void traverse(const string &S) {
    int u = 0;
    for (char'c
      u = next[u][gét(c)];
      cnt[u]++;
    dfs(0);
char str[1000010], pat[505];
int main() {
        freopen("in.txt","r",stdin);
 AC aho;
 int t,T;
scanf("%d",&T);
  for(int t=1;t<=T;t++) {</pre>
    int n
   scanf("%d",&n);
scanf("%s",str);
    for(int i=1;i<=n;i++) {</pre>
      scanf("%s",pat);
      aho.insert(pat);
    aho.build();
    aho.traverse(str);
    printf("Case `%d:\n",t);
    for(int i=0;i<n;i++) {
      printf("%d\n",aho.ans[i]);
    aho.clear();
  return 0;
```

8.2 hash

```
struct Hash {
  struct base
    string s; int b, mod;
    vector<int> hash, p;
    void init(string & s, int b, int mod) \{ // b >
     → 26, prime.
       s = s; b = b, mod = mod;
       hash_resize(s.size());
       p.resize(s.size());
       hash[0] = s[0] - 'A' + 1; p[0] = 1;
       for(int i = 1; i < s.size(); ++i) {
  hash[i] = (ll) hash[i - 1] * b % mod;
  hash[i] += s[i] - 'A' + 1;</pre>
         if(hash[i] >= mod) hash[i] -= mod;
p[i] = (ll) p[i - 1] * b % mod;
    int get(int l, int r) {
       int ret = hash[r]
       if(l) ret -= ([l] hash[l - 1] * p[r - l + 1] %
       if(ret < 0) ret += mod;
       return ret;
 } h[2];
```

```
void init(string &s)
    h[0].init(s, 29, 1e9+7);
h[1].init(s, 31, 1e9+9);
  pair<int, int> get(int l, int r) {
    return { h[0].get(l, r), h[1].get(l, r) };
|8.3 hash_segtree
                               - 1
#define INVALID CHAR
namespace strhash {
  int n:
  const int MAX = 100010;
  int ara[MAX]
  const int MOD[] = {1067737007, 1069815139};
  const int BASE[] = {982451653, 984516781};
  int BP[2][MAX], CUM[2][MAX];
  void init(char *str) {
    n = strlen(str);
    for(int i=0;i<n;i++) ara[i] = str[i]-'0'+1;
  void precal() {
    BP[0][0] = BP[1][0] = 1;
CUM[0][0] = CUM[1][0] = 1;
    for(int i=1; i<MAX; i++)
      BP[0][i] = (BP[0][i-1] * (long long) BASE[0])
      \rightarrow % MOD[0];
BP[1][i] = ( BP[1][i-1] * (long long) BASE[1] )

→ % MOD[1];

      CUM[0][i] = (CUM[0][i-1] + (long long) BP[0][i]
      CUM[1][i] = (CUM[1][i-1] + (long long) BP[1][i]
       → ) % MOD[1];
  struct node {
    int sz;
    int h[2];
node() {}
  } tree[4*MAX];
  int lazy[4*MAX];
  inline void lazyUpdate(int n,int st,int ed) {
    if(lazy[n]!=INVALID CHAR){
      tree[n].h[0] = (lazy[n] * (long long)

    □ CUM[0][ed-st]) % MOD[0];
tree[n].h[1] = (lazy[n] * (long long))

→ CUM[1][ed-st]) % MOD[1];
      if(st!=ed){
         lazy[2*n] = lazy[n];
         lazy[2*n+1] = lazy[n];
       lazy[n] = INVALID CHAR;
  inline node Merge(node a, node b) {
    node ret;
    ret.h[0] = ((a.h[0] * (long long) BP[0][b.sz])
     \rightarrow + b.h[0] ) % MOD[0];
    ret.h[1] = ((a.h[1])*(long long) BP[1][b.sz])
     \rightarrow + b.h[1] ) % MOD[1];
    ret.sz = a.sz + b.sz;
    return ret;
  inline void build(int n,int st,int ed) {
    lazy[n] = INVALÌD CHAŔ;
    if(st==ed) -
      tree[n].h[0] = tree[n].h[1] = ara[st];
```

```
tree[n].sz = 1;
      return;
    int mid = (st+ed)>>1;
    build(n+n,st,mid);
    build(n+n+1, mid+1,ed);
    tree[n] = Merge(tree[n+n], tree[n+n+1]);
 inline void update(int n,int st,int ed,int i,int

    j,int v)

    lazyUpdate(n,st,ed);
    if(st>j or ed<i) return;</pre>
    if(st>=i and ed<=i) {
      lazy[n] = v;
      lazyUpdate(n,st,ed);
      return;
    int mid = (st+ed)>>1;
    update(n+n,st,mid,i,j,v);
    update(n+n+1,mid+1,ed,i,j,v);
    tree[n] = Merge(tree[n+n], tree[n+n+1]);
 inline node query(int n,int st,int ed,int i,int j){
    lazyUpdate(n,st,ed);
    if(st>=i and ed<=j) return tree[n];</pre>
    int mid = (st+ed)/2;
    if(mid<i) return query(n+n+1,mid+1,ed,i,j);</pre>
    else if(mid>=j) return query(n+n,st,mid,i,j);
    else return Merge(query(n+n,st,mid,i,j),query(n+n+

→ 1,mid+1,ed,i,j));

8.4 kmp
// returns the longest proper prefix array of pattern p
// where lps[i]=longest proper prefix which is also
\rightarrow suffix of p[0...i]
vector<int> build lps(string p) {
 int sz = p.size();
  vector<int> lps;
 lps.assign(sz + 1, 0);
  int j = 0;
  lps[0] = 0;
 for(int i = 1; i < sz; i++) {
  while(j >= 0 && p[i] != p[j]) {
      if(j >= 1) j = lps[j - 1];
      else j = -1;
    lps[i] = j;
  return lps;
vector<int>ans;
// returns matches in vector ans in 0-indexed
void kmp(vector<int> lps, string s, string p) {
 int psz = p.size(), sz = s.size();
 int j = 0;
 for(int i = 0; i < sz; i++)
    while(j \ge 0 \&\& p[j] != s[i])
      if(j >= 1) j = lps[j - 1];
      else j = -1;
```

j = lps[j - 1];
// pattern found in string s at position i-psz+1

// after each loop we have j=longest common suffix

 \rightarrow of s[0..i] which is also prefix of p

if(i == psz) {

ans.push back(i - psz + 1);

```
8.5 manachar
vector<int> d1(n): // maximum odd length palindrome
```

vector<int> d1(n); // maximum odd length palindrome // here d1[i]=the palindrome has → d1[i]-1 right characters from if // e.g. for aba, d1[1]=2;for (int i = 0, l = 0, r = -1; i < n; i++) { int k = (i > r) ? 1 : min(d1[1 + r - i], r - i); while (0 <= i - k && i + k < n && s[i - k] == s[i + k]→ k]) { k++; d1[i] = k--;**if** (i + k > r) { l = i - k;r = i + k;vector<int> d2(n); // maximum even length palindrome // here d2[i]=the palindrome has → d2[i]-1 right characters from i // e.g. for abba, d2[2]=2; for (int i = 0, l = 0, r = -1; i < n; i++) { int k = (i > r) ? 0 : min(d2[l + r - i + 1], r - i +while $(0 \le i - k - 1) \& i + k < n \& s[i - k - 1]$ $\Rightarrow == s[i + k]) \{$ d2[i] = k--; $if(i + k > r) {$ l = i - k - 1;r = i + k;

8.6 palindromic_tree

```
const int A = 26;
const int N = 300010;
char s[N]; long long ans;
int last, ptr, nxt[N][A], link[N], len[N], occ[N];
void feed (int at)
  while (s[at - len[last] - 1] != s[at]) last =
     link[last];
  int ch = s[at] - 'a', temp = link[last];
  while (s[at - len[temp] - 1] != s[at]) temp =

→ link[temp];

 if (!nxt[last][ch]) {
  nxt[last][ch] = ++ptr, len[ptr] = len[last] + 2;
    link[ptr] = len[ptr] == 1 ? 2 : nxt[temp][ch];
  last = nxt[last][ch], ++occ[last];
int main() {
  len[1] = -1, len[2] = 0, link[1] = link[2] = 1, last
  for (int i = 1, n = strlen(s + 1); i \le n; ++i)
  → feed(i):
  for (int i = ptr; i > 2; --i) ans = max(ans, len[i]
  * 1LL * occ[i]), occ[link[i]] += occ[i];
  printf("%lld\n", ans);
  return 0;
```

```
8.7 persistant_trie
const int MAX = 200010;
const int B = 19;
|int root[MAX], ptr = 0;
struct node {
  int ara[2], sum;
  node() {
|void insert(int prevnode, int &curRoot, int val) {
  curRoot = ++ptr;
  int curnode = curRoot;
  for(int i = B; i >= 0; i --) {
    bool bit = val & (1 \ll i);
    tree[curnode] = tree[prevnode];
    tree[curnode].ara[bit] = ++ptr;
    tree[curnode].sum += 1
    prevnode = tree[prevnode].ara[bit];
    curnode = tree[curnode].ara[bit];
  tree[curnode] = tree[prevnode];
  tree[curnode].sum += 1;
int find xor max(int prevnode, int curnode, int x) {
  int ans = \overline{0}:
  for(int i = B; i >= 0; i - -) {
    bool bit = \dot{x} \& (1 << i);
    if(tree[tree[curnode].ara[bit ^ 1]].sum >

    tree[tree[prevnode].ara[bit ^ 1]].sum) {

      curnode = tree[curnode].ara[bit ^ 1]
      prevnode = tree[prevnode].ara[bit ^1];
      ans = ans | (1 << i);
    else {
      curnode = tree[curnode].ara[bit]
      prevnode = tree[prevnode].ara[bit];
  return ans:
void solve() {
  int n, q, L, R, K;
  cin >> n;
  for(int i=1;i<=n;i++) cin >> ara[i];
  for(int i=1;i<=q;i++) {
    cin >> L >> R >> K;
    cout << find xor max(root[L-1],root[R],K) << endl;</pre>
```

0.0 66

```
8.8 suffix_array
 // Everything is 0-indexed
char s[N]; // Suffix array will be built for this
int SA[N], iSA[N]; // SA is the suffix array, iSA[i]
stores the rank of the i'th suffix int cnt[N], nxt[N]; // Internal stuff bool bh[N], b2h[N]; // Internal stuff
int lcp[N]; // Stores lcp of SA[i] and SA[i + 1];
 \rightarrow lcp[n - 1] = 0
int lcpSparse[LOGN][N]; // lcpSparse[i][j] =
 \rightarrow min(lcp[j], ..., lcp[j - 1 + (1 << i)])
void buildSA(int n) {
  for (int i = 0; i < n; i++) SA[i] = i;
   sort(SA, SA + n, [](int i, int j) { return s[i] <

    s[i]; });
  for (int i = 0; i < n; i++) { bh[i] = i == 0 | | s[SA[i]] != s[SA[i - 1]];
     b2h[i] = 0;
   for (int h = 1; h < n; h <<= 1) {
```

```
int tot = 0;
    for (int i = 0, j; i < n; i = j) {
      j = i + 1;
      while (j < n && !bh[j]) j++;
      nxt[i] = j; tot++;
    } if (tot == n) break;
    for (int i = 0; i < n; i = nxt[i])
      for (int j = i; j < nxt[i]; j++) iSA[SA[j]] = i;
      cnt[i] = 0;
    cnt[iSA[n - h]]++;
b2h[iSA[n - h]] = 1;
    for (int i = 0; i < n; i = nxt[i])
      for (int j = i; j < nxt[i]; j++) {
        int s = SA[j] - h;
        if (s < 0) continue;
        int head = iSA[s];
iSA[s] = head + cnt[head]++;
b2h[iSA[s]] = 1;
      for (int j = i; j < nxt[i]; j++) {
        int s = SA[j] - h;
        if (s < 0 || | b2h[iSA[s]]) continue;</pre>
        for (int k = iSA[s] + 1; !bh[k] \&\& b2h[k];
         \rightarrow k++) b2h[k] = 0;
    for (int i = 0; i < n; i++) {
      SA[iSA[i]] = i;
      bh[i] [= b2h[i];
  for (int i = 0; i < n; i++) iSA[SA[i]] = i;
void buildLCP(int n) {
 for (int i = 0, k = 0; i < n; i++) {
    if (iSA[i] == n - 1) {
      k = 0;
      lcp[n - 1] = 0;
      continue;
    int j = SA[iSA[i] + 1];
    while (i + k < n \&\& j + k < n \&\& s[i + k] == s[j + k]
    lcp[i\dot{S}\dot{A}[i]]' = k;
    if (k) k--;
void buildLCPSparse(int n) {
  for (int i = 0; i < n; i++) lcpSparse[0][i] = lcp[i];</pre>
  for (int i = 1; i < LOGN; i++) {
    for (int j = 0; j < n; j++) {
      lcpSparse[i][j] = min(lcpSparse[i - 1][j],
          lcpSparse[i - 1][min(n - 1, j + (1 << (i -
       □ 1)))]);
pair<int, int> minLCPRange(int n, int from, int

→ minLCP) {
 int r = from;
  for (int i = LOGN - 1; i >= 0; i--) {
    int jump = 1 << i;
    if (r + jump < n and lcpSparse[i][r] >= minLCP) r

→ += jump;

  int l = from;
 for (int i = LOGN - 1; i >= 0; i--) {
    int jump = 1 << i;
    if (1 - jump >= 0 and lcpSparse[i][l - jump] >=

→ minLCP) l -= jump;
```

```
return make pair(l, r);
8.9 suffix automata
namespace sa{
  const int MAXN = 100005 << 1;</pre>
  const int MAXC = 26;
  char str[MAXN];
  int n, sz, last;
  int len[MAXN], link[MAXN], ed[MAXN][MAXC], cnt[MAXN];
bool terminal[MAXN];
  vector <int> G[MAXN];
  void init() {
   SET(ed[0]);
   len[0] = 0, link[0] = -1, sz = 1, last = 0,

    terminal[0] = false;

  inline int scale(char c) { return c-'a'; }
  void extend(char c) {
    int cur = sz++;
    terminal[cur] = false;
    cnt[cur] = 1;
    SET(ed[cur]);
len[cur] = len[last] + 1;
    int p = last;
    while (p != -1 \&\& ed[p][c]==-1) {
       ed[p][c] = cur;
      p = link[p];
    if (p == -1) link[cur] = 0;
    else {
       int q = ed[p][c];
       if (len[p] + 1 == len[q]) link[cur] = q;
       else {
         int clone = sz++;
len[clone] = len[p] + 1;
         memcpy(ed[clone],ed[q],sizeof(ed[q]));
         link[clone] = link[q];
         while (p != -1 && ed[p][c] == q) {
  ed[p][c] = clone;
           p = link[p];
         link[q] = link[cur] = clone;
         cnt[clone] = 0;
terminal[clone] = false;
     last = cur;
  void dfs(int s) {
    for(auto x : G[s]) dfs(x), cnt[s] += cnt[x];
  void build() {
    init();
    int n = strlen(str);
for(int i=0;i<n;i++) extend(scale(str[i]));</pre>
    for(int i=1;i<sz;i++) G[link[i]].pb(i);</pre>
    dfs(0);
    for(int i=0;i<sz;i++) G[i].clear();</pre>
    for(int i=last;i!=-1;i=link[i]) terminal[i] = true;
8.10 trie
#define N
                   200000
```

```
void init(){
  root = now = 1;
CLR(nxt),CLR(cnt);
inline int scale(char ch) { return (ch - 'a'); }
inline void Insert(char s[],int sz){
  int cur = root, to;
   for(int i=0 ; i < sz ; i++){
     to = scale(s[i]);
if( !nxt[cur][to] ) nxt[cur][to] = ++now;
     cur = nxt[cur][to];
   cnt[cur]++;
inline bool Find(char s[],int sz){
   int cur = root, to;
   for(int i=0 ; i<sz ; i++){
     to = scale(s[i]);
if(!nxt[cur][to]]) return false;
     cur = nxt[cur][to];
   return (cnt[cur]!=0);
inline void Delete(char s[],int sz){
  int cur = root, to;
   for(int i=0 ; i<sz ; i++){
  to = scale(s[i]) ;</pre>
     cur = nxt[cur][to];
   cnt[cur]--;
8.11 z_algo
 const int N = 100010;
 char s[N];
| int t, n, z[N];
int main() {
    scanf("%s", s);
    n = strlen(s), z[0] = n;
   int L = 0, R = 0;
   for (int i = 1; i < n; ++i) {
     if(i > R) {
       L = R = i:
       while (R < n \&\& s[R - L] == s[R]) ++R;
        z[i] = R - L; --R;
     } else {
       int k = i - L;
       if (z[k] < R - i + 1) z[i] = z[k];
       else {
         L = i:
         while (R < n \&\& s[R - L] == s[R]) ++R;
         z[i] = R - L; --R;
   }
```

#define N #define S int root,now;

int nxt[N][S], cnt[N];