

DU_Scolps

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Team Reference Document for NCPC 2024

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	3.1.7 Point	7		3 pragma		voia init() {		
	3.1.8 Segment Distance	5				len = 0`,cur = 0;		
		2		4 random	$\frac{17}{17}$	}		
	3.1.9 Segment Intersection	5	U	o viiii c	11	<pre>inline bool isBad(ll nm,ll nc) {</pre>		
	3.1.10 Side Of	$\frac{5}{5}$	7 N	ata a	17	<pre>inline bool isBad(ll nm,ll nc) { return ((C[len-1]-C[len-2])/(double)(M[len-2]-M[l])</pre>		
	3.2 3D	514		otes		on $ $ $ $ $ $ $ $ $ $ $ $ $ $ $ $		
	3.2.1 3D Convex Hull	5	- 1	1 Counting				
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	3.3 Circle	6		7.4.1 Triangles		inline veid addine/11 nm 11 nc) (
	3.3.1 Circle Intersection	6		7.4.2 Quadrilaterals	18	inline void addLine(ll nm,ll nc) {		
	3.3.2 Circle Polygon Intersection	6		7.4.3 Spherical coordinates	18	<pre>if(len == 0) M[len] = nm , C[len] = nc , ++len; else if(M[len-1] == nm) {</pre>		
	3.3.3 Circle Tangents	6	7	5 Sums		if(C[]on 1] ==) {		
	3.3.4 CircumCircle	6	7	6 Series	18	<pre>if(C[len-1] <= nc) return; else C[len-1] = nc;</pre>		
	3.4 Polygon	Ğl	7	7 Pythagorean Triples	18	etse ([ten-1] = nc;		
	3.4.1 Hull Diameter	6	7	8 Number Theory	18	else {		
	3.4.2 Line Hull Intersection	ĕГ	•	7.8.1 Primes				
	3.4.3 Polygon Center	žΙ		7.8.2 Estimates		<pre>while(len >= 2 &\& isBad(nm,nc))len; M[len] = nm , C[len] = nc , ++len;</pre>		
	3.4.4 Polygon Cut	٠		7.8.3 Perfect numbers	18			
	5.4.4 Torygon Cut	<u> </u>		7.8.4 Carmichael numbers	18 18	\ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \		
	3.5 Closest Pair	4		7.8.5 Mobius function	18			
	9.7 Minimum Englacing Cinels	4		7.8.6 Legendre symbol	$\bar{18}$	<pre>inline ll getY(int id , ll x) { return (M[id]*x + C[id]);</pre>		
	9.0 Delat La Delanasa	41		7.8.7 Jacobi symbol		return (M[10]*X + C[10]);		
	3.8 Point In Polygon	4		7.8.8 Primitive roots	18	}		
	3.9 near_pair	8		7.8.9 Discrete logarithm problem	18	<pre>inline ll sortedQuery(ll x) {</pre>		
	3.10 sweep	8		7.8.10 Pythagorean triples		<pre>if(cur >= len) cur = len-1;</pre>		
				7.0.10 I ymagorean urples	10	while(cur < len-1 && getY(cur+1,x) >= getY(cur,x)		
4	4 Graph	8		7.8.11 Postage stamps/McNuggets problem	18	→) cur++:		
_	4.1 2Sat	g		7.8.12 Fermat's two-squares theorem	18	→) cur++; return getY(cur,x);		
	4.2 Centroid_decomp	äl	7	9 Permutations		l }		
		ă		7.9.1 Factorial	18	inline ll TS(ll x) {		
	4.3 articulation_point	9				int low = 0, high = len-1 , mid ;		
	4.4 bcc	. 9T		7.9.3 Derangements	18	III		
		10		7.9.4 Burnside's lemma	18	WILLE III - LOW / I / 1		
		10	7	10 Partitions and subsets	18	mid = low + high >> 1;		
		10		7.10.1 Partition function	18	\mathbf{I} (get \mathbf{I} (mid, \mathbf{X}) < get \mathbf{I} (mid+1, \mathbf{X}) tow = mid + 1;		
	4.8 euler_path	10	7	11 General purpose numbers	18	else high = mid;		
	4.9 hld	11		7.11.1 Stirling numbers of the first kind	18	l }		
	4.10 hopcroft_karp	11		7.11.2 Eulerian numbers	18	<pre>return max(getY(low,x),getY(high,x));</pre>		
	4.11 hungarian			7.11.3 Stirling numbers of the second kind		<u> </u>		
	0	12						
		12		7.11.4 Bell numbers	10	CHT cht;		
	4.14 manhattan MST			7.11.6 Catalan numbers	19	cht.init();		
	4.15 mcmf		7	12 Inequalities	19			

void compute(int L, int R, int optL, int optR){

1.2 DC_Optimization

if(L > R) return;

int $M = L + R \gg 1$

```
9
```

```
pair<ll, int> best(1LL << 60, -1);
  for(int k = optL; k <= min(M, optR); k++){</pre>
    best = min(best, \{dp[prv][k] + C[k + 1][M], k\});
  dp[now][M] = best.ff;
  compute(L, M - 1, optL, best.ss);
  compute(M + 1, R, best.ss, optR);
1.3 SOS DP
for(int mask = 0; mask < (1<<N); ++mask){
         dp[mask][-1] = A[mask];
         for(int i = 0; i < N; ++i){
             if(mask \& (1 << i)) dp[mask][i] =
              \rightarrow dp[mask][i-1] + dp[mask \land (1<<i)][i-1];
             else dp[mask][i] = dp[mask][i-1];
         f[mask] = dp[mask][N-1];
    for(int i = 0; i < (1 << N); ++i) F[i] = A[i];
    for(int i = 0; i < N; ++i)
         for(int mask = 0; mask < (1 << N); ++mask){
             if(mask & (1<<i))
                  F[mask] += F[mask^(1<<i)];
1.4 dynamic_cht
//add lines with -m and -b and return -ans to
//make this code work for minimums.(not -x)
const ll inf = -(1LL << 62);</pre>
struct line {
  ll m, b;
  mutable function<const line*() > succ;
  bool operator < (const line& rhs) const {</pre>
    if (rhs.b != inf) return m < rhs.m:</pre>
    const line* s = succ();
    if (!s) return 0;
    ll x = rhs.m;
    return b - s->b < (s->m - m) * x;
struct CHT : public multiset<line> {
  bool bad(iterator y) {
    auto z = next(y)
    if (y == begin()) {
       if(z == end()) return 0;
       return y \rightarrow m == z \rightarrow m \&\& y \rightarrow b <= z \rightarrow b;
    auto x = prev(y);
    if (z == end()) return y -> m == x -> m \&\& y -> b
     \rightarrow <= x -> b;
    return 1.0 * (x -> b - y -> b) * (z -> m - y -> m)
     \Rightarrow >= 1.0 * (v -> b - z -> b) * (v -> m - x -> m);
  void add(ll m, ll b) {
    auto y = insert({'m, b });
y->succ = [ = ] { return next(y) == end() ? 0 :
     if (bad(y)) {
       erase(y);
      return;
    while (next(y) != end() \&\& bad(next(y)))
        erase(next(v));
    while (y != begin() \&\& bad(prev(y)))
     → erase(prev(y));
```

```
ll query(ll x) {
     auto 1 = *lower bound((line) {
      x, inf
    return l.m * x + l.b;
ľCHT* cht;
ll a[N], b[N]
|int32 t main() {
  ios base::sync with stdio(0);
  cin_tie(0):
  int n;
  cin >> n;
  for(int i = 0; i < n; i++) cin >> a[i];
  for(int i = 0; i < n; i++) cin >> b[i];
cht = new CHT();
  cht -> add(-b[0], 0);
  ll ans = 0;
  for(int i = 1; i < n; i++)
    ans = -cht \rightarrow query(a[i]);
     cht -> add(-b[i], -ans);
  cout << ans << nl;
  return 0;
|2 Data<sub>S</sub>tructure
2.1 2D_segtree
```

```
struct Point {
    int x, y, mx;
Point() {}
    Point(int x, int y, int mx) : x(x), y(y), mx(mx) {}
    bool operator < (const Point& other) const {</pre>
        return mx < other.mx;</pre>
struct Segtree2d {
    // I didn't calculate the exact size needed in
    terms of 2D container size.
// If anyone, please edit the answer.
    // It's just a safe size to store nodes for MAX *
        MAX 2D grids which won't cause stack overflow
    Point T[500000]; // TODO: calculate the accurate

→ space needed
    int n, m;
    // initialize and construct segment tree
    void init(int n, int m) {
         this \rightarrow n = n;
         this -> m = m:
         build(1, 1, 1, n, m);
    // build a 2D segment tree from data [ (a1, b1),
     // Tìme: O(n logn)
    Point build(int node, int a1, int b1, int a2, int
     → b2) {
         // out of range
        if (a1 > a2 or b1 > b2)
             return def();
         // if it is only a single index, assign value
             to node
        if (a1 == a2 and b1 == b2)
   return T[node] = Point(a1, b1, P[a1][b1]);
           split the tree into four segments
         T[node] = def();
         T[node] = maxNode(T[node], build(4 * node - 2,
         \rightarrow a1, b1, (a1 + a2) / 2, (b1 + b2) / 2 ));
```

```
T[node] = maxNode(T[node], build(4 * node - 1,
         (a1 + a2) / 2 + 1, b1, a2, (b1 + b2) / 2
    T[node] = maxNode(T[node], build(4 * node + 0,
        a1, (b1 + b2) / 2 + 1, (a1 + a2) / 2, b2)
    T[node] = maxNode(T[node], build(4 * node + 1,
         (a1 + a2) / 2 + 1, (b1 + b2) / 2 + 1, a2,
        b2)):
    return T[node];
// helper function for query(int, int, int, int);
Point query(int node, int al, int bl, int a2, int
\rightarrow b2, int x1, int y1, int x2, int y2) {
    // if we out of range, return dummy
    if (x1 > a2 \text{ or } y1 > b2 \text{ or } x2 < a1 \text{ or } y2 < b1
    \rightarrow or a1 > a2 or b1 > b2)
        return def();
    // if it is within range, return the node
    if (x1 \le a1 \text{ and } y1 \le b1 \text{ and } a2 \le x2 \text{ and } b2
    <= y2)
         return T[node];
    // split into four segments
    Point mx = def();
    mx = maxNode(mx, query(4 * node - 2, a1, b1,
         (a1 + a2) / 2, (b1 + b2) / 2, x1, y1, x2,
    mx = maxNode(mx, query(4 * node - 1, (a1 + a2))
        /2 + 1, b1, a2, (b1 + b2) /2, x1, y1,
     \equiv x2. v2) ):
    mx = maxNode(mx, query(4 * node + 0, a1, (b1 +
        b2) / 2 + 1, (a1 + a2) / 2, b2, x1, y1,
     \stackrel{\sim}{\rightarrow} x2, y2));
    mx = maxNode(mx, query(4 * node + 1, (a1 + a2))
        /2 + 1, (b1 + b2) /2 + 1, a2, b2, x1,

    y1, x2, y2));

    // return the maximum value
    return mx;
// query from range [ (x1, y1), (x2, y2) ]
// ˈtimeː O(logn)
Point query(int x1, int y1, int x2, int y2)
    return query(1, 1, 1, n, m, x1, y1, x2, y2);
// helper function for update(int, int, int);
Point update(int node, int al, int bl, int a2, int
→ b2, int x, int y, int value) {
    if (a1 > a2 \text{ or } b1 > b2)
        return def();
    if (x > a2 \text{ or } y > b2 \text{ or } x < a1 \text{ or } y < b1)
        return T[node];
    if (x == a1 \text{ and } y == b1 \text{ and } x == a2 \text{ and } y ==

→ b2)

         return T[node] = Point(x, y, value);
    Point mx = def();
    mx = maxNode(mx, update(4 * node - 2, a1, b1,
         (a1 + a2) / 2, (b1 + b2) / 2, x, y, value)
    ≒ );
    mx = maxNode(mx, update(4 * node - 1, (a1 +
         a2) / 2 + 1, b1, a2, (b1 + b2) / 2, x, y,

    value))

    mx = maxNode(mx, update(4 * node + 0, a1, (b1))
        + b2) / 2 + 1, (a1 + a2) / 2, b2, x, y,

    value))

    mx = maxNode(mx, update(4 * node + 1, (a1 +
        a2) / 2 + 1, (b1 + b2) / 2 + 1, a2, b2, x,
        y, value) );
    return T[node] = mx;
```

```
// update the value of (x, y) index to 'value'
    // Time: 0(logn)
    Point update(int x, int y, int value) {
        return update(1, 1, 1, n, m, x, y, value);
    // utility functions; these functions are virtual
    → because they will be overridden in child class
    virtual Point maxNode(Point a, Point b) {
        return max(a, b);
    // dummy node
   virtual Point def() {
    return Point(0, 0, -INF);
/* 2D Segment Tree for range minimum query; a override

→ of Segtree2d class */
struct Segtree2dMin : Segtree2d {
    // overload maxNode() function to return minimum
    Point maxNode(Point a, Point b) {
        return min(a, b);
    Point def() {
        return Point(0, 0, INF);
```

2.2 Segtree_beats_desc

```
Description: For update Ai Ai mod x and similar, keep
max in node and lazily update whenever min = max. For
→ update
Ai min(Ai, x) and similar, keep range max, second max
    in node and
lazily update whenever x > second max.
Time: O(\log^2 N), (\log N)
```

2.3 gp hash table

```
using namespace gnu pbds;
const int RANDOM = chrono::high resolution clock::now(
→ ).time since epoch().count();
using namespace __gnu_pbds;
struct chash
  const int RANDOM = (long
      long)(make unique<char>().get()) ^
      chrono::high resolution clock::now().time since
      epoch().count();
  static unsigned long long hash f(unsigned long long
    x += 0x9e3779b97f4a7c15;
   x = (x ^ (x >> 30)) * 0xbf58476d1ce4e5b9;
x = (x ^ (x >> 27)) * 0x94d049bb133111eb;
return x ^ (x >> 31);
  static unsigned hash combine(unsigned a, unsigned b)
      { return a * 31 + b; }
  ll operator()(ll x) const { return hash f(x)^RANDOM;
  → }
qp hash table<key, long long, chash> table;
```

2.4 iterative_segtree

```
const int N = 500010;
int n, a[N], tree[N << 1];</pre>
void init() {
```

```
for (int i = 0; i < n; ++i) tree[n + i] = a[i];</pre>
    for (int i = n - 1; i >= 0; --i)
        tree[i] = min(tree[i << 1], tree[i << 1 | 1]);
void update(int p, int v) {
    for (tree[p += n] = v; p > 1; p >>= 1)
        tree[p >> 1] = min(tree[p], tree[p ^{\uparrow} 1]);
|int query(int l, int r) {
    int ret = INT MAX:
    for (l += n, r += n; l < r; l >>= 1, r >>= 1) {
         if (l & 1) ret = min(ret, tree[l++]);
        if (r \& 1) ret = min(ret, tree[--r]);
    return ret;
2.5 mos algo
```

```
struct data{
  int l, r, id, bn;
data() {}
  data(int l, int r, int id){
    l = l, r = r, id = i\overline{d};
    bn = l / block sz;
  bool operator < (const data& other) const{</pre>
    if (bn != other.bn) return (bn < other.bn);</pre>
    return ((bn \& 1) ? (r < other.r) : (r > other.r));
int curL = 0, curR = -1;
for(int i = 0; i < 0.sz; i++){
  while(curL > Q[i].L){
    curl--; add(curl);
  while(curR < 0[i].R){</pre>
    curR++; add(curR);
  while(curL < Q[i].L){
    remove(curL); curL++;
  while(curR > Q[i].R){
    remove(curR); curŔ--;
```

2.6 ordered set

```
using namespace std:
using namespace gnu pbds;
typedef tree<
    int
    less < int > , // "less equal<int>," for multiset
    rb tree tag,
    tree order statistics node update > ordered_set;
ordered set OS;
```

2.7 persistant_segtree

```
const int MAX = 100010;
lint ncnt = 0;
|struct node {
  int sum:
  int left, right;
  node() {}
  node(int val) {
    sum = val:
    left = right = -1;
} tree[ ? ];
```

```
int ara[MAX];
int version[MAX];
void build(int n.int st.int ed) {
  if (st==ed) {
    tree[n] = node(ara[st]);
    return;
  int mid = (st+ed) / 2;
  tree[n].left = ++ncnt;
  tree[n].right = ++ncnt;
  build(tree[n].left, st, mid);
  build(tree[n].right, mid+1, ed);
  tree[n].sum = tree[tree[n].left].sum +

    tree[tree[n].right].sum;

void update(int prev,int cur,int st,int ed,int id, int
   val)
  if (id > ed or id < st) return;</pre>
  if (st == ed) {
    tree[cur] = node(val);
    return;
  int mid = (st+ed) / 2;
  if (id <= mid) ·
    tree[cur].right = tree[prev].right;
    tree[cur].left = ++ncnt;
    update(tree[prev].left,tree[cur].left, st, mid,
     \rightarrow id, val);
  else {
    tree[cur].left = tree[prev].left;
    tree[cur].right = ++ncnt;
    update(tree[prev].right, tree[cur].right, mid+1,
     → ed. id. val):
  tree[cur].sum = tree[tree[cur].left].sum +

→ tree[tree[cur].right].sum;

int query(int n,int st,int ed,int i,int j){
 if(st>=i && ed<=j) return tree[n].sum;</pre>
  int mid = (st+ed)/2;
  if(mid<i) return query(tree[n].right,mid+1,ed,i,j);</pre>
  else if(mid>=j) return

¬ query(tree[n].left,st,mid,i,j);
  else return query(tree[n].left,st,mid,i,j) +

¬ query(tree[n].right,mid+1,ed,i,j);

int main() {
 int n,q,l,r,k;
  sii(n,q);
  version[0] = ++ncnt;
  build(version[0],1,n);
  version[1] = ++ncnt;
  update(version[0], version[1], 1, n, id, val);
 query(version[0],1,n,id,id);
query(version[1],1,n,id,id);
  return 0;
```

2.8 segment_tree

```
int ara[MAX];
struct node {
 int sum:
} tree[4 * MAX];
int lazy[4 * MAX];
node Merge(node a, node b) {
```

```
node ret;
  ret.sum = a.sum + b.sum;
  return ret;
void lazyUpdate(int n, int st, int ed) {
  if(lazy[n] != 0){
    tree[n].sum += ((ed - st + 1) * lazv[n]);
    if(st != ed){
  lazy[2 * n] += lazy[n];
      lazy[2 * n + 1] += lazy[n];
    lazy[n] = 0;
void build(int n, int st, int ed) {
  lazy[n] = 0;
  if(st == ed){
    tree[n].sum = ara[st];
    return:
 int mid = (st + ed) / 2;
build(2 * n, st, mid);
build(2 * n + 1, mid + 1, ed);
  tree[n] = Merge(tree[2 * n], tree[2 * n + 1]);
void update(int n, int st, int ed, int i, int j, int
 - ν) {
  lazyUpdate(n, st, ed);
  if(st > j or ed < i) return;</pre>
  if(st >= i and ed <= j){
    lazy[n] += v;
    lazyUpdate(n, st, ed);
    return;
  int mid = (st + ed) / 2;
  update(2 * n, st, mid, i, j, v);
  update(2 * n + 1, mid+1, ed, i, j, v);
  tree[n] = Merge(tree[2 * n], tree[2 * n + 1]);
node query(int n, int st, int ed, int i, int j) {
  lazyUpdate(n, st, ed);
  if(st >= i and ed <= j) return tree[n];</pre>
  int mid = (st + ed) / 2;
  if (mid < i) return query (2 * n + 1, mid + 1, ed, i,
  else if(mid >= j) return query(2 * n, st, mid, i, j)
  else return Merge(query(2 * n, st, mid, i, j),
  \rightarrow query(2 * n + 1, mid + 1, ed, i, j));
2.9 sparse_table
```

```
int st[K + 1][MAXN];
void build() -
  std::copy(array.begin(), array.end(), st[0]);
  for (int i = 1; i <= K; i++)
    for (int j = 0; j + (1 << i) <= N; j++)
      st[i][j] = f(st[i - 1][j], st[i - 1][j + (1 <<
       \hookrightarrow (i - 1)));
```

3 Geometry

3.1 2D Primitive

A class for ordering angles (as represented by int points and a number so watch out for overflow if using int or ll. of rotations around the origin). Useful for rotational sweeping. Sometimes also represents points or vectors.

```
/* Usage:
 * vector<Angle> v = {w[0], w[0].t360() ...}; // sorted
* int j = 0; rep(i,0,n) { while (v[j] < v[i].t180()) + + f/ }
```

```
struct Angle {
  int x, y;
  int t;
  Angle(int x, int y, int t=0) : x(x), y(y), t(t) {}
  Angle operator-(Angle b) const { return {x-b.x, y-b.y,}
  int half() const {
   assert(x || y);
   return y < 0 \mid | (y == 0 \&\& x < 0);
  Angle t90() const { return \{-y, x, t + (half() \&\& x > \}
  Angle t180() const { return \{-x, -y, t + half()\}; }
  Angle t360() const { return \{x, y, t + 1\}; \}
|bool operator<(Angle a, Angle b) {
 // add a.dist2() and b.dist2() to also compare distandetypedef Point<double> P;
// Given two points, this calculates the smallest angle bereturn q0 + P((r-p0).cross(num), (r-p0).dot(num))/dp.di
// them, i.e., the angle that covers the defined line sedment.
pair<Angle, Angle> segmentAngles(Angle a, Angle b) {
 if (b < a) swap(a, b);
  return (b < a.t180() ? make pair(a, b) : make pair(b,</pre>
Angle r(a.x + b.x, a.y + b.y, a.t);
  if (a.t180() < r) r.t--;
  return r.t180() < a ? r.t360() : r;</pre>
Angle angleDiff(Angle a, Angle b) { // angle b - angle 43.1.6 Point Sort
 int tu = b.t - a.t; a.t = b.t;
 Returns the signed distance between point p and the
line containing points a and b. Positive value on left side
and negative on right as seen from a towards b. a==b
gives nan. P is supposed to be Point<T> or Point3D<T>
where T is e.g. double or long long. It uses products
```

in intermediate steps so watch out for overflow if using int or long long. Using Point3D will always give a non-/s negative distance. For Point3D, call .dist on the result of the cross product.

```
#include "Point.h"
template<class P>
double lineDist(const P& a, const P& b, const P& p) {
 return (double)(b-a).cross(p-a)/(b-a).dist();
```

3.1.3 Line Intersection

If a unique intersection point of the lines going through s1,e1 and s2,e2 exists {1, point} is returned. If no intersection point exists $\{0, (0,0)\}$ is returned and if infinitely many exists {-1, (0,0)} is returned. The wrong position will be returned if P is Point<ll> and the intersection point does not have integer coordinates. Prod- \(\sigma_1 \) ucts of three coordinates are used in intermediate steps

```
/* Usage:
  auto res = lineInter(s1,e1,s2,e2);
   if (res.first == 1)
      cout << "intersection point at " << res.second <<</pre>
```

```
#include "Point.h"
                                                          template<class P>
                                                          pair<int, P> lineInter(P s1, P e1, P s2, P e2) {
                                                           auto d = (e1 - s1).cross(e2 - s2);
                                                           if (d == 0) // if parallel
                                                          t; return {-(s1.cross(e1, s2) == 0), P(0, 0)};
                                                            auto p = s2.cross(e1, e2), q = s2.cross(e2, s1);
                                                            return \{1, (s1 * p + e1 * q) / d\};
                                                          3.1.4 Linear Transformation
Apply the linear transformation (translation, rotation po
                                                          and scaling) which takes line p0-p1 to line q0-q1 to point q0
                                                          #include "Point.h"
return make tuple(a.t, a.half(), a.y * (ll)b.x) < makePtupheebTtanbfbahafijonacomnst(Ph)pOy)const P& p1, const P&
                                                           P dp = p1-p0, dq = q1-q0, num(dp.cross(dq), dp.dot(dq))
                                                          3.1.5 On Segment
                                                         #include "Point.h"
                                                          template<class P> bool onSegment(P s, P e, P p) {
                                                            return p.cross(s, e) == 0 \& \& (s - p).dot(e - p) <= 0;
                                                          using namespace std;
                                                          typedef long long ll;
                                                          typedef pair <ll, ll> point;
                                                          #define x first
                                                          #define y second
                                                          int main() {
                                                           int n; cin >> n;
                                                            vector <point> p(n);
                                                            for (auto &it : p) scanf("%lld %lld", &it.x, &it.y);
                                                           sort(p.begin(), p.end(), [] (point a, point b) {
  return atan2l(a.y, a.x) < atan2l(b.y, b.x);</pre>
                                                            for (auto it : p) printf("%lld %lld\n", it.x, it.y);
                                                            return 0;
                                                          3.1.7 Point
```

```
// Class to handle points in the plane. T can be e.g.

→ double or long long. (Avoid int.)

                                                                                                                                              template <class T> int sqn(T x) \{ return (x > 0) - (x < 0) \}
                                                                                                                                               templaté<class T>
                                                                                                                                              struct Point {
                                                                                                                                               typedef Point P;
                                                                                                                                               T'x, y;
                                                                                                                                               explicit Point(T x=0, T y=0) : x(x), y(y) {}
                                                                                                                                               bool operator<(P p) const { return tie(x,y) <</pre>
                                                                                                                                               \begin{array}{c} \rightarrow & \text{tie}(p.x,p.y)\,; \\ \text{bool} & \text{operator} {==}(P\ p) \\ \end{array} \text{const} \ \{ \ \text{return} \\
                                                                                                                                                    tie(x,y)==tie(p.x,p.y); 
                                                                                                                                               P.operator+(P p) const { return P(x+p.x, y+p.y);
                                                                                                                                              epd operator-(P p) const { return P(x-p.x, y-p.y); }
P operator*(T d) const { return P(x*d, y*d); }
* // sweeps j such that (j-i) represents the number of#ፀይያዊዊላይቢናe oriented triangles with vertices at 0 and i | P operator/(T d) const { return P(x/d, y/d); }
```

```
5
```

```
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   dot(P p) const { return x*p.x + y*p.y; }
                                                                 if (onSegment(a, b, d)) s.insert(d);
 T cross(P'p) const { return x*p.y - y*p.x; }
                                                                 return {all(s)};
 T cross(P a, P b) const { return
     (a-*this).cross(b-*this); }
                                                              3.1.10 Side Of
 T dist2() const { return x*x + y*y;
 double dist() const { return sqrt((double)dist2()); }
 // angle to x-axis in interval [-pi, pi]
 double angle() const { return atan2(y, x); }
 P unit() const { return *this/dist(); } // makes
 P perp() const { return P(-y, x); } // rotates +90
  → degrees
 P normal() const { return perp().unit(); }
 // returns point rotated 'a' radians ccw around the
 P rotate(double a) const {
  return \dot{P}(x*\cos(a)-y*\sin(a),x*\sin(a)+y*\cos(a)); }
 friend ostream& operator<<(ostream& os, P p) {</pre>
  return os << "(" << p.x << "," << p.y << ")"; }
3.1.8 Segment Distance
Returns the shortest distance between point p and the
line segment from point s to e.
/* Usage:
 * Point<double> a, b(2,2), p(1,1);
    bool on Segment = segDist(a,b,p) < 1e-10;
 * Status: tested
#pragma once
#include "Point.h"
typedef Point<double> P;
double segDist(P\& s, P\& e, P\& p) {
  if (s==e) return (p-s).dist();
  return ((p-s)*d-(e-s)*t).dist()/d;
3.1.9 Segment Intersection
If a unique intersection point between the line segments
going from s1 to e1 and from s2 to e2 exists then it is
returned. If no intersection point exists an empty vec-
tor is returned. If infinitely many exist a vector with
2 elements is returned, containing the endpoints of the e2.
common line segment. The wrong position will be re-
turned if P is Point<ll> and the intersection point does
not have integer coordinates. Products of three coordi-
nates are used in intermediate steps so watch out for
overflow if using int or long long.
/* Usage:
 * vector<P> inter = segInter(s1,e1,s2,e2);
 * if (sz(inter)==1)
 * if (sz(inter)==1)
* cout << "segments intersect at " << inter[0] << end[∤] It also stores an outward-facing normal vector g
 * Status: stress-tested, tested on kattis:intersection|struct face {
#include "Point.h"
#include "OnSegment.h"
```

```
auto oa = c.cross(d, a), ob = c.cross(d, b),
     oc = a.cross(b, c), od = a.cross(b, d);
// Checks if intersection is single non-endpoint point vector<face> hull3(const vector<pt3> \&p) {
if (sgn(oa) * sgn(ob) < 0 \&\& sgn(oc) * sgn(od) < 0)
 return {(a * ob - b * oa) / (ob - oa)};
set<P> s;
if (onSegment(c, d, a)) s.insert(a);
if (onSegment(c, d, b)) s.insert(b);
if (onSegment(a, b, c)) s.insert(c);
```

```
Returns where p is as seen from s towards e. 1/0/-1 \Leftrightarrow \text{left/on}
                                                              lline/right. If the optional argument eps is given 0 is returned if p
                                                              is within distance eps from the line. P is supposed to be Point<T>
                                                              where T is e.g. double or long long. It uses products in intermediate
                                                              steps so watch out for overflow if using int or long long.
                                                               * bool left = sideOf(p1,p2,q)==1;
                                                               * Status: tested
                                                              #include "Point.h"
                                                              template<class P>
                                                              int sideOf(P s, P e, P p) { return sqn(s.cross(e, p));
                                                              template<class P>
                                                              int sideOf(const P& s, const P& e, const P& p, double eps) {
                                                                auto a = (e-s).cross(p-s);
                                                                double l = (e-s).dist()*eps;
                                                                return (a > l) - (a < -l);
                                                              3.2 3D
                                                              3.2.1 3D Convex Hull
                                                              #define ll long long
                                                              #define sz(x) ((int) (x).size())
#define all(x) (x).begin(), (x).end()
                                                              #define vi vector<int>
                                                              #define pii pair<int, int>
                                                              |\#define'rep(i, a, b)'for(int i = (a); i < (b); i++)|
                                                              using namespace std;
                                                              template<typename T>
  auto d = (e-s).dist2(), t = min(d,max(.0,(p-s).dot(e-<math>\frac{1}{2})), \frac{1}{2} minpq = priority queue<T, vector<T>, greater<T>>;
                                                              typedef long double ftype;
                                                              struct pt3 -
                                                                ftype x, y, z;
                                                                pt3(ftype x = 0, ftype y = 0, ftype z = 0) : x(x),
                                                                    y(y), z(z) \{ \}
                                                                pt3 operator-(const pt3 &o) const {
                                                                   return pt3(x - 0.x, y - 0.y, z - 0.z);
                                                                pt3 cross(const pt3 &o) const {
                                                                   return pt3(y * o.z - z * o.y, z * o.x - x * o.z, x
                                                                   \rightarrow * 0.V - V * 0.X);
                                                                ftype dot(const pt3 &o) const {
                                                                  return x * 0.x + y * 0.y + z * 0.z;
                                                              // A face is represented by the indices of its three
                                                                int a, b, c;
pt3 q;
template<class P> vector<P> segInter(P a, P b, P c, P d)/{ modify this depending on the coordinate sizes in
                                                                  vour use case
                                                              const ftype EPS = 1e-9;
                                                                int n = sz(p);
                                                                assert(n >= 3);
                                                                vector<face> f;
                                                                // Consider an edge (a->b) dead if it is not a CCW
                                                                     edge of some current face
                                                                // If an edge is alive but not its reverse, this is

→ an exposed edge.
```

```
// We should add new faces on the exposed edges.
vector<vector<bool>> dead(n, vector<bool>(n, true));
auto add face = [\&](int a, int b, int c)
  f.push_back({a, b, c, (p[b] - p[a]).cross(p[c] -
     p[a])});
  dead[a][b] = dead[b][c] = dead[c][a] = false;
// Initialize the convex hull of the first 3 points
// triangular disk with two faces of opposite
    orientation
add face(0, 1, 2);
add face(0, 2, 1);
rep(i, 3, n) {
   // f2 will be the list of faces invisible to the
   → added point p[i]
  vector<face> f2;
  for(face &F : f)
    if((p[i] - p[F.a]).dot(F.q) > EPS) {
      // this face is visible to the new point, so
          mark its edges as dead
      dead[F.a][F.b] = dead[F.b][F.c] =

    dead[F.c][F.a] = true;

    }else {
      f2.push back(F);
  // Add a new face for each exposed edge
  // Only check edges of alive faces for being
   → exposed.
  f.clear();
for(face &F : f2)
    int arr[3] = {F.a, F.b, F.c};
    rep(j, 0, 3)
      int a = arr[j], b = arr[(j + 1) % 3];
      if(dead[b][a]) {
        add face(b, a, i);
  f.insert(f.end(), all(f2));
return f;
```

3.2.2 **Point3D**

Class to handle points in 3D space. T can be e.g. double or long long.

```
template<class T> struct Point3D {
 typedef Point3D P;
 typedef const P& R;
 T x, y, z;
 explicit Point3D(T x=0, T y=0, T z=0) : x(x), y(y), z(z)
 bool operator<(R p) const {</pre>
   return tie(x, y, z) < tie(p.x, p.y, p.z); }
 bool operator==(R p) const {
   return tie(x, y, z) == tie(p.x, p.y, p.z); }
 P operator+(R p) const { return P(x+p.x, y+p.y, z+p.z);
 P operator-(R p) const { return P(x-p.x, y-p.y, z-p.z);
 P operator*(T d) const { return P(x*d, y*d, z*d); }
 P operator/(T d) const { return P(x/d, y/d, z/d); }
 T dot(R p) const { return x*p.x + y*p.y + z*p.z; }
 P cross(R p) const {
   return P(y*p.z - z*p.y, z*p.x - x*p.z, x*p.y - y*p.x)
 T dist2() const { return x*x + y*y + z*z; }
 double dist() const { return sqrt((double)dist2()); }
 //Azimuthal angle (longitude) to x-axis in interval [-r
```

```
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```

```
double phi() const { return atan2(y, x); }
  //Zenith angle (latitude) to the z-axis in interval [\mathbf{q}],
  double theta() const { return atan2(sqrt(x*x+y*y),z); |}
  P unit() const { return *this/(T)dist(); } //makes dist()=auto det = a * a - b;
  //returns unit vector normal to *this and p
  P normal(P p) const { return cross(p).unit(); }
  //returns point rotated 'angle' radians ccw around axis
  P rotate(double angle, P axis) const {
    return u*dot(u)*(1-c) + (*this)*c - cross(u)*s;
};
3.2.3 Polyhedron Volume
Description: Magic formula for the volume of a polyhedron. Faces
should point outwards.
template<class V, class L>
```

3.2.4 Spherical Distance

double v = 0;

return v / 6:

Returns the shortest distance on the sphere with radius radius between the points with azimuthal angles (longitude) f1 (ϕ_1) and f2 (ϕ_2) from x axis and zenith angles (latitude) t1 (θ_1) and t2 (θ_2) from z axis (0 = north pole). All angles measured in radians. The algorithm starts by converting the spherical coordinates to cartesian coordinates so if that is what you have you can use only the two last rows. dx*radius is then the difference between the two points in the x direction and d*radius is the total distance between the points.

```
double sphericalDistance(double f1, double t1,
   double f2, double t2, double radius) {
  double dx = \sin(t2)*\cos(f2) - \sin(t1)*\cos(f1);
 double dy = sin(t2)*sin(f2) - sin(t1)*sin(f1);
 double dz = cos(t2) - cos(t1);
 double d = sqrt(dx*dx + dy*dy + dz*dz);
 return radius*2*asin(d/2);
```

3.3 Circle 3.3.1 Circle Intersection

Computes the pair of points at which two circles intersect. Returns the center of the same circle. false in case of no intersection.

```
#include "Point.h"
typedef Point<double> P;
bool circleInter(P a, P b, double r1, double r2, pair<P, P>*
 if (a == b) { assert(r1 != r2): return false: }
  P \text{ vec} = b - a;
  double d2 = \text{vec.dist2}(), sum = r1+r2, dif = r1-r2,
         p = (d2 + r1*r1 - r2*r2)/(d2*2), h2 = r1*r1 -
  if (sum*sum < d2 || dif*dif > d2) return false;
  P mid = a + vec*p, per = vec.perp() * sqrt(fmax(0, h2))/ d2);
  *out = {mid + per, mid - per};
  return true;
```

3.3.2 Circle Polygon Intersection

Returns the area of the intersection of a circle with a ccw polygon. Time: O(n)

```
#include "Point.h"
typedef Point<double> P;
#define arg(p, q) atan2(p.cross(q), p.dot(q))
double circlePoly(P c, double r, vector<P> ps) {
  auto tri = [\&](Pp, Pq) {
```

```
auto r2 = r * r / 2;
 piPd=q-p;
    auto a = d.dot(p)/d.dist2(), b = (p.dist2()-r*r)/d.dist2(); break;
    if (det <= 0) return arg(p, q) * r2;</pre>
    auto s = max(0., -a-sqrt(det)), t = min(1., -a+sqrt(det));
   if (t < 0 \mid | 1 \le s) return arg(p, q) * r2;
    P u = p + d * s, v = p + d * t;
  auto sum = 0.0;
  rep(i,0,sz(ps))
    sum += tri(ps[i] - c, ps[(i + 1) % sz(ps)] - c);
  return sum;
3.3.3 Circle Tangents
```

Finds the external tangents of two circles, or internal if r2 is negated. (i,i) if along side (i,i+1), double signedPolyVolume(const V& p, const L& trilist) { | Can return 0, 1, or 2 tangents - 0 if one circle contains the other (or overlaps it, in the internal case, or if the circles are the same); 1 if the (i,j) if crossing sides (i,i+1) and (j,j+1). In the last case, if a corner ifor (auto i : trilist) v += p[i.a].cross(p[i.b]).dot(p[ircle]) are tangent to each other (in which case .first = .second and the tangent line is perpendicular to the line between the centers). .first and second give the tangency points at circle 1 and 2 respectively. To find the tangents of a circle with a point set r2 to 0.

```
#include "Point.h"
template<class P>
vector<pair<P, P>> tangents(P c1, double r1, P c2, doubl#dE≹in& cmp(i,j) sgn(dir.perp().cross(poly[(i)%n]-poly[(j
  P d = c2 - c1:
  double dr = r1 - r2, d2 = d.dist2(), h2 = d2 - dr * dr + dr + template + class P> int extrVertex(vector + P>+ poly, P dir)
  if (d2 == 0 | | h2 < 0) return {};
  vector<pair<P, P>> out;
  for (double sign : {-1, 1}) {
    P v = (d * dr + d.perp() * sqrt(h2) * sign) / d2;
    out.push back(\{c1 + v * r1, c2 + v * r2\});
  if (h2 == 0) out.pop back();
  return out:
3.3.4 CircumCircle
```

The circumcirle of a triangle is the circle intersecting all three vertices. ccRadius returns the radius of the circle going through points A, B and C and ccCenter returns



rep(i,0,2) {

return res;

```
#include "Point.h"
 typedef Point<double> P;
 double ccRadius (const P& A, const P& B, const P& C) {
 Oμετμέπ (B-A).dist()*(C-B).dist()*(A-C).dist()/
       abs((B-A).cross(C-A))/2;
   ccCenter(const P& A, const P& B, const P& C) {
p^*ppd_b^2; = C-A, c = B-A;
   return A + (b*c.dist2()-c*b.dist2()).perp()/b.cross(c)/2;
 3.4 Polygon
 3.4.1 Hull Diameter
 Returns the two points with max distance on a convex hull (ccw, no
 duplicate/collinear points).
```

```
#include "Point.h"
typedef Point<ll> P:
array<P, 2> hullDiameter(vector<P> S) {
 int n = sz(S), j = n < 2 ? 0 : 1;
  pair<ll, array<P, 2>> res({0, {S[0], S[0]}});
  rep(i,0,j)
   for (;; j = (j + 1) % n) {
```

```
res = \max(res, \{(S[i] - S[j]).dist2(), \{S[i], S[j]\}\}
    if ((S[(j + 1) % n] - S[j]).cross(S[i + 1] - S[i])
return res.second:
```

double s = sin(angle), c = cos(angle); P u = axis.ufit() return arg(p,u) * r2 + u.cross(v)/2 + arg(v,q) * r2 | Line-convex polygon intersection. The polygon must be ccw and have no collinear points. lineHull(line, poly) returns a pair describing the intersection of a line with the polygon:

```
3.4.2 Line Hull Intersection
(-1,-1) if no collision,
(i,-1) if touching the corner i,
is crossed, this is treated as happening on side (i, i+1). The points are
returned in the same order as the line hits the polygon. extrVertex
returns the point of a hull with the max projection onto a line.
Time: O(\log n)
#include "Point.h"
#define extr(i) cmp(i + 1, i) \geq 0 && cmp(i, i - 1 + n) <
  int n = sz(poly), lo = 0, hi = n;
  if (extr(0)) return 0;
  while (lo + 1 < hi) {
    int m = (lo + hi) / 2;
    if (extr(m)) return m:
    int ls = cmp(lo + 1, lo), ms = cmp(m + 1, m);
    (ls < ms \mid | (ls == ms \&\& ls == cmp(lo, m)) ? hi : lo)
  return lo;
#define cmpL(i) sqn(a.cross(poly[i], b))
template <class P>
array<int, 2> lineHull(P a, P b, vector<P>& poly) {
 int endA = extrVertex(poly, (a - b).perp());
 int endB = extrVertex(poly, (b - a).perp());
 if (cmpL(endA) < 0 \mid \mid cmpL(endB) > 0)
    return {-1, -1};
  array<int, 2> res;
```

int lo = endB, hi = endA, n = sz(poly);

if (res[0] == res[1]) **return** {res[0], -1};

case 0: return {res[0], res[0]};

case 2: **return** {res[1], res[1]};

(cmpL(m) == cmpL(endB) ? lo : hi) = m;

int m = ((lo + hi + (lo < hi ? 0 : n)) / 2) % n;

switch ((res[0] - res[1] + sz(poly) + 1) % sz(poly))

while ((lo + 1) % n != hi) {

res[i] = (lo + !cmpL(hi)) % n;

if (!cmpL(res[0]) && !cmpL(res[1]))

swap(endA, endB);

```
3.4.3 Polygon Center
Returns the center of mass for a polygon.
Time: O(n)
#include "Point.h"
typedef Point<double> P;
P polygonCenter(const vector<P>& v) {
  P res(0, 0); double A = 0;
  for (int i = 0, j = sz(v) - 1; i < sz(v); j = i++) {
    res = res + (v[i] + v[j]) * v[j].cross(v[i]);
    A += v[j].cross(v[i]);
  return res / A / 3;
3.4.4 Polygon Cut
Returns a vector with the vertices of a polygon with ev-
erything to the left of the line going from s to e cut away.
/* Usage:
 * vector<P> p = ...;
 * p = polygonCut(p, P(0,0), P(1,0));
 * Status: tested but not extensively
#include "Point.h"
#include "lineIntersection.h"
typedef Point<double> P;
vector < P > polygonCut(const vector < P > & poly, P s, P e) { | typedef pair < ld, ld > point; }
  vector<P> res;
  rep(i,0,sz(poly)) {
    P cur = poly[i], prev = i ? poly[i-1] : poly.back() point operator + (const point \delta a, const point \delta b) {
    bool side = s.cross(e, cur) < 0;
    if (side != (s.cross(e, prev) < 0))
      res.push back(lineInter(s, e, cur, prev).second);
    if (side)
      res.push back(cur);
  return res;
3.5 Closest Pair
Finds the closest pair of points.
Time: O(n \log n)
#include "Point.h"
typedef Point<ll> P;
pair<P, P> closest(vector<P> v) {
  assert(sz(v) > 1);
  set<P> S;
  sort(all(v), [](P a, P b) { return a.y < b.y; });
  pair<ll, pair<P, P>> ret{LLONG MAX, {P(), P()}};
  int j = 0;
  for (P p : v) {
    P d{1 + (ll)sqrt(ret.first), 0};
    while (v[j].y \ll p.y - d.x) S.erase(v[j++]);
    auto lo = S.lower bound(p - d), hi = S.upper_bound(pinldhe ld cross (point a, point b) {
    for (; lo != hi; ++lo)
      ret = min(ret, {(*lo - p).dist2(), {*lo, p}});
    S.insert(p);
  return ret.second;
3.6 Convex Hull
typedef long long ll;
typedef pair <ll, ll> point;
#define x first
#define y second
inline ll area (point a, point b, point c) {
```

```
return (b.x - a.x) * (c.y - a.y) - (b.y - a.y) *
  \rightarrow (c.x - a.x);
vector <point> convexHull (vector <point> p) {
  int n = p.size(), m = 0;
  if (n < 3) return p
  vector <point> hull(n + n);
  sort(p.begin(), p.end());
  for (int i = 0; i < n; ++i)
    while (m > 1 \text{ and area}(hull[m - 2], hull[m - 1],
       p[i]) <= 0) --m;
    hull[m++] = p[i];
  for (int i = n - 2, j = m + 1; i \ge 0; --i) {
    while (m >= j and area(hull[m - 2], hull[m - 1],
       p[i]) <= 0) --m;
    hull[m++] = p[i];
  hull.resize(m - 1); return hull;
3.7 Minimum Enclosing Circle
// Expected runtime: 0(n)
// Solves Gym 102299J
#define x first
```

```
using namespace std;
typedef long double ld:
#define y second
  return point(a.x + b.x, a.y + b.y);
point operator - (const point &a, const point &b) {
 return point(a.x - b.x, a.y - b.y);
point operator * (const point &a, const ld &b) {
 return point(a.x * b, a.y * b);
|point operator / (const point &a, const ld &b) {
 return point(a.x / b, a.y / b);
const ld EPS = 1e-8;
const ld INF = 1e20
const ld PI = acosl(-1);
|inline ld <mark>dist</mark> (point a, point b) {
  return hypotl(a.x - b.x, a.y - b.y);
inline ld sqDist (point a, point b) {
  return (a.x - b.x) * (a.x - b.x) + (a.y - b.y) *
   \rightarrow (a.y - b.y);
inline ld dot (point a, point b) {
  return a.x * b.x + a.y * b.y;
  return a.x * b.y - a.y * b.x;
|inline ld cross (point a, point b, point c) {
  return cross(b - a, c - a);
|inline point perp (point a) {
  return point(-a.y, a.x);
// circle through 3 points
pair <point, ld> getCircle (point a, point b, point c)
  pair <point, ld> ret;
  'ld den' = (ld) 2 * cross(a, b, c);
```

```
ret.x.x = ((c.y - a.y) * (dot(b, b) - dot(a, a))
  ret.x.y = ((b.x - a.x) * (dot(c, c) - dot(a, a)) -
  \rightarrow (c.x - a.x) * (dot(b, b) - dot(a, a))) / den;
  ret.y = dist(ret.x, a);
  return ret;
pair <point, ld> minCircleAux (vector <point> &s,

→ point a, point b, int n) {
  ld'lo = -INF, hi = INF;
  for (int i = 0; i < n; ++i)
    auto si = cross(b - a, s[i] - a);
    if (fabs(si) < EPS) continue;</pre>
    point m = getCircle(a, b, s[i]).x;
    auto cr = cross(b - a, m - a);
    si < 0? hi = min(hi, cr) : lo = max(lo, cr);
  d = 0 < 0 ? 0 : hi < 0 ? hi : 0;
  point c = (a + b) * 0.5 + perp(b - a) * v /

    sqDist(a, b);

  return {c, sqDist(a, c)};
pair <point, ld> minCircle (vector <point> &s, point
\rightarrow a, int n)
  random shuffle(s.begin(), s.begin() + n);
  point \bar{b} = s[0], c = (a + b) * 0.5;
  ld r = sqDist(a, c);
 for (int i = 1; i < n; ++i) {
  if (sqDist(s[i], c) > r * (1 + EPS))
     tie(c, r) = n == s.size() ? minCircle(s, s[i],

→ i) : minCircleAux(s, a, s[i], i);
  return {c, r};
pair <point, ld> minCircle (vector <point> s) {
  assert(!s.empty());
  if (s.size() == 1) return {s[0], 0};
  return minCircle(s, s[0], s.size());
int n; vector <point> p;
int main() {
 cin >> n;
  while (n--) {
   double x, y;
scanf("%lf %lf", &x, &y);
    p.emplace back(x, y);
  pair <point, ld> circ = minCircle(p);
  printf("%0.12f %0.12f %0.12f\n", (double) circ.x.x,
  return 0;
3.8 Point In Polygon
```

```
// Test if a point is inside a convex polygon in O(lg
→ n) time
// Solves SPOJ INOROUT
typedef long long ll;
typedef pair <ll, ll> point;
#define x first
#define y second
struct segment {
  point P1, P2;
  segment () {}
  segment (point P1, point P2) : P1(P1), P2(P2) {}
inline ll ccw (point A, point B, point C) {
  return (B.x - A.x) * (C.y - A.y) - (C.x - A.x) *
  \hookrightarrow (B.y - A.y);
```

if (l2 > r2)

swap(l2, r2);

(c.x - a.x);

return $max(l1, l2) \ll min(r1, r2) + EPS;$

return abs(s) < EPS ? 0 : s > 0 ? +1 : -1;

|bool intersect(const seg& a, const seg& b)

|bool operator<(const seg& a, const seg& b)

```
void rec(int l, int r) {
                                                            if (r - l <= 3) {
inline bool pointOnSegment (segment S, point P) {
                                                               for (int i = l; i < r; ++i) {
  ll x = P.x, y = P.y, x1 = S.P1.x, y1 = S.P1.y, x2 =
                                                                 for (int j = i + 1; j < r; ++j) {
     S.P2.x, y2 = S.P2.y;
                                                                   upd ans(a[i], a[j]);
  ll a = x - x1, b = y - y1, c = x2 - x1, d = y2 - y1,
     dot = a * c + b * d, len = c * c + d * d;
  if (x1 == x2 \text{ and } y1 == y2) return x1 == x and y1 == y2
                                                               sort(a.begin() + l, a.begin() + r, cmp y());
                                                               return;
  if (dot < 0 or dot > len) return 0
  return x1 * len + dot * c == x * len and y1 * len +
                                                             int m = (l + r) >> 1;
     dot * d == v * len:
                                                             int midx = a[m].x;
                                                             rec(l, m);
const int M = 17
                                                             rec(m, r);
const int N = 10010;
                                                             merge(a.begin() + l, a.begin() + m, a.begin() + m,
struct polygon {
                                                             \rightarrow a.begin() + r, t.begin(), cmp y());
  int n; // n > 1
                                                             copy(t.begin(), t.begin() + r - l, a.begin() + l);
  point p[N]; // clockwise order
                                                             int tsz = 0;
  polygon () {}
                                                            for (int i = l; i < r; ++i) {
   if (abs(a[i].x - midx) < mindist)</pre>
  polygon (int _n, point *T) {
    n = n:
                                                                 for (int j = tsz - 1; j >= 0 \&\& a[i].y - t[j].y
    for (int i = 0; i < n; ++i) p[i] = T[i];
                                                                 upd ans(a[i], t[j]);
  bool contains (point P, bool strictlyInside) {
                                                                 t[tsz++] = a[i];
    int lo = 1, hi = n - 1;
    while (lo < hi){
      int mid = lo + hi >> 1
      if (ccw(p[0], P, p[mid]) > 0) lo = mid + 1;
                                                           void solve(int n)
      else hi = mid;
                                                             t.resize(n):
    if (ccw(p[0], P, p[lo]) > 0) lo = 1;
                                                             sort(a.begin(), a.end(), cmp_x());
    if (!strictlyInside and
                                                            mindist = 1E20:
        pointOnSegment(segment(p[0], p[n - 1]), P))
                                                             rec(0, n);
        return 1;
    if (!strictlyInside and
        pointOnSegment(segment(p[lo], p[lo - 1]), P))
                                                          3.10 sweep
                                                          |const double EPS = 1E-9;
    if (lo == 1 or ccw(p[0], P, p[n - 1]) == 0) return
                                                          struct pt {
                                                           double x, v;
    return ccw(p[lo], P, p[lo - 1]) < 0;
                                                          struct seg {
                                                            pt p, q;
                                                            int id;
3.9 near_pair
                                                             double get y(double x) const {
struct pt {
                                                              if (abs(p.x - q.x) < EPS)
  int x, y, id;
                                                                 return p.y;
                                                               return p.y + (q.y - p.y) * (x - p.x) / (q.x - p.x);
struct cmp x {
  bool operator()(const pt & a, const pt & b) const {
    return a.x < b.x | | (a.x == b.x && a.y < b.y);
                                                          |bool intersect1d(double l1, double r1, double l2,
                                                              double r2) {
                                                            if (l1 > r1)
struct cmp_y {
                                                               swap(l1, r1);
```

bool operator()(const pt & a, const pt & b) const {

double dist = sqrt((a.x - b.x)*(a.x - b.x) + (a.y - b.x)

void upd ans(const pt & a, const pt & b) {

return a.y < b.y;

pair<int, int> best pair;

 \rightarrow b.y)*(a.y - b.y));

best pair = {a.id, b.id};

if (dist < mindist) {</pre>

mindist = dist:

};

vector<pt> a;

vector<pt> t;

double mindist;

```
int vec(const pt& a, const pt& b, const pt& c) {
  double s = (b.x - a.x) * (c.y - a.y) - (b.y - a.y) *
  return intersect1d(a.p.x, a.q.x, b.p.x, b.q.x) &&
    intersect1d(a.p.y, a.q.y, b.p.y, b.q.y) && vec(a.p, a.q, b.p) * vec(a.p, a.q, b.q) <= 0 &&
    vec(b.p, b.q, a.p) * vec(b.p, b.q, a.q) <= 0;
```

```
double x = max(min(a.p.x, a.q.x), min(b.p.x, b.q.x));
 return a.get y(x) < b.get y(x) - EPS;
struct event {
 double x;
 int tp, id;
 event() {}
 event(double x, int tp, int id) : x(x), tp(tp),
  → id(id) {}
 bool operator<(const event& e) const {
    if (abs(x - e.x) > EPS)
      return x < e.x;
    return tp > e.tp;
set<seg> s;
vector<set<seg>::iterator> where;
set<seg>::iterator prev(set<seg>::iterator it) {
 return it == s.begin() ? s.end() : --it;
set<seg>::iterator next(set<seg>::iterator it) {
 return ++it:
pair<int, int> solve(const vector<seq>& a) {
 int n = (int)a.size();
 vector<event> e;
 for (int i = 0; i < n; ++i)
   e.push back(event(min(a[i].p.x, a[i].q.x), +1, i));
   e.push back(event(max(a[i].p.x, a[i].q.x), -1, i));
 sort(e.begin(), e.end());
 s.clear();
 where.resize(a.size());
 for (size t i = 0; i < e.size(); ++i) {
   int id = e[i].id;
   if (e[i].tp == +1) {
     set<seg>::iterator nxt = s.lower bound(a[id]),
         prv = prev(nxt):
     if (nxt != s.end() && intersect(*nxt, a[id]))
        return make pair(nxt->id, id);
     if (prv != s.end() && intersect(*prv, a[id]))
        return make pair(prv->id, id);
     where[id] = s.insert(nxt, a[id]);
     set<seg>::iterator nxt = next(where[id]), prv =
         prev(where[id]);
     if (nxt != s.end() && prv != s.end() &&
         intersect(*nxt, *prv))
        return make pair(prv->id, nxt->id);
     s.erase(where[id]);
  return make pair(-1, -1);
```

4 Graph

4.1 2Sat

```
namespace sat{
 const int MAX = 200010:
 bool vis[MAX];
 vector <int> ed[MAX], rev[MAX]
 int n, m, ptr, dfs t[MAX], ord[MAX], par[MAX];
 inline int inv(int x){
   return ((x) \le n ? (x + n) : (x - n));
 void init(int vars){
   n = vars, m = vars << 1;
```

```
,
```

```
for (int i = 1; i \le m; i++){
    ed[i].clear();
rev[i].clear();
inline void add(int a, int b){
  ed[a].push back(b);
  rev[b].push back(a);
inline void OR(int a, int b){
 add(inv(a), b);
add(inv(b), a);
inline void AND(int a, int b){
  add(a, b);
  add(b, a);
void XOR(int a,int b){
  add(inv(b), a);
 add(a, inv(b));
add(inv(a), b);
  add(b, inv(a));
inline void XNOR(int a, int b){
 add(a,b);
add(b,a);
 add(inv(a), inv(b));
add(inv(b), inv(a));
inline void force true(int x){
  add(inv(x), x);
inline void topsort(int s){
 vis[s] = true;
  for(int x : rev[s]) if(!vis[x]) topsort(x);
  dfs t[s] = ++ptr;
inline void dfs(int s, int p){
  par[s] = p;
  vis[s] = true;
  for(int x : ed[s]) if (!vis[x]) dfs(x, p);
void build(){
 CLR(vis);
  ptr \= 0;
  for(int i=m;i>=1;i--)
    if (!vis[i]) topsort(i);
    ord[dfs t[i]] = i;
  ĆLR(vis);
  for (int i = m; i >= 1; i --){
    int x = ord[i];
    if (!vis[x]) dfs(x, x);
bool satisfy(vector <int>& res){
 build();
  CLR(vis):
  for (int i = 1; i \le m; i++){
    int x = ord[i];
    if (par[x] == par[inv(x)]) return false;
    if (!vis[par[x]]){
      vis[par[x]] = true;
      vis[par[inv(x)]] = false;
  res.clear();
  for (int i = 1; i \le n; i++){
    if (vis[par[i]]) res.push back(i);
  return true;
```

```
4.2 Centroid decomp
vector <int> ed[MAX]
bool isCentroid[MAX]
int sub[MAX], cpar[MAX], clevel[MAX];
int dis[20][MAX]:
void calcSubTree(int s,int p) {
  sub[s] = 1;
  for(int x : ed[s]) {
    if(x == p or isCentroid[x]) continue;
    calcSubTree(x,s);
    sub[s] += sub[x];
int getCentroid(int s,int p) {
  for(int x : ed[s]) {
   if(!isCentroid[x] && x!=p && sub[x]>(nn/2)) return

    getCentroid(x,s);

  return s;
void setDis(int s, int from, int p, int lev) {
  dis[from][s] = lev;
  for(int x : ed[s]) {
    if(x == p or isCentroid[x] ) continue;
    setDis(x, from, s, lev+1);
void decompose(int s,int p,int lev) {
  calcSubTree(s,p);
  nn = sub[s];
  int c = getCentroid(s,p);
  setDis(c,lev,p,0);
  isCentroid[c] = true;
  cpar[c] = \bar{p};
  clevel[c] = lev;
  for(int x : ed[c])
    if(!isCentroid[x]) decompose(x,c,lev+1);
int ans[MAX];
|inline void update(int v) {
  int u = v;
  while(u!=-1) {
    ans[u] = min(ans[u], dis[clevel[u]][v]);
    u = cpar[u];
|inline int query(int v) {
  int ret = INF;
  int u = v;
  while(u \stackrel{!}{=} -1) {
    ret = min(ret, dis[clevel[u]][v]+ans[u]);
    u = cpar[u];
  return ret;
int main()
  decompose(1,-1,0);
  for(int i=1; i<=n; i++) ans[i] = INF;
  update(v);
  query(v));
  return 0;
4.3 articulation_point
```

using namespace std;

```
const int N = 1e5 + 10;
vector<int> g[N];
int vis[N], low[N], cut[N], now = 0, n, m;
void dfs(int u, int p) {
 low[u] = vis[u] = ++now; int ch = 0;
 for(int v : g[ū]){
   if(v ^ p)
      if(vis[v]) low[u] = min(low[u], vis[v]);
      else {
        ch++; dfs(v, u);
low[u] = min(low[u], low[v]);
        if(p + 1 \&\& low[v]) >= vis[u]) cut[u] = 1;
        if(low[v] > vis[u]) {
          printf("Bridge %d -- %d\n", u, v);
   if(p == -1 \&\& ch > 1) cut[u] = 1;
void ArticulationPointAndBridge() {
 now = 0;
 for(int i = 0; i < n; i++) {
   if(!vis[i]) dfs(i, -1);
```

4.4 bcc

```
// clear ed[] every test case
// tot -> total number of components
// bcc[i] contains the nodes of the i'th component
// any self loop or multiple edge?
const int MAX = ?
vector <int> ed[MAX];
bool cut[MAX];
int tot, Time, low[MAX], st[MAX];
vector <int> bcc[MAX];
stack <int> S;
void popBCC(int s,int x) {
 cut[s] = 1;
 bcc[++tot].pb(s);
  while(bcc[tot].back() ^ x) {
    bcc[tot].pb(S.top());
    S.pop();
void dfs(int s, int p = -1) {
  S.push(s);
  int ch = 0
  st[s] = low[s] = ++Time;
  for(int x : ed[s]) {
    if(!st[x]) {
      ch++:
      dfs(x,s);
      low[s] = min(low[s], low[x])
      if(p = -1 \text{ and } low[x] > = st[s]) popBCC(s,x);
      else if(p == -1) if(ch > 1) popBCC(s,x);
    else if(p != x) low[s] = min(low[s],st[x]);
  if(p == -1 \&\& ch > 1) cut[s] = 1;
void processBCC(int n) {
  for(int i=1;i<=n;i++) bcc[i].clear();</pre>
  CLR(st); CLR(cut);
  Time = tot = 0:
  for(int i=1; i<=n; i++) {
    if(!st[i]) {
      dfs(i,-1)
      if(!S.empty()) ++tot;
      while(!S.empty()) {
```

```
bcc[tot].push back(S.top());
         S.pop();
4.5 bridge_tree
const int MAXN = ?:
const int MAXE = ?;
struct edges {
  int u,v;
} ara[MAXE];
vector <int> ed[MAXN];
vector <int> isBridge[MAXN];
vector <int> brTree[MAXN];
bool vis[MAXN];
int st[MAXN], low[MAXN], Time = 0;
int cnum;
int comp[MAXN];
void findBridge(int s,int par) {
  int i,x,child = 0,i;
  vis[s] = 1;
 Time++;

st[s] = low[s] = Time;

st[s] = size(
  for(i=0; i<ed[s].size(); i++) {</pre>
    x = ed[s][i];
    if(!vis[x]) {
       child++
       findBridge(x,s);
       low[s] = min(low[s],low[x]);
      if(low[x] > st[s]) {
        isBridge[s][i] =
        j = lower bound(ed[x].begin(),ed[x].end(),s)-e_
          \rightarrow d[x].begin();
        isBridge[x][j] = 1;
    else if(par!=x)
      low[s]' = min(low[s], st[x]);
void dfs(int s) {
 int i,x;
vis[s] = 1;
comp[s] = cnum;
  for(i=0; i<ed[s].size(); i++) {</pre>
    if(!isBridge[s][i]) {
   x = ed[s][i];
      if(!vis[x]) dfs(x);
void processBridge(int n,int m) {
  CLR(vis);
  for(int i=1; i<=n; i++) if(!vis[i]) findBridge(i,-1);</pre>
  cnum = 0;
  CLR(vis);
for(int i=1; i<=n; i++) {
    dfs(i);
  n = cnum;
  for(int i=1; i<=m; i++) {
    if(comp[ara[i].u] != comp[ara[i].v])
      brTree[comp[ara[i].u]].pb(comp[ara[i].v]);
       brTree[comp[ara[i].v]].pb(comp[ara[i].u]);
```

```
|int main() {
  int n,m,u,v;
  scanf("%d %d",&n,&m);
for(int i=1; i<=m; i++) {
    sii(u,v);
    ed[u].pb(v);
    ed[v].pb(u);
    isBridge[u].pb(0);
    isBridge[v].pb(0);
    ara[i].u = u;
ara[i].v = v;
  for(int i=1; i<=n; i++) sort(all(ed[i]));</pre>
  processBridge(n,m);
  return 0;
4.6 dinic
namespace dinic {
  using T = int;
  const int MAXN = 5010;
  int n, src, snk, work[MAXN];
  T dist[MAXN];
  struct Edge{
    int to, rev pos;
    T c, f;
  vector <Edge> ed[MAXN];
  void init(int n, int src, int snk) {
    n = n, src = src, snk = snk;
    for(int i=1;i<=n;i++) ed[i].clear();</pre>
  inline void addEdge(int u, int v, T c, T rc = 0) {
    Edge a = \{v, (int)ed[v].size(), c, 0\};
    Edge b = \{u, (int)ed[u].size(), rc, 0\};
    ed[u].push back(a);
    ed[v].push_back(b);
  bool dinic bfs() {
    SET(dist);
    dist[src] = 0;
    queue <int> q;
    q.push(src);
    while(!q.empty()){
      int u = q.front(); q.pop();
      for(Edge &e : ed[u]){
        if(dist[e.to] == -1 \text{ and } e.f < e.c) 
          dist[e.to] = dist[u] + 1;
           q.push(e.to);
    return (dist[snk]>=0);
  T dinic dfs(int u, T fl){
    if (u == snk) return fl;
for (; work[u] < (int)ed[u].size(); work[u]++) {
   Edge &e = ed[u][work[u]];</pre>
      if (e.c <= e.f) continue;</pre>
      int v = e.to;
      if (dist[v] = dist[u] + 1)
        T df = dinic dfs(v, min(fl, e.c - e.f));
        if (df > 0){
           e.f += df:
           ed[v][e.rev pos].f -= df;
           return df;
```

```
return 0;
 † solve() {
    T ret `= 0;
    while (dinic bfs()) {
      CLR(work);
      while (T delta = dinic dfs(src, INF)) ret +=

    delta;

    ŕeturn ret;
int main() {
 int n, m, u, v, c;
  cin >> n >> m;
  dinic::init(n, 1, n);
  while(m--) {
    cin >> u >> v >> c;
    dinic::addEdge(u, v, c, c);
  cout << dinic::solve() << '\n';</pre>
  return 0;
4.7 dsu_on_tree
void calcSubSize(int s,int p) {
 sub[s] = 1;
  for(int x : G[s]) {
    if(x==p) continue;
    calcSubSize(x,s);
    sub[s] += sub[x];
void add(int s,int p,int v,int bigchild = -1) {
 freq[color[s]] += v;
  for(int x : G[s]) {
    if(x==p || x==bigchild) continue;
    add(x,s,v);
void dfs(int s,int p,bool keep) {
  int bigChild = -1;
  for(int x : G[s]) {
    if(x==p) continue;
    if(bigChild==-1 || sub[bigChild] < sub[x] )</pre>

→ biaChild = x:

  for(int x : G[s]) {
    if(x==p || x==bigChild) continue;
    dfs(x,s,0);
  if(biqChild!=-1) dfs(bigChild,s,1);
  add(s,p,1,bigChild);
 if(keep==0)
    add(s,p,-1);
```

4.8 euler_path

```
vector <int> ed[MAX+5], sltn;
int inDeg[MAX+5], outDeg[MAX+5];
bool vis[MAX+5];
void dfs(int nd) {
  vis[nd] = true;
  while(ed[nd].size()) {
    int v = ed[nd].back();
    ed[nd].pop_back();
    dfs(v);
}
```

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```
sltn.pb(nd);
int findEuler (int n) {
  int src , snk , ret = 1;
bool found_src = false, found_snk = false;
  CLR(inDeg); CLR(outDeg);
  for(int u = 1; u <= n; u++) {
  for(int i = 0; i<ed[u].size(); i++) {</pre>
       int v = ed[u][i];
       outDeg[u]++;
       inDeg[v]++;
  int diff;
  for(int i = 1; i<=n; i++)
    diff = outDeg[i] - inDeg[i];
    if(diff == 1) {
       if(found src) return 0;
       found sr\overline{c} = true;
       src = i;
    else if (diff == -1) {
       if(found snk) return 0;
       found snk = true;
       snk = i;
    else if(diff != 0) return 0;
  if(!found src) {
    ret = 2;
    for(int i = 1; i <= n; i++) {
  if( outDeg[i] ) {</pre>
         found src = true;
         src = i;
         break;
  if(!found src) return ret;
  CLR(vis);
  sltn.clear();
  dfs(src);
  for(int i = 1; i<=n; i++) {
   if(outDeg[i] && !vis[i]) return 0;</pre>
  for(int i = (int)sltn.size()-1; i>=0; i--)
   → `printf("%d ",sltn[i]);
  puts("");
  return ret;
```

4.9 hld

```
const int N = 3e4 + 5;
vector<int> G[N];
int sz[N], H[N], P[N];
void dfs(int cur, int h)
    sz[cur] = 1;
H[cur] = h;
    for(int& to : G[cur])
        G[to].erase(find(G[to].begin(), G[to].end(),

    cur));
P[to] = cur;
        dfs(to, h + 1);
sz[cur] += sz[to];
         if(sz[to] > sz[G[cur][0]]) swap(G[cur][0], to);
int base[N], pos[N], head[N];
int ptr = 0;
void hld(int cur)
```

```
pos[cur] = ++ptr;
    base[ptr] = cur;
    for(int to : G[cur]) {
        head[to] = (to' == G[cur][0] ? head[cur] : to);
        hld(to);
segtree ST;
|int query(int u, int v)
    int ret = 0;
    while(head[u] != head[v])
        if(H[head[u]] > H[head[v]]) swap(u, v);
        ret += ST.query(pos[head[v]], pos[v]);
        v = P[head[v]]:
    if(H[u] > H[v]) swap(u, v);
    ret += ST.query(pos[u], pos[v]);
    return ret;
void update(int u, int val) {
    ST.update(pos[u], val);
void build(int n, int root)
    ptr = 0;
    dfs(root, 0);
    head[root] = root;
    hld(root);
    ST = seqtree(n);
   clear graph
   call build
   Prob : sum of values from u to v
```

4.10 hopcroft_karp

```
struct HopcroftKarp {
  const int N, M;
  std::vector<std::vector<int>> adj left;
  std::vector<int> matchL, matchR;
  HopcroftKarp(int N, int M, const

    std::vector<std::pair<int, int>>& edge)
      : N(N), M(M), matchL(N, -1), matchR(M, -1),
       → adj left(N) {
    for (auto [l, r] : edge)
  adj_left[l].push_back(r);
  int maxmatching() {
    int sz = 0:
    for (bool updated = true; updated;) {
      updated = false;
      static std::vector<int> root(N), prev(N), qq(N);
      static int qi, qj;
      // std::queue<int> q;
      qi = qj = 0;
std::fill(root.begin(), root.end(), -1),
      std::fill(prev.begin(), prev.end(), -1);
      for (int i = 0; i < N; i++)
        if (matchL[i] == -1)
    qq[qj++] = i, root[i] = i, prev[i] = i;
/ g.push(i), root[i] = i;
      while (qi < qj) {
         int u = qq[qi++];
         // int u = q.front(); q.pop();
         if (matchL[root[u]] != -1) continue;
         for `(int v : adj_left[u]) {
           if (\text{matchR}[v] = -1) {
             while (v != -1)
```

```
matchR[v] = u, std::swap(matchL[u], v),
            \rightarrow u = prev[u];
         updated = true, sz++;
         break:
       if (prev[matchR[v]] == -1)
         v = matchR[v], prev[v] = u, root[v] =

    root[u], qq[qj++] = v;
// v = matchR[v], prev[v] = u, root[v] =

        \rightarrow root[u], q.push(v);
  }
return sz;
```

4.11 hungarian

```
// Given NN matrix A[i][j]. Calculate a permutation
    p[i] that minimize A[i][p[i]].
template <typename T>
pair <T, vector <int>>> Hungarian (int n, int m, T
\leftarrow c[N][N]
 vector <T> v(m), dist(m);
  vector < int > L(n, -1), R(m, -1);
  vector <int> index(m), prev(m);
  auto residue = [&] (int i, int j) {return c[i][j] -

    ∨[j];};
  iota(index.begin(), index.end(), 0);
  for (int f = 0; f < n; ++f)
    for (int j = 0; j < m; ++j) {
      dist[j] = residue(f, j), prev[j] = f;
    \hat{T} w; int i, j, l, s = 0, t = 0;
    while (true) {
      if (s == t) {
        l = s, w = dist[index[t++]];
        for (int k = t; k < m; ++k) {
   j = index[k]; T h = dist[j];</pre>
          if (h <= w) {
            if (h < w) t = s, w = h;
index[k] = index[t], index[t++] = j;
        for (int k = s; k < t; ++k) {
          j = index[k];
          if (R[j] < 0) goto augment;
      int q = index[s++], i = R[q];
      for (int k = t; k < m; ++k) {
          = index[k];
        \dot{T} h = residue(i, j) - residue(i, q) + w;
        if (h < dist[j]) -
          dist[j] = h, prev[j] = i;
          if (h == w) {
             if (R[j] < 0) goto augment;
             index[k] = index[t], index[t++] = j;
  augment:
    for (int k = 0; k < l; ++k) v[index[k]] +=
        dist[index[k]] - w;
    do
      R[j] = i = prev[j], swap(j, L[i]);
    } while (i ^ f);
   ret = 0;
 for (int i = 0; i < n; ++i) ret += c[i][L[i]];
```

```
return {ret, L};
4.12 kuhn
namespace bpm{
  const int L = 105;
  const int R = 105;
  vector <int> G[L];
  int matchR[R], matchL[L], vis[L], it;
  void init(int n)
    SET(matchL), SET(matchR), CLR(vis);
    for(int i=1;i<=n;i++) G[i].clear();</pre>
  inline void addEdge(int u,int v) { G[u].pb(v); }
  bool dfs(int s) {
  vis[s] = it;

    for(auto x : G[s]) {
      if( matchR[x] == -1 or (vis[matchR[x]] != it and
         dfs(matchR[x])) )
        matchL[s] = x; matchR[x] = s;
        return true;
    return false;
  int solve() {
    int cnt = 0;
    for(int i=1;i<=n;i++) {</pre>
      if(dfs(i)) cnt++, it++;
    return cnt;
```

```
4.13 lca
// Don't forget to clear ed after test case ends(vt,

→ cost are cleared inside)

using namespace std;
const int MAX = 100010;
int LG;
int dep[MAX], par[MAX][21];
vector <int> ed[MAX];
void dfs(int s, int p, int d) {
  dep[s] = d, par[s][0] = p;
  for(int x : ed[s]) {
    if(x == p) continue;
    dfs(x, s, d+1);
void preprocess(int root, int n) {
  LG = \lg(n);
  memset(par, -1, sizeof(par));
  dfs(root, -1, 0);
  for(int j=1; j<=LG; j++) {</pre>
    for(int i=1;i<=n;i++) {</pre>
      if(par[i][j-1] != -1) par[i][j] =
       → par[par[i][j-1]][j-1];
int getLCA(int u, int v) {
  if(dep[u] < dep[v]) swap(u, v);</pre>
  for(int i=LG;i>=0;i--) {
    if(dep[u] - (1 << i) >= dep[v]) u = par[u][i];
  if(u == v) return u;
  for(int i=LG;i>=0;i--) {
    if (par[u][i] != -1 and par[u][i] - par[v][i]) {
```

```
u = par[u][i], v = par[v][i];
  return par[u][0];
4.14 manhattan MST
using namespace std;
using ll = long long;
struct UnionFind {
    vector<int> UF:
    int cnt;
    UnionFind(int N) : UF(N, -1),
, □ cnt(N) {}
    int find(int v) { return UF[v] < 0 ? v : UF[v] =</pre>
                                  , find(UF[v]); }
    bool ioin(int v. int w) {
        if ((v = find(v)) == (w = find(w))) return
            false:
        if (UF[v] > UF[w]) swap(v, w);
        UF[v] += UF[w];
        UF[w] = v;
        cnt - - :
        return true;
    bool connected(int v, int w) {
        return find(v) == find(w);
    int getSize(int v) { return -UF[find(v)]; }
template <class T>
|struct KruskalMST {
    using Edge = tuple<int, int, T>;
    T mstWeight;
    vector<Edge> mstEdges;
    UnionFind_uf;
    KruskalMST(int V, vector<Edge> edges) :

→ mstWeight(),

  uf(V) {
        sort(edges.begin(), edges.end(), [&](const
            Edge &a, , const Edge &b)
            return get<2>(a) < get<2>(b);
        for (auto &&e : edges) {
            if (int(mstEdges.size()) >= V - 1) break;
            if (uf.join(get<0>(e), get<1>(e))) {
                mstEdges.push back(e);
                mstWeight += \overline{q}et < 2 > (e);
template <class T>
struct ManhattanMST : public, MruskalMST<T> {
    using Edge = typename KruskalMST<T>::Edge;
    static vector<Edge>
        generateCandidates(vector<pair<T, T>>, P) {
        vector<int> id(P.size())
        iota(id.begin(), id.end(), 0);
        vector<Edge> ret;
        ret.reserve(P.size() * 4);
        for (int h = 0; h < 4; h++)
            sort(id.begin(), id.end(), [&](int i, int
                 return P[i].first - P[j].first <
                 → P[j].second - P[i].second;
            map<T, int> M;
            for (int i : id) {
                auto it = M.lower bound(-P[i].second);
                for (; it != M.en\overline{d}(); it =
```

```
int j = it->second;
                     T dx = P[i].first - P[j].first, dy
                     if (dy > dx) break;
                     ret.emplace back(i, j, dx + dy);
                M[-P[i].second] = i;
            for (auto &&p : P) {
                if (h % 2)
                     p.first = -p.first;
                else'
                     swap(p.first, p.second);
        return ret:
    ManhattanMST(const vector<pair<T, T>> &P)
        : KruskalMST<T>(P.size'()

    generateCandidates(P)) {}
int main() {
    int N;
    cin >> N;
    vector<pair<ll, ll>> P(N);
for (auto &&p : P) cin >> p.first >> p.second;
    ManhattanMST mst(P)
    cout << mst.mstWeight << '\n';</pre>
    for (auto &&[v, w, weight] : mst.mstEdges) cout <<

    v << □ ,. □ << w << '\n'; return 0;
</pre>
```

4.15 mcmf

```
namespace mcmf {
 using T = int;

const T INF = ?; // 0x3f3f3f3f or

→ 0x3f3f3f3f3f3f3f3f3f3fLL
 const int MAX = ?; // maximum number of nodes
 int n, src, snk;
T dis[MAX], mCap[MAX];
 int par[MAX], pos[MAX];
 bool vis[MAX];
 struct Edge{
    int to, rev_pos;
   T cap, cost, flow;
 vector <Edge> ed[MAX];
 void init(int n, int src, int snk) {
   n = n, src = src, snk = snk;
    for(int i=1;i<=n;i++) ed[i].clear();
 void addEdge(int u, int v, T cap, T cost) {
    Edge a = \{v, (int)ed[v].size(), cap, cost, 0\};
    Edge b = \{u, (int)ed[u].size(), 0, -cost, 0\};
    ed[u].pb(a);
    ed[v].pb(b);
 inline bool SPFA(){
    CLR(vis);
    for(int i=1; i<=n; i++) mCap[i] = dis[i] = INF;</pre>
    queue <int> q;
   dis[src] = 0;
vis[src] = true;
    q.push(src);
    while(!q.empty()){
      int u = q.front();
      q.pop();
      vis[u] = false;
      for(int i=0; i<(int)ed[u].size(); i++) {</pre>
```

```
Edge &e = ed[u][i];
      int v = e.to;
      if (e.cap > e.flow \&\& dis[v] > dis[u] + e.cost) { | 5.2 | FloorSum |}
        dis[v] = dis[u] + e.cost;
        par[v] = u;
        pos[v] = i;
        mCap[v] = min(mCap[u], e.cap - e.flow);
        if(!vis[v]) {
          vis[v] = true;
          q.push(v);
   }
 return (dis[snk] != INF);
inline pair <T, T> solve() {
 T F = 0, C = 0, f;
 int u. v:
 while(SPFA()){
   u = snk;
   f = mCap[u];
    F += f;
    while(u!=src){
      v = par[u];
      ed[v][pos[u]].flow += f; // edge of v-->u
      ed[u][ed[v][pos[u]].rev pos].flow -= f;
    \dot{C} += dis[snk] * f;
 return make pair(F,C);
```

5 Math 5.1 FWHT

```
const int N = 1 << 20:
// apply modulo if necessary
void fwht xor(int *a, int n, int dir = 0) {
    for (int h = 1; h < n; h <<= 1) {
        for (int i = 0; i < n; i += h << 1)
             for (int j = i; j < i + h; ++j) {
                 int x = a[j], y = a[j + h];
                 a[j] = x + y, a[j + h] = x - y;
if (dir) a[j] >>= 1, a[j + h] >>= 1;
void fwht or(int *a, int n, int dir = 0) {
    for (int h = 1; h < n; h <<= 1) {
        for (int i = 0; i < n; i += h << 1)
             for (int j = i; j < i + h; ++j) {
                 int x = a[j], y = a[j + h];

a[j] = x, a[j + h] = dir ? y - x : x +
                  }
void fwht and(int *a, int n, int dir = 0) {
    for (int h = 1; h < n; h <<= 1) {
        for (int i = 0; i < n; i += h << 1)
             for (int j = i; j < i + h; ++j) {
                 int x = a[j], y = a[j + h];
                 a[j] = dir ? x - y : x + y, a[j + h] =

→ y;

    }
```

```
// O(log a) sum^n floor(ax + b / c) = f
long long FloorSumAP(long long a, long long b, long
    long c, long long n){
   if(!a) return (b / c) * (n + 1);
  if(a >= c or b >= c) return ( (n * (n + 1) ) / 2) *
       (a / c) + (n + 1) * (b / c) + FloorSumAP(a % c,
  \overrightarrow{b} % c, c, n);
long long m = (a * n + b) / c;
   return m * n - FloorSumAP(c, c - b - 1, a, m - 1);
 1/\sqrt{0(\log a)} sum^n x * floor(ax + b / c) = q, sum^n
 \rightarrow floro(ax + b / c)^2 = h
struct dat {
  long long f, g, h;
  dat(long long f = 0, long long q = 0, long long h =
   \rightarrow 0) : f(f), q(q), h(h) {};
long long mul(long long a, long long b){
  return (a * b) % MOD;
dat query(long long a, long long b, long long c, long
  if(!a) return {mul(n + 1, b / c), mul(mul(mul(b / c,
   \rightarrow n), n + 1), inv2), mul(mul(n + 1, b / c), b /c)};
   long long f, g, h;
   dat nxt;
  if(a >= c or b >= c){
  nxt = query(a % c, b % c, c, n);
     f = (nxt.f + mul(mul(mul(n, n + 1), inv2), a / c)
     \rightarrow + mul(n + 1, b / c)) % MOD;

g = (nxt.g + mul(a / c, mul(mul(n, n + 1), mul(2 *
         n + 1, inv6))) + mul(mul(b / c, mul(n, n +
     h = 1), inv2) \% MOD;

h = (nxt.h + 2 * mul(b / c, nxt.f) + 2 * mul(a / c)
          c, nxt.g) + mul(mul(a / c, a / c), mul(mul(n,
         n + 1), mul(2 * n + 1, inv6))) + <math>mul(mul(b / a))
      \stackrel{\sim}{\sim} c, b / c), n + 1) + mul(mul(a / c, b / c),
      \stackrel{\rightharpoonup}{=} mul(n, n + 1)) ) % MOD;
     return {f, q, h};
   long long m = (a * n + b ) / c;
   nxt = query(c, c - b - 1, a, m - 1);
   f = (mul(m, n) - nxt.f) % MOD;
   g = mul(mul(m, mul(n, n + 1)) - nxt.h - nxt.f,
  h = (mul(n, mul(m, m + 1)) - 2 * nxt.q - 2 * nxt.f -

→ f) % MOD:

  return {f, g, h};
5.3 NOD
N = input()
primes = array containing primes till 10^6
for all p in primes :
     if p*p*p > N:
         break
     count =
     while N divisible by p:
         N = N/p
         count = count + 1
     ans = ans * count
if N is prime:
     ans = ans * 2
|else if N is square of a prime:
     ans = ans * 3
else if N != 1
```

ans = ans $^{-}$ * 4

5.4 Pollard Rho

```
we#include <bits/stdc++.h>
using namespace std;
typedef long long ll;
typedef unsigned long long ull;
namespace Rho {
 ull mul (ull a, ull b, ull mod) {
    ll ret = a * b - mod * (ull) (1.L / mod * a * b);
    return ret + mod * (ret < 0) - mod * (ret >= (ll)
  ull bigMod (ull a, ull e, ull mod) {
    ull ret = 1;
    while (e) -
      if (e \& 1) ret = mul(ret, a, mod);
      a = mul(a, a, mod), e >>= 1;
    return ret;
  bool isPrime (ull n) {
    if (n < 2 \text{ or } n \% 6 \% 4 != 1) \text{ return } (n | 1) == 3;
    ull a[] = \{2, 325, 9375, 28178, 450775, 9780504,

→ 1795265022}

    ull s = builtin ctzll(n - 1), d = n >> s;
    for (ull \overline{x}: a) {
      ull p = bigMod(x % n, d, n), i = s;
      while (p != 1 \text{ and } p != n - 1 \text{ and } x % n \text{ and } i--)
       \rightarrow p = mul(p, p, n);
      if (p != n - 1 and i != s) return 0;
    return 1;
  ull pollard (ull n) {
    auto f = [\&] (ull x) {return mul(x, x, n) + 1;};
    ull x = 0, y = 0, t = 0, prod = 2, i = 1, q;
    while (t++ \% 40 \text{ or } \gcd(\text{prod}, n) == 1) {
      if (x == y) x = ++i, y = f(x);
      if ((q = mul(prod, max(x, y) - min(x, y), n)))
       \rightarrow prod = q;
      x = f(x), y = f(f(y));
    return gcd(prod, n);
  vector <ull> factor (ull n) {
    if (n == 1) return {};
    if (isPrime(n)) return {n};
    ull x = pollard(n);
    auto l = factor(x), r = factor(n / x);
l.insert(l.end(), r.begin(), r.end());
    return l;
int t; ll n;
int main() {
 cin >> t;
  while (t--)
    scanf("%lld", &n);
vector <ull> facs = Rho::factor(n);
sort(facs.begin(), facs.end());
    printf("%d", (int) facs.size());
    for (auto it : facs) printf(" %llu", it);
    puts("");
  return 0;
```

```
14
```

```
5.5 catalan
//Recursive
const int MOD = ....
const int MAX =
int catalan[MAX];
void init()
void init() +
      catalan[0] = catalan[1] = 1;
     for (int i=2; i<=n; i++) {
    catalan[i] = 0;
    for (int j=0; j < i; j++) {
        catalan[i] += (catalan[j] *</pre>

    catalan[i-j-1]) % MOD;

                 if (catalan[i] >= MOD) {
   catalan[i] -= MOD;
//Analytical formula:
ans = ncr(2*n,n) - ncr(2*n,n-1) = ncr(2*n,n)/(n+1)
```

```
//r[i][j]= inverse of p[i] modulo p[j] //ans= x[0]+x[1]*p[0]+x[2]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*(p[0]*p[1])+...+x[k-1]*
\rightarrow ]*p[1]*p[2]*...*p[k-2])
//ans %= ((p[0]*p[1]*p[2]*...*p[k-1])
  for (int i = 0; i < k; ++i) {
                              \hat{x}[i] = a[i];
                              for (int j = 0; j < i; ++j) {
    x[i] = r[j][i] * (x[i] - x[j]);
                                                             x[i] = x[i] % p[i];
                                                             if (x[i] < 0)
x[i] += p[i];
  ll mul= p[0], res=x[0], tot=1;
  F(i,0,k) tot *= p[i];
  F(i,1,k){
       res+= x[i]*mul;
res %= tot;
        mul *= p[i];
   res %= mul;
  return res;
```

5.7 derangement

```
int derangement(int n)
if(!n) return n;
if(n \le 2) return n-1;
return (n-1)*(derangement(n-1) + derangement(n-2));
```

5.8 diophantine

```
void print solution(int a, int b, int c)
   int x, y;
   if (a == 0 \& \& b == 0)  {
       if (c == 0)
            cout << "Infinite Solutions Exist" << endl;
       else {
            cout << "No Solution exists" << endl:
   int qcd = qcd extend(a, b, x, y);
   if (c % qcd != 0) {
        cout<< "No Solution exists"<< endl;</pre>
   else {
```

```
cout << "x = " << x * (c / gcd)<< ", y = " <<
          \rightarrow y * (c / gcd)<< endl;
5.9 discrete_log
// Returns minimum x for which a ^ x % m = b % m, a

→ and m are coprime.

|int solve(int a, int b, int m) {
    a \% = m, b \% = m;
    int n = sqrt(m) + 1;
    int an = 1;
for (int i = 0; i < n; ++i)
    an = (an * 1ll * a) % m;</pre>
    unordered map<int, int> vals;
    for (int \overline{q} = 0, cur = b; q \le n; ++q) {
         vals[cur] = q;
         cur = (cur * 111 * a) % m;
    for (int p = 1, cur = 1; p <= n; ++p) {
         cur = (cur * 1ll * an) % m;
         if (vals.count(cur)) {
             int ans = n * p - vals[cur];
             return ans;
    return -1;
5.10 factorial_mod_p
// O(log p(n)) gives me n! % p for large n, p
int factmod(int n, int p) {
    vector<int> f(p);
    f[0] = 1;
    for (int i = 1; i < p; i++)
         f[i] = f[i-1] * i % p;
    int res = 1;
    while (n > 1) {
         if ((n/p) % 2)
res = p - res;
         res = res * f[n%p] % p;
         n /= p;
```

5.11 fft

return res;

```
typedef complex<double> base;
#define PI acos(-1)
void fft(vector<base> &a, bool invert){
    int n = (int)a.size();
    for (int i = 1, j = 0; i < n; ++i){
         int bit = n >> 1;
         for (; j >= bit; bit >>= 1) j -= bit;
         j += bit;
        if (i < j)swap(a[i], a[j]);
    for (int len = 2; len <= n; len <<= 1){
    double ang = 2 * PI / len * (invert ? -1 : 1);</pre>
         base wlen(cos(ang), sin(ang));
        for (int i = 0; i < n; i + len){
             base w(1);
             for (int j = 0; j < len / 2; ++j){
                  base u = a[i + j], v = a[i + j + len /
                 a[i + j] = u + v;
                 a[i + j + len / 2] = u - v;
                 w *= wlen:
```

```
if (invert) for (int i = 0; i < n; ++i) a[i] /= n;</pre>
void multiply(const vector<int> &a, const vector<int>

→ &b, vector<int> &res){
    vector<base> fa(a.begin(), a.end()), fb(b.begin(),
      → b.end());
     size t n = 1;
     while (n < max(a.size(), b.size())) n <<= 1;
     n <<= 1;
    fa.resize(n), fb.resize(n);
fft(fa, false), fft(fb, false);
     for (size t i = 0; i < n; ++i) fa[i] *= fb[i];</pre>
     fft(fa, true); res.resize(n);
for (size_t i = 0; i < n; ++i) res[i] =</pre>

    int(fa[i].real() + 0.5);
```

```
5.12 gauss_eliminition
const double EPS = 1e-9;
const int INF = 2; // it doesn't actually have to be

→ infinity or a big number
int gauss (vector < vector<double> > a, vector<double>
    & ans) {
    int n = (int) a.size();
    int m = (int) a[0].size() - 1;
    vector<int> where (m, -1);
     for (int col=0, row=0; col<m \&\& row<n; ++col) {
         int sel = row;
         for (int i=row: i<n: ++i)
             if (abs (a[i][col]) > abs (a[sel][col]))
                  sel = i
         if (abs (a[sel][col]) < EPS)
             continue;
         for (int i=col; i<=m; ++i)
    swap (a[sel][i], a[row][i]);</pre>
         where[col] = row;
         for (int i=0; i<n; ++i)
             if (i != row) {
                  double c = a[i][col] / a[row][col];
for (int j=col; j<=m; ++j)</pre>
                      a[i][j] -= a[row][j]^* c;
         ++row:
    ans.assign (m, 0);
    for (int i=0; i<m; ++i)
         if (where[i] != -1)
             ans[i] = a[where[i]][m] / a[where[i]][i];
    for (int i=0; i<n; ++i) {
         double sum = 0;
         for (int j=0; j<m; ++j)
    sum += ans[j] * a[i][j];</pre>
         if (abs (sum - a[i][m]) > EPS)
             return 0;
     for (int i=0; i<m; ++i)</pre>
         if (where[i] == -1)
             return INF;
//modular
int gauss (vector < bitset<N> > a, int n, int m,
    bitset<N> & ans) {
    vector<int> where (m, -1);
    for (int col=0, row=0; col<m && row<n; ++col) {
         for (int i=row; i<n; ++i)</pre>
              if (a[i][col]) {
                  swap (a[i], a[row]);
                  break;
```

```
if (! a[row][col])
            continue;
        where[col] = row;
        for (int i=0; i<n; ++i)
            if (i != row && a[i][col])
                `a[i] ^= a[row];
   }
        // The rest of implementation is the same as
//rank
const double EPS = 1E-9;
int compute rank(vector<vector<double>> A) {
    int n = A.size();
    int m = A[0].size();
    int rank = 0;
    vector<bool> row selected(n, false);
    for (int i = 0; \bar{i} < m; ++i) {
        int j;
        for (j = 0; j < n; ++j) {
            if (!row selected[j] && abs(A[j][i]) > EPS)
        if (j != n) {
            ++rank;
            row selected[j] = true;
            for (int p = i + 1; p < m; ++p)
                A[j][p] /= A[j][i];
            for (int k = 0; k < n; ++k)
                if (k != j \&\& abs(A[k][i]) > EPS) {
                    for (int p = i + 1; p < m; ++p
                        A[k][p] -= A[j][p] * A[k][i];
            }
        }
    return rank;
```

5.13 gen_all_k_combs

```
vector<int> ans;
void gen(int n, int k, int idx, bool rev) {
   if (k > n || k < 0)</pre>
        return;
    if (!n)
        for (int i = 0; i < idx; ++i) {
            if (ans[i])
                 cout \ll i + 1;
        cout << "\n";
        return;
    ans[idx] = rev;
    gen(n-1, k-rev, idx + 1, false);
    ans[idx] = !rev;
    gen(n-1, k-!rev, idx + 1, true);
void all combinations(int n, int k) {
    ans.resize(n);
    gen(n, k, 0, false);
```

5.14 integrate_adaptive

```
/*Description: Fast integration using an adaptive

→ Simpsons rule.
Usage: double sphereVolume = quad(-1, 1, [](double x)
return quad(-1, 1, [&](double y) {
return quad(-1, 1, [&](double z) {
return x*x + y*y + z*z < 1; });});*/
```

```
typedef double d;
#define S(a, b) (f(a) + 4 * f((a + b) / 2) + f(b)) *
\rightarrow (b - a) / 6
template <class F>
d rec(F\& f, d a, d b, d eps, d S) {
    d c = (a + b) / 2;

d S1 = S(a, c), S2 = S(c, b), T = S1 + S2;

if (abs(T - S) <= 15 * eps || b - a < 1e-10)

return T + (T - S) / 15;

return rec(f, a, c, eps / 2, S1) + rec(f, c, b,

→ eps / 2. S2):

template <class F>
d quad(d a, d b, F f, d eps = 1e-8)
     return rec(f, a, b, eps, S(a, b));
5.15 matrix_expo
ll\ mod=(1e9)+7;
|struct Matrix{
     int row, col;
     vector<vector<ll>>> mat;
     Matrix(int x, int y){
          row=x;
         col=y;
         mat.assign(row, vector<ll>(col,0));
    Matrix operator *(Matrix &other){
         assert(col==other.row);
         Matrix product(row,other.col);
         for(int i=0;i<row;i++){</pre>
              for(int j=0;j<col;j++){
                   for(int k=0; k<other.col; k++){</pre>
                        product.mat[i][k]=(product.mat[i][,
                             k]+(mat[i][j]*other.mat[j][k])
                         ⇒ %mod)%mod:
              }
         return product;
|Matrix expo(Matrix &m, ll n){
     assert(m.row==m.col);
     Matrix ret(m.row,m.col)
     for(int i=0;i<m.row;i++) ret.mat[i][i]=1;</pre>
     while(n){
         if(n\&1) ret=ret*m;
         n/=2;
m=m*m;
     return ret;
5.16 next_lexicographical_k_comb
     int k = (int)a.size();
```

const int root = 15311432;

```
const int k = 1 \ll 23;
int root 1;
vector<int> rev;
void pre(int sz){
     root 1 = bigmod(root, mod - 2, mod);
     if (rev.size() == sz) return;
     rev.resize(sz);
     rev[0] = 0;
     int lg_n = __builtin_ctz(sz);
     for (int i = 1; i < \overline{sz}; ++i)
     \rightarrow rev[i]=(rev[i>>1]>>1)|((i&1)<<(lq n-1));
|void fft(vector<int> &a, bool inv){
    int n = a.size();
    for (int i = 1; i < n - 1; ++i) if (i < rev[i])

    swap(a[i], a[rev[i]]);

    for (int len = 2; len <= n; len <<= 1) {
   int wlen = inv ? root_1 : root;</pre>
         for (int i = len; i < k; i <<= 1) wlen = 1ll *

→ wlen * wlen % mod;

         for (int st = 0; st < n; st += len){</pre>
              int w = 1;
              for (int j = 0; j < len / 2; j++){
                  int ev = a[st + j];
                  int od = 1ll * a[st + j + len / 2] * w

→ % mod;

                  a[st + j]' = ev + od < mod ? ev + od :
                   \rightarrow ev + od - mod;
                  a[st + j + len / 2] = ev - od >= 0?
                   \rightarrow ev - od : ev - od + mod;
                  w = 111 * w * wlen % mod;
         }
    if (inv){
         int n = bigmod(n, mod - 2, mod);
         for (int \&x : a) x = 111 * x * n 1 % mod;
|vector<int> mul(vector<int> &a, vector<int> &b){
    int n = a.size(), m = b.size(), sz = 1;
    while (sz < n + m - 1) sz <<= 1;
    vector<int> x(sz), y(sz), z(sz);
    for (int i = 0; i < sz; ++i){
    x[i] = i < n ? a[i] : 0;
    y[i] = i < m ? b[i] : 0;</pre>
     pre(sz);fft(x, 0);fft(y, 0);
     for (int i = 0; i < sz; ++i) z[i] = 111 * x[i] *
     \rightarrow y[i] % mod;
    fft(z, 1); z.resize(n + m - 1);
     return z:
```

5.18 seg_sieve

return modint<MOD>(c < 0 ? c + MOD : c);</pre>

this->value += other.value; if (this->value >= MOD)

inline modint<MOD> &operator+=(modint<MOD> other) {

```
this->value -= MOD;
                                                                                                                                  if (spf[i] == 0) phi[i] = i - 1, spf[i] = i,
             f = (a - (a \% p) + p);
                                                                     return *this;
                                                                                                                                  \rightarrow prime[sz++] = i;
        f = max(p * p, f);
                                                                                                                                  for (int j = 0; j < sz && i * prime[j] < T &&
        for (unsigned j = f; j \le b; j += p)
                                                                 inline modint<MOD> &operator-=(modint<MOD> other) {
                                                                                                                                      prime[j] <= spf[i]; j++)</pre>
             notPrime[j - a] = true;
                                                                                                                                      spf[i * prime[j]] = prime[j];
if (i % prime[j] == 0)
                                                                     this->value -= other.value;
                                                                     if (this->value < 0)</pre>
    if (a == 1)
                                                                         this->value += MOD:
                                                                                                                                           phi[i * prime[j]] = phi[i] * prime[j];
        notPrime[0] = 1:
                                                                     return *this;
                                                                                                                                           phi[i * prime[j]] = phi[i] * (prime[j]
                                                                 inline modint<MOD> &operator*=(modint<MOD> other) {

→ - 1);
5.19 stirling
                                                                     this->value = (int64 t)this->value *
                                                                                                                                  }
that no box is empty.

→ other.value

                                                                                                                              dp[0] = 0;
int stirling2(int n, int k)
                                                                     if (this->value < 0) this->value += MOD;
                                                                                                                              for (int i = 1; i < T; i++) dp[i] = dp[i - 1] +
                                                                     return *this:
                                                                                                                              → phi[i] % mod;
 if(n < k)
                                                                                                                             inv = 1; // q(1)
 return 0;
                                                                 inline modint<MOD> operator-() const { return
 if(k == 1)
 return 1
                                                                     modint<MOD>(this->value ? MOD - this->value :
                                                                                                                         mint p c(long long n) {
 if(dp[n][k] == dp[n][k])
                                                                                                                             if^{-}(n \% 2 == 0) return n / 2 \% mod * ((n + 1) %
                                                                     0); }
                                                                 modint<MOD> pow(uint64_t k) const {
 return dp[n][k];

→ mod) % mod;

 return dp[n][k] = stirling2(n-1,k-1) +
                                                                                                                              return (n + 1) / 2 % mod * (n % mod) % mod;
                                                                     modint<MOD> x =

    stirling2(n-1,k)*k;

                                                                                  v = 1:
                                                                                                                         mint p q(long long n) {
                                                                     for (; k; k \Rightarrow = \bar{1}) {
                                                                                                                              return n % mod;
is empty
                                                                         if (k & 1) y *= x;
int stirling1(int n, int k)
                                                                                                                         mint solve(long long x)
                                                                         x *= x;
                                                                                                                             if (x < T) return dp[x];</pre>
 dp[n][k] = stirling1(n-1,k-1) + stirling(n-1,k)*n-1;
                                                                     return y;
                                                                                                                              if (mp.find(x) != mp.end()) return mp[x];
                                                                                                                              mint ans = 0:
                                                                                                                             for (long long i = 2, last; i <= x; i = last + 1) {
    last = x / (x / i);
    ans += solve(x / i) * (p_g(last) - p_g(i - 1));</pre>
                                                                 modint<MOD> inv() const { return pow(MOD - 2); }
5.20 stirling_number_of_the_second_kind
                                                                 → // MOD must be a prime
   1 / k! * sum (-1)^i nCr(k, i) * (k - i) ^ n
                                                                 inline modint<MOD> operator/(modint<MOD> other)
ll f(int n, int k) {
                                                                     const { return *this * other.inv(); }
                                                                                                                             ans = p c(x) - ans;
    ll res = 0;
                                                                 inline modint<MOD> operator/=(modint<MOD> other) {
                                                                                                                             ans /= inv:
    for (int i = 0; i < k; ++i) {

→ return *this *= other.inv(); }

                                                                                                                             return mp[x] = ans;
        if (i \& 1) res = (res - nCr(k, i) * bp(k - i,
                                                                 inline bool operator==(modint<MOD> other) const {
         \rightarrow n, mod) % mod + mod) % mod;

→ return value == other.value; }

                                                                                                                            // namespace Dirichlet
        else res = (res + nCr(k, i) * bp(k - i, n,
                                                                 inline bool operator!=(modint<MOD> other) const {

→ mod) % mod) % mod;

                                                                    return value != other.value; }
                                                                 inline bool operator<(modint<MOD> other) const {
    if (res < 0) res += mod:
                                                                    return value < other.value; }</pre>
    return res * ifac[k] % mod;
                                                                                                                         5.22 totient
                                                                 inline bool operator>(modint<MOD> other) const {
                                                                  → return value > other.value; }
                                                                                                                         int phi(int n) {
5.21 sum of totient
                                                                                                                             int result = n:
                                                             template <int32 t MOD>
                                                                                                                              for (int i = 2; i * i <= n; i++) {
using namespace __gnu_pbds;
const int N = 3e5 + 9, mod = 998244353;
                                                             modint<MOD> operator*(int64 t value, modint<MOD> n) {
                                                                                                                                  if (n % i == 0) {
                                                             → return modint<MOD>(value) * n: }
                                                                                                                                      while (n % i == 0)
template <const int32 t MOD>
                                                             template <int32 t MOD>
                                                                                                                                           n /= i;
struct modint {
                                                                                                                                      result -= result / i;
                                                             modint<MOD> operator*(int32 t value, modint<MOD> n) {
    int32 t value;
                                                             → return modint<MOD>(value % MOD) * n; }
    modin\overline{t}() = default;
                                                             template <int32 t MOD>
    modint(int32 t value ) : value(value ) {}
                                                                                                                              if (n > 1)
                                                            istream &operator>>(istream &in, modint<MOD> &n) {
    inline modint<MOD> operator+(modint<MOD> other)
                                                                                                                                  result -= result / n;

→ return in >> n.value: }

                                                                                                                              return result;
     template <int32 t MOD>
        int32 t c = this->value + other.value;
                                                            ostream &operator<<(ostream &out, modint<MOD> n) {
        return modint<MOD>(c >= MOD ? c - MOD : c);
                                                                                                                         |void phi 1 to n(int n) {
                                                             → return out << n.value; }</pre>
                                                                                                                             vector<int> phi(n + 1);
    inline modint<MOD> operator-(modint<MOD> other)
                                                            using mint = modint<mod>;
                                                                                                                             phi[0] = 0;
phi[1] = 1;
                                                            namespace Dirichlet {
                                                            // solution for f(x) = phi(x)
        int32 t c = this->value - other.value;
                                                                                                                              for (int i = 2; i <= n; i++)
phi[i] = i;
        return modint<MOD>(c < 0 ? c + MOD : c);
                                                            const int T = 1e7 + 9;
                                                             long long phi[T];
                                                                                                                              for (int i = 2; i \le n; i++) {
                                                            gp hash table<long long, mint> mp;
    inline modint<MOD> operator*(modint<MOD> other)
                                                                                                                                  if (phi[i] == i) {
                                                            |\underline{m}\underline{i}\underline{n}t dp[T], inv;
     for (int j = i; j <= n; j += i)
        int32 t c = (int64_t)this->value * other.value
                                                            int sz, spf[T], prime[T];
                                                                                                                                           phi[j] -= phi[j] / i;
                                                            void init() {
```

memset(spf, 0, sizeof spf);

for (int i = 2; i < T; i++) {

phi[1] = 1;sz = 0:

}

16

5.23 xor basis int basis[d]; // basis[i] keeps the mask of the vector → whose f value is i void insertVector(int mask) { for (int i = 0; i < d; i++) { ((mask & 1 << i) == 0) continue; (!basis[i]) { // If there is no basis vector with the i'th bit set, then insert this vector into the basis basis[i] = mask; return; mask ^= basis[i]; // Otherwise subtract the basis vector from this vector

6 Misc 6.1 interval_container

```
/* Description: Add and remove intervals from a set of
    disjoint intervals. Will merge the added interval
   with any overlapping intervals in
the set when adding. Intervals are [inclusive,
    exclusive).
Time: O(\log N)
set<pii>::iterator addInterval(set<pii>& is, int L,
    int R) {
    if (L == R) return is.end();
auto it = is.lower_bound({L, R}), before = it;
    while (it != is.end() && it->first <= R) {
        R = max(R, it->second);
before = it = is.erase(it);
    if (it != is.begin() && (--it)->second >= L) {
        `L = min(L, it->first);
R = max(R, it->second); is.erase(it);
    return is.insert(before, {L, R});
void removeInterval(set<pii>& is, int L, int R) {
    if (L == R) return;
    auto it = addInterval(is, L, R):
    auto r2 = it->second;
    if (it->first == L)
        is.erase(it);
         (int\&)it->second = L:
    if (R != r2) is.emplace(R, r2);
```

6.2 interval_cover

```
template <class T>
vi cover(pair<T, T> G, vector<pair<T, T>> I) {
    vi S(sz(I)), R;
iota(all(S), 0);
sort(all(S), [&](int a, int b) { return I[a] <
    int at = 0;
    while (cur < G.second) { // (A)
         pair<T, int> mx = make pair(cur, -1);
         while (at < sz(I) \&\& I[S[at]].first <= cur) { 
 <math>mx = max(mx, make\_pair(I[S[at]].second, 

    S[at]));
at++;
         if (mx.second == -1) return {};
         cur = mx.first;
         R.push back(mx.second);
```

```
return R;
```

6.3 pragma

```
// Praamas
#pragma comment(linker, "/stack:200000000")
#pragma GCC optimize("03,unroll-loops")
#pragma GCC target("avx,avx2,fma")
```

6.4 random

```
// shuffle(v.begin(), v.end(),

→ default random engine(rnd(1, 1000)));
mt19937 rng(chrono::steady clock::now().time since epo
   ch().count());
rnd(ll l, ll r)
    return uniform int distribution<ll>(l, r) (rng);
```

6.5 vimre

```
imap jk <Esc>
|set nu
|set mouse=a
set autoindent
set tabstop=4
set shiftwidth=4
    smartindent
    relativenumber
set
    laststatus=2
set hlsearch
let mapleader = " "
nnoremap <leader>s :w<Enter>
|nnoremap <leader>y ggVG"+y<CR>
syntax on
vnoremap <leader>/ :s!^!//!<CR> :noh <CR>
vnoremap <leader>u :s!^//!!<CR>
nnoremap <leader>/ :s!^!//!<CR> :noh <CR>
nnoremap <leader>u :s!^//!!<CR>
```

7 Notes

7.1 Counting

1. Lucas Theorem

For non-negative integers m and n and a prime p, the following congruence relation holds: :

$$\begin{pmatrix} m \\ n \end{pmatrix} \equiv \prod_{i=0}^{k} \begin{pmatrix} m_i \\ n_i^i \end{pmatrix} \pmod{p},$$

where:

$$m = m_k p^k + m_{k-1} p^{k-1} + \dots + m_1 p + m_0,$$

and:

$$n = n_k p^k + n_{k-1} p^{k-1} + \dots + n_1 p + n_0$$

are the base p expansions of m and n respectively. This uses the convention that $\begin{pmatrix} m \\ n \end{pmatrix} = 0$ if $m \le n$.

2. Stirling Numbers of the first kind

S(n,k) counts the number of permutations of n elements with k disjoint cycles.

$$S(n,k) = (n-1) \cdot S(n-1,k) + S(n-1,k-1)$$
 (2)

where, S(0,0) = 1, S(n,0) = S(0,n) = 0

$$\sum_{k=0}^{n} S(n,k) = n! \tag{3}$$

3. Stirling Numbers of the second kind

 $S(n,k) \cdot k! = \text{number of ways to color } n \text{ nodes using colors from } 1 \text{ to } k$ such that each color is used at least once.

An r-associated Stirling number of the second kind is the number of ways to partition a set of n objects into k subsets, with each subset containing at least r elements. It is denoted by $S_r(n,k)$ and obeys the recurrence relation.

$$S_r(n+1,k) = kS_r(n,k) + \binom{n}{r-1}S_r(n-r+1,k-1)$$
 (4)

4. Bell Numbers

Counts the number of partitions of a set.

$$B_{n+1} = \sum_{k=0}^{n} \left(\frac{n}{k}\right) \cdot B_k \tag{5}$$

 $B_n = \sum_{k=0}^n S(n,k)$, where S(n,k) is stirling number of second kind.

5. Some identities

Vandermonde's Identify:
$$\sum_{k=0}^{r} {m \choose k} {n \choose r-k} = {m+n \choose r}$$

Hockey-Stick Identify:
$$n, r \in N, n > r, \sum_{i=r}^{n} {i \choose r} = {n+1 \choose r+1}$$

Involutions: permutations such that p^2 = identity permutation. $a_0 = a_1 = 1$ and $a_n = a_{n-1} + (n-1)a_{n-2}$ for n > 1. 7.2 Fibonacci

Let A,B and n be integer numbers

$$k = A - B$$

$$F_A F_B = F_{k+1} F_A^2 + F_k F_A F_{A-1} \tag{6}$$

$$\sum_{i=0}^{n} F_i^2 = F_{n+1} F_n \tag{7}$$

$$\sum_{i=0}^{n} F_i F_{i+1} = F_{n+1}^2 - (-1)^n \tag{8}$$

$$\sum_{i=0}^{n} F_i F_{i-1} = \sum_{i=0}^{n-1} F_i F_{i+1}$$
 (9)

$$GCD(F_m, F_n) = F_{GCD(m,n)}$$
(10)

$$\sum_{0 \le k \le n} \binom{n-k}{k} = Fib_{n+1} \tag{11}$$

$$\gcd(F_n, F_{n+1}) = \gcd(F_n, F_{n+2}) = \gcd(F_{n+1}, F_{n+2}) = 1 \tag{12}$$

7.3 Notes

7.4 Geometry

7.4.1 Triangles

Circumradius: $R=\frac{abc}{4A}$, Inradius: $r=\frac{A}{s}$ Length of median (divides triangle into two equal-area triangles):

 $m_a = \frac{1}{2}\sqrt{2b^2 + 2c^2 - a^2}$

Length of bisector (divides angles in two):
$$s_a = \sqrt{bc \left[1 - \left(\frac{a}{b+c}\right)^2\right]}$$

Law of tangents:
$$\frac{a+b}{a-b} = \frac{\tan \frac{\alpha+\beta}{2}}{\tan \frac{\alpha-\beta}{2}}$$

7.4.2 Quadrilaterals

With side lengths a, b, c, d, diagonals e, f, diagonals angle θ , area A and magic flux $F = b^2 + d^2 - a^2 - c^2$:

$$4A = 2ef \cdot \sin\theta = F \tan\theta = \sqrt{4e^2f^2 - F^2}$$

For cyclic quadrilaterals the sum of opposite angles is 180° , ef = ac+bd, and $A=\sqrt{(p-a)(p-b)(p-c)(p-d)}$.

7.4.3 Spherical coordinates

$$x = r \sin \theta \cos \phi \qquad r = \sqrt{x^2 + y^2 + z^2}$$

$$y = r \sin \theta \sin \phi \qquad \theta = a\cos(z/\sqrt{x^2 + y^2 + z^2})$$

$$z = r \cos \theta \qquad \phi = a\tan(2(y, x))$$

1.3 Sums
$$1^2 + 2^2 + 3^2 + \dots + n^2 = \frac{n(2n+1)(n+1)}{6}$$

$$1^3 + 2^3 + 3^3 + \dots + n^3 = \frac{n^2(n+1)^2}{4}$$

$$1^4 + 2^4 + 3^4 + \dots + n^4 = \frac{n(n+1)(2n+1)(3n^2 + 3n - 1)}{30}$$

$$\sum_{i=1}^n i^m = \frac{1}{m+1} \left[(n+1)^{m+1} - 1 - \sum_{i=1}^n \left((i+1)^{m+1} - i^{m+1} - (m+1)i^m \right) \right]$$

$$\sum_{i=1}^{n-1} i^m = \frac{1}{m+1} \sum_{k=0}^m \binom{m+1}{k} B_k n^{m+1-k}$$

$$\sum_{i=1}^n kx^k = (x - (n+1)x^{n+1} + nx^{n+2})/(x-1)^2$$
7.6 Series

$$e^{x} = 1 + x + \frac{x^{2}}{2!} + \frac{x^{3}}{3!} + \dots, (-\infty < x < \infty)$$

$$\ln(1+x) = x - \frac{x^{2}}{2} + \frac{x^{3}}{3} - \frac{x^{4}}{4} + \dots, (-1 < x \le 1)$$

$$\sqrt{1+x} = 1 + \frac{x}{2} - \frac{x^{2}}{8} + \frac{2x^{3}}{32} - \frac{5x^{4}}{128} + \dots, (-1 \le x \le 1)$$

$$\sin x = x - \frac{x^{3}}{3!} + \frac{x^{5}}{5!} - \frac{x^{7}}{7!} + \dots, (-\infty < x < \infty)$$

$$\cos x = 1 - \frac{x^{2}}{2!} + \frac{x^{4}}{4!} - \frac{x^{6}}{6!} + \dots, (-\infty < x < \infty)$$

$$(x+a)^{-n} = \sum_{k=0}^{\infty} (-1)^{k} \binom{n+k-1}{k} x^{k} a^{-n-k}$$

7.7 Pythagorean Triples

The Pythagorean triples are uniquely generated by

$$a = k \cdot (m^2 - n^2), b = k \cdot (2mn), c = k \cdot (m^2 + n^2),$$

with m > n > 0, k > 0, $m \perp n$, and either m or n even.

7.8 Number Theory

7.8.1 Primes

p = 962592769 is such that $2^{21} \mid p-1$, which may be useful. For $x^2 + y^2 = 2z^2$ is equivalent to $(\frac{x+y}{2})^2 + (\frac{x-y}{2})^2 = z^2$. hashing use 970592641 (31-bit number), 31443539979727 (45-bit), 7.8.11 Postage stamps/McNuggets problem 3006703054056749 (52-bit). There are 78498 primes less than

Primitive roots exist modulo any prime power p^a , except for p = 2, a > 02, and there are $\phi(\phi(p^a))$ many. For p=2,a>2, the group \mathbb{Z}_{2a}^{\times} is 1=ab-a-b. instead isomorphic to $\mathbb{Z}_2 \times \mathbb{Z}_{2^{a-2}}$.

7.8.2 Estimates

 $\sum_{d|n} d = O(n \log \log n).$

n < 1e7, 2000 for n < 1e10, 200 000 for n < 1e19.

7.8.3 Perfect numbers

n > 1 is called perfect if it equals sum of its proper divisors and 1. Even n is perfect iff $n = 2^{p-1}(2^p - 1)$ and $2^p - 1$ is prime (Mersenne's). No odd perfect numbers are yet found.

7.8.4 Carmichael numbers

A positive composite *n* is a Carmichael number $(a^{n-1} \equiv 1 \pmod{n})$ for all $a: \gcd(a, n) = 1$), iff n is square-free, and for all prime divisors p of 7.9.2 Cycles n, p-1 divides n-1.

7.8.5 Mobius function

 $\mu(1) = 1$. $\mu(n) = 0$, if n is not squarefree. $\mu(n) = (-1)^s$, if n is the product of s distinct primes. Let f, F be functions on positive integers. If for all $n \in N$, $F(n) = \sum_{d \mid n} f(d)$, then $f(n) = \sum_{d \mid n} \mu(d) F(\frac{n}{d})$, and vice versa. $\phi(n) = \sum_{d|n} \mu(d) \frac{n}{d}. \quad \sum_{d|n} \mu(d) = 1.$

If f is multiplicative, then $\sum_{d|n} \mu(d) f(d) = \prod_{p|n} (1 - f(p))$, Permutations of a set such that none of the elements appear in their $\sum_{d|n} \mu(d)^2 f(d) = \prod_{p|n} (1 + f(p)).$

7.8.6 Legendre symbol

If p is an odd prime, $a \in \mathbb{Z}$, then $\left(\frac{a}{p}\right)$ equals 0, if p|a; 1 if a is a quadratic residue modulo p; and -1 otherwise. Euler's criterion: 7.9.4 Burnside's lemma $\left(\frac{a}{p}\right) = a^{\left(\frac{p-1}{2}\right)} \pmod{p}.$

7.8.7 Jacobi symbol

If $n = p_1^{a_1} \cdots p_k^{a_k}$ is odd, then $\left(\frac{a}{n}\right) = \prod_{i=1}^k \left(\frac{a}{n_i}\right)^{k_i}$.

If the order of g modulo m (min n > 0: $g^n \equiv 1 \pmod{m}$) is $\phi(m)$, then If f(n) counts "configurations" (of some sort) of length n, we can ignore g is called a primitive root. If Z_m has a primitive root, then it has $\phi(\phi(m))$ distinct primitive roots. Z_m has a primitive root iff m is one of 2, 4, p^k , $2p^k$, where p is an odd prime. If Z_m has a primitive root g, then for all a coprime to m, there exists unique integer $i = \text{ind}_{\mathcal{G}}(a)$ modulo $\phi(m)$, such that $g^i \equiv a \pmod{m}$. ind_{σ}(a) has logarithm-like properties: ind(1) = 0, ind(ab) = ind(a) + ind(b).

If p is prime and a is not divisible by p, then congruence $x^n \equiv a \pmod{p}$ has $\gcd(n, p-1)$ solutions if $a^{(p-1)/\gcd(n, p-1)} \equiv 1 \pmod{p}$, and no solutions otherwise. (Proof sketch: let g be a primitive root, and $g^i \equiv a \pmod{p}, g^u \equiv x \pmod{p}. x^n \equiv a \pmod{p} \text{ iff } g^{nu} \equiv g^i \pmod{p}$ iff $nu \equiv i \pmod{p}$.)

7.8.9 Discrete logarithm problem

Find x from $a^x \equiv b \pmod{m}$. Can be solved in $O(\sqrt{m})$ time and space with a meet-in-the-middle trick. Let $n = \lceil \sqrt{m} \rceil$, and x = ny - z. Equation becomes $a^{ny} \equiv ba^z \pmod{m}$. Precompute all values that the RHS can take for $z = 0, 1, \dots, n-1$, and brute force y on the LHS, each time 7.11.1 Stirling numbers of the first kind checking whether there's a corresponding value for RHS.

7.8.10 Pythagorean triples

Integer solutions of $x^2 + y^2 = z^2$ All relatively prime triples are given by: $x = 2mn, y = m^2 - n^2, z = m^2 + n^2$ where $m > n, \gcd(m, n) = 1$ and $m \not\equiv n \pmod{2}$. All other triples are multiples of these. Equation

Let a, b be relatively-prime integers. There are exactly $\frac{1}{2}(a-1)(b-1)$ numbers not of form ax + by $(x, y \ge 0)$, and the largest is (a-1)(b-1)

7.8.12 Fermat's two-squares theorem

Odd prime p can be represented as a sum of two squares iff $p \equiv$ 1 (mod 4). A product of two sums of two squares is a sum of two The number of divisors of n is at most around 100 for n < 5e4, 500 for squares. Thus, n is a sum of two squares iff every prime of form p = 4k + 3 occurs an even number of times in *n*'s factorization.

7.9 Permutations

ractoriai									
	n	123	4 5	6	7	8	9	10	
	n!	12.6	24.12	0.720	5040	40320	362880	3628800	
	n	11	12	13	14	. 18	5 16	<u> </u>	
	n!	4.0e7	4.8e8	6.2e	98.7e	$10 \ 1.3e$	$212\ 2.1\epsilon$	13 3.6e14	
	n	20	25	30	40	50 1	.00 1:	50 171	
	n!	2e18	2e25	$3\mathrm{e}32$ 8	3e47 3	3e64 9e	e157~6e2	262 >DBL_MAX	

Let $g_S(n)$ be the number of *n*-permutations whose cycle lengths all belong to the set S. Then

$$\sum_{n=0}^{\infty} g_S(n) \frac{x^n}{n!} = \exp\left(\sum_{n \in S} \frac{x^n}{n}\right)$$

7.9.3 Derangements

original position.

$$D(n) = (n-1)(D(n-1) + D(n-2)) = nD(n-1) + (-1)^n = \left\lfloor \frac{n!}{e} \right\rfloor$$

Given a group G of symmetries and a set X, the number of elements of *X up to symmetry* equals

$$\frac{1}{|G|} \sum_{g \in G} |X^g|,$$

where X^g are the elements fixed by g(g.x = x).

rotational symmetry using $G = \mathbb{Z}_n$ to get

$$g(n) = \frac{1}{n} \sum_{k=0}^{n-1} f(\gcd(n,k)) = \frac{1}{n} \sum_{k|n} f(k)\phi(n/k)$$

7.10 Partitions and subsets

7.10.1 Partition function

Number of ways of writing n as a sum of positive integers, disregarding the order of the summands.

$$p(0) = 1, \ p(n) = \sum_{k \in \mathbb{Z} \setminus \{0\}} (-1)^{k+1} p(n - k(3k - 1)/2)$$

$$p(n) \sim 0.145/n \cdot \exp(2.56\sqrt{n})$$

$$\frac{n}{p(n)} \mid 0.1234567892050100$$

$$\frac{n}{p(n)} \mid 1.1235711152230627 \sim 2e5 \sim 2e8$$

7.11 General purpose numbers

Number of permutations on n items with k cycles.

$$c(n,k) = c(n-1,k-1) + (n-1)c(n-1,k), \ c(0,0) = 1$$

$$\sum_{k=0}^{n} c(n,k)x^{k} = x(x+1)...(x+n-1)$$

c(8,k) = 8,0,5040,13068,13132,6769,1960,322,28,1c(n,2) = 0.0.1.3.11.50.274.1764.13068.109584...

7.11.2 Eulerian numbers

Number of permutations $\pi \in S_n$ in which exactly k elements are greater than the previous element. k j:s s.t. $\pi(j) > \pi(j+1)$, k+1 j:s s.t. $\pi(i) \ge i, k \text{ i:s s.t. } \pi(i) > i.$

$$E(n,k) = (n-k)E(n-1,k-1) + (k+1)E(n-1,k)$$

$$E(n,0) = E(n,n-1) = 1$$

$$E(n,k) = \sum_{i=0}^{k} (-1)^{i} \binom{n+1}{i} (k+1-j)^{n}$$

7.11.3 Stirling numbers of the second kind

Partitions of n distinct elements into exactly k groups.

$$S(n,k) = S(n-1,k-1) + kS(n-1,k)$$

$$S(n,1) = S(n,n) = 1$$

$$S(n,k) = \frac{1}{k!} \sum_{j=0}^{k} (-1)^{k-j} \binom{k}{j} j^n$$

7.11.4 Bell numbers

Total number of partitions of n distinct elements. B(n) =1,1,2,5,15,52,203,877,4140,21147,... For p prime,

$$B(p^m + n) \equiv mB(n) + B(n+1) \pmod{p}$$

7.11.5 Bernoulli numbers

$$\sum_{j=0}^{m} {m+1 \choose j} B_j = 0. \quad B_0 = 1, B_1 = -\frac{1}{2}. \ B_n = 0, \text{ for all odd } n \neq 1.$$

7.11.6 Catalan numbers

$$C_n = \frac{1}{n+1} \binom{2n}{n} = \binom{2n}{n} - \binom{2n}{n+1} = \frac{(2n)!}{(n+1)!n!}$$

$$C_0 = 1, \ C_{n+1} = \frac{2(2n+1)}{n+2}C_n, \ C_{n+1} = \sum C_i C_{n-i}$$

 $C_n = 1, 1, 2, 5, 14, 42, 132, 429, 1430, 4862, 16796, 58786, \dots$

- sub-diagonal monotone paths in an $n \times n$ grid.
- strings with *n* pairs of parenthesis, correctly nested.
- binary trees with with n + 1 leaves (0 or 2 children).
- ordered trees with n+1 vertices.
- ways a convex polygon with n+2 sides can be cut into triangles const int N=?; by connecting vertices with straight lines.
- permutations of [n] with no 3-term increasing subseq.

7.12 Inequalities

7.12.1 Titu's Lemma

For positive reals $a_1, a_2, ..., a_n$ and $b_1, b_2, ..., b_n$,

$$\frac{a_1^2}{b_1} + \frac{a_2^2}{b_2} + \ldots + \frac{a_n^2}{b_n} \ge \frac{a_1 + a_2 + \ldots + a_n^2}{b_1 + b_2 + \ldots + b_n}$$

Equality holds if and only if $a_i = kb_i$ for a non-zero real constant k.

7.13 Games

7.13.1 Grundy numbers

For a two-player, normal-play (last to move wins) game on a graph (V, E): $G(x) = \max(\{G(y) : (x, y) \in E\})$, where $\max(S) = \min\{n \ge 0 : n \not\in E\}$ S). x is losing iff G(x) = 0.

- 7.13.2 Sums of games
 Player chooses a game and makes a move in it Grundy number of a position is xor of grundy numbers of positions in summed
 - Player chooses a non-empty subset of games (possibly, all) and makes moves in all of them A position is losing iff each game is in a losing position.
 - Player chooses a proper subset of games (not empty and not all), and makes moves in all chosen ones. A position is losing iff grundy numbers of all games are equal.
 - Player must move in all games, and loses if can't move in some game A position is losing if any of the games is in a losing position.

7.13.3 Misère Nim

A position with pile sizes $a_1, a_2, ..., a_n \ge 1$, not all equal to 1, is losing iff $a_1 \oplus a_2 \oplus \cdots \oplus a_n = 0$ (like in normal nim.) A position with n piles of size 1 is losing iff n is odd.

7.14 NumberTheory

$$\sum_{k=1}^{n} \frac{1}{\gcd(k,n)} = \sum_{d|n} \frac{1}{d} \cdot \phi\left(\frac{n}{d}\right) = \frac{1}{n} \sum_{d|n} d \cdot \phi(d)$$
 (13)

$$\sum_{k=1}^{n} \frac{k}{\gcd(k,n)} = \frac{n}{2} \cdot \frac{1}{n} \cdot \sum_{d|n} d \cdot \phi(d)$$
 (14)

$$\sum_{k=1}^{n} \frac{n}{\gcd(k,n)} = 2 * \sum_{k=1}^{n} \frac{k}{\gcd(k,n)} - 1, \text{ for, } n > 1$$
 (15)

$$\sum_{i=1}^{n} \sum_{j=1}^{n} [\gcd(i,j) = 1] = \sum_{d=1}^{n} \mu(d) \left\lfloor \frac{n}{d} \right\rfloor^{2}$$
 (16)

$$\sum_{i=1}^{n} \sum_{j=1}^{n} \gcd(i,j) = \sum_{d=1}^{n} \phi(d) \left\lfloor \frac{n}{d} \right\rfloor^{2}$$
(17)

$$\sum_{i=1}^{n} \sum_{j=1}^{n} i \cdot j[\gcd(i,j) = 1] = \sum_{i=1}^{n} \phi(i)i^{2}$$
 (18)

$$\sum_{i=1}^{n} \sum_{j=1}^{n} \operatorname{lcm}(i,j) = \sum_{l=1}^{n} \left(\frac{\left(1 + \left\lfloor \frac{n}{l} \right\rfloor\right) \left(\left\lfloor \frac{n}{l} \right\rfloor\right)}{2} \right)^{2} = \sum_{d|l} \mu(d) l d$$
 (19)

8 String

8.1 aho_corasick

using namespace std;

```
const int A = ?;
struct AC {
  int nd, pt;
  int next[N][A], link[N], out link[N], cnt[N], ans[N];
  vector <int> ed[N], out[N];
  AC(): nd(0), pt(0) \{ node(); \}
  int node() {
    memset(next[nd], 0, sizeof next[nd]);
link[nd] = out_link[nd] = cnt[nd] = 0;
ed[nd].clear();
out[nd].clear();
    return nd++:
  void clear() {
    nd = pt = 0;
    node();
  inline int get(char c) { return c - 'a'; }
  void insert(const string &T) {
    int u = 0;
     for (char'c : T) {
   if (!next[u][get(c)]) next[u][get(c)] = node();
       u = next[u][qet(c)];
     ans[pt] = 0;
    out[u].push back(pt++);
  void build() {
    queue <int> q;
     for (q.push(0); !q.empty(); ) {
       int u = a.front():
       q.pop();
       int c = 0; c < A; ++c) {
  int v = next[u][c];
  if (!y) next[u][c] = next[link[u]][c];</pre>
            link[v] = u ? next[link[u]][c] : 0;
```

```
out link[v] = out[link[v]].empty() ?
                  out link[link[v]] : link[v];
             ed[link[v]].push back(v);
             q.push(v);
  void dfs(int s) {
     for(int x : ed[s]) dfs(x), cnt[s] += cnt[x];
     for(int e : out[s]) ans[e] = cnt[s];
  void traverse(const string &S) {
     for (char c : S) {
   u = next[u][get(c)];
       cnt[u]++;
     dfs(0);
char str[1000010], pat[505];
int main() {
           freopen("in.txt","r",stdin);
 int t,T;
scanf("%d",&T);
for(int t=1;t<=T;t++) {</pre>
    int n;
scanf("%d",&n);
scanf("%s",str);
for(int i=1;i<=n;i++) {</pre>
       scanf("%s",pat);
aho.insert(pat);
     aho.build();
     aho.traverse(str);
printf("Case %d:\n",t);
for(int i=0;i<n;i++) {</pre>
       printf("%d\n",aho.ans[i]);
     aho.clear();
  return 0:
```

```
8.2 hash
struct Hash {
  struct base {
    string s; int b, mod;
     vector<int> hash, p;
     void init(string & s, int b, int mod) \{ // b >
     → 26, prime.
       s = s; b = b, mod = mod;
       hash_resize(s.size());
       p.resize(s.size());
       hash[0] = s[0] - 'A' + 1; p[0] = 1;
       for(int i = 1; i < s.size(); ++i) {
    hash[i] = (ll) hash[i - 1] * b % mod;
    hash[i] += s[i] - 'A' + 1;
    if(hash[i] >= mod) hash[i] -= mod;
    p[i] = (ll) p[i - 1] * b % mod;
    int get(int l, int r) {
       int ret = hash[r];
       if(l) ret -= (ll) hash[l - 1] * p[r - l + 1] %
       if(ret < 0) ret += mod;
       return ret;
  void init(string &s) {
    h[0].init(s, 29, 1e9+7);
h[1].init(s, 31, 1e9+9);
  pair<int, int> get(int l, int r) {
    return { h[0].get(l, r), h[1].get(l, r) };
8.3 hash_segtree
```

```
#define INVALID CHAR
namespace strhash {
  const int MAX = 100010;
  int ara[MAX]
  const int MOD[] = {1067737007, 1069815139};
  const int BASE[] = {982451653, 984516781};
  int BP[2][MAX], CUM[2][MAX];
  void init(char *str) {
   n = strlen(str);
    for(int i=0;i<n;i++) ara[i] = str[i]-'0'+1;</pre>
  void precal() {
   BP[0][0] = BP[1][0] = 1;
   CUM[0][0] = CUM[1][0] = 1;
    for(int i=1;i<MAX;i++)</pre>
      BP[0][i] = (BP[0][i-1] * (long long) BASE[0])
         _%_MOD[0];
      BP[1][i] = (BP[1][i-1] * (long long) BASE[1])
      → % MOD[1];
      CUM[0][i] = (CUM[0][i-1] + (long long) BP[0][i]
         ) % MOD[0];
      CUM[1][i] = (CUM[1][i-1] + (long long) BP[1][i]
       → ) % MOD[1];
  struct node {
   int sz:
    int h[2]
   node()
  } tree[4*MAX];
  int lazy[4*MAX];
  inline void lazyUpdate(int n,int st,int ed) {
   if(lazy[n]!=INVALID CHAR){
```

```
tree[n].h[0] = (lazy[n] * (long long)

→ CUM[1][ed-st]) % MOD[1];

    if(st!=ed){
      lazy[2*n] = lazy[n];
     lazy[2*n+1] = lazy[n];
    lazy[n] = INVALID CHAR;
inline node Merge(node a, node b) {
 node ret;
  ret.h[0] = ( (a.h[0] * (long long) BP[0][b.sz] )

→ + b.h[0] ) % MOD[0];
  ret.h[1] = ((a.h[1] * (long long) BP[1][b.sz])
  \rightarrow + b.h[1] ) % MOD[1];
  ret.sz = a.sz + b.sz;
 return ret:
inline void build(int n.int st.int ed) {
  lazy[n] = INVALÍD CHAŘ;
  if(st==ed) {
    tree[n].h[0] = tree[n].h[1] = ara[st];
    tree[n].sz = 1;
    return;
  int mid = (st+ed)>>1;
  build(n+n,st,mid);
  build(n+n+1,mid+1,ed);
 tree[n] = Merge(tree[n+n], tree[n+n+1]);
inline void update(int n,int st,int ed,int i,int
lazyUpdate(n,st,ed);
 if(st>j or ed<i) return;</pre>
 if(st>=i and ed<=j) {
    lazy[n] = v;
    lazyUpdate(n,st,ed);
    return:
 int mid = (st+ed)>>1;
 update(n+n,st,mid,i,j,v);
 update(n+n+1,mid+1,ed,i,j,v);
 tree[n] = Merge(tree[n+n],tree[n+n+1]);
inline node query(int n,int st,int ed,int i,int j){
  lazyUpdate(n,st,ed);
  if(st>=i and ed<=j) return tree[n];</pre>
  int mid = (st+ed)/2;
  if(mid<i) return query(n+n+1,mid+1,ed,i,j);</pre>
  else if(mid>=j) return query(n+n,st,mid,i,j);
  else return Merge(query(n+n,st,mid,i,j),query(n+n+

→ 1,mid+1,ed,i,j));
```

```
8.4 kmp
// returns the longest proper prefix array of pattern p
// where lps[i]=longest proper prefix which is also
\rightarrow suffix of p[0...i]
vector<<mark>int</mark>> build lps(string p) {
  int sz = p.size();
  vector<int> lps;
  lps.assign(sz + 1, 0);
  int j = \bar{0};
  lps[0] = 0;
  for(int i = 1; i < sz; i++) {
```

```
20
    while(j >= 0 \&\& p[i] != p[j]) {
      if(j >= 1) j = lps[j - 1];
      else j = -1;
    lps[i] = j;
  return lps;
vector<int>ans;
// returns matches in vector ans in 0-indexed
void kmp(vector<int> lps, string s, string p) {
  int psz = p.size(), sz = s.size();
  int j = 0;
 for(int i = 0; i < sz; i++) {
  while(j >= 0 && p[j] != s[i])
      if(j >= 1) j = lps[j - 1];
      else j = -1;
     ++:
    if(j == psz) {
        = lps[j - 1];
      j = lps[] - 1];
// pattern found in string s at position i-psz+1
      ans.push back(i - psz + 1);
    // after each loop we have j=longest common suffix
     \rightarrow of s[0..i] which is also prefix of p
```

8.5 manachar

```
vector<int> d1(n); // maximum odd length palindrome

→ centered at i

                     // here d1[i]=the palindrome has

→ d1[i]-1 right characters from i
                     // e.g. for aba, d1[1]=2;
for (int i = 0, l = 0, r = -1; i < n; i++) {
 int k = (i > r) ? 1 : min(d1[1 + r - i], r - i);

while (0 <= i - k && i + k < n && s[i - k] == s[i + k]

    k]) {

   k++;
 d1[i] = k--;
 if (i + k > r) {
   l = i - k;
   r = i + k;
vector<int> d2(n); // maximum even length palindrome
// here d2[i]=the palindrome has

→ d2[i]-1 right characters from i
                     // e.g. for abba, d2[2]=2;
for (int i = 0, l = 0, r = -1; i < n; i++) {
 int k = (i > r) ? 0' : min(d2[l + r' - i + 1], r - i + 1]
 while (0 \le i - k - 1) \& i + k < n \& s[i - k - 1]
  \rightarrow == s[i + k]) {
   k++;
 d2[i] = k--;
 if(i + k > r) {
   l'= i - k - 1:
    r = i + k;
```

8.6 palindromic_tree

```
const int A = 26:
const int N = 300010;
char s[N]; long long ans;
```

```
int last, ptr, nxt[N][A], link[N], len[N], occ[N];
void feed (int at)
  while (s[at - len[last] - 1] != s[at]) last =
     link[last];
  int ch = s[at] - 'a', temp = link[last];
  while (s[at - len[temp] - 1] != s[at]) temp =

    link[temp];
  if (!nxt[last][ch]) {
   nxt[last][ch] = ++ptr, len[ptr] = len[last] + 2;
    link[ptr] = len[ptr] == 1 ? 2 : nxt[temp][ch];
  last = nxt[last][ch], ++occ[last];
int main() {
  len[1] = -1, len[2] = 0, link[1] = link[2] = 1, last
  for (int i = 1, n = strlen(s + 1); i \le n; ++i)
   → feed(i);
  for (int i = ptr; i > 2; --i) ans = max(ans, len[i]
   * 1LL * occ[i]), occ[link[i]] += occ[i];
  printf("%lld\n", ans);
  return 0;
8.7 persistant_trie
const int MAX = 200010:
const int B = 19;
int root[MAX], ptr = 0;
struct node {
  int ara[2], sum;
node() {}
} tree[MAX * (B+1)];
void insert(int prevnode, int &curRoot, int val) {
  curRoot = ++ptr;
  int curnode = curRoot;
  for(int i = B; i >= 0; i--) {
  bool bit = val & (1 << i);
  tree[curnode] = tree[prevnode];
  tree[curnode].ara[bit] = ++ptr;</pre>
    tree[curnode].sum += 1
    prevnode = tree[prevnode].ara[bit];
    curnode = tree[curnode].ara[bit];
  tree[curnode] = tree[prevnode];
  tree[curnode].sum += 1;
int find xor max(int prevnode, int curnode, int x) {
  int ans = 0;
  for(int i = B; i >= 0; i--) {
    bool bit = x \& (1 << i);
    if(tree[tree[curnode].ara[bit ^ 1]].sum >
     curnode = tree[curnode].ara[bit ^ 1]; sum) {
  curnode = tree[curnode].ara[bit ^ 1];
  prevnode = tree[prevnode].ara[bit ^ 1];
       ans = ans \mid (1 << i);
    else {
      curnode = tree[curnode].ara[bit]
      prevnode = tree[prevnode].ara[bit];
  return ans;
void solve() {
  int n, q, L, R, K;
  cin >> n;
  for(int i=1;i<=n;i++) cin >> ara[i];
  for(int i=1;i<=q;i++) {
    cin >> L >> R >> K:
    cout << find xor max(root[L-1],root[R],K) << endl;
```

```
8.8 suffix_array
 // Everything is 0-indexed
char s[N]; // Suffix array will be built for this
int SA[N], iSA[N]; // SA is the suffix array, iSA[i]
stores the rank of the i'th suffix int cnt[N], nxt[N]; // Internal stuff bool bh[N], b2h[N]; // Internal stuff
|<mark>int</mark> lcp[N]; // Stores lcp of SA[i] and SA[i + 1];
 \rightarrow lcp[n - 1] = 0
int lcpSparse[LOGN][N]; // lcpSparse[i][j] =
 \rightarrow min(lcp[j], ..., lcp[j - 1 + (1 << i)])
void buildSA(int n) {
  for (int i = 0; i < n; i++) SA[i] = i;
   sort(SA, SA + n, [](int i, int j) { return s[i] <
   \hookrightarrow s[j]; \});
  for (int i = 0; i < n; i++) {
  bh[i] = i == 0 || s[SA[i]] != s[SA[i - 1]];</pre>
     b2h[i] = 0;
   for (int h = 1; h < n; h <<= 1) {
     int tot = 0;
     for (int i = 0, j; i < n; i = j) {
       j = i + 1;
       while (j < n && !bh[j]) j++;
       nxt[i] = j; tot++;
     } if (tot == n) break;
     for (int i = 0; i < n; i = nxt[i])</pre>
       for (int j = i; j < nxt[i]; j++) iSA[SA[j]] = i;</pre>
       cnt[i] = 0;
     b2h[iSA[n - h]] = 1;
     for (int i = 0; i < n; i = nxt[i])
       for (int j = i; j < nxt[i]; j++) {
         int s = SA[j] - h;
         if (s < 0) continue;
         int head = iSA[s];
         iSA[s] = head + cnt[head]++;
         b2h[iSA[s]] = 1;
       for (int j = i; j < nxt[i]; j++) {</pre>
         int s = SA[j] - h;
         if (s < 0 || !b2h[iSA[s]]) continue;</pre>
         for (int k = iSA[s] + 1; !bh[k] && b2h[k];
              k++) b2h[k] = 0;
     for (int i = 0; i < n; i++) {
       $A[i$A[i]]_=_i
       bh[i] [= b2h[i];
   for (int i = 0; i < n; i++) iSA[SA[i]] = i;
void buildLCP(int n) {
  for (int i = 0, k = 0; i < n; i++) {
     if (iSA[i] == n - 1) {
       k = 0;
       lcp[n - 1] = 0;
       continue;
     int j = SA[iSA[i] + 1];
     while (i + k < n \&\& j + k < n \&\& s[i + k] == s[j + k]
     lcp[iSA[i]] = k;
     if (k) k--;
```

```
void buildLCPSparse(int n) {
 for (int i = 0; i < n; i++) lcpSparse[0][i] = lcp[i];</pre>
 for (int i = 1; i < LOGN; i++) {
    for (int j = 0; j < n; j++) {
     lcpSparse[i][j] = min(lcpSparse[i - 1][j],
          lcpSparse[i - 1][min(n - 1, j + (1 << (i -
      □ 1)))]);
pair<int, int> minLCPRange(int n, int from, int

→ minLCP) {
 int r = from;
 for (int i = LOGN - 1; i >= 0; i --) {
   int jump = 1 << i;
   if (r + jump < n and lcpSparse[i][r] >= minLCP) r
    int l = from;
 for (int i = LOGN - 1; i >= 0; i --) {
    int jump = 1 << i;
   if (l - jump >= 0 and lcpSparse[i][l - jump] >=

→ minLCP) l -= iump:

  return make pair(l, r);
```

8.9 suffix_automata

```
namespace sa{
 const int MAXN = 100005 << 1;
 const int MAXC = 26;
 char str[MAXN];
 int n, sz, last;
int len[MAXN], link[MAXN], ed[MAXN][MAXC], cnt[MAXN];
bool terminal[MAXN];
 vector <int> G[MAXN];
 void init() {
   SET(ed[0]);
    len[0] = 0, link[0] = -1, sz = 1, last = 0,

    terminal[0] = false;

 inline int scale(char c) { return c-'a'; }
 void extend(char c) {
   int cur = sz++;
    terminal[cur] = false;
    cnt[cur] = 1;
    SET(ed[cur]);
len[cur] = len[last] + 1;
    int p = last;
   while (p != -1 \&\& ed[p][c]==-1) {
      ed[p][c] = cur;
      p = link[p];
   if (p == -1) link[cur] = 0;
   else {
     int q = ed[p][c];
      if (len[p] + 1 == len[q]) link[cur] = q;
      else {
        int clone = sz++;
        len[clone] = len[p] + 1;
        memcpy(ed[clone],ed[q],sizeof(ed[q]));
        link[clone] = link[q];
        while (p != -1 \&\& ed[p][c] == q) {
          ed[p][c] = clone;
          p = link[p];
```

```
link[q] = link[cur] = clone;
         cnt[clone] = 0;
         terminal[clone] = false;
     last = cur;
  void dfs(int s) {
    for(auto x : G[s]) dfs(x), cnt[s] += cnt[x];
  void build() {
    init();
    int n = strlen(str);
    for(int i=0;i<n;i++) extend(scale(str[i]));</pre>
    for(int i=1;i<sz;i++) G[link[i]].pb(i);</pre>
    for(int i=0;i<sz;i++) G[i].clear();</pre>
    for(int i=last;i!=-1;i=link[i]) terminal[i] = true;
8.10 trie
#define N
#define S
                   200000
26
int root, now;
int nxt[N][S], cnt[N];
void init(){
  root = now = 1;
CLR(nxt),CLR(cnt);
inline int scale(char ch) { return (ch - 'a'); }
inline void Insert(char s[],int sz){
  int cur = root, to;
  for(int i=0 ; i < sz ; i++){
    to = scale(s[i]) ;
if( !nxt[cur][to] ) nxt[cur][to] = ++now;
    cur = nxt[cur][to];
  cnt[cur]++;
inline bool Find(char s[],int sz){
  int cur = root, to;
  for(int i=0; i<sz; i++){
  to = scale(s[i]);
  if(!nxt[cur][to]) return false;</pre>
    cur = nxt[cur][to];
  return (cnt[cur]!=0);
inline void Delete(char s[],int sz){
  int cur = root, to;
for(int i=0; i<sz; i++){
  to = scale(s[i]);</pre>
    cur = nxt[cur][to];
  cnt[cur]--;
```

```
8.11 z_algo
const int N = 100010;
char s[N];
int t, n, z[N];
int main() {
    scanf("%s", s);
    n = strlen(s), z[0] = n;
    int L = 0, R = 0;
    for (int i = 1; i < n; ++i) {
        if (i > R) {
            L = R = i;
            while (R < n && s[R - L] == s[R]) ++R;
        z[i] = R - L; --R;</pre>
```

```
} else {
   int k = i - L;
   if (z[k] < R - i + 1) z[i] = z[k];
   else {
      L = i;
      while (R < n && s[R - L] == s[R]) ++R;
      z[i] = R - L; --R;
   }
}</pre>
```

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