

Lambda Calculus

Abstract Syntax

$X = \{x, y, z, \dots\}$ — variables

$C = \{X, Y, Z, \dots\}$ — constructors

$\Lambda =$	X	— variable
	$ \quad C$	— constructor
	$ \quad \lambda X. \Lambda$	— abstraction
	$ \quad \Lambda \Lambda$	— application

Simple Typing á la Curry

Types:

$$\mathcal{T} = A \mid \mathcal{T} \rightarrow \mathcal{T}$$

where A is a set of primitive types.

Environment $\Gamma : X \rightarrow \mathcal{T}$ is an assignment of types to (some) variables.

Typing judgements:

$$\Gamma \vdash x : \Gamma x \quad [\text{VAR}]$$

$$\frac{\Gamma, x : \tau \vdash M : \rho}{\Gamma \vdash \lambda x. M : \tau \rightarrow \rho} \quad [\text{ABSTR}]$$

$$\frac{\Gamma \vdash M : \tau \rightarrow \rho, \Gamma \vdash N : \tau}{\Gamma \vdash M N : \rho} \quad [\text{APP}]$$