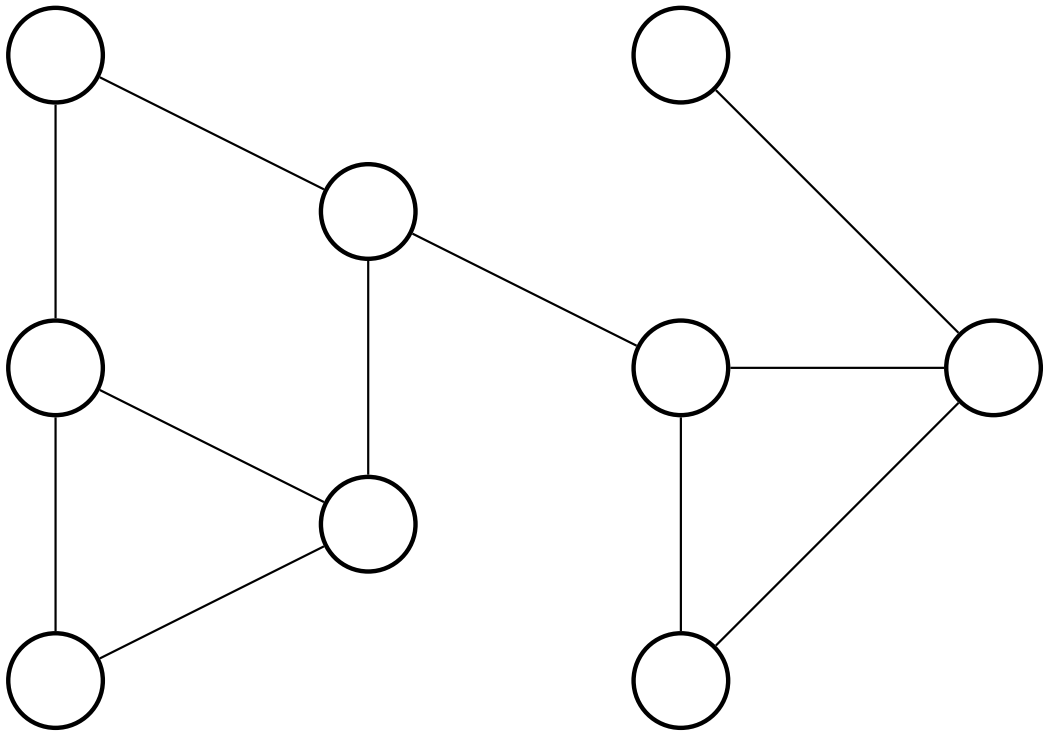
Dýptarleit og breiddarleit

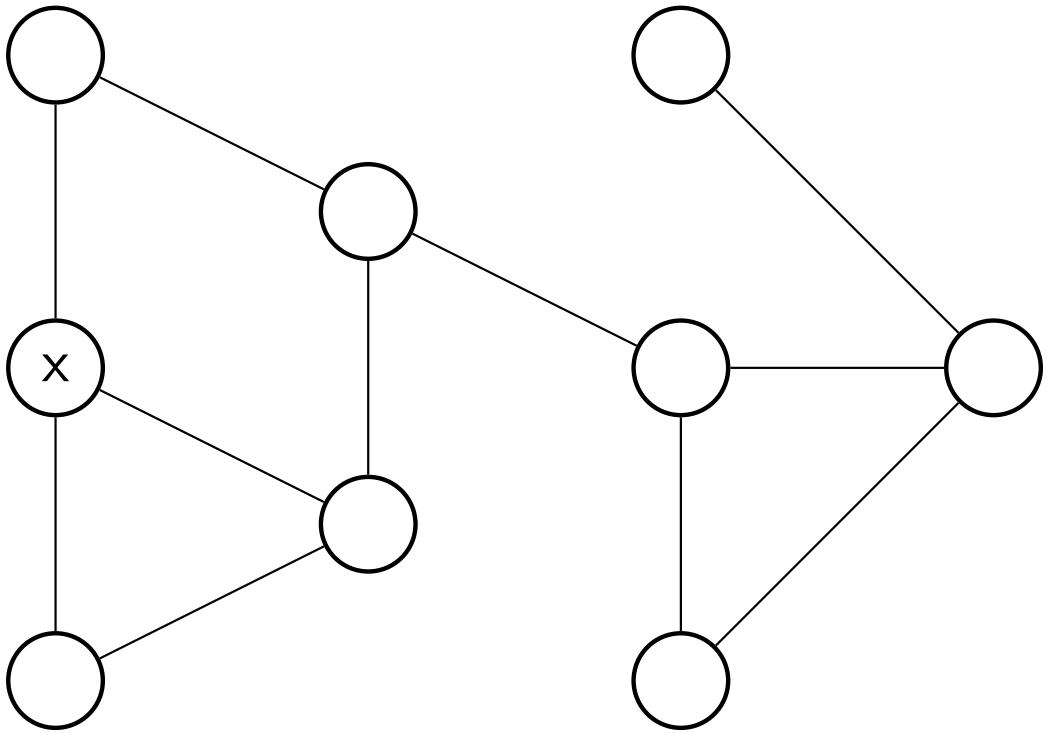
Bergur Snorrason

28. febrúar 2023

- Hvernig ítrum við í gegnum alla hnúta nets.
- Þetta má að sjálfsögðu gera á marga vegu, en algengt er að notast við annað að tvennu:
  - *Dýptarleit* (e. *deapth-first search*).
  - *Breiddarleit* (e. *breadth-first search*).
- Báðar aðferðir byggja á því að byrja í einhverjum hnút og heimsækja svo nágranna hans.

- Dýptarleit byrjar í einhverjum hnút.
- Sá hnútur er kallaður *upphafshnúturinn*.
- Í hverju skrefi heimsækir leitin einhvern nágranna hnútsins sem hefur ekki verið heimsóttur áður í leitinni.
- Ef allir nágrannar hafa verið heimsóttir þá er farið til baka og nágrannar síðasta hnúts eru skoðaðir.
- Tökum dæmi.

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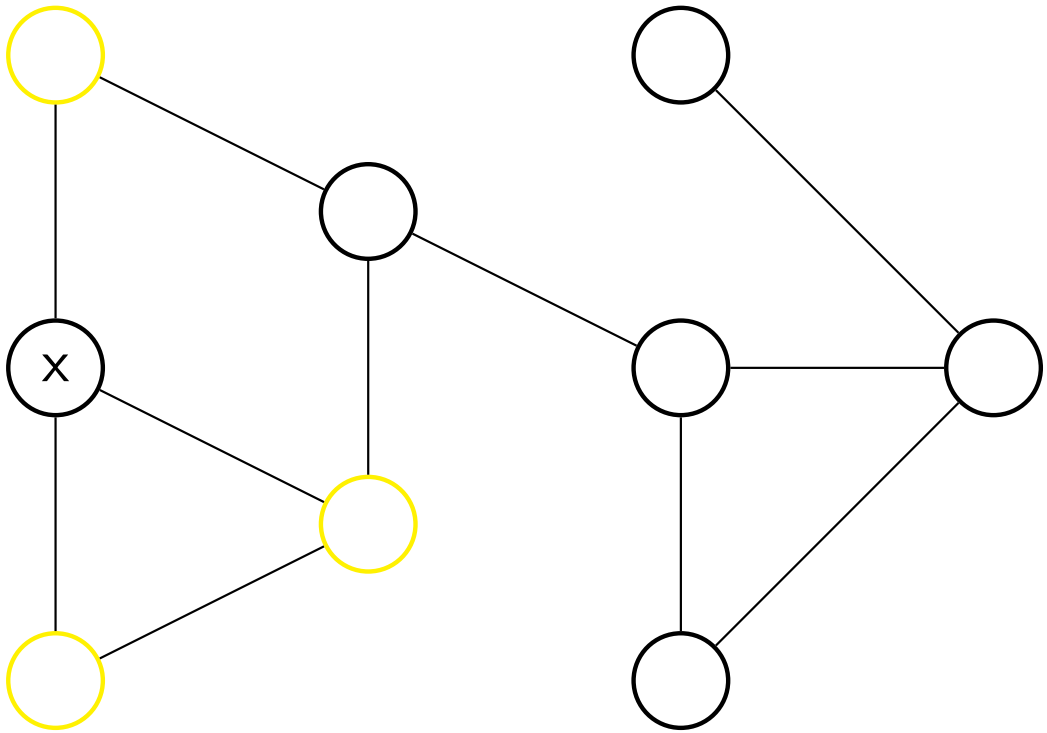
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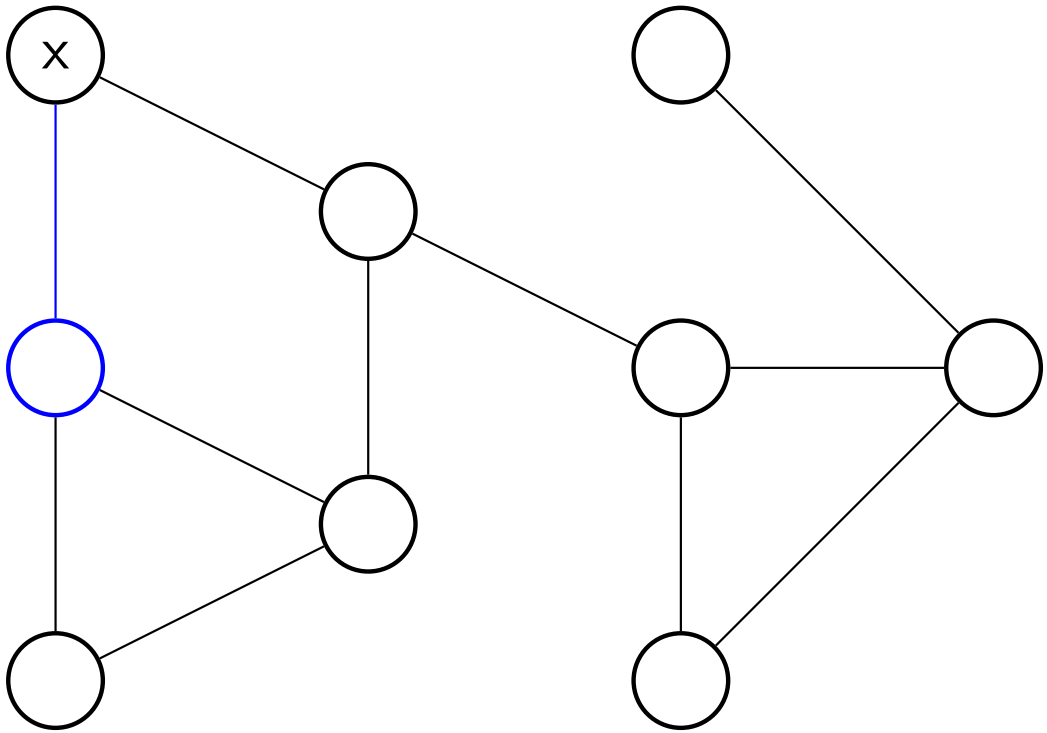
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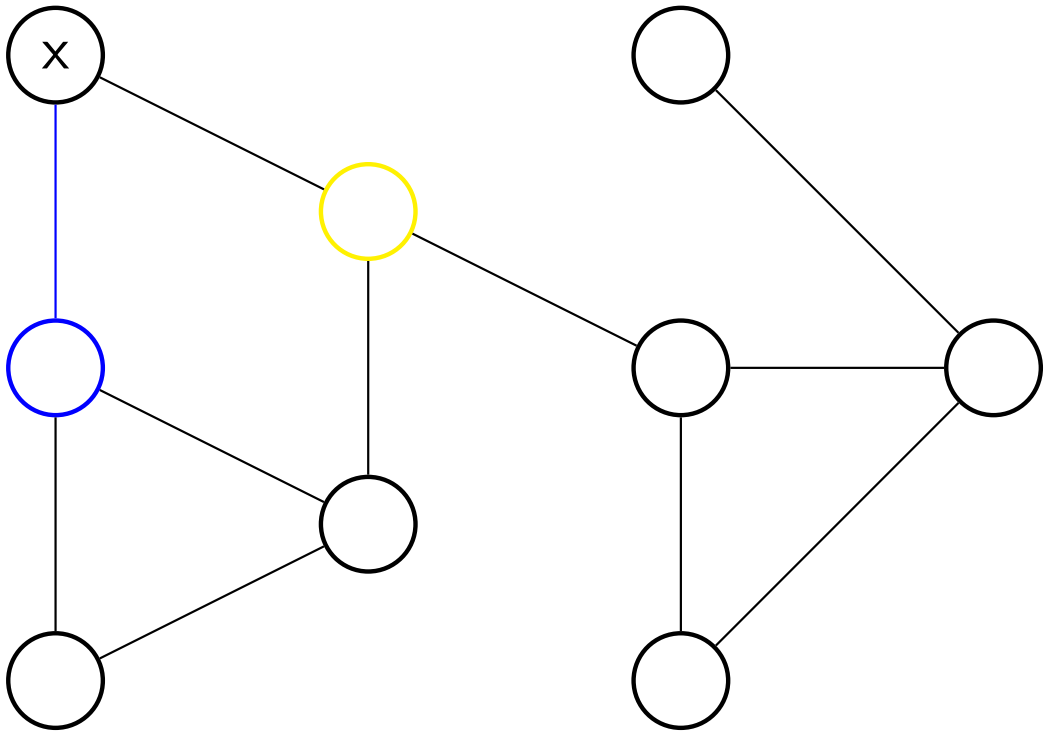
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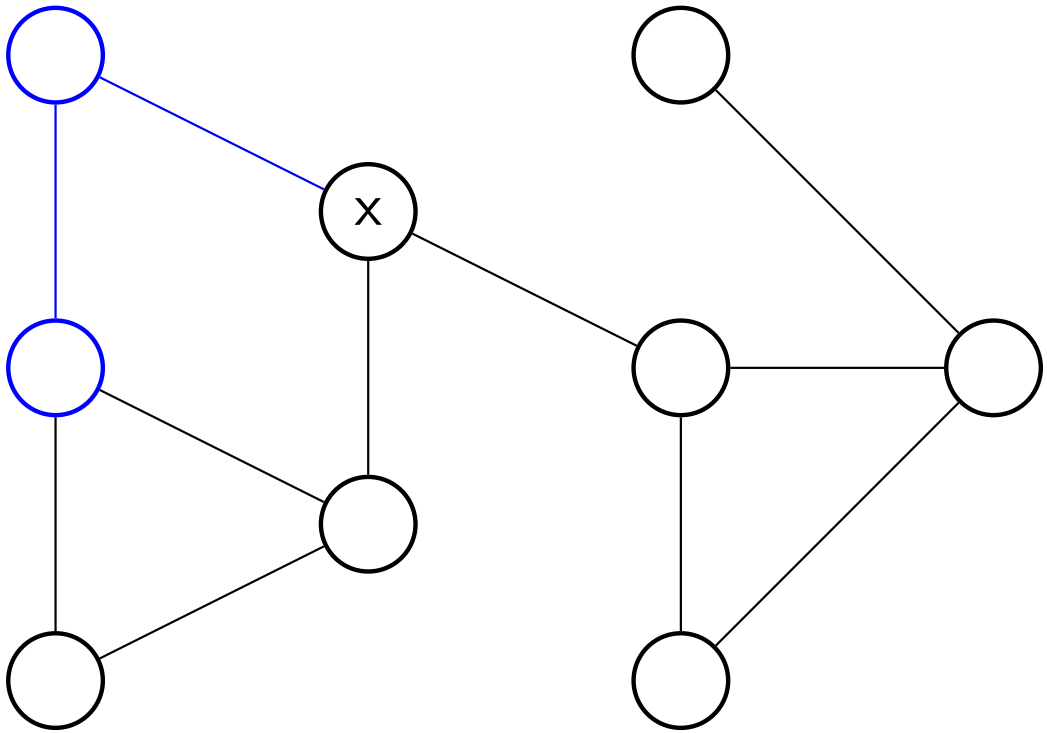
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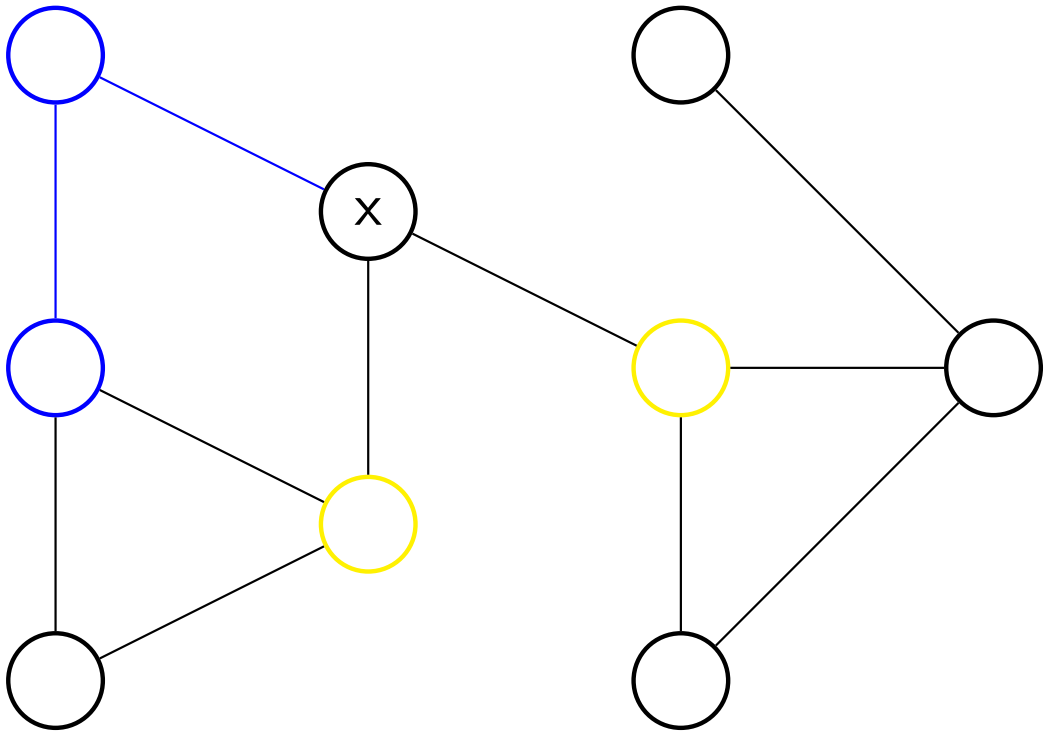
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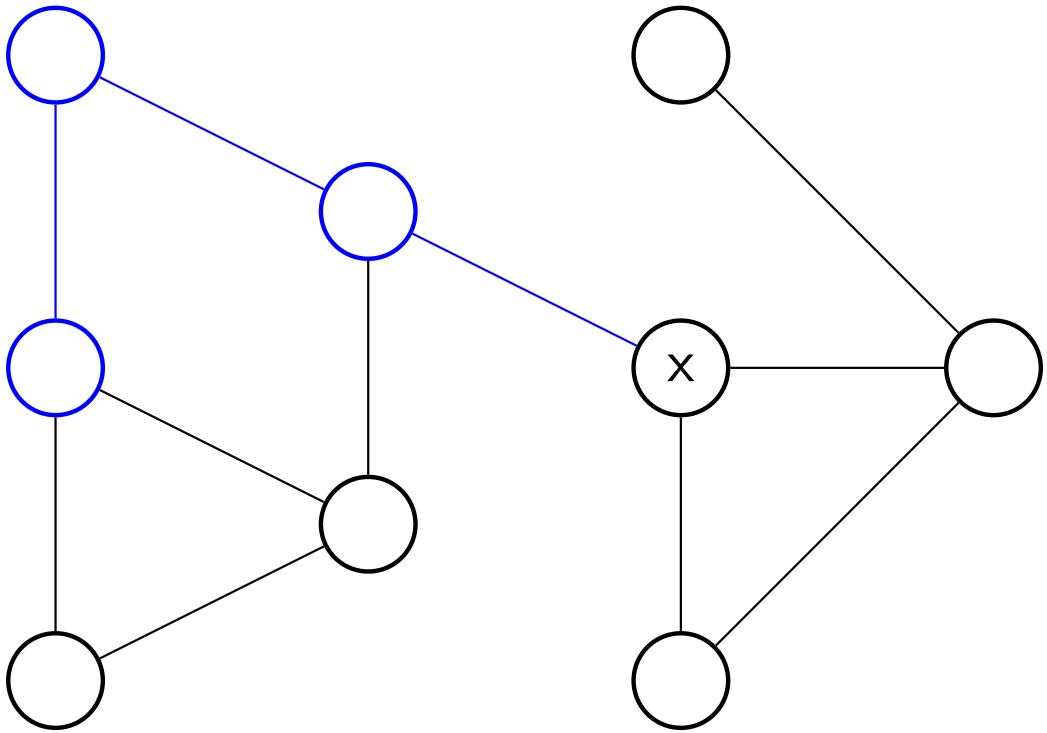
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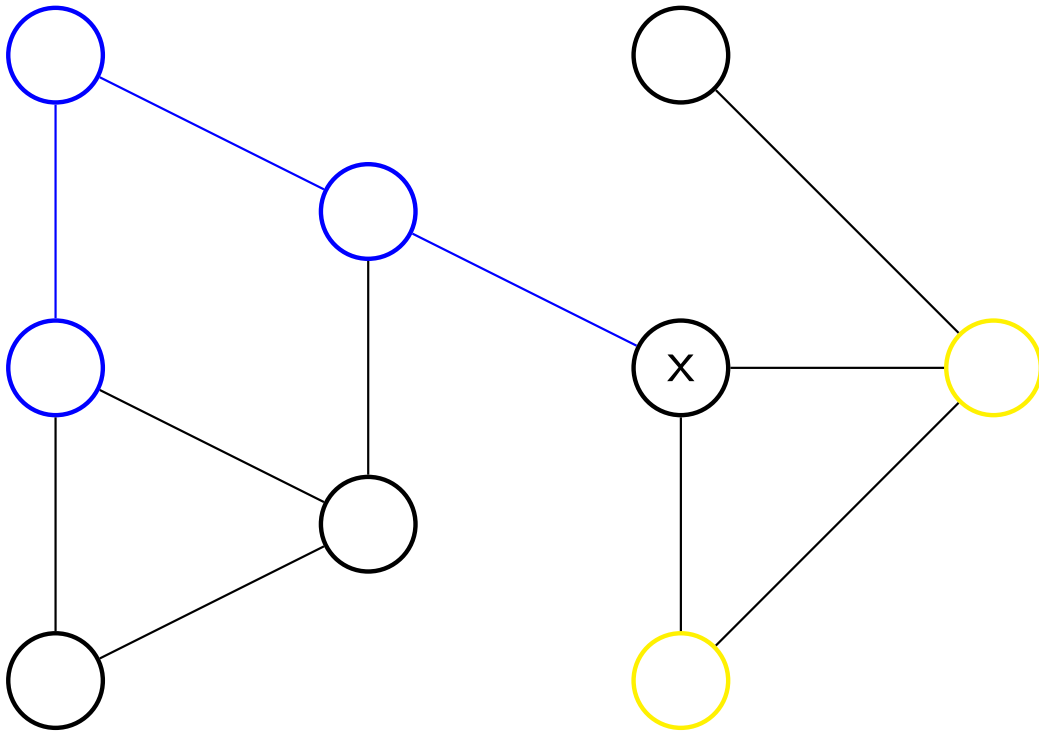
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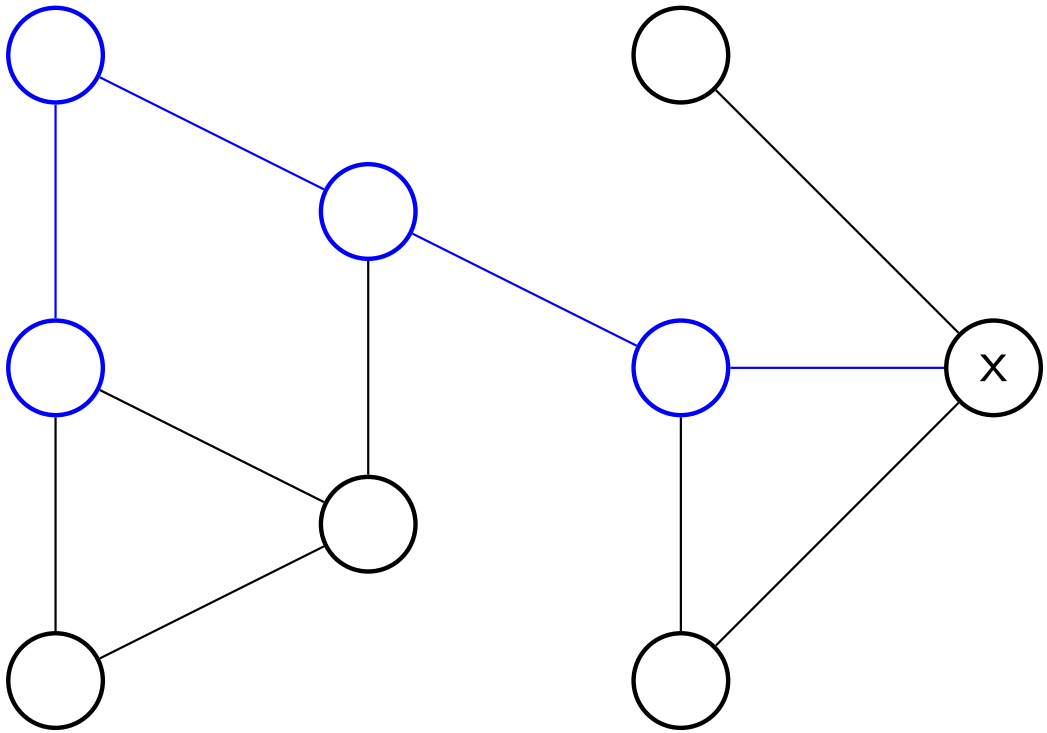
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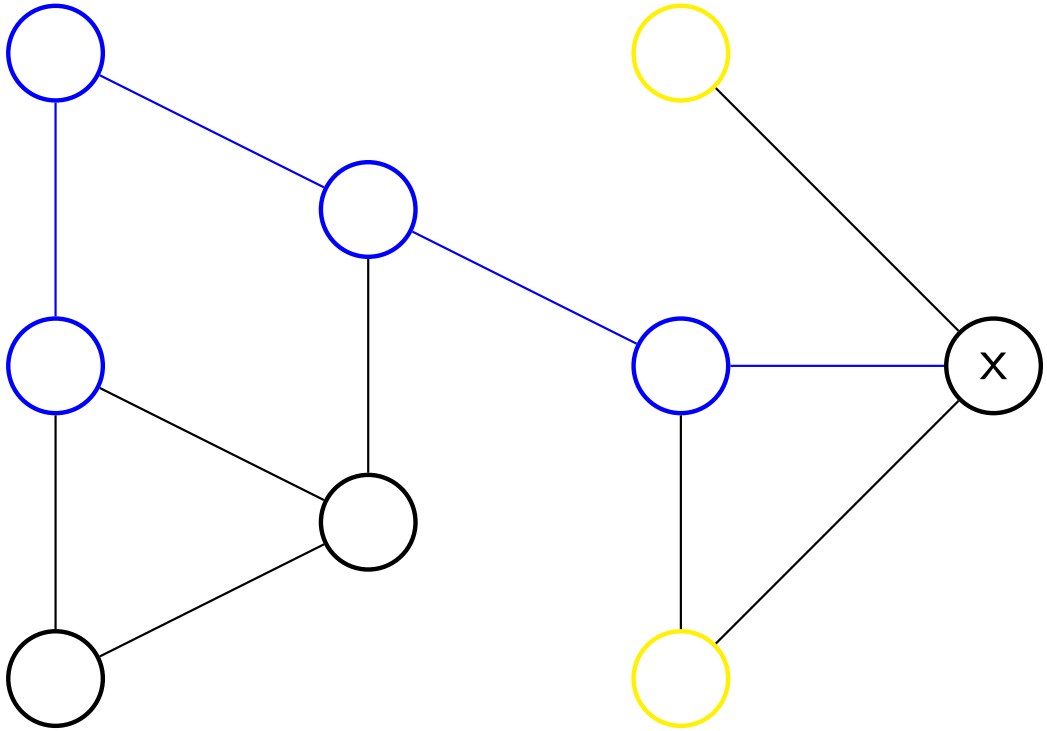
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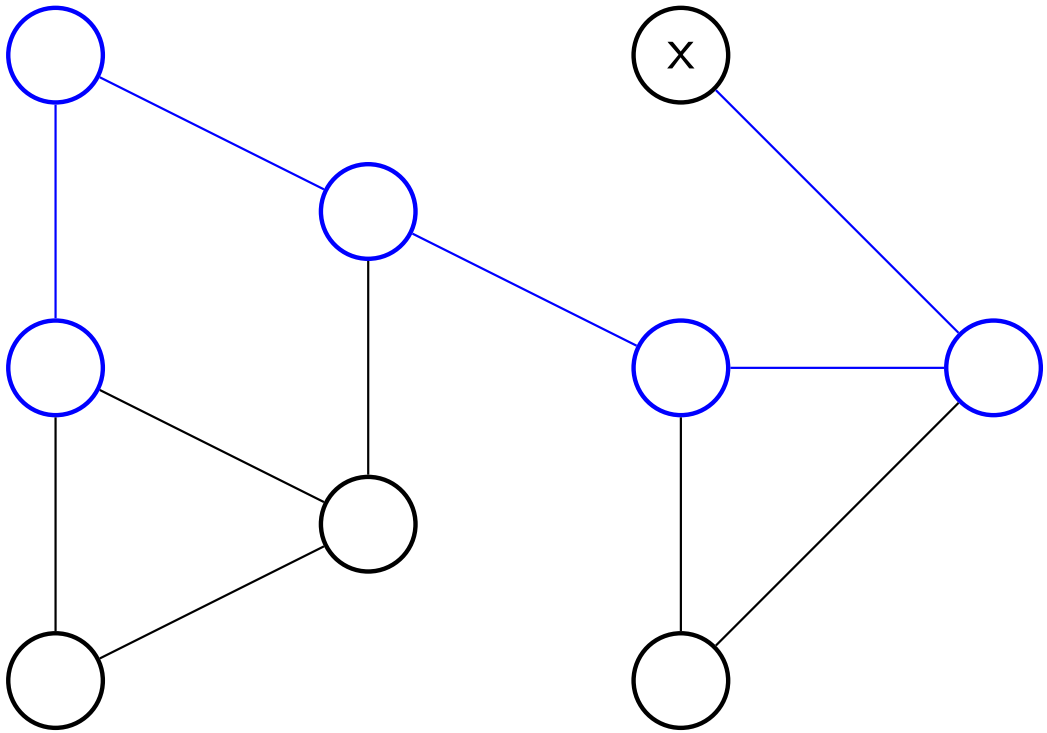
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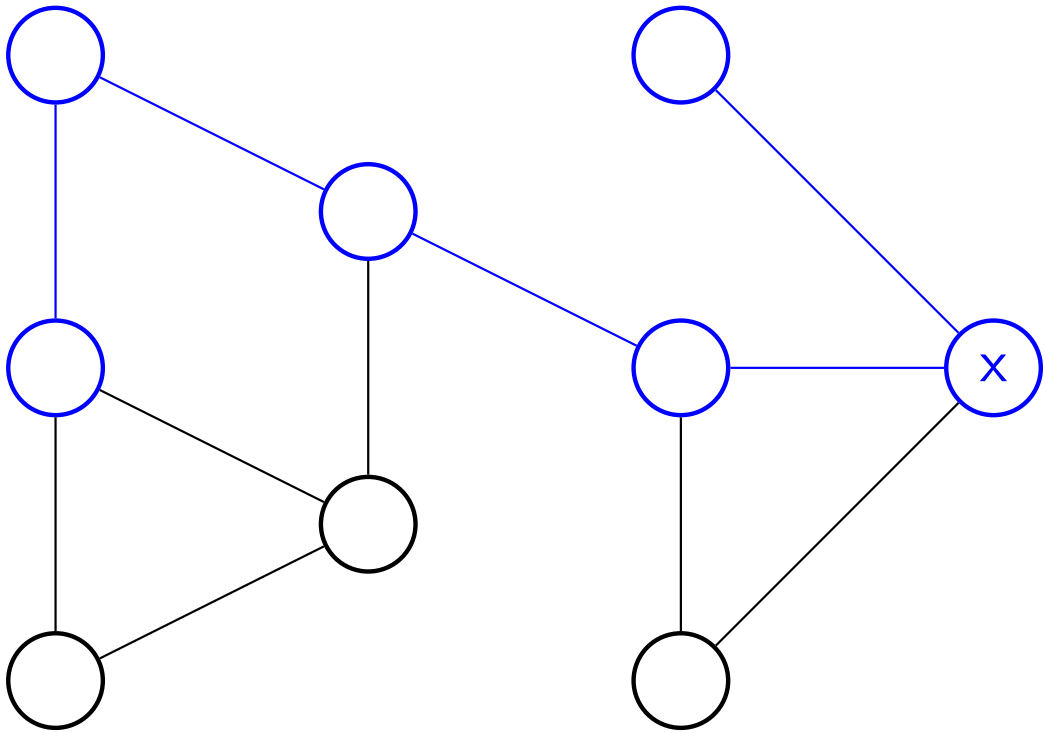
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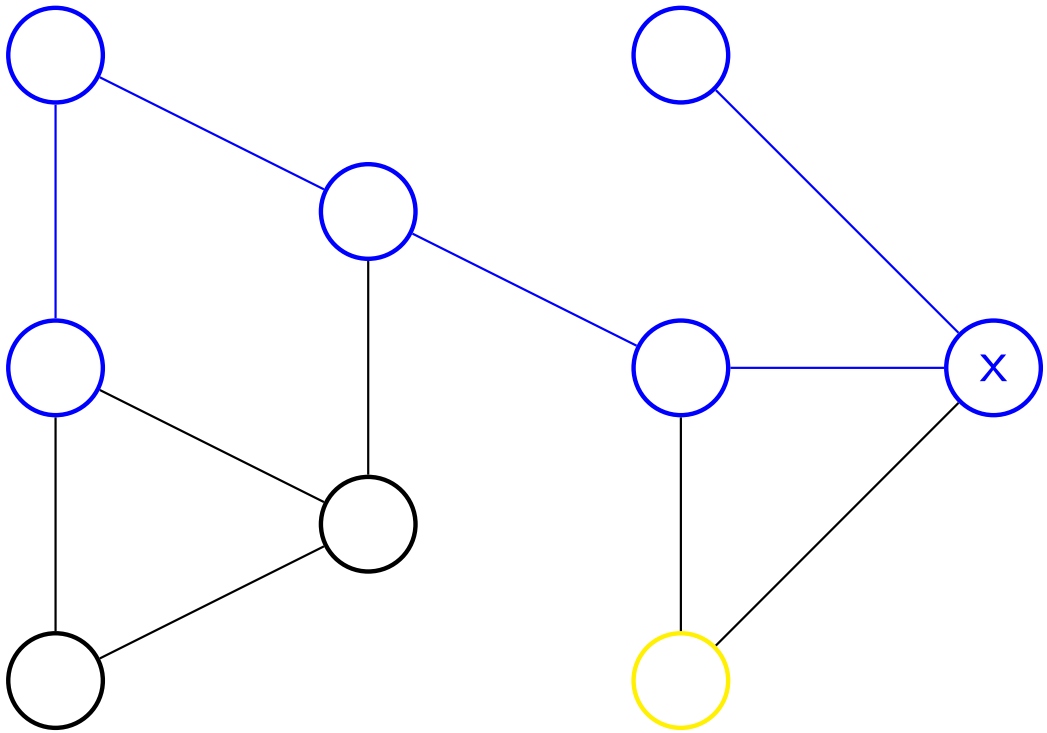
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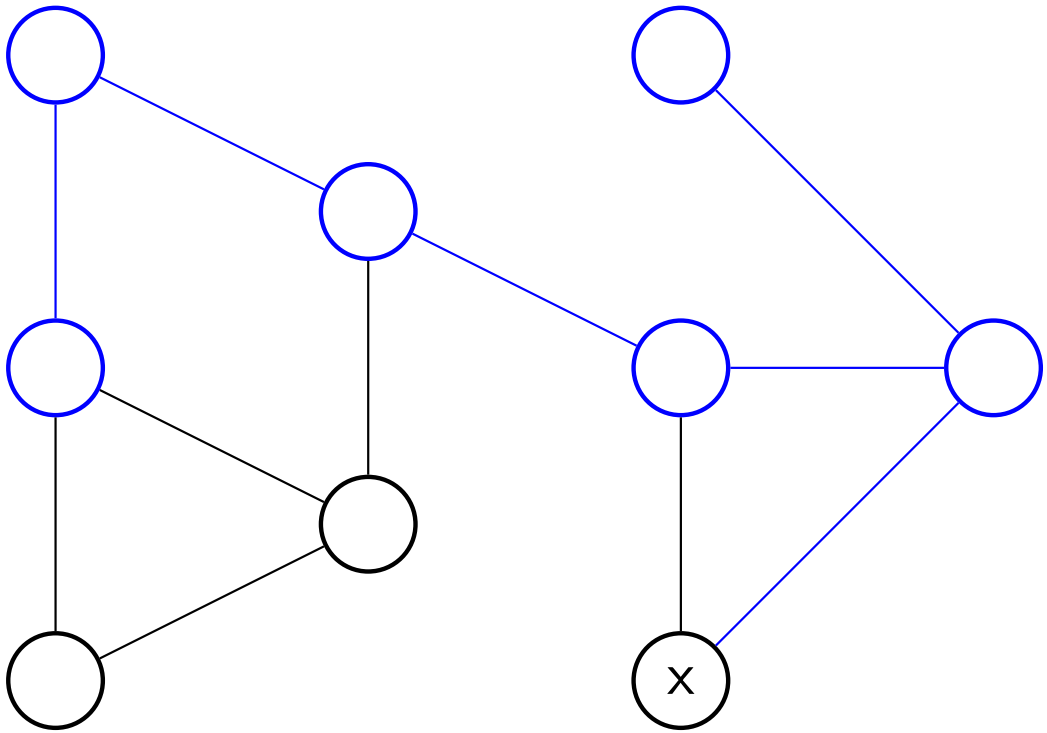
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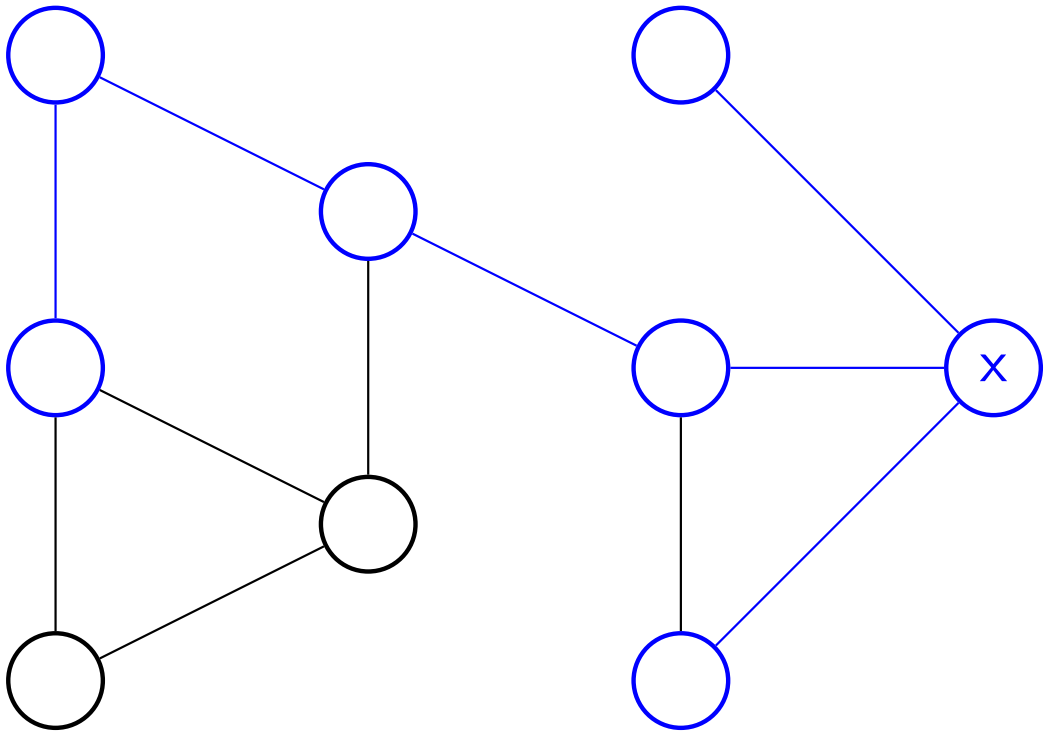
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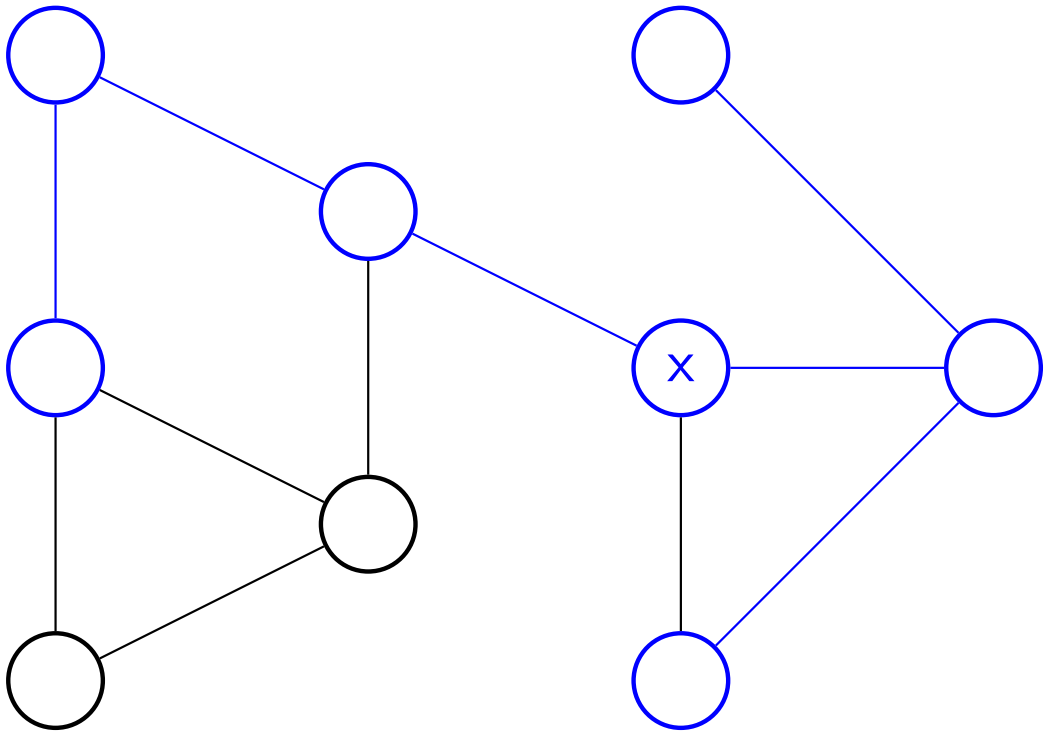
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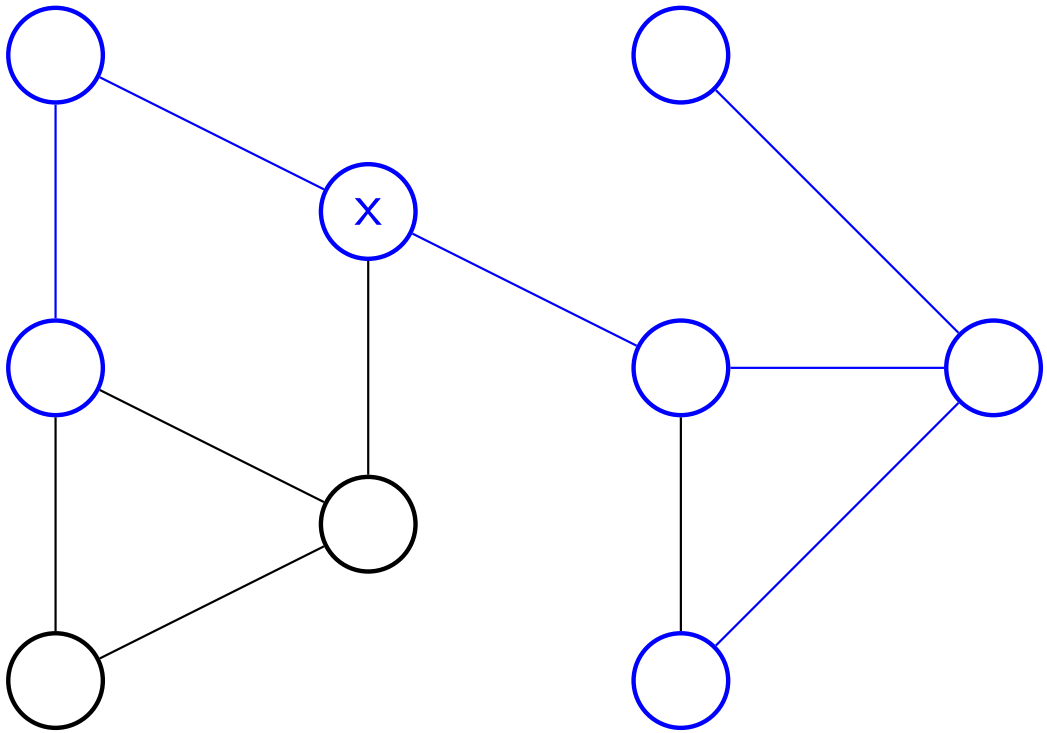
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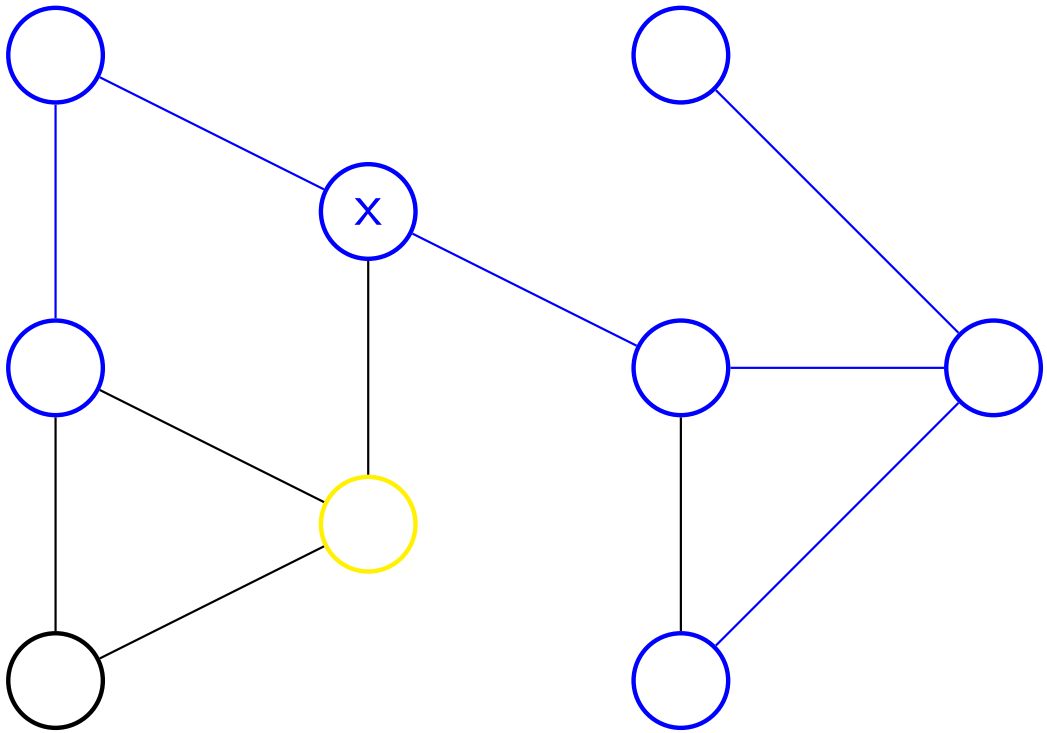
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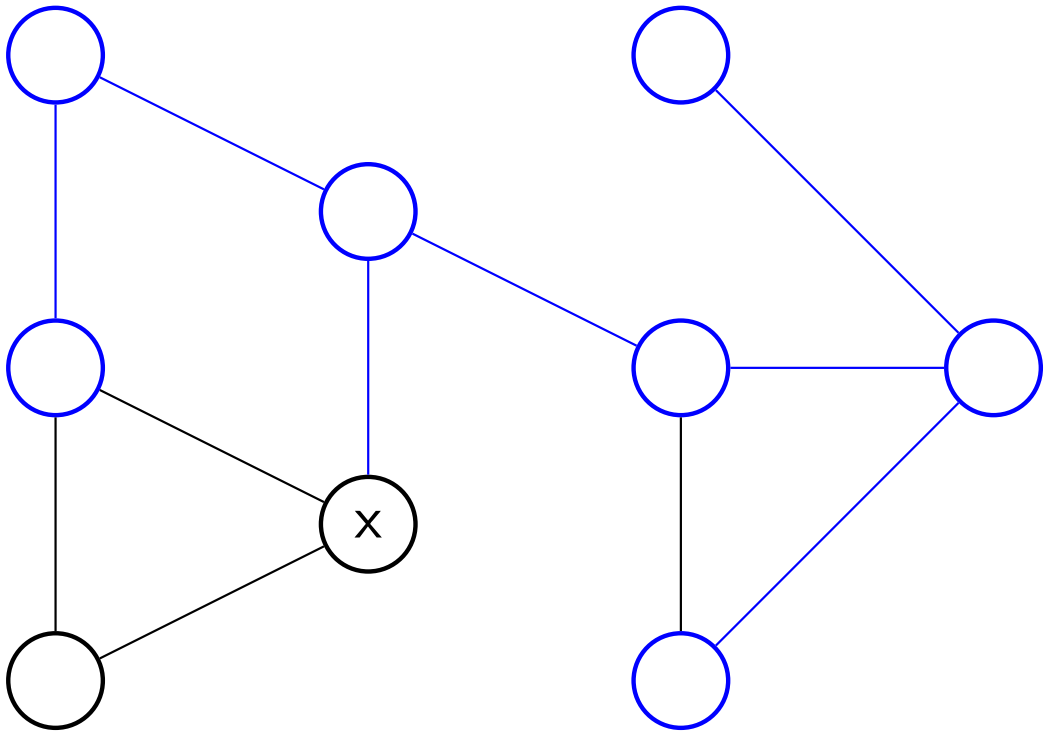
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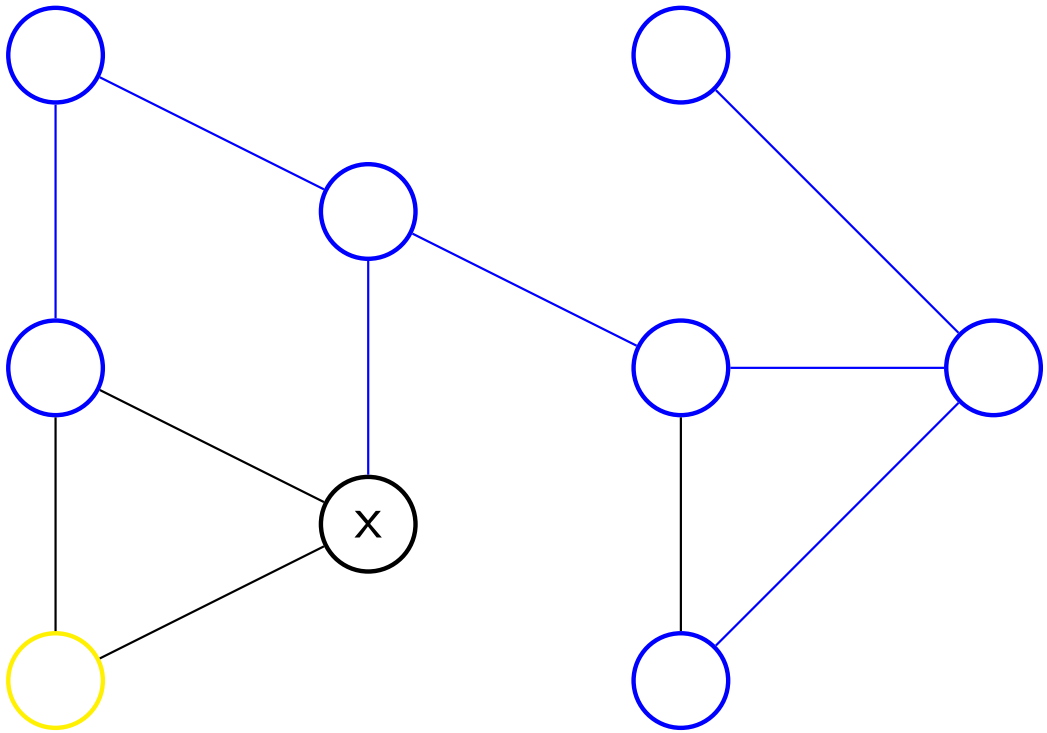
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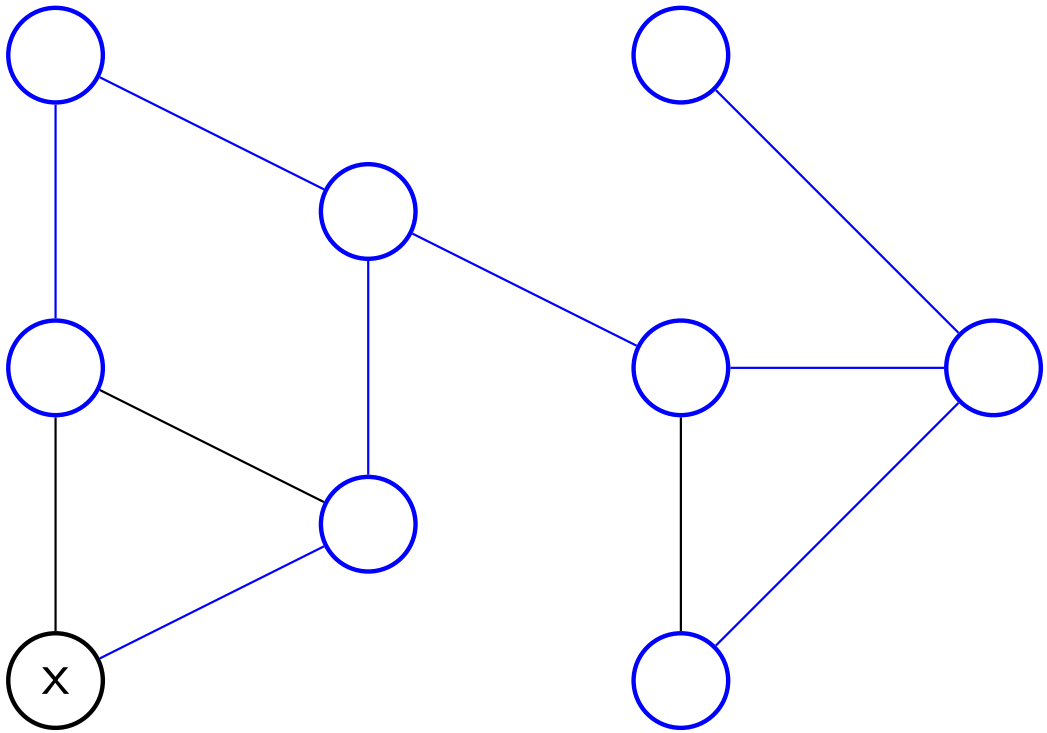
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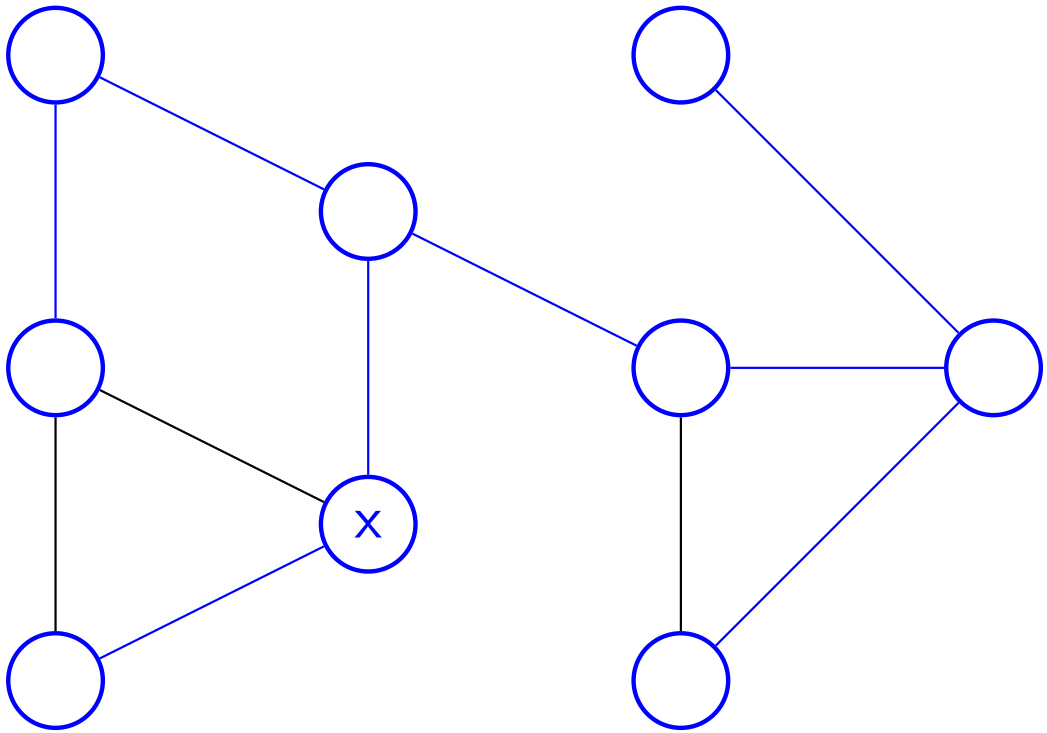
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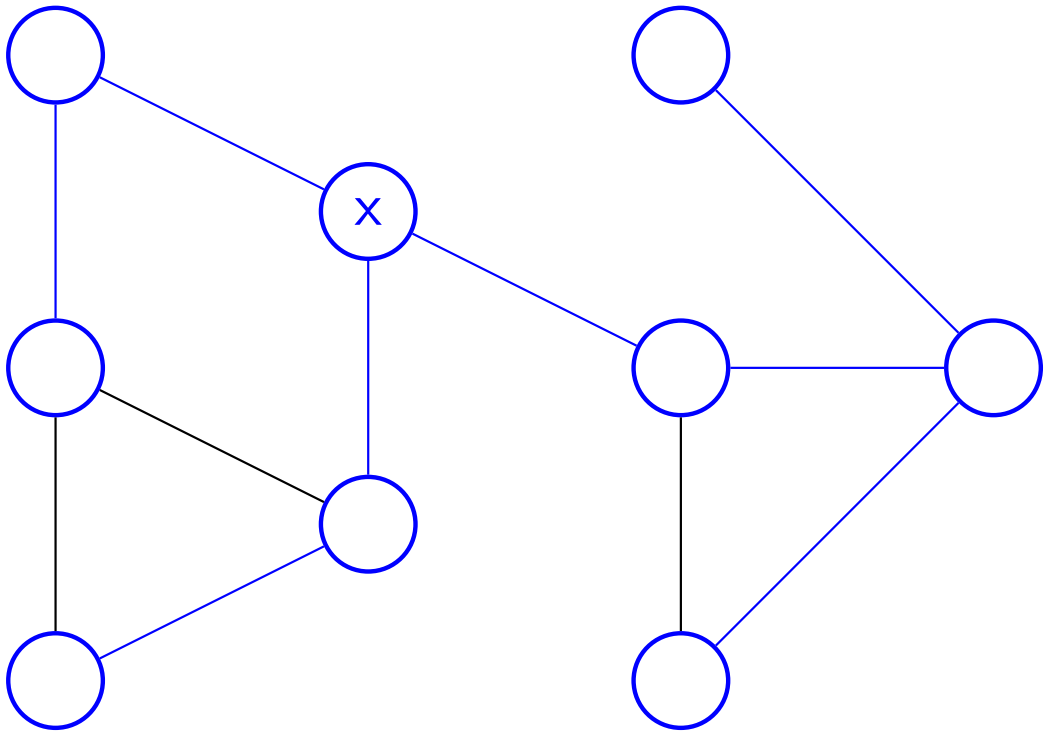
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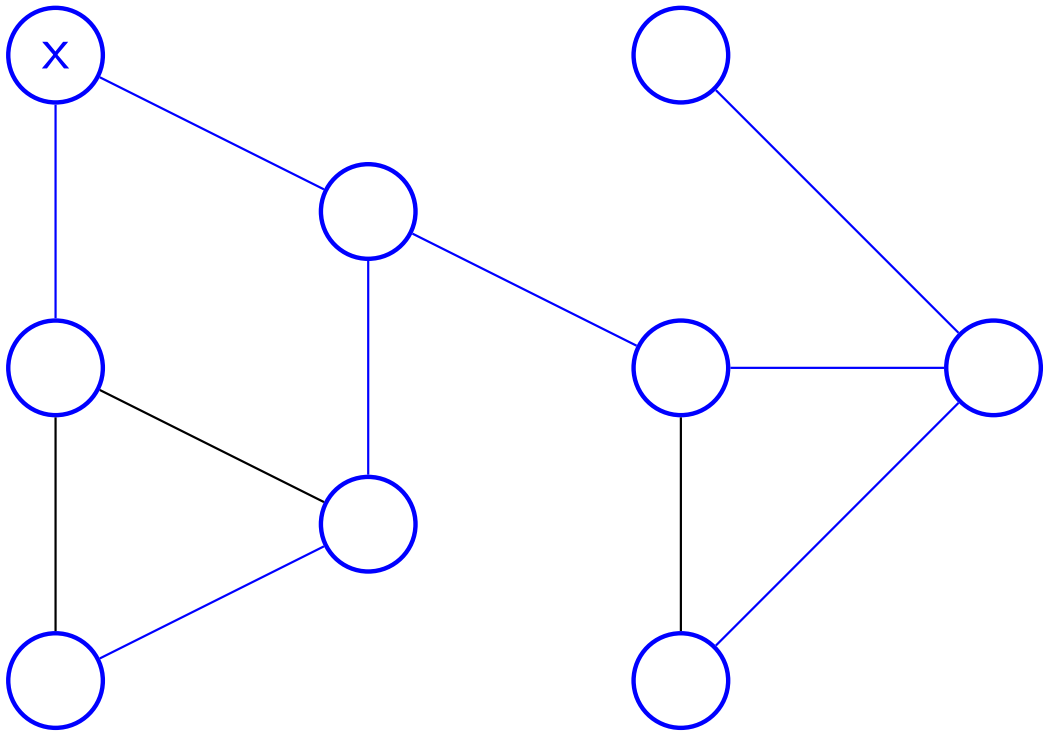
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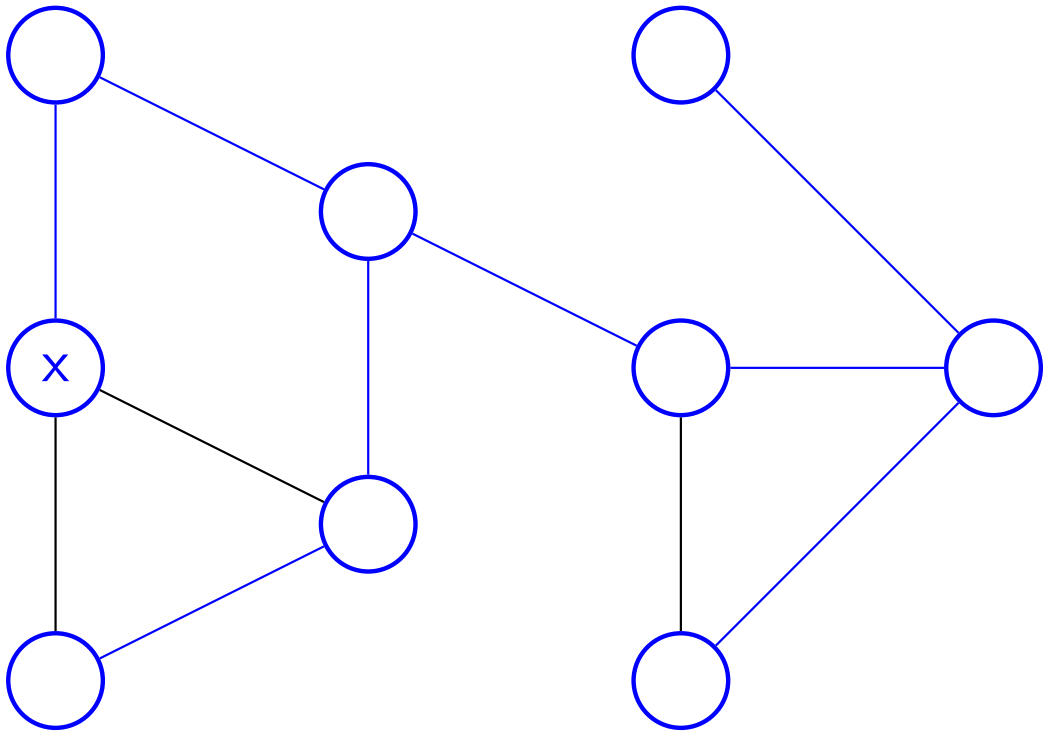
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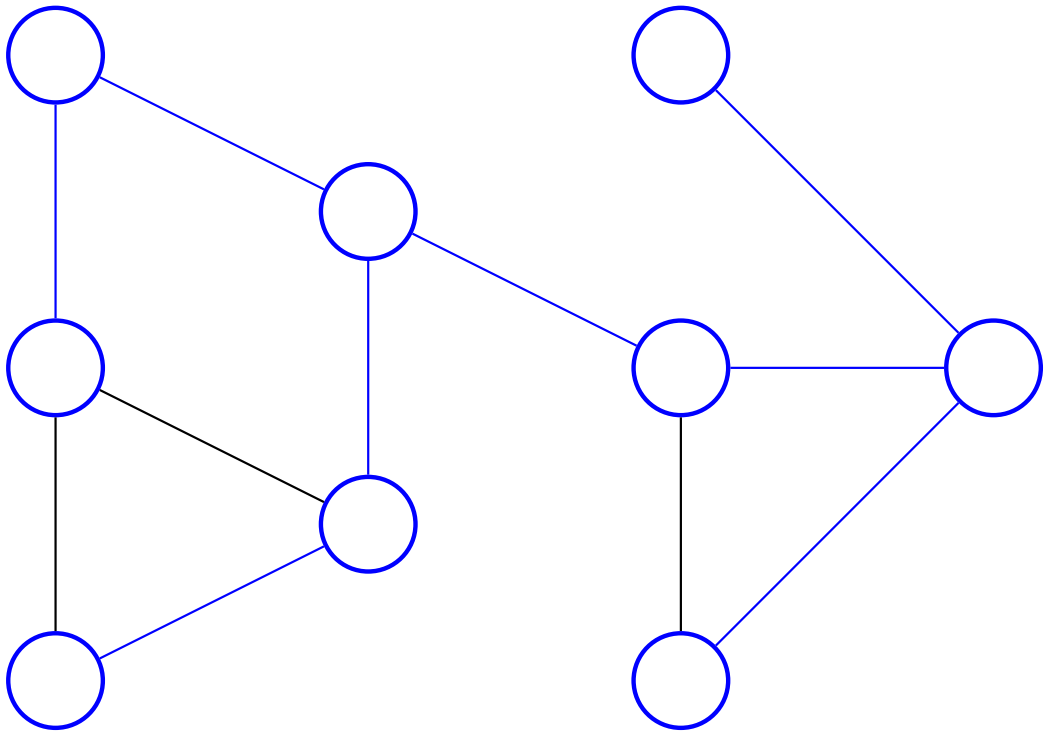


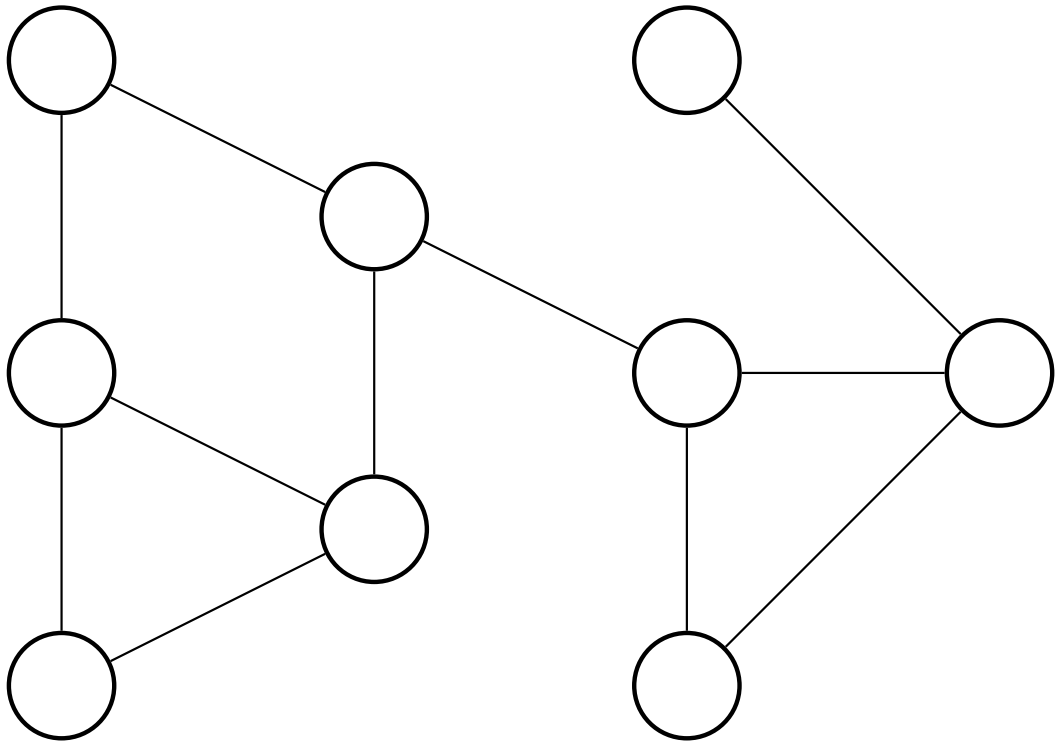
- Þegar kemur að því að útfæra dýptarleit er oftast notast við endurkvæmni.
- Endurkvæmnin sér sjálkrafa um að „fara til baka”.

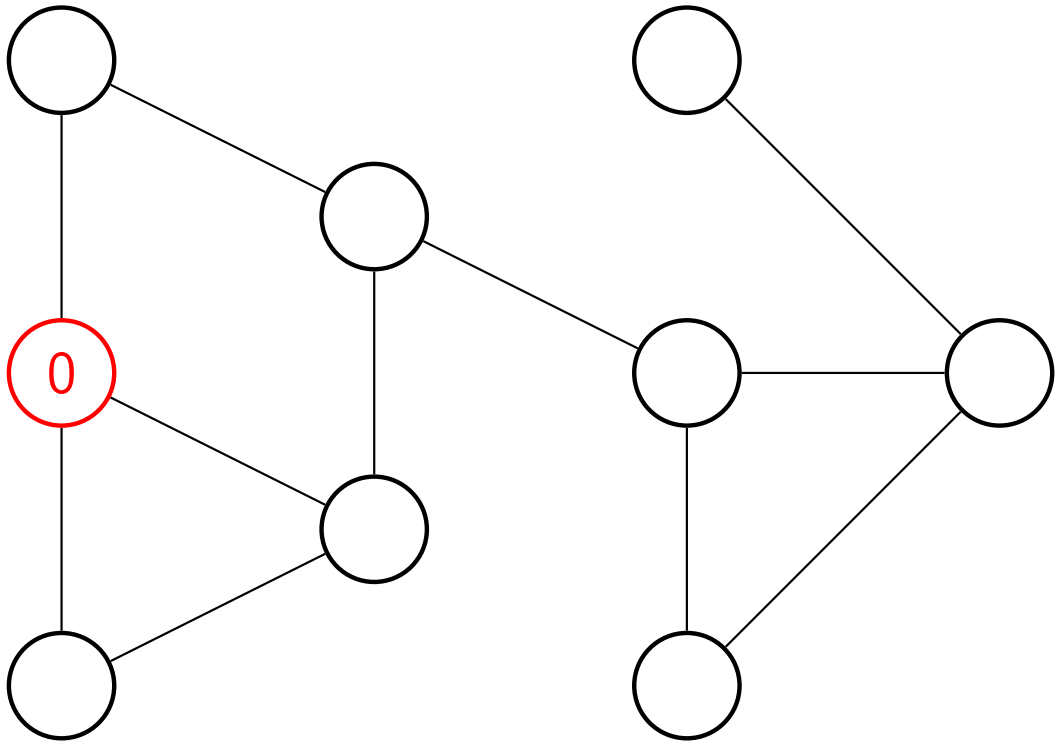
```
 6 vi v;
 7 void dfs(vvi& g, int x)
 8 {
 9     int i;
10     printf("Vid erum i nodu %d\n", x + 1);
11     v[x] = 1;
12     for (i = 0; i < g[x].size(); i++) if (v[g[x][i]] == 0)
13         dfs(g, g[x][i]);
14 }
```

- Eftir kall á `dfs(g, 0)` segir `v[j]` okkur hvort til sé vegur frá hnúti 0 til hnúts j.

- Tökum eftir að leitin heimsækir hvern hnút í mesta lagi einu sinni og ferðast eftir hverjum legg í mesta lagi tvisvar (einu sinni í stefndu neti).
- Svo tímaflækjan er $\mathcal{O}(E + V)$.
- Við getum í rauninni ekki beðið um betri tímaflækju.
- Við munum alltaf þurfa að skoða alla hnúta og ef við skoðum ekki alla leggi þá erum við að hunsa uppbyggingu netsins.

- Við byrjum á að merkja alla hnúta sem „ósnerta", nema við merkjum einn hnút sem „séðan".
- Sá hnútur er kallaður *upphafshnúturinn*.
- Við endurtökum svo eftirfarandi þar til engir „séðir" hnútar eru eftir:
  - Veljum þann „séða" hnút sem við sáum fyrst.
  - Merkjum alla „ósnerta" nágranna hans sem „séða".
  - Merkjum upprunalegu hnútinn „kláraðann".
- Tökum dæmi.
- Við munum merkja „séða" hnúta með hvenær við sáum þá.

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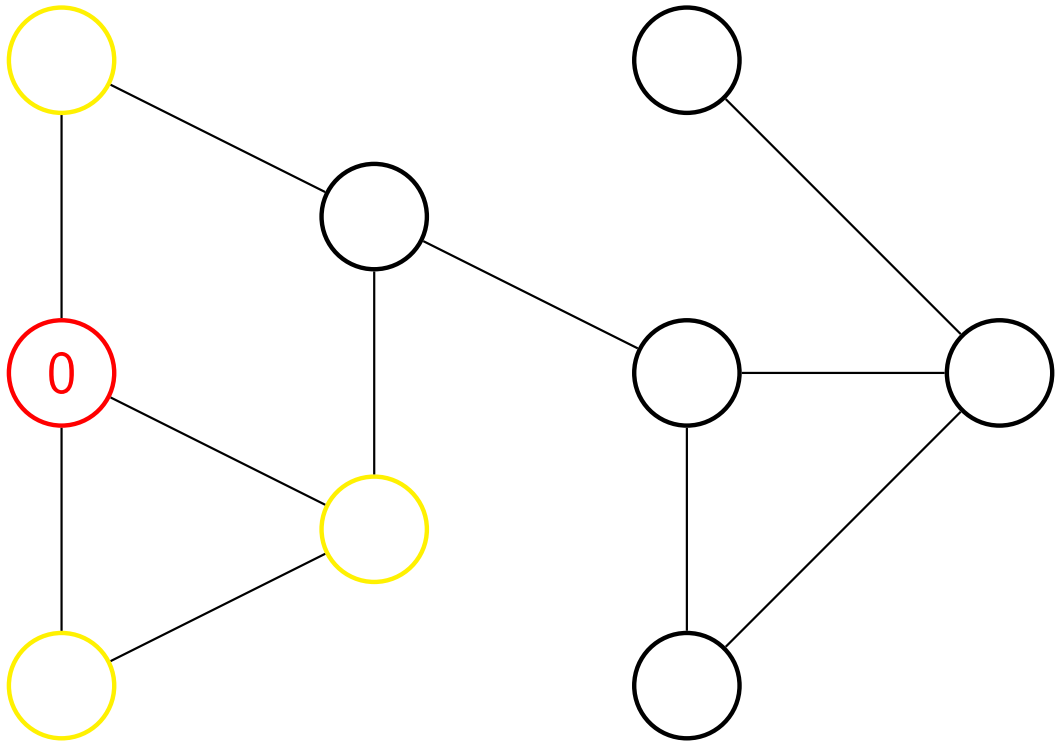
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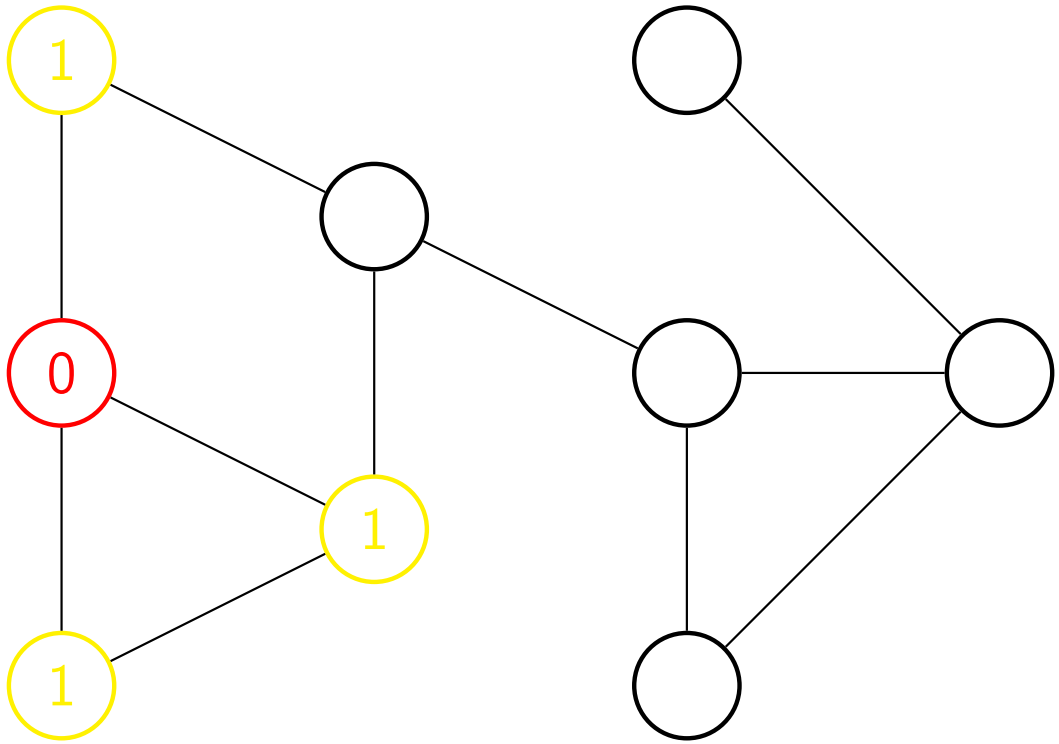
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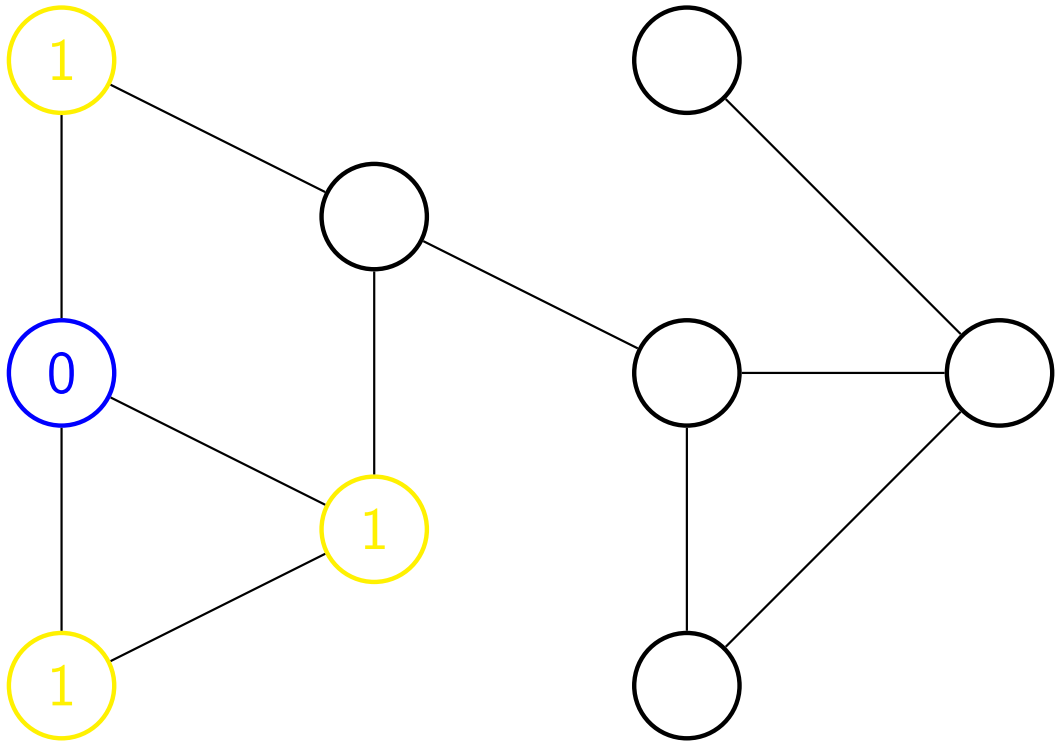
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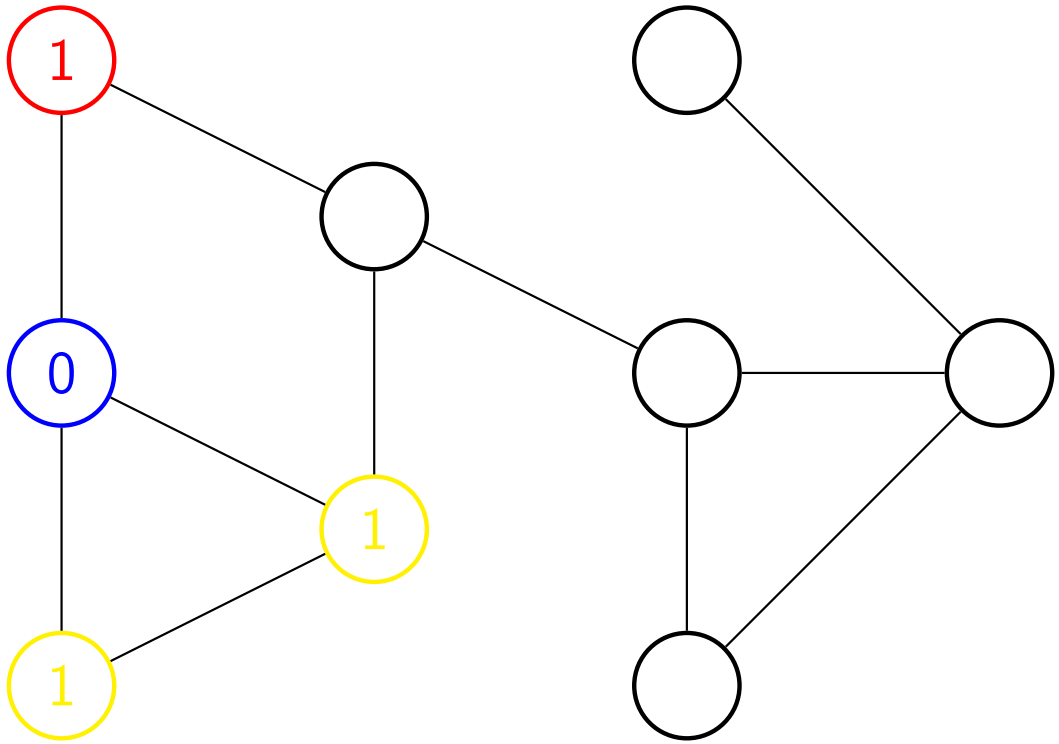
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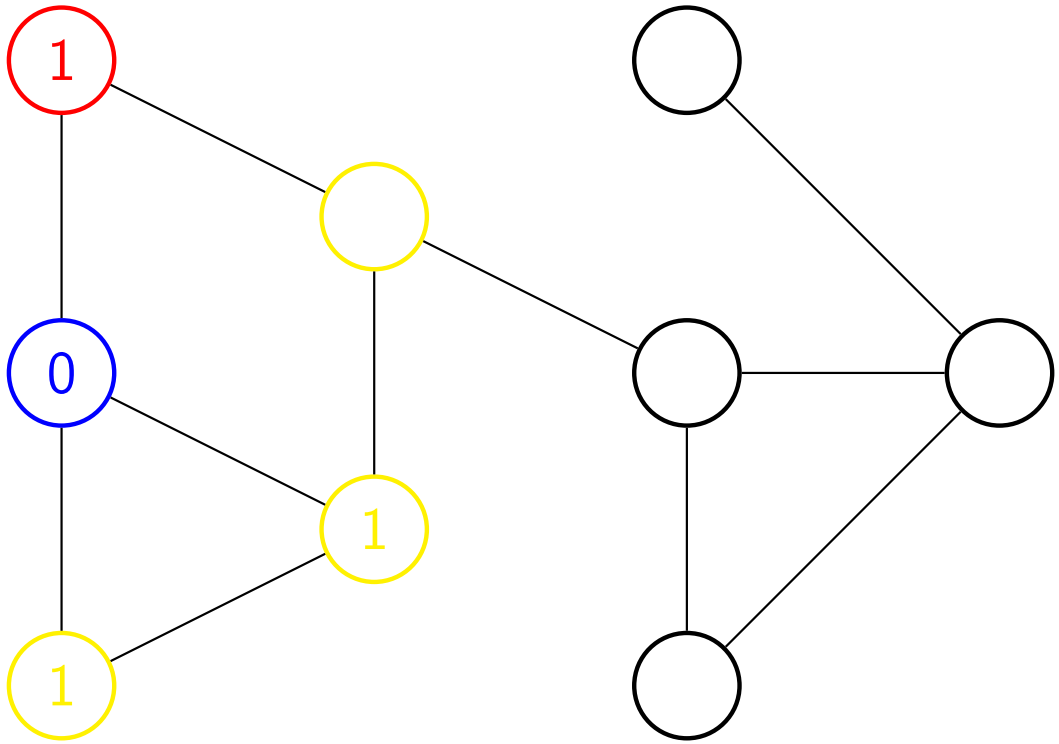
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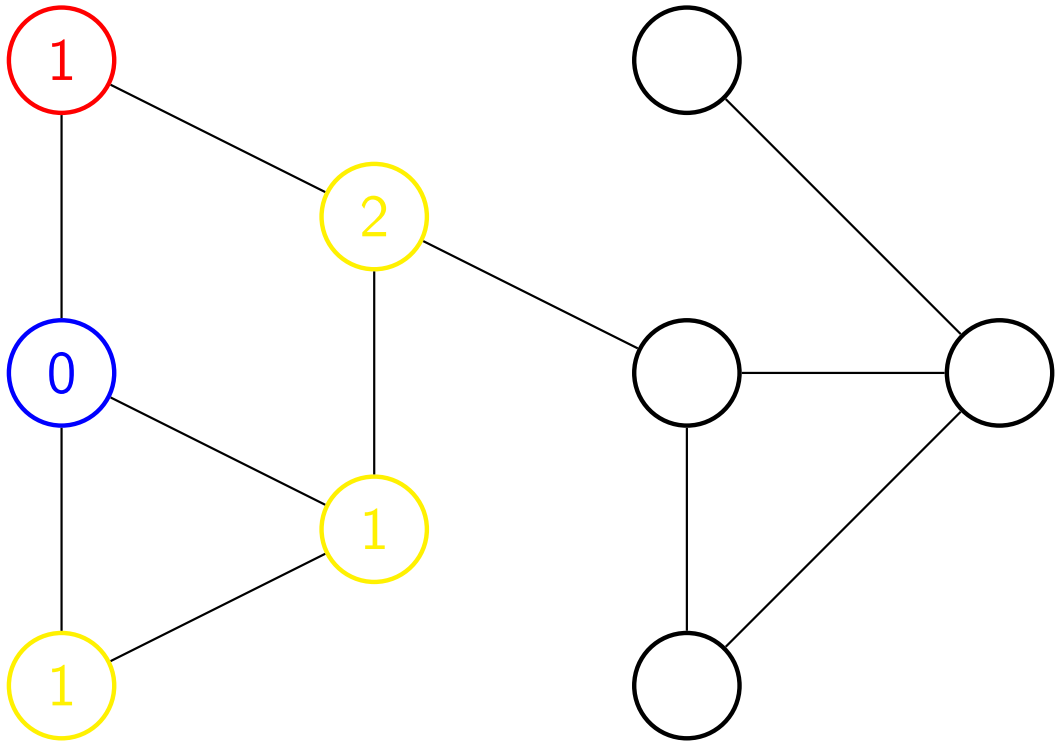
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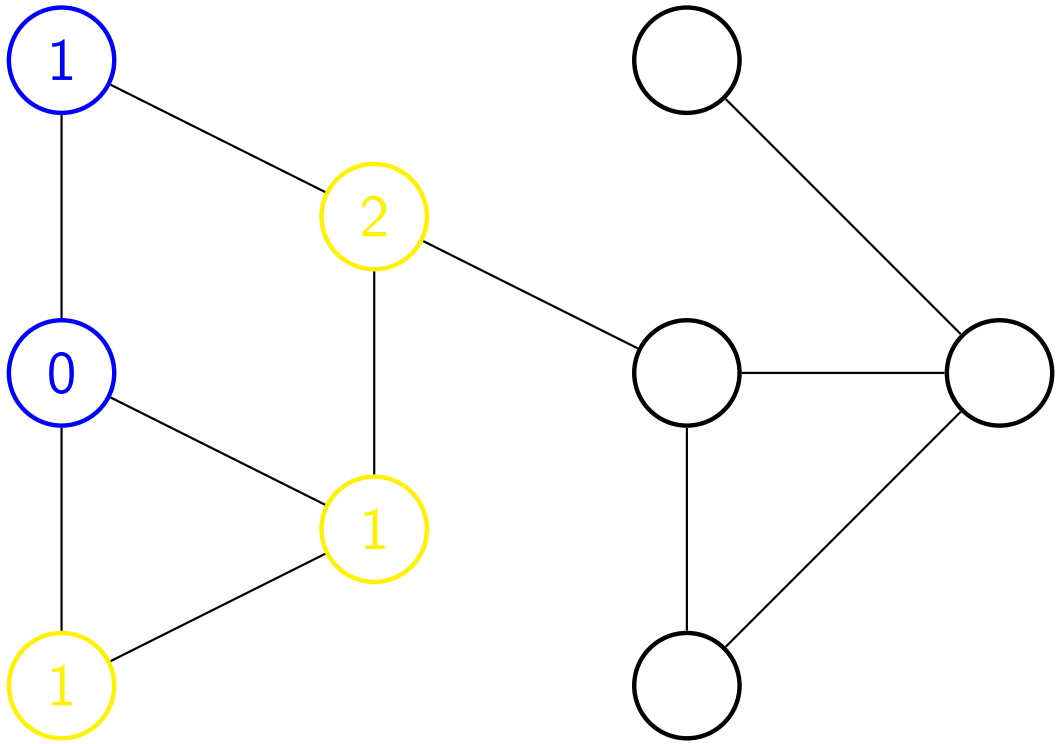
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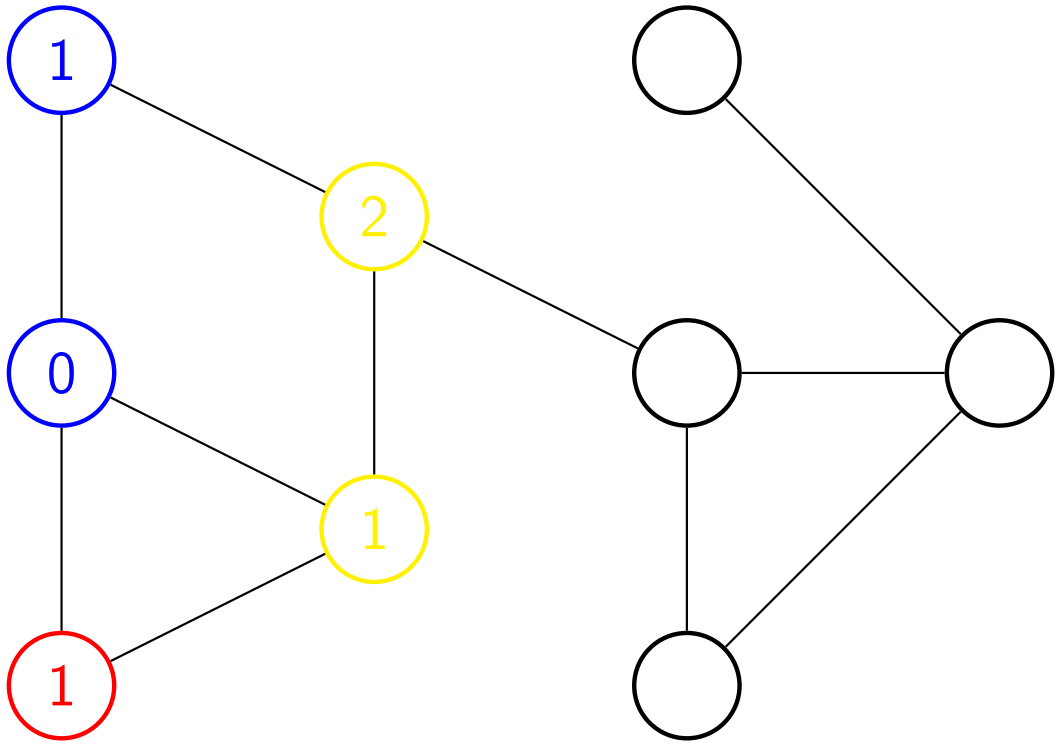
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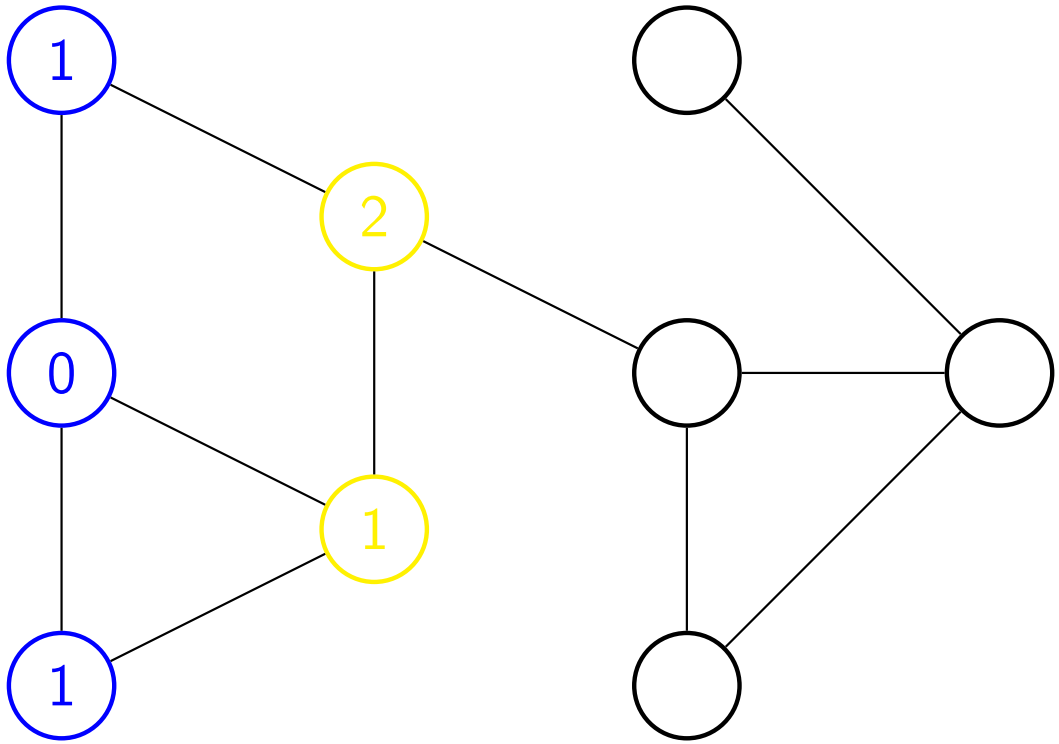
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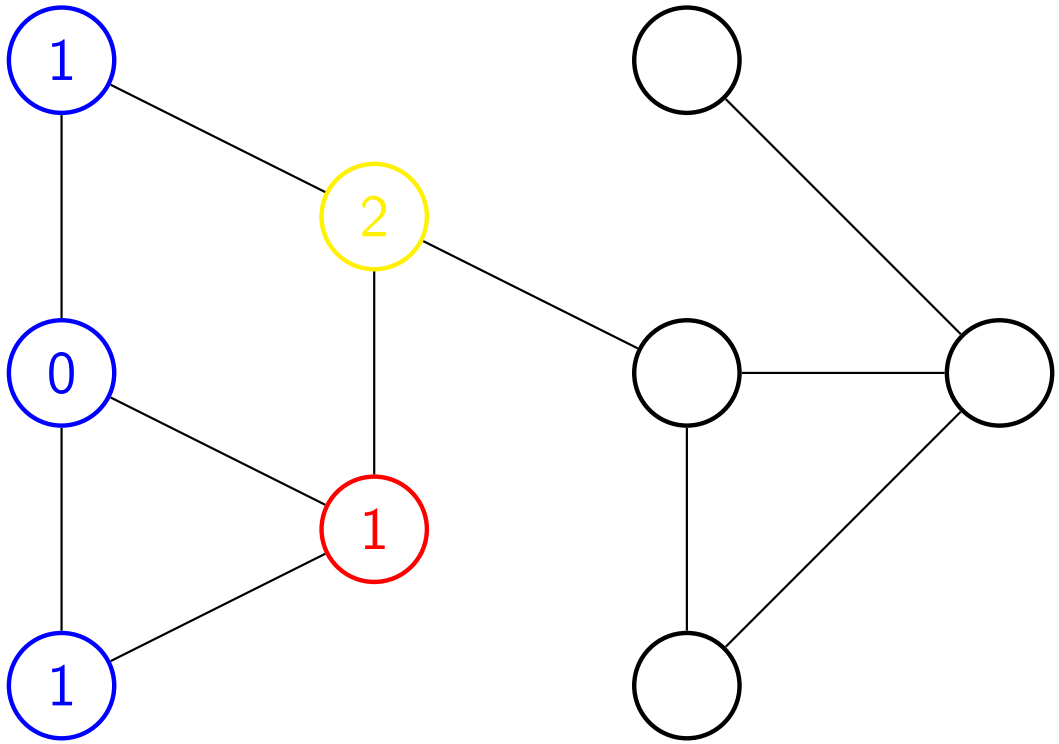
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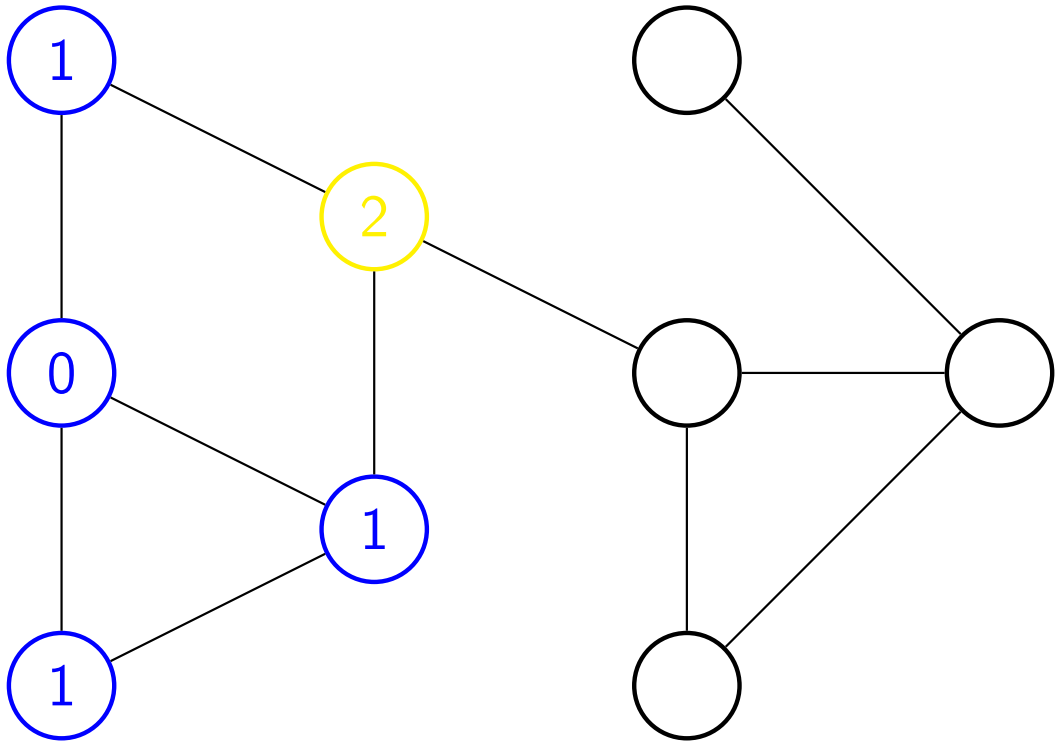
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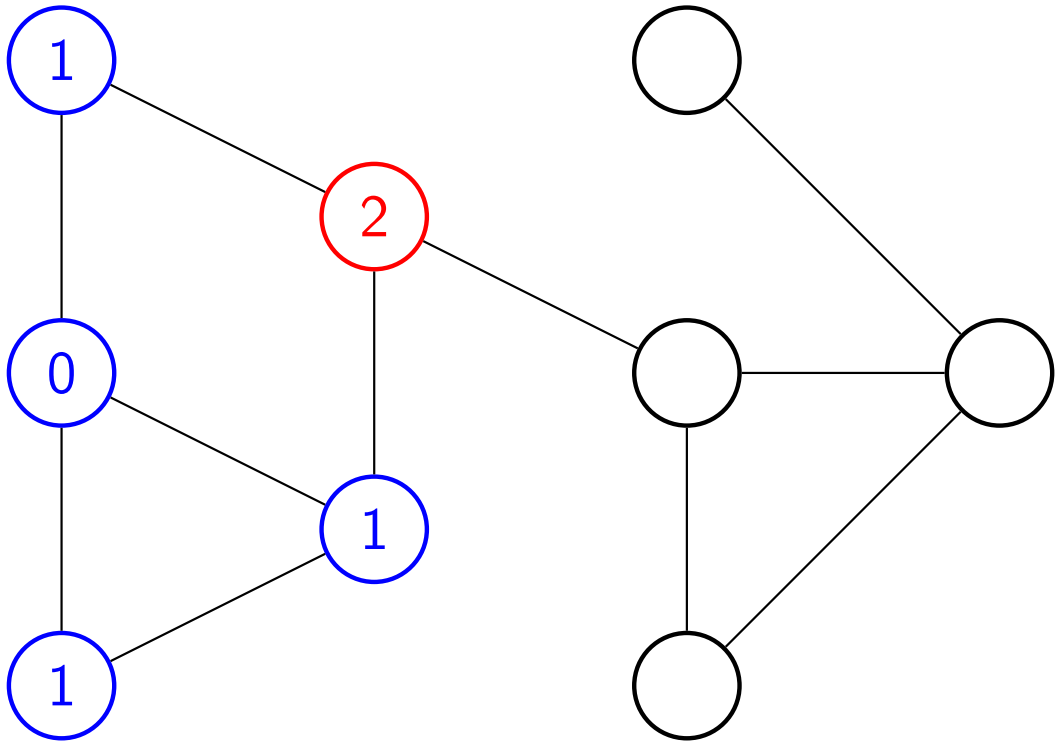
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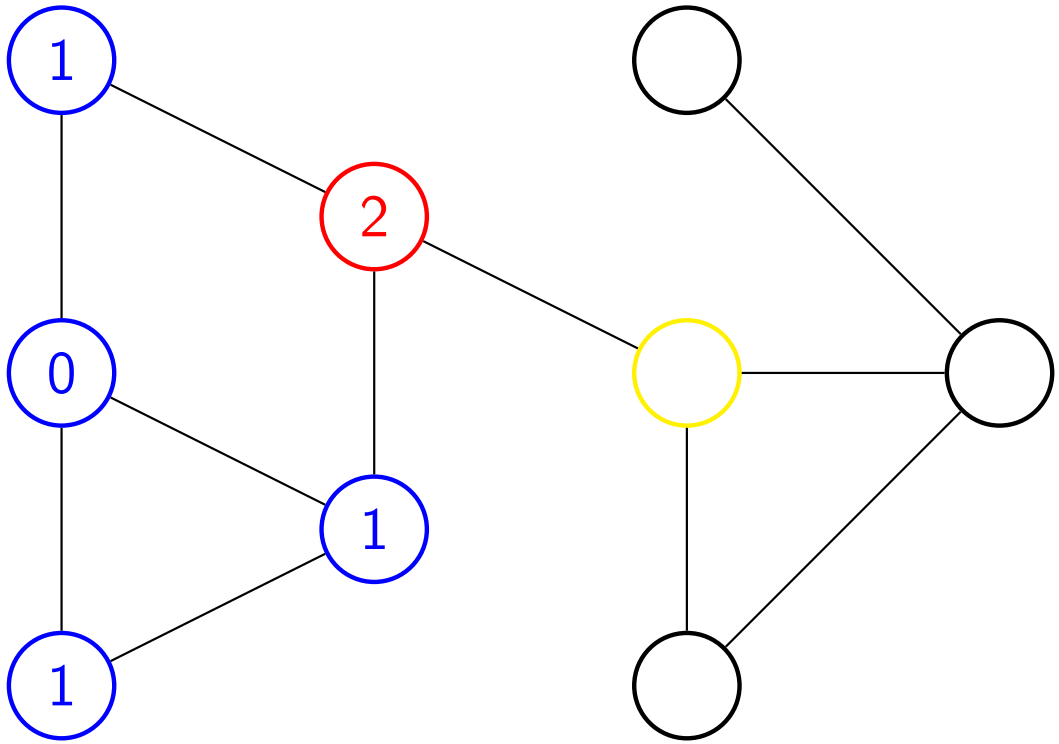
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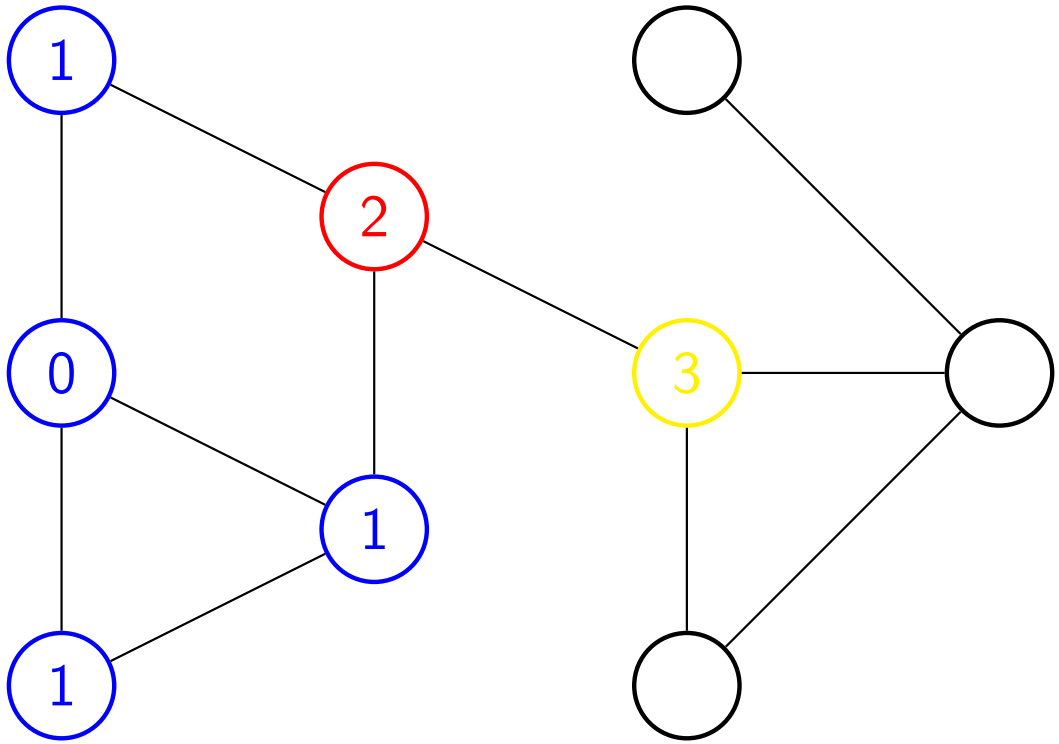
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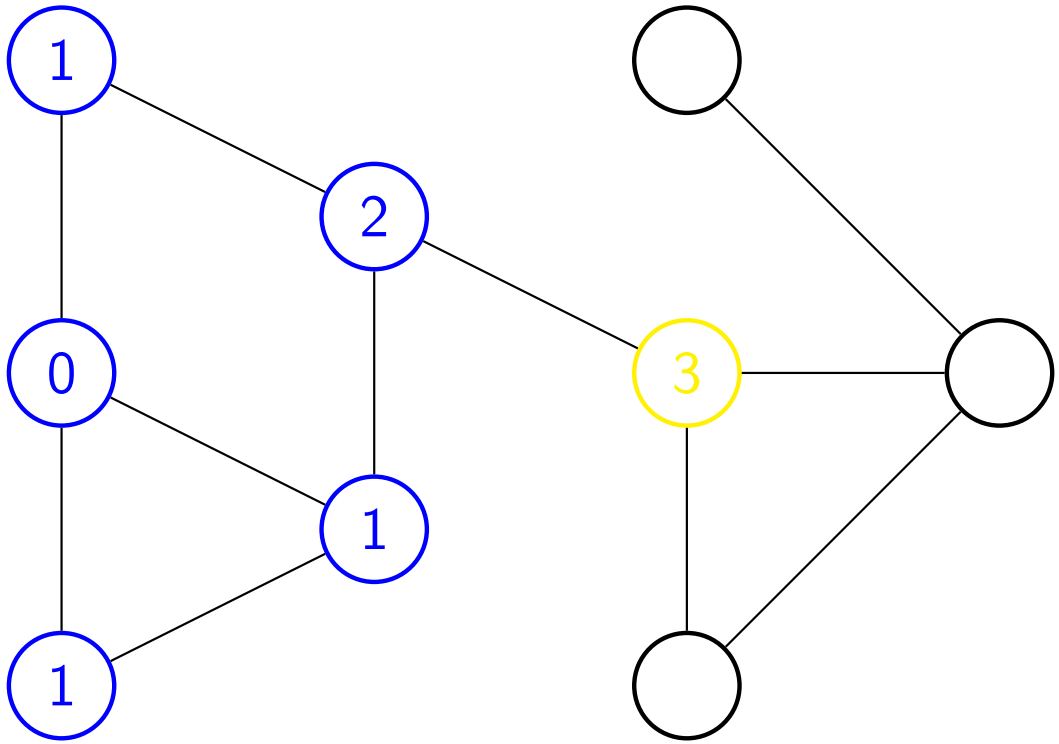
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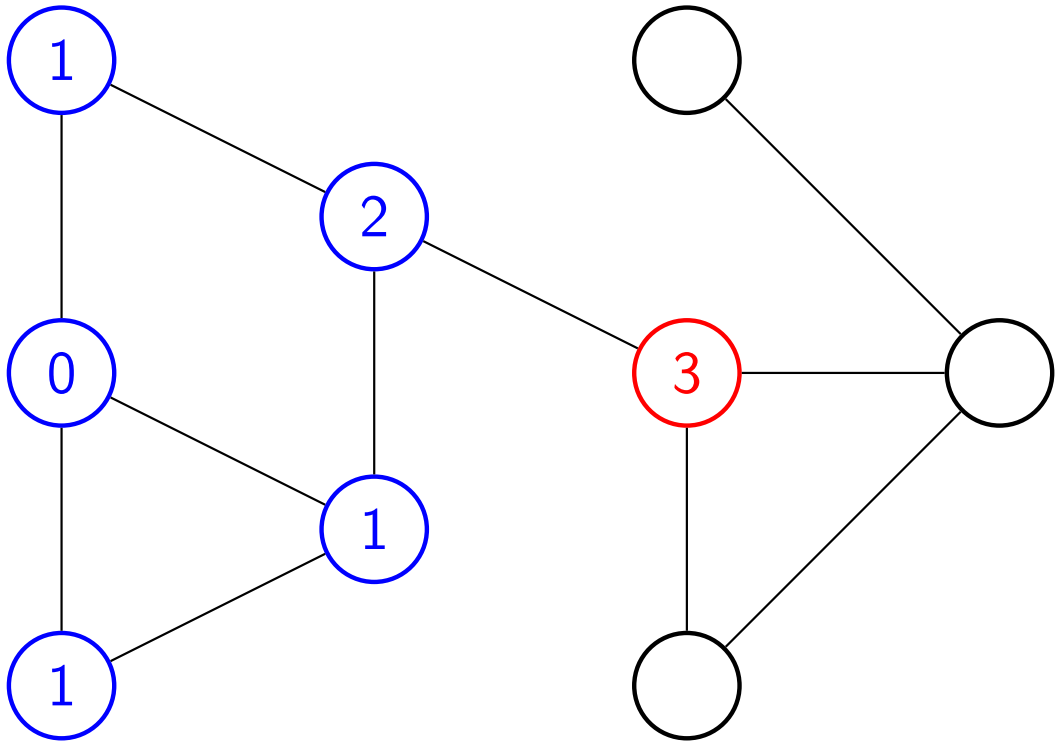
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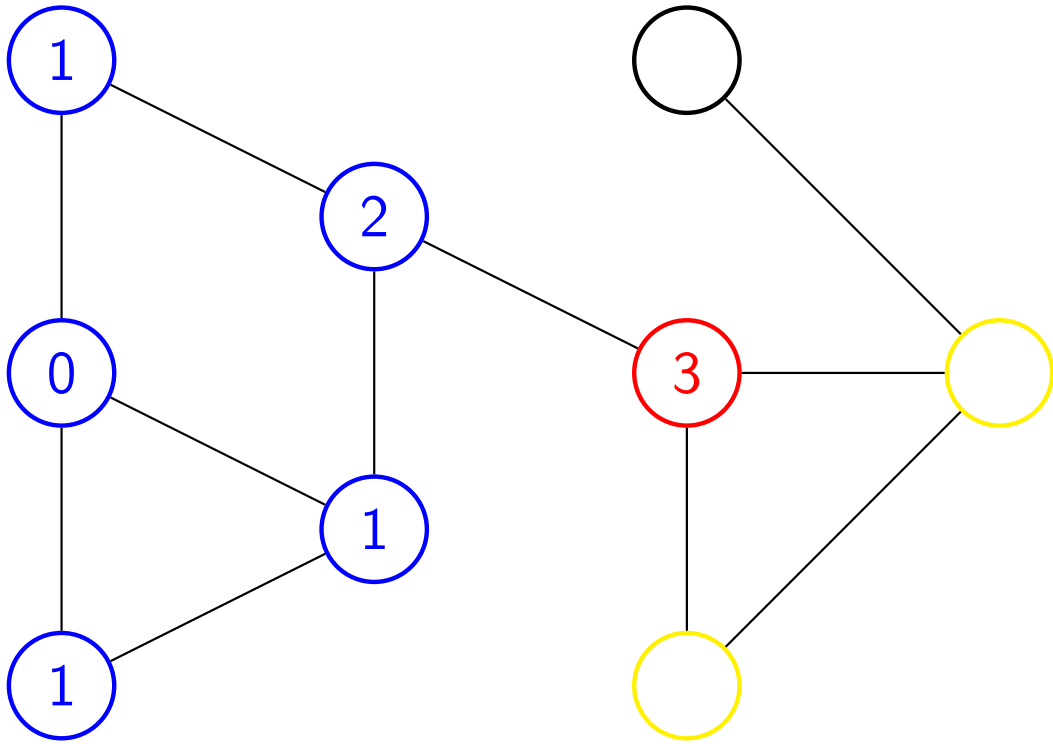
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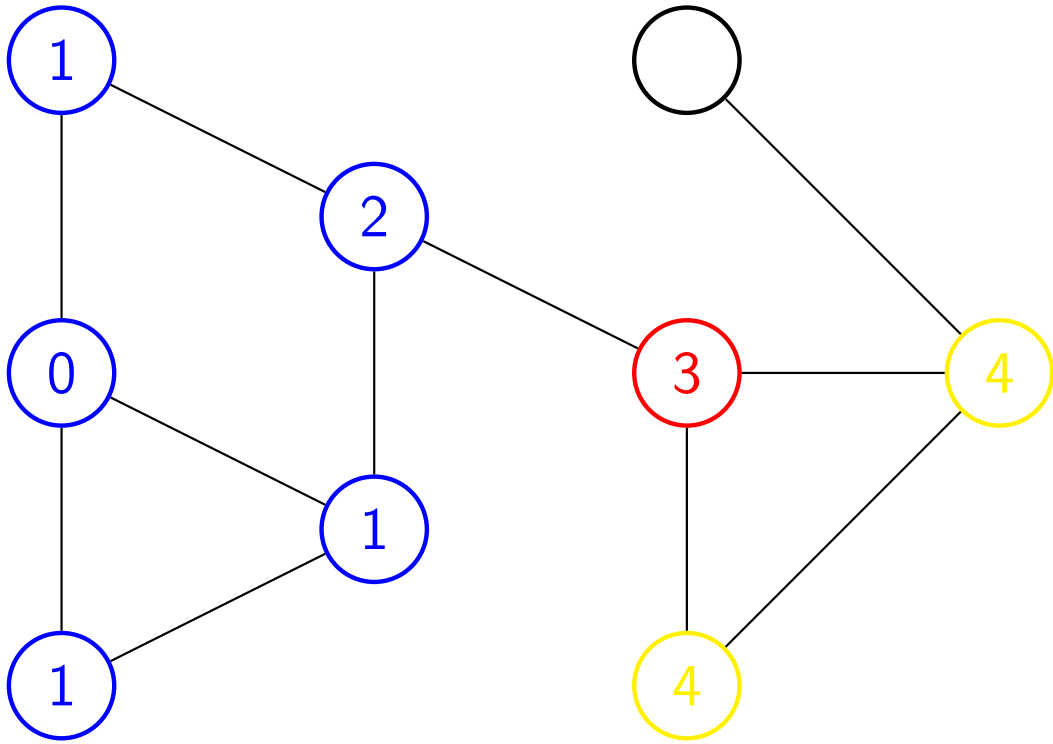
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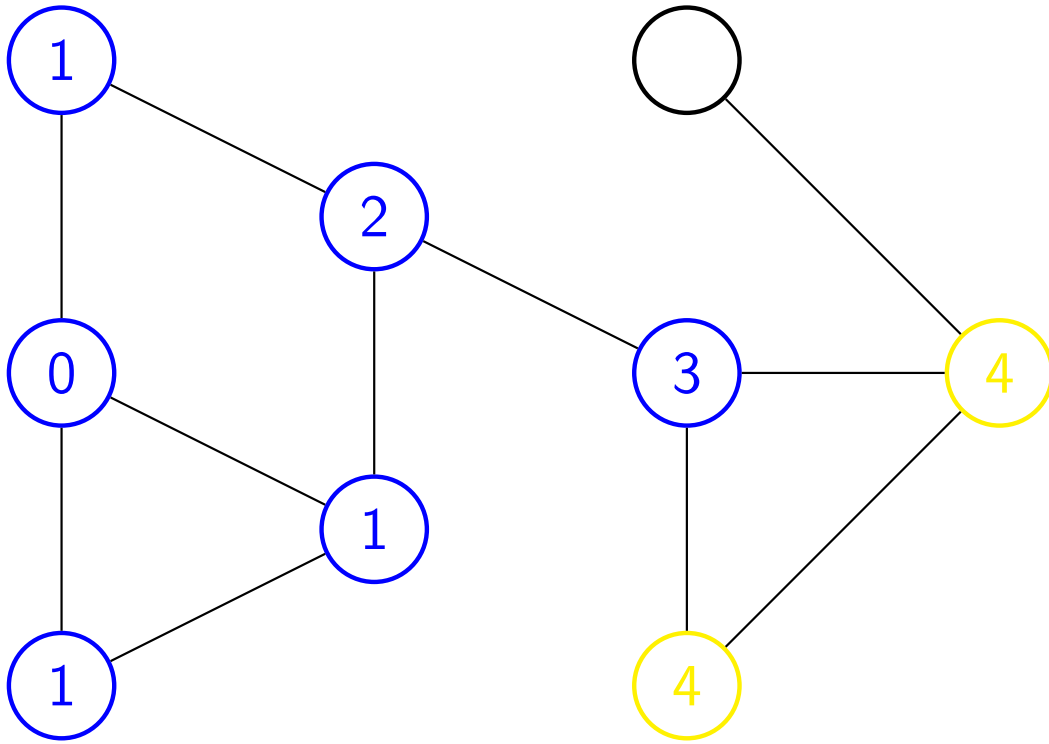
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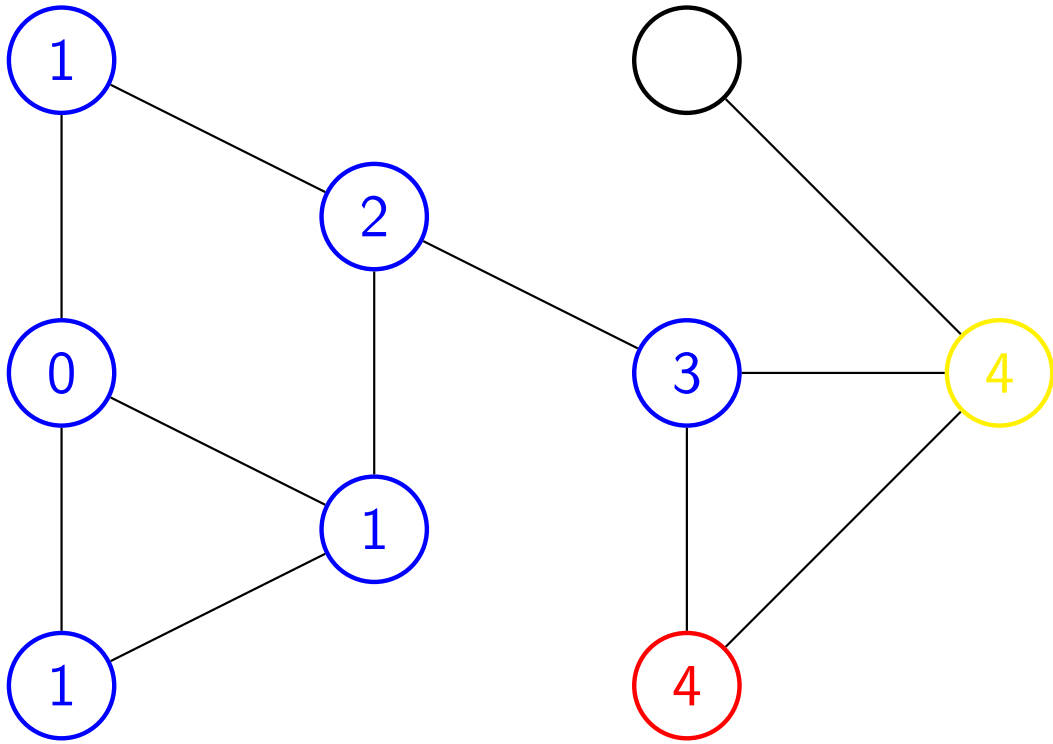
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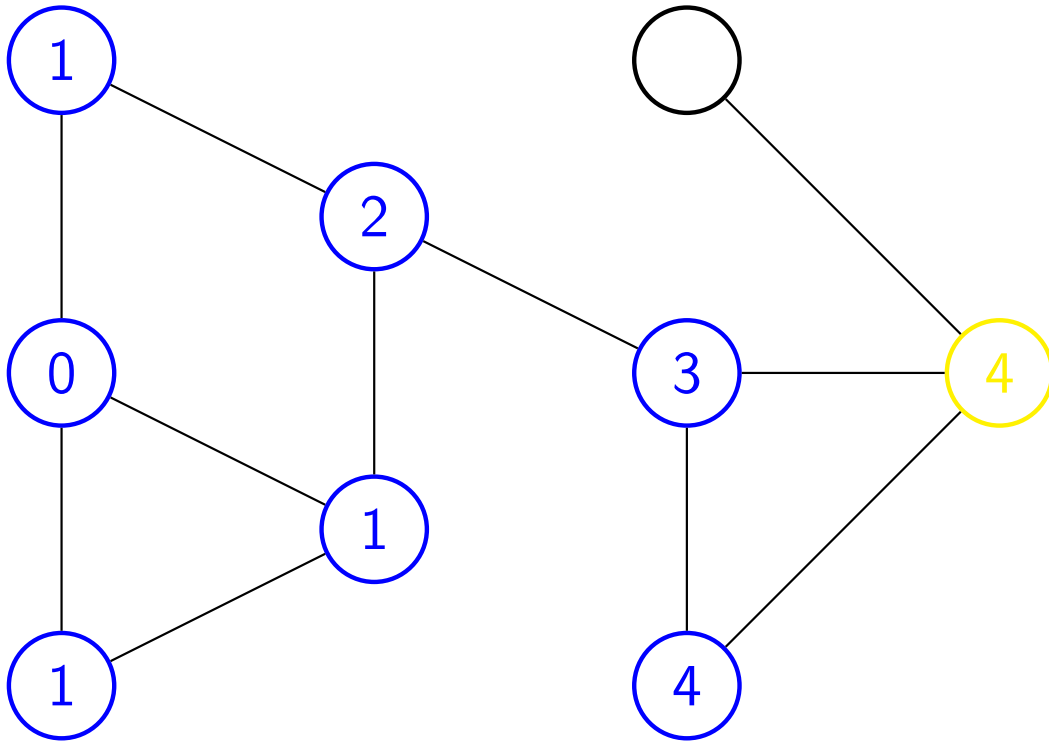
1
2
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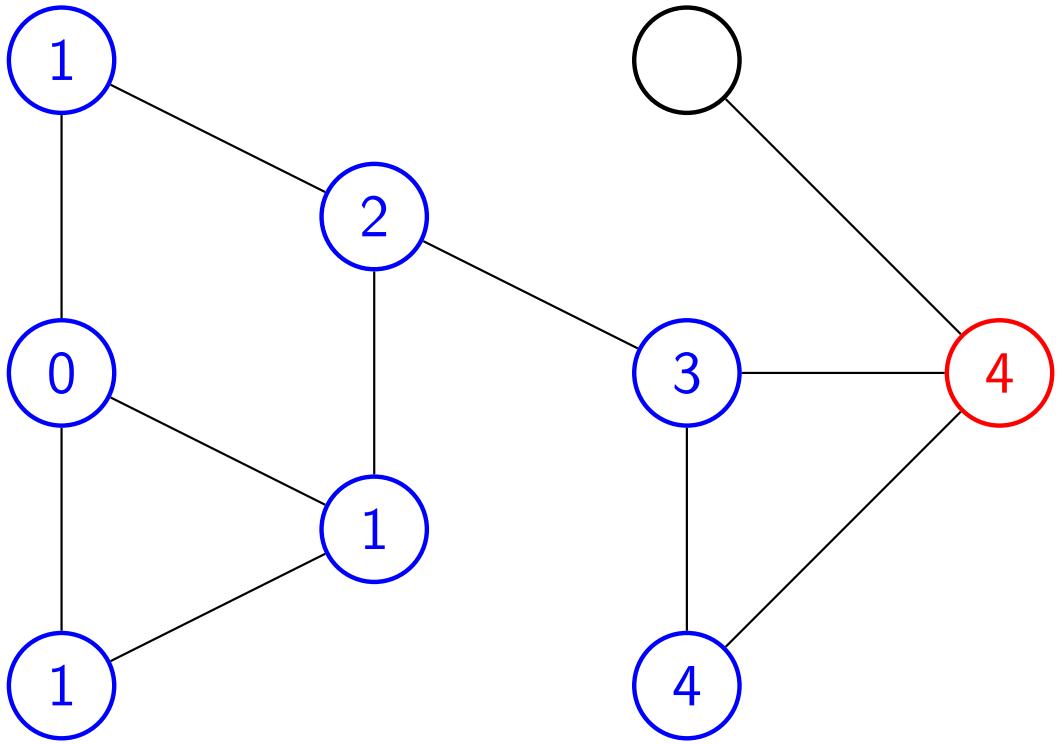
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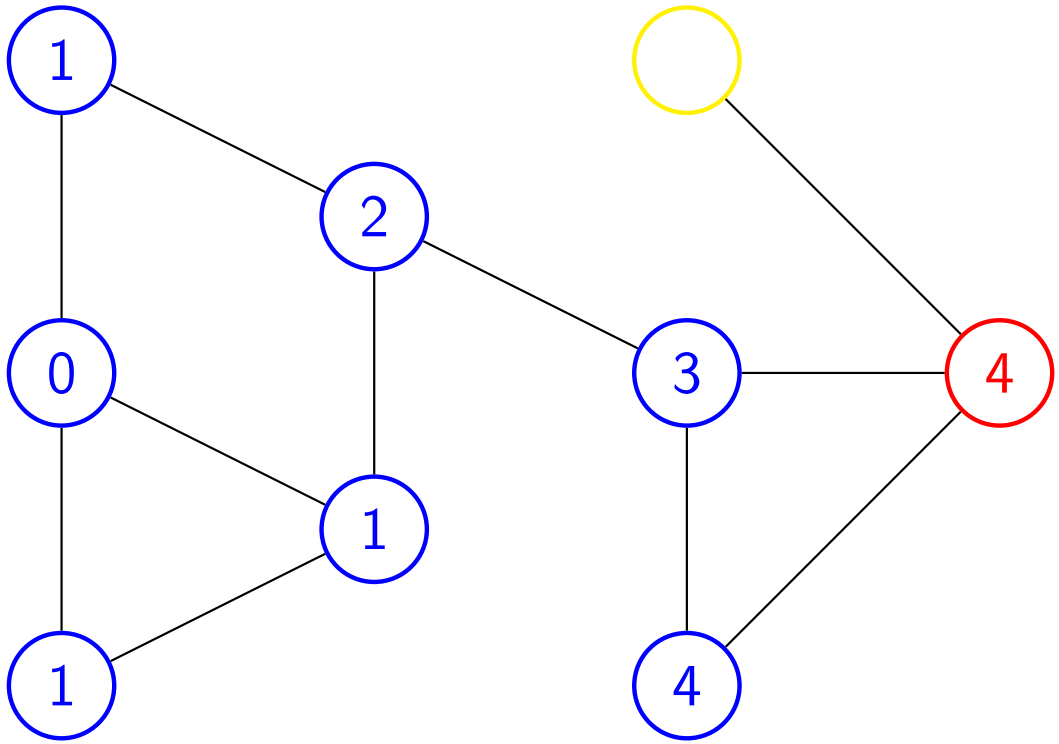
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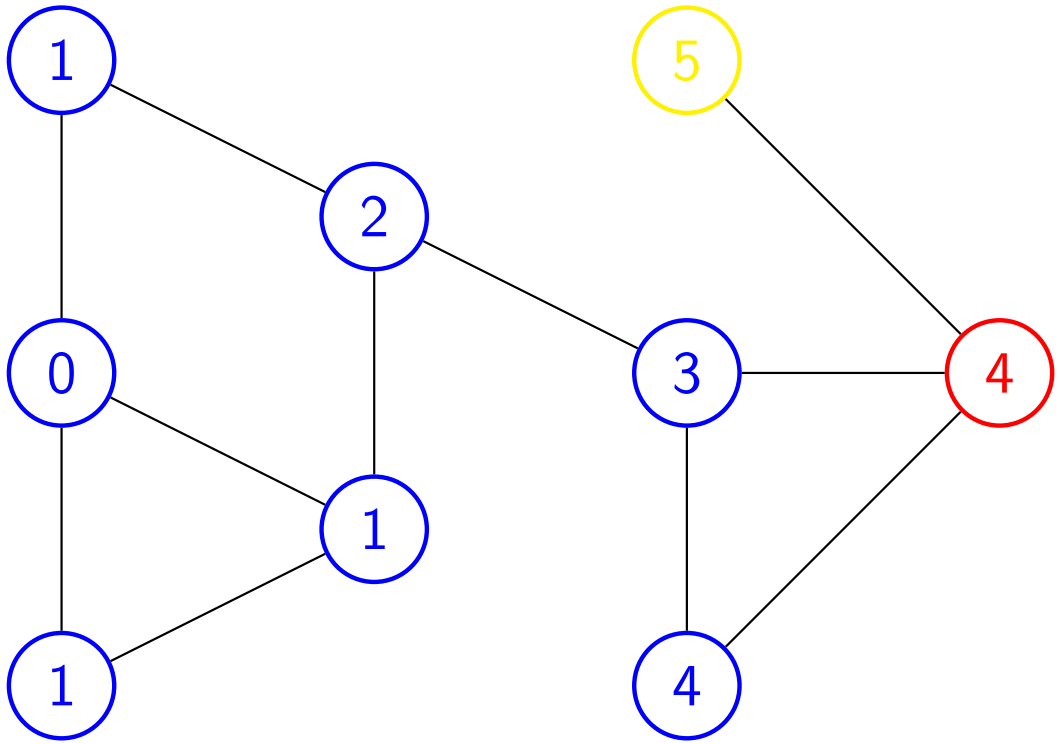
1
5
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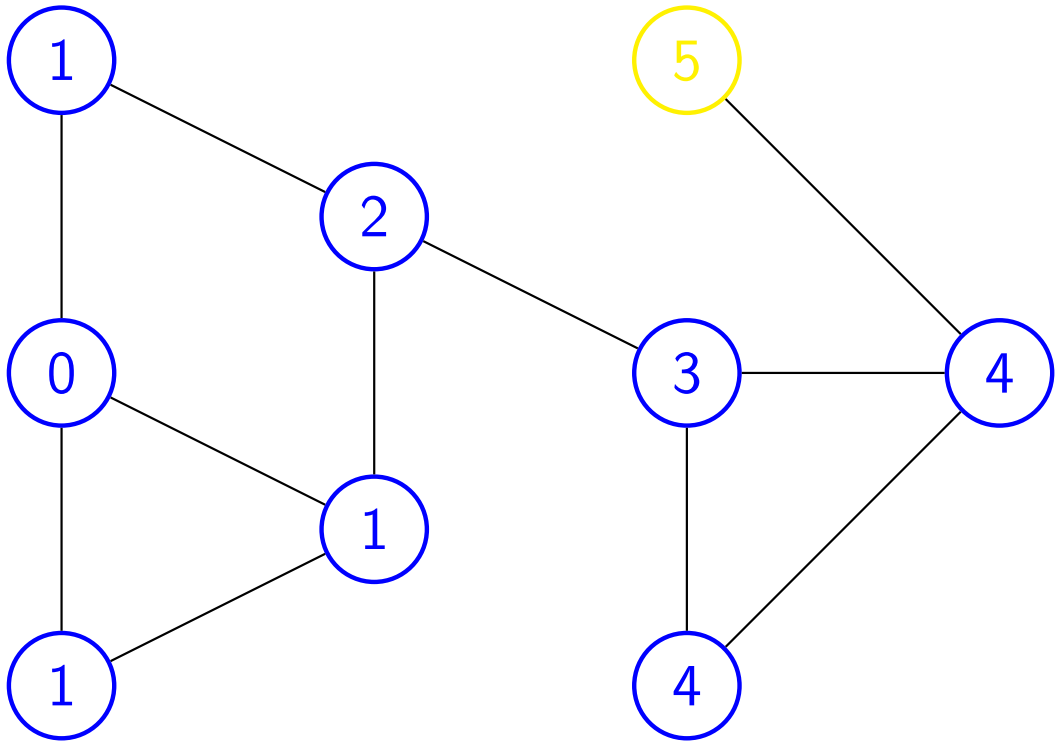
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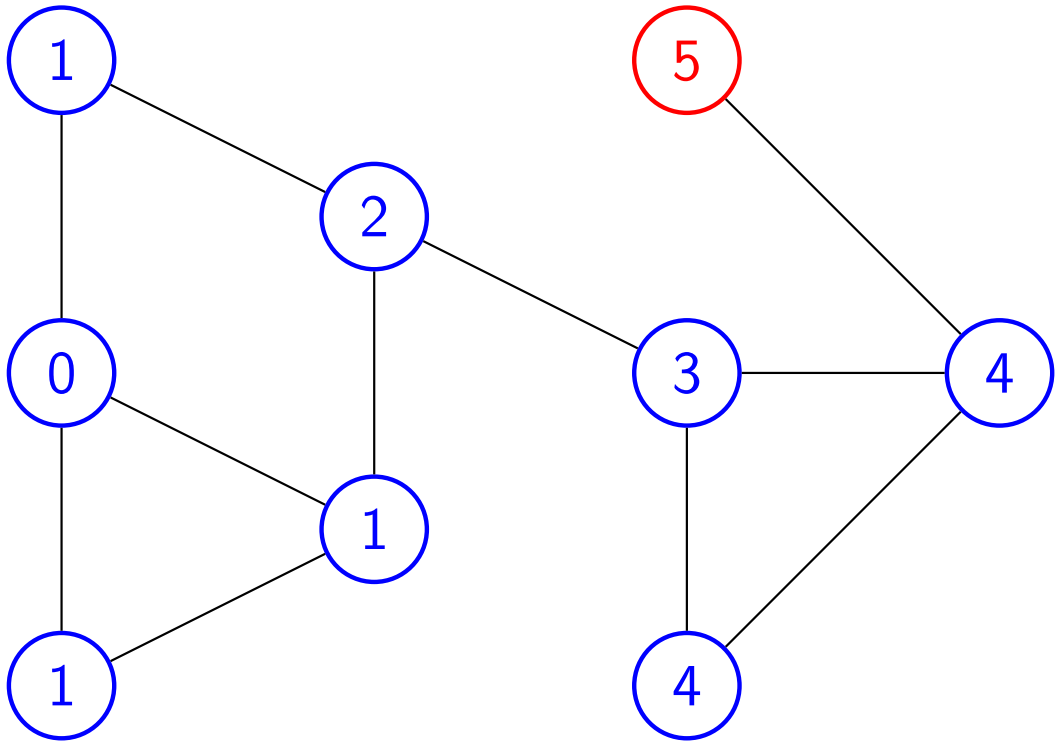
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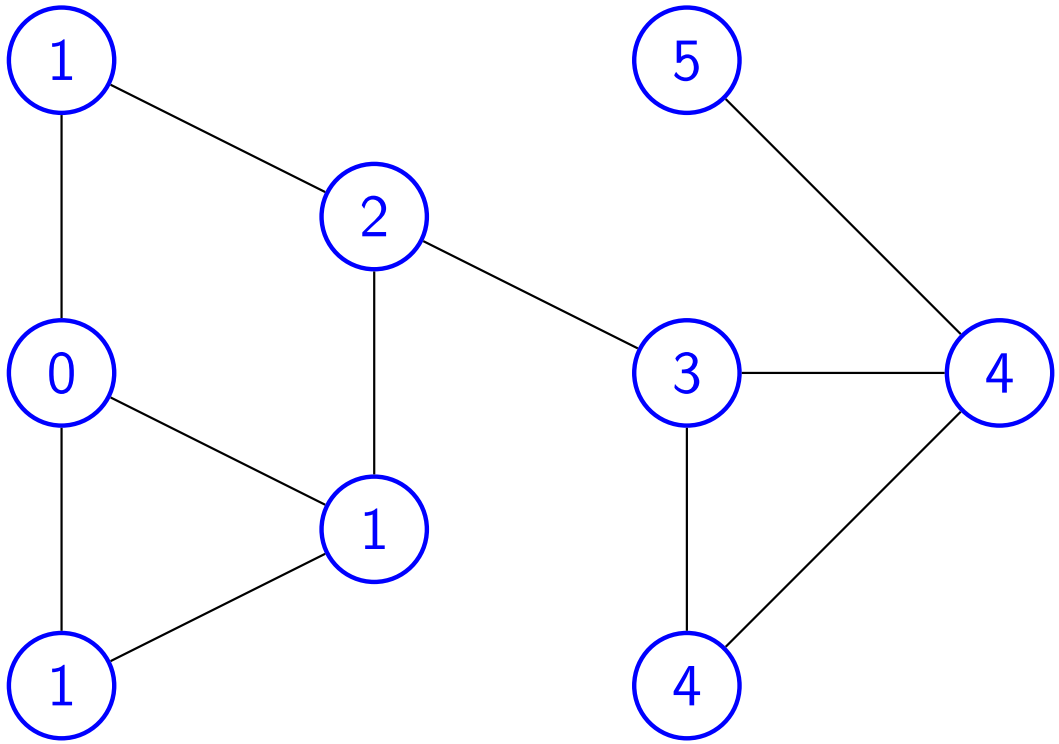
1
5
2
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1
4

- Við munum halda utan um „séða" hnúta með biðröð.
- Við byrjum því á að setja upphafshnútinn okkar í biðröðina.
- Við tökum svo hnút úr biðröðinni, setjum alla „óséða" nágranna hans í biðröðina og höldum áfram þangað til biðröðin er tóm.

```
18      vi d(n, -1);
19      queue<int> q;
20      q.push(0);
21      d[0] = 0;
22      while (q.size() > 0)
23      {
24          int x = q.front();
25          q.pop();
26          for (i = 0; i < g[x].size(); i++) if (d[g[x][i]] == -1)
27          {
28              q.push(g[x][i]);
29              d[g[x][i]] = d[x] + 1;
30          }
31      }
```

- Við segjum að hnútar u og v séu fjarlægð k frá hvorum öðrum ef stysti vegurinn frá u til v er af lengd k.
- Við segjum líka að það séu k skref á milli hnútanna.
- Ef enginn vegur er á milli hnútanna segjum við að lengdin á milli þeirra sé ∞.
- Mikilvægur eiginleiki breiddarleitar er að hún heimsækir fyrst þá hnúta sem eru næst upphafshnútnum.
- Með öðrum orðum, ef u er k_1 skref frá upphafshnútnum og v er k_2 skref frá upphafshnútunum, $k_1 \neq k_2$, þá heimsækir breiddarleit u á undan v þá og því aðeins að $k_1 < k_2$.
- Við getum því notað breiddarleit til að finna fjarlægðina frá upphafshnútnum að öllum öðrum hnútum.

- Líkt og í dýptarleit þá heimsækjum við hvern hnút í mesta lagi einu sinni og ferðumst eftir hverjum legg í mesta lagi tvisvar (einu sinni í stefndu neti).
- Svo tímaflækjan er aftur $\mathcal{O}(E + V)$.

- Báðar leitirnar segja okkur til hvaða hnúta megi komast frá upphafshnútnum og gera það með sömu tímaflækju.
- Breiddarleit gefur okkur einnig fjarlægð allra hnúta frá upphafshnútnum.
- Í dýptarleit getum við unnið áfram með gögnin eftir endurkvæma kallið okkar, sem býður upp á mikla fjölbreyttni.
- Dýptarleit má því finna í reikniritum sem finna grannröð neta, tengipunkta og brýr (þetta verður allt skilgreint seinna).

- Tökum dæmi.
- Fyrsta lína inntaksins inniheldur tvær heiltölur, r og c.
- Inntakið inniheldur síðan r strengir, allir af lengd c.
- Strengirnar byrjar og enda allir á `'X'` ásamt því að fyrsti og síðasti strengurinn inniheldur bara stafinn `'X'`.
- Annars innihalda strengirnir bara stafina `'X'`, `'.'` og eitt stykki `'O'`.

```
11 14
XXXXXXXXXXXXXX
X.X..........X
X.X.XXXX.X.X.X
X.X.....XXXX.X
XOXXXXXX...X.X
X..X.X...X.X.X
X...XX.XXX.X.X
X....X...X.X.X
XXXX.XXX.X.X.X
X........X...X
XXXXXXXXXXXXXX
```

- Við viljum svo prenta sama borð, nema í stað bókstafana á að koma hversu fá skref við þurfum að taka til frá `'O'` til að komast þangað ef við megum ferðast upp, niður, til hægri og til vinstri, en ekki á reitunum með `'X'`.
- Fyrir þá reiti sem við komumst ekki á prentum við `-1`.

- Sem dæmi hefur inntakið

```
11 14
XXXXXXXXXXXXXX
X.X..........X
X.X.XXXX.X.X.X
X.X.....XXXX.X
XOXXXXXX...X.X
X..X.X...X.X.X
X...XX.XXX.X.X
X....X...X.X.X
XXXX.XXX.X.X.X
X........X...X
XXXXXXXXXXXXXX
```

úttakið

```
-1 -1 -1 -1 -1 -1 -1 -1 -1 -1 -1 -1 -1 -1
-1  3 -1 47 46 45 44 43 42 41 40 39 38 -1
-1  2 -1 48 -1 -1 -1 -1 43 -1 41 -1 37 -1
-1  1 -1 49 50 51 52 53 -1 -1 -1 -1 36 -1
-1  0 -1 -1 -1 -1 -1 -1 21 22 23 -1 35 -1
-1  1  2 -1 -1 -1 18 19 20 -1 24 -1 34 -1
-1  2  3  4 -1 -1 17 -1 -1 -1 25 -1 33 -1
-1  3  4  5  6 -1 16 15 14 -1 26 -1 32 -1
-1 -1 -1 -1  7 -1 -1 -1 13 -1 27 -1 31 -1
-1 11 10  9  8  9 10 11 12 -1 28 29 30 -1
-1 -1 -1 -1 -1 -1 -1 -1 -1 -1 -1 -1 -1 -1
```

- Hvernig tengist þetta dæmi efni vikunnar?
- Við getum túlkað þessa mynd sem net.
- Ímyndum okkur að hver auður reitur sé hnútur.
- Við tengjum svo aðliggjandi auða hnúta með leggjum.
- Þar sem við viljum finna fjalægðir frá tilteknum hnút til allra annara hnúta notum við breiddarleit.
- Fjöldi hnúta í netinu er alltaf minni en $r \cdot c$ og fjöldi leggja er alltaf minni en $2 \cdot r \cdot c$.
- Svo þetta reiknirit er $\mathcal{O}(E + V) = \mathcal{O}(r \cdot c)$.

```cpp
int isin(int x, int y, int r, int c)
{
    if (x >= r || x < 0 || y >= c || y < 0) return 0;
    return 1;
}

int main()
{
    int i, j, r, c, x, y, g[4][2] = {{-1, 0}, {1, 0}, {0, -1}, {0, 1}};
    cin >> r >> c;
    string a[r];
    for (i = 0; i < r; i++) cin >> a[i];
    for (i = 0; i < r; i++) for (j = 0; j < c; j++) if (a[i][j] == 'O')
        x = i, y = j;
    int d[r][c];
    for (i = 0; i < r; i++) for (j = 0; j < c; j++) d[i][j] = -1;
    queue<ii> q;
    q.push(ii(x, y)), d[x][y] = 0;
    while (q.size() > 0)
    {
        ii p = q.front(); q.pop();
        x = p.first, y = p.second;
        for (i = 0; i < 4; i++)
        {
            int xx = x + g[i][0], yy = y + g[i][1];
            if (a[xx][yy] == 'X' || !isin(xx, yy, r, c) || d[xx][yy] != -1)
                continue;
            q.push(ii(xx, yy)), d[xx][yy] = d[x][y] + 1;
        }
    }
    for (i = 0; i < r; i++)
    {
        for (j = 0; j < c; j++)  printf("%2d ", d[i][j]);
        printf("\n");
    }
    return 0;
}
```

- Tökum annað dæmi.
- Ykkur er gefið net og tiltekinn hnút u.
- Prentið alla einfalda vegi í netinu sem byrja í u.
- Munið að vegur er einfaldur ef hann heimsækir aldrei sama hnútinn tvisvar.

- Við getum leyst þetta með því að breita dýptarleitar útfærslunni okkar lítillega.
- Til að koma í veg fyrir að heimsækja hnút oftar en einu sinni í dýpterleit merkjum við hann og heimsækjum ekki merkta hnúta.
- Við munum ennþá þurfa að merkja hnúta því við erum að leita að einföldum vegum (almennt er ekki takmarkaður fjöldi vega í neti).
- Munurinn er að við munum merkja hnút þegar við sjáum hann, halda áfram endurkvæmt til ómerktra nágranna hans og afmerkja hann svo.

```
 7 void dfs(vvi& g, int x)
 8 {
 9     int i;
10     v[x] = 1;
11     p.push_back(x);
12     for (i = 0; i < p.size(); i++) printf("%d ", p[i] + 1);
13     printf("\n");
14     for (i = 0; i < g[x].size(); i++) if (v[g[x][i]] == 0)
15         dfs(g, g[x][i]);
16     p.pop_back();
17     v[x] = 0;
18 }
```

- Þetta forrit prentar alla einfalda vegi sem byrja í hnút u.
- Til að hámarka fjölda slíkra vega getum við búið til net þar sem öll pör hnúta eru nágrannar.
- Þá myndi þetta forrit prenta allar umraðanir sem byrja á x.
- Tímaflækjan er því $\mathcal{O}(V!)$.