Tasks
1) Bare Metal (no 05)
- some real-time systems are constructed with no 05 - this can be done using a "superloop"
- Ph trul pro to + 1Ha
task 2 executes at 0.5 Hz  main  task 1 (just a function)  task 1 (just a function)
5uper loop
busy wait ~15
busy wast

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e	+			
٥.	sp	in	H	
S	pi	n 3	++	)
	(C)			

#include < X/X stabool.h> void task 2 (void); of do some work void task 2 (void); of thet actually tosks y int main (wird) { while (frue) {
task 1(); task2(); busy ugit Sfor (unt32t spin =0; spin < 17000 coo task1 (); for (un+32 t spn=6; spin < 17000000; 3 Treat-time programs don't terminate task 1 () and task 2 () are synchronous - they execute periodically shortcomings:
-timing is approximate
- cannot respond quickly to asynchronous
events (IRQs) - a better method is to keep time with the sectick interrupt

main

entry Sectick ISR

This example expands on the previous one by adding an asynchronaus fush button input

false

tout to pushed?

False

tout on a synchronaus

false

tout on input

- the push button is labelled INTO in the lab

#include <LPC17xx.h > #include <stdbool.h > #include <stdint.h > uint 32\_t ms Tick = 0; millisecond tick count bool button Pushed = false; Her void EINT3 IRQ Handler (void) ENTO (push button) ISR buttonPushed =true; LPC\_GPIOINT -> IO2 Int C(r /= (1<<10); - clear turren void task 1 (void); void task 2 (word); task 3 (void);

int main (void) {
 uint32 t t1 Next = 0; } when to run each task
 uint32 t t2 Next = 0; }

setup INTO();
Systick\_ Config (System Core Clerk / 1000);

interrupt every 1 ms while (true) {
if (ms Ticks >= t1 Next) fails when ms Ticks
over flows If (ms Ticks - t 1 Next >= (1000) task 18(); t1 Next += 1000; add 1s 3
if(msTicks-t2Nex+>= \$2000)
task2();
t2Nex++=2000; add 25 if (button hished) {
button hished = false;
task 3 ();
}

- this example is not a true superloop; it is an example of a "roll-your-own" scheduler





