

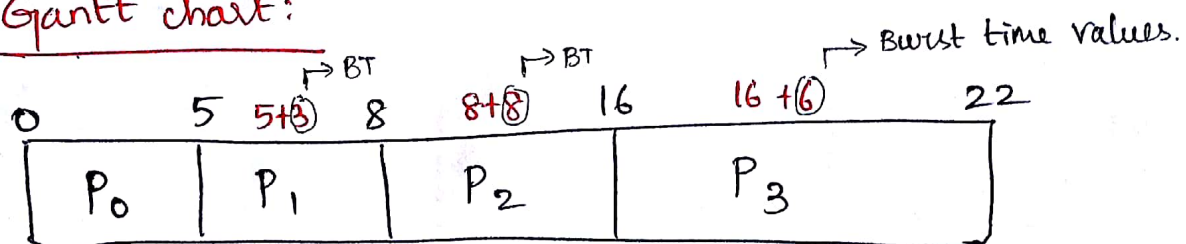
CPU SCHEDULING ALGORITHMS:

- * FCFS → first come, first serve
- * SJF → shortest job first.
- * SRT → Shortest Remaining Time.
- * Priority based scheduling
- * Round robin based scheduling.

FCFS:

Process	Arrival time (AT)	Burst time (BT)	Completion time (CT)	CT-AT	TAT-BT
				Turnaround time (TAT)	Waiting time (WT)
P ₀	0	5	5	5	0
P ₁	1	3	8	7	4
P ₂	2	8	16	14	6
P ₃	3	6	22	19	13

Gantt chart:



* completion time values are found by using Gantt chart.

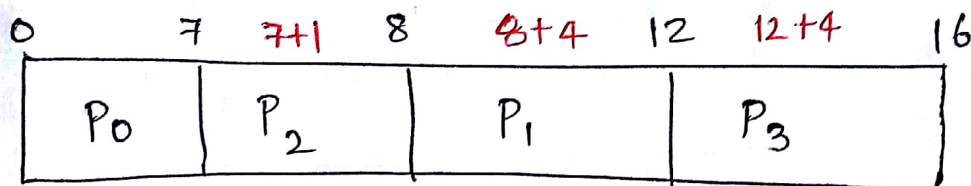
$$\begin{aligned}\text{Avg (TAT)} &= \frac{\text{tot (TAT)}}{\text{no. of process}} = \frac{5 + 7 + 14 + 19}{4} \\ &= \frac{45}{4} = 11.25 \text{ msec} \\ \text{Avg (WT)} &= \frac{0 + 4 + 6 + 13}{4} = \frac{23}{4} = 5.75 \text{ msec}\end{aligned}$$

SJF: (check notes for explanation) \Rightarrow AT \rightarrow Arrival time
 BT \rightarrow Burst time
 CT \rightarrow Completion time
 TAT \rightarrow Turnaround time
 WT \rightarrow Waiting time

- * Shortest Job first
- * Non-preemptive

Process	AT	BT	CT	TAT	WT
P ₀	0	7	7	7	0
P ₁	2	4	12	10	6
P ₂	4	1	8	4	3
P ₃	5	4	16	11	7
				tot: 32	tot: 16

Gantt chart:



$$\text{Avg (TAT)} = \frac{32}{4} = 8 \text{ msec.}$$

$$\text{Avg (WT)} = \frac{16}{4} = 4 \text{ msec.}$$

- \because P₀ & P₁ are Shortest processes.
- \because P₀ completed.
- \because P₁ P₂ P₃ are remaining in Queue
- \because $\frac{P_1 \ P_2 \ P_3}{4 \ 1 \ 4} \Rightarrow$ BT values
- \because In the above processes check BT values.
- \because P₂ is small.
- \because remaining P₁ P₃ with same BT values
- \because Take small process as P₁ \rightarrow first
P₃ \rightarrow Last.

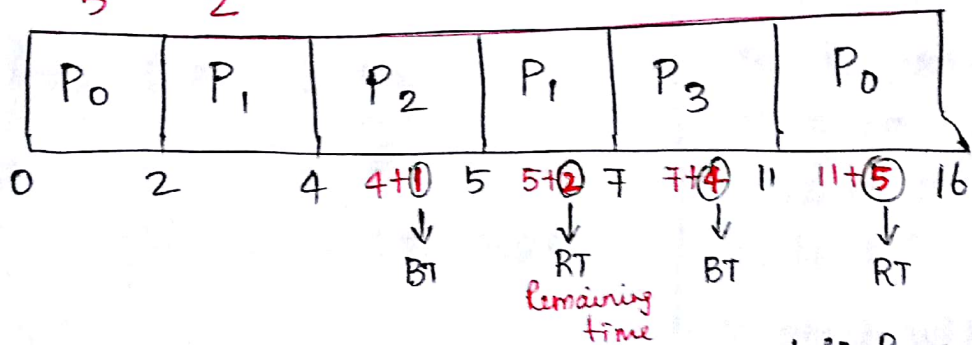
SRT:

* Shortest Remaining Time

* preemptive. (A process is pause the process to switching over to another process depending upon some criteria)

Process	AT	BT	CT	TAT	WT
P ₀	0	7	16	16	9
P ₁	2	4	7	5	1
P ₂	4	1	5	1	0
P ₃	5	4	11	6	2

Remaining Time for P₀ & P₁ (shortest processes)



$$\text{Avg (TAT)} = \frac{28}{4}$$

$$= \boxed{7 \text{ msec.}}$$

$$\text{Avg (WT)} = \frac{12}{4}$$

$$= \boxed{3 \text{ msec.}}$$

∵ P₀ & P₁ are the two shortest processes.

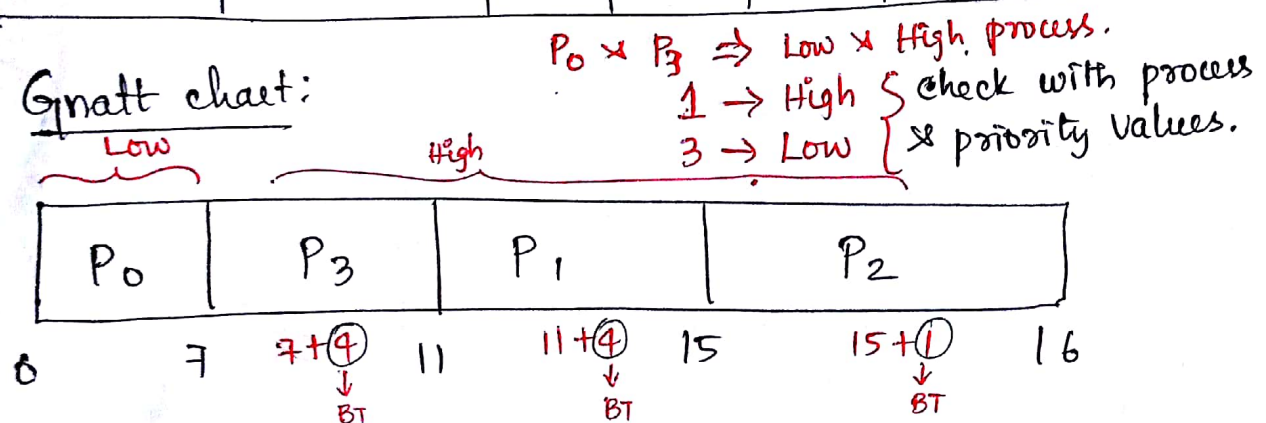
∵ As given in the name we need to find remaining time (RT) for P₀ & P₁

⇒ P₀ = BT → 7 } 7-2
 AT → 2 } 5
 P₁ = BT → 4 } 4-2
 AT → 2 } 2
 ∵ other P₂ & P₃ are calculated with BT.

Priority Based Scheduling:

Process	Priority	AT	BT	CT-BT		TAT-BT
				CT	TAT	WT
P ₀	3	0	7	7	7	0
P ₁	2	2	4	15	13	9
P ₂	4	4	1	16	12	11
P ₃	1	5	4	11	6	2

Gantt chart:



$$\text{Avg (TAT)} = \frac{38}{4}$$

$$= 9.5 \text{ msec.}$$

$$\text{Avg (WT)} = \frac{22}{4}$$

$$= 5.5 \text{ msec.}$$

∴ first finish off the P₀ process. since it is Low priority.

∴ Next High priority with P₁ P₂ P₃ processes.

∴ start with (P₃) with BT value 4

↳ P₃ : 1 → High

∴ next P₁ with BT value 4

↳ P₁ : 2 → High

∴ Next P₀ → we completed already.

∴ Next P₂ with BT value 1

↳ P₂ : 4 → High

Round Robin based scheduling:

Quantum value = 3 msec.

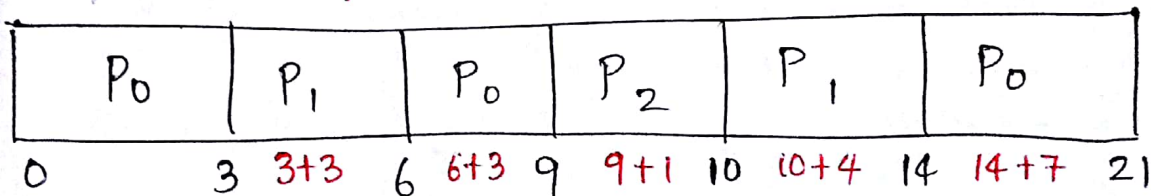


Also called as time slice.

Process	AT	BT	CT - AT		TAT - BT
			CT	TAT	WT
P ₀	0	7	21	21	14
P ₁	2	4	14	12	8
P ₂	4	1	10	6	5

Quantum: Maximum time CPU given to the process at a single point of time.

Short process



$$\text{Avg (TAT)} = \frac{39}{4}$$

$$= 9.75 \text{ msec}$$

$$\text{Avg (WT)} = \frac{27}{4}$$

$$= 6.75 \text{ msec}$$

∵ P₀ & P₁ are short process add with the Quantum values.

∵ and after that all the three processes are added with the Burst time (BT) values.