Assignment (3) Compsci 10m3 Sepanta Kamali Q1) Show that  $\lfloor \frac{n}{2} \rfloor \lceil \frac{n}{2} \rceil = \lfloor \frac{n^2}{4} \rfloor$  for all  $n \in \mathbb{Z}$ casel: niseven Casez: nisodd n=2K+1, KEZ n=ZK, KEZ n = K and because KEZ  $\frac{N}{N} = \frac{2}{2}K+1 = K+\frac{1}{2}$ therefore:  $\lfloor \frac{n}{2} \rfloor = K$ ,  $\lfloor \frac{n}{2} \rfloor = K$  $\begin{bmatrix} \frac{1}{2} \end{bmatrix} = \begin{bmatrix} K + \frac{1}{2} \end{bmatrix} = K$ therefore  $\lfloor \frac{n}{2} \rfloor \lceil \frac{n}{2} \rceil = K^2$ [2] = K+1 So L2][2] = K(K+1) and as  $K = \frac{n}{2}$  so  $K^2 = \frac{n^2}{4}$ we know that n=ZK+1 and we know that it K is 20 N3 = (SK+1) = AK3+ AK+1 an integer, then K2 is also n2 4K2+4K+1 = K2++1 an in teger so [n2]= n2 and we know that  $K^2$  and K are in tegers so  $\lfloor \frac{n^2}{4} \rfloor = \lfloor K^2 + K + \frac{1}{4} \rfloor = K^2 + K$ The left and right sides are equal so it's proved. = K(K+1) So the lettand right sides are equal. Therefore, we can conclude that [ ] [ n ]= [ n ] for all integers n.

$$\begin{array}{c} (3z) \quad a_{n} = 8a_{n-1} - 16a_{n-2} \\ a) \quad a_{n-1} = 1 \quad a_{n-2} = 1 \\ a_{n-1} = 1 \quad a_{n-2} = 1 \\ a_{n} \quad a_{n-2} = 1 \\ a_{n} \quad a_{n-1} = 2^{n-1}, \quad a_{n-2} = 2^{n-2} \\ a_{n} \quad a_{n} \quad a_{n-1} - 16a_{n-2} \\ a_{n} = 8a_{n-1} - 16a_{n-2} \\ a_{n} = 9k2 - 16k2 \\ a_{n} = 9k2 - 16k2 \\ a_{n} = 2k2 - 16k2 \\ a_{n} = 2k2 - 2k2 \\ a_{$$

b) 
$$\sum_{j=0}^{2} (j-1)(2j^{2}+1)$$
 $j=10$ 

$$\sum_{j=10}^{2} (j-1)(2j^{2}+1) = \sum_{j=10}^{2} (2j^{2}+2j^{2}+1)$$
 $j=10$ 

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So we have 
$$\mathcal{E} \log (\frac{K+1}{K})$$

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 $\mathcal{E} \log (\frac{K+1}{K}) = \mathcal{E} \log (K+1) + (\log (1+1) - \log K)$ 
 $\mathcal{E} \log (\frac{K+1}{K}) = \mathcal{E} \log (1+1)$ 
 $\mathcal{E} \log (1+1) = \log$ 

b) if  $A = \begin{pmatrix} 1 & 3 \\ -5 & 2 \end{pmatrix}$  then  $A^2 = \begin{pmatrix} 1 & 7 \\ 25 & 4 \end{pmatrix}$  False  $A^{2} = A \times A \longrightarrow \begin{bmatrix} 1 & 3 \\ -5 & 2 \end{bmatrix} \begin{bmatrix} 1 & 3 \\ -5 & 2 \end{bmatrix} = \begin{bmatrix} -14 & 9 \\ -15 & -11 \end{bmatrix}$ C) if A is 6x4 matrix and B is unts matrix then AB has 16 entries False we can only multiply matricis A and B it the number of the columns of A is equal to number of rows of B So if A is man, B has to be nxp and the AB matrix would be a mxp matrix. So in this case as A is 6x4 and B is 4x3, the AB matris would be 6x5 and has 30 entries. det lorgest\_prod (arr, N:real) arr={a,a,-,an} for i= 1 to n-1 0 For j = i+1 to no if ai \*aj < N : t = max(t, a; \*aj)return t the big-o of this algorithm is O(n2) as there are two nested for loops inside of each other.

Q9) prove that 
$$\frac{6n^2 + 4n^5 - 4}{7n^2 - 3}$$
 is  $O(n^3)$ 

We can break  $6n^2 + 4n^5 - 4$  into separate terms:

 $7n^2 - 3$ 

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As the term  $4n^5 + 6n^5 + 4n^5 - 4$ 

Fastest, the big -0 of the entire function is approaximately equal to the big -0 at the term  $4n^5 + 6n^2 - 3$ 

Also when grows, the -3 in the  $6n^2 - 3n^2 - 3$ 

Alianminator doesn't affect the big -0 as much as  $6n^2 - 3n^2 - 3n$ 

Q10)

0) 
$$Hn) = 1+2+3+...+(n^2-1)+n^2$$
  $O(n^4)$ 

b)  $Hn) = 3n^4$   $O(n^4)$ 

Q11)

Q11)

Q1)

 $x^3 + 7n^2 + 3$ 
 $2n+1$ 
 $2n+1$ 

$$\frac{\chi^{3}}{2\chi+1} \geqslant \frac{\chi^{2}}{3\chi}$$

$$\frac{\chi^{3}}{2\chi+1} \geqslant \frac{\chi^{2}}{3\chi} \qquad \chi > 0 \qquad C = \frac{1}{3}, K = 0$$

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$$\frac{3\pi}{2}$$

So we conclude that 
$$\frac{\chi^3}{2n+1}$$
 is  $\frac{\pi}{2}$ 

As 
$$x^3 + 7n + 3$$
; s both  $O(x^2)$  and  $O(x^2)$ ,















10000 < logn (3n (n2 logn (n3+88n2+3 (n.2)