

- Lemma 1:

If T is a tree with n vertices and m edges, then $m=n-1$

- i. Proof:
- ii. Base case: if $m = 0$, then T , being connected has only 1 vertice.
- iii. $\forall m > 0: P(0), \dots, P(m-1) \Rightarrow P(m)$
- iv. let $m > 0$, assume for any tree T' with $|E(T')| < m$ that
- v. $|E(T')| = |V(T')| - 1$
- vi. must show: if T has m edges and n vertice, then $m=n-1$
- vii. Pick any edge $e \in E(T)$ and remove it. This results in two subtrees: T_1 and T_2 each with fewer than m edges:
- viii. $|E(T_i)| < m$ for $i = 1, 2$
- ix. By the induction hypothesis, we have $|E(T_i)| = |V(T_i)| - 1$ for $i = 1, 2$
- x. since no vertices were removed, $|V(T_1)| + |V(T_2)| = n$
- xi. Therefore, $m = |E(T_1)| + |E(T_2)| + 1 = |V(T_1)| - 1 + |V(T_2)| - 1 + 1$
- xii. $= |V(T_1)| + |V(T_2)| - 1$
- xiii. $= n-1$ as required

- Lemma 2:

If G is acyclic with n vertices, m edges, and k connected components, then $m=n-k$

- i. Proof:
- ii. Let $T_1, T_2, T_3, \dots, T_k$ denote the connected components of G , which are necessarily trees.
- iii. Let m_i, n_i denote #vertices and #edges in T_i ($1 \leq i \leq k$)
- iv. By Lemma 1: $|E(T_i)| = |V(T_i)| - 1$ ($1 \leq i \leq k$) ie $m=n-1$
- v. SO, $m = \sum_{i=1}^k m_i = \sum_{i=1}^k (n_i - 1)$
- vi. $= \sum_{i=1}^k (n_i) - \sum_{i=1}^k (1) = n - k$

- Lemma 3:

If G is connected with n vertices and m edges, then $m \geq n - 1$

- i. Proof
- ii. if $m=0$, G being connected, can have only 1 vertice, so $n=1$
 $m=0 \geq 0=n-1$

- iii. $\forall m > 0: P(0), P(1), \dots, P(m-1) \Rightarrow P(m)$
- iv. let $m > 0$ be chosen arbitrarily
- v. Assume for any connected Graph G' that $|E(G')| \geq |V(G')| - 1$
- vi. Must show: if G is connected with m edges, n vertices, then $m \geq n-1$
- vii. Pick any $e \in E(G)$ and remove it. We have 2 cases:
 - a) Case 1: $G-e$ is connected
 - Note $G-e$ has n vertices and $m-1$ edges
 - By the induction hypothesis: $m-1 \geq n-1$
 - Therefore, $m \geq n \geq n-1$
 - $m \geq n-1$
 - b) Case 2: $G-e$ is NOT connected
 - Claim: $G-e$ consists of two connected components: G_1 and G_2
 - Note: $|E(G_i)| < m \quad i = 1, 2$
 - So, by Induction Hypothesis, we have $|E(G_i)| \geq |V(G_i)| - 1 \quad i = 1, 2$
 - also $|V(G_1)| + |V(G_2)| = n$ since no vertices were removed
 - Thus, $m = |E(G_1)| + |E(G_2)| + 1 \geq |V(G_1)| - 1 + |V(G_2)| - 1$ by ind hyp
 - $= |V(G_1)| + |V(G_2)| - 1$
 - $= n-1$
 - $= m \geq n-1$
 - in Each case $m \geq n-1$

- Lemma 4:

If G is a graph with n vertices, m edges, and k connected components, then $m \geq n-k$

- i. Proof:
- ii. Let G_1, G_2, \dots, G_k be the connected components of G
- iii. By Lemma 3: Let m_i, n_i be #edges and #vertices in G_i ($1 \leq i \leq k$)
- iv. $m_i \geq n_i - 1$
- v. then: $m = \sum_{i=1}^k (m_i) \geq \sum_{i=1}^k (n_i - 1) = \sum_{i=1}^k (n_i) - \sum_{i=1}^k (1)$
- vi. $= n-k$
- vii. therefore $m \geq n-k$