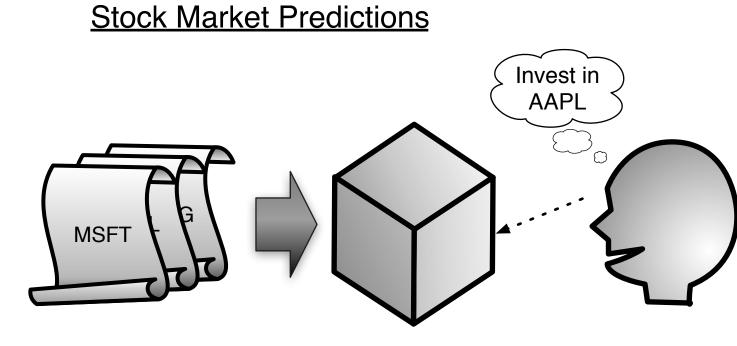
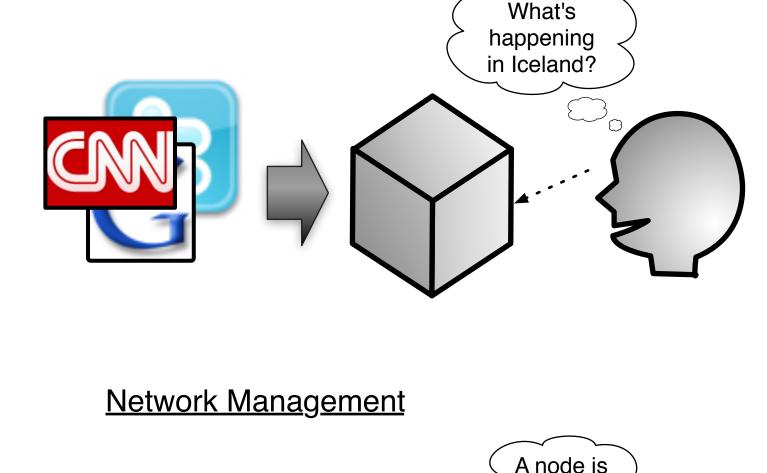
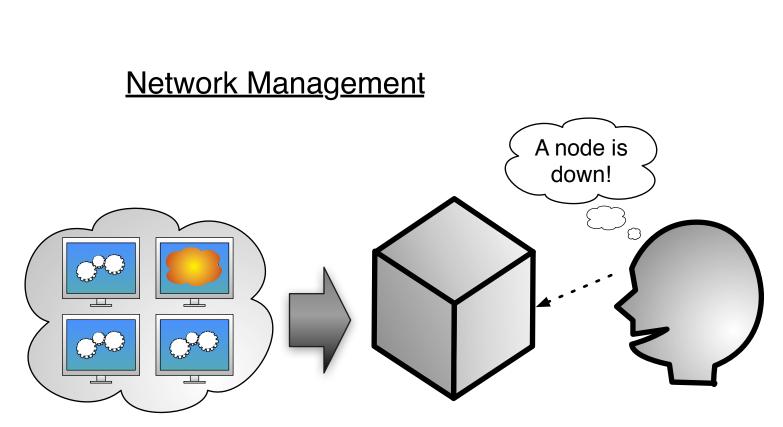


Trend Monitoring







## STREAM MONITORING

**Approach** 

**Stream Processors** 

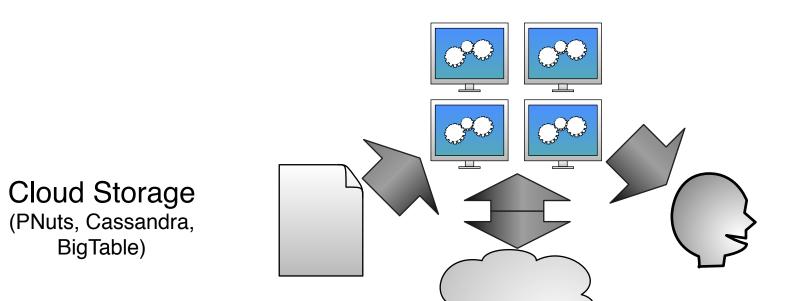
(Streambase, Cayuga)

#### <u>Design</u>

# **Limitations**

- Limited Storage - Limited Processing Power

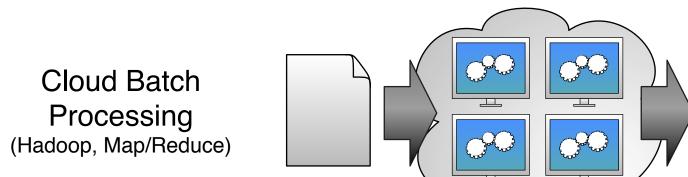
Limited Query Complexity\*



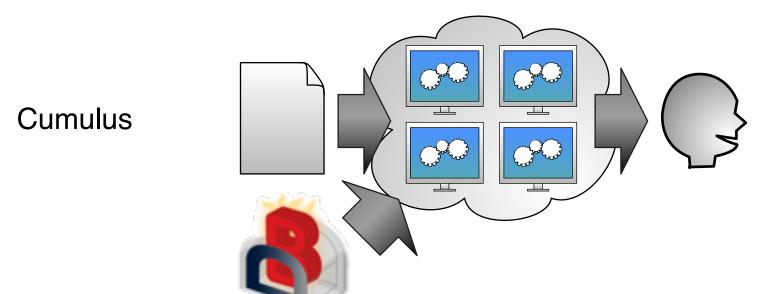
- High Development Costs - Consistency Limitations

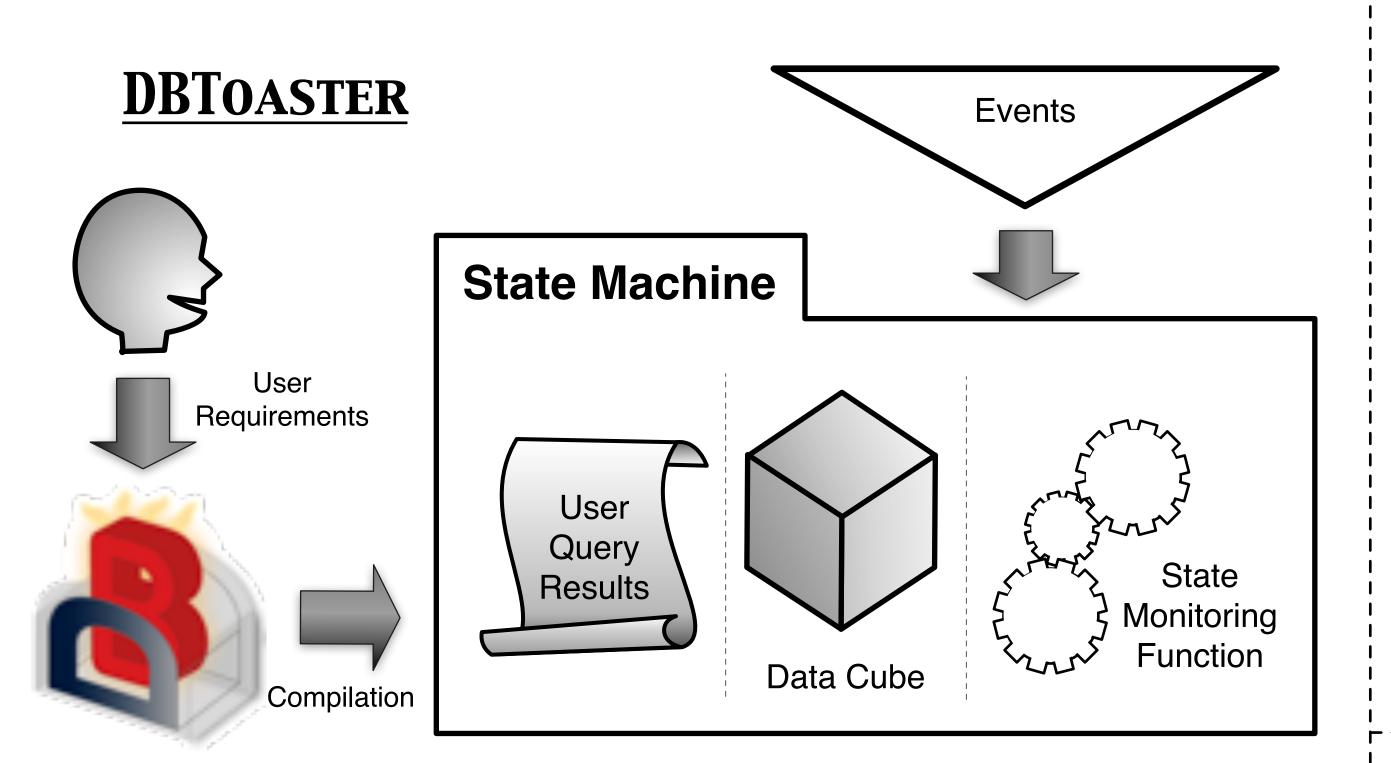
- High Network

Overheads

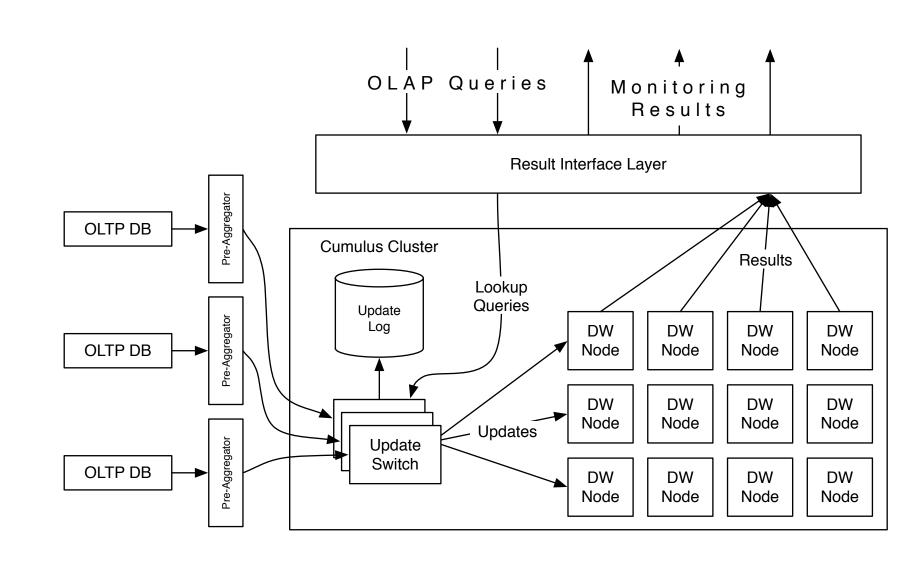


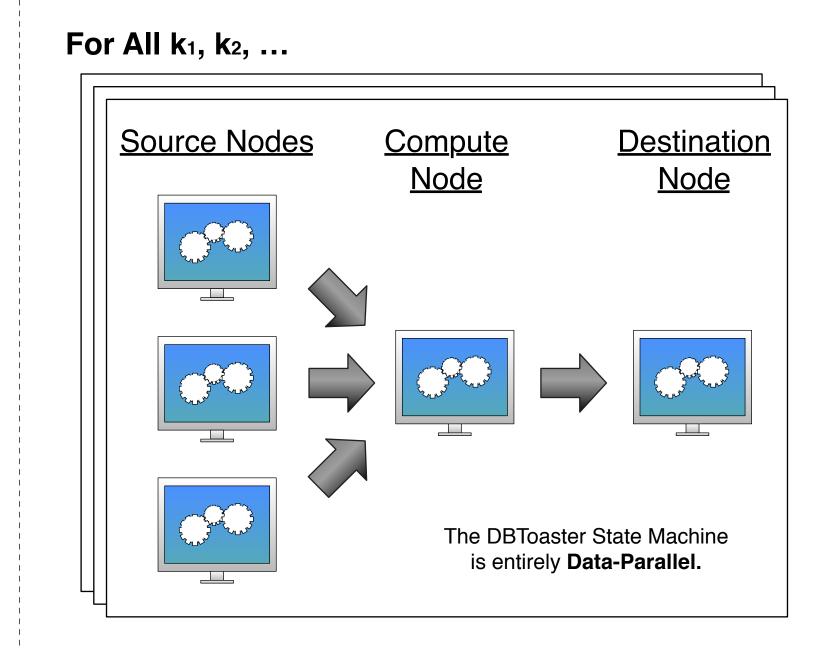
- High Latency - High Network Overheads





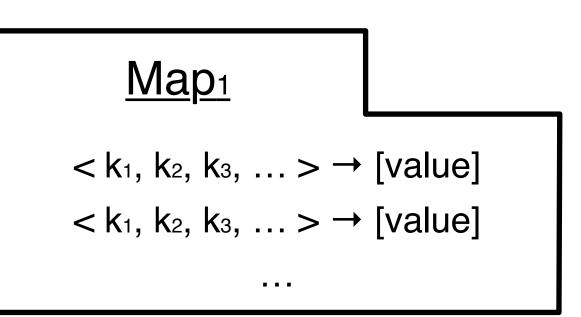
# **CUMULUS ARCHITECTURE**

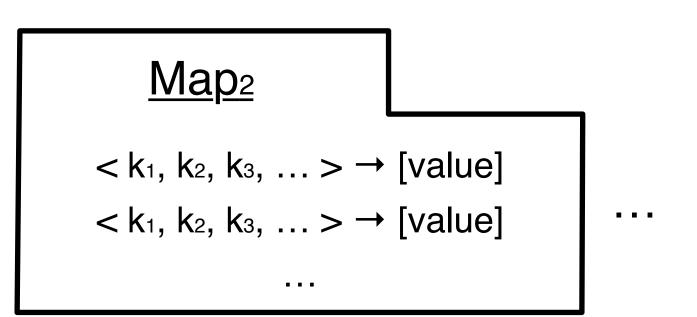




### THE DBT STATE MACHINE

**State: Multikey Maps** 





## **Events: Triggers**

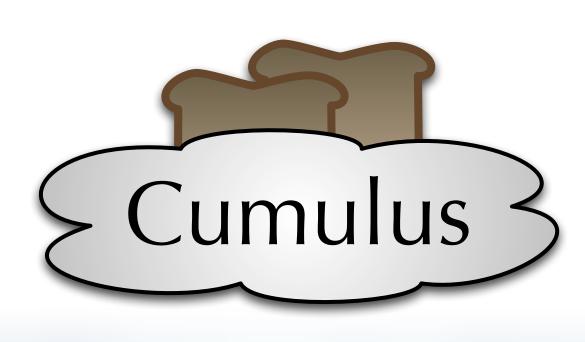
ON Event(param<sub>1</sub>, param<sub>2</sub>, ...){ (e.g., ON TABLE INSERT) statement 1 statement 2

# **Distribution Challenge:**

Ensure triggers run on consistent view of the state.

#### **Statements**

for all  $(k_1, k_2, ...)$  { update Map<sub>1</sub>[ $k_1$ , param<sub>1</sub>, ...] +=  $f(\underline{k}, \underline{param}, \underline{Map_2}, ...)$ 



Oliver Kennedy Aleksandar Vitorovic Christoph Koch

The D.A.T.A. Lab (EPFL)

**CONSISTENCY** 

#### **State Machine Requirements**

#1 : Events are atomic.

#2 : Events must be serialized (in any arbitrary, but consistent order)

Serialized execution requires nodes to know about all events in the system.

Problem: Serialized execution is not scalable!

**Solution:** Allow out of order execution (Separate Chronological and Logical clocks)

Order Conflict  Chronological Order  Logical Order		<u>Fix</u>
Read	Write	Corrective Updates (Read Log)
Write	Read	Versioned Maps (Write Log)
Read	Read	No Correction Needed
Write	Write	Updates are Deltas

