Mini-Project (2)

Step 1: Design & Implement our protocol

Introduction:
Our group built the protocol on the top of UDP socket connection. We implemented connection-oriented, reliable, pipelined protocol which includes flow control and congestion control mechanisms as well. Since UDP does not inherently provide these features, we added mechanisms to ensure reliable communication, similar to TCP.

Connection-oriented and reliable communication: We simulate a **3-way TCP handshake** on top of the UDP protocol to make our protocol connection-oriented. We have a packet class which simlates the TCP packet structure.
Firstly, we create a UDP socket between the client and the server. Then, the client will make a packet with initial sequence number x, set the TCP SYN bit to true, and send the packet to the server. Once the server receives this packet, the server will make a packet with setting TCP SYN bit and ACK bit to true, indicating it is SYNACK, and an initial sequence number decided by the server, an ack number which is x+1.
Next, the server will send the packet it made to the client. When the client receives this packet indicating SYNACK, it shows the server is alive, the connection on the client side is established. The client then will make a packet with ack number to ack the SYNACK, send this packet to the server. If the server receives the correct ACK number, it indicates the client is alive and the connection is established.
Additionally, a timer is set for the SYN packet sent by the client and the SYNACK packet sent by the server. If either of these packets is not acknowledged within the timeout period, it will be resent for maximum 3 times, ensuring reliable communication during the handshake.

Pipelined protocol:
We chose to implement **Go-back-N** to ensure data reliability.
Sender:
1. Sender can have up to N (window size) un-acked packets in pipeline, representing the number of unacked packets in the pipeline.
2. Sender sets timer for oldest un-acket packet. When timer expires, sender will retransmit all un-acked packets.
3. Sends packets sequentially, keep track of the sequence number. The next sequence number is the current sequence number plus the length of its payload.
Receiver:
1. Only sends cumulative ack for the last packet received in order.
2. Does not ack packet if there's a gap, discard out of order packets.
2. Will resend the last ack number to sender if there is a gap.

Flow control:
Our protocol implements flow control using receiver window size (rwnd). The receiver controls sender, and sender will not overflow. The receiver advertises its available buffer size (rwnd) in every acknowledgment packet header.The sender adjusts its transmission rate to make sure the number of outstanding packets does not exceed the receiver's rwnd, so that we ensure the sender does not overwhelm the receiver.

Congestion control:
We simulate the **TCP Tahoe** congestion control.
Slow start: The congestion window (cwnd) starts at 1 packet and doubles the size with every successful acknowledgment until it reaches the threshold (ssthresh) or a packet loss is detected.
Congestion avoidance: When cwnd exceeds ssthresh, the window size grows linearly instead of exponentially.
Loss indicated by timeout or 3 duplicate acks: When detect packet loss, the cwnd is reset to 1 and ssthresh is set to half of the current cwnd.

Socket type: UDP socket, based on UDP, build TCP like protocol.

Packet structure: A Packet class was implemented to simulate a TCP-like packet.
Each packet contains:

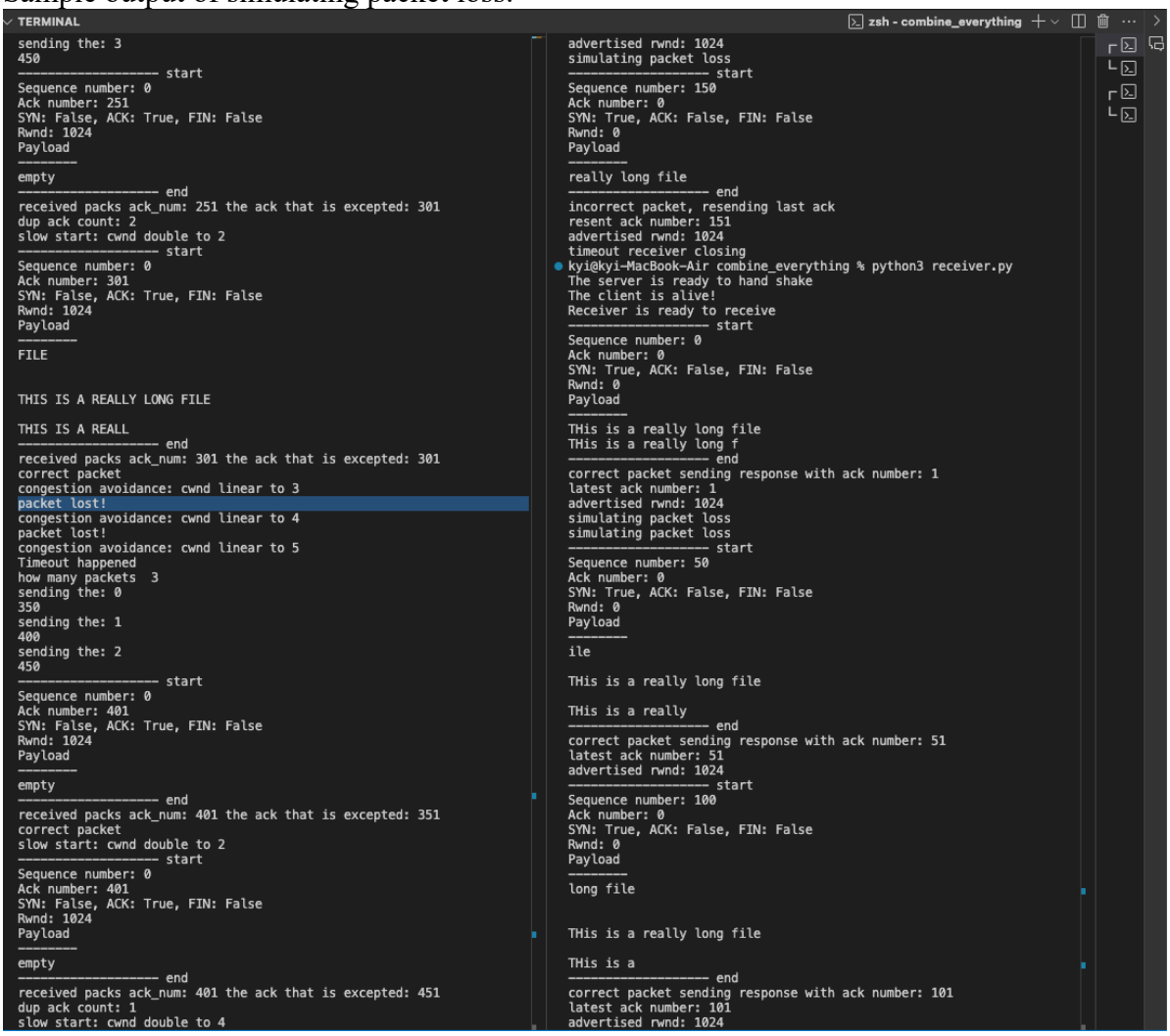
- Sequence number: For tracking order.
- Acknowledgment number: For reliable communication.
- Flags: SYN, ACK, FIN flags for connection control.
- Payload: the data being transmitted. We will get the payload data from a text file.
- Receiver window size (rwnd): Advertised by the receiver for flow control.

Step 2: Analyze Traffic

Simulate packet loss:
Where to do it: in the function that sends packets, before sending packets to the receiver.
How to do it: we used a random function to simulate packet loss by randomly decide whether or not to drop a packet. We set the loss rate is 25%, if generating a random number that is less than 0.25, we will drop this packet. If the generated number is greater than 0.25, than we will send this packet over network.

```
if random.random() > 0.25: # simulate packets being dropped on the way back
    response_packet = Packet.from_bytes_to(ack_packet)
    print(response_packet)
    print('received packs ack_num:', response_packet.ack_num, 'the ack that is excepted:', sndpkt[0].sequence_num+1)

    # if it is a dup ack, increase the counter
    if response_packet.ack_num == prev_ack_num:
        dupAckCount+=1
        print(f"dup ack count: {dupAckCount}")
    else: # if it is a new ack, reset the new ack counter
        dupAckCount=0

    # TCP tahoe, 3 dup acks, reset the cwnd, resend pkts
    if dupAckCount==3:
        # retransimit
        for i, pkt in enumerate(sndpkt):
            print('sending the:', i)
            print(pkt.sequence_num)
            senderSocket.sendto(pkt.change_to_bytes(),(receiverName, receiverPort))
        cwnd = 1
        ssthresh = max(cwnd/2, 1)
        dupAckCount = 0
        set_timer = time.time()
    else: # pop out the acked packet
        for i, pkt in enumerate(sndpkt):
            if (response_packet.ack_num == pkt.sequence_num+1): # checking if it the correct
                print('correct packet')
                sndpkt = sndpkt[i+1:] # pop out the acked packet
                set_timer = time.time()
                break

    prev_ack_num = response_packet.ack_num
else:
    print("packet lost!")
```

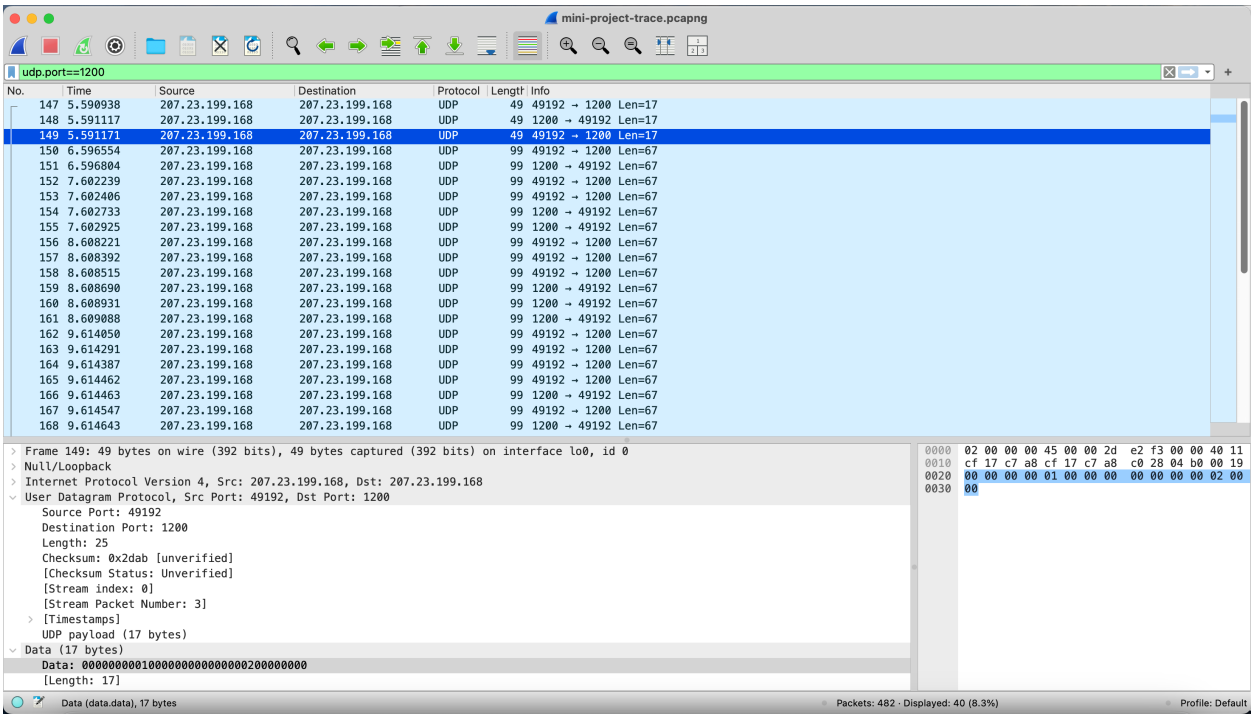
Sample output of simulating packet loss:



Proof of connection-oriented:

In Wireshark:

The first 3 packets are the packets being sent in the handshake part, these packets only have 17 bytes (header) and no payload. The first packet is from sender and includes TCP SYN message, the second packet is from receiver and includes SYNACK, and the third packet is from sender and includes ACK number.



The screenshot shows a connection-oriented protocol through the exchange of ACK numbers during the handshake process. The sender sends a SYN packet with a sequence number, and the receiver replies with an ACK (SYNACK), acknowledging the client's sequence number and sending its own. Then, sender sends an ACK to acknowledge SYNACK. This back-and-forth of sequence and acknowledgment numbers confirms reliable communication, typical of connection-oriented protocols like TCP.

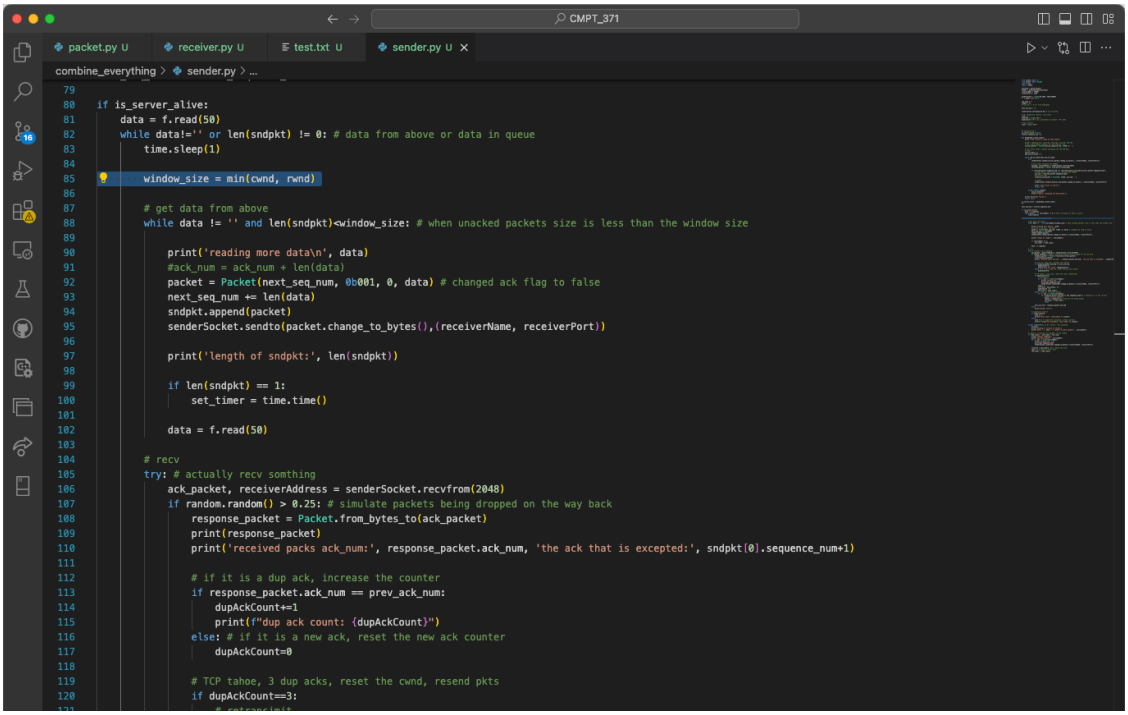
```
handshaking > sender.py > ...
15  # syn, ack, fin
16
17  def handshake_client_side():
18      print ("The client is ready to hand shake")
19
20      #order: sequence_num, syn_flag, ack_flag, ack_num, payload
21      # the client makes a packet, set SYN flag to true
22      initial_packet = Packet(initial_sequence_num, 0b100, 0, '')
23
24      # the client sends a packet including the TCP SYN msg
25      # timer
26      num_of_sends = 1
27      max_num_of_sends = 4
28
29      while num_of_sends<=max_num_of_sends:
30          try:
31              print("The initial sequence number from sender is: ", initial_sequence_num)
32              clientSocket.sendto(initial_packet.change_to_bytes(), (serverName, serverPort))
33
34              # the client receives the SYNACK
35              message, serverAddress = clientSocket.recvfrom(2048)
36              received_packet = Packet.from_bytes_to(message)
37
38              if received_packet.flags==0b110 and received_packet.ack_num==initial_packet.sequence_num+1:
39                  # sends back a ACK indicates it receives the SYNACK
40                  print("The SYN and ACK flags are true, received SYNACK.")
41                  ack_num = received_packet.sequence_num+1
42                  print("The SYNACK number is: ", ack_num)
43                  receiver_ack_packet = Packet(0, 0b010, ack_num, '')
44
PROBLEMS  OUTPUT  DEBUG CONSOLE  TERMINAL  PORTS
TERMINAL                                              zsh - handshaking
kyi@kyi-MacBook-Air handshaking % python3 receiver.py
The server is ready to hand shake
The initial sequence number from receiver is:  10000
The SYN flag is true, receiver starts to establish connection
The ACK number for SYN message is:  20001
The ACK flag is true, received ACK

The received ACK number is:  10001
The client is alive!

kyi@kyi-MacBook-Air handshaking % python3 sender.py
The client is ready to hand shake
The initial sequence number from sender is:  20000
The SYN and ACK flags are true, received SYNACK.
The SYNACK number is:  10001
The server is alive!
kyi@kyi-MacBook-Air handshaking %
```

Proof of flow control:

In receiver: set buffer size and used buffer size, put rwnd in the packet header and send it to sender:

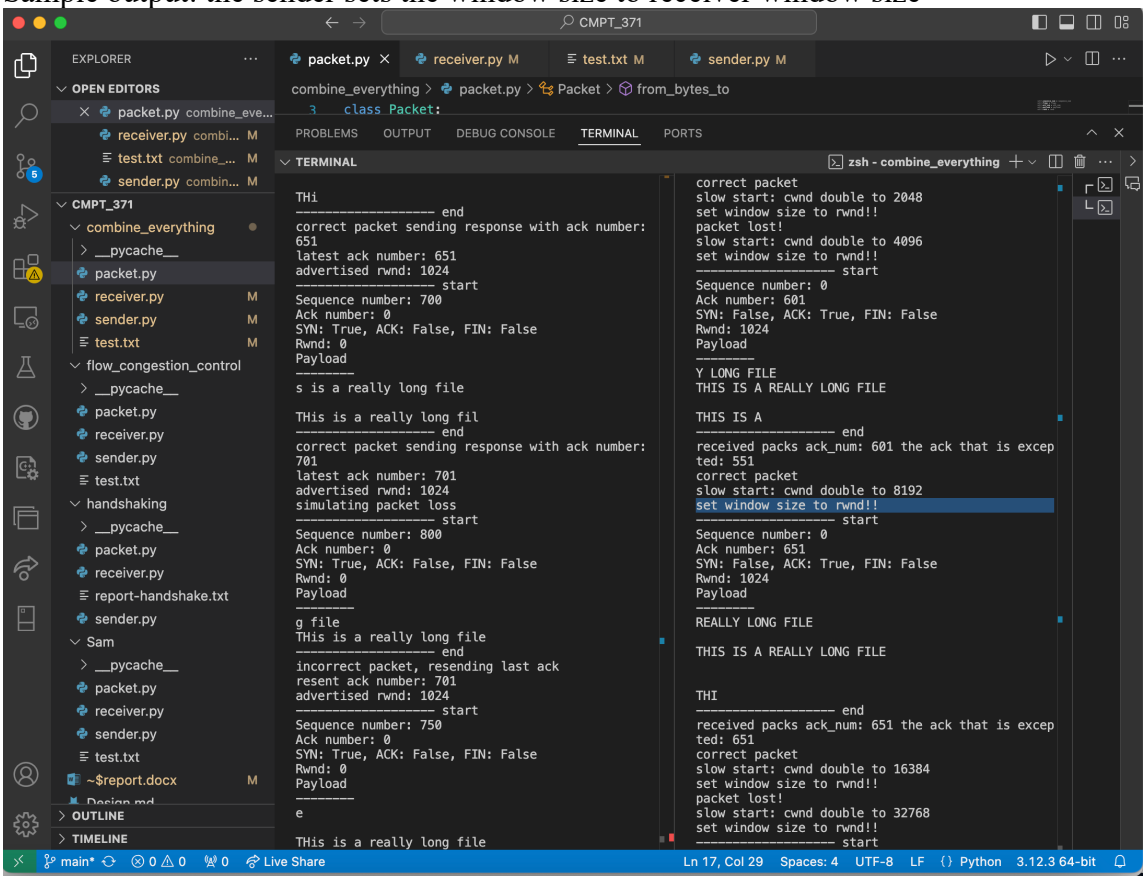
```
CMPT_371
packet.py U  receiver.py U  test.txt U  sender.py U
combine_everything > receiver.py > buffer_size
    # flow control
    buffer_size = 1024
    buffer_used = 0

    flag = True
    while flag:
        try:
            message, senderAddress = receiverSocket.recvfrom(2048)
        except:
            print('timeout receiver closing')
            break

        packet = Packet.from_bytes_to(message)

        if random.random() > 0.25: # pretend to receive packet
            print(packet)

            # check if there is enough buffer
            if buffer_used+len(packet.payload)<=buffer_size:
                if nextseqnum == packet.sequence_num:
                    # if it is he next packet respond with the correct ack
                    # else respond with the previous ack and keep doing that
                    # until the right packet is sent
                    payload = packet.payload
                    sndpkt = Packet(0, 0b010, nextseqnum+1, packet.payload.upper())
                    print('correct packet sending response with ack number:', nextseqnum+1)
                    nextseqnum += len(packet.payload)
                    last_ack_num = packet.sequence_num+1
                    print('latest ack number:', last_ack_num)
                else:
                    print('incorrect packet, resending last ack')
                    sndpkt = Packet(0, 0b010, last_ack_num, 'empty')
                    print('resent ack number:', sndpkt.ack_num)
            else:
                print("buffer full")
                sndpkt = Packet(0, 0b010, last_ack_num, 'empty')
            sndpkt.rwnd = buffer_size-buffer_used
            print('advertised rwnd:', sndpkt.rwnd)
            receiverSocket.sendto(sndpkt.change_to_bytes(),senderAddress)
        else:
            print('simulating packet loss')
```

In sender: get the rwnd from the receiver, set the window size to rwnd or cwnd (smaller one), so that the sender won't overwhelm receiver:



```
if is_server_alive:
    data = f.read(50)
    while data!='' or len(sndpkt) != 0: # data from above or data in queue
        time.sleep(1)

        window_size = min(cwnd, rwnd)

        # get data from above
        while data != '' and len(sndpkt)<window_size: # when unacked packets size is less than the window size

            print('reading more data\n', data)
            #ack_num = ack_num + len(data)
            packet = Packet(next_seq_num, 0b001, 0, data) # changed ack flag to false
            next_seq_num += len(data)
            sndpkt.append(packet)
            senderSocket.sendto(packet.change_to_bytes(),(receiverName, receiverPort))

            print('length of sndpkt:', len(sndpkt))

            if len(sndpkt) == 1:
                set_timer = time.time()

            data = f.read(50)

        # recv
        try: # actually recv somthing
            ack_packet, receiverAddress = senderSocket.recvfrom(2048)
            if random.random() > 0.25: # simulate packets being dropped on the way back
                response_packet = Packet.from_bytes_to(ack_packet)
                print(response_packet)
                print('received packs ack_num:', response_packet.ack_num, 'the ack that is excepted:', sndpkt[0].sequence_num+1)

                # if it is a dup ack, increase the counter
                if response_packet.ack_num == prev_ack_num:
                    dupAckCount+=1
                    print(f"dup ack count: {dupAckCount}")
                else: # if it is a new ack, reset the new ack counter
                    dupAckCount=0

                # TCP tahoe, 3 dup acks, reset the cwnd, resend pkts
                if dupAckCount==3:
```

Sample output: the sender sets the window size to receiver window size

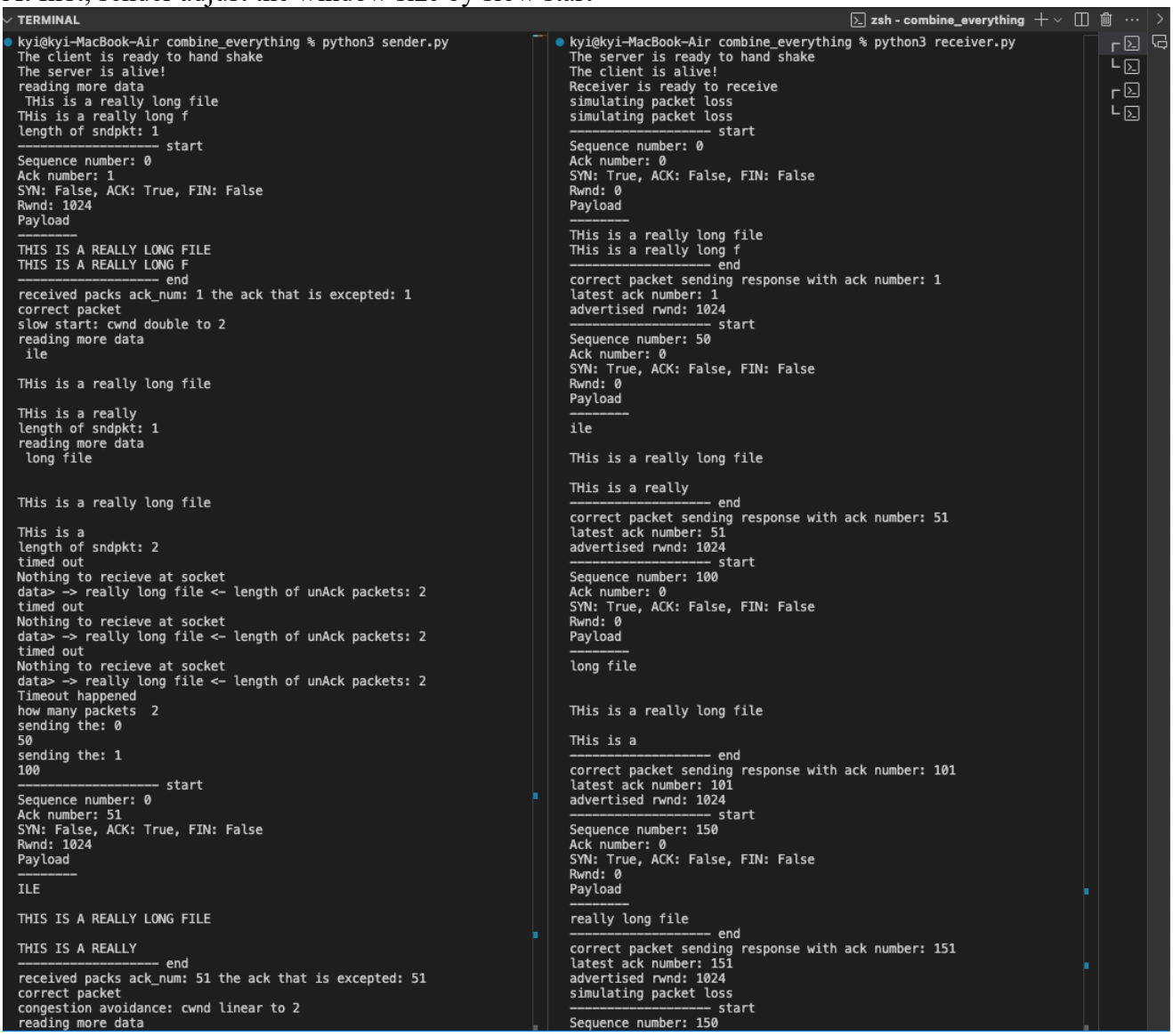


```
THi
-------------------- end
correct packet sending response with ack number:
651
latest ack number: 651
advertised rwnd: 1024
-------------------- start
Sequence number: 700
Ack number: 0
SYN: True, ACK: False, FIN: False
Rwnd: 0
Payload
---------
s is a really long file

THis is a really long fil
-------------------- end
correct packet sending response with ack number:
701
latest ack number: 701
advertised rwnd: 1024
simulating packet loss
-------------------- start
Sequence number: 800
Ack number: 0
SYN: True, ACK: False, FIN: False
Rwnd: 0
Payload
---------
g file
THis is a really long file
-------------------- end
incorrect packet, resending last ack
resent ack number: 701
advertised rwnd: 1024
-------------------- start
Sequence number: 750
Ack number: 0
SYN: True, ACK: False, FIN: False
Rwnd: 0
Payload
---------
e

THis is a really long file
```

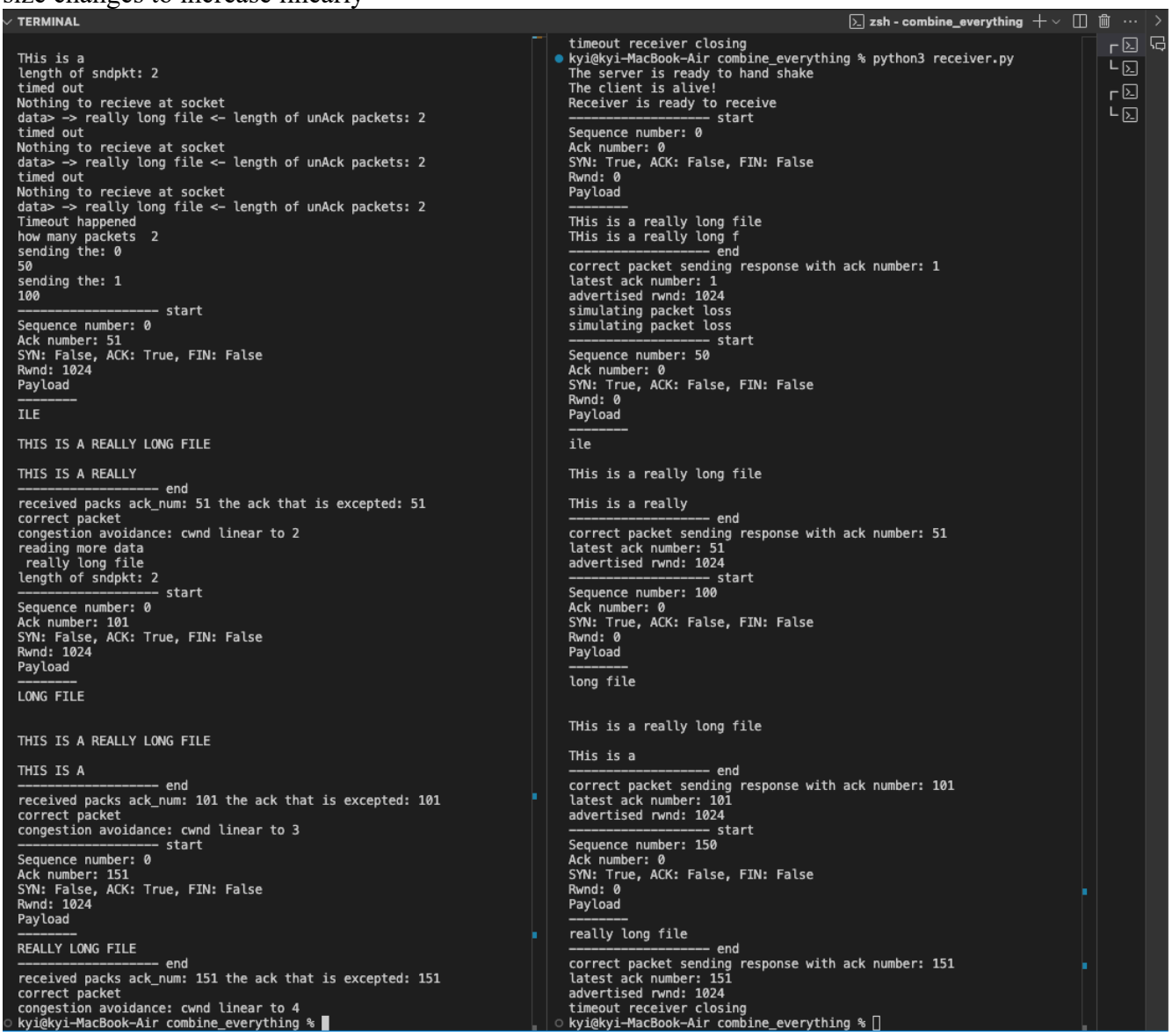
```
correct packet
slow start: cwnd double to 2048
set window size to rwnd!!
packet lost!
slow start: cwnd double to 4096
set window size to rwnd!!
-------------------- start
Sequence number: 0
Ack number: 601
SYN: False, ACK: True, FIN: False
Rwnd: 1024
Payload
---------
Y LONG FILE
THIS IS A REALLY LONG FILE

THIS IS A
-------------------- end
received packs ack_num: 601 the ack that is excep
ted: 551
correct packet
slow start: cwnd double to 8192
set window size to rwnd!!
-------------------- start
Sequence number: 0
Ack number: 651
SYN: False, ACK: True, FIN: False
Rwnd: 1024
Payload
---------
REALLY LONG FILE

THIS IS A REALLY LONG FILE


THI
-------------------- end
received packs ack_num: 651 the ack that is excep
ted: 651
correct packet
slow start: cwnd double to 16384
set window size to rwnd!!
packet lost!
slow start: cwnd double to 32768
set window size to rwnd!!
-------------------- start
```

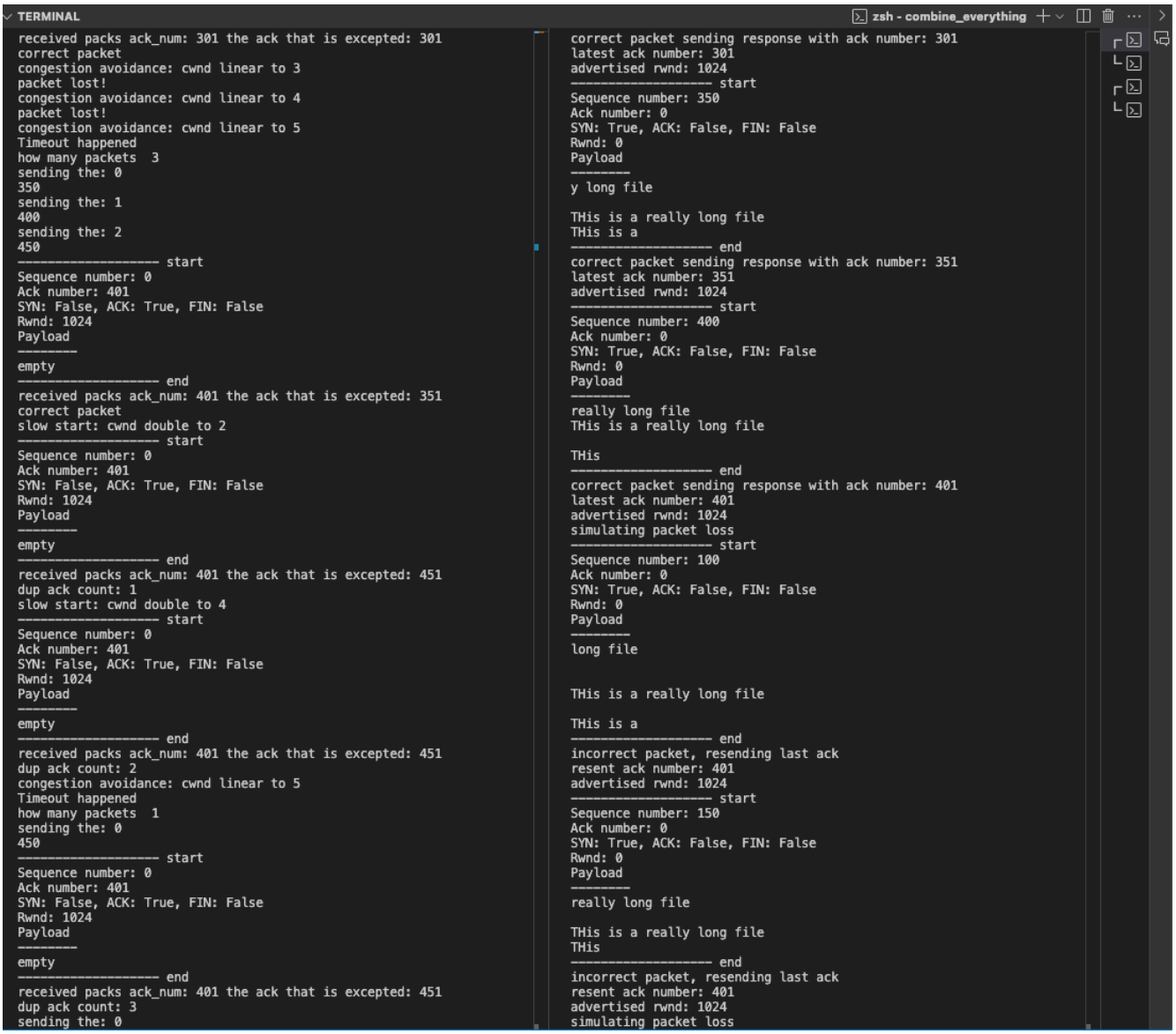
Proof of congestion control:

At first, sender adjust the window size by slow start

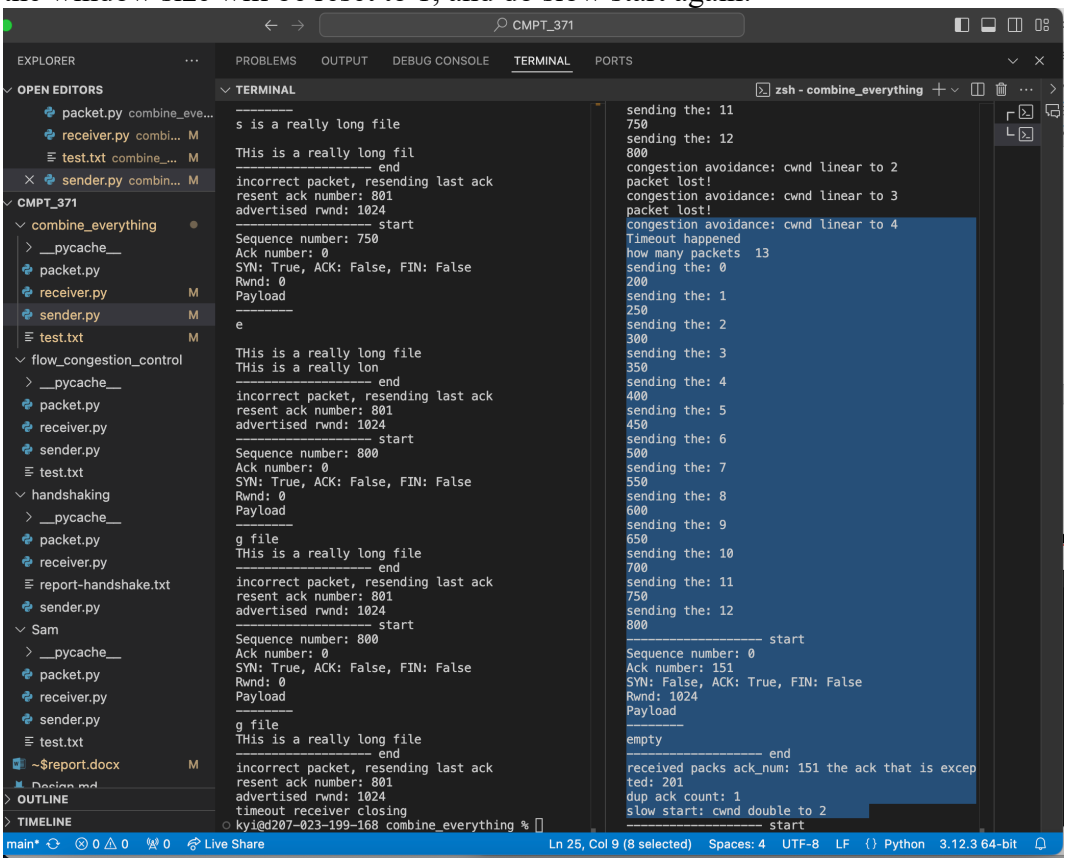


When the window size is greater or equal to ssthresh, the congestion avoidance happens, window size changes to increase linearly




TERMINAL
zsh - combine_everything
received packs ack_num: 301 the ack that is excepted: 301
correct packet
congestion avoidance: cwnd linear to 3
packet lost!
congestion avoidance: cwnd linear to 4
packet lost!
congestion avoidance: cwnd linear to 5
Timeout happened
how many packets 3
sending the: 0
350
sending the: 1
400
sending the: 2
450
---------- start
Sequence number: 0
Ack number: 401
SYN: False, ACK: True, FIN: False
Rwnd: 1024
Payload
----
empty
---------- end
received packs ack_num: 401 the ack that is excepted: 351
correct packet
slow start: cwnd double to 2
---------- start
Sequence number: 0
Ack number: 401
SYN: False, ACK: True, FIN: False
Rwnd: 1024
Payload
----
empty
---------- end
received packs ack_num: 401 the ack that is excepted: 451
dup ack count: 1
slow start: cwnd double to 4
---------- start
Sequence number: 0
Ack number: 401
SYN: False, ACK: True, FIN: False
Rwnd: 1024
Payload
----
empty
---------- end
received packs ack_num: 401 the ack that is excepted: 451
dup ack count: 2
congestion avoidance: cwnd linear to 5
Timeout happened
how many packets 1
sending the: 0
450
---------- start
Sequence number: 0
Ack number: 401
SYN: False, ACK: True, FIN: False
Rwnd: 1024
Payload
----
empty
---------- end
received packs ack_num: 401 the ack that is excepted: 451
dup ack count: 3
sending the: 0
correct packet sending response with ack number: 301
latest ack number: 301
advertised rwnd: 1024
---------- start
Sequence number: 350
Ack number: 0
SYN: True, ACK: False, FIN: False
Rwnd: 0
Payload
----
y long file
THis is a really long file
THis is a
---------- end
correct packet sending response with ack number: 351
latest ack number: 351
advertised rwnd: 1024
---------- start
Sequence number: 400
Ack number: 0
SYN: True, ACK: False, FIN: False
Rwnd: 0
Payload
----
really long file
THis is a really long file
THis
---------- end
correct packet sending response with ack number: 401
latest ack number: 401
advertised rwnd: 1024
simulating packet loss
---------- start
Sequence number: 100
Ack number: 0
SYN: True, ACK: False, FIN: False
Rwnd: 0
Payload
----
long file
THis is a really long file
THis is a
---------- end
incorrect packet, resending last ack
resent ack number: 401
advertised rwnd: 1024
---------- start
Sequence number: 150
Ack number: 0
SYN: True, ACK: False, FIN: False
Rwnd: 0
Payload
----
really long file
THis is a really long file
THis
---------- end
incorrect packet, resending last ack
resent ack number: 401
advertised rwnd: 1024
simulating packet loss

When timeout or sender receive duplicate acks for 3 times, the ssthresh will be halved down, and the window size will be reset to 1, and do slow start again:



In Wireshark, it shows the timeout happens, it did congestion control: