Why MapReduce?

CMPT 732, Fall 2017

Why did MapReduce get created the way it was?

Why is the Hadoop cluster infrastructure (YARN, HDFS) structured the way it is?

This is a good time for a little context...

MapReduce History

The original publication: <u>MapReduce: Simplified Data Processing on Large Clusters</u>, 2004.

Highlights:

- "Many real world tasks are expressible in this model, as shown in the paper."
- "Programs written in this functional style are automatically parallelized and executed on a large cluster of commodity machines. The run-time system takes care of the details of partitioning..., scheduling..., failures..., communication."
- "Our abstraction is inspired by the *map* and *reduce* primitives present in Lisp and many other functional languages. We realized that most of our computation involved applying a *map* operation... and then applying a *reduce* operation..."
- "map: (k1,v1) → list(k2,v2) reduce: (k2,list(v2)) → list(v2)"
- "A cluster consists of hundreds or thousands of machines, and therefore machine failures are common."

The Hadoop implementation of MapReduce (that we have been using) was inspired by this paper. ... and it seems to be a fairly faithful implementation.

MapReduce lets do computation in parallel in a *very* scalable way. Compared to other options at the time, it was very easy and flexible.

But now it seems very restrictive.

How would you implement an iterative algorithm (gradient descent, logistic regression, PageRank, etc) with MapReduce?

Probably something like:

```
bool converged = False;
int iter = 0
while (!converged ) {
    Job job = Job.getInstance(conf, "gradient descent");
    :
    TextInputFormat.addInputPath(job, new Path('tempfile' + (iter-1)));
    TextOutputFormat.setOutputPath(job, new Path('tempfile' + iter));
    job.waitForCompletion(true);
    converged = ...;
    iter++;
}
```

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Fault Tolerance

Computers fail. In a cluster, something failing starts to be more and more likely...

If each machine fails with probability 0.1%, then...

| Machines | Failure | |
|------------|--------------------|--|
| in Cluster | Probability | |
| 1 | 0.1% | |
| 10 | 1.0% | |
| 100 | 9.5% | |
| 1000 | 63.2% | |
| 10,000 | 99.996% | |

Lesson: In a cluster, failure is a real possibilty we have to deal with.

What happens if a cluster node disappears?

HDFS handles faults by keeping multiple copies of each block. Number of copies is controled by dfs.replication: default 3.

This also makes it easier to schedule jobs local to data: there are more options.

What happens if one of the nodes holding a particular block fails? You'll have to find on on Assignment 5A.

What if all of them disappear simultaneously? You can try that too.

Spoilers: (1) it heals; (2) everything else keeps working.

For YARN, what happens if some work is happening and the node dies? The task tight have done nothing. It might have written incomplete output (to a file/database/wherever).

Again, you can experiment in 5A.

Remember that our jobs are broken up into smaller pieces: map or reduce tasks, Spark tasks on one partition.

If the task writing part-00005 fails, then it should be safe to delete it and start that task on another node.

But what if our output is a database or something else? What if getting the prerequisites of part-00005 is a huge amount of work?

These depend on the details: you should occasionally ask yourself "what if part of this fails?"

Where Your Data Might Be

We have seen several places our data might be stored (permanently or temporarily) while we're working on it.

There are massive speed differences between them: having some (approximate) idea of the tradeoffs is important.

| Tech | Size | Latency (cycles) | Throughput (B/s) | |
|----------------|-------------|-------------------|--------------------------------|--------|
| Registers | few kB | 1 | | |
| L1 cache | <100 kB | 4 <u>*</u> | 250 G-1 T <u>**</u> | |
| RAM | 10s GB | 50-100 <u>**</u> | 20 G-50 G <u>*</u> | |
| SSD | 100s GB | 10 ⁶ * | 500M (per disk) <u>*</u> | |
| Spinning HD | TB | 10 ⁷ * | 100 M (per disk) <u>*</u> | |
| Network | PB | $10^7 + disk$ | 120 M (shared?) <u>*</u> | |
| (Good consumer | -grade hard | lware, ~2017. E | verything approximate, obvious | ısly.) |

MapReduce's habit of putting map output onto disk seems like an obvious bottleneck.

Keeping data in memory should be much faster (if it fits). I sure would be nice to have a tool that did that...

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