COMPARISON OF MILP SOLVERS WITH USING RELATED-KEY DIFFERENTIAL ATTACK

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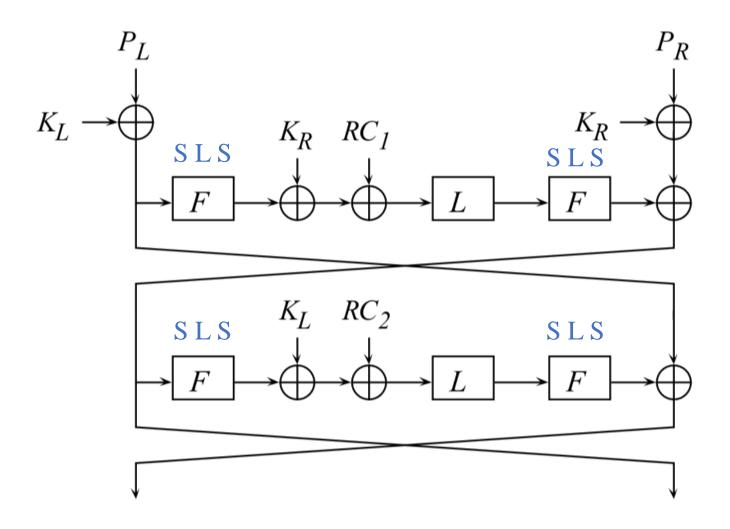


Overview



- 1. ITUbee Block Cipher
- 2. MILP Model of Related-Key Differential Attack
- 3. GLPK Solver
- 4. HiGHS Solver
- 5. Gurobi Solver
- 6. CPLEX Solver
- 7. Comparison





ITUbee Block Cipher



The encryption operation proceeds as follows:

- 1. $X_1 \leftarrow P_L \oplus K_L$ and $X_0 \leftarrow P_R \oplus K_R$.
- 2. for i = 1...20 do
 - (a) if $i \in \{1, 3, ..., 19\}$
 - i. $RK \leftarrow K_R$.
 - (b) else
 - i. $RK \leftarrow K_L$.
 - (c) $X_{i+1} \leftarrow X_{i-1} \oplus F(L(RK \oplus RC_i \oplus F(X_i)))$. Note that 16-bit round constant is added to the rightmost 16 bits.
- 3. $C_L \leftarrow X_{20} \oplus K_R$ and $C_R \leftarrow X_{21} \oplus K_L$.

ITUbee Block Cipher



- F(X) = S(L(S(X))).
- -S(a||b||c||d||e) = s[a]||s[b]||s[c]||s[d]||s[e]| where a,b,c,d,e are 8-bit values and s is the S-box used in AES [10].
- $-L(a||b||c||d||e) = (e \oplus a \oplus b)||(a \oplus b \oplus c)||(b \oplus c \oplus d)||(c \oplus d \oplus e)||(d \oplus e \oplus a).$

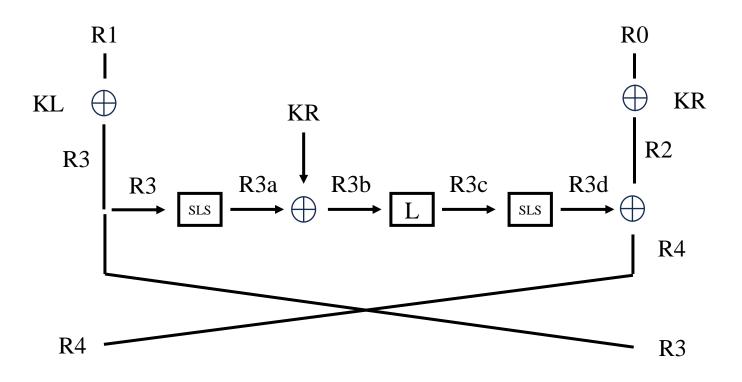
ITUBEE offers 80-bit security. Also, the cipher provides the same security level against attacks in related key models as well as in single key attack models. In addition, the cipher has no weak keys.



4.3 Related Key Differential Attacks

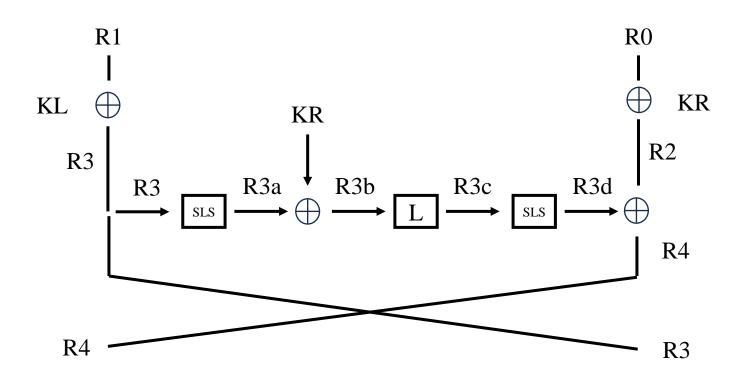
In the related key attack model, the adversary is able to collect plaintext and ciphertext pairs under related keys. In our algorithm, we divide the master key into two parts K_L and K_R and we use these parts between F functions in the rounds. It is trivial to see that when there is a difference in the key used between F functions, then at least one of the F functions will be active. In the best case for the adversary, the difference will be only one part of the key K_L or K_R . As a result, in two consecutive rounds there exists at least one active F function in the case of the related key attack. The probabilities of differentials for the F function is less than 2^{-17} as given in the Section 4.1. Thus, for 10 consecutive rounds the probabilities will be less than $(2^{-17})^5 = 2^{-85}$ which is not usable in an attack.





Minimize: R3 + R3a + R3c + R3d





Add Constraint : $KL + KR \ge 1$



XOR:

$$\begin{split} x_{in_1}^{\oplus} + x_{in_2}^{\oplus} + x_{out}^{\oplus} &\geq 2d^{\oplus} \quad , \\ d^{\oplus} &\geq x_{in_1}^{\oplus} \quad , \\ d^{\oplus} &\geq x_{in_2}^{\oplus} \quad , \\ d^{\oplus} &\geq x_{out}^{\oplus} \quad . \end{split}$$

```
def xor(p,in1,in2,out,dummy):
    for i in range(len(in1)):
        p.add_constraint(in1[i]+in2[i]+out[i]-2*dummy[i] >=0)
        p.add_constraint(dummy[i]- in1[i] >= 0)
        p.add_constraint(dummy[i]- in2[i] >= 0)
        p.add_constraint(dummy[i]- out[i] >= 0)
```



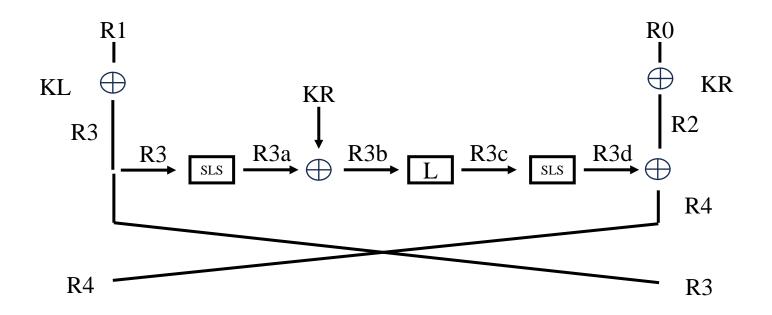
L Function:

```
pp = Polyhedron(vertices=[
[ 0,0,0,0,0,0,0,0,0,0],
[ 0,0,0,0,1,1,0,0,1,1 ],
[ 0,0,0,1,0,0,0,1,1,1 ],
[ 0,0,0,1,1,1,0,1,0,0 ],
[ 0,0,0,1,1,1,0,1,1,1 ],
[ 0,0,1,0,0,0,1,1,1,0 ],
[ 0,0,1,0,1,1,1,1,0,1 ],
[ 0,0,1,0,1,1,1,1,1,1,
[ 0,0,1,1,0,0,1,0,0,1 ],
[ 0,0,1,1,1,1,1,0,1,0 ],
[ 0,0,1,1,1,1,1,0,1,1 ],
[ 0,0,1,1,0,0,1,1,1,1 ],
[ 0,0,1,1,1,1,1,1,1,1,],
[ 0,0,1,1,1,1,1,1,1,0 ],
```

MILP Model of Related-Key Differential Attack



```
KL-KR
                     {0: 0.0, 1: 1.0, 2: 0.0, 3: 1.0, 4: 0.0} {0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0}
R1-R0
                     {0: 0.0, 1: 1.0, 2: 0.0, 3: 1.0, 4: 0.0} {0: 1.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 1.0}
                     {0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0} {0: 1.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 1.0}
R3-R2
R3a-R3b-R3c-R3d
                     {0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0} {0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0} {0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0} {0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0}
                     {0: 1.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 1.0} {0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0}
R4-R3
                     {0: 0.0, 1: 1.0, 2: 0.0, 3: 1.0, 4: 0.0} {0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0} {0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0} {0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0} {0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0}
R4a-R4b-R4c-R4d
R5-R4
                     {0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0} {0: 1.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 1.0}
                     {0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0} {0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0} {0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0} {0: 0.0, 4: 0.0} {0: 0.0, 4: 0.0} {0: 0.0, 4: 0.0} {0: 0.0, 4: 0.0}
R5a-R5b-R5c-R5d
                     {0: 1.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 1.0} {0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0}
R6-R5
R6a-R6b-R6c-R6d
                     {0: 0.0, 1: 1.0, 2: 0.0, 3: 1.0, 4: 0.0} {0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0} {0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0} {0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0} {0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0}
                     {0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0} {0: 1.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 1.0}
R7-R6
                     {0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0} {0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0} {0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0} {0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0}
R7a-R7b-R7c-R7d
R8-R7
                     {0: 1.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 1.0} {0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0}
R8a-R8b-R8c-R8d
                     {0: 0.0, 1: 1.0, 2: 0.0, 3: 1.0, 4: 0.0} {0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0} {0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0} {0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0}
R9-R8
                     {0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0} {0: 1.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 1.0}
                     {0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0} {0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0} {0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0} {0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0} {0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0}
R9a-R9b-R9c-R9d
R10-R9
                     {0: 1.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 1.0} {0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0}
R10a-R10b-R10c-R10d {0: 0.0, 1: 1.0, 2: 0.0, 3: 1.0, 4: 0.0} {0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0} {0: 0.0, 1: 0.0, 2: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0} {0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0}
R11-R10
                     {0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0} {0: 1.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 1.0}
```



```
Toplam aktif s-kutusu: 16.0
KL-KR {0: 0.0, 1: 1.0, 2: 0.0, 3: 1.0, 4: 0.0} {0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0}
R1-R0 {0: 0.0, 1: 1.0, 2: 0.0, 3: 1.0, 4: 0.0} {0: 1.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 1.0}
R3-R2 {0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0} {0: 1.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 1.0}
R3a-R3b-R3c-R3d {0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0} {0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0}
R4-R3 {0: 1.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 1.0} {0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0}
```



GLPK (GNU Linear Programming Kit)



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Introduction to GLPK

The GLPK (<u>GNU Linear Programming Kit</u>) package is intended for solving large-scale linear programming the form of a callable library.

GLPK supports the GNU MathProg modeling language, which is a subset of the AMPL language.

The GLPK package includes the following main components:

- · primal and dual simplex methods
- · primal-dual interior-point method
- · branch-and-cut method
- · translator for GNU MathProg
- application program interface (API)
- stand-alone LP/MIP solver



HiGHS - high performance software for linear optimization

Open source serial and parallel solvers for large-scale sparse linear programming (LP),

mixed-integer programming (MIP), and quadratic programming (QP) models

Solvers

LP

HiGHS has implementations of the three main solution techniques for LP. HiGHS will choose the most appropriate technique for a given problem, but this can be over-ridden by setting the option solver.

Simplex

HiGHS has efficient implementations of both the primal and dual simplex methods, although the dual simplex solver is likely to be faster and is more robust, so is used by default. The novel features of the dual simplex solver are described in



HiGHS - high performance software for linear optimization

Open source serial and parallel solvers for large-scale sparse linear programming (LP),

mixed-integer programming (MIP), and quadratic programming (QP) models

Interior point o

HiGHS has one interior point (IPM) solver based on the preconditioned conjugate gradient method, as discussed in

Implementation of an interior point method with basis preconditioning, Mathematical Programming Computation, 12, 603-635, 2020. DOI: 10.1007/s12532-020-00181-8.

This solver is serial. An interior point solver based on direct factorization is being developed.

Setting the option solver to "ipm" forces the IPM solver to be used

Primal-dual hybrid gradient method

HiGHS includes the cuPDLP-C primal-dual hybrid gradient method for LP (PDLP). Currently this only runs on CPU, so it is unlikely to be competitive with the HiGHS interior point or simplex solvers. Enabling HiGHS to run PDLP on a GPU is work in progress.



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PROBLEM	SOLVER	TIME(sec)	SPEED
Related-Key Differential	GLPK	257	x (base)
Cryptanalysis of ITUbee	HiGHS	231	1.1x
algoirthm (8 round)			
48649 constraints	Gurobi	37	6.9x
320 variables	CPLEX	16	16x

CPU: 11th Gen Intel(R) Core(TM) i7-1165G7 @ 2.80GHz



PROBLEM	SOLVER	TIME(sec)	SPEED
Related-Key Differential	GLPK	1610	x (base)
Cryptanalysis of ITUbee	HiGHS	321	5x
algoirthm (10 round)	Gurobi	34	47x
60801 constraints	Guiobi	J -1	777
390 variables	CPLEX	40	40x

CPU: 11th Gen Intel(R) Core(TM) i7-1165G7 @ 2.80GHz



PROBLEM	SOLVER	TIME(sec)	SPEED
Related-Key Differential	GLPK	4543	x (base)
Cryptanalysis of ITUbee	HiGHS	1780	2.5x
algoirthm (12 round)			
72953 constraints	Gurobi	48	94x
460 variables	CPLEX	53	85x

CPU: 11th Gen Intel(R) Core(TM) i7-1165G7 @ 2.80GHz

Comparison





CPU: 11th Gen Intel(R) Core(TM) i7-1165G7 @ 2.80GHz

Referances



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