USAGE OF MIXED INTEGER LINEAR PROGRAMMING IN CRYPTANALYSIS OF BLOCK CIPHERS

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Block Ciphers

Block ciphers designed to securely encrypt and decrypt fixed-size blocks of data by transforming plaintext into ciphertext using a symmetric key.
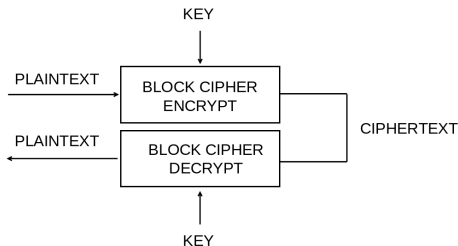

Figure 1: Block Cipher Encrytion and Decryption

Cryptanalysis

The main purpose of cryptanalysis:

- Prove that an algorithm is secure against known attacks in the open literature
- Find better complexities for those attacks and contribute to the security of the algorithm
- Break it ! (breaking an algorithm means obtaining the key with a lower complexity than the algorithm claims.)

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Mixed-Integer Linear Programming (MILP)

Goal: Optimize (maximize or minimize) a linear function.

$$c_1 x_1 + c_2 x_2 + \cdots + c_n x_n$$

Constraints:

$$a_{11} x_1 + a_{12} x_2 + \cdots + a_{1n} x_n \leq b_1$$

$$\vdots$$

$$a_{m1} x_1 + \cdots + a_{mn} x_n \leq b_m$$

$$x_1, x_2, \ldots, x_p \in \mathbb{Z}, \quad x_{p+1}, \ldots, x_n \geq 0$$

- x_i: decision variables
- c_i: objective function weights
- a_{ij}: constraint coefficients
- b_j: limits of constraints

A company produces two products: **A** and **B**. The goal: **maximize profit** under production limits.

Maximize profit:

$$Z = 40x + 30y$$

- Profit per unit of A $= 40$
- Profit per unit of B $= 30$

MILP Example: Constraints

Constraints:

- **Labor:** $2x + y \leq 100$ (100 hours available)
- **Machine:** $x + 3y \leq 90$ (90 hours available)
- **Demand:** $x \geq 20$ (min. 20 units of A)
- **Non-negativity:** $x, y \geq 0$

MILP example

We will use Sagemath and GLPK solver for modeling.

```
1 p.<x,y> = MixedIntegerLinearProgram(maximization=True,
      solver="GLPK")
2 p.set_objective(40*x[0] + 30*y[0])
3 p.add_constraint(2*x[0] + y[0] == 100)
4 p.add_constraint(x[0] + 3*y[0] == 90)
5 p.add_constraint(x[0] >= 20)
6 p.add_constraint(y[0] >= 0)
7 print(p.solve())
8 print(p.get_values(x, y))
```

- Maximum profit: 2160
- A : 42
- B : 16

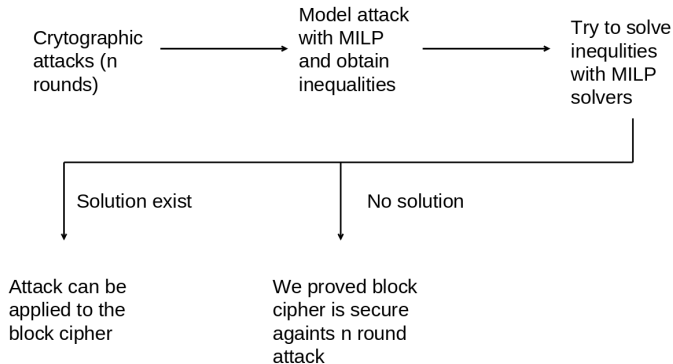# Cryptanalysis and MILP

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Differential Cryptanalysis

Idea: Study how input differences between two plaintexts affect output differences in ciphertexts.

Definition

For two plaintexts P and P':

$$\Delta P = P \oplus P'$$

Differential Cryptanalysis: Probability

- Input difference: ΔX
- Output difference: ΔY

In an ideal cipher: - For n-bit input,

$$\Pr[\Delta X \to \Delta Y] = \frac{1}{2^n}$$

In practice: - Some input differences lead to output differences with much higher probability.

Differential Characteristics

Figure 2: Concept of constructing differential characteristic

Differential Cryptanalysis

$\Delta X/\Delta Y$	0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F
0	16	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
1	0	0	0	0	0	0	0	0	4	4	4	4	0	0	0	0
2	0	4	0	4	0	4	4	0	0	0	0	0	0	0	0	0
3	0	0	0	0	0	0	0	0	2	2	2	2	2	2	2	2
4	0	0	4	0	0	0	2	2	0	0	0	4	2	2	0	0
5	0	0	4	0	0	0	2	2	0	0	4	0	2	2	0	0
6	0	2	0	2	2	0	0	2	2	0	2	0	0	2	2	0
7	0	2	0	2	2	0	0	2	0	2	0	2	2	0	0	2
8	0	0	0	0	4	4	0	0	0	0	0	0	2	2	2	2
9	0	0	0	0	4	4	0	0	0	0	0	0	2	2	2	2
A	0	0	0	0	0	4	4	0	2	2	2	2	0	0	0	0
B	0	4	0	4	0	0	0	0	0	0	0	0	2	2	2	2
C	0	0	4	0	0	0	2	2	4	0	0	0	0	0	2	2
D	0	0	4	0	0	0	2	2	0	4	0	0	0	0	2	2
E	0	2	0	2	2	0	0	2	0	2	0	2	0	2	2	0
F	0	2	0	2	2	0	0	2	2	0	2	0	2	0	0	2

Table 1: Differential Distribution Table (DDT) of PRINCE algorithm

Steps of Differential Cryptanalysis

1. **Select Differentials:** Use DDT to find $(\Delta X, \Delta Y)$ pairs with high probability.

2. **Collect Pairs:** Generate many plaintext pairs P, P' with $P \oplus P' = \Delta X$.

3. **Encrypt:** Process these pairs through the cipher.

4. **Analyze Outputs:** Check if observed ciphertext differences match expected ΔY.

5. **Key Recovery:** Use biases to guess round subkeys (esp. in the last rounds).
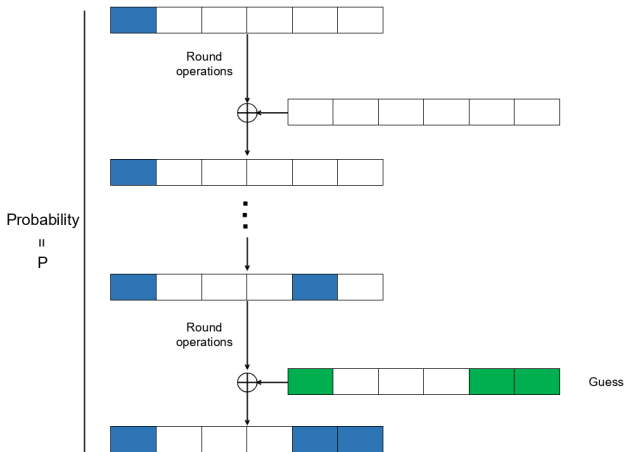
Figure 3: Differential Cryptanalysis

Adds differentials also to key sechedule.
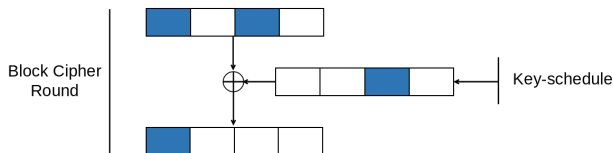
Figure 4: Related-key Differential Cryptanalysis

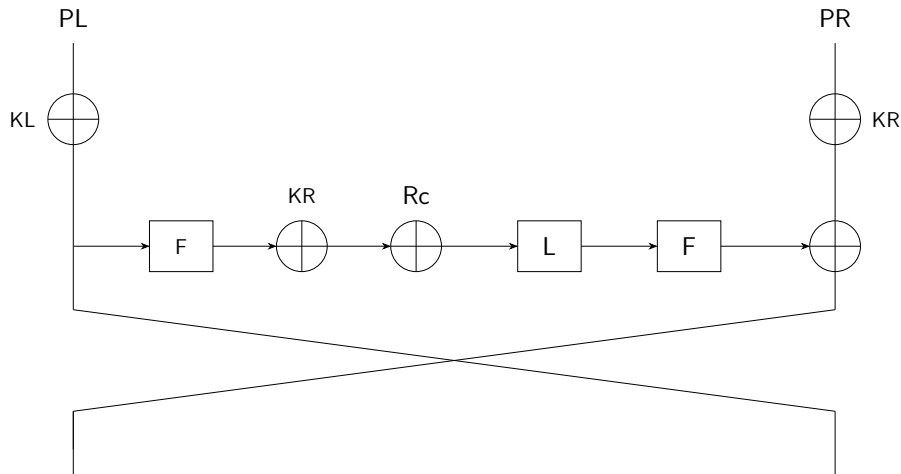# Table of Contents

Figure 5: ITUbee round structure

ITUbee Block Cipher

L Function

$$L(a, b, c, d, e) = (e \oplus a \oplus b, \; a \oplus b \oplus c, \; b \oplus c \oplus d, \; c \oplus d \oplus e, \; d \oplus e \oplus a)$$

F function

$$F(X) = S(L(S(X)))$$

$$S(a \parallel b \parallel c \parallel d \parallel e) = s[a] \parallel s[b] \parallel s[c] \parallel s[d] \parallel s[e]$$

ITUbee Block Cipher

Key Schedule

- 80-bit key is split into two 40-bit halves:

$$K = K_L \parallel K_R$$

- Odd rounds use K_R
- Even rounds use K_L
- No dedicated key schedule

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MILP Modelling of XOR Operation

Following [2], let x_{in1} and x_{in2} be the input differences, and x_{out} be the output difference of the XOR. The branch number is 2.

We introduce a binary variable d:

- $d = 0$ if $x_{in1} = x_{in2} = x_{out} = 0$
- $d = 1$ otherwise

The XOR is modelled with:

$$x_{in1} + x_{in2} + x_{out} \geq 2d$$
$$d \geq x_{in1}$$
$$d \geq x_{in2}$$
$$d \geq x_{out}$$

MILP Modelling of Linear Transformation

Similarly, for a linear transformation with:

- Inputs: $x_{in1}, x_{in2}, \ldots, x_{inM}$
- Outputs: $x_{out1}, x_{out2}, \ldots, x_{outM}$

Let B be the branch number and d the dummy variable (as in XOR modelling).

The model is:

$$x_{in1} + \cdots + x_{inM} + x_{out1} + \cdots + x_{outM} \geq B \cdot d$$

and

$$d \geq x_{in_i}, \quad d \geq x_{out_j} \quad \forall i, j$$

MILP Representation of S-boxes

Goal: minimize the number of **active S-boxes** to find the best differential trails.

Key idea:

- Active input ($\neq 0$) $\Rightarrow$ active output
- Passive input ($= 0$) $\Rightarrow$ passive output

We only track **activity** with a binary variable — no internal structure is modelled.

$L(a, b, c) = (a \oplus b, a \oplus c, b \oplus c)$ where a,b.c are 2 bit values

H-representation example

Input Differences	Output Differences
0 0 0	0 0 0
0 0 1	0 1 1
0 1 0	1 0 1
0 1 1	1 1 0
0 1 1	1 1 1
1 0 0	1 1 0
1 0 1	1 0 1
1 0 1	1 1 1
1 1 0	0 1 1
...	...
1 1 1	1 0 1
1 1 1	1 1 0
1 1 1	1 1 1

Table 2: Input and Output Differences of L function

H-representation example

No.	Inequality
1	$(0, 0, -1, 0, 0, 0)x + 1 \geq 0$
2	$(0, -1, 0, 0, 0, 0)x + 1 \geq 0$
3	$(-1, 0, 0, 0, 0, 0)x + 1 \geq 0$
4	$(0, 0, 0, -1, 0, 0)x + 1 \geq 0$
5	$(0, 0, 0, 0, -1, 0)x + 1 \geq 0$
6	$(0, 0, 0, 0, 0, -1)x + 1 \geq 0$
7	$(0, -1, 1, 0, 0, 1)x + 0 \geq 0$
8	$(0, 0, 0, -1, 1, 1)x + 0 \geq 0$
...	...
18	$(-1, 1, 0, 1, 0, 0)x + 0 \geq 0$
19	$(1, -1, -1, 1, 1, -1)x + 1 \geq 0$
20	$(1, -1, 0, 1, 0, 0)x + 0 \geq 0$
21	$(-1, 1, -1, 1, -1, 1)x + 1 \geq 0$
22	$(-1, -1, 1, -1, 1, 1)x + 1 \geq 0$

Table 3: H-representation of L function

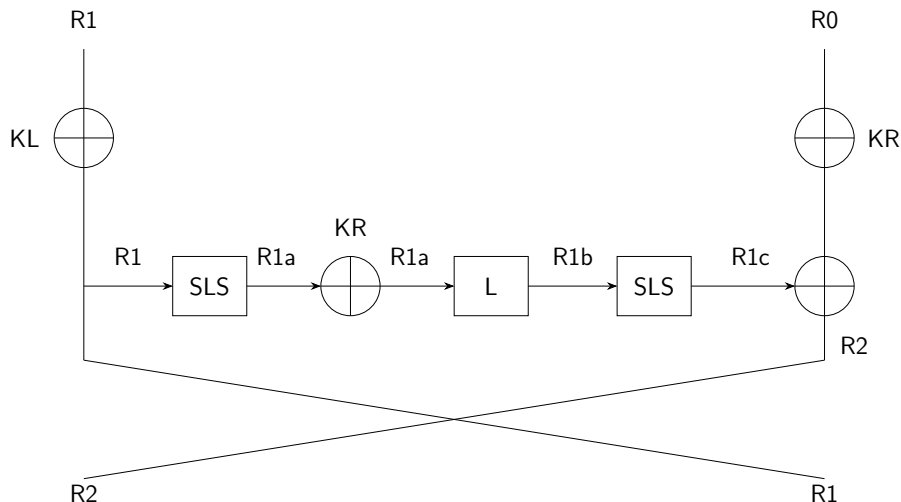# Differential cryptanalysis model with MILP

Figure 6: ITUbee MILP differential cryptanalysis sketch

- Number of Constrains: 18169
- Number of Variables: 85
- Minimum number of active s-box for 3 round: 16.0
- Best probability for s-box: 2^{-6}
- Best probability for 3 round differential characteristic will be $(2^{-6})^{16} = 2^{-96}$

which is not usable for differential attack.

Differential cryptanalysis model with MILP

Stage	Binary Index	Characteristic
1. round	R1	{0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 1.0}
input	R0	{0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0}
1. round	R1a	{0: 1.0, 1: 0.0, 2: 0.0, 3: 1.0, 4: 1.0}
middle values	R1b	{0: 0.0, 1: 1.0, 2: 1.0, 3: 0.0, 4: 1.0}
	R1c	{0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 1.0}
2. round	R2	{0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 1.0}
input	R1	{0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 1.0}
2. round	R2a	{0: 1.0, 1: 0.0, 2: 0.0, 3: 1.0, 4: 1.0}
middle values	R2b	{0: 0.0, 1: 1.0, 2: 1.0, 3: 0.0, 4: 1.0}
	R2c	{0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 1.0}
3. round	R3	{0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0}
input	R2	{0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 1.0}
3. round	R3a	{0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0}
middle values	R3b	{0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0}
	R3c	{0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0}
4. round	R4	{0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 1.0}
input	R3	{0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0}

Table 4: Best 3 round differential characteristic of ITUbee algorithm
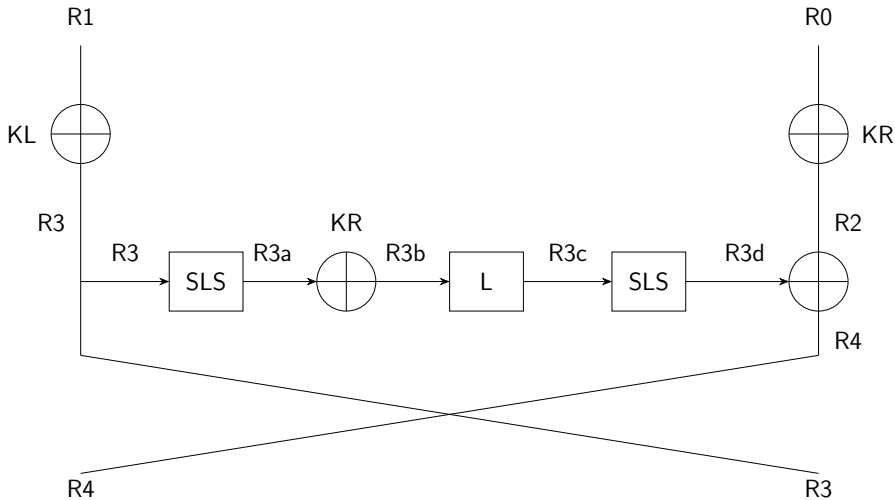
Figure 7: ITUbee MILP Related-Key Differential Cryptanalysis Sketch

- Number of Constrains: 48649
- Number of Variables: 320
- Minimum number of active s-box for 8 round: 16.0
- Best probability for s-box: 2^{-6}
- Best probability for 8 round differential characteristic will be $(2^{-6})^{16} = 2^{-96}$

which is not usable for related-key differential attack.

Stage	Binary Index	Characteristic
Key	KL	{0: 0.0, 1: 1.0, 2: 0.0, 3: 1.0, 4: 0.0}
	KR	{0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0}
1. round	R1	{0: 0.0, 1: 1.0, 2: 0.0, 3: 1.0, 4: 0.0}
input	R0	{0: 1.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 1.0}
1. round	R2	{0: 1.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 1.0}
middle values	R3-3a-3b-3c-3d	{0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0}
2. round	R4	{0: 1.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 1.0}
input	R3	{0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0}
2. round	R4a	{0: 0.0, 1: 1.0, 2: 0.0, 3: 1.0, 4: 0.0}
middle values	R4b-4c-4d	{0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0}
...	...	...
7. round	R9	{0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0}
input	R8	{0: 1.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 1.0}
7. round middle values	R9a-9b-9c-9d	{0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0}
8. round	R10	{0: 1.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 1.0}
input	R9	{0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0}
8. round	R10a	{0: 0.0, 1: 1.0, 2: 0.0, 3: 1.0, 4: 0.0}
middle values	R10b-10c-10d	{0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0}
9. round	R11	{0: 0.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 0.0}
input	R10	{0: 1.0, 1: 0.0, 2: 0.0, 3: 0.0, 4: 1.0}

Table 5: Best 8 round related-key differential characteristic of ITUbee algorithm

Our contribution

Stage	Differential	Linear	Related-key Differential
[1]	3 round	3 round	10 round
Our Result	3 round	3 round	**8 round**

Table 6: Result Comparison

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References

[1] Ferhat Karakoç, Hüseyin Demirci, and A Emre Harmancı. "ITUbee: a software oriented lightweight block cipher". In: *Lightweight Cryptography for Security and Privacy: Second International Workshop, LightSec 2013, Gebze, Turkey, May 6-7, 2013, Revised Selected Papers 2*. Springer. 2013, pp. 16–27.

[2] Nicky Mouha et al. "Differential and linear cryptanalysis using mixed-integer linear programming". In: *Information Security and Cryptology: 7th International Conference, Inscrypt 2011, Beijing, China, November 30–December 3, 2011. Revised Selected Papers 7*. Springer. 2012, pp. 57–76.