

OPERATING SYSTEM ASSIGNMENT - 4

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PART-B

Q6 Distributed Deadlock Detection:

Given fragments:

- $S_1 : P_1 \rightarrow P_2, P_3 \rightarrow P_4$
- $S_2 : P_2 \rightarrow P_5, P_5 \rightarrow P_6$
- $S_3 : P_6 \rightarrow P_1$

a) Global wait-for graph
combine edges (all):

$P_1 \rightarrow P_2$

$P_2 \rightarrow P_5$

$P_5 \rightarrow P_6$

$P_6 \rightarrow P_1$ (cycle)

$P_3 \rightarrow P_4$

b) Deadlock detection

The cycle:

$P_1 \rightarrow P_2 \rightarrow P_5 \rightarrow P_6 \rightarrow P_1$

Hence, deadlock exists. Processes in deadlock:

P_1, P_2, P_5, P_6

c) Distributed algorithm to detect it
A suitable algorithm:

Chandy - Misra - Haas Probe - Based Deadlock

Detection sends Algorithm

This algorithm sends PROBE messages across sites to detect cycles in distributed wait-for graph.

Q4

Distributed file system performance:

given:

Local access = 5 ms

Remote access = 25 ms

Probability = 0.3 \rightarrow local = 0.7

a) Expected access time:

$$E = (0.7 \times 5) + (0.3 \times 25)$$

$$= 3.5 + 7.5 = 11 \text{ ms}$$

b) Caching strategy:

Client-side caching with LRU (Least Recently Used)

JUSTIFYING:

Users usually access the same files repeatedly. Caching them locally reduces remote accesses \rightarrow significantly lowering avg. access time and network load.

Q5

Check Pointing Strategy:

given:

full checkpoint = 200 ms

incremental = 50 ms

RPO = 1 sec

period = 10 sec

a) Optimal mix:

 \rightarrow Take 1 full checkpoint at the start. \rightarrow then take incremental checkpoints every 1 second \rightarrow total 9 incremental checkpoints.

Total Overhead :

full : 200ms

incremental : $9 \times 50 = 450$ ms

total : 650 ms over 10 seconds

b) Reasoning :

→ $RPO = 1s$ → system must checkpoints at least once every 1 second

- full checkpoint is expensive; incremental is cheaper
- taking only incremental checkpoints avoids high overhead while staying within RPO.
- one full checkpoint ensures a clean baseline for recovery and consistency.

Q9

Case Study: Global E-Commerce Platform

a) Distributed scheduling flash sales:

- sudden traffic spikes across continents
- hotspot servers overloaded while others idle.
- maintaining fairness across regions.
- handling network latency & replication delays.
- Ensuring fast task migration b/w servers

Suitable load-balancing algorithm :

Consistent hashing with dynamic load redistribution.

BEST CHOICE:

Work stealing algorithm Idle server "steal" tasks from overloaded server → perfect for unpredictable flash sale workloads

- b) Fault tolerance strategy (RTO and RPO focused)
- use geo-replicated active-active clusters with:
- synchronous replication inside a region (low RPO)
 - asynchronous cross-region replication (low latency)
 - automatic failover using health checks + DNS-based global load balancing.

ENSURES:

- low RPO:
Data loss minimised since updates are replicated frequently.
- low ~~R~~ RTO:
Traffic can switch to a healthy region within seconds.

EXAMPLES:

- multi-region distributed file system.
- checkpoint + log-based recovery in each region.
- stateless microservices + replicated databases.

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