

Design and Analysis of Computer Algorithms

CS 6363.005: Homework #1

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Problem 1

Part (a)

$$\Theta(\log_2(n))$$

Part (b)

$$\Theta(n^2)$$

Part (c)

The *RecursiveAlg* was designed to find duplicate item in the array $A[1...n]$, by dividing the array in two part, comparing every item from each half, and then keeping doing the operation recursively.

Problem 2

Part (a)

Proof: Let $nodes(n)$ be the nodes number of a good binary tree *gbt*, n is the leaves number. We are required to prove $\forall n > 0, nodes(n) = 2n - 1$.

Base Case: $n = 1$: The only node in the tree is the leaf itself. The tree has $2n - 1 = 1$ node. Hence the claim holds true for $n = 1$.

Induction step: Let $k > 1$ be an arbitrary natural number.

Let us assume the induction hypothesis: for every *gbt* with i leaves, $0 < i \leq k$, assume $nodes(i) = 2i - 1$. We will prove $nodes(i + 1) = 2(i + 1) - 1$.

Since the tree is a *gbt*, every non-leaf node point two distinct *gbt*'s. If we want to add one leaves to the tree, we need to add two leaves to a leaf node, making the original leaf node a non-leaf node. Hence two nodes are added to the *gbt*, which mean

$$nodes(k + 1) = nodes(k) + 2 = 2k - 1 + 2 = 2(k + 1) - 1 \quad (1)$$

Thus establishes the claim for $k + 1$.

By the principle of mathematical induction, the claim holds for all $n \in \mathbb{N}$. \square

Part (b)

It does not hold that a *bbt* with $n > 0$ leaves, has $2n - 1$ nodes in total.

Consider removing one leaf from a *gbt* only. The leaf's parent node now has one child point to a *bbt*, which is Null (i.e. the empty tree). Now there are $n - 1$ leaves with $2n - 2$ nodes, which indicates $nodes(n - 1) = 2(n - 1)$ for a *bbt*.

Thus, the upper bound possible is $2n$.

Problem 3

The Asymptotic Bounds are as follow:

(a)

$$\sum_{i=1}^n \frac{n}{i} = \Theta(n \log(n))$$

(b)

$$\sum_{i=1}^n i^3 = \Theta\left(\frac{n^2(n+1)^2}{4}\right) = \Theta(n^4)$$

(c)

$$\begin{aligned} \sum_{i=1}^n \log\left(\frac{n}{i}\right) &= \sum_{i=1}^n (\log(n) - \log(i)) \\ &= \sum_{i=1}^n (\log(n)) - \sum_{i=1}^n (\log(i)) \\ &= n \log(n) - \log n! \\ &\approx n \log n - (n \log n - n) \\ &= \Theta(n) \end{aligned}$$

(d)

$$\frac{2}{n} + \sqrt{n} + \log^3 n = \Theta\left(n^{\frac{1}{2}}\right)$$

(e)

$$Bits(n) = \log_2(10^n) = n \log_2(10) = \Theta(n)$$

(f)

$$P(100) = \Theta(1)$$

Problem 4

- $n^{\frac{7}{(2n)}}$
- $\cos n + 2$
- $\lg^* \left(\frac{n}{8}\right), \lg^* \left(2^{2^{2^n}}\right)$
- $\lg^{(201)} n$
- $1 + \lg \lg \lg n$
- $12 + \lfloor \lg \lg(n) \rfloor$
- $\sqrt{\lg n}$
- $\lg n$
- $(\lg^* n)^{\lg n}$
- $n^{3/2 \lg n}$
- $n^{1/\lg \lg n}$
- $n^{\frac{1}{125}}$
- $n, 2^{\lg n}$
- $n(\lg n)^4$
- $\sqrt[n]{n^e}$
- $\sum_{i=1}^n i$
- $n^{2.1}$
- $\sum_{i=1}^n i^2$
- n^5
- $n^{\lg \lg n}$
- $(\lg(2 + n))^{\lg n}$
- $\left(1 + \frac{1}{154}\right)^{15n}$

Problem 5

The pseudo-code:

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1: procedure P(n)
2:   if  $n = 0$  then                                ▷ If there is no space left.
3:     return 1
4:   else if  $n = 1$  then                                ▷ If there is only 1 space left.
5:     return 1
6:   else                                ▷ If there are more than 1 spaces, then there are two ways to park.
7:     return  $P(n - 1) + P(n - 2)$ 
8:   end if
9: end procedure

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Algorithm 1: Count Number of Distinct Strings

Problem 6

Part (a) $T(n) = 2T(n/2) + n^4$

$T(n) = 2T(n/2) + n^4$ indicates that at the i recursive level, the total operations are:

$$f_i(n) = 2^i \times \left(\frac{n}{2^i}\right)^4 = \frac{n^4}{2^{3i}}$$

The depth of the recursion tree is $\lg n$, hence, the total running time is:

$$T(n) = \sum_{i=0}^{\lg n} f_i(n) = \sum_{i=0}^{\lg n} \frac{n^4}{2^{3i}}$$

Now consider the ratio of successive level sums:

$$r = \frac{f_{i+1}(n)}{f_i(n)} = \frac{\frac{n^4}{2^{3(i+1)}}}{\frac{n^4}{2^{3i}}} = \frac{1}{8} < 1$$

Which means the total running time $T(n)$ is mainly decided by the $T(\text{root})$, i.e.

$$T(n) = \Theta(n^4) \tag{2}$$

Part (b) $T(n) = 16T(n/4) + n^2$

$T(n) = 16T(n/4) + n^2$ indicates that at the i recursive level, the total operations are:

$$f_i(n) = 16^i * \left(\frac{n}{4^i}\right)^2 = n^2$$

which is not related to the level i . As the depth of the recursion tree is $\log_4 n$, hence, the total running time is:

$$T(n) = \sum_{i=0}^{\log_4 n} f_i(n) = \sum_{i=0}^{\log_4 n} n^2 = n^2 \log_4 n = \Theta(n^2 \log n) \quad (3)$$

Part (c) $T(n) = 2T(n/3) + T(n/4) + n$

Assume $T(n) = 2T(n/3) + T(n/4) + n$ indicates that at the i recursive level, the total operations are $f_i(n)$. We can write the $f_i(n)$ as:

$$\begin{aligned} f_0(n) &= n, \\ f_1(n) &= \left(\frac{2}{3} + \frac{1}{4}\right)n = \left(\frac{11}{12}\right)n, \\ f_2(n) &= \left(\frac{4}{3^2} + \frac{4}{3 \times 4} + \frac{1}{4^2}\right)n = \left(\frac{121}{144}\right)n = \left(\frac{11}{12}\right)^2 n, \\ f_3(n) &= \left(\frac{8}{3^3} + \frac{12}{3^2 \times 4} + \frac{6}{3 \times 4^2} + \frac{1}{4^3}\right)n = \left(\frac{1331}{1728}\right)n = \left(\frac{11}{12}\right)^3 n, \\ &\dots \end{aligned}$$

According the pattern of $f_i(n)$ above, it is easy to conclude and prove by induction that the general form of $f_i(n)$ is:

$$f_i(n) = \left(\frac{11}{12}\right)^i n$$

Now consider the ratio of successive level sums:

$$r = \frac{f_{i+1}(n)}{f_i(n)} = \frac{\left(\frac{11}{12}\right)^{i+1} n}{\left(\frac{11}{12}\right)^i n} = \frac{11}{12} < 1$$

Which means that the total running time $T(n)$ is mainly decided by the $T(\text{root})$, i.e.

$$T(n) = \Theta(n) \quad (4)$$

Part (d) $T(n) = T(n-1) + \sqrt{n}$

$T(n) = T(n-1) + \sqrt{n}$ indicates that at the i recursive level, the total operations are:

$$f_i(n) = \sqrt{n-i}$$

The depth of the recursion tree is n , hence, the total running time is:

$$T(n) = \sum_{i=0}^n f_i(n) = \sum_{i=0}^n \sqrt{n-i} = \sum_{i=1}^n \sqrt{n}$$

Now consider the ratio of successive level sums:

$$r = \frac{f_{i+1}(n)}{f_i(n)} = \frac{\sqrt{n-i+1}}{\sqrt{n-i}} < 1$$

Which means that the total running time $T(n)$ is mainly decided by the $T(\text{root})$, i.e.

$$T(n) = \Theta(\sqrt{n}) \tag{5}$$

Part (e) $T(n) = 3T(n/2) + 5n$

$T(n) = 3T(n/2) + 5n$ indicates that at the i recursive level, the total operations are:

$$f_i(n) = 3^i \times \left(\frac{5n}{2^i}\right) = \left(\frac{3}{2}\right)^i 5n$$

The depth of the recursion tree is $\lg n$, hence, the total running time is:

$$T(n) = \sum_{i=0}^{\lg n} \left(\frac{3}{2}\right)^i 5n = 5n \sum_{i=0}^{\lg n} \left(\frac{3}{2}\right)^i$$

Now consider the ratio of successive level sums:

$$r = \frac{f_{i+1}(n)}{f_i(n)} = \frac{\left(\frac{3}{2}\right)^{i+1} 5n}{\left(\frac{3}{2}\right)^i 5n} = \frac{3}{2} > 1$$

Which means the total running time $T(n)$ is mainly decided by the $T(\text{leaf})$, the last recursion's running time i.e.

$$T(n) = 5 \times 3^{\lg n} = \Theta(n^{\lg 3}) \tag{6}$$

Part (f) $T(n) = T(\sqrt{n}) + 7$

$T(n) = T(\sqrt{n}) + 7$ indicates that at the i recursive level, the total operations are:

$$f_i(n) = 7$$

The depth of the recursion tree is n , hence, the total running time is:

$$T(n) = \sum_{i=0}^{\sqrt{n}} f_i(n) = \sum_{i=0}^{\sqrt{n}} 7$$

In this case, the ratio of successive level sums is apparently $r = 1$, which means that the total running time $T(n)$ is mainly decided by the depth of the recursion tree.

Assume the base case is $n = 1$, $\sqrt[2^i]{n}$, in which i is a natural number representing the i th recursion, will not reach 1, though $\lim_{i=0}^{\infty} \sqrt[2^i]{n} = 1$. Thus, the recursion will not reach base case, i.e.

$$T(n) = \Theta(\infty) \tag{7}$$