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Design and Analysis of Computer Algorithms

CS 6363.005: Homework #2

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Problem 1 Inversions (25 points)

Algorithm 1 Count Inversion of two sorted arrays.

```

1: procedure COUNT_INVERSIONS( $A[1 \dots n + m]$ ,  $B[1 \dots n]$ ,  $C[1 \dots m]$ )
2:    $InversionsCount = 0$ 
3:    $Bindex = 1$ 
4:    $Cindex = 1$ 
5:   for  $Aindex = 1$  to  $n + m$  do
6:     if  $Cindex > m$  then ▷  $C$  is exhausted.
7:        $A[Aindex] = B[Bindex]$ 
8:        $Bindex ++$ 
9:     else if  $Bindex > n$  then ▷  $B$  is exhausted.
10:       $A[Aindex] = C[Cindex]$ 
11:       $Cindex ++$ 
12:     else if  $B[Aindex] < C[Cindex]$  then ▷  $B$  is smaller than  $C$ .
13:        $A[Aindex] = B[Bindex]$ 
14:        $Bindex ++$ 
15:     else ▷  $C$  is smaller than  $B$ , count inversions.
16:        $A[Aindex] = C[Cindex]$ 
17:        $Cindex ++$ 
18:        $InversionsCount = InversionsCount + (n - Aindex)$ 
19:     end if
20:   end for
21:   return  $InversionsCount$ 
22: end procedure

```

Algorithm 2 Sort and Count Inversions

```
1: procedure SORT_AND_COUNT( $A[1 \dots n]$ )
2:   if  $n > 1$  then
3:     Sort_and_Count( $A[1 \dots \lfloor n/2 \rfloor]$ )
4:     Sort_and_Count( $A[\lfloor n/2 \rfloor + 1 \dots n]$ )
5:     Count_Inversions( $\overline{A}[1 \dots n]$ ,  $A[1 \dots \lfloor n/2 \rfloor]$ ,  $A[\lfloor n/2 \rfloor + 1 \dots n]$ )
6:
```
