# Master Theorem.sage

February 22, 2022

#### 1 Master Theorem Recurrences

### 1.1 Maxwell Kapral

```
<sub>
     <i>powered by SageMath in Jupyter</i>
</sub>
```

#### 1.2 Of the Form:

$$T\left(n\right) = a T\left(\frac{n}{b}\right) + \mathcal{O}\left(n^{k} \log^{p}\left(n\right)\right)$$

Where n is the size of the problem, a is the number of subproblems in the recursion  $a \ge 1$ ,  $\frac{n}{b}$  is the size of each subproblem, b > 1,  $k \ge 0$ ,  $p \in \mathbb{R}$ 

 $\begin{array}{l} \text{expr is a lazy, magic variable. If expr is some } f\left(n\right), \text{ then } T\left(n\right) \in \mathcal{O}\left(f\left(n\right)\right) \text{ if } f\left(n\right) \in \Omega\left(n^{\log_b(a) + c_1}\right) \\ \text{and } a \cdot f\left(\frac{n}{h}\right) \leq c_2 \cdot f\left(n\right) \text{ where } c_1 > 0 \text{ and } c_2 < 1, \text{ for all } n \text{ large enough } (+\infty). \end{array}$ 

```
[1]: from sage.rings.asymptotic.term_monoid import OTermMonoid, ExactTermMonoid
     from sage.rings.asymptotic.term_monoid import DefaultTermMonoidFactory as T_M
     from sage.rings.asymptotic.growth_group import GrowthGroup
     n = SR.var('n')
     def mastertheorem(a, b, k, p, expr=n):
         assert(a >= 1), "a not greater than or equal to 1"
         assert(b > 1), "b not greater than 1"
         assert(k >= 0), "k not greater than or equal to 0"
         assert(p in RR), "p must be a real number"
         if expr != n.factorial():
             G_G = GrowthGroup('(RR_+)^n * n^RR * log(n)^RR')
             bigO = OTermMonoid(T_M, growth_group=G_G, coefficient_ring=RR)
             bigTheta = ExactTermMonoid(T_M, growth_group=G_G, coefficient_ring=RR)
             if expr == n:
                 if p != 0:
                     expr = (n**k)*(log(n)**p)
                 else:
                     expr = n**k
```

```
ans = None
      eps = 0.01
       if bigO(n**(log(a, b) - eps)).can_absorb(bigO(expr)):
           ans = n**log(a, b)
       if ans is None:
           if a >= b**k:
               if p > -1:
                   for i in range(0, p+1, 1):
                       if bigTheta((n**log(a, b))*(log(n))**i).
→can_absorb(bigTheta(expr)):
                           ans = (n**log(a, b))*(log(n)**(i+1))
                           break
               elif p == -1:
                   ans = (n**log(a, b))*log(log(n), b)
               elif p < -1:
                   ans = n**log(a, b)
      if ans is None:
           c1 = 0.01
           if bigO(expr).can_absorb(bigO(n**(log(a, b) + c1))):
               c2 = 0.99
               if bigO(c2*expr).can_absorb(bigO(a*(expr/b))):
                   ans = expr
  else:
      ans = factorial(n)
   #########
   # OUTPUT #
   #########
  if a == 1:
      if p == 0:
           show(LatexExpr(r"T\left(n\right)=T\left(\dfrac{n}{"+latex(b)+
                          r"}\right)+"+latex(expr.simplify())))
      elif p == 1:
           if k > 0:
               show(LatexExpr(r"T\left(n\right)=T\left(\dfrac{n}{"+latex(b)+
                              r"}\right)+"+latex(expr.simplify())))
           else:
               show(LatexExpr(r"T\left(n\right)=T\left(\dfrac{n}{"+latex(b)+
                              r"}\right)+\log{n}"))
      elif p == (-1):
           show(LatexExpr(r"T\eft(n\right)=T\eft(\dfrac{n}{"+latex(b)+}
                          r"}\right)+"+latex(expr.simplify())))
      elif p > 1:
           if k > 0:
               show(LatexExpr(r"T\left(n\right)=T\left(\dfrac{n}{"+latex(b)+
                              r"}\right)+"+latex(expr.simplify())))
           else:
```

```
show(LatexExpr(r"T\left(n\right)=T\left(\dfrac{n}{"+latex(b)+
                           r"}\right)+\log^{"+latex(p)+r"}{n}"))
    elif p < (-1):
        show(LatexExpr(r"T\left(n\right)=T\left(\dfrac{n}{"+latex(b)+
                       r"}\right)+"+latex(expr.simplify())))
else:
   if p == 0:
        show(LatexExpr(r"T\left(n\right)="+latex(a)+
                       r"\thinspace T\left(\dfrac{n}{"+latex(b)+
                       r"}\right)+"+latex(expr.simplify())))
    elif p == 1:
        if k > 0:
            show(LatexExpr(r"T\left(n\right)="+latex(a)+
                           r"\thinspace T\left(\dfrac{n}{"+latex(b)+
                           r"}\right)+"+latex(expr.simplify())))
        else:
            show(LatexExpr(r"T\left(n\right)="+latex(a)+
                           r"\thinspace T\left(\dfrac{n}{"+latex(b)+
                           r"}\right)+\log{n}"))
    elif p == (-1):
        show(LatexExpr(r"T\left(n\right)="+latex(a)+
                       r"\thinspace T\left(\dfrac{n}{"+latex(b)+
                       r"}\right)+"+latex(expr.simplify())))
    elif p > 1:
        if k > 0:
            show(LatexExpr(r"T\left(n\right)="+latex(a)+
                           r"\thinspace T\left(\dfrac{n}{"+latex(b)+
                           r"}\right)+"+latex(expr.simplify())))
        else:
            show(LatexExpr(r"T\left(n\right)="+latex(a)+
                           r"\thinspace T\left(\dfrac{n}{"+latex(b)+
                           r"}\right)+\log^{"+latex(p)+r"}{n}"))
    elif p < (-1):
        show(LatexExpr(r"T\left(n\right)="+latex(a)+
                       r"\thinspace T\left(\dfrac{n}{"+latex(b)+
                       r"}\right)+"+latex(expr.simplify())))
print('') # Newline
show(LatexExpr(r"T\left(n\right)\in\Theta\left("+latex(ans.simplify())+
               r"\right)"))
```

## 2 Examples

2.0.1 The cost  $T\left(n\right)$  of a merge sort on a list of n numbers is governed by the following recurrence relation:

$$T\left( n\right) =2\,T\left( \frac{n}{2}\right) +n-1,\quad T\left( 1\right) =0$$

[2]: # a = 2, b = 2, k = 1, p = 0 mastertheorem(2, 2, 1, 0)

$$T\left(n\right) = 2T\left(\frac{n}{2}\right) + n$$

 $T\left(n\right)\in\Theta\left(n\log\left(n\right)\right)$ 

2.0.2 Prove that

$$T\left(n\right)\in\mathcal{O}\left(n^{2}\right)$$

2.0.3 For

$$T\left(n\right)=3\,T\left(\frac{n}{5}\right)+n^{2}\quad n>5$$

$$T(5) = 51$$
  $n = 5$ 

[3]: # a = 3, b = 5, k = 2, p = 0mastertheorem(3, 5, 2, 0)

$$T\left(n\right)=3\,T\left(\frac{n}{5}\right)+n^2$$

 $T\left( n\right) \in\Theta\left( n^{2}\right)$ 

2.0.4 Express the Running Time of the Following Recurrence Relation

$$T\left(n\right)=5\,T\left(\frac{n}{3}\right)+n^{4}\quad n>1$$

$$T(1) = 1$$
  $n = 1$ 

[4]: # a = 5, b = 3, k = 4, p = 0mastertheorem(5, 3, 4, 0)

$$T\left(n\right) = 5T\left(\frac{n}{3}\right) + n^4$$

$$T\left(n\right)\in\Theta\left(n^{4}\right)$$

[5]: # a = 3, b = 2, k = 2, p = 0mastertheorem(3, 2, 2, 0)

$$T\left( n\right) =3\,T\left( \frac{n}{2}\right) +n^{2}$$

 $T\left(n\right)\in\Theta\left(n^{2}\right)$ 

[6]: # a = 4, b = 2, k = 2, p = 0mastertheorem(4, 2, 2, 0)

$$T\left(n\right)=4\,T\left(\frac{n}{2}\right)+n^{2}$$

 $T\left(n\right)\in\Theta\left(n^{2}\log\left(n\right)\right)$ 

[7]:  $\# a = 1, b = 2, k = 0, p = 0, expr = 2^n$ mastertheorem(1, 2, 1, 0, (2\*\*n))

$$T\left( n\right) =T\left( \frac{n}{2}\right) +2^{n}$$

 $T(n) \in \Theta(2^n)$ 

[8]: # a = 2, b = 2, k = 1, p = -1mastertheorem(2, 2, 1, -1)

$$T\left(n\right) = 2\,T\left(\frac{n}{2}\right) + \frac{n}{\log(n)}$$

 $T\left(n\right) \in \Theta\left(\frac{n \log(\log(n))}{\log(2)}\right)$ 

[9]: # a = 16, b = 4, k = 1, p = 0, expr = n!mastertheorem(16, 4, 1, 0, factorial(n))

$$T\left(n\right) = 16T\left(\frac{n}{4}\right) + n!$$

$$T(n) \in \Theta(n!)$$

[10]: # a = 6, b = 3, k = 2, p = 1 mastertheorem(6, 3, 2, 1)

$$T(n) = 6T\left(\frac{n}{3}\right) + n^2\log(n)$$

 $T(n) \in \Theta(n^2 \log(n))$ 

[11]: # a = 4, b = 2, k = 0, p = 1 mastertheorem(4, 2, 0, 1)

$$T\left(n\right) = 4T\left(\frac{n}{2}\right) + \log n$$

 $T\left(n\right) \in \Theta\left(n^2\right)$ 

[12]: # a = 2, b = 2, k = 1, p = 1 mastertheorem(2, 2, 1, 1)

$$T\left(n\right) = 2T\left(\frac{n}{2}\right) + n\log\left(n\right)$$

$$T(n) \in \Theta\left(n\log\left(n\right)^2\right)$$

[13]: # a = 2, b = 4, k = 0.51, p = 0 mastertheorem(2, 4, 0.51, 0)

$$T\left(n\right) = 2T\left(\frac{n}{4}\right) + n^{0.51}$$

$$T\left(n\right)\in\Theta\left(n^{0.51}\right)$$

[14]: # a = sqrt(2), b = 2, k = 0, p = 1 mastertheorem(sqrt(2), 2, 0, 1)

$$T\left(n\right) = \sqrt{2}\,T\left(\frac{n}{2}\right) + \log n$$

$$T\left(n\right)\in\Theta\left(\sqrt{n}\right)$$

[15]: # a = 3, b = 3, k = 1/2, p = 0mastertheorem(3, 3, 1/2, 0)

$$T\left(n\right) = 3T\left(\frac{n}{3}\right) + \sqrt{n}$$

 $T(n) \in \Theta(n)$ 

[16]: # a = 3, b = 2, k = 1, p = 0mastertheorem(3, 2, 1, 0)

$$T\left( n\right) =3\,T\left( \frac{n}{2}\right) +n$$

 $T\left(n\right) \in \Theta\left(n^{\frac{\log(3)}{\log(2)}}\right)$ 

[17]: # a = 2, b = 7, k = 6, p = 0 mastertheorem(2, 7, 6, 0)

$$T\left(n\right) = 2T\left(\frac{n}{7}\right) + n^6$$

 $T\left(n\right) \in \Theta\left(n^6\right)$ 

[18]: # a = 16, b = 4, k = 2, p = 3 mastertheorem(16, 4, 2, 3)

$$T\left(n\right)=16\,T\left(\frac{n}{4}\right)+n^{2}\log\left(n\right)^{3}$$

 $T\left(n\right)\in\Theta\left(n^{2}\log\left(n\right)^{4}\right)$ 

[19]: # a = 9, b = 2, k = 3, p = 0 mastertheorem(9, 2, 3, 0)

$$T\left(n\right)=9\,T\left(\frac{n}{2}\right)+n^{3}$$

 $T\left(n\right) \in \Theta\left(n^{\frac{2\log(3)}{\log(2)}}\right)$ 

[20]: # a = 27, b = 3, k = 4, p = 0mastertheorem(27, 3, 4, 0)

$$T\left(n\right) = 27T\left(\frac{n}{3}\right) + n^4$$

$$T\left(n\right)\in\Theta\left(n^{4}\right)$$

[21]: # a = 16, b = 2, k = 4, p = 5mastertheorem(16, 2, 4, 5)

$$T\left(n\right) = 16 \, T\left(\frac{n}{2}\right) + n^4 \log\left(n\right)^5$$

$$T\left(n\right)\in\Theta\left(n^{4}\log\left(n\right)^{6}\right)$$

[22]: # a = 1, b = 2, k = 1, p = -1mastertheorem(1, 2, 1, -1)

$$T\left(n
ight) = T\left(rac{n}{2}
ight) + rac{n}{\log(n)}$$

$$T\left(n
ight) \in \Theta\left(rac{n}{\log(n)}
ight)$$

[23]: # a = 16, b = 2, k = 4, p = -1 mastertheorem(16, 2, 4, -1)

$$T\left(n\right) = 16\,T\left(\frac{n}{2}\right) + \frac{n^4}{\log(n)}$$

$$T\left(n
ight) \in \Theta\left(rac{n^4\log(\log(n))}{\log(2)}
ight)$$

[24]: # a = 16, b = 2, k = 4, p = -4 mastertheorem(16, 2, 4, -4)

$$T\left(n\right) = 16T\left(\frac{n}{2}\right) + \frac{n^4}{\log(n)^4}$$

$$T(n) \in \Theta(n^4)$$

 ${\bf 2.0.5} \quad {\bf The \ Camerini \ Minimum \ Bottleneck \ Spanning \ Tree \ Algorithm}$ 

$$T\left(n\right) = \, T\left(\frac{n}{2}\right) + \mathcal{O}\left(n\right)$$

[25]: # a = 1, b = 2, k = 1, p = 0 mastertheorem(1, 2, 1, 0)

$$T\left(n\right) = T\left(\frac{n}{2}\right) + n$$

 $T\left( n\right) \in\Theta\left( n\right)$