

**University of Applied Sciences** 

## Automatic Interaction Diagram Generation of Vue.js-based Web Applications

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at

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1. Supervisor: Prof. Dr. Gefei Zhang

2. Supervisor: Prof. Dr.-Ing. Hendrik Gärtner

submitted by: B.Sc. Anton Karakochev

matriculation number: 553324

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#### Abstract

This thesis aims to further explore the concept of interaction diagrams for scenario testing introduced by [ZZ19]. It is applied to a different framework (Vue.js) ... - automatically generate based on vue js code - lists - objects - computed property - scenarios in Gherkin generated

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## Chapter 1

## Introduction

1 p motivation, etc

## Chapter 2

## Fundamentals and State of The Art

#### 2.1 State of the Art

TODO

# 2.2 Scenario Testing of AngularJS-based Single Page Web Applications

Zhang and Zhao [ZZ19] present a method with the goal of achieving better undestanding of AngularJS-based single page applications (SPAs) and also devised a way to specify test coverage criteria based on it. At the center of the proposed method are interaction diagrams, which are used to model the overall data and control flow of an application. [ZZ19]

#### 2.2.1 Abstract Syntax

Zhang and Zhao [ZZ19] model a Angular JS-based SPA as a tuple (T, C, D, E), where

- T is a HTML template, consisting of a set of HTML tags (widgets)  $(T = \{h\})$
- C is a controller (view-model), written in JavaScript. It is modeled as a tuple (V, F, \$scope), where F and V are top level variables and functions respectively and  $\$scope \in V$  is a distinguished element of V. Further V(\$scope) and F(\$scope) denote all variables and functions of \$scope respectively.  $W = V \setminus \{\$scope\}$  denotes top level variables not in scope. Additionally  $init \in F$  is defined as an initialization function
- D is a set of data bindings between HTML tags and variable properties of  $\$scope\ D \subseteq \{(h, V(\$scope) \cup F(\$scope))\}$ . Given  $d = (n, o)\ source(d) = n$  and target(d) = o. For Two-way bindings  $D' \subseteq D$  and  $\forall d \in D'\ target(d) \in V(\$scope)$ .
- E a set of event handler bindings between HTML tags and function properties of  $scope: E \subset \{(h, F(scope))\}$ . In addition, for each function  $f \in F \cup F(scope)$ ,  $R(f) \subseteq V \cup V(scope)$

and  $W(f) \subseteq V \cup V(\$scope)$  are defined as the values that the given function reads from and writes to.  $Inv(f) \in F$  are defined as the functions invoked by f. [ZZ19]

#### 2.2.2**Interaction Diagrams**

Zhang and Zhao [ZZ19] define interaction diagrams as a directed graph (N, E) where the set of nodes N is defined as the union of  $N_H$  (HTML tag nodes),  $N_{\$scope}$  (TODO name),  $N_js$  (TODO name).

```
TODO double check, write as text N_H = \{n_h | (h, v) \in D\}
N_{\$scope} = \{n_v | (h, v) \in D\} \cup \{n_e | (h, e) \in E\}
N_j s = \{n_v | v \in W\} \cup \{n_f | f \in F\}
n_{init} is distinguished by an incoming arrow without a starting vertex
Edges:
bindings
e_d = (target(d), source(d)) E_{data} = \{e_d | d \in D\}
additionally if d \in D' also create e'_t = (source(t), target(t)) and E'_d ata = \{e_d | d \in D'\}
for events (E_{event}) (h, f) \in E,
write E_W(f, v) where f \in F \cup F(\$scope), v \in W(f)
read E_R(v, f) where f \in F \cup F(\$scope), v \in R(f)
invoked E_{Inv} (f, v) where f \in F \cup F(\$scope), v \in I(f)
E_{init} default values of widgets for each h \in T where \nexists v | (h, v) \in D create an edge (h, v)
```

explained in a lot of detail in [ZZ19, p. 9]

#### 2.2.3**Testing and Interactions**

Zhang and Zhao [ZZ19] define an interaction as a round of user input including updates to the widgets by the application. Interaction can be triggered explicitly by the user (by invoking an event handler) or implicitly while the user is updating data. [ZZ19]

Given the interaction diagrams as described in 2.2.2 it is possible to derive which widgets get updated by a user input action or set up by the initial function. Zhang and Zhao [ZZ19] define it formally as follows:

Given a node  $n \in N_H \cup \{init\}$ , we say a node  $m \in N_H$  reacts to n iff

- 1.  $\exists n_0, n_1, n_2, \dots, n_k \in \mathbb{N}, n_0 = n, n_k = m \text{ such that for each } 0 \le i < k \ (n_i, n_{i+1}) \in \mathbb{N}$ E, and
- 2.  $\forall n_p, 1$

We write l(n) for the set of all nodes representing the widgets that react to n. This set contains the widgets that are automatically updated upon user input, and thus constitute an interaction.

For example, in order for the widget n, which was clicked by the user, to update the widget m, m must be reachable from n by following the directed edges of the interaction diagram and only the first edge can be an event-handling edge.

What is crucial is that the interactions l(n) define an upper bound of what can be updated, i.e. what might get updated. Nevertheless, this information is sufficient in order to be able to define coverage criteria [ZZ19].

#### 2.2.4 Coverage Criteria

Interactions should not be tested in isolation and in order for tests to make sence, interactions as preconditions are required [ZZ19]. In order to define coverage criteria, Zhang and Zhao [ZZ19] extend their notation, as described in 2.2.2, by defining -  $\mathcal{I} = \{w \in T | l(w) \neq \emptyset\}$  all widgets, that result updates.

A sequence of user interactions, including the initial function is referred to as a scenario  $A = (a_0, a_1, \ldots, a_n)$  where  $a_0 = init$  and  $\forall 0 < k \le n, a_k \in \mathcal{I}$ . The widgets, to which a scenario reacts, are equal to the widgets to which the last widget in the scenario reacts -  $l(A) = l(a_n)$ .

The set of scenarios is generated by starting with the initial scenario, containing only the initial function  $S_0 = \{(init)\}$  and prolonging it iteratively by each widget, where the user can take an action. This is terminated once all  $i \in \mathcal{I}$  are included in at least one scenario. Formally: Define  $A \oplus x = a_0, a_1, \ldots, a_n, x$  For  $n > 0, S_{n+1} = \{p \oplus x | p \in S_n, x \in l(p) \cap \mathcal{I}\}$ 

Based on the scenario sets Zhang and Zhao [ZZ19] define the following coverage criteria:

- Each set  $S_n$  of test scenarios should be tested.
- For each given  $S_n$ , each  $p \in S_n$  should be tested.
- For each given p, each  $w \in l(p)$  should be tested. That is, there should be a test case for each widget that may be modified after the scenario p.

### 2.3 Scenario Testing

Scenario testing, was originally introduced in Kaner [Kan03] and later as Kaner [Kan13]. The author defines scenarios as hypothetical stories, which aid a person in understanding a complex system or problem. Scenario tests are tests, which are based on such scenarios. [Kan13, p. 1] Further, [Kan03, pp. 2–5] defines five characteristics, which make up a good scenario test as follows: A Scenario test must be

- based on a story based on a description of how the program is being used
- motivating stakeholders have interest in this test succeeding and would see to it's resolution
- credible probable to happen in the real world
- complex complex use, data or environment

• easy to evaluate - it should be easy to tell if the test succeded or failed based on the results

Kaner [Kan13] describes the biggest advantages of scenario testing to be - understanding and learning the product in early stages of development(1), connecting of testing and requirement documentations(2), exposing shortcomings in delivering of desired benefits(3), exploration of expert use of the program(4), expose requirement related issues(5).

### 2.4 Behavior-Driven Development

Behavior-Driven Development (BDD), pioneered by North [Nor06] is a software development process, that combines principles from Test-Driven Development and Domain-Driven design [EE04].

Its main goal is to specify a system in terms of its functionality (i.e. it's behaviors) with a simple domain-specific language (DSL) making use of English-like sentences. This stimulates collaboration between developers and non-technical stakeholders and further results in a closer connection between acceptance criteria for a given function and matching tests used for its validation.

BDD splits a user story into multiple scenarios, each formulated in the form of *Given*, *When*, *Then* statements, respectively specifying the prerequisite/context, event and outcomes of a scenario.

[TODO] example here? cut shorter

At present ... there based on the division of behavior descriptions and behaviors. Such as Jest/Jasmine combine behavior descriptions and behaviors into one, whereas as Cucumber uses a DSL named Gherkin to specify the behavior descriptions and provides a set of tools to generate behaviors.

#### 2.4.1 Gherkin Language

#### 2.5 Model-View-ViewModel

Model-View-ViewModel (MVVM) is a design pattern, which helps in creating a clear separation between business and presentation logic and User interface (UI) of an application. [Bri17, pp. 7–9]

In MVVM there are three core components - the view, model and view model. Those components are clearly separated from each other - the view is aware of the view model and the view model is aware of the model. However, this does not hold in reverse - the model is unaware of the view model and the view model is unaware of the view.



Figure 2.1: MVVM design pattern overview, adapted from [Bri17, p. 7]

#### 2.5.1 View

The view is what the user sees. It is responsible for the structure, layout and appearance of the application.

#### 2.5.2 View-Model

The view model implements event handlers and properties, to which the view can bind to. It also notifies the view of any changes to the underlying data. It defines the functionality, offered by the UI, but the view determines how it is presented.

#### 2.5.3 Model

The model encapsulates the data of the application and validation its logic.

### 2.6 Vue.js

Vue.js [vue21a] is a progressive front end framework for building user interfaces and single-page applications based on the MVVM design pattern described in 2.5 [Mac18] [21a].

#### 2.6.1 Components

At the core of Vue.js are components, which are small, self-contained, composable and often reusable custom elements. Almost any type of application can be represented as a tree of components [21a].

In more concrete terms, a Vue.js component is a single file with the extension of .vue, which consists of a template, script and optional style part. The template is a HTML-based template, which can be parsed specification compliant browsers and HTML parsers. It can contain other components or html elements and is equivalent to the view in MVVM.

The *code* section of a Vue.js component includes the view-model of the component. It has a special json object *data*, which is equivalent to the MVVM model. The *script* part of a computed includes css-like styles.

data binding is a general technique that binds data sources from the provider and consumer together and synchronizes them.

#### 2.6.2 Reactivity

... enables data binding

#### 2.6.3 Directives

Vue.js enables one way bindings(from source - data to target - component or html tag) via the v-bind (line X,Y) or moustache syntax (line Z). Bindings can contain expressions (line X,Y).

Two way binding can be achieved using the v-model directive (line,X,Y,Z). Event handlers can be bound by the method name or also expressions.

#### 2.6.4 Data Binding

Vue.js provides support for various forms of Data Binding via a special syntax. Both the data and computed objects of a *Vue.js* component are reactive

Via a special syntax Vue.js - one way - two way - event bindings - inline expressions - computed properties

#### 2.6.5 Vue.js directives

#### 2.6.6 Structure of a Component

```
(template, code etc.)
bindings two way, one way
g - data - computed properties
template part code part bindings
```

#### 2.7 ESLint

ESLint [21b] is a linting tool (linter) for ECMAScript/JavaScript. Linters are static code analysis tool, which can be used to flag and potentially automatically fix common code issues and enforce consistent code styling.

#### 2.7.1 Architecture

At a very high level, ESLint consists of



Figure 2.2: ESLint Architecture taken from [21c]

#### 2.8 Rules

At the core of ESLint are rules. Rules are extensible pieces of code, bundled as plugins, which can be used to verify various aspects of code. An example would be a rule, which checks for matching closing paranthesis. Each rule consists of a *metadata* object and a *create* function. The metadata object includes metadata such as documentation strings, the type of the rule and whehter it is fixable or not.

Based on type, rules can be either suggestions, problems or layout. Suggestions indicate some

improvement, but are not required and would not cause the linting to fail. Problems on the other hand would result in a linting failure. Layouts are rules that care mainly about the formatting of code, such as whitespaces, semicolons, etc.

If the fixable property is specified, it indicates that the errors reported by this rule can be automatically fixed. This can be applied via the --fix command line option. It has two possible values - code or whitespace indicating the type of fixes, that this rule would apply. For example in Integrated development environments (IDEs) fixable code errors would show a fix shortcut displayed next to them and whitespace rules could be applied when saving the file.

The *create* function of rules takes as arguments a *context* and returns an object of methods which are called by ESLint for each node based on the Visitor pattern while traversing the Abstract syntax tree (AST). ESLint provides a very powerful matching mechanism for specifying what nodes to match called selectors [21d] inspired by estools [est21a].

TODO cut shorter

TODO custom architecture image

TODO what are selectors

#### 2.8.1 AST Explorer

An incredibly useful tool when working with ASTs is AST Explorer, developed by Kling [Kli21]. It enables the exploration of syntax tree generated by various parsers and also includes the vue-eslint-parser [vue21b]

#### 2.8.2 ESTree AST

By default ESLint uses the [esl21] parser to parse JavaScript source code into an AST as defined by ESTree specification [est21b]. When Parsing .vue files ESLint uses this parser for the code inside the <script> tag.

#### 2.8.3 ESLint Parser Vue AST

In order to parse the *<template>* section of .vue files, ESLint uses the vue-eslint-parser [vue21b]. This parser outputs an AST compliant with their own AST specification, defined in [vue21c].

TODO add here from ast def

#### 2.8.4 Selectors

Chapter 3

Concept

### 3.1 Parsing Vue.js

#### 3.1.1 Assumptions

It is assumed that the Vue.js code, for which interaction diagrams are going to be generated, compiles and does not contain syntactical errors. No checks are performed in order to verify that. Naturally, logical errors are not an issue.

#### 3.1.2 Limitations

In order to be able to generate interaction diagrams, which capture every aspect of Vue.js, the generation must be directly based on an AST, which covers every possible syntax, such as [vue21b].

The approach proposed here only includes the following features of Vue.js:

- Event handlers (including anonymous method syntax and method reference syntax)
- Any one or two-way binding expressions  $(v-model,\ v-bind,\ "moustache",\ v-if)$  excluding v-else
- *v-for* statements for lists, excluding iterating through properties of an object or iteration with index (property zipped with index)
- distinguishing between properties and computed properties
- complex object and lists (non-nested) models
- methods, including the resolution of arguments, they have been called with (excluding methods called with other methods as arguments)

#### 3.1.3 AST

```
calls: calledMethod*;
16
      calledMethod: calledMethodIdentifier '(' calledArgs ')';
      accesedVariable: identifier;
      calledArgs: (calledMethod | accesedVariable)*;
      bindings: binding*;
      binding: tag bindingSource+;
      bindingSource: (accesedVariable | calledMethod) (EVENT_BINDING |
          ONE_WAY_BINDING)
                    accesedVariable TWO_WAY_BINDING;
      tag: name tagId loc;
      tagId: LINE '_' COLUMN '_' LINE '_' COLUMN;
      name: UNICODE | identifier;
      loc: start end;
      start: LINE COLUMN;
      end: LINE COLUMN;
      calledMethodIdentifier: methodDefinitionIdentifier | id* NAME_IDENTIFIER;
      methodDefinitionIdentifier: THIS NAME_IDENTIFIER;
      thisIdentifier: THIS identifier;
      identifier: NAME_IDENTIFIER id*;
      id: NUMERIC_INDEX | GENERIC_INDEX | NAME_IDENTIFIER;
      //terminals, tokens
42
      LINE: [0-9]+;
      COLUMN: [0-9]+;
44
      EVENT_BINDING: 'event';
      TWO_WAY_BINDING: 'two-way';
      ONE_WAY_BINDING: 'one-way';
      GENERIC_INDEX: 'i' | 'j' | 'k' | 'l' | 'm' | 'n';
      THIS: 'this';
      NUMERIC_INDEX: [0-9]+;
      NAME_IDENTIFIER: JS_IDENTIFIER;
      JS_IDENTIFIER: (UNICODE | '$' | '_') (UNICODE | '$' | '_' | [0-9])*;
```

#### UNICODE: [\u0000-\uFFFF];

A Vue.js SPA, including all the necessary information for 3.1.2, can be defined using the above grammar.

The application consists of bindings methodDefinitions a createdMethod, topLevelProperties and computedProperties.

The topLevelProperties represent the data object of the Vue.js script tag. Each property will be represented flattened, as a list of identifiers and prefixed with this, in order to indicate it belongs to the top level data object. For example  $problem:\{a:0, b:0\}$  will be represented as follows:

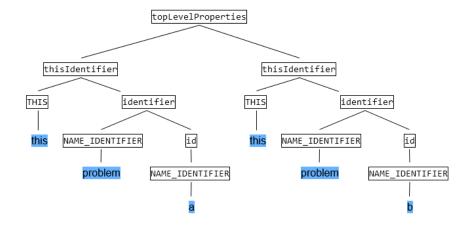


Figure 3.1: AST for top-level example

Bindings can be obtained from the Vue.js template.

Each binding consists of an HTML tag, an a list of binding sources for that tag - pairs of variable or method call and a binding type. The binding type represents the type of the binding - either event, one-way or two-way. Two-way bindings are only valid with properties, whereas for events and one-way bindings, both method calls and properties are possible, since in Vue.js a binding source could be an expressions defined as an inline anonymous functions (  $< div \ v-if="value" == true"/>$ ). The binding sources are a list, since a tag could have multiple different bound properties, or a bound expression. The information about how exactly the properties are bound, if it is the same type of binding, is discarded.

Method calls include the parameters they have been called with - other methods or just variables. It is also possible to call methods with binary expressions - those are represented as a special method, which takes 2 parameters - the left and right side operators of the binary expression. Expressions with multiple terms can be represented as multiple binary expressions. This representation loses information such as the order of operations, but since we are only interested in which properties are being accessed, this loss does not pose an issue.

A special case is accessing lists. For example  $< div \ v - bind = "problems[0].a"/>$  would result in the following:



Figure 3.2: AST for example

v-for statements, are substituted:

results in

```
subjects[i].problems[0]
```

which in term produces the following AST:

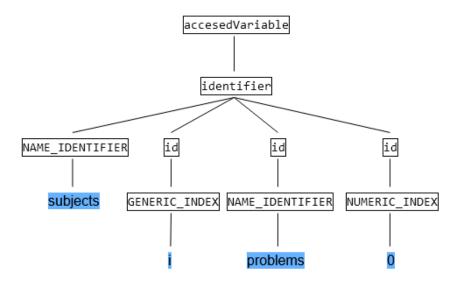


Figure 3.3: AST for example

Nested lists are also possible and would result in multiple generic indices being added.

Each tag includes its location in the source code (starting and ending line and column), which can also be used as an identifier. Tags also have humanly readable names, which are either equal to the text of the tag, if it exists, or to the identifier of the first binding.

methodDefinitions include all methods definitions from the method object of the view-model. Each method definition consists of the following

- identifier an identifier, equal to this followed by the name of the method
- arguments names of arguments, each of which is a simple name identifiers
- reads variable it reads from
- writes varaibles it writes to
- calls method calls, including arguments, same as for bindings

computedProperties are similar to methodDefinitions with the exception that they do not have arguments. Albeit bad practise, it is still possible for computed properties to have side effects and therefore they were modelled as methods.

### 3.2 Interaction Diagram Generation

The simplified Vue.js AST can be used to create a directed graph, which will represent the interaction diagrams. It is hard to directly generate this graph, therefore the capabilities of a directed, compounded graph will be leveraged and later on converted to a directed graph.

Vertices in this graph have the following properties

- 1. Globally Unique Identifier (GUID) used to reference and globally identify the vertex
- 2. label the name of the vertex, which is going to be displayed
- 3. type the type of the vertex (data, tag or method). Additionally for data vertices: numeric, generic or undefined (representing simple data vertices)
- 4. loc defined only on tag verices. Their location in the source code
- 5. parent defined only on vertices of type 'data'. A GUID of another vertex, used for a child/parent relationship (compound graph).

Edges in the graph are directed and each have a label property, which is one of 'event', 'calls' or 'simple'.

The core idea of the algorithm is to generate vertices only for nodes which are being accessed instead of the whole application. A second pass of the data is also needed to add additional edges for lists.

#### 3.2.1 Variable Identifiers

Variable identifiers are represented by *identifier* and *thisIdentifier* in the AST 3.1.3. For the *this* and for each *id* node in the *identifier* or *thisIdentifier* a vertex is created in order. Those vertices are connected using unidrection edges, labeled with 'data' and also each vertex (excluding the first one) has its parent set to the previous. There is one exception to this process - When accessed from *write* of a *method*, nodes of type *GENERIC\_INDEX* are omitted. The reason behind this will be explained in this section 3.2.3;

Each vertex has a GUID equal to the value of its terminal symbol (NUMERIC\_INDEX, GENERIC\_INDEX or NAME\_IDENTIFIER), concatenated with the value of the previous vertix. The label of those vertices are equal to the terminal symbol in case of NAME\_IDENTIFIER and in case of GENERIC\_INDEX and NUMERIC\_INDEX, combined with the label of the previous vertex using square braces. Set the type of each vertex to 'data'. Add the type 'numeric' to vertices created from NUMERIC\_INDEX nodes and 'generic' to vertices created from GENERIC\_INDEX nodes.

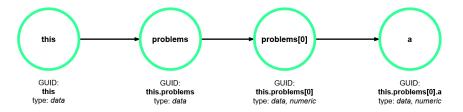


Figure 3.4: Example Graph obtained for the identifier this . problems [0]. a

#### 3.2.2 Object representation

Using the representation for identifiers in the previous section, objects will result in being displayed dynamically, based on which properties are accessed. Nodes and edges are created on a 'create if non-existent' basis. In the example below, if this .problem.b is accessed after this .problem.a it will only result in the creation of the node a and edge this .problem -> this .problem.b.

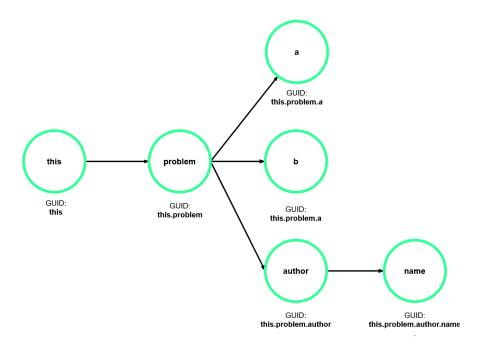


Figure 3.5: Example Graph for the following object accesses - this.problem.a, this.problem.b, this.problem.author.name

Furthermore, updates can be formulated nicely with the above representation. If *problem* were to be changed, it would result in a cascade update of all properties. If *author* would be change, it would only result in a cascading change in *name*.

#### 3.2.3 List representation

Lists will be represented based on the template in 3.6 for a list named P.

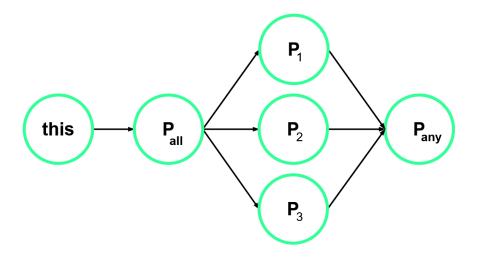


Figure 3.6: Generic List representation

Concrete elements, which are accessed, are detoned as  $P_{\langle index \rangle}$  and additionally a vertex  $P_{all}$ , which can be used to update all elements of a list and their properties, is created. Another vertex  $P_{any}$  is also created, which can be used to observe once any vertex of  $P_1, P_2, P_3, \ldots P_n$  changes. If  $P_1$  were to be updated by any method, it would not result in updates to any of  $P_2, P_3, \ldots P_n$ 

The same construct can also be leveraged when it comes to properties of list elements. Each top level property of that element will have an all vertex, connected to the  $P_{all}$  node of the list.

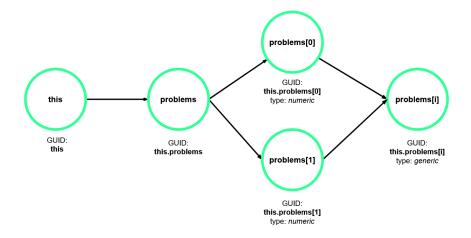


Figure 3.7: Concrete example of a list representation

#### 3.2.4 Method representation

Methods have two related AST nodes - *methodDefinition*, representing the definition of a method and *calledMethod* representing a call of a method.

First it should be determined if a vertex needs to be created for an *calledMethod* node. This is done by looking up based on the name of the *calledMethod* in *methodDefinition*, ignoring *THIS*. If the lookup is successful a vertex should be created as described below. If not, it should be checked

if the method is a method call on a top level property instance. This can be done by comparing if it starts in the same way as one of *topLevelProperties*. If that's the case, it is assumed, that it mutates whole property and the method should instead be treated as a write operation. If both of the above fail, the called method does not belong to context and is of no interest.

The next step is to resolve the names of the arguments it has been called with *calledArgs*. Every argument, that can neither be found in *computedProperties* nor *topLevelProperties* is replace by a fixed word such as OTHER or \*. In order to obtain the GUID of the vertex, the name of the *methodDefinitionIdentifier* is taken, THIS is excluded, and concatened with the the resolved arguments, which are joined with , and surrounded with brackets. The *label* of this node is equal to its name, excluding THIS from arguments.

The vertex for the method call is now completed. Multiple calls of this method with the same arguments will all result in the same vertex.

Now vertices for nodes the method interacts with, based on its *methodDefinition*, have to be created. Those include the variables it reads - *reads*, and writes - *writes* and methods it calls - *calls*.

Firstly, the arguments from the *methodDefinition* need to be substituted with the resolved arguments the method was actually called with and update all *reads*, *writes calls* referening them. All of them, which do not start with *THIS* can be discarded, since they do not belong to the context. Once filtered out, create a list of vertices for each variable in *reads* and *writes* as described in 3.2.1 and connect the most precise of those (the last of each list) to the method vertex. For the vertices resulting from *writes*, this edge has a label of 'writes' and the property vertex as the source and method vertex as the sink. For the vertices resulting from *reads*, this edge has the method vertex as the source and property vertex as the sink. Finally the process described in this section is repeated recursively for each *calledMethod* node in *calls* and an edge labeled 'calls' is added from the current method vertex to the resulting ones.

#### Computed property representation

Computed properties are represented similarly to methods, except they cannot have arguments, so no substitution of arguments is required. When defining their *label* and GUID both are equal to the *methodDefinitionIdentifier*. reads, writes and calls are computed in the same manner as methods.

#### Combining it all together

Interaction diagrams can be generated from the simplified Vue.js AST in the following way:

For each binding in bindings: - for each tag, bindingSource in binding:

Create a vertex for tag, with a GUID tagId and label name and type 'tag'.

if the *binding* is an *accessedVariable*, determine if it is a computed property by looking it up in *computedProperties* and if so, treat it as a computed property, and create vertices as described in 3.2.4. Otherwise determine if it as top level property, by doing a lookup on *topLevelProperties*,

treat it as a proeprty and create vertices for it as desribed in 3.2.1. In either cases, connect it to the tag vertex, based on the binding type. If the accessed Variable is neither, it does not belong to context and can be discarded.

if the *binding* is an *calledMethod* create a vertices for it as described in 3.2.4. Connect it to the *tag* vertex, based on the binding type.

Based on binding type, the following edges are created: A) If the binding type is an event binding, create an edge with the tag vertex as a source and the binding vertex as sink and label it 'event'. B) If the binding type is one-way, create an edge with the binding vertex as source and the tag vertex as sink. C) If the binding type is two-way, create both edges - A) and B).

For the initial method - createdMethod, create a vertex with GUID and name equal to created and create vertices for its reads, writes and calls analogous to methods as described in 3.2.4.

Once all of the above is done, additional edges will need to be added for the *all* vertices of properties of elements inside lists 3.2.3. Also the edges to the *any* vertex will be missing.

'numeric' vertices to the 'generic' one.

This is achievable by first finding all vertices of type 'generic' or 'numeric' and obtaining the parent of each of them. Those parents form a subset of all vertices, that have have 'generic' or 'numeric' vertices as children and adittionally other properties (properties on list elements, with  $type(v) \neq generic, type(v) \neq numeric$ ). For each of those parent vertices p: Firstly, connect each 'numeric' vertex to the 'generic' one. Recursively connect each child in the tree of the 'numeric' vertex to the child of the tree of the 'generic' vertex with the same name. If either does not exist, no edge is created. Do the same for p and all 'numeric' vertices and the 'generic' vertex.

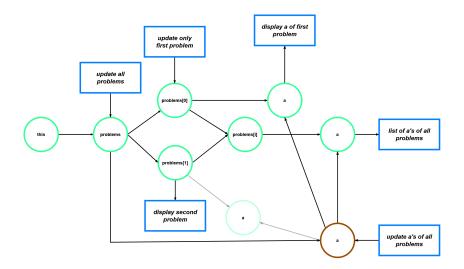


Figure 3.8: Example Graph including list elements with properties

#### 3.3 Scenario Generation

In order to generate scenarios in Gherkin, interactions can be sliced in a similar manner as described by Zhang and Zhao [ZZ19] and summarized in 2.2.2.

Let N denote the set of all nodes in the graph and E denote all edges in the graph. Let  $n \in N$ ,  $m \in N$  be any two nodes in the graph and  $(n, m) \in E$  represent an edge from n to m. Let type(n) be a function, that returns the type of a node and label(n, m) be a function that returns the label of the edge from n to m. Let  $E_{out}(n)$  be a function, which returns all outgoing edges of n. Let  $E_{in}(n)$  be a function, which returns all incoming edges to n.

Let  $N_I$  denote all nodes, that the user can interact with. A node  $n \in N$  is also in  $N_I$  if  $\exists e \in E_{out}(n)$  where label(e) = event. Let  $N_H$  denone all html tag nodes. A node  $n \in N$  is also in  $N_H$  if type(n) = tag. Given a node  $n_H \in N_H \cup created$ , a node  $m_H \in N_H$  reacts to  $n_H$  iff

1.  $\exists n_0, n_1, n_2, \dots, n_k \in N, n_0 = n_H, n_k = m_H \text{ such that for each } 0 \le i < k \ (n_i, n_{i+1}) \in E, \text{ and } label(n_i, n_{i+1}) \ne event \text{ and if } label(n_i, n_{i+1}) \ne calls \ \forall n_{i+1_{in}} E_{in}(n_{i+1}) \ label(n_{i+1_{in}}) \ne event$ 

Analogous to [ZZ19] let l(n) denotes all nodes, that react to n. A sequence of user interactions, starting with the initial function is referred to as a scenario  $A = (a_0, a_1, \ldots, a_n)$  where  $a_0 = created$  and  $\forall 0 < k \le n, a_k \in N_I$ . Define the HTML tags, to which a scenario reacts, to be equal to the tags to which the last tag in the scenario reacts  $l(A) = l(a_n)$ . Define a function that returns the last element in a scenario -  $last(A) = a_n$ 

The set of scenarios is generated by starting with the initial scenario, containing only the initial function  $S_0 = \{(created)\}$ . It is then prolonged by all tags  $n \in N_I$ , representing that the user can click anywhere. For further steps, only tags, that can be updated are included, so additionally  $n \in l(A)$  must hold. The newly included tag should also not be the same as the last element of the scenario, which means that also  $n \neq last(A)$  must hold. Formally: Define  $A \oplus x = a_0, a_1, \ldots, a_n, x$ . Then

$$S_{n} = \begin{cases} \{(created)\} & \text{if } n = 0 \\ \{p \oplus x | p \in S_{n-1}, x \in \mathcal{I}\} & \text{if } n = 1 \\ \{p \oplus x | p \in S_{n-1}, x \in l(p) \cap \mathcal{I}, x \neq last(p)\} \end{cases}$$
(3.1)

This is repeated up to k times, where k is a constant set by the user.

A Gherkin scenarios templated can then be obtained by the following template:

Scenario -  $n_0, n_1 \dots n_k$ 

Given -  $n_0, n_1 ... n_{k-1}$ 

When -  $n_k$ 

Then -  $l(n_k)$ 

## Chapter 4

# ${\bf Implementation}$

### 4.1 Project Structure and Overview

The application is written mostly in TypeScript using npm as a package manager and Node.js as a runtime environment.

Notable dependencies are TypeScript, for stricter syntax and types, lodash[21e] for enrichening of collections, graphlib[dag21] for the interaction diagrams graph, babel[21f] as a transcompiler, eslint[21b], estree[est21b] and eslint-plugin-vue[vue21b] for parsing Vue.js code, d3-graphviz[Jac21] in comination with light-server[Che21] for visualization of the generated interaction diagrams. The full list of dependencies can be found in the package.json of the project.

#### 4.1.1 Project structure

The main source files of the project and their tests are included in the *src* directory and structured in several packages, each corresponding to a step in the process, except common, which is shared among all steps. Each will be described in more detail in the following sections. Each package includes a models directory, which includes the data types defined and used in this section. The *web* directory contains code used to view the resulting diagram in the browser. The *scripts* directory includes helper bash scripts, *results* holds snapshots of the results throughout development (with the latest being the current) and *resources* hold various additional files needed.

```
root:.
     -- resources
        - output
      - results
      - scripts
        src
         -- common
10
            |-- models
12
         -- main.ts
         -- generator
            - models
         -- parsing
            - builders
            |-- models
            - visitors
```

### 4.2 Parsing Vue.js

Instead of implementing a parser, which directly outputs the simplified Vue.js AST described in 3.1.3 the capabilities of ESLint - Pluggable JavaScript linter [21b] and [vue21b] were used. The source files, which handle the parsing reside in the parsing directory. The ESLinter class provides a wrapper around the Node.js API of ESLint. Custom visitors are implemented in order to extract the necessary nodes from the AST of ESLint. Each visitor has a matching file in models, which holds the models specific to that visitor, and a builder, which keeps track of the visited nodes and builds the result data type. There are a total of three visitors - one for top level properties, another for bindings and the last one for method defintions, computed properties and the created method.

#### 4.2.1 Common Data Types

Below are the common data types used by all visitors.

```
export type Identifier = This | NameIdentifier | NumericIndex
GenericIndex;
4 interface BaseIdentifier {
    readonly name: string;
6
  export interface This extends BaseIdentifier {
    name: "this";
    discriminator: Identifier Type. THIS;
10 }
  export interface NumericIndex extends BaseIdentifier {
    discriminator: IdentifierType.NUMERIC INDEX;
12
14 export interface GenericIndex extends BaseIdentifier {
    discriminator: IdentifierType.GENERIC INDEX;
16
  export interface NameIdentifier extends BaseIdentifier {
    discriminator: IdentifierType.NAME IDENTIFIER;
```

```
export type Entity = Method | Property;

sexport interface Property {
   id: Identifiers;
   discriminator: EntityType.PROPERTY;
}

rexport interface Method {
   id: Identifiers;
   args: ReadonlyArray<Entity>;
   discriminator: EntityType.METHOD;
}
```

Discriminators are used to be able to differentiate between the types using type guards. The enums themselves are omitted here (*IdentifierType*, *EntityType*). The definitions here are not exactly the same as in the AST 3.1.3 - some constraints are omitted, such as method names having to end on a a *NameIdentifier*. This will be given, since the parsed code would be invalid javascript otherwise.

#### 4.2.2 Top Level Properties

The result of the top level properties has the following data type.

```
export type TopLevelProperties = Array<Property>;

export interface TopLevelPropertiesResult {
  topLevel: TopLevelProperties;
}
```

The top level properties visitor is the simplest of all since it only reacts to the top level data node inside the *script* object of the Vue.js SPA, which is a *ObjectExpression* 2.8.3. It can be selected via the following selector:

```
"ExportDefaultDeclaration > ObjectExpression >
Property[key.name = data] ReturnStatement > ObjectExpression"(node){
    ...
}
```

In natural language the selector reads: "Select ObjectExpression nodes, which have a direct parent ReturnStatement, that has an indirect parent Property with a property key.name equal to data and a direct parent ObjectExpression with a direct parent ExportDefaultDeclaration.

For each of the properties 2.8.3 of the *ObjectExpression* the name of the key (identifier) is stored. If the property is an object (value of *ObjectExpression* 2.8.3) it is concatenated with the previously obtained key. Finally all obtained properties are prefixed with 'this'.

#### 4.2.3 Bindings

The result of the bindings visitor has the following data type.

```
export enum BindingType {
      EVENT = "event", ONE_WAY = "one-way", TWO_WAY = "two-way",
    }
    export interface Tag {
      id: string;
      loc: Location;
      name: string;
      position?: string;
    }
    export interface BindingValue {
12
      item: Entity;
      bindingType: BindingType;
14
    }
16
    export type Binding = { tag: Tag; values: BindingValue[] };
    export interface BindingsResult {
      bindings: Binding[];
```

The ESlint AST nodes, which are interesting when parsing the bindings are *VElement 2.8.3*, *Identifier 2.8.3*, *MemberExpression 2.8.3* and *CallExpression 2.8.3*.

A identifier 3.1.3, abstracted in []common/identifier.ts can be a single Identifier or a MemberExpression, which can contain other Identifier nodes or MemberExpression nodes. Property Identifiers are extracted by finding the root MemberExpression or Identifier and traversing it. It is easy to determine if a MemberExpression or Identifier is the root - its parent is anything but a MemberExpression.

A CallExpression contains information about the name of the method and the arguments it has been called with, both in the form of nested MemberExpression and CallExpression nodes. Once again, only the root CallExpression node is extracted and converted to a calledMethod 3.1.3, abstracted in the Method interface in shared. ts.

VElement nodes represent any HTML tag, matching a tag 3.1.3 abstracted in the Tag interface in codetemplate-bindings.ts and contain information about the location of the tag and potentially a VText 2.8.3 node, which will be set as its name if it is present. If not present, the name of the tag is equal to the name of the first binding. Therefore, information about tags is extracted once a VELement is exited, since all bindings will be known at this point.

Further, the binding type has to be determined. This can be extracted based on the VAttribute 2.8.3. Event bindings have a VAttribute with a key.name.name equal to 'on', two-way bindings equal to 'model' and everything else is interpreted as one-way bindings. This includes moustache statements, v-bind, v-if bindings and all other except v-for statements. This filter be achieved via the powerful :not in combination with :matches selectors:

```
: not (: matches (
    VAttribute[key.name.name=on],
    VAttribute[key.name.name=model],

VAttribute[key.argument.name=key],
    VAttribute[key.name.name=for]))
```

With all the above, for example to match all two-way bindings and pass them on to the builder can be done via

```
"VAttribute[key.name.name=model] > VExpressionContainer

:matches(MemberExpression, Identifier, CallExpression)"(
    node) {
    if (utils.isRootNameOrCallExpression(node) &&
        utils.notArgument(node))
        builder.identifierOrExpressionNew(node, BindingType.TWO_WAY);
},
```

Bindings also need to substitute v-for statements. This is done by substituting the left side of the v-for statement with its right side and a generic index in all bindings that use it.

#### 4.2.4 Method Definitions

```
export interface MethodDefinition {
   id: Identifiers;
   args: ReadonlyArray<Property>;
   reads: ReadonlyArray<Property>;
   writes: ReadonlyArray<Property>;
   calls: ReadonlyArray<Method>;
}
s export type MethodDefinitions = Array<MethodDefinition>;

export interface MethodsResult {
   init?: MethodDefinition;
   computed: MethodDefinitions;
   methods: MethodDefinitions;
}
```

All method definition like structures (computed properties, created) and methods are parsed by the visitor defined in *methods.js*. Analogous to how the top level *data* object is selected 4.2.2, the *methods*, *created* and *computed* objects can be selected. The name of the method including its arguments can be extracted from by *Property[value.type=FunctionExpression]* nodes. Using this information, one can have three selectors, one of each type, to determine what is being defined. For example for regular methods:

```
"ExportDefaultDeclaration > ObjectExpression > Property
[key.name = methods] Property[value.type=FunctionExpression]"(node) {
   builder.newMethod(node, MethodType.METHOD);
},
```

Further the properties read, written and methods called need to be extracted. Methods called can be obtained by selecting CallExpression nodes. Properties written to can be obtained from the left side of a AssignmentExpression 2.8.3. There does not seem be an easy way to select all properties read from. Therefore all accessed identifiers are first stored and everything except reads, that can have an identifier (object properties, variable declarations, variables written to and names of called methods) is subtracted from it, in order to obtain the variables that the method reads from.

The following code can be used to obtain all variables written to by the current method-like in scope.

```
"ExportDefaultDeclaration > ObjectExpression >
2 :matches(Property[key.name = methods], Property[key.name = created],
Property[key.name = computed]) AssignmentExpression"(node) {
   builder.identifierOrExpressionNew(node.left, AccessType.WRITES);
},
```

#### 4.2.5 Output

Combinging all of the above, the following data structure is output.

```
1 export class Result {
    fileName: string;
3 topLevel: TopLevelPropertiesResult;
    methods: MethodsResult;
5 bindings: BindingsResult;
....
7 }
```

### 4.3 Interaction Diagram Generation

The generation of the interaction diagram graph from the result class from 4.2.5 is done in the *Transformer* class.

The resolution of methods is abstracted in the *MethodResolver* class. It produces a *ResolvedMethodDefintition* for each called method in bindings and the initial method. In order to prevent duplicate resolution of methods and wasting of resources a *MethodCache* is introduced. The *Transformer* does not use the *MethodResolver* directly, but instead accesses it via the *MethodCache*. The cache includes directly called (bound to) and indirectly called (calls of methods), for each of which vertices will have to be created. Each *ResolvedMethodDefintition* has the following data type:

```
export enum GeneralisedArgument {
    METHOD = "method", OTHER = "other",
  }
4 export type ResolvedArgument =
    | Property | Generalised Argument . METHOD | Generalised Argument . OTHER;
6 export interface ResolvedMethodDefintition {
    id: Identifiers;
    args: ReadonlyArray<ResolvedArgument>;
    reads: ReadonlyArray<Property>;
    writes: ReadonlyArray<Property>;
    calls: ReadonlyArray<CalledMethod>;
12
  export interface CalledMethod {
    id: Identifiers;
    args: ReadonlyArray<ResolvedArgument>;
16
```

As the underlying structure for the graph graphlib is used and wrapped in an own class - Extended Graph. It creates vertices on a 'create if not exist' basis by first looking up to see if the vertex exists in the graph, and if it does, does not add it again. Presence of edges is not checked, if an edge is added again, the previous one is simply overwritten. There can only be one edge per direction between two nodes, since no multigraph is used. Nodes and Edges in the graph have the following structure, as specified in 3.2:

```
export enum EdgeType {
   SIMPLE = "simple", EVENT = "event", CALLS = "calls",
}
export interface Edge {
   source: Node;
   sink: Node;
   label: EdgeType;
```

8

```
export enum NodeType {
   TAG = "tag", DATA = "data", METHOD = "method", INIT = "init",
  export type Node = TagNode | DataNode | MethodNode | InitNode;
6 interface BaseNode {
    id: string;
    name: string;
  }
10 export interface TagNode extends BaseNode {
    loc: Location;
12
    discriminator: NodeType.TAG;
  }
14 export interface MethodNode extends BaseNode {
    discriminator: NodeType.METHOD;
16
  export interface InitNode extends BaseNode {
    discriminator: NodeType.INIT;
  }
20 export interface DataNode extends BaseNode {
    parent?: string;
    type: IdentifierType;
    discriminator: NodeType.DATA;
24
```

The algorithm for generating the interaction diagrams graph as described in 3.2.4 is implemented in the *Transformer* class.

First the init method is resolved by querying the MethodCache and afterwards each of the bindings. The order is not important. If the bindings are properties, the vertices for them are created directly. If they are methods or computed properties, they are resolved by the MethodCache and the correct edges based on the binding type created. The above is done in the addInit() and addBindings() methods.

At this stage no vertices are created for the reads, writes, and calls properties of the resolved methods. This happens in the addIndirectlyCalledMethods() method, after all bindings and the init method have been resolved by taking all methods stored in the MethodCache and creating the appropriate vertices.

Lastly in the addEdgesForLists method additional edges will be added for the all vertices of properties of elements inside lists.

#### 4.4 Scenario Generation

The scenario generation is implemented in the *scenario*. ts file and currently outputs the obtained scenarios, tags, nodes react to and the gherkin scenarios templates to the console. The scenarios are generated exactly as described in 3.3

### 4.5 Display of Interaction Diagram

In order to display the interaction diagrams it is displayed on a single html page, utilizing the d3-graphviz [Jac21] library.

Event edges are displayed as thick bold lines, calls edges are displayed as solid lines and all other as dashed lines.

Html tags are separated from all other nodes in a subgraph in order to be able to more easily identify them. In addition to their name, they also include the line(s) on which they are in the source code as a subline. The initial function has a grey background so it can be identified right away. A button to download the displayed diagram as a .png image is also provided in the top left.

### 4.6 Usage

The application can be started using

```
npm run generate — [file to parse] [output graph path] [depth]
```

where file to parse is the path to the .vue file to be parsed and output graph path is the path to which the output interaction diagram graph will be written (JSON object). Depth is the factor k in 3.3. A web server can be startet in order to view the results in a browser at localhost:8000 via

```
npm run serve
```

In order to create a snapshot for each file in resources / test - files.

```
npm run create-results
```

Can be used. This snapshot includes the output interaction diagram graph as data.json, the generated scenarios as a scenario.txt, index.html, which can be used to display the interaction diagram and a .vue file for which all of the above was generated.

All snapshots can be viewed using

```
npm run results
```

and navigating to localhost:8001

## Chapter 5

# Testing and Evaluation

- 1. zhang code, generated
- 2. example with lists 3. complex list(add, sub) with focus on updates

#### 5.1 Scenario

Zhang and Zhao [ZZ19] use an example of a small SPA titled "Math for kids", which generates simple math addition problems and keeps track of a small statistic, which includes the number of correctly and incorrectly answered problems. The same application has been implemented and extended in Vue.js and will be used to display the capabities of the interaction diagram and scenario generation application.

### 5.2 Math for Kids



(a) Math for Kids in Angular JS by Zhang and Zhao [ZZ19]  $\,$ 



(b) Own implementation of Math for Kids in Vue.js

The source code can be found in resources/test-files/test.vue and also in A.2.1.

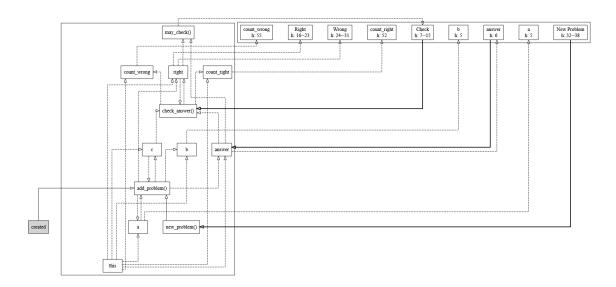


Figure 5.2: Math for Kids in Vue.js generated interaction diagram

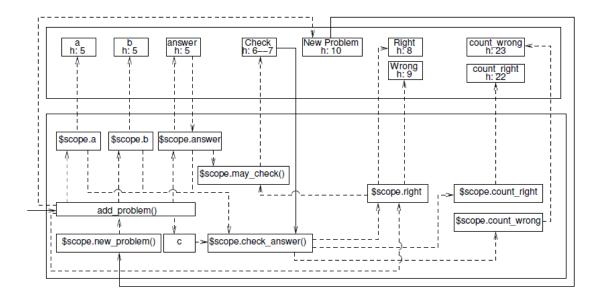


Figure 5.3: Math for Kids in AngularJS interaction diagram by Zhang and Zhao [ZZ19]

When comparing the generated diagram 5.2 to the original by Zhang and Zhao [ZZ19] 5.3 there are some differences.

Due to different modelling there is a *this* vertex and connections from it to top level variables, which does not exist in the diagram by Zhang and Zhao [ZZ19]. In the generated diagram there is an edge of type 'event' between the *answer* tag and *answer* property. This is due to the way that two-way bindings were decided to be represented.

There are two differences between the add\_problem node in the generated diagram and the one by zhang2019scenario. In Zhang and Zhao [ZZ19] there is an edge from the add\_problem method to the New Problem tag, which is missing in the own generated version. This seems like a mistake in [ZZ19], since this edge should not exist.

The other difference is that in [ZZ19] there is no edges from add\_problem to right (write relation), however there should be, inside add\_problem this.right = undefined.

```
l(created) -> a, b, answer, Check, Right, Wrong
l(answer) -> answer, Check
l(Check) -> Right, Wrong, Check, count_right, count_wrong
l(New Problem) -> a, b, answer, Check, Right, Wrong
```

The generated reaction based on the graph differ for l(created), all other sets are exactly the same as the ones by Zhang and Zhao [ZZ19] (albeit in different order). This may stem from the fact that the edges of  $add\_problem$  to  $\$scope.\ right$  is missing in [ZZ19], however it is later correctly factored in l(NewProblem).

Zhang and Zhao [ZZ19] define  $l(add_problem())$  (the equivalent of l(created)) as  $l(add_problem() = a, b)$  where it should be  $l(add_problem() = Check, Right, Wrong, a, b, answer)$  since the answer property is updated based on the diagram.

One could argue, that the version in [ZZ19] is correct, since answer is set to undefined, which

is the same as the initial value of the variable, so it would not trigger an update as part of the *init* method. The interaction diagram generator does not perform this check. If  $l(add_problem())$  were to be called later in the application (when the NewProblem button is clicked) it would indeed set a new value to answer, which is correctly reflected by Zhang and Zhao [ZZ19].

#### 5.2.1 Scenarios

Gherkin scenarios of up to 4 actions are generated, which can be seen in the following figures. The program outputs the scenarios as plain text to the console, but here they are displayed in a nicer way. The caption of each figure is an example of how a text could be generated based on the template scenario output by the application.

Some scenarios seem a bit repetitive, but it is up to *The Three Amigos* - The product owner, tester and developer [21g] to decide which templates to discard. Since interaction diagrams model what might be updated, it will probably be reasonable to discard some scenarios. The fact that everything which might get updated can also be leveraged in another way - negative criteria can be definied (assert tag/component X did not change).

```
Scenario:[created]

When 'created'
Then 'a'
And 'b'
And 'answer'
And 'Check'
And 'Right'
And 'Wrong'
```

(a) Scenario: initialization. When the application is created, then 'a' and 'b' should show random numbers and 'Check' should be disabled and 'Right' and Wrong should be invisible.

(c) Scenario: clicking on check without typing an answer. Given the application has been created, when I click 'Check' then 'Right' and 'Wrong' should be invisible and 'count\_right' and 'count\_wrong' should have the same values and 'Check' should be disabled.

(b) Scenario: typing an answer. Given the application has been created, when I type an answer then 'answer' should display it and 'Check' should be enabled.

(d) Scenario: obtaining a new problem at application start. Given the application has been created, when I click 'New Problem' then 'a' and 'b' should have random values and 'answer' should be empty and 'Check' should be disabled and 'Right' and 'Wrong' should be invisible.

(e) Scenario: checking my answer. Given the application has been created and 'answer' contains my answer, when I click 'Check' then 'Right' should be visible if the answer was right 'Wrong' should be visible if the answer was wrong and 'Check' should be disabled and 'count\_right' should be incremented by one if the answer was right and 'count\_wrong' should be incremented by one if the answer did not change.

(g) Scenario: requesting a new problem and clicking on check without answering it. Given the application has been created and I requested a new problem, when I click 'Check' then 'Right' and 'Wrong' should be invisible and 'count\_right' and 'count\_wrong' should have the same values and 'Check' should be disabled.

(f) Scenario: typing an answer to a new problem. Given the application has been created and a new problem was obtained, when I type an answer then 'answer' should display it and 'Check' should be enabled.

(h) Scenario: requesting a new problem, answering it and checking my answer. Given the application has been created and I requested a new problem and answered it, when I click 'Check' then 'Right' should be visible if the answer was right 'Wrong' should be visible if the answer was wrong and 'Check' should be disabled and 'count\_right' should be incremented by one if the answer was right and 'count\_wrong' should be incremented by one if the answer did not change.

### 5.3 Diagram with list

which can be found in resources / test - files / list. vue

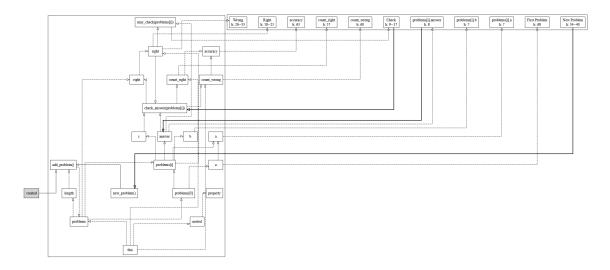


Figure 5.5: TODO

## Chapter 6

### Conclusion

- lists recursive use directly the ast of eslint for true solution
- betterr output not just console
- also present if else and nested methods idea?

interaction diagram simple, parsing harder, need to diferentiate blocks... hmm actually that's fairly straight forward

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### Appendix A

### Appendix

#### A.1 Source Code

Together with this thesis a .zip file is submitted, which contains the source code.

### A.2 Math for Kids

#### A.2.1 Source Code

```
<template>
      < div >
        <div style="margin-left: 5px; margin-top: 20px">
           <form @submit.prevent>
             < span > \{\{ a \}\} < / span > + < span > \{\{ b \}\} < / span > = \}
             <span><input class="question" v-model="answer" /></span>
             <button
                class="btn btn-primary"
                style = "margin-left: 5px"
                : disabled = "!may check()"
                type="submit"
                @ click="check_answer()"
               Check
14
             </button>
             < button
16
                class="btn btn-success"
                style = "margin - left: 5px"
                disabled = "1"
               v-if = "right = true"
```

```
Right
22
              </button>
              <button
24
                class="btn btn-danger"
                style="margin-left: 5px"
                disabled = "1"
                v-if = "right = false"
28
                Wrong
30
              </button>
              <button
32
                 class="btn btn-info"
                style="margin-left: 5px"
                 @click="new problem()"
36
                New Problem
              </button>
38
           </form>
         </div>
40
         <hr />
         < div style = "margin-left: 5px">
42
           <table>
              <tr>
44
                <strong>Statistics&nbsp;</strong>
                <\!td\!>\!<\!span\ class\!=\!"badge\ badge\!-\!success"\!>\!Right<\!/\!span\!>\!<\!/td\!>
46
                <\!td\!>\!<\!span\ class\!=\!"badge\ badge\!-\!danger"\!>\!Wrong\!<\!/span\!>\!<\!/td\!>
              </\mathrm{tr}>
48
              <tr>
                50
                <td>
                  <span> {{ count_right }}</span>
52
                <td>
                  <span> {{ count_wrong }}</span>
                56
              </t\,r>
           < / table >
58
         </div>
```

```
</div>
    </template>
62
    <script>
    export default {
64
      name: "HelloWorld",
      props: {
66
        msg: String,
      },
68
      data() {
        return {
          a: 0,
          b: 0,
72
          c: 0,
          answer: undefined,
74
           right: undefined,
          count right: 0,
76
          count_wrong: 0,
        };
78
      },
      created: function () {
        this.add problem();
      },
82
      methods: {
        add problem() {
84
           let \max = 100;
           this.c = Math.floor(Math.random() * (max - 1)) + 1;
86
           this.a = Math.floor(Math.random() * (this.c -2)) + 1;
           this.b = this.c - this.a;
           this.answer = undefined;
90
           this.right = undefined;
        },
92
        check_answer() {
           this.right = this.c === parseInt(this.answer);
94
           this.right ? (this.count_right += 1) : (this.count_wrong += 1);
        },
96
        may_check() {
           // answer non-empty and right undefined
```

```
return (
              !(this.answer === undefined || this.answer === "") &&
100
             this.right === undefined
           );
102
         },
         new_problem() {
104
           this.add_problem();
         },
106
       },
     };
108
    </script>
    <style scoped>
110
    .index {
       padding: 0 20px 0 5px;
112
    }
114
     .question {
       font-color: white;
116
       width: 2em;
       background-color: #3ff;
118
       margin-left: 5px;
       margin-right: 5px;
120
    }
    </style>
```