

CODE OPTIMIZATION

(Improve target code)

- ✓ Run faster
- ✓ Take less space

Machine independent code optimization
(Does not require properties of target machine)

NOT REQUIRED
Machine instruction
Registers allocation

Dead Code elimination + Variable propagation

```
c = a * b;  
x = a;  
for (i=1; i<=N; i++)  
{  
    d = (x * b) + 1;  
}
```

```
c = a * b;  
for (i=1; i<=N; i++)  
{  
    d = c + i;  
}
```

Code motion: Bring loop invariant statements outside

```
for (i=1; i<=N; i++)  
{  
    x = 10;  
    y = x + i;  
}
```

```
x = 10;  
for (i=1; i<=N; i++)  
{  
    y = x + i;  
}
```

```
for (i=1; i<=N; i++)  
{  
    x = 10;  
    y = x + 3;  
}
```

```
x = 10;  
y = x + 3;
```

eliminate loop.

(dead code ~~elimination~~
elimination)

Reducing strength

```
i = 1;  
while (i < 10)  
{  
    y = i * 4;  
    i++;  
}
```

'*' is replaced
by '+'

```
i = 1;  
t = 4;  
while (t < 40)  
{  
    y = t;  
    t = t + 4;  
}
```

Booth Algorithm

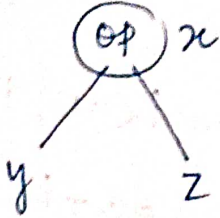
```
i = 1;  
t = 0;  
while (t < 40)  
{  
    t = t + 4;  
}
```

y = t;

DIRECTED ACYCLIC GRAPH (DAG)

Case-1

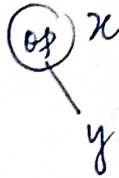
$$x = y \text{ op } z$$



Case-2

$$x = \text{op } y$$

or $x = x \text{ op } y$

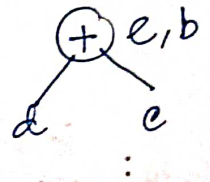


Case-3

$$x = y$$



eg. $e = d \times c$
 $b = e$

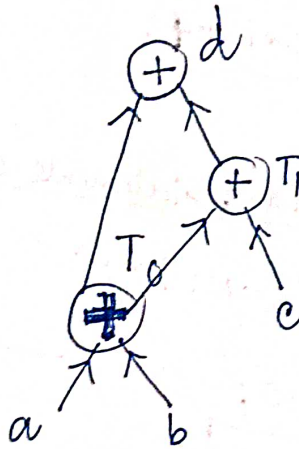


Eg.-1

$$T_0 = a + b$$

$$T_1 = T_0 + c$$

$$d = T_0 + T_1$$



Eg.2

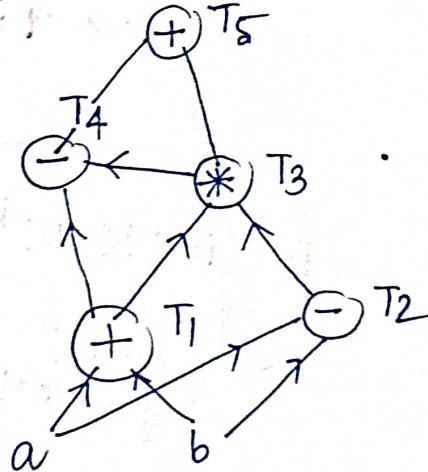
$$T_1 = a + b$$

$$T_2 = a - b$$

$$T_3 = T_1 * T_2$$

$$T_4 = T_1 - T_3$$

$$T_5 = T_4 + T_3$$



Eg.3 ① $a = b * c$ ✓

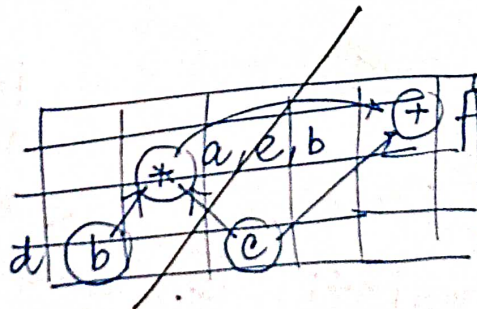
② $d = b$ ✓

③ $e = d * c$

④ $b = e$

⑤ $f = b + e$

⑥ $g = f + d$

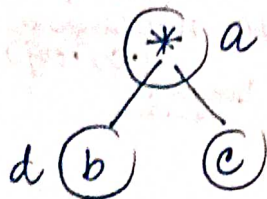


$$a = b * c$$



$$a = b * c$$

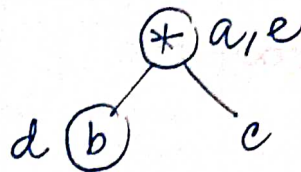
$$d = b$$



$$a = b * c$$

$$d = b$$

$$e = d * c$$

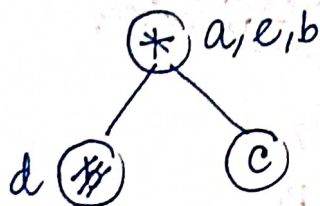


$$a = b * c$$

$$d = b$$

$$e = d * c$$

$$b = e$$



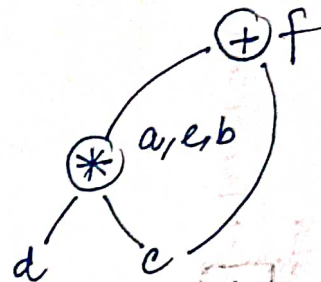
$$a = b * c$$

$$d = b$$

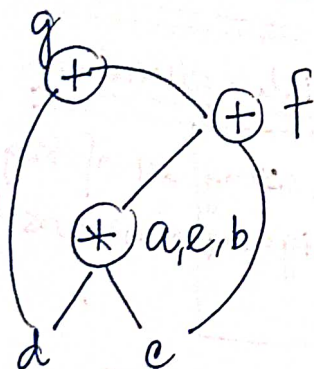
$$e = d * c$$

$$b = e$$

$$f = b + c$$



$$g = f + d$$



$$a = d * c$$

$$f = a + c$$

$$g = d + f$$

VARIABLE
PROPAGATION

```

c = a * b;
x = a;
for (i = 1; i <= N; i++)
{
    d = x * b + i;
}
    
```

$c = a * b;$

$x = a;$

for ($i = 1; i \leq N; i++$)

{ $d = c + i;$

}

DEAD CODE
ELIMINATION

Example of
reducing strength
* elimination

```

c = a * b;
for (i = 1; i <= N; i++)
{
    d = c + i;
}
    
```


Peephole optimization

Redundant Instruction Elimination

```
int f(int x)
{
    int y, z;
    y = 10;
    z = x + y;
    return z;
}
```

Vars: x, y, z
instruction: 04

```
int f(int x)
{
    int y;
    y = 10;
    y = y + x;
    return y;
}
```

```
int f(int x)
{
    int y = 10;
    return x + y;
}
```

```
int f(int x)
{
    return x + 10;
}
```

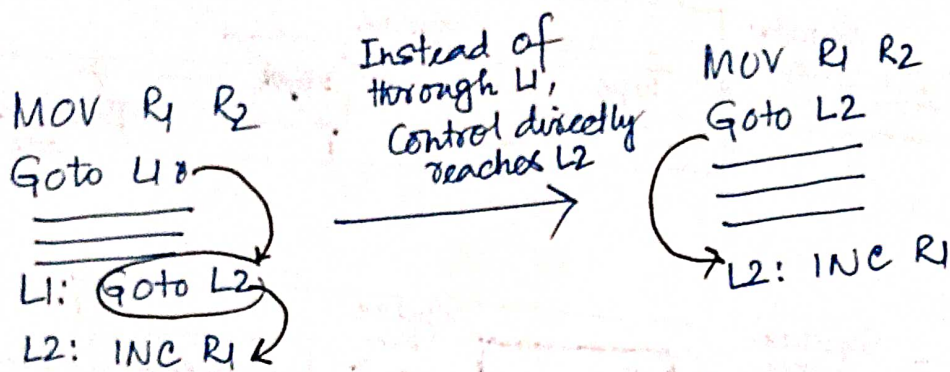
Vars: x
instruction: 01

Unreachable Code

```
void f int f(int x)
{
    return x + 10;
    printf("value of %.d", x + 10);
}
```

Unreachable code
(can be removed)

✓ Flow of control optimization



Algebraic Simplification

$a = a + 0$ — can be removed.

$a = a + 1$ → $\frac{\text{INC } a}{a++}$ ✓ (increment a)

$x = 2y$ → $x = y + y$ because + is lightweight than *.