

## LAB-4

### Integer multiplication by Divide-and-Conquer

Karatsuba  $(x, y)$

// To compute product of 2 nos.

// Input: Two non-negative int  $x$  &  $y$

// output: Product of  $x$  &  $y$

if  $x < 10$  or  $y < 10$

return  $x * y$

else

$n = \text{max length of } (x, y)$

$a, b = \text{first and second half of } x$

$c, d = \text{first and second half of } y$

compute  $p = a + b$  and  $q = c + d$

recursively compute  $ac = a \cdot c$

$bd = b \cdot d$  and  $pq = p \cdot q$

compute  $ad + bc = pq - ac - bd$

compute  $10^n \cdot ac + 10^{n/2} \cdot (ad + bc) + bd$

### TC of Karatsuba Algorithm

This algorithm recursively computes 3 products each of size half of input, adds in  $O(n)$  for combining the localised optimal solutions.

So recurrence relation will be

$$T(n) = 3T(n/2) + O(n)$$

Applying master theorem,

$$a = 3$$

$$b = 2$$

$$d = 1$$

$$3 > 2^1 \quad 3 > 2$$

$$\begin{aligned} \therefore T(n) &= O(n^{\log_2 3}) \\ &= O(n^{1.585}) \end{aligned}$$

The Brute-force to calculate the product takes  $O(n^2)$  time as each digit is multiplied with all other digits.

$\therefore$  Karatsuba multiplication is efficient approach as compared to brute-force technique.



## Testcases (Integer multiplication)

### Positive

- 1)  $x = 12$        $y = 34$   
 $xy = 408$
- 2)  $x = 123$        $y = 456$   
 $xy = 56088$
- 3)  $x = 789$        $y = 123$   
 $xy = 97047$
- 4)  $x = 12345$        $y = 67890$   
 $xy = 838102050$
- 5)  $x = '1203456789'$        $y = 9876543210$   
 $xy = 121932631112635269$

### Negative

- 1)  $x = -12$        $y = 34$   
 $xy = -408$
- 2)  $x = 123$        $y = -456$   
 $xy = -56088$
- 3)  $x = -789$        $y = -123$   
 $xy = 97047$
- 4)  $x = -12345$        $y = 67890$   
 $xy = -838102050$
- 5)  $x = -987654321$        $y = -123456789$   
 $xy = 121932631112635269$

## Counting Inversions Using Divide and Conquer

\* This f<sup>n</sup> is for recursive division step  
SortCountInv (A[n]) :

// Recursively count inversions in left & right subarray  
// Input: array A of n distinct integers.  
// Output: sorted array B and no. of inversions of A.

if  $n == 0$  ||  $n == 1$

return (A, 0) ;

else

(C, leftInv) = SortCountInv (first half of A)

(D, rightInv) = SortCountInv (second half of A)

(B, splitInv) = Merge and SortCountInv (C, D)

return (B, leftInv + rightInv + splitInv)

### Merge Sort CountInv

// To return the sorted array B and Counting Inversions

// Input: Sorted array C and D

// Output: Sorted array B and no. of counting Inversions.

// ASSUMPTION : n is even

$i = 1$      $j = 1$  ,    splitInv = 0

for  $k = 1$  to  $n$  do

if  $C[i] < D[j]$  ,

$B[k] = C[i]$  ;  $i++$

else

$B[k] = D[j]$  ;  $j++$

splitInv +=  $\left(\frac{n}{2} - i + 1\right)$

return (B, splitInv)



## Time Complexity for Divide and Conquer

Divide: It computes the middle of the subarray which takes constant time  $O(1)$

Conquer: We recursively solve 2 subproblems, each of size  $n/2$ , which contributes to  $2T(n/2)$

Combine: Here, the entire merge procedure takes  $O(n)$  such that  $c(n) = O(n)$

$$\therefore T(n) = \begin{cases} \Theta(1) & , \text{ if } n=1 \\ 2T(n/2) + \Theta(n) & , \text{ if } n>1 \end{cases}$$

$$T(n) = \begin{cases} c & , \text{ if } n=1 \\ 2T(n/2) + cn & , \text{ if } n>1 \end{cases}$$

Using master theorem

$$T(n) = aT(n/b) + \Theta(n^d)$$

$$a = 2, \quad b = 2, \quad d = 1$$

$$e = b^d$$

$$\therefore T(n) = O(n^d \log n)$$

$$\therefore T(n) = O(n \log n)$$

## Counting Inversion using Brute-Force Technique

CountInv ( $A[0, 1, \dots, n]$ )

// Count the total no. of inversion in given array

// Input: array A of n distinct integers

// Output: The number of inversions of A.

numInv = 0

for  $i = 1$  to  $n-1$  do

for  $j = i+1$  to  $n$  do

if ( $A[j] > A[i]$ )

numInv ++;

return numInv;

## Time Complexity of Brute-Force

The Summation eq<sup>n</sup> is

$$T(n) = \sum_{i=1}^{n-1} \sum_{j=i+1}^n 1$$

$$= \sum_{i=1}^{n-1} (n - i + 1 - 1)$$

$$= \sum_{i=1}^{n-1} (n - i)$$

$$= \sum_{i=1}^{n-1} 1 = \frac{n(n-1)}{2}$$

$$= \frac{3}{2}n^2 - \frac{3}{2}n$$



## Testcases (Counting Inversion)

### Positive

- 1) codes = [84205, 55881, 26415, 66611, 64035]  
o/p = Student 1: Brute-force inversions = 6,  
Optimised inversions = 6
- 2) codes = [21648, 86929, 82205, 65035, 80421]  
o/p = Student 2: Brute-force inversions = 5,  
Optimised inversions = 5
- 3) codes = [99915, 52246, 49670, 17129, 73238]  
o/p = Student 5: Brute-force inversions = 7,  
Optimised inversions = 7
- 4) codes = [52327, 47445, 12180, 91815, 20328]  
o/p = Student 7: Brute-force inversions = 6,  
Optimised inversions = 6
- 5) codes = [75649, 85312, 72434, 11379, 91516]  
o/p = Student 8: Brute-force inversions = 5,  
Optimised inversions = 5

### Negative

- 1) codes = [26122, 65915, 24881, 61651, -96661]  
o/p: Invalid course code "-96661" found.
- 2) codes = [F#)3h, 80374, 66321, 27208, 63188]  
o/p: Student 4: Non-numeric invalid course code  
"F#)3h" found.

3) codes = [-21203, rig%1, 100-i, 94596, nd]0%5]  
o/p: student 6: non-numeric invalid code "-21203"  
found.

4) codes = [-57974, 76465, 73778, 95316, C.^wr]  
o/p: student 9: Non-numeric invalid code "-57974"  
found.

5) codes = [97569, 24732, 65650, }([j+, 79904]  
o/p: student 10: Non-numeric invalid code "}([j+ "  
found.