Dynamic Programming Problems Solved Today

1. Minimum Falling Path Sum

Problem:

Given a square matrix matrix, find the minimum sum of a falling path from the top row to the bottom. A falling path can move straight down, diagonally left, or diagonally right.

Approaches:

1 Greedy Attempt (Incorrect for general case):

- Started from the minimum element in the first row.
- Moved to the smallest adjacent element in the row below.
- Flaw: Greedy doesn't guarantee global minimum.

Bottom-Up DP (Correct)

- Modify the matrix in place from bottom to top.
- For each cell, add the minimum of the three possible paths from the row below.

Clue:

Use the i+1 row to compute the optimal choice for row i. Start from the second-last row and move upward.

✓ Code:

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class Solution {

public:

```
int minFallingPathSum(vector<vector<int>>& matrix) {
  int n = matrix.size();
  for (int i = n - 2; i >= 0; i--) {
    for (int j = 0; j < n; j++) {
      int down = matrix[i + 1][j];
    }
}</pre>
```

```
int left = (j > 0) ? matrix[i + 1][j - 1] : INT_MAX;
int right = (j < n - 1) ? matrix[i + 1][j + 1] : INT_MAX;
matrix[i][j] += min({down, left, right});
}

return *min_element(matrix[0].begin(), matrix[0].end());
}
};</pre>
```

3 Top-Down Recursive + Memoization (Correct)

Clue:

Use recursion from row 0 with memoization table dp[i][j] storing the result of the best path starting from cell (i, j).

Code:

```
cpp
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class Solution {
public:
    int n;
    vector<vector<int>>> dp;

int solve(vector<vector<int>>> & matrix, int i, int j) {
    if (j < 0 || j >= n) return INT_MAX;
    if (i == n - 1) return matrix[i][j];
    if (dp[i][j] != INT_MAX) return dp[i][j];

    int down = solve(matrix, i + 1, j);
    int left = solve(matrix, i + 1, j - 1);
```

```
int right = solve(matrix, i + 1, j + 1);

return dp[i][j] = matrix[i][j] + min({down, left, right});
}

int minFallingPathSum(vector<vector<int>>& matrix) {
    n = matrix.size();
    dp = vector<vector<int>>(n, vector<int>(n, INT_MAX));
    int result = INT_MAX;

for (int j = 0; j < n; j++) {
    result = min(result, solve(matrix, 0, j));
    }
    return result;
}</pre>
```

Key Takeaways:

Clues for Identifying DP:

- Overlapping subproblems? => Yes → Use recursion + memoization.
- Optimal substructure? => Yes → Use bottom-up tabulation.
- Multiple paths/states converging to a goal → Matrix DP often applies.
- The problem wants a minimum/maximum across a path or choices → Think DP.

Common Pitfalls to Avoid:

- X Using INT_MAX directly in additions without checking overflow.
- X Assuming greedy works for global minimum problems.

• X Forgetting bounds when accessing matrix[i][j - 1] or j + 1.

1. Max Frequency Difference in Substring

Topic: Sliding Window + Frequency Count

• Problem:

Given a string s and integer k, find a substring of length \geq k where the difference freq[a] - freq[b] is maximized with:

- a has odd frequency
- b has even frequency

Approaches Tried:

- Brute force with nested loops over all substrings.
- Frequency count using a vector of size 5 (since only digits '0' to '4').

Issues:

• X TLE: Brute force over all substrings is $O(n^2 * 5)$ — too slow for $n \approx 30,000$.

Clue:

Use a sliding window or optimized prefix frequency table with hash map trick. Keep substring length $\geq k$.

2. Coin Change (Recursive)

Topic: DP – Recursion with Pruning

Problem:

Find the minimum number of coins needed to make up a given amount.

Approaches Tried:

- Pure recursion with backtracking.
- Used a priority queue to track minimum count.

Issues:

• X TLE & Overflow: curr + coins[i] exceeded INT_MAX.

Clue:

Use memoization to cache (i, curr) pairs.

Avoid pushing into priority queue inside recursion — not efficient.

3. Partition Linked List

Topic: Linked List Manipulation

Problem:

Reorder linked list such that all nodes with value < x come before nodes $\ge x$.

Issues:

- X Infinite loop / missing curr = curr->next in some branches.
- New nodes were unnecessarily created; original structure can be reused.

Clue:

Use two dummy lists: one for $\leq x$, one for $\geq x$. Connect them at the end.

4. Minimum Falling Path Sum

Topic: DP on Grid

Problem:

Find the path with minimum sum starting from any cell in the first row and moving down with left-down, down, and right-down moves.

Approaches:

- 1. X Greedy based on smallest in first row incorrect globally.
- 2. Sottom-Up DP modifying matrix.
- 3. Recursive + Memoization (top-down).
- 4. X Overflow when using INT_MAX + matrix[i][j].

• Clue:

Use DP table or modify input matrix. Never add to INT_MAX directly.

Summary Table:

Problem	Topic	Status	Key Issue/Clue
Max Frequency Difference in Substring	Sliding Window + Freq	X TLE	Needs optimized counting with window or prefix
Coin Change (Recursive)	Recursion + DP	X Overflow	Needs memoization & overflow handling
Partition Linked List	Linked List	Fixed	Dummy lists, remember to update pointers
Minimum Falling Path Sum	Grid DP	✓	Bottom-up DP or top-down memoized recursion

1. Max Frequency Difference in Substring

Problem:

Find a substring of length \geq k in s where:

- One character a has **odd frequency**.
- Another character b has even frequency.
 Maximize freq[a] freq[b].
- Status: X TLE on brute force approach
- Clue:

Instead of checking all substrings, use a **sliding window** or **prefix frequency** to optimize.

No fully optimized code attempted today — here's a future direction:

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```
int maxDifference(string s, int k) {
  int n = s.size();
  int maxDiff = INT_MIN;

for (int i = 0; i <= n - k; i++) {
   vector<int> count(5);
```

```
for (int j = i; j < n; j++) {
      count[s[j] - '0']++;
      if ((j - i + 1) >= k) {
        int maxOdd = INT_MIN, minEven = INT_MAX;
       for (int f = 0; f < 5; f++) {
         if (count[f] % 2 == 1) maxOdd = max(maxOdd, count[f]);
         if (count[f] % 2 == 0 && count[f] != 0) minEven = min(minEven, count[f]);
       }
        if (maxOdd != INT_MIN && minEven != INT_MAX) {
          maxDiff = max(maxDiff, maxOdd - minEven);
       }
     }
   }
  }
  return (maxDiff == INT_MIN) ? -1 : maxDiff;
}
```

2. Coin Change – Recursive Version

• Problem:

Find minimum number of coins to make a given amount.

- Status: X TLE / Overflow in pure backtracking
- Optimized Recursive Code with Memoization:

```
cpp
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class Solution {
public:
```

```
int solve(vector<int>& coins, int amount, unordered_map<int, int>& dp) {
    if (amount == 0) return 0;
    if (amount < 0) return INT_MAX;
    if (dp.count(amount)) return dp[amount];
   int ans = INT_MAX;
   for (int coin: coins) {
     int res = solve(coins, amount - coin, dp);
     if (res!=INT_MAX) ans = min(ans, res + 1);
   }
    return dp[amount] = ans;
 }
  int coinChange(vector<int>& coins, int amount) {
    unordered_map<int, int> dp;
   int result = solve(coins, amount, dp);
    return (result == INT_MAX) ? -1 : result;
 }
};
```

3. Partition Linked List

Problem:

Reorder the linked list such that nodes $\leq x$ come before nodes $\geq x$.

- Status: Successfully debugged
- Final Code:

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```
class Solution {
public:
  ListNode* partition(ListNode* head, int x) {
    ListNode* before = new ListNode(-1);
   ListNode* after = new ListNode(-1);
   ListNode* b = before;
    ListNode* a = after;
   while (head) {
     if (head->val < x) \{
        b->next = head;
       b = b - next;
     } else {
       a->next = head;
       a = a - next;
     }
     head = head->next;
   }
   a->next = nullptr;
    b->next = after->next;
   return before->next;
 }
};
```

4. Minimum Falling Path Sum

• Problem:

Find the minimum falling path sum from top to bottom of a square matrix.

Status: Solved with both bottom-up and recursive memoized DP

Bottom-Up Tabulation:

```
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class Solution {
public:
  int minFallingPathSum(vector<vector<int>>& matrix) {
    int n = matrix.size();
    for (int i = n - 2; i \ge 0; i = 0; i = 0) {
      for (int j = 0; j < n; j++) {
        int down = matrix[i + 1][j];
        int left = (j > 0) ? matrix[i + 1][j - 1] : INT_MAX;
        int right = (j < n - 1)? matrix[i + 1][j + 1]: INT_MAX;
        matrix[i][j] += min({down, left, right});
      }
    }
    return *min_element(matrix[0].begin(), matrix[0].end());
 }
};
```

Recursive + Memoization:

```
cpp
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class Solution {
public:
  int n;
  vector<vector<int>> dp;
```

```
int dfs(vector<vector<int>>& mat, int i, int j) {
  if (j < 0 || j >= n) return INT_MAX;
  if (i == n - 1) return mat[i][j];
  if (dp[i][j] != INT_MAX) return dp[i][j];
  int down = dfs(mat, i + 1, j);
  int left = dfs(mat, i + 1, j - 1);
  int right = dfs(mat, i + 1, j + 1);
  return dp[i][j] = mat[i][j] + min({down, left, right});
}
int minFallingPathSum(vector<vector<int>>& mat) {
  n = mat.size();
  dp = vector<vector<int>>(n, vector<int>(n, INT_MAX));
  int result = INT_MAX;
  for (int j = 0; j < n; j++) {
    result = min(result, dfs(mat, 0, j));
  }
  return result;
}
```

};

Problem	Topic	Status	Key Fix
Max Frequency Difference in Substring	Sliding Window	X TLE	Optimize with prefix or freq tracking
Coin Change (Recursive)	Recursion + Memo	~	Memoization + Overflow check
Partition Linked List	Linked List	<u>~</u>	Use dummy nodes
Minimum Falling Path Sum	Grid DP	<u>~</u>	Bottom-Up / Recursive DP