

# Binary Tree Algorithms - Advanced Patterns

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## 1. Burn a Binary Tree from a Leaf Node

### Problem

Given a binary tree and a start node, simulate infection/burning of all nodes minute by minute. Return the time it takes to infect the entire tree.

### Approach

1. Mark Parents using BFS
2. Do BFS from the start node and infect unvisited neighbors (left, right, parent)
3. Track minutes as levels

### Code

```
int amountOfTime(TreeNode* root, int start) {
    unordered_map<TreeNode*, TreeNode*> parent;
    queue<TreeNode*> q;
    TreeNode* source = nullptr;

    // Step 1: Mark parents
    q.push(root);
    while (!q.empty()) {
        TreeNode* node = q.front(); q.pop();
        if (node->val == start) source = node;
        if (node->left) parent[node->left] = node, q.push(node->left);
        if (node->right) parent[node->right] = node, q.push(node->right);
    }

    // Step 2: BFS from source
    unordered_map<TreeNode*, bool> visited;
    q.push(source);
    visited[source] = true;
    int time = -1;

    while (!q.empty()) {
```

```

    int n = q.size();
    while (n-- > 0) {
        TreeNode* curr = q.front(); q.pop();
        for (TreeNode* neighbor : {curr->left, curr->right, parent[curr->parent]}) {
            if (neighbor && !visited[neighbor]) {
                visited[neighbor] = true;
                q.push(neighbor);
            }
        }
    }
    time++;
}
return time;
}

```

## 2. Distance K from Target Node

### Problem

Find all nodes at distance K from a target node.

### Approach

1. Build Parent Map using BFS
2. Do a Level Order BFS starting from target
3. Stop when level == K

### Code

```

vector<int> distanceK(TreeNode* root, TreeNode* target, int k) {
    unordered_map<TreeNode*, TreeNode*> parent;
    queue<TreeNode*> q;
    q.push(root);

    // Step 1: Mark Parents
    while (!q.empty()) {
        TreeNode* node = q.front(); q.pop();
        if (node->left) parent[node->left] = node, q.push(node->left);
        if (node->right) parent[node->right] = node, q.push(node->right);
    }
}

```

```

// Step 2: BFS from target
unordered_map<TreeNode*, bool> visited;
q.push(target);
visited[target] = true;
int dist = 0;

while (!q.empty()) {
    if (dist == k) break;
    int n = q.size();
    while (n-- > 0) {
        TreeNode* curr = q.front(); q.pop();
        for (TreeNode* neighbor : {curr->left, curr->right, parent[curr->left], parent[curr->right]}) {
            if (neighbor && !visited[neighbor]) {
                visited[neighbor] = true;
                q.push(neighbor);
            }
        }
    }
    dist++;
}

vector<int> res;
while (!q.empty()) res.push_back(q.front()->val), q.pop();
return res;
}

```

## 3. Construct Binary Tree from Traversals

### A. Preorder + Inorder → Binary Tree

#### Approach

1. The first element in preorder is root
2. Find this root in inorder → elements to left = left subtree, right = right subtree
3. Recursively build left and right

#### Code

```

TreeNode* build(vector<int>& preorder, int preStart, int inStart, int inEnd) {
    if (inStart > inEnd) return nullptr;

```

```

TreeNode* root = new TreeNode(preorder[preStart]);
int inIndex = inMap[root->val];
int leftSize = inIndex - inStart;

root->left = build(preorder, preStart + 1, inStart, inIndex - 1, inMap);
root->right = build(preorder, preStart + 1 + leftSize, inIndex + 1, inMap);
return root;
}

TreeNode* buildTree(vector<int>& preorder, vector<int>& inorder) {
    unordered_map<int, int> inMap;
    for (int i = 0; i < inorder.size(); i++) inMap[inorder[i]] = i;
    return build(preorder, 0, 0, inorder.size() - 1, inMap);
}

```

## B. Postorder + Inorder → Binary Tree

### Approach

1. The last element in postorder is root
2. Similar partitioning using inorder
3. Recursively build right subtree first (postorder: Left → Right → Root)

### Code

```

TreeNode* build(vector<int>& postorder, int postEnd, int inStart, int inEnd, unordered_map<int, int> inMap) {
    if (inStart > inEnd) return nullptr;

    TreeNode* root = new TreeNode(postorder[postEnd]);
    int inIndex = inMap[root->val];
    int rightSize = inEnd - inIndex;

    root->right = build(postorder, postEnd - 1, inIndex + 1, inEnd, inMap);
    root->left = build(postorder, postEnd - rightSize - 1, inStart, inIndex, inMap);
    return root;
}

TreeNode* buildTree(vector<int>& inorder, vector<int>& postorder) {
    unordered_map<int, int> inMap;
    for (int i = 0; i < inorder.size(); i++) inMap[inorder[i]] = i;
    return build(postorder, postorder.size() - 1, 0, inorder.size() - 1, inMap);
}

```

## 4. Serialize and Deserialize a Binary Tree (BFS)

### Serialize: Level Order with Nulls

```
string serialize(TreeNode* root) {
    if (!root) return "#";
    queue<TreeNode*> q;
    q.push(root);
    string result;

    while (!q.empty()) {
        TreeNode* node = q.front(); q.pop();
        if (!node) result += "# ", continue;
        result += to_string(node->val) + " ";
        q.push(node->left);
        q.push(node->right);
    }
    return result;
}
```

## 5. Morris Traversal (Inorder and Preorder)

### Inorder (Left → Root → Right)

#### Code

```
vector<int> morrisInorder(TreeNode* root) {
    vector<int> res;
    TreeNode* curr = root;

    while (curr) {
        if (!curr->left) {
            res.push_back(curr->val);
            curr = curr->right;
        } else {
            TreeNode* pred = curr->left;
            while (pred->right && pred->right != curr) pred = pred->right;

            if (!pred->right) {
                pred->right = curr;
            }
        }
    }
    return res;
}
```

```

        curr = curr->left;
    } else {
        pred->right = nullptr;
        res.push_back(curr->val);
        curr = curr->right;
    }
}
}
return res;
}

```

## Preorder (Root → Left → Right)

### Code

```

vector<int> morrisPreorder(TreeNode* root) {
    vector<int> res;
    TreeNode* curr = root;

    while (curr) {
        if (!curr->left) {
            res.push_back(curr->val);
            curr = curr->right;
        } else {
            TreeNode* pred = curr->left;
            while (pred->right && pred->right != curr) pred = pred->right

            if (!pred->right) {
                res.push_back(curr->val); // Visit before threading
                pred->right = curr;
                curr = curr->left;
            } else {
                pred->right = nullptr;
                curr = curr->right;
            }
        }
    }
    return res;
}

```

## Key Notes

- **Diameter  $\neq$  Burn Time always**
  - Burn time is from a specific node outward (like BFS)
  - Diameter is longest path in tree (from any node to any node)
- **Morris Traversal** is threading based and uses  $O(1)$  space