

Referential Transparency (same input -> same output) means we can memoize!!

```
Algorithm isValidAvI(T)
        for each p in T.positions() do
                 setHeight(p, -1) // initialize the attribute in the tree to special value
        height(T, T.root())
         res <- isValidAvlHelper(T, T.root())
1
        return res
Algorithm isValidAvlHelper(T, p)
n
        if T.isExternal(p) then
                 return true
n
        IRes <- isValidAvlHelper(T, T.leftChild(p))</pre>
n
        rRes <- isValidAvlHelper(T, T.rightChild(p))
n
        lh <- height(T, T.leftChild(p))</pre>
n
        rh <- height(T, T.rightChild(p))</pre>
n
         pRes <- (abs(lh-rh) <= 1)
n
        return IRes /\ rRes /\ pRes
n
Algorithm height(T, p) // memoized version of height (simplifies code of isValidAvlHelper)
        if T.isExternal(p)) then
                 return 0
         h <- getHeight(p)
         if h < 0 then // has the height already been computed for internal node
                 lh <- height(T, T.leftChild(p))</pre>
                 rh <- height(T, T.rightChild(p))</pre>
                 h \leftarrow max(lh, rh) + 1
                 setHeight(p, h)
         return h
```

abc

– а, b, с

ab, ac, bc

abc

 2^3

cababc

abdbc

longest common subsequence is based on aligning the two strings so the maximum number of characters match

cababc.

 $\underline{a}\underline{b}d\underline{c}\underline{b}$ LCS length = 3

caba.bc

 $.\underline{a}\underline{b}\underline{d}\underline{c}\underline{b}.\underline{LCS}\underline{length} = 3$

cababc

 $.\underline{a}\underline{b}dcb\underline{LCS}length = 2$

$$L_{i,j} = MAX \{ L_{i-1, j-1} + S(a_i,b_j), L_{i, j-1} + 0, L_{i-1, j} + 0 \}$$

Longest Common Sequence

		0	1	2	3	4	5	6	7	8	9	10	11
			G	Α	Α	Т	Т	С	Α	G	Т	Т	Α
0		0	0	0	0	0	0	0	0	0	0	0	0
1	G	م.	1	1	1	1	1	1	1	1	1	1	1
2	G	0	چر	_1	1	1	1	1	1	2	2	2	2
3	Α	0	1	ν,	2~	2	2	2	2	2	2	2	3
4	Т	0	1	2	2	3.	3	3	3	3	3	3	3
5	С	0	1	2	2	3	3	4.	4_	4	4	4	4
6	G	0	1	2	2	3	3	4	4	5.	_5_	-5 -	5
7	Α	0	1	2	3	3	3	4	5	5	5	5	6

ATTGACTTAA.G

A . . G . C . T . A G G

 $\mathsf{G} \; . \; \mathsf{A} \; \mathsf{A} \; \mathsf{T} \; \mathsf{T} \; \mathsf{C} \; \mathsf{A} \; \mathsf{G} \; \mathsf{T} \; \mathsf{T} \; \mathsf{A}$

 $\mathsf{G}\;\mathsf{G}\;\mathsf{A}\;.\;\mathsf{T}\;.\;\mathsf{C}\;.\;\mathsf{G}\;.\;.\;\mathsf{A}$

```
L_{i,j} = MAX \{ L_{i-1, i-1} + S(a_i,b_i), L_{i,i-1, L_{i-1, j}} \}
Non-memoized version:
Algorithm LCS(S1, S2):
  Input: Strings S1 and S2
  Output: Length of the LCS of S1 and S2
        m <- S1.size()
        n <- S2.size()
        return LCShelper(S1, S2, m, n)
Algorithm LCShelper(S1, S2, m, n):
  Input: Strings S1 and S2 with at least m and n elements, respectively
  Output: Length of the LCS of S1[1..m] and S2[1..n]
        if n = 0 then
            return 0
        else if m = 0 then
            return 0
        else if S1[m] = S2[n] then
            return LCShelper (S1, S2, m -1, n -1) + 1
        else
            return MAX ( LCShelper (S1, S2, m, n -1), LCShelper (S1, S2, m -1, n) )
Memoized version:
Algorithm LCS(S1, S2):
  Input: Strings 51 and 52
  Output: Length of the LCS of S1 and S2
        m <- S1.size()
        n <- S2.size()
        L <- new array[m+1, n+1] // create and initialize a table
        for i <- 0 to m do
            for j <- 0 to n do
                L[i,j] < -1
        return LCShelper(S1, S2, m, n, L)
Algorithm LCShelper(S1, S2, m, n, L):
  Input: Strings S1 and S2 with at least m and n elements, respectively
  Output: Length of the LCS of S1[1..m] and S2[1..n]
   if L[m,n] < 0 then // has the LCS already been computed
        if n = 0 then
            L[m,n] < 0
        else if m = 0 then
            L[m,n] < 0
        else if S1[m] = S2[n] then
            L[m,n] \leftarrow LCShelper(S1, S2, m-1, n-1) + 1
        else
            L[m,n] \leftarrow MAX (LCShelper (S1, S2, m, n-1), LCShelper (S1, S2, m-1, n))
    return L[m,n]
```