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## ISYE 4133 Assignment 5

I worked with Nicholas Loprinzo on this homework.

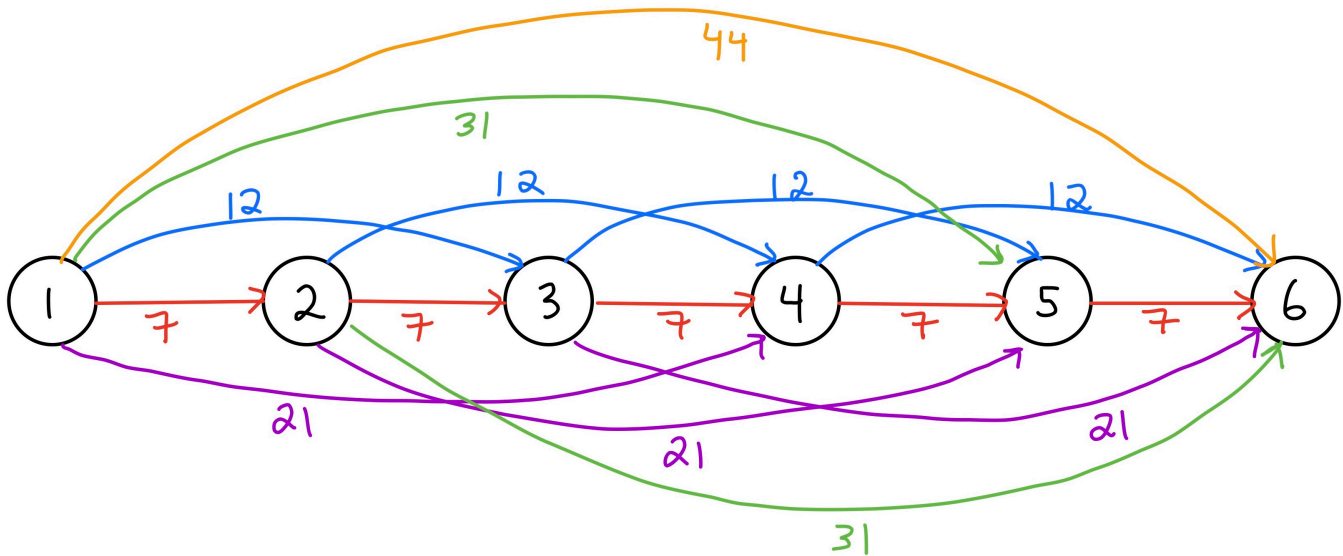
### 1

I used the following reference for the idea to model this problem as a shortest path problem:

<https://www.lix.polytechnique.fr/~liberti/teaching/isic/isc612-06/exercices.pdf>

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The possibilities for this renewal plan can be modeled by the following graph:



The nodes 1 to 5 represent the start of each year. The 6th node represents the end of the 5th year or end of the planning period. For each node  $i < 6$  and  $i < j$ , arc  $(i, j)$  represents the event that machinery bought at the beginning of the  $i$ -th year is sold at the beginning of the  $j$ -th year. The cost (in thousands of dollars)  $c_{ij}$  along each arc  $(i, j)$  is given by:

$$c_{ij} = 12 + \left( \sum_{k=i}^{j-1} m_k \right) - g_j$$

where  $m_k$  is the maintenance cost in the  $k$ -th year and  $g_j$  is the gain from selling the machinery at the beginning of the  $j$ -th year.

A renewal plan which minimizes the total operation cost is the shortest path from node 1 to 6 because the cost of the path is the cost of the plan.

Before formulating a linear integer program for this problem, we can use Dijkstra's algorithm to find the best renewal plan. We have:

$$U = \{1, 2, 3, 4, 5, 6\}$$

Node	$P_v$	Predecessor of $U$
1	0	$\emptyset$
2	$\infty$	UD
3	$\infty$	UD
4	$\infty$	UD
5	$\infty$	UD
6	$\infty$	UD

$$U = \{2, 3, 4, 5, 6\}$$

Node	$P_v$	Predecessor of $U$
1	0	$\emptyset$
2	7	1
3	12	1
4	21	1
5	31	1
6	44	1

$$U = \{3, 4, 5, 6\}$$

Node	$P_v$	Predecessor of $U$
1	0	$\emptyset$
2	7	1
3	12	1
4	19	2
5	28	2
6	38	2

$$U = \{4, 5, 6\}$$

Node	$P_v$	Predecessor of $U$
1	0	$\emptyset$
2	7	1
3	12	1
4	19	2
5	24	3
6	33	3

$$U = \{5, 6\}$$

Node	$P_v$	Predecessor of $U$
1	0	$\emptyset$
2	7	1
3	12	1
4	19	2
5	24	3
6	31	4

$$U = \{6\}$$

Node	$P_v$	Predecessor of $U$
------	-------	--------------------

1	0	$\emptyset$
2	7	1
3	12	1
4	19	2
5	24	3
6	31	4

The shortest path has cost \$31,000:

$$1 \rightarrow 2 \rightarrow 4 \rightarrow 6$$

OR

$$1 \rightarrow 3 \rightarrow 4 \rightarrow 6$$

OR

$$1 \rightarrow 3 \rightarrow 5 \rightarrow 6$$

The following is the integer program:

Nodes  $V = \{1, 2, 3, 4, 5, 6\}$

Arcs  $A = \{(1, 2), (1, 3), (1, 4), (1, 5), (1, 6), (2, 3), (2, 4), (2, 5), (2, 6), (3, 4), (3, 5), (3, 6), (4, 5), (4, 6), (5, 6)\}$

$$x_{ij} = \begin{cases} 1 & \text{if } i \rightarrow j \text{ is a path } p \\ 0 & \text{otherwise} \end{cases}$$

$$\min \sum_{(i,j) \in A} c_{ij} x_{ij}$$

$$\sum_{i:(1,i) \in A} x_{1,i} = 1$$

$$\sum_{i:(i,6) \in A} x_{i,6} = 1$$

$$\forall j \in V \setminus \{1, 6\} \quad \sum_{i:(i,j) \in A} x_{ij} - \sum_{i:(j,i) \in A} x_{ji} = 0$$

$$x_{ij} \in \{0, 1\} \forall i, j$$

The following is the solution in Gurobi:

```
In [2]: V = set([1,2,3,4,5,6])
c = np.array([7,12,21,31,44,7,12,21,31,7,12,21,7,12,7])
m = Model()
X = m.addVars([i for i in permutations(V,2) if i[1]>i[0]],vtype = GRB.
    BINARY, name = 'X')
m.addConstr(sum(X[1,i] for i in V if i!=1) == 1)
m.addConstr(sum(X[i,6] for i in V if i!=6) == 1)

j = V-{1,6}

for v in j:
    m.addConstr(sum(X[i,v] for i in V if i<v) - sum(X[v,i] for i in V
    if i>v) == 0)

m.setObjective(np.dot(c,X.values()), GRB.MINIMIZE)

m.optimize()

printOptimal(m)
```

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Optimize a model with 6 rows, 15 columns and 30 nonzeros

Variable types: 0 continuous, 15 integer (15 binary)

Coefficient statistics:

Matrix range [1e+00, 1e+00]

Objective range [7e+00, 4e+01]

Bounds range [1e+00, 1e+00]

RHS range [1e+00, 1e+00]

Found heuristic solution: objective 44.0000000

Presolve removed 6 rows and 15 columns

Presolve time: 0.00s

Presolve: All rows and columns removed

Explored 0 nodes (0 simplex iterations) in 0.01 seconds

Thread count was 1 (of 8 available processors)

Solution count 2: 31 44

Optimal solution found (tolerance 1.00e-04)

Best objective 3.1000000000000e+01, best bound 3.1000000000000e+01, gap 0.0000%

Variable Values:

X[1,2] 0.0

X[1,3] 1.0

X[1,4] 0.0

X[1,5] 0.0

X[1,6] 0.0

X[2,3] 0.0

X[2,4] 0.0

X[2,5] 0.0

X[2,6] 0.0

X[3,4] 0.0

X[3,5] 1.0

X[3,6] 0.0

X[4,5] 0.0

X[4,6] 0.0

X[5,6] 1.0

Objective Value: 31.0

The following is the linear program relaxation:

$$\begin{aligned}
 & \min \sum_{(i,j) \in A} c_{ij} x_{ij} \\
 & \sum_{i:(1,i) \in A} x_{1,i} = 1 \\
 & \sum_{i:(i,6) \in A} x_{i,6} = 1 \\
 & \forall j \in V \setminus \{1,6\} \quad \sum_{i:(i,j) \in A} x_{ij} - \sum_{i:(j,i) \in A} x_{ji} = 0 \\
 & x_{ij} \geq 0 \quad \forall i, j
 \end{aligned}$$

The following is the solution in Gurobi:

```

In [3]: V = set([1,2,3,4,5,6])
c = np.array([7,12,21,31,44,7,12,21,31,7,12,21,7,12,7])
m = Model()
X = m.addVars([i for i in permutations(V,2) if i[1]>i[0]], name = 'X')
m.addConstr(sum(X[1,i] for i in V if i!=1) == 1)
m.addConstr(sum(X[i,6] for i in V if i!=6) == 1)

j = V-{1,6}

for v in j:
    m.addConstr(sum(X[i,v] for i in V if i<v) - sum(X[v,i] for i in V
if i>v) == 0)

m.setObjective(np.dot(c,X.values()), GRB.MINIMIZE)

m.optimize()

printOptimal(m)

```

Optimize a model with 6 rows, 15 columns and 30 nonzeros

Coefficient statistics:

Matrix range [1e+00, 1e+00]

Objective range [7e+00, 4e+01]

Bounds range [0e+00, 0e+00]

RHS range [1e+00, 1e+00]

Presolve removed 2 rows and 9 columns

Presolve time: 0.01s

Presolved: 4 rows, 6 columns, 12 nonzeros

Iteration	Objective	Primal Inf.	Dual Inf.	Time
0	1.9000000e+01	2.000000e+00	0.000000e+00	0s
2	3.1000000e+01	0.000000e+00	0.000000e+00	0s

Solved in 2 iterations and 0.01 seconds

Optimal objective 3.100000000e+01

Variable Values:

X[1,2] 1.0

X[1,3] 0.0

X[1,4] 0.0

X[1,5] 0.0

X[1,6] 0.0

X[2,3] 0.0

X[2,4] 1.0

X[2,5] 0.0

X[2,6] 0.0

X[3,4] 0.0

X[3,5] 0.0

X[3,6] 0.0

X[4,5] 0.0

X[4,6] 1.0

X[5,6] 0.0

Objective Value: 31.0

The linear program relaxation will always give the integer solution for free.



## 2

$x_j$ : amount of widgets produced in time period  $j$ .  $\forall j = 1, \dots, K$

$s_j$ : amount of widgets held in inventory over in time period  $j$ .  $\forall j = 1, \dots, K$

$$\min \sum_{j=1}^K x_j c_j + h_j s_j$$

$$x_j + s_{j-1} = d_j + s_j \quad \forall j$$

$$p_j \geq 0 \quad \forall j$$

$$s_j \geq 0 \quad \forall j$$

$$s_0 = 0$$

## 3 (a)

$x_j$ : if item  $j$  is included in the subset.  $\forall j = 1, \dots, n$

$$\min \sum_j c_j x_j$$

$$\sum_j a_j x_j \leq B$$

$$x_j \in \{0, 1\}$$

## 3 (b)

The following table outlines the penalties for the  $i$ -th unit of increase in  $B$ :

$i$	penalty $p_i$ for the $i$ -th unit
1	2
2	3
3	3
4	6
5	6
6	10
7	10
8	10

9	10
10	10

B can be increased by a certain amount,  $S = \sum_i a_i x_i - B$

$$y_i = \begin{cases} 1 & \text{if } S \text{ is } \geq i \\ 0 & \text{otherwise} \end{cases} \quad \forall i = 1, \dots, 10$$

For example, if B increases by 6, then  $y_1 = y_2 = y_3 = y_4 = y_5 = y_6 = 1$  and then this will increase costs by  $1(2) + 1(3) + 1(3) + 1(6) + 1(6) + 1(10) = 30$

The cost function is altered to include this added penalty cost ( $\sum p_i y_i$ ).

The following logic can be used to create the constraints:

If  $S - 1 \geq 0$  then  $y_1 = 1$

If  $S - 2 \geq 0$  then  $y_2 = 1$

If  $S - 3 \geq 0$  then  $y_3 = 1$

And so on...

The constraints become:

$$-(y_1 - 1) \leq Mz_1, \quad S - 1 \leq M(1 - z_1)$$

$$-(y_2 - 1) \leq Mz_2, \quad S - 2 \leq M(1 - z_2)$$

$$-(y_3 - 1) \leq Mz_3, \quad S - 3 \leq M(1 - z_3)$$

The following is the integer linear program:

$$\begin{aligned} \min \quad & \sum_j c_j x_j + \sum_i p_i y_i \\ & \sum_j a_j x_j - B \leq S \\ & -(y_i - 1) \leq Mz_i \quad \forall i = 1, \dots, 10 \\ & S - i \leq M(1 - z_i) \quad \forall i = 1, \dots, 10 \\ & x_j \in \{0, 1\} \\ & y_i \in \{0, 1\} \\ & S \geq 0 \\ & z_i \in \{0, 1\} \end{aligned}$$

### 3 (c)

For this change, the penalty table might look like this:

$i$	penalty $p_i$ for the $i$ -th unit
1	2
2	1
3	.9
4	.8
5	.7
6	.6
7	.5
8	.4
9	.3
10	.2

Because the only thing that changes in this problem is that the  $p_i$  values are different, the whole formulation remains the same.

Based on the above example non-increasing penalty table, if B increases by 6, then

$y_1 = y_2 = y_3 = y_4 = y_5 = y_6 = 1$  and then this will increase costs by  
 $1(2) + 1(1) + 1(.9) + 1(.8) + 1(.7) + 1(.6) = 6$

## 4 (a)

$w_j$ : number of cores per program

$q_j$ : number of programs desired

$x_{ij}$ : number of programs with  $w_j$  number of cores

Let  $F = \{P = (a_{P,1}, \dots, a_{P,n}) : \sum_{i=1}^n w_j a_{P,j} \leq W, a_{P,j} \in Z_+ \forall 1 \leq j \leq n\}$  be the set of all valid patterns.

Then solve:

$$\begin{aligned} \min \quad & \sum_{P \in F} x_P \\ \text{s.t.} \quad & \sum_{P \in F} x_P a_{P,j} = b_j \quad \forall j = 1, \dots, n \\ & x_P \geq 0 \quad \forall P \in F \end{aligned}$$

```
In [5]: w = np.array([20,30,32,16,5])
        q = np.array([1000,1500,2000,2500,1200])
```

Get all possible patterns:

```
In [6]: L = min(w)
        R = 64
        n = len(w)

        m = Model()

        m.params.PoolSolutions = 400

        m.params.PoolSearchMode = 2

        y = m.addVars(n, vtype = GRB.INTEGER, name = 'Y')

        m.addConstr(np.dot(w, y.values()) <= R)

        m.setObjective(0, GRB.MINIMIZE)

        m.optimize()

        printOptimal(m)
```

Changed value of parameter PoolSolutions to 400  
Prev: 10 Min: 1 Max: 2000000000 Default: 10

Changed value of parameter PoolSearchMode to 2

Prev: 0 Min: 0 Max: 2 Default: 0

Optimize a model with 1 rows, 5 columns and 5 nonzeros

Variable types: 0 continuous, 5 integer (0 binary)

Coefficient statistics:

Matrix range [5e+00, 3e+01]

Objective range [0e+00, 0e+00]

Bounds range [0e+00, 0e+00]

RHS range [6e+01, 6e+01]

Found heuristic solution: objective 0.0000000

Presolve time: 0.00s

Presolved: 1 rows, 5 columns, 5 nonzeros

Variable types: 0 continuous, 5 integer (0 binary)

Root relaxation: objective 0.000000e+00, 0 iterations, 0.00 seconds

Nodes		Current Node			Objective Bounds				
Work	Expl	Unexpl	Obj	Depth	IntInf	Incumbent	BestBd	Gap	It/N
ode Time									
0s	0	0	-	0		0.00000	0.00000	0.00%	-
Optimal solution found at node 0 - now completing solution pool...									
0s	0	0	-	0		0.00000	0.00000	0.00%	-
0s	0	0	-	0		0.00000	0.00000	0.00%	-
0s	0	2	-	0		0.00000	0.00000	0.00%	-

Explored 187 nodes (0 simplex iterations) in 0.02 seconds

Thread count was 8 (of 8 available processors)

Solution count 94: 0 0 0 ... 0

No other solutions better than 1e+100

Optimal solution found (tolerance 1.00e-04)

Best objective 0.000000000000e+00, best bound 0.000000000000e+00, gap 0.0000%

Variable Values:

Y[0] -0.0

Y[1] -0.0

Y[2] -0.0

Y[3] -0.0

Y[4] -0.0

Objective Value: 0.0

```
In [7]: NumSolns = m.solCount

solutions = np.empty([n, NumSolns])

for x in range(NumSolns):
    m.params.outputFlag = 0
    m.params.SolutionNumber = x
    for i in range(n):
        solutions[i,x] = m.Xn[i]
```

There are 94 possible patterns.

Now solve the LP:

```
In [9]: m = Model()

x = m.addVars(NumSolns, vtype = GRB.CONTINUOUS, name = 'x')

for i in range(solutions.shape[0]):
    m.addConstr(np.dot(solutions[i,:],x.values()) >= q[i])

m.setObjective(sum(x.values()), GRB.MINIMIZE)

m.optimize()

printOptimal(m)
```

Optimize a model with 5 rows, 94 columns and 181 nonzeros

Coefficient statistics:

Matrix range [1e+00, 1e+01]

Objective range [1e+00, 1e+00]

Bounds range [0e+00, 0e+00]

RHS range [1e+03, 2e+03]

Presolve removed 0 rows and 72 columns

Presolve time: 0.00s

Presolved: 5 rows, 22 columns, 46 nonzeros

Iteration	Objective	Primal Inf.	Dual Inf.	Time
0	0.0000000e+00	2.775000e+03	0.000000e+00	0s
5	2.8083333e+03	0.000000e+00	0.000000e+00	0s

Solved in 5 iterations and 0.01 seconds

Optimal objective 2.808333333e+03

Objective Value: 2808.3333333333335

Now solved with Column Generation:

```
In [11]: def solveRestrictedLP(I):

    m = Model()

    m.params.outputFlag = 0

    x = m.addVars(len(I), vtype = GRB.CONTINUOUS, name = 'X')

    for i in range(solutions.shape[0]):
        m.addConstr(np.dot(solutions[i,I],x.values()) >= q[i])

    m.setObjective(sum(x.values()), GRB.MINIMIZE)

    m.optimize()

    y = m.Pi

    return(m.X, m.VBasis, m.objVal, y)
```

```
In [12]: def solveReducedCostLP(y):

    m = Model()

    m.params.outputFlag = 0

    a = m.addVars(solutions.shape[0], vtype = GRB.INTEGER, name = 'A')

    m.addConstr(np.dot(w,a.values()) <= 64)

    m.setObjective((1 - np.dot(a.values(),y)) , GRB.MINIMIZE)

    m.optimize()

    for i in range(solutions.shape[1]):
        if (list(solutions.T[i]) == list(np.floor(m.X))):
            newpattern = i

    return(m.X, m.objVal, newpattern)
```

```
In [13]: w = np.array([20,30,32,16,5])
q = np.array([1000,1500,2000,2500,1200])

I = []

for i in range(NumSolns):
    if sum(solutions[:,i]) == 1:
        I.append(i)
```



```
In [14]: obj = -10000
while obj < 0:
    X, VBasis, objVal, y = solveRestrictedLP(I)
    x, obj, newpattern = solveReducedCostLP(y)
    I.append(newpattern)
    print('-----')
    print(obj)
    print(objVal)
```

```
-----
-11.0
8200.0
-----
-3.0
7100.0
-----
-2.0
5225.0
-----
-1.0
4558.333333333333
-----
-1.0
3808.333333333335
-----
0.0
2808.333333333333
```

## 4 (b)

More constraints must be added to account for the RAM and bandwidth requirements.

```
In [15]: w = np.array([20,30,32,16,5])
q = np.array([1000,1500,2000,2500,1200])
ram = np.array([24,8,16,20,5])
bandwidth = np.array([0,2,1,.05,2.8])
```

```
In [16]: n = len(w)

m = Model()

m.params.PoolSolutions = 400

m.params.PoolSearchMode = 2

y = m.addVars(n, vtype = GRB.INTEGER, name = 'Y')

m.addConstr(np.dot(w, y.values()) <= 64)
m.addConstr(np.dot(ram, y.values()) <= 32)
m.addConstr(np.dot(bandwidth, y.values()) <= 4)

m.setObjective(0, GRB.MINIMIZE)

m.optimize()

printOptimal(m)
```

Changed value of parameter PoolSolutions to 400

Prev: 10 Min: 1 Max: 2000000000 Default: 10

Changed value of parameter PoolSearchMode to 2

Prev: 0 Min: 0 Max: 2 Default: 0

Optimize a model with 3 rows, 5 columns and 14 nonzeros

Variable types: 0 continuous, 5 integer (0 binary)

Coefficient statistics:

Matrix range [5e-02, 3e+01]

Objective range [0e+00, 0e+00]

Bounds range [0e+00, 0e+00]

RHS range [4e+00, 6e+01]

Found heuristic solution: objective 0.0000000

Presolve time: 0.00s

Presolved: 3 rows, 5 columns, 14 nonzeros

Variable types: 0 continuous, 5 integer (3 binary)

Root relaxation: objective 0.000000e+00, 0 iterations, 0.00 seconds

Nodes		Current Node			Objective Bounds				
Work	Expl	Unexpl	Obj	Depth	IntInf	Incumbent	BestBd	Gap	It/N
ode	Time								
0	0	-	0	0.00000	0.00000	0.00%	-		
0s									
Optimal solution found at node 0 - now completing solution pool...									
0	0	-	0	0.00000	0.00000	0.00%	-		
0s									
0	0	-	0	0.00000	0.00000	0.00%	-		

```

0s
      0      2      -      0      0.00000      0.00000      0.00%      -
0s

Explored 27 nodes (0 simplex iterations) in 0.01 seconds
Thread count was 8 (of 8 available processors)

Solution count 14: 0 0 0 ... 0
No other solutions better than 1e+100

Optimal solution found (tolerance 1.00e-04)
Best objective 0.000000000000e+00, best bound 0.000000000000e+00, gap
0.0000%
```

Variable Values:

```

Y[0] -0.0
Y[1] -0.0
Y[2] -0.0
Y[3] -0.0
Y[4] -0.0
```

Objective Value: 0.0

```

In [17]: NumSolns = m.solCount

solutions = np.empty([n, NumSolns])

for x in range(NumSolns):
    m.params.outputFlag = 0
    m.params.SolutionNumber = x
    for i in range(n):
        solutions[i,x] = m.Xn[i]
```

```

In [18]: m = Model()

x = m.addVars(NumSolns, vtype = GRB.CONTINUOUS, name = 'X')

for i in range(solutions.shape[0]):
    m.addConstr(np.dot(solutions[i,:],x.values()) >= q[i])

m.setObjective(sum(x.values()), GRB.MINIMIZE)

m.optimize()

printOptimal(m)
```

Optimize a model with 5 rows, 14 columns and 19 nonzeros

Coefficient statistics:

Matrix range [1e+00, 2e+00]

Objective range [1e+00, 1e+00]

Bounds range [0e+00, 0e+00]

RHS range [1e+03, 2e+03]

Presolve removed 0 rows and 3 columns

Presolve time: 0.00s

Presolved: 5 rows, 11 columns, 17 nonzeros

Iteration	Objective	Primal Inf.	Dual Inf.	Time
0	0.0000000e+00	6.450000e+03	0.000000e+00	0s
6	4.5000000e+03	0.000000e+00	0.000000e+00	0s

Solved in 6 iterations and 0.01 seconds

Optimal objective 4.500000000e+03

Variable Values:

X[0] 0.0

X[1] 0.0

X[2] 800.0

X[3] 0.0

X[4] 0.0

X[5] 1000.0

X[6] 1200.0

X[7] 0.0

X[8] 0.0

X[9] 500.0

X[10] 0.0

X[11] 0.0

X[12] 1000.0

X[13] 0.0

Objective Value: 4500.0

```
In [19]: def solveRestrictedLP(I):

    m = Model()

    m.params.outputFlag = 0

    x = m.addVars(len(I), vtype = GRB.CONTINUOUS, name = 'X')

    for i in range(solutions.shape[0]):
        m.addConstr(np.dot(solutions[i,I],x.values()) >= q[i])

    m.setObjective(sum(x.values()), GRB.MINIMIZE)

    m.optimize()

    y = m.Pi

    return(m.X, m.VBasis, m.objVal, y)
```

```
In [20]: def solveReducedCostLP(y):

    m = Model()

    m.params.outputFlag = 0

    a = m.addVars(solutions.shape[0], vtype = GRB.INTEGER, name = 'A')

    m.addConstr(np.dot(w,a.values()) <= 64)
    m.addConstr(np.dot(ram,a.values()) <= 32)
    m.addConstr(np.dot(bandwidth,a.values()) <= 4)

    m.setObjective((1 - np.dot(a.values(),y)) , GRB.MINIMIZE)

    m.optimize()

    for i in range(solutions.shape[1]):
        if (list(solutions.T[i]) == list(np.floor(m.X))):
            newpattern = i

    return(m.X, m.objVal, newpattern)
```

```
In [21]: w = np.array([20,30,32,16,5])
q = np.array([1000,1500,2000,2500,1200])

I = []

for i in range(NumSolns):
    if sum(solutions[:,i]) == 1:
        I.append(i)
```

```
In [22]: obj = -10000
while obj < 0:
    X, VBasis, objVal, y = solveRestrictedLP(I)
    x, obj, newpattern = solveReducedCostLP(y)
    I.append(newpattern)
    print('-----')
    print(obj)
    print(objVal)
```

```
-----
-1.0
8200.0
-----
-1.0
7200.0
-----
-1.0
6950.0
-----
-1.0
5750.0
-----
-0.5
5350.0
-----
-0.5
5100.0
-----
0.0
4500.0
```