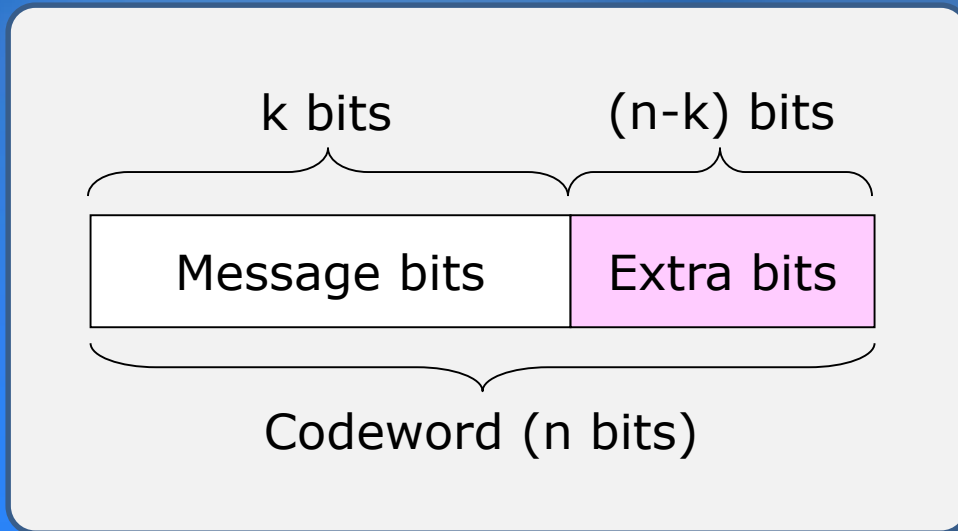


# Block Codes

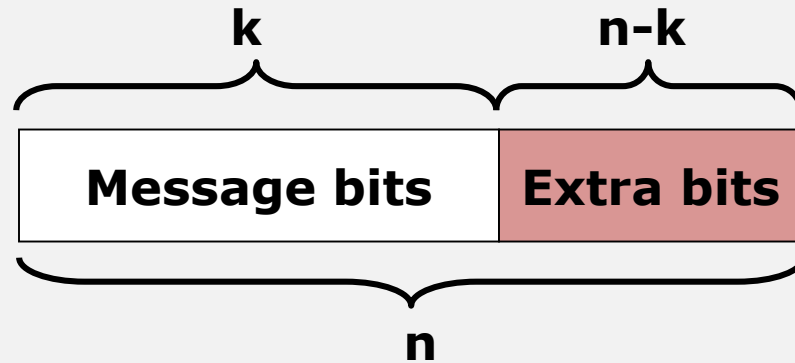
# $(n,k,d)$ Block Codes



- Split the message into  $k$ -bit blocks
- Create a codeword by adding  $(n-k)$  extra bits to each block.
  - The extra bits are computed based on the message bits.
  - Thus, they contain no new information.
- $d$  = minimum Hamming Distance between codewords
- Sometimes we drop the  $d$  and indicate only  $(n,k)$

# Code Rate

- **Code rate:** the fraction of sent bits that contain useful information (i.e. the message).
- For the  $(n,k,d)$  block code

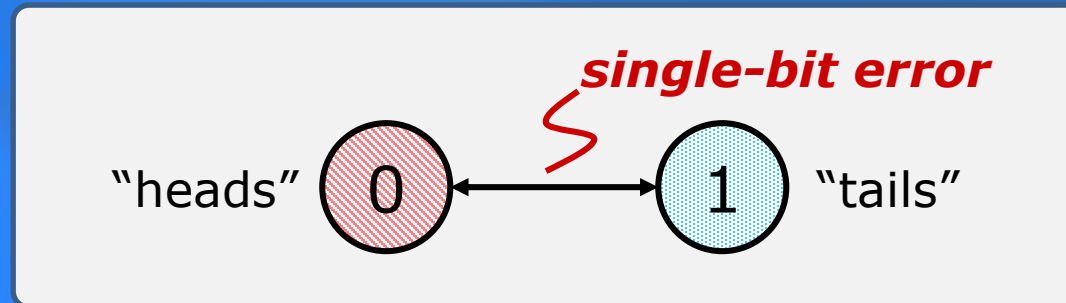


$$\text{code rate} = \frac{k}{n}$$

- Related terms
  - **Gross bit rate:** rate that all bits are sent  $= \frac{F_s}{\text{SPB}}$ 
    - Also called the data signaling rate
  - **Net bit rate:** rate that useful bits are sent  $= \text{code rate} \times \text{gross bit rate}$

# Hamming Distance

- The **Hamming Distance** between two codewords is the number of bit positions where the corresponding bits are different.
- For example
  - The Hamming distance between (00) and (10) is 1.
  - The Hamming distance between (00000000) and (11110011) is 6.
- The Hamming distance measures the number of bit errors it takes transform one codeword to another.
  - For example, if we use no coding, each bit is represented by one of two code words ("0" and "1").
  - Since the Hamming distance is 1, a single-bit error changes one code word the other.



# Error Detection vs Correction

- **Error detection**
  - We can detect errors
  - But, we don't know how to fix them
- **Error correction**
  - We can detect errors
  - And, we can correct them
- **For a given code, the receiver can choose whether to use the code to detect errors or to correct them.**

# The Error Detection/Correction Capability

- The minimum Hamming distance determines the maximum number of bit errors the receiver can detect or correct.
- If the minimum Hamming distance is  $d$ , the receiver can either
  - Detect but not correct errors in at most  $d-1$  bits of each codeword
  - OR
  - Detect and correct errors in at most  $(d-1)/2$  bits of each code word
- For example, if  $d = 3$ , the receiver can either
  - Detect 1 or 2 bit errors in each codeword.
  - Detect and correct 1 bit errors in each codeword.
    - If a 2 bit error does occur, it will be detected, but incorrectly corrected.