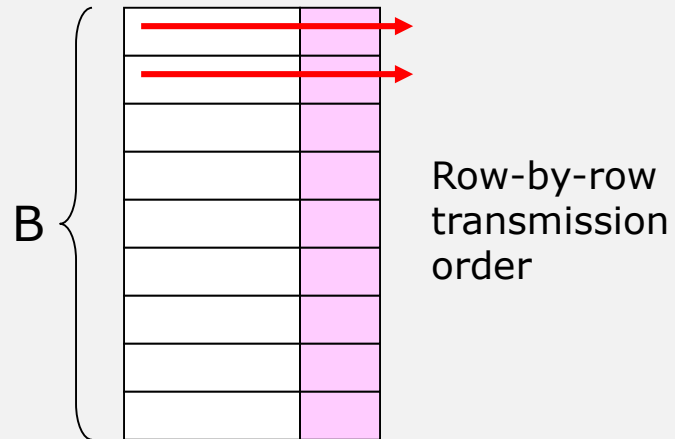


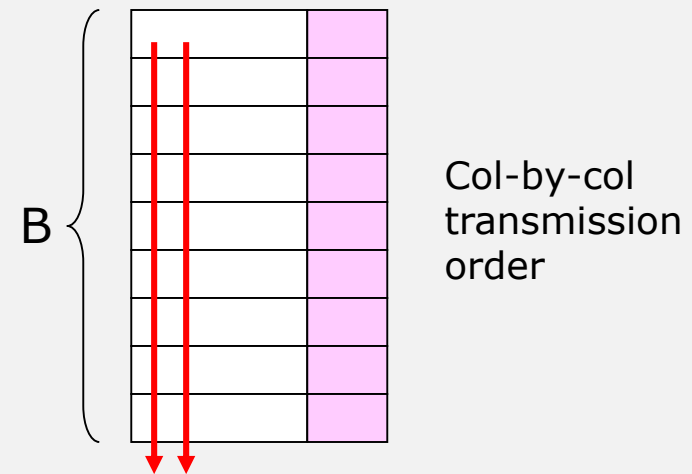
Burst Error Correction

Burst Errors

Correcting errors in a single bit or a few bits is nice, but in many situations errors come in bursts many bits long (e.g. damage to storage media, burst of interference on wireless channel, etc.).



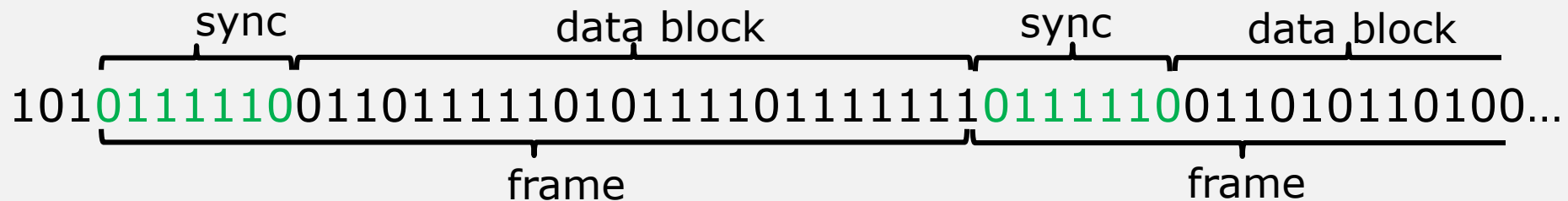
Problem: Bits from a particular codeword are transmitted sequentially, so a B-bit burst produces multi-bit errors.



Solution: **interleave bits** from B different codewords. Now a B-bit burst produces 1-bit errors in B different codewords.

Framing

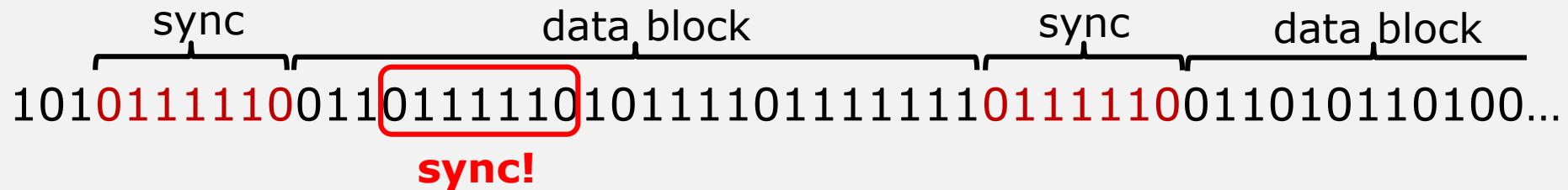
- Looking at a received bit stream, **how do we know where a block of interleaved codewords begins?**
- Possible solutions
 - Physical indication, e.g.
 - beginning of disk sector
 - separate control channel
 - **Frame sync sequence**
 - a unique bit pattern placed to mark start of a block (like start bit)
 - example sync sequence: **0111110** (5 1's framed by 0's)



- However, there is something that we need to take care of before adding the sync sequence.

Bit Stuffing: Motivation

- Bit pattern used for the sync sequence cannot appear elsewhere in frame, otherwise our search for the start of the frame will get confused.



- To deal with this, we preprocess the data block by bit stuffing, to make sure the sync sequence does not appear.

Bit Stuffing

- If the sync sequence is [0111110], one way to guarantee the sync sequence never appears in the data block is to make sure there are never five 1's in a row.
- Bit stuffing
 - Scan through the data block
 - Whenever the sender finds *four* consecutive 1's, it stuffs (adds) a 0 bit into the outgoing stream.
 - The 0 is always stuffed whether or not the next bit is 1.
- Unstuffing
 - When the receiver sees *four* consecutive 1's, it unstuffs (removes) the next bit (which will be a 0).

Bit Stuffing Example

Input Stream



01101111110011101111111100000

Stuffed Stream



01101111**0**11001111**0**01111**0**1111**0**00000

Stuffed bits

Unstuffed Stream



01101111110011101111111100000

Example: Channel Coding

011011101101

Take an input message

↓
0110
1110
1101

Break the message into k-bit blocks ($k = 4$)

↓
01101111
11100101
11010110

Add $(n-k)$ parity bits to form n-bit codewords ($n = 8$)

↓
011111110001100111101110
↙ ↓ ↓ ↓ ↘
01111011100011001111001110

Interleave bits from B codewords ($B = 3$)

Bit-stuff

↓
011111001111011100011001111001110

Add sync pattern ([0111110])

Example: Error Correction

0111110011110111100011001111001110

bit errors →

011111001111011100100001111001110

01111011100100001111001110

011111110010000111101110

01100111

11110101

11000110

0 1 0

1 0 1

1 1

1 1 0

1 1 1

0 1

1 1 0

0 0 1

1 0

0110

1110

1101

Sent bit stream

Received bit stream (with errors)

Search for and remove sync pattern

Destuff the frame

**De-interleave to form B n-bit codewords
(B=3, n=8)**

Perform error correction

Extract the k=4 corrected message bits