

Relational Design Theory

Motivation & overview

Designing a database schema

- Usually many designs possible
- Some are (much) better than others!
- How do we choose?

Often use higher-level design tools, but ...

- Some designers go straight to relations
- Useful to understand why tools produce certain schemas
- Very nice theory for relational database design

Example: College application info.

- SSN and name
- Colleges applying to
- High schools attended (with city)
- Hobbies

Apply(SSN, sName, cName, HS, HScity, hobby)

123 Ann from PAHS (P.A.) and GHS (P.A.) plays tennis and trumpet and applied to Stanford, Berkeley, and MIT

```
123 Ann Stanford PAHS P.A. tennis
123 Ann Berkeley PAHS P.A. tennis
123 Ann Berkeley PAHS P.A. trumpet
:: CHS:
```

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Design "anomalies"

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Design "anomalies"

- Redundancy
- Update anomaly

update facts trumpet cornet

123 Ann from PAHS (P.A.) and GHS (P.A.) plays tennis and trumpet and applied to Stanford, Berkeley, and MIT

Design "anomalies"

- Redundancy
- Update anomaly
- Deletion anomaly



- **Example: College application info.**
 - SSN and name
 - Colleges applying to
 - High schools attended (with city)
 - Hobbies

```
Student(SSN, sName)
Apply(SSN, cName) hobby)
HighSchool (SSN, HS), Houly )
```

Design by decomposition

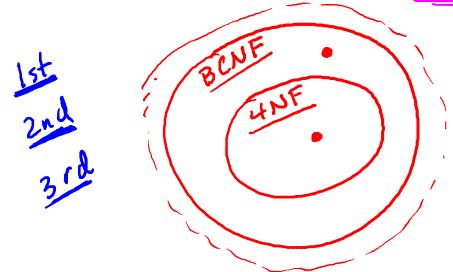
- Start with "mega" relations containing everything
- Decompose into smaller, better relations with same info.
- Can do <u>decomposition</u> automatically

Automatic decomposition

- "Mega" relations + properties of the data
- System decomposes based on properties
- Final set of relations satisfies normal form
 - No anomalies, no lost information

Properties and Normal Forms

- \checkmark Functional dependencies \Rightarrow Boyce-Codd Normal Form
- \checkmark + Multivalued dependences \Rightarrow Fourth Normal Form



Functional Dependencies and BCNF

Apply(SSN, sName, cName)

Redundancy; Update & Deletion Anomalies

- Storing SSN-sName pair once for each college <</p>

Functional Dependency (SN) -> sName

- Same SSN always has same sName
- Should store each SSN's sName only once

Boyce-Codd Normal Form If $A \rightarrow B$ then A is a key

Decompose: Student(SSN, SName) Apply(SSN, cName)

Multivalued Dependencies and 4NF

Apply(SSN, cName, HS)

- Redundancy; Update & Deletion Anomalies
- Multiplicative effect C colleges, H high schools tuples
- Not addressed by BCNF: No functional dependencies

Multivalued Dependency SSN ->> CName SSN ->> HS

- Given SSN has every combination of cName with HS
- Should store each cName and each HS for an SSN once

Fourth Normal Form If A >>> B) then (A is a key

Decompose: Apply(SSN, cName) HighSchool(SSN, HS)

Design by decomposition

- "Mega" relations + properties of the data
- System decomposes based on properties
- Final set of relations satisfies normal form
 - No anomalies, no lost information
- Functional dependencies ⇒ Boyce-Codd Normal Form
- Multivalued dependences ⇒ Fourth Normal Form