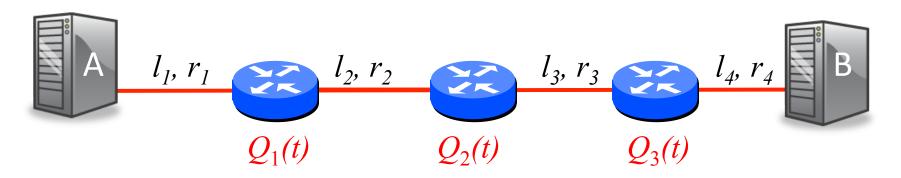
# CS144 An Introduction to Computer Networks

#### **Packet Switching**

Guaranteed Delay



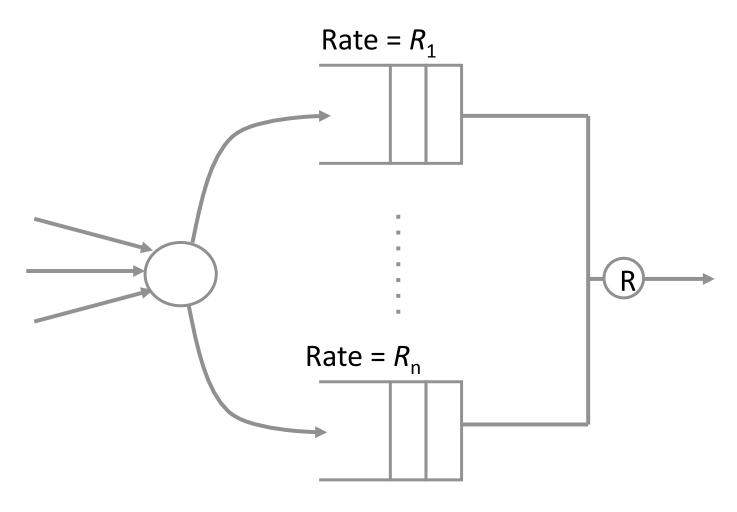
### Delay guarantees: Intuition



End-to-end delay, 
$$\tau = \sum_{i} \left( \frac{p}{r_i} + \frac{l_i}{c} + Q_i(t) \right)$$

If we know the upper bound of  $Q_1(t)$ ,  $Q_2(t)$  and  $Q_3(t)$ , then we know the upper bound of the end-to-end delay.

### Delay guarantees: Intuition



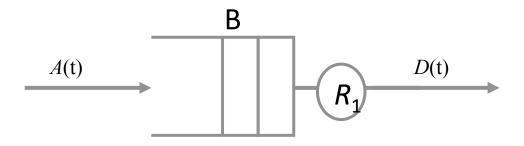
# So how can we control the delay of packets?

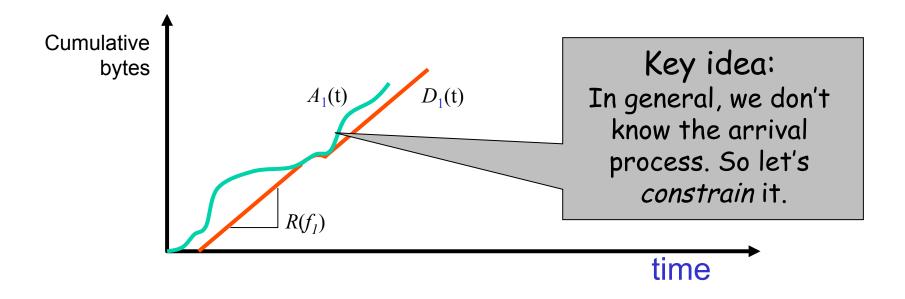
### What we already know how to control:

- 1. The <u>rate</u> at which a queue is served (WFQ).
- 2. The <u>size</u> of each queue.

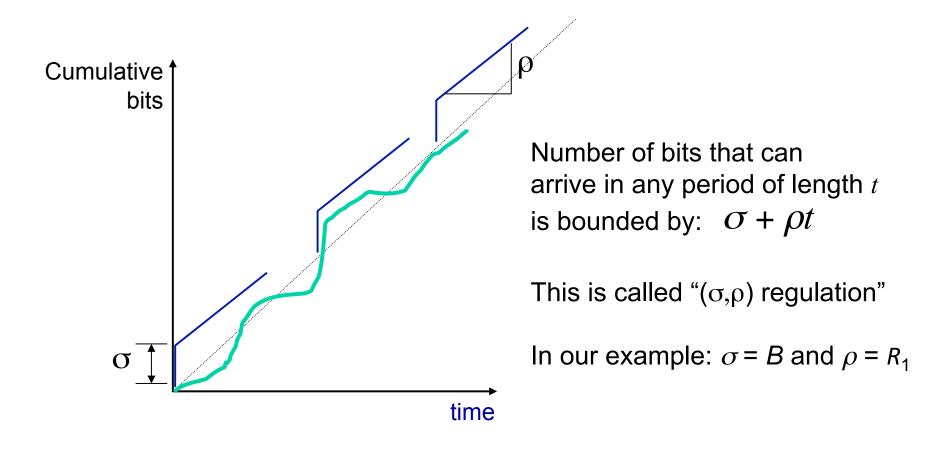
How do we make sure no packets are dropped?

### Zooming in on one queue

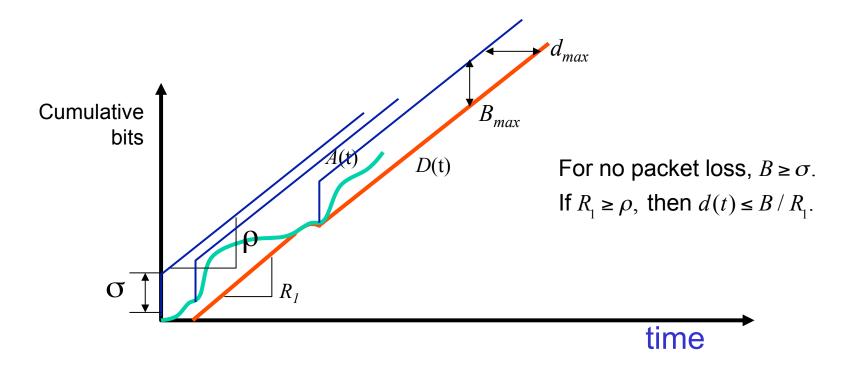




## Constraining traffic

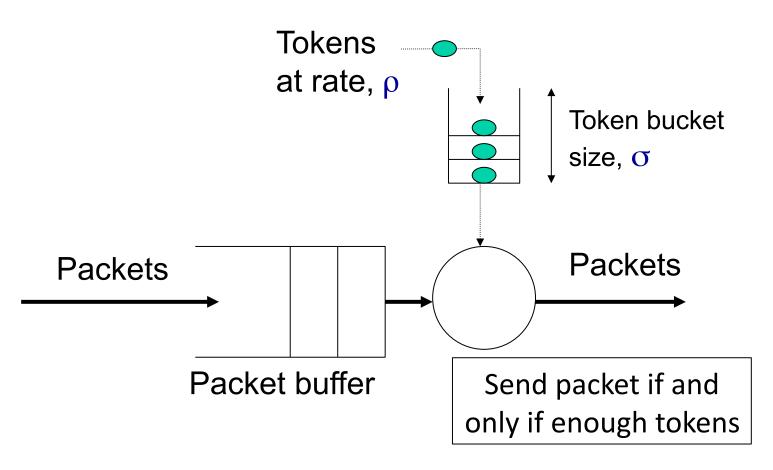


# (σ,ρ)-constrained Arrivals and Minimum Service Rate

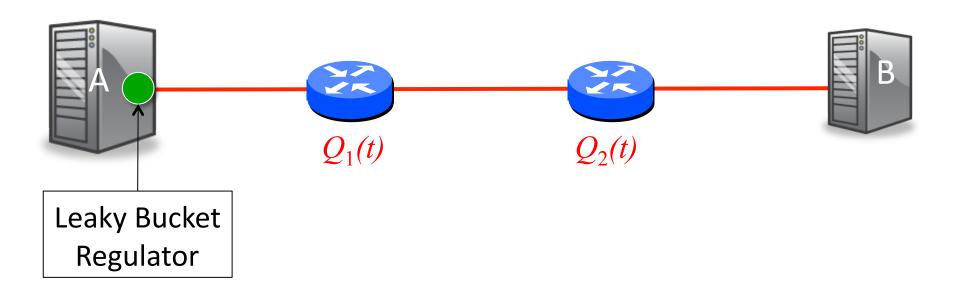


If flows are leaky-bucket constrained, and routers use WFQ, then end-to-end delay guarantees are possible.

### The leaky bucket regulator

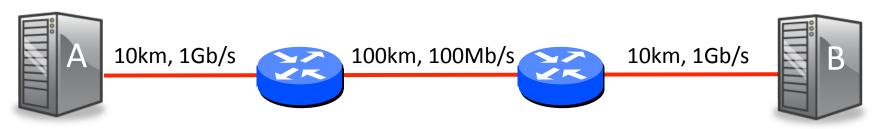


### Putting it all together



### An example

In the network below, an application wants a rate of 10Mb/s and an end to end delay of less than 5ms for 1000byte packets.



### In practice

While it is technically possible to do so, very few networks actually control end to end delay.

#### Why?

- It is complicated to make work, requiring coordination.
- In most networks, a combination of over-provisioning and priorities work well enough.

### Summary

If we know the size of a queue and the rate at which it is served, then we can bound the delay through it.

We can pick the size of the queue, and WFQ lets us pick the rate at which it is served.

Therefore, we just need a way to prevent packets being dropped along the way. For this, we use a leaky bucket regulator.

We can therefore bound the end to end delay.