## 3. Longest Substring Without Repeating Characters

This is a sliding window technique problem. The technique involves maintaining an optimization in a fixed or variable size window which is continuous in nature. (Subarray or substring).

In the problem, given a string we need to find the largest substring of unique character. So for example: In abcdaa, abcd is the largest substring of unique characters whose length is 4.

To solve this problem, first let us find the brute force way to solve it, we run a nested loop inside another loop, and form substring length until we get a character equal to the outer loop character. this is a  $O(n^2)$  solution.

But using sliding window we eliminate the need for two loops. This can be solved in n time complexity. The solution is as follows:

- Maintain left and right pointer.
- Keep the last occurred index of a character in an unordered map. (Hashmap key value pair).
- 3. We run a for loop on the string.
- 4. Inside we check if we dont find the key in the hashmap already, then we increase the current\_length by 1 and put the key into the hashmap with value as the index.
- Else if we find the key in the hashmap, if the key is already present with a range of I, we do update I to 1 plus the last occurance of the current key and set value for current key in the key value pair.
- Keep a max and current length count and return the max\_len after the end of the function.

The code:

```
class Solution {
 int lengthOfLongestSubstring(string s) {
    int I = 0;
    int r = 0;
    int max_len = 0;
    int curr_len = 0;
    unordered_map<char, int> pos_map;
    for(int i=0;i<s.size();i++){
      if(pos_map.find(s[i]) != pos_map.end()&&pos_map[s[i]]>=l){
        l = pos_map[s[i]] + 1;
         pos_map[s[i]] = i;
        r++;
        curr_len = r - l;
        max_len = max(max_len,curr_len);
        else{
         pos_map[s[i]] = i;
        curr_len = r-l;
        max_len = max(max_len,curr_len);
      return max_len;
```

ab Chaa

I for a char

I for a