

## Exercise Sheet: Dijkstra

### Data Structures and Algorithms (X\_400614)

1. Execute Dijkstra's algorithm on the graph pictured in Figure 1, starting at the vertex A. In case of ties, the vertex with the lower letter is handled first.
  - a. List the vertices in the order in which they are deleted from the priority queue and for each the shortest distance from A to the vertex.  
**A,0 D,3 B,4 G,4 C,6 E,6 F,7**
  - b. Draw the shortest path tree that results.

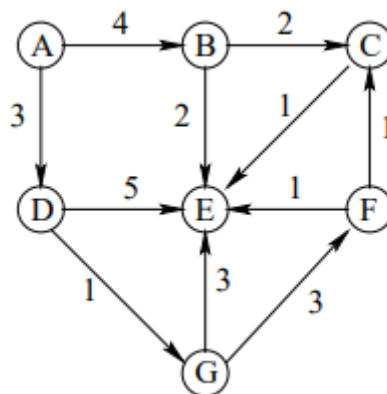
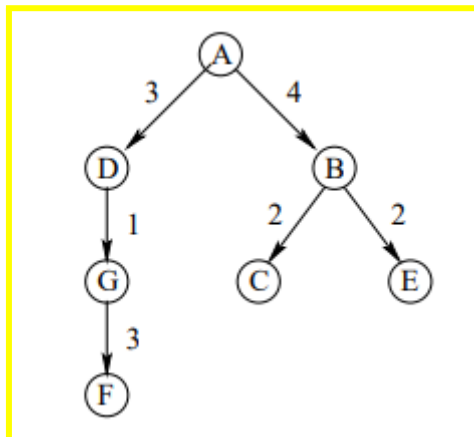


Figure 1.

2. Consider the graph in Figure 2.
  - a. Execute Dijkstra's algorithm on this graph, starting from s. After each step, note down:
    - i. The upper bounds  $d[u]$  (which denote what is currently believed to be the shortest distance to each node), for  $u \in V$ , between s and each node  $u$  computed so far,
    - ii. the set  $M$  of all nodes for which the minimal distance has been correctly computed so far,
    - iii. and the predecessor  $p(u)$  for each node in  $M$ .

**When we choose s:  $d[s] = 0$ ,  $d[x] = d[v] = d[w] = d[y] = d[t] = \infty$ ,  $M = \{s\}$ , there is no  $p(s)$ .**

**When we choose x:  $d[s] = 0$ ,  $d[x] = 2$ ,  $d[v] = 20$ ,  $d[w] = d[y] = d[t] = \infty$ ,**

$M = \{s, x\}$ , there is no  $p(s)$ ,  $p(x) = s$ .

When we choose  $y$ :  $d[s] = 0$ ,  $d[x] = 2$ ,  $d[v] = 10$ ,  $d[y] = 3$ ,  $d[w] = d[t] = \infty$ ,  $M = \{s, x, v, y\}$ , there is no  $p(s)$ ,  $p(x) = s$ ,  $p(v) = x$ ,  $p(y) = x$ .

When we choose  $t$ :  $d[s] = 0$ ,  $d[x] = 2$ ,  $d[v] = 10$ ,  $d[y] = 3$ ,  $d[w] = 20$ ,  $d[t] = 6$ ,  $M = \{s, x, v, y, t\}$ , there is no  $p(s)$ ,  $p(x) = s$ ,  $p(v) = x$ ,  $p(y) = x$ ,  $p(t) = y$ .

When we choose  $v$ :  $d[s] = 0$ ,  $d[x] = 2$ ,  $d[v] = 10$ ,  $d[y] = 3$ ,  $d[w] = 20$ ,  $d[t] = 6$ ,  $M = \{s, x, v, y, t\}$ , there is no  $p(s)$ ,  $p(x) = s$ ,  $p(v) = x$ ,  $p(y) = x$ ,  $p(t) = y$ .

When we choose  $w$ :  $d[s] = 0$ ,  $d[x] = 2$ ,  $d[v] = 10$ ,  $d[y] = 3$ ,  $d[w] = 15$ ,  $d[t] = 6$ ,  $M = \{s, x, v, y, t, w\}$ , there is no  $p(s)$ ,  $p(x) = s$ ,  $p(v) = x$ ,  $p(y) = x$ ,  $p(t) = y$ ,  $p(w) = v$ .

- b. Replace the weight of edge  $(x,y)$  with  $-1$  and redo the algorithm. Does the algorithm still compute correctly? If not, where does it break?

The algorithm works correctly.

When we choose  $s$ :  $d[s] = 0$ ,  $d[x] = d[v] = d[w] = d[y] = d[t] = \infty$ .

When we choose  $x$ :  $d[s] = 0$ ,  $d[x] = 2$ ,  $d[v] = 20$ ,  $d[w] = d[y] = d[t] = \infty$ .

When we choose  $y$ :  $d[s] = 0$ ,  $d[x] = 2$ ,  $d[v] = 10$ ,  $d[y] = 1$ ,  $d[w] = d[t] = \infty$ .

When we choose  $t$ :  $d[s] = 0$ ,  $d[x] = 2$ ,  $d[v] = 10$ ,  $d[y] = 1$ ,  $d[t] = 4$ ,  $d[w] = 18$ .

When we choose  $v$ :  $d[s] = 0$ ,  $d[x] = 2$ ,  $d[v] = 10$ ,  $d[y] = 1$ ,  $d[t] = 4$ ,  $d[w] = 18$ .

When we choose  $w$ :  $d[s] = 0$ ,  $d[x] = 2$ ,  $d[v] = 10$ ,  $d[y] = 1$ ,  $d[t] = 4$ ,  $d[w] = 15$ .

- c. Now, additionally change the weight of edge  $(v,y)$  to  $-10$ . Show that in this case the algorithm doesn't work correctly. That is, show that there is a vertex  $u \in V$  such that  $d[u]$  is not equal to minimum distance from  $s$  to  $u$  after the execution of the algorithm.

The algorithm doesn't work correctly, for example, the distance from  $s$  to  $y$  is  $0$ , but the algorithm computes exactly the same values of  $d[\cdot]$  as in part b), so  $d[y] = 1$ .

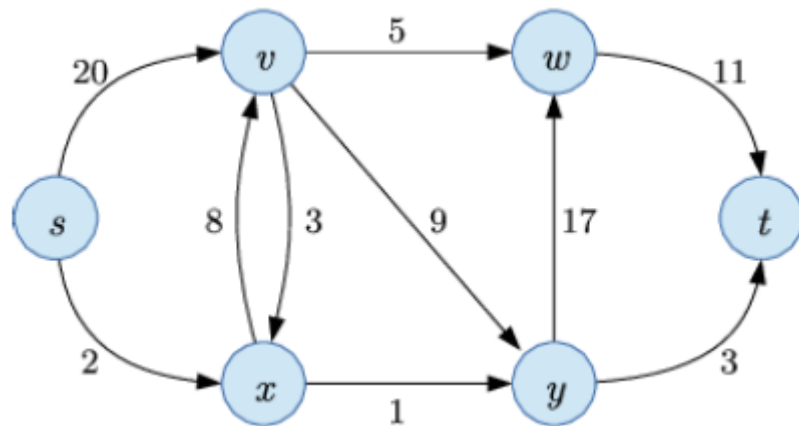


Figure 2.

3. A robot moves in the environment in Figure 3. It starts from the node labeled *Start* and needs to reach the node labeled *End*. The environment is continuous and the scale is supplied on the figure.

- a. Considering the robot as a point, what is the shortest path from *Start* to *End*? Compute the Euclidean distance between the nodes. Draw the graph and apply Dijkstra's algorithm to find the shortest path from the node *Start* to the node *End*.

Initially only Start has distance 0. In the first step we compute the distance from Start to 1, which is  $\sqrt{5}$ , and from Start to 2, which is  $\sqrt{29}$ . These are set as the distances for 1 and 2.

In the next step we choose 1 and see that the distance from 1 to 3 is  $\sqrt{26}$ .

Thus, the distance from Start to 3 is set to  $\sqrt{5} + \sqrt{26}$ .

The next step will consider 2, and set the distance from 2 to end to  $\sqrt{29} + \sqrt{17}$ .

The distance from Start to 3 is unchanged as Start - 1 - 3 is shorter than Start - 2 - 3.

Finally, considering 3 doesn't change any distances as  $\sqrt{29} + \sqrt{17} < 2 * \sqrt{5} + \sqrt{26}$ .

Shortest path	Start	1	2	3	End
Init	0	inf	inf	inf	inf
Start	–	$\sqrt{5}$	$\sqrt{29}$	inf	inf
1	–	–	$\sqrt{29}$	$\sqrt{5} + \sqrt{26}$	inf
2	–	–	–	$\sqrt{5} + \sqrt{26}$	$\sqrt{29} + \sqrt{17}$
3	–	–	–	–	$\sqrt{29} + \sqrt{17}$

The shortest path is Start-2-End.

- b. Can you find the same path faster using another algorithm? If so, which algorithm is better suited for this problem, and why?

Yes, A\* is better. It will start exploring in the right direction, whereas Dijkstra is not aware of the directionality of the shortest path.

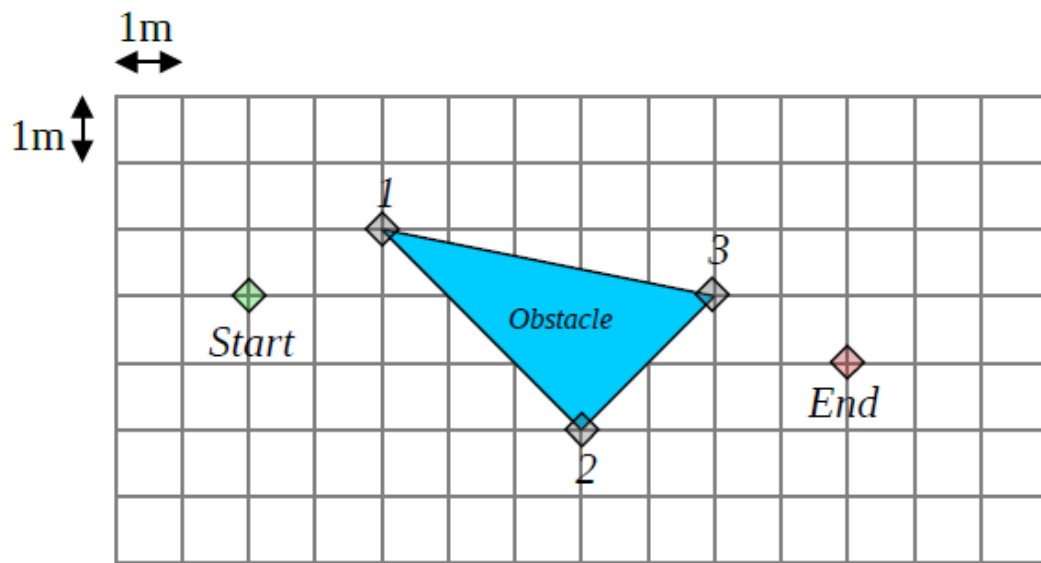


Figure 3.