

Exercise Sheet: Dijkstra

Data Structures and Algorithms (X_400614)

1. Execute Dijkstra's algorithm on the graph pictured in Figure 1, starting at the vertex A. In case of ties, the vertex with the lower letter is handled first.
 - a. List the vertices in the order in which they are deleted from the priority queue and for each the shortest distance from A to the vertex.
A,0 D,3 B,4 G,4 C,6 E,6 F,7
 - b. Draw the shortest path tree that results.

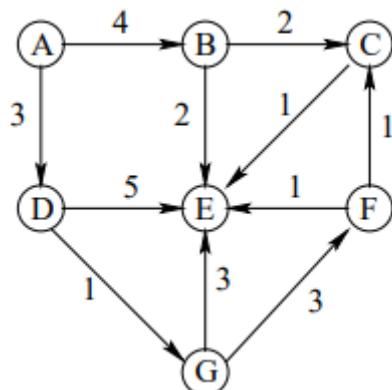
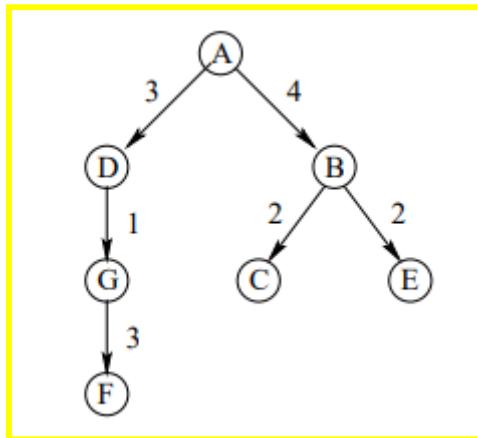


Figure 1.

2. Consider the graph in Figure 2.
 - a. Execute Dijkstra's algorithm on this graph, starting from s. After each step, note down:
 - i. The upper bounds $d[u]$ (which denote what is currently believed to be the shortest distance to each node), for $u \in V$, between s and each node u computed so far,
 - ii. the set M of all nodes for which the minimal distance has been correctly computed so far,
 - iii. and the predecessor $p(u)$ for each node in M .
- When we choose s: $d[s] = 0$, $d[x] = d[v] = d[w] = d[y] = d[t] = \infty$, $M = \{s\}$, there is no $p(s)$.**
- When we choose x: $d[s] = 0$, $d[x] = 2$, $d[v] = 20$, $d[w] = d[y] = d[t] = \infty$,**

$M = \{s, x\}$, there is no $p(s)$, $p(x) = s$.

When we choose y : $d[s] = 0$, $d[x] = 2$, $d[v] = 10$, $d[y] = 3$, $d[w] = d[t] = \infty$, $M = \{s, x, v, y\}$, there is no $p(s)$, $p(x) = s$, $p(v) = x$, $p(y) = x$.

When we choose t : $d[s] = 0$, $d[x] = 2$, $d[v] = 10$, $d[y] = 3$, $d[w] = 20$, $d[t] = 6$, $M = \{s, x, v, y, t\}$, there is no $p(s)$, $p(x) = s$, $p(v) = x$, $p(y) = x$, $p(t) = y$.

When we choose v : $d[s] = 0$, $d[x] = 2$, $d[v] = 10$, $d[y] = 3$, $d[w] = 20$, $d[t] = 6$, $M = \{s, x, v, y, t\}$, there is no $p(s)$, $p(x) = s$, $p(v) = x$, $p(y) = x$, $p(t) = y$.

When we choose w : $d[s] = 0$, $d[x] = 2$, $d[v] = 10$, $d[y] = 3$, $d[w] = 15$, $d[t] = 6$, $M = \{s, x, v, y, t, w\}$, there is no $p(s)$, $p(x) = s$, $p(v) = x$, $p(y) = x$, $p(t) = y$, $p(w) = v$.

- b. Replace the weight of edge (x,y) with -1 and redo the algorithm. Does the algorithm still compute correctly? If not, where does it break?

The algorithm works correctly.

When we choose s : $d[s] = 0$, $d[x] = d[v] = d[w] = d[y] = d[t] = \infty$.

When we choose x : $d[s] = 0$, $d[x] = 2$, $d[v] = 20$, $d[w] = d[y] = d[t] = \infty$.

When we choose y : $d[s] = 0$, $d[x] = 2$, $d[v] = 10$, $d[y] = 1$, $d[w] = d[t] = \infty$.

When we choose t : $d[s] = 0$, $d[x] = 2$, $d[v] = 10$, $d[y] = 1$, $d[t] = 4$, $d[w] = 18$.

When we choose v : $d[s] = 0$, $d[x] = 2$, $d[v] = 10$, $d[y] = 1$, $d[t] = 4$, $d[w] = 18$.

When we choose w : $d[s] = 0$, $d[x] = 2$, $d[v] = 10$, $d[y] = 1$, $d[t] = 4$, $d[w] = 15$.

- c. Now, additionally change the weight of edge (v,y) to -10 . Show that in this case the algorithm doesn't work correctly. That is, show that there is a vertex $u \in V$ such that $d[u]$ is not equal to minimum distance from s to u after the execution of the algorithm.

The algorithm doesn't work correctly, for example, the distance from s to y is 0 , but the algorithm computes exactly the same values of $d[\cdot]$ as in part b), so $d[y] = 1$.

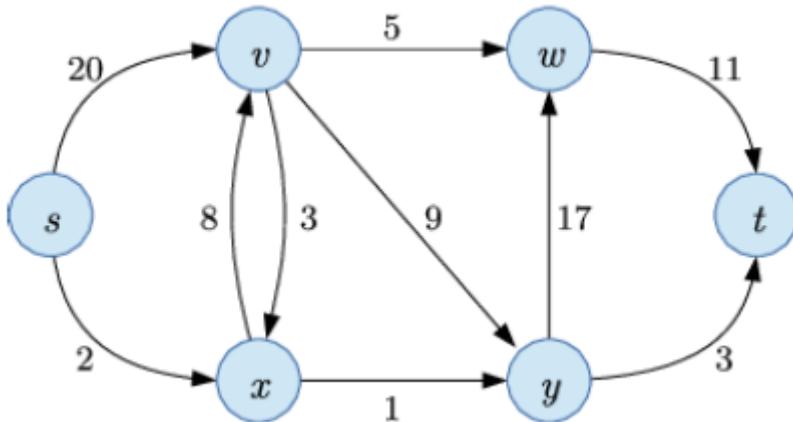


Figure 2.

3. A robot moves in the environment in Figure 3. It starts from the node labeled *Start* and needs to reach the node labeled *End*. The environment is continuous and the scale is supplied on the figure.

- a. Considering the robot as a point, what is the shortest path from *Start* to *End*? Compute the Euclidean distance between the nodes. Draw the graph and apply Dijkstra's algorithm to find the shortest path from the node *Start* to the node *End*.

Initially only Start has distance 0. In the first step we compute the distance from Start to 1, which is $\sqrt{5}$, and from Start to 2, which is $\sqrt{29}$. These are set as the distances for 1 and 2.

In the next step we choose 1 and see that the distance from 1 to 3 is $\sqrt{26}$. Thus, the distance from Start to 3 is set to $\sqrt{5} + \sqrt{26}$.

The next step will consider 2, and set the distance from 2 to end to $\sqrt{29} + \sqrt{17}$. The distance from Start to 3 is unchanged as Start - 1 - 3 is shorter than Start - 2 - 3.

Finally, considering 3 doesn't change any distances as $\sqrt{29} + \sqrt{17} < 2 * \sqrt{5} + \sqrt{26}$.

Shortest path	Start	1	2	3	End
Init	0	inf	inf	inf	inf
Start	-	sqrt 5	sqrt 29	inf	inf
1	-	-	sqrt 29	sqrt 5 + sqrt 26	inf
2	-	-	-	sqrt 5 + sqrt 26	sqrt 29 + sqrt 17
3	-	-	-	-	sqrt 29 + sqrt 17

The shortest path is Start-2-End.

- b. Can you find the same path faster using another algorithm? If so, which algorithm is better suited for this problem, and why?

Yes, A* is better. It will start exploring in the right direction, whereas Dijkstra is not aware of the directionality of the shortest path.

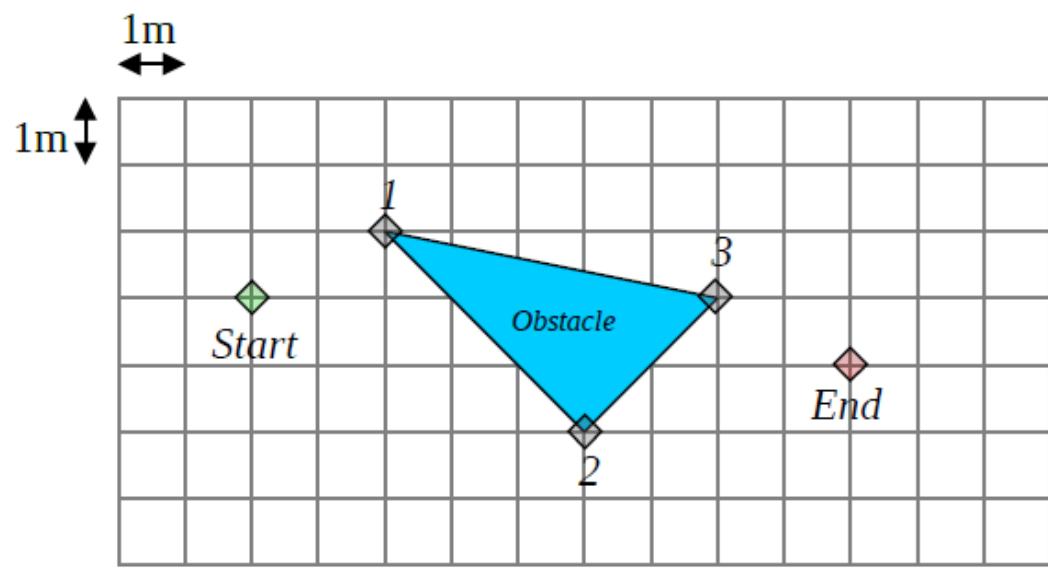


Figure 3.