

OSCORE-capable Proxies

draft-tiloca-core-oscore-capable-proxies-07

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Recap

- › **A CoAP proxy (P) can be used between client (C) and server (S)**
 - A security association might be required between C and P --- use cases in next slides
- › **Good to use OSCORE between C and P**
 - Especially, but not only, if C and S already use OSCORE end-to-end
- › **This is not defined and not admitted in OSCORE (RFC 8613)**
 - C and S are the only considered “OSCORE endpoints”
 - It is forbidden to double-protect a message, i.e., both over $C \leftrightarrow S$ and over $C \leftrightarrow P$
- › **This started as an Appendix of *draft-tiloca-core-groupcomm-proxy***
 - Agreed at IETF 110 [1] and at the June 2021 CoRE interim [2] to have a separate draft

[1] <https://datatracker.ietf.org/doc/minutes-110-core-202103081700/>

[2] <https://datatracker.ietf.org/doc/minutes-interim-2021-core-07-202106091600/>

Contribution

› Twofold update to RFC 8613

1. Define the use of OSCORE in a communication leg including a proxy

- › Between origin client/server and a proxy; or between two proxies in a chain
- › Not only an origin client/server, but also an intermediary can be an “OSCORE endpoint”

2. Explicitly admit nested OSCORE protection – “OSCORE-in-OSCORE”

- E.g., first protect end-to-end over $C \leftrightarrow S$, then further protect the result over $C \leftrightarrow P$
- Typically, at most 2 OSCORE “layers” for the same message
 - › 1 end-to-end + 1 between two adjacent hops
- Possible to seamlessly apply 2 or more OSCORE layers to the same message

› Focus on OSCORE, but the same applies “as is” to Group OSCORE

Several use cases

- › **Section 2.1, CoAP group communication through a proxy [3]**
 - The proxy identifies the client before forwarding
- › **Section 2.2, Observe multicast notifications with Group OSCORE [4]**
 - The client securely provides the Ticket Request to the proxy
- › **Sections 2.3 and 2.4, OMA Lightweight Machine-to-Machine (LwM2M)**
 - The LwM2M Client uses the LwM2M Server as proxy towards External Application Servers
 - The LwM2M Server uses the LwM2M Gateway as reverse proxy towards External End Devices
- › **Further use cases are listed in Section 2.5**
 - Transport indication through trusted proxies – *draft-ietf-core-transport-indication*
 - CoAP performance measurements involving on-path probes – *draft-ietf-core-coap-pm*
 - EST over OSCORE through a CoAP-to-HTTP proxy – *draft-ietf-ace-coap-est-oscore*
 - OSCORE-protected “onion forwarding”, a la TOR – *draft-amsuess-t2trg-onion-coap*
 - Proxies as entry point to a firewalled network

[3] <https://datatracker.ietf.org/doc/draft-tiloca-core-groupcomm-proxy/>

[4] <https://datatracker.ietf.org/doc/draft-ietf-core-observe-multicast-notifications/>

Message processing

› **Stable and well defined**

- No need for an explicit signaling method to guide the message processing
- High-level general algorithm, fitting a client, proxy or server as a message processor

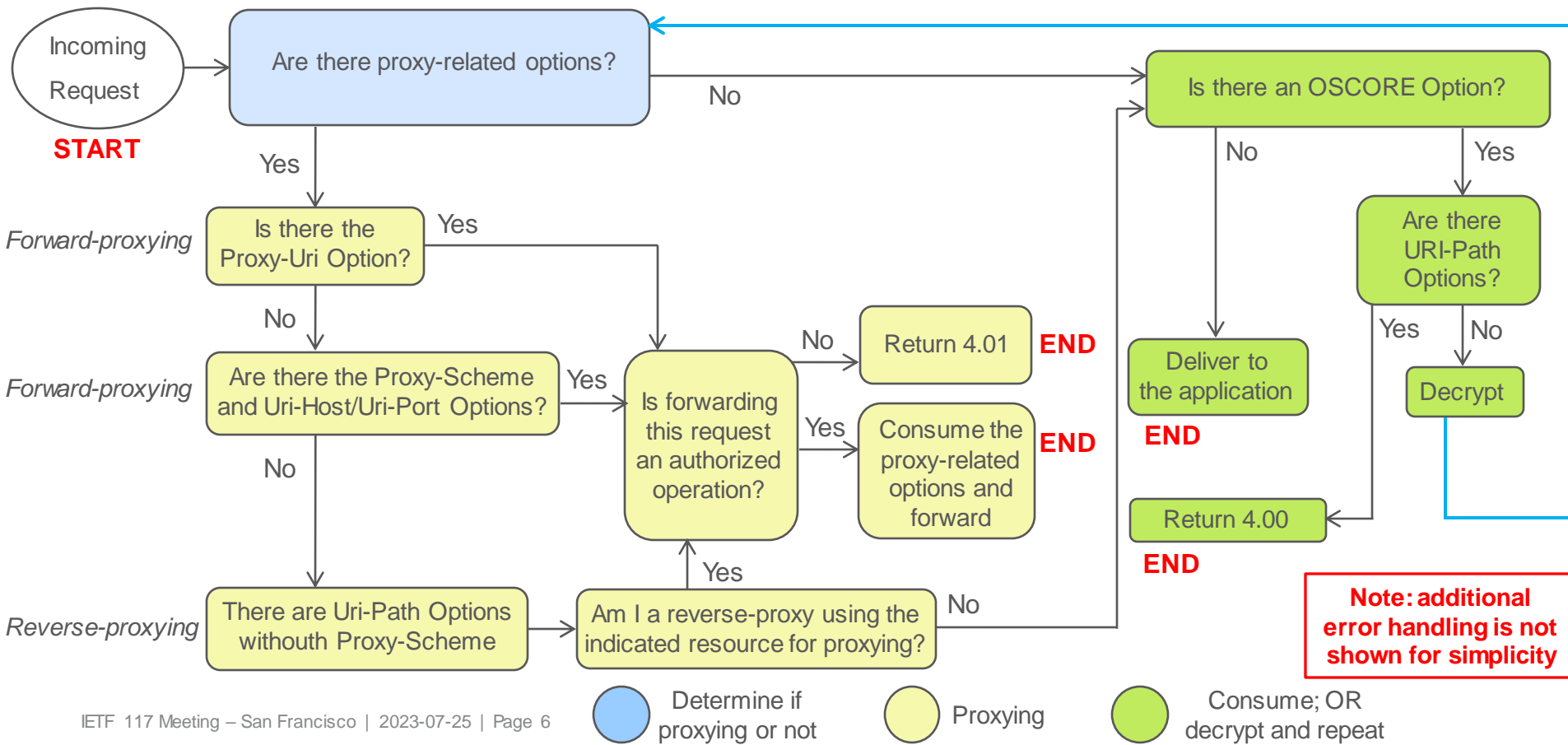
› **Dedicated Section 3.2-3.5 for each message and direction**

- Outgoing request, incoming request (most delicate), outgoing response, incoming response

› **Recent clarifications suggested by Christian (already in version -06)**

- A proxy performs authorization checks, before forwarding an incoming request
- For example, based on the OSCORE Security Context locally used for decryption

Processing an incoming request



Recent updates (already in v -06)

› **OSCORE protection of CoAP options in outgoing messages**

- If a CoAP option is originally defined as class U or I for OSCORE ...
- ... when should it be fully protected anyway, as if it was of class E?

› **Improved general rules**

- Now better covering corner cases and class I options
- A good sanity check was the Request-Hash option [5]
 - › Processed as class I in responses
 - › Expected to be elided from responses, but still possible to send it on the wire

› **Current rule formulation – Section 3.1**

- Three general cases, all phrased as “Any CoAP option such that ...”.
- The rationale is to encrypt as many options as possible. If there is a match, encrypt!
- Multiple examples are provided for each case

Recent updates (already in v -06)

› **New Section 5 – Guidelines on establishing OSCORE Security Contexts**

- Generally agnostic of the used establishment method

› **For OSCORE**

- Guidelines for the client using EDHOC [6], first with the proxy, then with the origin server
- Reference to the possible to use, optimized EDHOC workflow [7]

› **For Group OSCORE**

- Expected between origin client and servers, which rely on the Group Manager

Updates in version -07

› **Section 2.5: mentioned two additional use cases**

- CoAP performance measurements involving on-path probes – *draft-ietf-core-coap-pm*
- EST over OSCORE through a CoAP-to-HTTP proxy – *draft-ietf-ace-coap-est-oscore*
- Both those documents include an informative reference to this document

› **Section 5: clarifications when Group OSCORE is used – Thanks Christian!**

- New text: no need for a proxy to be in the same OSCORE group of the origin client and server
- Not forbidden altogether anymore, as it was in version -06; it might be fine and desirable

› **Section 7: added first security considerations**

- Inherited from OSCORE, Group OSCORE, and the specifically used CoAP options
- Any OSCORE endpoint in the chain enjoys the same, usual (Group) OSCORE properties

Updates in version -07

› Added new Section 6, on using SCHC (RFC 8724)

- The compression of CoAP headers with SCHC is already defined in RFC 8824
- When OSCORE is used, an inner and outer compression are performed on a message
- *draft-tiloca-schc-8824-update* spells out (but does not extend) how this works with proxies
 - › The inner compression is end-to-end; the outer compression is hop-by-hop

› Section 6 generalizes the above for nested OSCORE protections

- No changes to the core mechanics of SCHC

› When a sender endpoint produces an outgoing message, it performs:

- One inner compression for each applied OSCORE layer
 - › Each for the OSCORE endpoint intended to decrypt and strip that OSCORE layer
- Exactly one outer compression, after all the inner compressions
 - › Intended for the (next-hop towards the) recipient origin endpoint

Updates in version -07

› **Appendix A: improved notation in the 5 examples of message exchange**

- Easier to see what is encrypted, and which OSCORE Security Context is used
- Easier to follow the sub-steps taken by each endpoint
- Two examples consider EDHOC [6], one of which with the optimized workflow [7]

› **Removed old Appendix B**

- It included early specification notes towards “OSCORE-protected onion forwarding”
- Discussed and agreed to remove it during the CoRE interim meeting on 2023-06-07
- The removed and already revised content is now an Experimental document in T2TRG [8]

› **Added new Appendix B**

- State diagram in ASCII art, showing the message processing of incoming requests
- Based on the earlier slide 6, plus additional error handling; consistent with Section 3.3

[6] <https://datatracker.ietf.org/doc/draft-ietf-lake-edhoc/>

[7] <https://datatracker.ietf.org/doc/draft-ietf-core-oscore-edhoc/>

[8] <https://datatracker.ietf.org/doc/draft-amsuess-t2trg-onion-coap/>

Summary

› **Proposed update to RFC 8613**

- Define the use of OSCORE in a communication leg including a proxy
- Explicitly admit nested OSCORE protection: “OSCORE-in-OSCORE” / “Matryoscore”
- Seamlessly usable also with Group OSCORE

› **Stable approach and mechanics**

- Signaling-free message processing; CoAP options are protected as much as possible

› **Additional material**

- On establishing of an OSCORE Security Context with proxies
- On compressing CoAP headers with SCHC (RFC 8824)
- Examples of message exchanges, also considering key establishment with EDHOC

› **The authors believe that version -07 is ready for a WG Adoption Call**

Thank you!

Comments/questions?

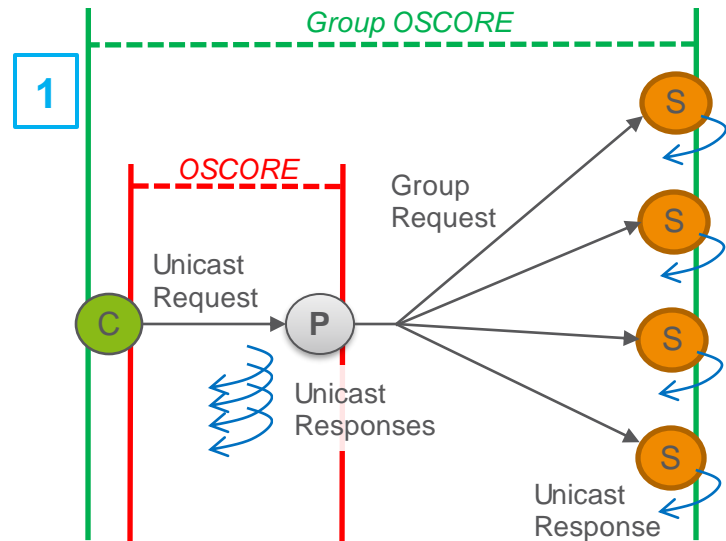
<https://gitlab.com/crimson84/draft-tiloca-core-oscore-to-proxies>

Backup

Use cases

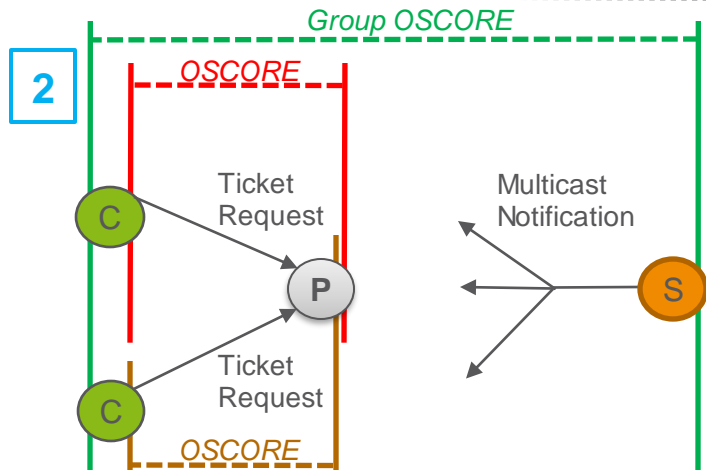
1. CoAP Group Communication with Proxies

- *draft-tiloca-core-groupcomm-proxy*
- CoAP group communication through a proxy
- P must identify C through a security association



2. CoAP Observe Notifications over Multicast

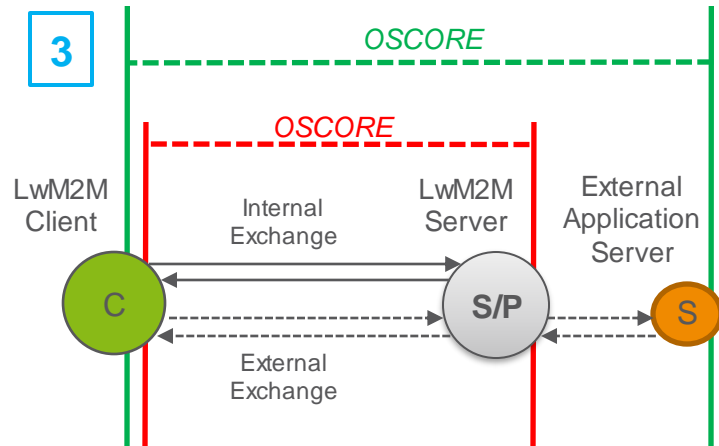
- *draft-ietf-core-observe-multicast-notifications*
- If Group OSCORE is used for end-to-end security ...
- ... C provides P with a Ticket Request obtained from S
- That provisioning should be protected over $C \leftrightarrow P$



Use cases

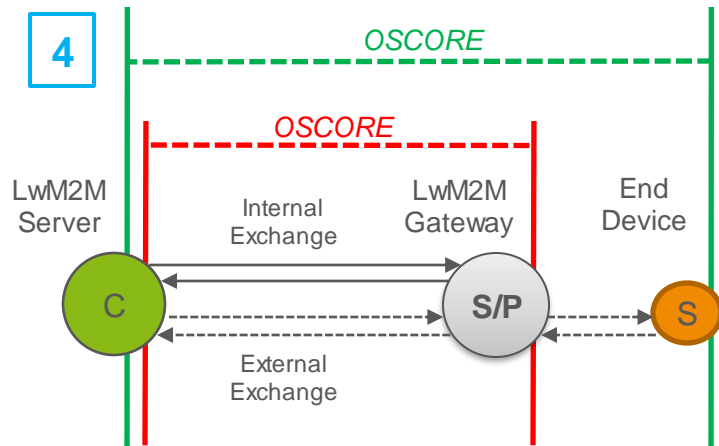
3. LwM2M Client and external Application Server

- From the *L2wM2M Transport Bindings* specification:
 - › OSCORE can be used between a LwM2M endpoint and a non-LwM2M endpoint, via the LwM2M Server
- The LwM2M Client may use OSCORE to interact:
 - › With the LwM2M Server (LS), as usual; and
 - › With an external Application Server, via LS acting as proxy



4. Use of the LwM2M Gateway (from David Navarro)

- It provides the LwM2M Server with access to:
 - a) Resources at the LwM2M Gateway
 - b) Resources at external End Devices, through the LwM2M Gateway, via dedicated URI paths
- In case (b), the LwM2M Gateway acts, at its core, as a reverse-proxy



Use case 3 – LwM2M

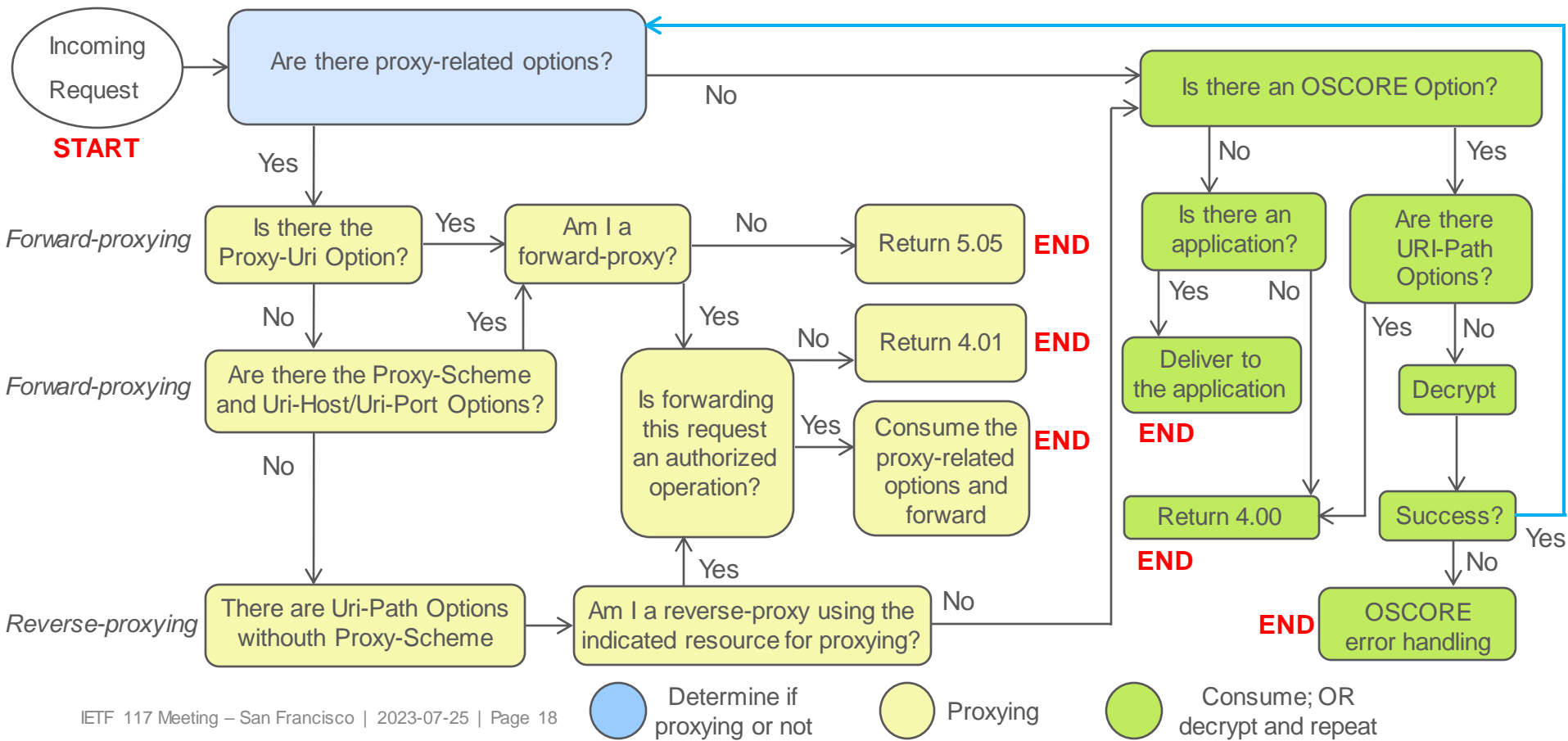
› OMA LwM2M Client and External Application Server

– *Lightweight Machine to Machine Technical Specification – Transport Binding*

OSCORE MAY also be used between LwM2M endpoint and non-LwM2M endpoint, e.g., between an Application Server and a LwM2M Client via a LwM2M server. Both the LwM2M endpoint and non-LwM2M endpoint MUST implement OSCORE and be provisioned with an OSCORE Security Context.

- The LwM2M Client may register to and communicate with the LwM2M Server using OSCORE
- The LwM2M Client may communicate with an External Application Server, also using OSCORE
- The LwM2M Server would act as CoAP proxy, forwarding traffic outside the LwM2M domain

Processing an incoming request



Encryption of Class U/I Options

