# AEAD Key Usage Limits in OSCORE Key Update for OSCORE

draft-hoeglund-core-oscore-key-limits-01

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#### Draft Overview (1/2)

- OSCORE (RFC8613) uses AEAD algorithms to provide security
  - Confidentiality and Integrity
- > Need to follow limits in key usage and failed decryptions, before rekeying
  - Otherwise, it is possible to break the security properties of the AEAD algorithm
  - Reference draft-irtf-cfrg-aead-limits-03
- > (1) AEAD limits and their impact on OSCORE
  - Defining appropriate limits for OSCORE
  - Originally starting from the same assumptions in TLS
  - Revisited based on John Mattsson's input at the April CoRE interim
    - https://datatracker.ietf.org/meeting/110/materials/slides-110-saag-analysis-of-usage-limits-of-aead-algorithms-00.pdf

#### Draft Overview (2/2)

- > (2) Updates to OSCORE
  - Counters in the Security Context: key encryption use (q) and invalid decryptions (v)
  - Necessary steps to take during message processing (counting)
  - Update the keys when the limits are exceeded (rekeying)
- > (3) Defined a new method for rekeying OSCORE (**new**)
  - Loosely inspired by Appendix B.2 of OSCORE
  - Goal: renew the Master Secret and Master Salt; derive new keys from those
  - Achieves Perfect Forward Secrecy

### Key Limits (1/2)

- Selected fixed values for 'q', 'v', and 'l'
  - $-q = 2^20$ ,  $v = 2^20$  and  $l = 2^8$  // 'l' is the max message length in cipher blocks
  - Based on earlier discussions and John Mattsson's presentation
- > Table with 'IA' and 'CA' probabilities based on those values
  - These are based on the formulas in the CFRG document

±		L
Algorithm name	IA probability	CA probability
AEAD_AES_128_CCM AEAD_AES_128_GCM AEAD_AES_256_GCM AEAD_CHACHA20_POLY1305	2^-68 2^-99 2^-99 2^-75	2^-70 2^-89 2^-89 -
+	+	

Figure 1: Probabilities for algorithms based on chosen q, v and l values.

#### Integrity Advantage (IA):

Probability of breaking integrity properties

# Confidentiality Advantage (CA): Probability of breaking confidentiality properties

#### Key Limits (2/2)

- > Specific look at AEAD\_AES\_128\_CCM\_8
  - Due to short Tag length the limits can be most problematic here
- > Table with 'IA' and 'CA' probabilities for various values of 'q', 'v' and 'I'

4		L			L
'q', 'v' and 'l'	IA probability	CA probability	'q', 'v' and 'l'	IA probability	CA probability
q=2^20, v=2^20, l=2^8   q=2^15, v=2^20, l=2^8   q=2^10, v=2^20, l=2^8   q=2^20, v=2^15, l=2^8   q=2^15, v=2^15, l=2^8   q=2^10, v=2^15, l=2^8   q=2^10, v=2^15, l=2^8   q=2^20, v=2^10, l=2^8   q=2^10, v=2^10, l=2^8   q=2^10, v=2^10, l=2^8	2^-44 2^-44 2^-49 2^-49 2^-49 2^-54 2^-54	2^-70 2^-80 2^-90 2^-70 2^-80 2^-90 2^-70 2^-80 2^-80 2^-90	q=2^20, v=2^20, 1=2^6   q=2^15, v=2^20, 1=2^6   q=2^10, v=2^20, 1=2^6   q=2^20, v=2^15, 1=2^6   q=2^15, v=2^15, 1=2^6   q=2^10, v=2^15, 1=2^6   q=2^20, v=2^10, 1=2^6   q=2^15, v=2^10, 1=2^6   q=2^10, v=2^10, 1=2^6	2^-44 2^-44 2^-49 2^-49 2^-49 2^-54 2^-54	2^-74 2^-84 2^-94 2^-74 2^-84 2^-94 2^-74 2^-84 2^-94
		L	. +		

Figure 2: Probabilities for AEAD\_AES\_128\_CCM\_8 based on chosen q, v and l values.

- > From the considered values the best triple is  $(q = 2^20, v = 2^10, l = 2^8)$ 
  - The question is if an IA of 2^-54 is good enough?

#### OSCORE Key Update method (1/3)

- Defined a new method for rekeying OSCORE
  - Client and server exchange two nonces R1 and R2
  - UpdateCtx() function for deriving new OSCORE Security Context using the nonces
  - Current Sec Ctx (to renew) ==> Intermediate Sec Ctx ==> New Sec Ctx

#### > Properties

- Only one intermediate Security Context is derived
- The ID Context does not change
- Can be initiated by either the client or server
- It is robust and secure against a peer rebooting
- The procedure completes in one round-trip (after that, the new context can be used)
- Compatible with possible prior key establishment through the EDHOC protocol

#### OSCORE Key Update method (2/3)

- › Key update messages are OSCORE-protected and self-evident
- OSCORE Option: defined the use of flag bit 1 to signal presence of flag bits 8-15
- Defined flag bit 15 -- 'd' -- to indicate:
  - This is a OSCORE key update message
  - "id detail" is specified (length + value); used to transport a nonce for the key update

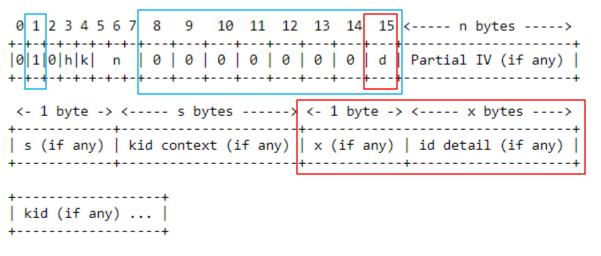


Figure 3: The OSCORE option value, including 'id detail'

### OSCORE Key Update method (3/3)

```
Client
                (initiator)
                                      (responder)
Generate R1
CTX 1 =
  updateCtx(R1,
            CTX OLD)
                           Request #1
Protect with CTX 1
                       OSCORE Option:
                                             CTX 1 =
                                               updateCtx(R1,
                         d flag: 1
                                                      CTX OLD)
                         ID Detail: R1
                                             Verify with CTX 1
                          . . .
                                             Generate R2
                                             CTX NEW =
                                               updateCtx(R1|R2,
                                                      CTX OLD)
                           Response #1
                                             Protect with CTX NEW
CTX NEW =
                       OSCORE Option:
  updateCtx(R1|R2,
            CTX OLD)
                         d flag: 1
                         ID Detail: R1|R2
Verify with CTX NEW
Discard CTX OLD
```

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```
updateCtx( N, CTX IN ) {
 CTX OUT
              // The new Security Context
 MSECRET NEW // The new Master Secret
 MSALT NEW // The new Master Salt
 if <the original Security Context was established through EDHOC> {
   EDHOC-KevUpdate( N )
   // This results in updating the key PRK 4x3m of the EDHOC session,
   // i.e., PRK 4x3m = Extract(N, PRK <math>4x3m)
   MSECRET NEW = EDHOC-Exporter( "OSCORE Master Secret", key length
     = EDHOC-KDF(PRK 4x3m, TH 4, "OSCORE Master Secret", key length
   MSALT NEW = EDHOC-Exporter( "OSCORE Master Salt", salt length )
     = EDHOC-KDF( PRK 4x3m, TH 4, "OSCORE Master Salt", salt length )
 else {
   Master Secret Length = < Size of CTX IN.MasterSecret in bytes >
   MSECRET NEW = HKDF-Expand-Label(CTX IN.MasterSecret, Label,
                                   N, Master Secret Length)
               = HKDF-Expand(CTX IN.MasterSecret, HkdfLabel,
                             Master Secret Length)
   MSALT NEW = N;
 < Derive CTX OUT using MSECRET NEW and MSALT NEW,
   together with other parameters from CTX IN >
 Return CTX OUT:
```

#### Summary and Next Steps

- > Twofold update to OSCORE
  - Tracking and reacting to defined key limits, to preserve security of the AEAD cipher
  - New efficient key update procedure (rekeying) with Perfect Forward Secrecy
    - Initially planned as a separate draft
    - > Preference at the April CoRE interim to have it in this document
- Take and adopt feedback on the new key limits and especially for CCM\_8
- Main next steps are tracked as Gitlab issues in [1]
- > Need for reviews, on both key limits and key update

[1] https://gitlab.com/rikard-sics/draft-hoeglund-oscore-rekeying-limits/

# Thank you!

## Comments/questions?

https://gitlab.com/rikard-sics/draft-hoeglund-oscore-rekeying-limits/