

21 2013. / 2014.

1. adresa 32-bitna

C1: 3 linije po 1B $s = n \cdot b = 3B$

$$w(p) = \log_2(b) = 0$$

$$w(l) = \log_2(n) = 3$$

$$w(o) = 32 - 3 = 29$$

0	
1	✓ +
2	✓
3	
4	✓
5	✓
6	✓
7	✓

1. 0x01

o=0 i=1 PROMAŠAJ

2. 0x86

1 0 0 0 0 1 1 0

i=6

gjena = 25T

C1: 2T

C2: 3T

C3: 5T

3. 0xD4

1 1 0 1 0 1 0 0

i=4 PROMAŠAJ

4. 0x01

i=1 POGODAK

ukupno 2 = C1: 502T

5. 0x37

i=7 PROMAŠAJ

6. D5

i=5 PROMAŠAJ

7. A2

i=2 PROMAŠAJ

8. 0x41

i=1 PROMAŠAJ, ZAMJENA

9. 0x02

i=2 PROMAŠAJ, ZAMJENA

10. 0x8c

i=11 0 0 1 1 PROMAŠAJ, ZAMJENA

11. 0x23

i=1 PROMAŠAJ, ZAMJENA

0 0 1 0 1 0 0 1

o=5

12. 0xDD

1 1 0 1 1 1 0 1

PROMAŠAJ, ZAMJENA

i=5 o=

2. LA = 32b
 stranica = 4KB
 opisnici = 4B
 program = 1GB $\times 2$

4432

a) jednorazinsko stranicenje

$$VP = \log_2(4KB) = 12 \quad VS = 32 - 12 = 20 \text{ bita}$$

$$2GB = \frac{2 \cdot 1024 \cdot 1024 \cdot 1024}{24 \cdot 1024} = \frac{2^{20}}{2^4} = 524288 \text{ opisnika pokrene}$$

$$\text{ukupno zauzede: } 2^{12} \cdot 4B = 2^{12} \cdot 2^2 B = 2^{14} B = 2MB$$

b) dvorazinsko stranicenje, stranicni imenik 256 stavki, po 6B.

$$VP = 12 \text{ bita} \quad VS1 = \log_2 256 = 8 \text{ bita}$$

$$VS2 = 12 \text{ bita}$$

$$\text{tablica ima: } 2^{12} = 128$$

$$\text{ukupno zauzede: } 256 \cdot 6B + 4096 \cdot 128 \cdot 4B = \underline{2MB}$$

3.

F	CPIA	CPIB	CPIc	CPIA
---	------	------	------	------

ukupno zauzeće: $256 \cdot 6B + 4096 \cdot 12B \cdot 4B = 2 \text{ MB}$

3.

	f	CPI_A	CPI_B	CPI_C	CPI_D
$P1$	1.5	1	2	4	4
$P2$	2	2	2	2	3
$P3$	2.5	3	1	4	1

$P = 10^6$ instrukcije

$A=10\%, B=20\%, C=50\%, D=20\%$

$$CPI_1 = 0.1 \cdot 1 + 0.2 \cdot 2 + 0.5 \cdot 4 + 0.2 \cdot 4 = 3.3$$

$$CPI_2 = 0.1 \cdot 2 + 0.2 \cdot 2 + 0.5 \cdot 2 + 0.2 \cdot 3 = 2.2$$

$$CPI_3 = 0.1 \cdot 3 + 0.2 \cdot 1 + 0.5 \cdot 4 + 0.2 \cdot 1 = 2.7$$

$$n_c = n_i \cdot CPI$$

$$n_{c1} = n_i \cdot CPI_1 = 3 \ 300 \ 000$$

$$t_{cpi1} = n_i \cdot CPI \cdot T_{cpi} = \frac{n_{c1}}{f_1} = 2.2 \text{ ms}$$

$$n_{c2} = n_i \cdot CPI_2 = 2 \ 200 \ 000$$

$$t_{cpi2} = \frac{n_{c2}}{f_2} = 1.1 \text{ ms}$$

$$n_{c3} = n_i \cdot CPI_3 = 2 \ 700 \ 000$$

$$t_{cpi3} = \frac{n_{c3}}{f_3} = 1.08 \text{ ms}$$

$$T_1 = \frac{1}{f} =$$

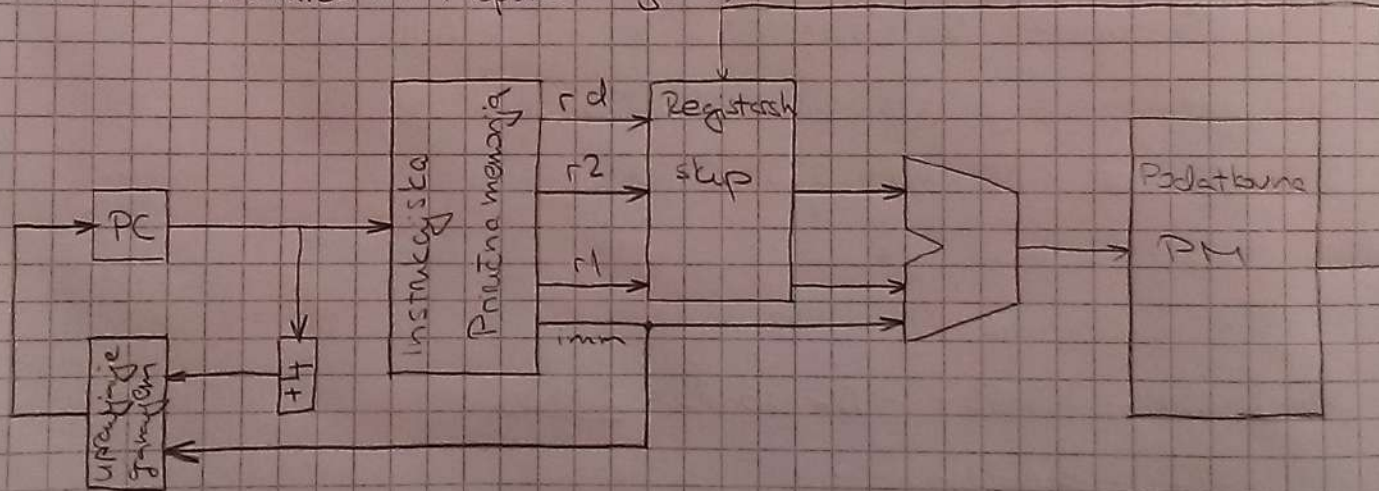
$P3$ će najbrže izvesti zadani program.

4. LW \$t5, -16(\$t5)
SW \$t5, -16(\$t5)
add \$t5, \$t5, \$t5

IF	ID	EX	MEM	WB		
	IF	ID	EX	MEM	WB	
		IF	ID	EX	MEM	WB

1. i 2. instrukcja	RAW hazard
2. i 3. instrukcja	WAR hazard

Posto je protačna arhitektura, nema izvateja izvan redoslijeda
pa stoga u ovoj arhitekturi ne postoji WAW i WAR hazard
(i bez dinamičkim rasporednjem)



b)

IF	ID	EX	MEM	WB					
	3	3	3	IF	ID	EX	MEM	WB	
					IF	ID	EX	MEM	WB

c)

IF	ID	EX	MEM	WB					
	IF	○	○	○	ID	EX	MEM	WB	
					IF	ID	EX	MEM	WB

6. LDA 31524
STA 31342

①
 A 1000 100
 2
 W
 D 34 2
 ← RET

IF	○	○	○	ID	EX	MEM	WB
IF	ID	EX	MEM	WB			

5. $\vec{C} = s(\vec{A} - \vec{B})$

VEKTORSKI KOD:

MOVE R0, #4 - inicijalizacija broja 4

LD R1, 0(\$s) - učitaj skalar

LOOP: LDV V0, 0(\$a) - učitaj \vec{A}

LDV V1, 0(\$b) - učitaj \vec{B}

LDV V2, 0(\$c) - učitaj \vec{C}

SUBV V2, V0, V1

MULVS V2, R1, V2

STV V2, 0(\$c)

ADD \$a, \$a, #8

ADD \$b, \$b, #8

ADD \$c, \$c, #8

SUB R0, R0, #1

BNE R0, #0, LOOP

6. LDA \$1324 // op. kod 34 Little endian - niži bajt, niža adresa
 STA \$3142 // op. kod A2

\$1320 : 34
 \$1321 : 24
 \$1322 : 13
 \$1323 : A2
 \$1324 : 42
 \$1325 : 31

