# **Computer Security Capstone**

Project 4: Capture The Flag (CTF)

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#### Goal

 Understand the exploitation of basic programming bugs, Linux system knowledge, and reverse-engineering

- You will learn about
  - □ Solving basic CTF problems
  - □ Investigating C/Linux functions deeply instead of simply using them
  - □ What buggy codes are and how they can be exploited

#### What is CTF?

- A traditional outdoor game
  - ☐ Two teams each have a flag
  - □ Objective: to capture the other team's flag



From Wikipedia

- In computer security, it is a type of cryptosport: a computer security competition
  - ☐ Giving participants experience in securing a machine
  - Required skills: reverse-engineering, network sniffing, protocol analysis, system administration, programming, etc.
  - □ How?
    - A set of challenges is given to competitors
    - Each challenge is designed to give a "Flag" when it is countered

# A CTF Example

A toy CTF

\$ python -c 'v = input(); print("flag:foobar") if v == "1" else print("failed")'

- ☐ You should enter "1" to pass the *if* statement and get the flag (flag:foobar)
- □ Otherwise, "failed" is obtained

### Requirements

- Linux/Unix environment is required
  - ☐ Connecting to our CTF servers for all the tasks

- You are NOT allowed to team up: one student one team
  - □ Discussions are allowed between teams, but any collaboration is prohibited
- TA: Cheng-I Hu

#### How to Proceed?

- Connecting to each CTF server: nc <ip> <port>
  - □ ip: 140.113.207.245
  - □ port is given at each problem
  - ☐ The program of each problem runs as a service at the server
  - ☐ You can do whatever you are allowed to do
- You can use python with pwntools, too

# How to Proceed? (Cont.)

- For each CTF problem, you should
  - □ analyze its given executable files or source code files
  - □ interact with the server to get a flag
  - ☐ The flag format: CSC2025{[a-zA-Z0-9\_]+}
- You will need to submit the programs
  - ☐ run the programs when you demo

#### What If Get Stuck?

- Learn to use "man" in UNIX-like systems
  - ☐ If you don't know something, ask "man"
  - □ e.g., what is man?
    - \$ man man
- Learn to find answers with FIRST-HAND INFORMATION/REFERENCE
  - □ Google is your best friend (Using ENGLISH KEYWORDS!!)
  - ☐ First-hand information: Wikipedia, cppreference.com, devel mailing-list, etc.
  - ☐ First-hand reference: papers, standards, spec, man, source codes, etc.
  - □ Second-hand information: blog, medium, ptt, reddit, stackoverflow post, etc.

#### Two Tasks

- Task I: Basic CTF problems (70%)
- Task II: Advance CTF problems (30%)
- Download all given executable and source files from e3
  - □ CTF Server using ubuntu 24.04 (for some problem to calculate address)

#### Task I: Basic CTF Problems

■ Task I-1: Password Checker (20%)

■ Task I-2: Secure Random (20%)

● Task I-3: Simple Shell (15%)

■ Task I-4: Simple ROP (15%)

#### Task I-1: Password Checker

- Goal: Learn how type conversion works in C/C++
- Server port: 30170

- Hints
  - □ <u>Implicit conversions of type</u>

#### Task I-2: Secure Random

- Goal: learn about the glibc PRNG
- Server port: 30171

- Hints
  - □ Is the random function really random?
    - Make sure you have time synchronization in your environment!!

# Task I-3: Simple Shell

- Goal: learn to identify basic logic flaw and buffer overflow in source codes
- Server port: 30172

- Hints
  - ☐ Inspect the code, where buffer overflow can occur?
  - What can you modify?
  - □ Inspect the impact of overflow by using gdb
    - You may want to install gdb extensions like gef

# Task I-4: Simple ROP

- Goal: Given buffer overflow, try to find a way to open up a shell for remote command execution!!
- Server port: 30173

- Hints
  - ☐ Inspect the code, where buffer overflow can occur?
  - □ With NX enabled, you cannot write shell code for buffer overflow
  - □ Stack buffer overflow
  - Return-oriented programming
  - You may want to use tools like ROPgadget to find gadget for ROP

#### Task II: Advance CTF Problems

• Task II-1: ret2Flag (10%)

● Task II-2: Simple RTOS (10%)

Task II-3: Hard ROP(10%)

# Task II-1: Ret2Flag

- Goal: Learn exploit buffer overflow to control program flow
- Server port: 30174

- Hints
  - □ Inspect the code, where buffer overflow can occur?
  - How can you bypass <u>canary</u> protection?
  - How can you find the function address?
    - You may want to find address that related with putFlag
  - ☐ Try to leak the information you need!!!

# Task II-2: Simple RTOS

- Goal: learn to identify dangerous function usage
- Server port: 30175

- Hints
  - ☐ How do you use printf normally?
    - Which conversion specifier can modify variable?
  - How can you return to the function you want?

#### Task II-3: Hard ROP

- Goal: Try to ROP with libc gadget!!
- Server port: 30176
- Hints
  - ☐ First, try to leaking every thing you need
  - ☐ Try to use libc gadget

# Important: How to Prepare Your Program?

- Must provide a Makefile which compiles your source codes into senven executable file
- You can use any language and library you want
  - ☐ Use your environment to demo
  - □ Do not hardcode the flag in your program
- Test requirements for your program
  - ☐ Do not need user interaction to get flag
    - For online tasks, you can only input server IP and port
    - For local tasks, you can only input file path
  - Must print flag to stdout

# Important: How to Demo Your Program?

- Download your code from e3
- Run make if needed
- Run your executables
- Ask some questions about your code
- Binary file for all task will not change
  - You can hardcode some symbol address if you need
  - ☐ FLAG during demo will change to avoid hardcode the flag

# **Project Submission**

- Due date: 5/28 11:55 p.m.
- Submission rules
  - □ Put all your files into a directory and name it using your student ID(s)
  - □ Zip the directory and upload the zip file to New e3
  - ☐ A sample of the zip file: 1234567.zip 1234567
    - Makefile (if needed)

    - L ...
    - (Please have a studentID folder in your zip)
  - ☐ If files are not in a directory after unzip, 10 points will be deducted.

# Questions?

#### Useful Info

- command
  - □ checksec
  - □ readelf
- pwntools
  - □ connect to server and control what content will be sent to it
  - ☐ generate shellcode, attach gdb ...etc.
- gdb
  - □ normal plugins: pwngdb / gef /peda
  - dynamic analysis of the program

```
rbp →
func:
         push rbp
         mov rbp, rsp
                                   Call fun = push next_rip
         sub rsp, 0x30
                                              jmp func
         move eax, 0x0
                                                                   rsp \rightarrow
         leave
         ret
main:
         call func
rip \rightarrow
         mov eax, 0x0 // address 0x4005a0
          ...
```

high address

Stack frame of main

rbp → func: push rbp mov rbp, rsp Call fun = push next\_rip sub rsp, 0x30 jmp func move eax, 0x0 leave rsp  $\rightarrow$ ret main: call func rip  $\rightarrow$ mov eax, 0x0 // address 0x4005a0 ...

high address Stack frame of main 0x4005a0 (return address)

rbp → func: rip 👈 push rbp mov rbp, rsp sub rsp, 0x30 move eax, 0x0 leave rsp  $\rightarrow$ ret main: call func mov eax, 0x0 // address 0x4005a0 ...

high address

0x4005a0 (return address)

Stack frame of main

rbp → func: push rbp mov rbp, rsp rip 👈 sub rsp, 0x30 move eax, 0x0 leave ret rsp → main: call func mov eax, 0x0 // address 0x4005a0 ...

high address Stack frame of main 0x4005a0 (return address) old rbp

```
func:
          push rbp
         mov rbp, rsp
         sub rsp, 0x30
rip 👈
          move eax, 0x0
          leave
         ret
                                                           rbp \rightarrow rsp \rightarrow
main:
         call func
          mov eax, 0x0 // address 0x4005a0
          ...
```

high address

Stack frame of main

0x4005a0 (return address)

old rbp

```
func:
        push rbp
        mov rbp, rsp
        sub rsp, 0x30
rip 👈
        move eax, 0x0
        leave
        ret
                                                            rbp →
main:
        call func
        mov eax, 0x0 // address 0x4005a0
                                                            rsp ->
         ...
```

high address Stack frame of main 0x4005a0 (return address) old rbp Local variables of func()

...

#### Example: Stack frame during a function call

```
func:
        push rbp
                                leave = mov rsp, rbp
        mov rbp, rsp
                                       pop rbp
        sub rsp, 0x30
        move eax, 0x0
rip 👈
        leave
        ret
main:
        call func
        mov eax, 0x0 // address 0x4005a0
```

high address Stack frame of main 0x4005a0 (return address) old rbp Local variables of func()

rbp →

rsp -

...

#### Example: Stack frame during a function call

```
func:
         push rbp
                                   leave = mov rsp, rbp
         mov rbp, rsp
                                           pop rbp
         sub rsp, 0x30
         move eax, 0x0
rip 👈
         leave
         ret
                                                           rbp \rightarrow rsp \rightarrow
main:
         call func
         mov eax, 0x0 // address 0x4005a0
```

# high address Stack frame of main

0x4005a0 (return address)

old rbp

Local variables of func()

```
rbp →
func:
        push rbp
        mov rbp, rsp
        sub rsp, 0x30
                                   ret = pop rip
        move eax, 0x0
        leave
                                                            rsp →
rip 👈
        ret
main:
        call func
        mov eax, 0x0 // address 0x4005a0
         ...
```

high address Stack frame of main 0x4005a0 (return address) old rbp Local variables of func()

rbp → func: push rbp mov rbp, rsp sub rsp, 0x30 move eax, 0x0 rsp → leave ret main: call func mov eax, 0x0 // address 0x4005a0 rip 💙 ...

high address

Stack frame of main

0x4005a0 (return address)

old rbp

Local variables of func()

# Common Security Protection in Binary

- Canary
  - Put canary value before old rbp and return address

- PIE/ALSR
  - ☐ Randomize the address space of a process
    - The offset between different symbol still the same!!

# **Common Security Protection in Binary**

- Relro
  - □ Lazy binding option for program
    - Full Relro will have GOT table read only before calling main function
    - Partial Relro make GOT table writable and resolve symbol after calling main function

- NX
  - Making stack not executable.
    - Make shellcode not able to run on stack.

# Common Security Protection in Binary

- You may check the protection mechanism in binary using checksec
  - □ slimm609/checksec: Checksec

```
(env) huroy@build-server:~/csc2025-project4/hard_rop$ checksec --file hard rop
  '/home/huroy/csc2025-project4/hard_rop/hard_rop'
   Arch:
               amd64-64-little
               Full RELRO
   RELRO:
   Stack:
              Canary found
              NX enabled
   NX:
              PIE enabled
  PIE:
               Enabled
   SHSTK:
               Enabled
   IBT:
   Stripped:
               No
   Debuginfo: Yes
```

#### Example: Stack frame with canary

#### func:

push rbp mov rbp, rsp sub rsp, 0x30 rax,QWORD PTR fs:0x28 QWORD PTR [rbp-0x8],rax



rax,QWORD PTR [rbp-0x8] rax,QWORD PTR fs:0x28 call <\_\_stack\_chk\_fail@plt> leave

ret

rbp - 8 → **→** 

rbp

rsp

high address

Stack frame of main

0x4005a0 (return address)

old rbp

Canary, first bytes is \x00

Local variables of func()