

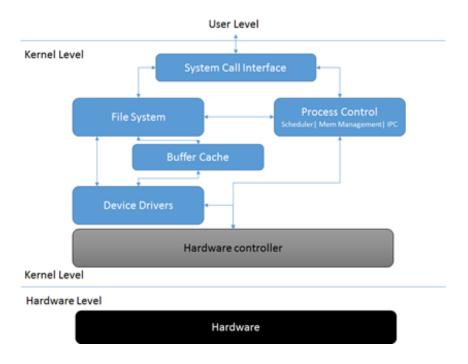
MODULE 2: UNIX Files and File System





Unix File System Architecture







The above figure shows the relationships between the major file-system related components in both user space and the kernel.



- User space contains the applications (for this example, the user of the file system) and the GNU C Library (glibc), which provides the user interface for the file system calls (open, read, write, close). The system call interface acts as a switch, funneling system calls from user space to the appropriate endpoints in kernel space.
- Everything below the system call interface and above the hardware is the Kernel. The kernel provides the file system, the process control subsystem, the memory management and the operating system functions over system calls.
- System calls for file manipulation are for example: creat, open, read, write, close, link, unlink and trunc.



There are three main categories of I/O in UNIX: Block devices, character device, and sockets for network communication.

- Block devices: This category includes disks and tape. A block device is essential to isolate disk details from the rest of the kernel. Block devices are accessible thru the system file (example: /dev/sda the first scsi drive on the first scsi bus)
- Character devices: this category includes printers, terminals (example: /dev/null a bit bucket or bottomless sink for data)







Sockets: This category of device helps connecting to databases, and serving web pages as well as network utilities such as rlogin for remote login, http for hypertext transfer protocol and ftp for file transfer.

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• Almost all the files under /dev are device files. Whereas reading and writing to a regular file stores data on a disk or other filesystem, accessing a device file communicates with a driver in the kernel, which generally in turn communicates with a piece of hardware (a hardware device, hence the name).





```
ls -la /dev | more
total 4
drwxr-xr-x 23 root root
                                   4520 Oct 13 14:20 .
                                   4096 Oct 1 10:30 ..
drwxr-xr-x 33 root root
                               10, 54 Oct 13 14:20 acpi thermal rel
crw-----
            1 root root
                                    235 Oct 13 14:20 autofs
            1 root root
drwxr-xr-x
            2 root
                   root
                                     320 Oct 13 14:20 block
drwxr-xr-x 2 root root
                                     80 Oct 13 14:20 bsg
            1 root disk
                               10,
                                    234 Oct 13 14:20 btrfs-control
crw-rw----
                                     60 Oct 13 14:20 bus
drwxr-xr-x
            3 root root
lrwxrwxrwx
            1 root root
                                      3 Oct 13 14:20 cdrom -> sr0
                                      3 Oct 13 14:20 cdrw -> sr0
lrwxrwxrwx
            1 root root
drwxr-xr-x
            2 root root
                                   4520 Oct 16 10:23 char
            1 root root
                                      1 Oct 13 14:21 console
crw----
lrwxrwxrwx
            1 root root
                                     11 Oct 13 14:20 core -> /proc/kcore
                                     60 Oct 13 14:20 cpu
drwxr-xr-x
            2 root root
                               10,
                                     59 Oct 13 14:20 cpu dma latency
            1 root root
crw-----
                               10.
                                    203 Oct 13 14:20 cuse
            1 root root
crw----
drwxr-xr-x
            8 root root
                                    160 Oct 13 14:20 disk
```









Basics of Files

- From a user perspective in a Unix system, everything is treated as a file. Even such devices such as printers and disk drives.
- All files in the Unix file system can be loosely categorized into 3 types, specifically:
 - Ordinary files
 - Directory files
 - Device files





- Text file or Binary file

Both can be executable files (.sh) with suitable permissions. Ex: scripts, documentation, GIF, JPEG, etc.

- How to create? using editors (vi editor)
- How to delete? using rm command

Directory Files

- Can contain any kind of files

What is "." and ".."??

- How to create? using mkdir command (mkdir /tmp/demofolder)
- How to delete? using rmdir command (rmdir /tmp/demofolder)







Unix File System

- The Unix file system is a methodology for logically organizing and storing large quantities of data such that the system is easy to manage.
- A file can be informally defined as a collection of (typically related) data, which can be logically viewed as a stream of bytes (i.e. characters).
- A file is the smallest unit of storage in the Unix file system.





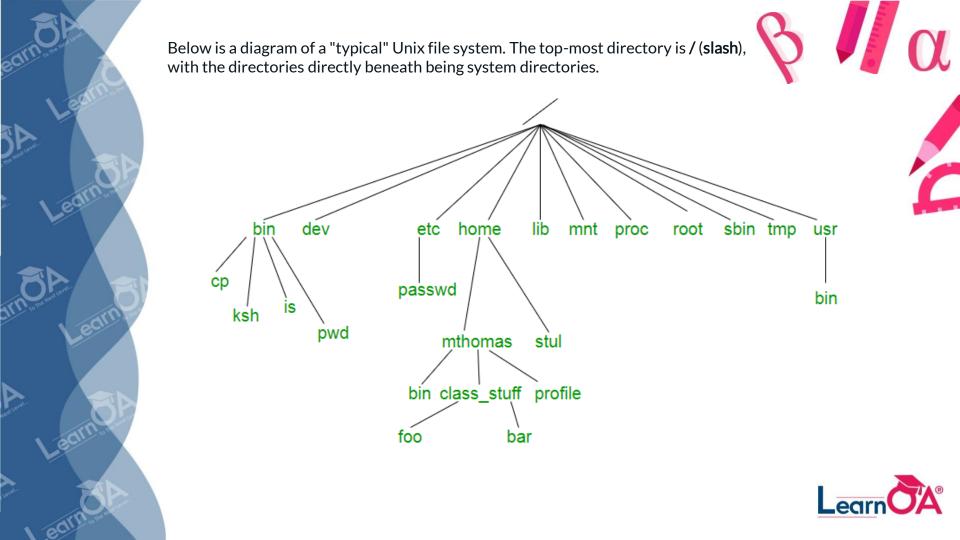
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By contrast, a file system consists of files, relationships to other files, as well as the attributes of each file. File attributes for a generic operating system might include (but are not limited to):

- a file type (i.e. what kind of data is in the file)
- a file name (which may or may not include an extension)
- a physical file size
- a file owner
- file protection/privacy capability
- file time stamp (time and date created/modified)

The Unix file system has a hierarchical (or tree-like) structure with its highest level directory called root (denoted by /, pronounced slash).







Following system files (i.e. directories) are present in most Unix filesystems:

- bin short for binaries, is the directory where many commonly used executable commands reside
- dev contains device specific files
- etc contains system configuration files
- home contains user directories and files
- lib contains all library files
- mnt contains device files related to mounted devices







- proc contains files related to system processes
- root the root user's' home directory (note this is different than /)
- sbin system binary files reside here. If there is no sbin directory on your system, these files most likely reside in etc
- tmp storage for temporary files which are periodically removed from the filesystem
- usr also contains executable commands





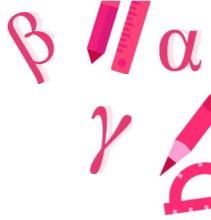


- We store your information in a file, and the operating system stores the information about a file in an inode, sometimes called as inode number.
- 'An inode is metadata of the data'.
- Whenever a user or a program needs access to a file, the operating system first searches for the exact and unique inode (inode number), in a table called as an inode table. In fact the program or the user who needs access to a file, reaches the file with the help of the inode number found from the inode table.





Inode Structure(contd.)



. (DOT)	9700524
(DOT DOT)	9699798
file1.txt	9700532
file2.txt	9700533

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file & folder names

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Corresponding inodes

An inode is a list of information related to a particular object (either a file, a directory or a symbolic link). The inode stores the type of object, permissions, ownership and the location where the data is stored.

Lets see an example directory listing

_ ls -ia 9700524 . 9700532 file1.txt 9700534 file3.txt 9699798 . 9700533 file2.txt 9700535 file4.txt





Each inode, an entry in the inode table, is a data structure which the system uses to store the following information about a file:

- Type of file (ordinary, directory or special file).
- Access permissions for the file owner, the owner's group members and others (i.e. the general public).
- Number of links to the file.
- File owner's user and group IDs
- File size in bytes.
- The disk addresses of the data blocks where the contents of the file are actually stored.
- Time of last access (read or executed), time of last modification (i.e. written) and time which the inode itself was last changed.





Inode Structure(contd.)

File: important/file1.txt

Size: 0 Blocks: 0 IO Block: 4096 regular empty file

Device: 803h/2051d Inode: 9700532 Links: 1

Access: (0644/-rw-r--r-) Uid: (0/ root) Gid: (0/ root)

Access: 2019-10-15 17:09:20.130714453 +0530 Modify: 2019-10-15 17:09:20.130714453 +0530 Change: 2019-10-15 17:09:20.130714453 +0530

Birth: -





Links in Unix

- On the UNIX command line, the tool In abbreviates the term link.
- It allows you to create an additional reference to a file, or directory.
- It does that by adding an additional name of an entry in the file allocation table of the file system.
- It is an essential tool in unix and it is part of the package coreutils.

```
_ dpkg -S `which ln`
coreutils: /bin/ln
```







Links on unix file system are of two types:

Hardlink

- Linking files by reference
- System maintains a count of the number of links
- Does not work across file system

Softlink or Symbolic link

- Linking files by name
- No counter is maintained
- Work across file system







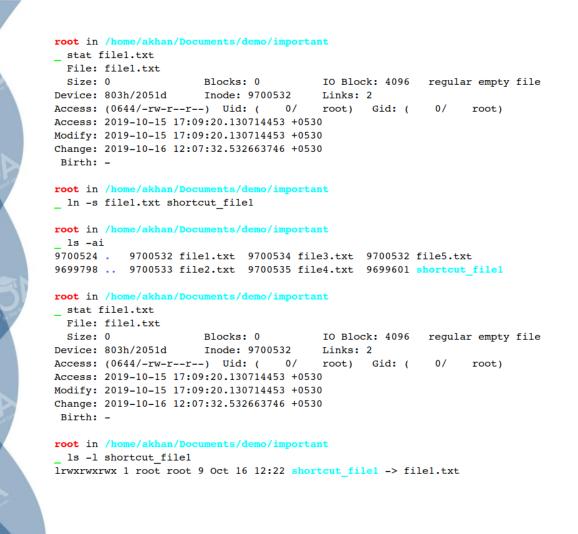
```
root in /home/akhan/Documents/demo/important
file1.txt file2.txt file3.txt file4.txt
root in /home/akhan/Documents/demo/important
 stat file1.txt
  File: file1.txt
 Size: 0
                       Blocks: 0
                                          IO Block: 4096
                                                           regular empty file
Device: 803h/2051d
                       Inode: 9700532
                                          Links: 1
Access: (0644/-rw-r--r--) Uid: ( 0/
                                          root)
                                                  Gid: (
                                                                  root)
Access: 2019-10-15 17:09:20.130714453 +0530
Modify: 2019-10-15 17:09:20.130714453 +0530
Change: 2019-10-16 12:07:14.200475126 +0530
Birth: -
root in /home/akhan/Documents/demo/important
ln file1.txt file5.txt
root in /home/akhan/Documents/demo/important
ls -ai
9700524 .
           9700532 file1.txt 9700534 file3.txt 9700532 file5.txt
9699798 .. 9700533 file2.txt 9700535 file4.txt
root in /home/akhan/Documents/demo/important
 stat file1.txt
 File: file1.txt
                                                           regular empty file
  Size: 0
                       Blocks: 0
                                          IO Block: 4096
Device: 803h/2051d
                       Inode: 9700532
                                          Links: 2
Access: (0644/-rw-r--r--) Uid: (
                                          root) Gid: (
                                                                  root)
Access: 2019-10-15 17:09:20.130714453 +0530
Modify: 2019-10-15 17:09:20.130714453 +0530
Change: 2019-10-16 12:07:32.532663746 +0530
 Birth: -
```

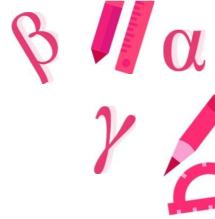


A hard link is a UNIX path name for a file.

It is created with In command.







Symbolic link is also a means of referencing a file.

It is created with In -s command.





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Hardlink	Softlink
Does not create a new inode	Create a new inode
Cannot link directories unless it is done by root	Can link directories
Cannot link files across file systems	Can link files across file system
Increase hard link count of the linked inode	Does not change hard link count of the linked inode





Following are commonly used file related commands in linux:

- cat display entire file(s) in the terminal window
- cd change the current directory
- cp copy file(s)
- chmod change file(s) protection modes
- lpr send file(s) to the line printer







- mkdir create a new directory
- more nicely view the contents of file(s)
- less opposite of more
- mv move (and/or rename) file(s)
- pwd print absolute pathname of current working directory
- rm remove file(s) rmdir remove an empty directory

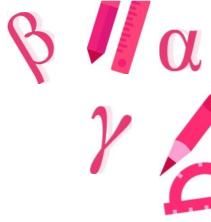






Comparing files in Unix

```
root in /home/akhan/Documents/demo/important
   whatis cmp
                     - compare two files byte by byte
cmp (1)
root in /home/akhan/Documents/demo/important
 cmp file1.txt file2.txt
file1.txt file2.txt differ: byte 16, line 1
root in /home/akhan/Documents/demo/important
 whatis comm
                     - compare two sorted files line by line
comm (1)
root in /home/akhan/Documents/demo/important
 comm file1.txt file2.txt
hello from file1
        hello from file2
root in /home/akhan/Documents/demo/important
  whatis diff
                  - compare files line by line
diff (1)
root in /home/akhan/Documents/demo/important
diff file1.txt file2.txt
1c1
< hello from file1
> hello from file2
```



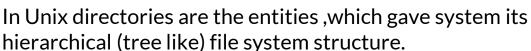
To get detailed information on the use of the command use man pages.

man cmp

man is the system's manual pager.







Directory file contains some sequential entries. 3 things are important:

- 1. **Byte offset:** This is the number which gives location of files inode in that directory.
- 2. **Inode number:** Inodes entry is made in the directory file along with its name.
- 3. **File names:** These are the names of files located under present directory.







and

Directories in Unix(contd.)





The first two file names in every directory are (".") and ("..")

- where "." Represents present or current directory
- ".." Represents parent directory file.

Each entry in a directory file is of 16-byte size

- where 14 bytes for 14 characters in file name
- and 2 bytes for inode number





Access permissions to the users for directory file:

- **1. READ -** Allow the user to see contents i.e. files in that directory
- **2. WRITE** Allows the user to create or delete subdirectories and files in that directory
- **3. EXECUTE** Allows the user to search the directory i.e allow the user to open subdirectories in that directory

For kernel writing permission for directory means, to actually write in directory file.







Mada		_	_	E11				
Mode	Links	Owner	Group	File si	ze	mod	lified	Filename
\bigcup		\prod	\int	\bigcap		ĺ		
dr-xr-xr-x	387	root	root	0	Oct	13	14:20	proc
drwxr-xr-x	48	akhan	akhan	4096	Oct	16	12:19	root
drwxr-xr-x	45	root	root	1420	Oct	16	10:23	run
drwxr-xr-x	2	root	root	12288	Oct	11	12:47	sbin
drwxr-xr-x	5	root	root	4096	Jul	24	12:47	snap
drwxr-xr-x	2	root	root	4096	Jun	24	2016	srv
drwxr-xr-x	5	root	root	4096	May	16	13:47	statping
dr-xr-xr-x	13	root	root	0	Oct	16	13:20	sys
drwxrwxrwt	: 10	root	root	4096	Oct	16	13:24	tmp
-rw-rr	- 1	root	root	0	Oct	16	13:24	tmp_information
drwxr-xr-x	13	root	root	4096	Jul	26	10:11	usr
drwxr-xr-x	16	root	root	4096	Jul	2	13:11	var
lrwxrwxrwx	1	root	root	30	Oct	1	10:30	<pre>vmlinuz -> boot/vmlinuz-4.15.0-65-generic</pre>
lrwxrwxrwx	1	root	root	30	Oct	1	10:30	<pre>vmlinuz.old -> boot/vmlinuz-4.15.0-64-generic</pre>

Last



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Following are the attributes associated with files and directories:

- file name: the name associated with the file (recall, this can be any type of file)
- modification date: the date the file was last modified, i.e. a "timestamp". If the file has not been modified within the last year (or six months for Linux), the year of last modification is displayed.
- size: the size of the file in bytes (i.e. characters)
- group: associated group for the file







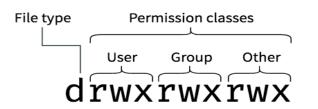


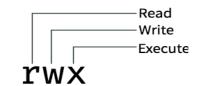
- owner: the owner of the file
- number of links: the number of other links associated with this file
- permission modes: the permissions assigned to the file for the owner, the group and all others

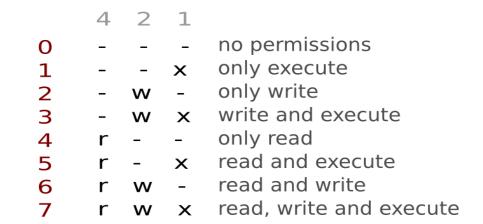


Files and Directories Permissions













- 9 permission bits used to determine 3 types of accesses: READ, WRITE, EXECUTE.
- Permission can be set based on GROUP, OWNER, ANYONE ELSE
- Use chmod command to change permission
 - Binary 001 for EXECUTE
 - Binary 010 for WRITE
 - Binary 100 for READ







Files and Directories Permissions (contd.)

Examples:

```
chmod 777 /tmp/dir
chmod u+rwx,g+rwx,o+rwx /tmp/dir
chmod u+rwx,o+rx /tmp/dir
chmod g-rx /tmp/dir
chmod 700 /tmp/dir
```













Linux kernel uses algorithm *namei* (*name* to *i*node) is used for converting a path to an inode number. The kernel parses the path by accessing each inode in the path and finally returning the inode of the required file.

- If path name string starts from root then working inode=root node otherwise working inode=current directories inode
- While there are more components in path string read next component from input parameter verify,
 - ✓ whether working inode is of directory
 - ✓ whether the access permissions for it are ok











- If working inode is root inode and component is "..", then control is send back to beginning of loop, because ".." means parent directory and "/" has no parent directory otherwise.
- Read the directory file (using working inode) by repetitive use of bmap,bread,brealease algorithms.
- If component matches any entry in directory file get inode number for matched component release working inode now working inode=inode of matched component and go to while loop again else means if component doesn't match return no inode is there
- Return working inode



```
Algorithm: namei
Input: pathname
Output: locked inode
    if (path name starts from root)
            working inode = root inode (algorithm: iget);
    else
            working inode = current directory inode (algorithm: iget);
    while (there is more path name)
            read next path name component from input;
            verify that working inode is of a directory and access permissions are OK;
            if (working inode is of root and component is "..")
                    continue;
            read directory (working inode) by repeated use of algorithms: bmap, bread, brelse;
            if (component matches an entry in the directory (working inode)
                    get inode number for matched component;
                    release working inode (algorithm: iput);
                    working inode = inode of matched component (algorithm: iget);
            else
                    return (no inode) // component not in the directory
    return (working inode);
```







- Consider, entered path is "/etc/passwd"
- Kernel starts parsing pathname
- Working inode=root
- Then, search root for the directory /etc
- Kernel access data in root block by block to find etc
- Now working inode=etc (iget)
- Search etc block by block to find passwd
- Release an inode of etc (iput), set working inode=passwd (iget)
- Since no more component in pathname
- Return working inode







Open Files and File Descriptor in Unix

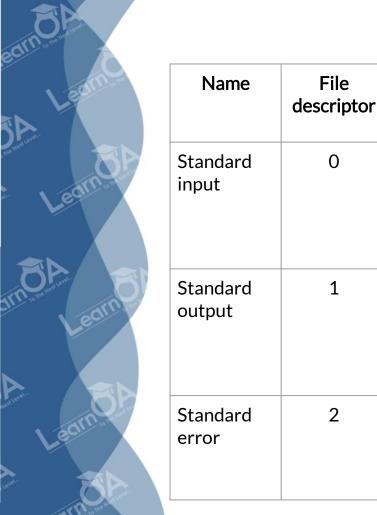
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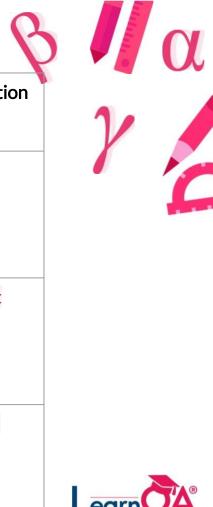


- A file descriptor is a number that uniquely identifies an open file in a operating system. It describes a data resource, and how that resource may be accessed.
- When the UNIX kernel starts any process for example, grep, ls, or a shell - it sets up several places for that process to read from and write to. These places are called "open files".
- The kernel gives each file a number called a file descriptor.
- On a Unix-like operating system, the first three file descriptors, by default, are STDIN (<u>standard input</u>), STDOUT (standard output), and STDERR (standard error).

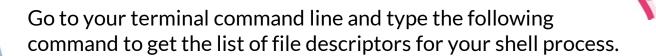




Name	File descriptor	Description	Abbreviation
Standard input	0	The default data stream for input, for example in a command pipeline. In the terminal, this defaults to keyboard input from the user.	stdin
Standard output	1	The default data stream for output, for example when a command prints text. In the terminal, this defaults to the user's screen.	stdout
Standard error	2	The default data stream for output that relates to an error occurring. In the terminal, this defaults to the user's screen.	stderr







```
_ ls -l /proc/$$/fd
total 0
lrwx----- 1 root root 64 Oct 16 19:42 0 -> /dev/pts/4
lrwx----- 1 root root 64 Oct 16 19:42 1 -> /dev/pts/4
lrwx----- 1 root root 64 Oct 16 19:42 2 -> /dev/pts/4
lrwx----- 1 root root 64 Oct 16 19:42 2 -> /dev/pts/4
```

What is file descriptor 255 for?

As an alternative connection to the tty in case fd 1 (/dev/stdout) and fd 0 (/dev/stdin) get blocked.

\$\$ is the process ID of the current shell instance.

```
_ echo $$
25306
```





Other Types of File Systems

Some of the other common file system types are:

- FAT
- FAT32
- FAT64
- NTFS







- FAT stands for "File Allocation Table".
- The file allocation table is used by the operating system to locate files on a disk.
- A file may be divided into many sections and scattered around the disk due to fragmentation.
- FAT keeps track of all pieces of a file.
- FAT provides quick access to files.
- The speed of file access depends on file type, file size, partition size, fragmentation and number of files in a folder.







FAT32 File System

- FAT32 is an advanced version of FAT file system.
- It can be used on drives from 512 MB to 2TB in size.
- One of the most important features of FAT and FAT32 is that they offer compatibility with operating systems other than Windows 2000 also.
- FAT32 provides good file access in partition sizes less than 500 MB or greater than 2 GB.
- FAT32 provides better disk space utilization.





- NTFS stands for "New Technology File System".
- Windows 2000 professional fully supports NTFS.
- NTFS is highly reliable.
- It is recoverable file system.
- It uses transaction logs to update the file and folders logs automatically.
- The system also has a great amount of fault tolerance. It means that if transaction fails due to power or system failure, the logged transactions are used to recover the data.









Thank You!

