

# Computer Architecture Assignment

Fast Adders  
(Carry look Ahead Adder)

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## Carry look Ahead Adder:

Different types of Digital systems are constructed from very few types of basic network configurations such as AND gate, NAND gate, Or gate, etc... These elementary circuits are used over and over again in various topological combinations. In addition to performing logic, digital systems must also store binary numbers. For these memory cells, also known as FLIP-FLOP's are designed. To perform some functions such as binary addition. Hence, to perform such functions, combinations of logic gates and FLIP-FLOPs are designed over a single-chip IC. These IC's form the practical building blocks of the Digital systems. One of such building blocks used for binary addition is the Carry Look-ahead Adder.

### What is a Carry Look-ahead Adder?

A digital computer must contain circuits which can perform arithmetic operations such as addition, subtraction, multiplication, and division. Among these, addition and subtraction are the basic operations whereas multiplication and division are the repeated addition and subtraction respectively.

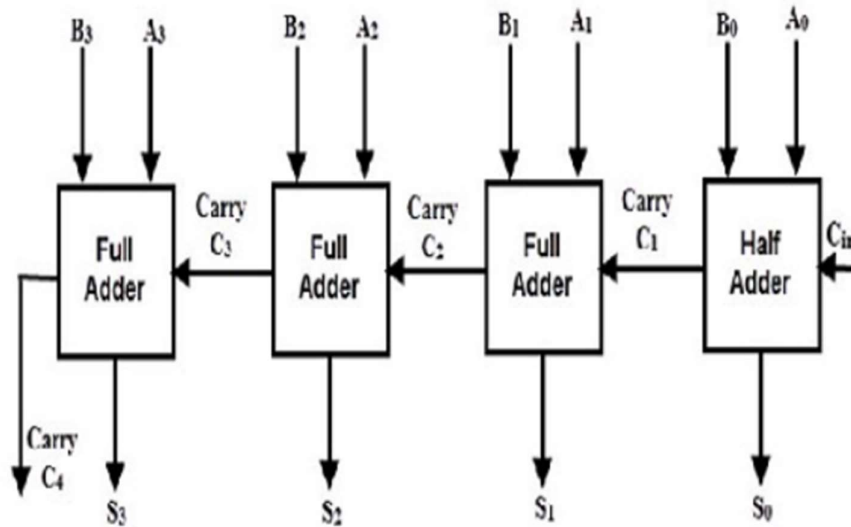
To perform these operations 'Adder circuits' are implemented using basic logic gates. Adder circuits are evolved as Half-adder, Full-adder, Ripple-carry Adder, and Carry Look-ahead Adder.

Among these Carry Look-ahead Adder is the faster adder circuit. It reduces the propagation delay, which occurs during addition, by using more complex hardware circuitry. It is designed by transforming the ripple-carry Adder circuit such that the carry logic of the adder is changed into two-level logic.

## 4-Bit Carry Look-ahead Adder

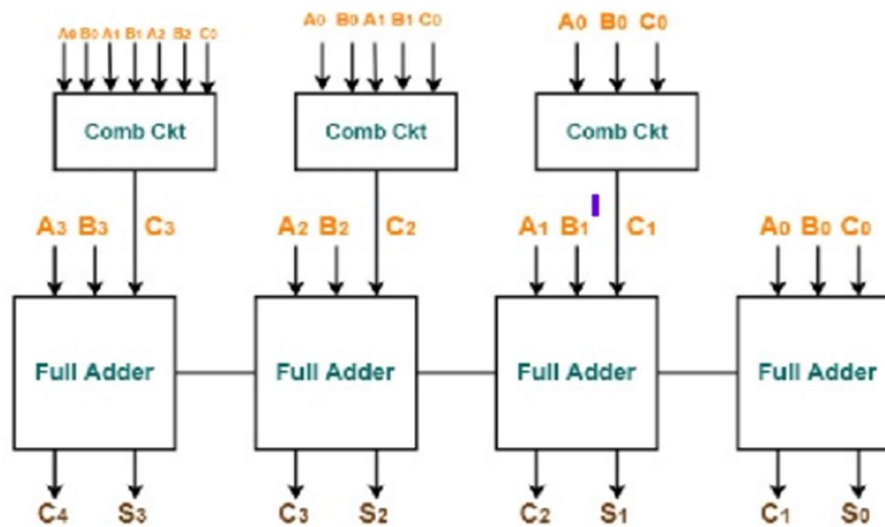
In parallel adders, carry output of each full adder is given as a carry input to the next higher-order state. Hence, these adders it is not possible to produce carry and sum outputs of any state unless a carry input is available for that state.

So, for computation to occur, the circuit has to wait until the carry bit propagated to all states. This induces carry propagation delay in the circuit.



*4-bit-Ripple-Carry-Adder*

The propagation delay of the adder is calculated as “the propagation delay of each gate times the number of stages in the circuit”. For the computation of a large number of bits, more stages have to be added, which makes the delay much worse. Hence, to solve this situation, Carry Look-ahead Adder was introduced.



4-bit-Carry-Look-ahead-Adder-Logic-Diagram

In this adder, the carry input at any stage of the adder is independent of the carry bits generated at the independent stages. Here the output of any stage is dependent only on the bits which are added in the previous stages and the carry input provided at the beginning stage. Hence, the circuit at any stage does not have to wait for the generation of carry-bit from the previous stage and carry bit can be evaluated at any instant of time.

## Truth Table of Carry Look-ahead Adder

For deriving the truth table of this adder, two new terms are introduced – Carry generate and carry propagate. Carry generate  $G_i = 1$  whenever there is a carry  $C_{i+1}$  generated. It depends on  $A_i$  and  $B_i$  inputs.  $G_i$  is 1 when both  $A_i$  and  $B_i$  are 1. Hence,  $G_i$  is calculated as  $G_i = A_i \cdot B_i$ .

Carry propagated  $P_i$  is associated with the propagation of carry from  $C_i$  to  $C_{i+1}$ . It is calculated as  $P_i = A_i \oplus B_i$ . The truth table

of this adder can be derived from modifying the truth table of a full adder.

Using the  $G_i$  and  $P_i$  terms the Sum  $S_i$  and Carry  $C_{i+1}$  are given as below –

- $S_i = P_i \oplus G_i$ .
- $C_{i+1} = C_i.P_i + G_i$ .

Therefore, the carry bits  $C_1$ ,  $C_2$ ,  $C_3$ , and  $C_4$  can be calculated as

- $C_1 = C_0.P_0 + G_0$ .
- $C_2 = C_1.P_1 + G_1 = (C_0.P_0 + G_0).P_1 + G_1$ .
- $C_3 = C_2.P_2 + G_2 = (C_1.P_1 + G_1).P_2 + G_2$ .
- $C_4 = C_3.P_3 + G_3 = C_0.P_0.P_1.P_2.P_3 + P_3.P_2.P_1.G_0 + P_3.P_2.G_1 + G_2.P_3 + G_3$ .

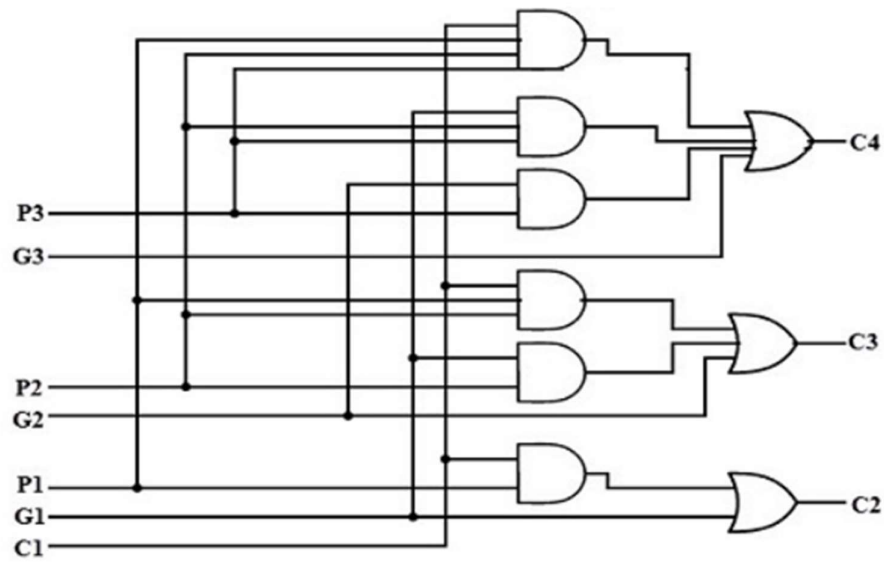
It can be observed from the equations that carry  $C_{i+1}$  only depends on the carry  $C_0$ , not on the intermediate carry bits.

A	B	$C_i$	$C_{i+1}$	Condition
0	0	0	0	No carry generate
0	0	1	0	
0	1	0	0	
0	1	1	1	No carry propagate
1	0	0	0	
1	0	1	1	
1	1	0	1	Carry generate
1	1	1	1	

Carry-Look-ahead-Adder-Truth-Table

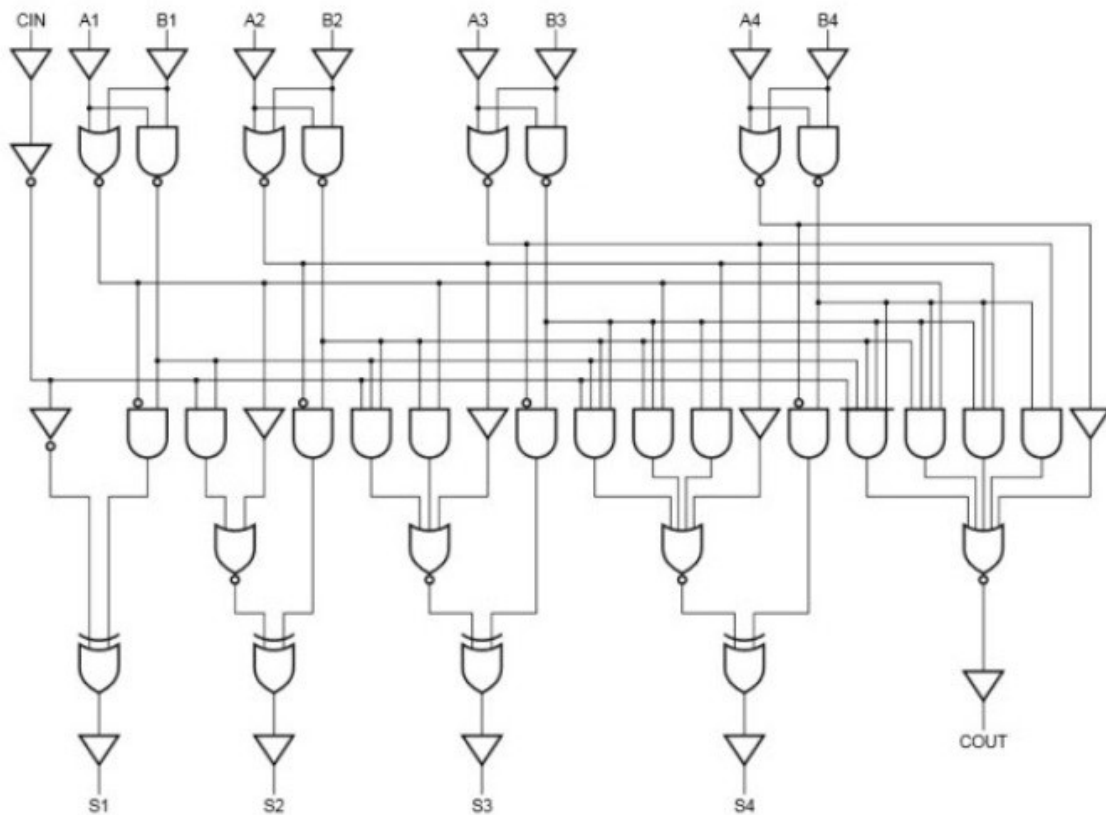
## Circuit Diagram

The above equations are implemented using two-level combinational circuits along with AND, OR gates, where gates are assumed to have multiple inputs.



*Carry-Output-Generation-Circuit-of-Carry-Look-ahead-Adder*

The Carry Look-ahead Adder circuit for 4-bit is given below.



8-bit and 16-bit Carry Look-ahead Adder circuits can be designed by cascading the 4-bit adder circuit with carry logic.

## *C program to implement Carry look Ahead Adder*

```
/* C Program For Implementation Of Look Ahead Carry Adder */
```

```
#include <stdio.h>
```

```
#include <math.h>
```

```
#include <stdlib.h>
```

```
#include <string.h>
```

```
int get1(int a)
```

```
{
```

```
char ch='B';
```

```
if(a==1)
```

```
ch='A';
```

```
do
```

```
{
```

```
printf("\n\tENTER VALUE OF %c:",ch);
```

```
scanf("%d",&a);
```

```
if(a<=0)
```

```
printf("\n\t!!INVALID NUMBER.ENTER VALUE (0< A)!");
```

```
}while(a<=0);
```

```
return(a);
```

```
}
```

```
int and(int a,int b)
```

```
{  
  
int c;  
  
if(a< b)  
  
c=a;  
  
else  
  
c=b;  
  
return (c);  
  
}
```

```
int or(int a,int b)
```

```
{  
  
int x;  
  
if(a>b)  
  
x=a;  
  
else  
  
x=b;  
  
return x;  
  
}
```

```
int exor(int a,int b)
```

```
{  
  
int x;
```



```
if(a==b)
```

```
x=0;
```

```
else
```

```
x=1;
```

```
return x;
```

```
}
```

```
void add()
```

```
{
```

```
int i=7,A,B,a,b,cin,num;
```

```
int n1[8],n2[8],cg[8],cp[8],sum[8];
```

```
for(i=0;i<=7;i++)
```

```
{
```

```
n1[i]=0; // Num 1
```

```
n2[i]=0; // Num 2
```

```
cg[i]=0; // Gi
```

```
cp[i]=0; // Pi
```

```
sum[i]=0; // Sum
```

```
}
```

```
A = a = get1(1);
```

```
B = b = get1(0);
```

```
i=7;
```

```

do
{
n1[i]=a%2;

a=a/2;

n2[i]=b%2;

b=b/2;

i--;

}while((a!=0)|| (b!=0));

i=0;

printf("\n\t\t Binary Form");

printf("\n\t A = %d : ",A);

for(i=0;i<=7;i++)

printf("%d ",n1[i]);

printf("\n\t B = %d : ",B);

for(i=0;i<=7;i++)

printf("%d ",n2[i]);

cin=0;

for(i=7;i>=0;i--)

{

sum[i]=exor(cin,exor(n1[i],n2[i])); // Sum Pi (+) Bi

```

```

cg[i]=and(n1[i],n2[i]); // Gi = Ai . Bi

cp[i]=or(n1[i],n2[i]); // Pi = Ai (+) Bi

cin=or(cg[i],and(cp[i],cin)); // Cin =Gi + PiCi
}

printf("\n\n\t\t SUM: ");

num=0;

for(i=0;i<=7;i++)

{

printf(" %d",sum[i]);

num=num + (sum[i]*pow(2,7-i));

}

printf("\n\n\t\t SUM: %d + %d= %d\n",A,B,num);

printf("\t\t The Carry Is : %d\n\n",cin);

}

void main()

{

int ch,a,b,c,d;

while(1)

{

M: printf("***** MENU FOR LOOK AHEAD CARRY ADDER *****");

printf("\n\t\t1.ADDITION OF TWO NUMBER");

```

```
printf("\n\t\t2.EXIT\n");

printf("*****");

printf("\n\t\tEnter Your Option:");

scanf("%d",&ch);

switch(ch)

{

case 1:

add();

break;

case 2:

exit(0);

break;

default:

printf("ERROR!!!!!!!!!! INVALID ENTRY...\n");

printf("Back To Main Menu\n\n");

goto M;

}

}

}
```

# Output:

```
Output Clear
/tmp/NpTF3xcCU8.o
***** MENU FOR LOOK AHEAD CARRY ADDER *****
1.ADDITION OF TWO NUMBER
2.EXIT
*****
Enter Your Option:1
ENTER VALUE OF A:12
ENTER VALUE OF B:5
Binary Form
A = 12 : 0 0 0 0 1 1 0 0
B = 5 : 0 0 0 0 0 1 0 1

SUM: 0 0 0 1 0 0 0 1

SUM: 12 + 5= 17
The Carry Is : 0

***** MENU FOR LOOK AHEAD CARRY ADDER *****
1.ADDITION OF TWO NUMBER
2.EXIT
*****
Enter Your Option:1
ENTER VALUE OF A:3
ENTER VALUE OF B:-2
!INVALID NUMBER. ENTER VALUE (0< A)!
ENTER VALUE OF B:5
```

## Output

Clear

```
***** MENU FOR LOOK AHEAD CARRY ADDER *****
1.ADDITION OF TWO NUMBER
2.EXIT
*****
Enter Your Option:1
ENTER VALUE OF A:3
ENTER VALUE OF B:-2
!INVALID NUMBER.ENTER VALUE (0< A)!
ENTER VALUE OF B:5
Binary Form
A = 3 : 0 0 0 0 0 1 1
B = 5 : 0 0 0 0 0 1 0 1

SUM: 0 0 0 0 1 0 0 0

SUM: 3 + 5= 8
The Carry Is : 0

***** MENU FOR LOOK AHEAD CARRY ADDER *****
1.ADDITION OF TWO NUMBER
2.EXIT
*****
Enter Your Option:1
ENTER VALUE OF A:11
ENTER VALUE OF B:52
Binary Form
```

```
Output Clear
SUM: 0 0 0 0 1 0 0 0

SUM: 3 + 5= 8
The Carry Is : 0

***** MENU FOR LOOK AHEAD CARRY ADDER *****
1.ADDITION OF TWO NUMBER
2.EXIT
*****
Enter Your Option:1
ENTER VALUE OF A:11
ENTER VALUE OF B:52
Binary Form
A = 11 : 0 0 0 0 1 0 1 1
B = 52 : 0 0 1 1 0 1 0 0

SUM: 0 0 1 1 1 1 1 1

SUM: 11 + 52= 63
The Carry Is : 0

***** MENU FOR LOOK AHEAD CARRY ADDER *****
1.ADDITION OF TWO NUMBER
2.EXIT
*****
Enter Your Option:2
```

## Mips Code:

.data

result: .word 0

.text

.globl main

main:

# Prompt user for input

```
li $v0, 4
la $a0, prompt
syscall

# Read the first number

li $v0, 5
syscall
move $t0, $v0

# Prompt user for the second number

li $v0, 4
la $a0, prompt
syscall

# Read the second number

li $v0, 5
syscall
move $t1, $v0
```



# Perform the addition using fast adder algorithm

add \$t2, \$t0, \$t1    #  $t2 = t0 + t1$

addi \$t3, \$t0, 1    #  $t3 = t0 + 1$

addi \$t4, \$t1, -1    #  $t4 = t1 - 1$

and \$t5, \$t3, \$t4    #  $t5 = t3 \& t4$

srl \$t6, \$t5, 1    #  $t6 = t5 \gg 1$

add \$t7, \$t2, \$t6    #  $t7 = t2 + t6$

# Store the result in memory

sw \$t7, result

# Display the result to the user

li \$v0, 4

la \$a0, result\_msg

syscall

lw \$a0, result

li \$v0, 1

syscall

# Exit the program

li \$v0, 10

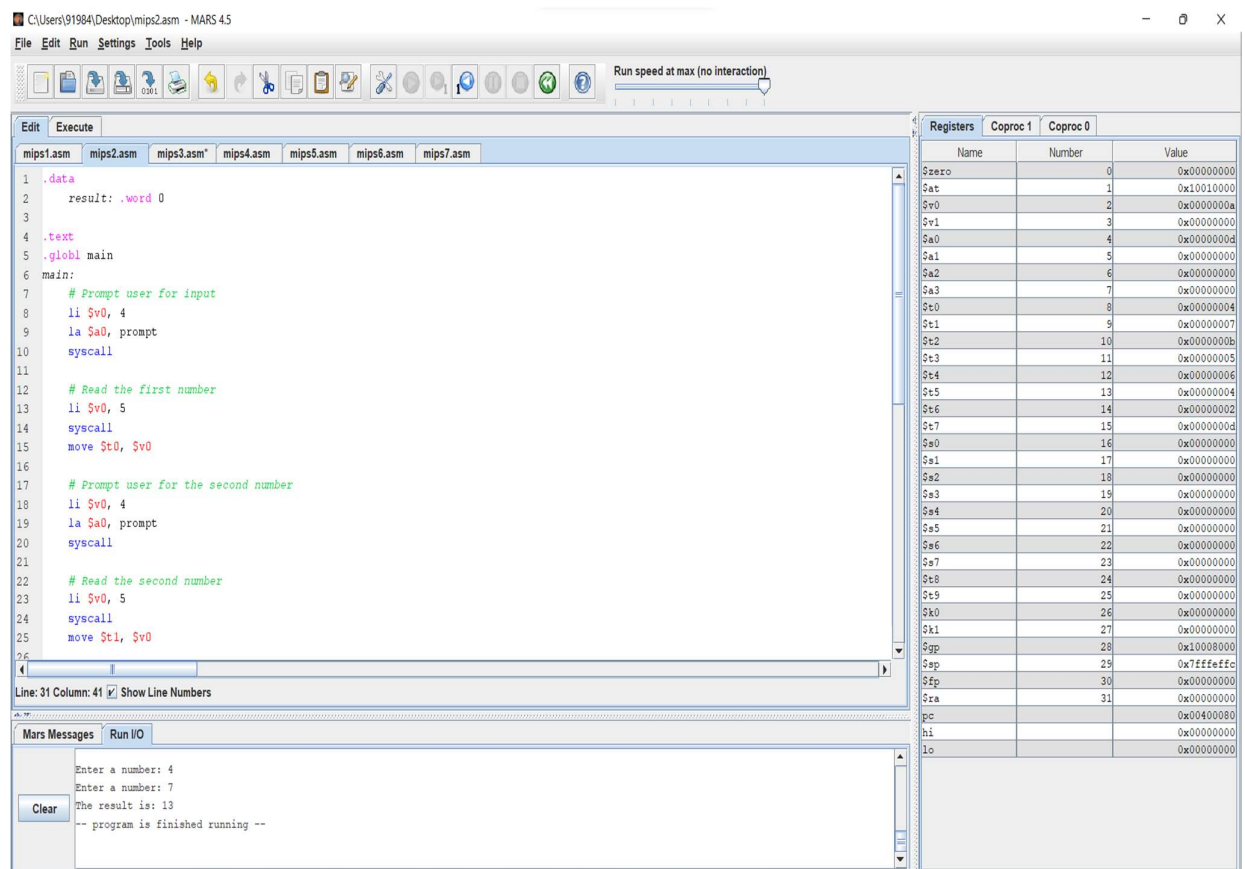
syscall

.data

prompt: .asciiz "Enter a number: "

result\_msg: .asciiz "The result is: "

OUTPUT:



C:\Users\91984\Desktop\mips2.asm - MARS 4.5

File Edit Run Settings Tools Help

Run speed at max (no interaction)

Edit Execute

mips1.asm mips2.asm mips3.asm mips4.asm mips5.asm mips6.asm mips7.asm

```

24 syscall
25 move $t1, $v0
26
27 # Perform the addition using fast adder algorithm
28 add $t2, $t0, $t1 # t2 = t0 + t1
29 addi $t3, $t0, 1 # t3 = t0 + 1
30 addi $t4, $t1, -1 # t4 = t1 - 1
31 and $t5, $t3, $t4 # t5 = t3 & t4
32 srl $t6, $t5, 1 # t6 = t5 >> 1
33 add $t7, $t2, $t6 # t7 = t2 + t6
34
35 # Store the result in memory
36 sw $t7, result
37
38 # Display the result to the user
39 li $v0, 4
40 la $a0, result_msg
41 syscall
42 lw $a0, result
43 li $v0, 1
44 syscall
45
46 # Exit the program
47 li $v0, 10
48 syscall
49
50 .data
51 prompt: .asciiz "Enter a number: "
52 result_msg: .asciiz "The result is: "

```

Line: 52 Column: 38 Show Line Numbers

Mars Messages Run IO

Clear

Enter a number: 4  
Enter a number: 7  
The result is: 12

Registers Coproc 1 Coproc 0

Name	Number	Value
\$zero	0	0x00000000
\$at	1	0x10010000
\$v0	2	0x0000000a
\$v1	3	0x00000000
\$a0	4	0x0000000d
\$a1	5	0x00000000
\$a2	6	0x00000000
\$a3	7	0x00000000
\$t0	8	0x00000004
\$t1	9	0x00000007
\$t2	10	0x0000000b
\$t3	11	0x00000005
\$t4	12	0x0000000e
\$t5	13	0x00000004
\$t6	14	0x00000002
\$t7	15	0x0000000d
\$a0	16	0x00000000
\$a1	17	0x00000000
\$a2	18	0x00000000
\$a3	19	0x00000000
\$a4	20	0x00000000
\$a5	21	0x00000000
\$a6	22	0x00000000
\$a7	23	0x00000000
\$t8	24	0x00000000
\$t9	25	0x00000000
\$k0	26	0x00000000
\$k1	27	0x00000000
\$gp	28	0x10008000
\$sp	29	0x7ffffcfc
\$fp	30	0x00000000
\$ra	31	0x00000000
pc		0x00400080
hi		0x00000000
lo		0x00000000

C:\Users\91984\Desktop\mips2.asm - MARS 4.5

File Edit Run Settings Tools Help

Run speed at max (no interaction)

Edit Execute

Text Segment

Bkpt	Address	Code	Basic	Source
	0x00400000	0x24020004	addiu \$2,\$0,0x00000004	8: li \$v0, 4
	0x00400004	0x3c011001	lui \$1,0x00000001	9: la \$a0, prompt
	0x00400008	0x24240004	ori \$4,\$1,0x00000004	
	0x0040000c	0x00000000	syscall	10: syscall
	0x00400010	0x24020005	addiu \$2,\$0,0x00000005	13: li \$v0, 5
	0x00400014	0x00000000	syscall	14: syscall
	0x00002402	addu \$8,\$0,\$2		15: move \$t0, \$v0
	0x00400018	0x24020004	addiu \$2,\$0,0x00000004	18: li \$v0, 4
	0x0040001c	0x3c011001	lui \$1,0x00000001	19: la \$a0, prompt
	0x00400020	0x34240004	ori \$4,\$1,0x00000004	
	0x00400024	0x00000000	syscall	20: syscall
	0x00400028	0x24020005	addiu \$2,\$0,0x00000005	23: li \$v0, 5

Data Segment

Address	Value (+0)	Value (+4)	Value (+8)	Value (+c)	Value (+10)	Value (+14)	Value (+18)	Value (+1c)
0x10010000	0x00000002	0x65746e45	0x20612072	0x626d756e	0x203a7265	0x56655400	0x73657220	0x20746e75
0x10010020	0x203a7369	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010040	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010060	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010080	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100a0	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100c0	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x100100e0	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010100	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000
0x10010120	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000	0x00000000

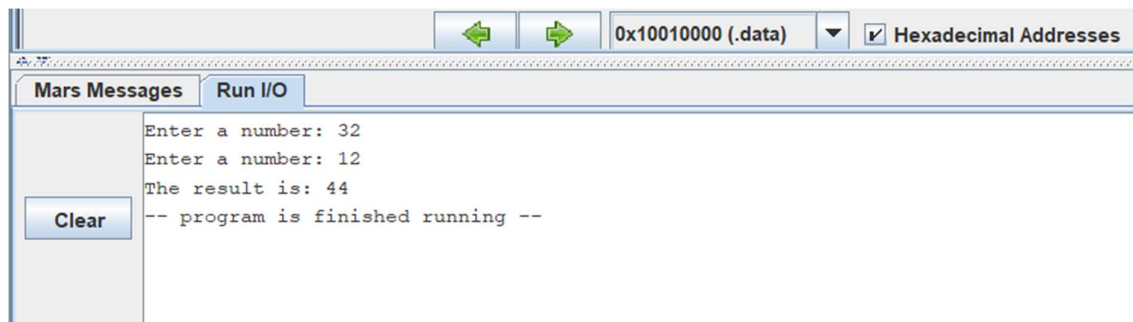
Mars Messages Run IO

Clear

Enter a number: 32  
Enter a number: 12  
The result is: 44  
-- program is finished running --

Registers Coproc 1 Coproc 0

Name	Number	Value
\$zero	0	0x00000000
\$at	1	0x10010000
\$v0	2	0x0000000a
\$v1	3	0x00000000
\$a0	4	0x0000000d
\$a1	5	0x00000000
\$a2	6	0x00000000
\$a3	7	0x00000000
\$t0	8	0x00000020
\$t1	9	0x0000000c
\$t2	10	0x0000002e
\$t3	11	0x00000021
\$t4	12	0x0000000b
\$t5	13	0x00000001
\$t6	14	0x00000000
\$t7	15	0x0000002e
\$a0	16	0x00000000
\$a1	17	0x00000000
\$a2	18	0x00000000
\$a3	19	0x00000000
\$a4	20	0x00000000
\$a5	21	0x00000000
\$a6	22	0x00000000
\$a7	23	0x00000000
\$t8	24	0x00000000
\$t9	25	0x00000000
\$k0	26	0x00000000
\$k1	27	0x00000000
\$gp	28	0x10008000
\$sp	29	0x7ffffcfc
\$fp	30	0x00000000
\$ra	31	0x00000000
pc		0x00400080
hi		0x00000000
lo		0x00000000



## Advantages of Carry Look Ahead Adder

The advantages of carry look ahead adder are-

- It generates the carry-in for each full adder simultaneously.
- It reduces the propagation delay.

## Disadvantages of Carry Look Ahead Adder-

The disadvantages of carry look ahead adder are-

- It involves complex hardware.
- It is costlier since it involves complex hardware.
- It gets more complicated as the number of bits increases.

## Applications

The **carry lookahead adder applications** are:

- Carry lookahead adders operating with high speed are employed as integrated circuits so that it is simple to integrate adder in many circuits. Also, the increase in the count of gates is even moderate when implemented for higher bits.

- When CLA's are used for high-bit calculations, the device offers more speed whereas the circuit complexity also increases. Usually, these are used for 4-bit modules so that they are integrated together for high-bit computations.
- On a regular basis, carry-lookahead adders are used in boolean computations.