

Assignment

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Question : Consider communication over a memoryless binary symmetric channel using a (7, 4) Hamming code. Each transmitted bit is received correctly with probability $(1 - \epsilon)$, and flipped with probability ϵ . For each codeword transmission, the receiver performs minimum Hamming distance decoding, and correctly decodes the message bits if and only if the channel introduces at most one bit error.

For $\epsilon = 0.1$, the probability that a transmitted codeword is decoded correctly is _____ (rounded off to two decimal places). (rounded off to two decimal places).

Solution: Given that, Let X be a random variable defined in the Table I;

RV	Value	Description
n	7	The total number of bits
ϵ	0.1	Probability of error in transimitted bit
X	$0 \leq X \leq 7$	The number of bit errors in transmission

TABLE I

RANDOM VARIABLE X DECLARATION

Then, $X \sim \text{Bin}(n, p)$ where

$$n = 7 \quad p = \epsilon = 0.1 \quad (1)$$

the pmf of X is given by

$$p_X(k) = {}^7C_k(\epsilon)^k(1 - \epsilon)^{7-k} \quad (2)$$

the cdf of X is given by

$$F_X(k) = \sum_{i=0}^k {}^7C_i(\epsilon)^i(1 - \epsilon)^{7-i} \quad (3)$$

From equation (3), the probability of getting one or less error is given by

$$F_X(1) = \sum_{i=0}^1 {}^7C_i(\epsilon)^i(1 - \epsilon)^{7-i} \quad (4)$$

$$= {}^7C_0(\epsilon)^0(1 - \epsilon)^7 + {}^7C_1(\epsilon)^1(1 - \epsilon)^6 \quad (5)$$

$$= (1 - \epsilon)^7 + 7(\epsilon)^1(1 - \epsilon)^6 \quad (6)$$

From (1) and (6),

$$F_X(1) = (1 - 0.1)^7 + 7(0.1)^1(1 - 0.1)^6 \quad (7)$$

$$= 0.85 \quad (8)$$

\therefore the probability that a transmitted codeword is decoded correctly is 0.85.

Gaussian

Let parameters be defined in the Table II;

RV	Value	Description
$\mu = np$	1.25	Mean of Binomial distribution
$\sigma^2 = npq$	0.9375	Variance of Binomial distribution

TABLE II
PARAMETERS

Central limit theroem states that , Z be a random variable

$$Z \approx \frac{X - \mu}{\sigma} \quad (9)$$

Z converges to normal distribution for large value of n

$$f(x) = \frac{1}{\sqrt{2\pi}} e^{-\frac{x^2}{2}} \quad (10)$$

Q function is defined

$$Q(x) = \int_x^\infty f(x) dx \quad (11)$$

then CDF of X is:

$$F_X(k) = \int_{-\infty}^k f(x) dx \quad (12)$$

$$= 1 - \int_x^\infty f(x) dx \quad (13)$$

$$= 1 - Q(x) \quad (14)$$

From (1) and (14),

$$F_X(1) = 1 - Q(1) \quad (15)$$

$$= 0.84134 \quad (16)$$

\therefore the probability that a transmitted codeword is decoded correctly is 0.84134. The CDF plot is given in fig1

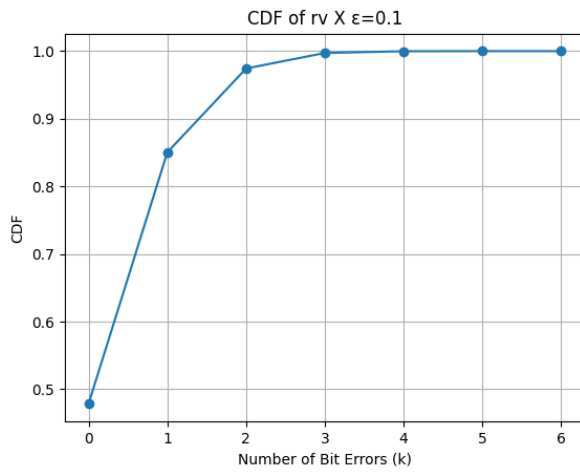


Fig. 1. CDF plot