

## Linked List :

Linked List is an Linear Data Structure where elements (nodes) are connected with each other.

Each Node has two parts:

1) Data → Actual value stored

2) Next → Link/Reference the points to the next node

Eg.

5	200
---	-----

6	null
---	------

100

↓

address of value of  
the node the node

next

Part to

Point

next

node

Basically  
Elements

single.

A linked list has 2

- \* Head → Starting Element

- \* Last / Tail → Final Element

In a Linked List Last Always  
points to next or null.

25/10/20

## Advantages :

- 1) The list can grow or shrink  
[Dynamic Memory]
- 2) Easy Insertions / Deletions.

## Disadvantage :

- 1) Each Node Requires additional Memory for the pointer

- 2) Accessing an Element has  $O(n)$  while array is  $O(1)$ .

## Types of Linked List

1) Single

2) Doubly

3) Circular

- 20) Find & Delete  $N^{\text{th}}$  Node from the List [Important X]

## Basic Problems

for Single Linked List.

- 1) Traversal ✓
- 2) Insertion ✓
- 3) Deletion ✓ [ updating an List ] ✓  
Not written
- 4) Search an Element ✓
- 5) Length of List ✓
- 6) Middle Node ✓
- 7) Find Loop & Remove ✓
- 8) Reverse a List ✓
- 9) Palindrome / Not ✓
- 10) Remove Duplicates ✓
- 11) Merge List (Sorted) ✓
- 12) Sorting of a List ✓
- 13) Delete Node without head / start. ✓
- 14) Intersection of two List ✓
- 15) Reverse Nodes with specified Index
- 16) Delete Nodes greater than value
- 17) Rotate a List ✓
- 18) Add two List & two Numbers of List
- 19) Split a List.

# Single Linked List

Structure:

Class Node {

int data;

Node next;

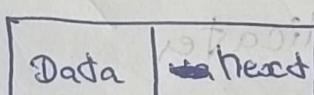
Node (int data) {

this.data = data;

this.next = null;

}

This will help the operations to create an Element (Node with next and Data Value).



Traversal:

Traversing through the node.

public void print()

if (start == null)

print ("Empty");

Node dptr = start;

while (dptr != null) {

print (dptr.data);

dptr = dptr.next;

}

## ★ Insertion :

### 1) At Beginning :

```
insert Begin (int data) {
```

```
    Node nn = new Node (data);
```

```
    if (start == null) {
```

```
        start = nn;
```

```
        return;
```

```
} else {
```

```
    nn.next = start;
```

```
    start = nn;
```

```
}
```

```
// Create a node (nn)
```

```
// nn's next should be start
```

```
so, it will come at start.
```

```
// Change the start to nn.
```

### 2) insert at End :

```
insert End (int data) {
```

```
    Node nn = new Node (data);
```

```
    if (start == null) {
```

```
        start = nn;
```

```
        return;
```

```
}
```

else {

    Node tptr = start;

    for ( ; tptr != null ; tptr = tptr.next);

        tptr.next = nn;

}

// Iterate to the Last Node.

// Change the Last Node's next to nn.

\* Insert At Position :

Insert At (int index, int data) {

~~if (start == null)~~

    if (index == 0) {

        insert At Begin (data);

    } else {

        Node nn = new Node (data);

        Node tptr = start;

        for (int i = 0 ; i < index - 1 &&

            ; tptr != null , tptr = tptr.next)

            if (tptr == null)

                System.out.print ("Invalid

                Index");

            nn.next = tptr.next;

            tptr.next = nn

}

} // Iterate till the index at change

★ Search an Element or Value:

```
Search(int value) {  
    if (start == null) {  
        return -1;  
    } else {  
        Node tptr = start; int i = 0;  
        for (; tptr != null && tptr.data != val; i++);  
        return i;  
    }  
}  
  
// Check for an Empty List  
// Iterate the Node and the  
// index till the Node (tptr)  
// reaches the val  
// Also check whether the  
// tptr is null, if it is null  
// return -1  
// Else Return the index that  
// has been iterated.
```

★ Length of the List :

length() {

    Node \*ptr;

    int Count = 0;

    ptr = head;

    while (ptr != null)

    {

        Count++;

        ptr = ptr->next;

    }

    return Count

}

// Iterate the Node till the  
end, and Count.

★ Reverse Program :

reverse() {

    Node prev = null;

    Node curr = start;

    Node next = null;

    while (curr != null) {

        next = curr->next;

        curr->next = prev;

        prev = curr;

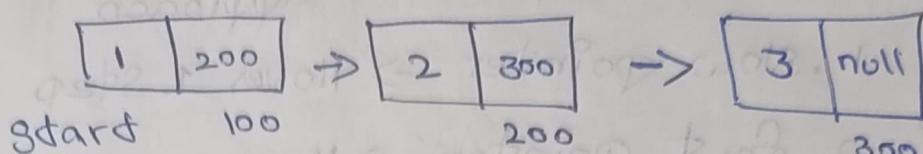
        curr = next;

    }

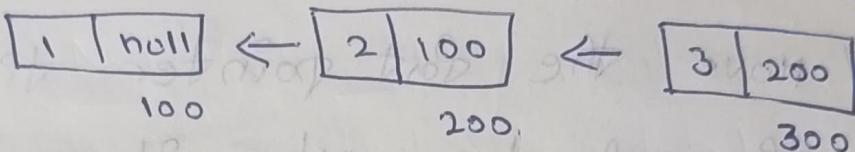
    head = prev

}

## \* Reversing a List:



**Reverse**



// use three Nodes

// Prev, curr, next

// iterate the curr.

// store the current's next in next

// change the curr.next to prev

// change the current to next.

At start, prev = null, curr = start,

next = null

Curr	Prev	Next	curr.next
100	null	200	100.next = null
200	100	300	200.next = 100
300	200	null	300.next = 200
null	300	null	Reverse Completed.

// Prev will point at last hence  
change start to prev

\* Middle Node : ~~best~~ no prior knowledge  
// use two pointer approach  
// 1. slow pointer → moves 1 step  
// 2. fast pointer → moves 2 steps  
// Move both the pointers at  
at time.  
// Move the fast pointer till the  
the end value - when it reaches  
null, The data pointed by slow is  
middle value

middle()

Node slow = head;

Node fast = head;

if (head == null) {

System.out.println("List is Empty");

return;

else {

while (fast != null && fast.next != null)

{

fast = fast.next.next;

slow = slow.next;

}

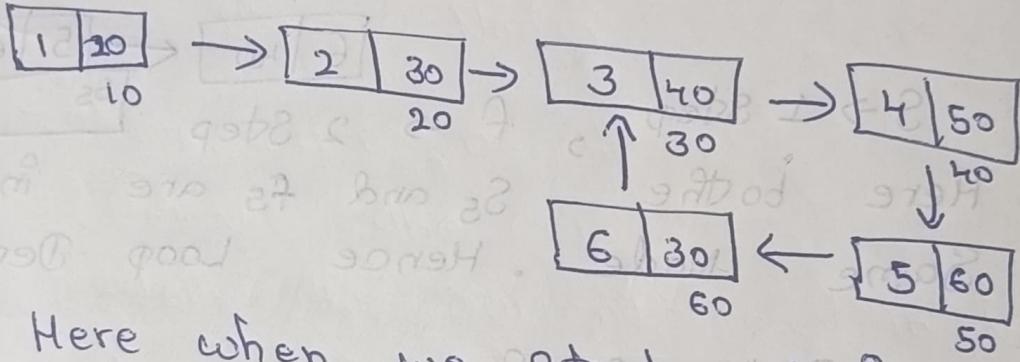
return slow;

}

## Finding Loop & Removing it:

// In a SLL Loop occurs when  
the last node points the next  
to any of its previous Nodes.  
[Last must point to null]

// Example:



// Here when we start printing  
the node  
The values from addresses 30 to 60  
Repeats, such that

1 2 3 4 5 6 3 4 5 6 3 4 5 6 ...

Algorithm: [Finding Loop]

Same as Middle use two  
pointer approach.

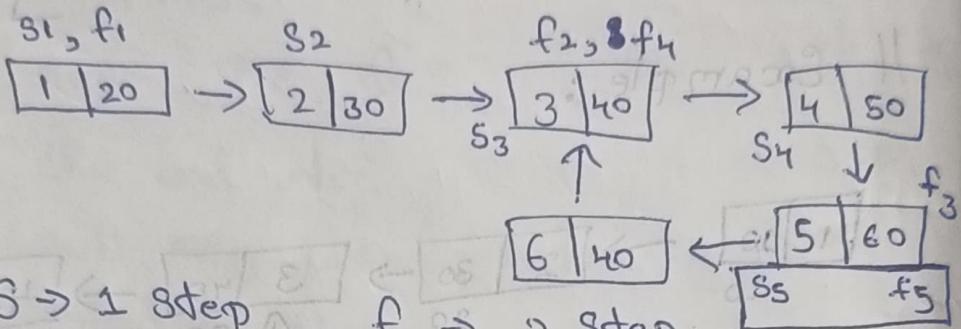
// slow  $\rightarrow$  moves 1 step

// fast  $\rightarrow$  moves 2 steps.

If slow and fast meets at

a certain point [then Loop exists]

Else Loop does not exists.



$S \rightarrow 1$  step,  $f \rightarrow 2$  step

Here both the  $s_5$  and  $f_5$  are in  
Same  $\rightarrow$  Node. Hence Loop Detected

loop Detect()

Node slow = head;

Node fast = head;

while (fast != null && fast.next != null)

{

...  
slow = slow.next;

fast = fast.next.next;

if (slow == fast) {

return true // has Loop

}

return false

}

## Algorithm : Loop Removal.

// use 4 pointers:

// slow = as same as Detect

// fast = as same as Detect

// prev = To keep track of slow's prev

// ptr = To remove Loop

// If Loop is detected (fast == slow)

Start moving the ptr and slow

same [1 step]. They will meet

at a point [Loop start] use prev to  
remove the Loop.

---

RemoveLoop () {

Node fast = head;

Node slow = head;

Node prev = null;

Node ptr = head;

while (fast != null && fast.next != null) {

    fast = fast.next.next;

    prev = slow;

    slow = slow.next;

    if (slow == fast) {

        ptr = ptr.next;

        prev = slow;

        slow = slow.next;

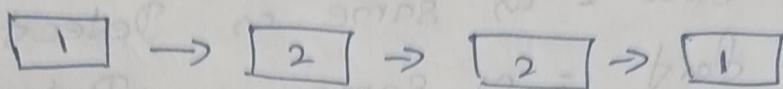
        prev.next = null;

    return head;

} - 3

★ List is Palindrome (Not :)

Eg. [Palindrome List]



① USING Stack :

Palindrome (Node head) {

    Node curr = head;

    while (curr != null) {

        Stack.push(curr.data);  
        curr = curr.next;

}

    while (head != null) {

        int c = stack.pop();

        if (head.data != c) {

            return false;

}

        head = head.next;

}

    return true;

}

// Create a stack and push the elements, Check whether the popped and head data is equal.

②

Optimal Solution: [splitting and checking]

is Palindrome (head) {

if (head == null || head.next == null)  
return true;

Node slow = head, fast = head;

while (fast != null && fast.next != null) {

fast = fast.next.next;

slow = slow.next;

Node right = reverse (slow);

Node left = ~~slow~~ head;

while (right != null) {

if (right.data != left.data) {

return false;

}

right = right.next;

left = left.next;

}

return true;

}

|| Split the List into 2

|| reverse the 2<sup>nd</sup> half (right)

|| Check whether Both the lists are

Equal.

Remove Duplicates:

remove (Node head) {

Node dptr = head;

while (dptr != null && dptr.next != null)

{ but head = word

if (dptr.val == dptr.next.val)

{ head, head.next = word

{ next.word = word

dptr.next = dptr.next.next;

}

else

{

dptr = dptr.next;

}

return head;

{

II Traverse through the List

II If dptr(data) & dptr's next(data)  
is equal remove one of the

Node

Eg: ① → ① → ②

Input

① → ②

Output

.page 3

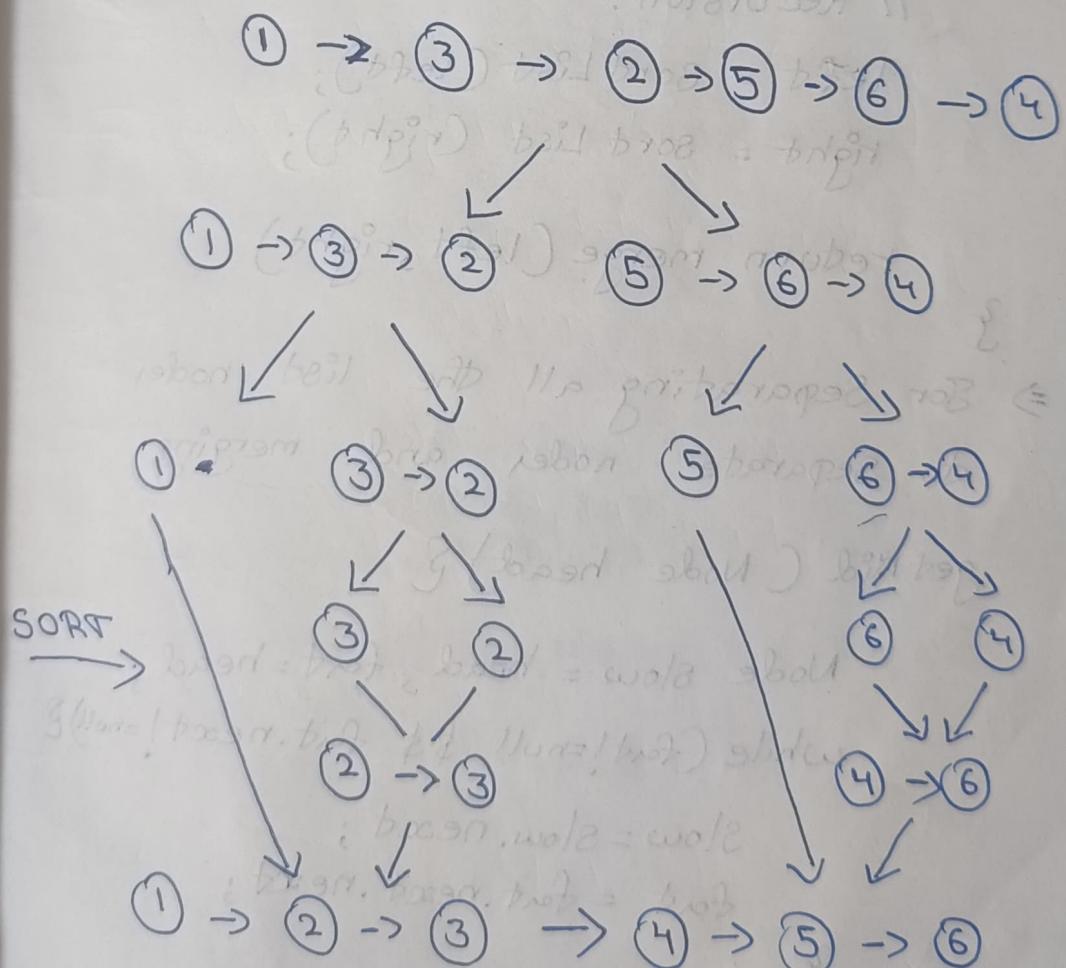
# ★ Sorting of an Linked List :

## \* MERGE SORT

Algorithm :

- 1) Find the Middle Node.
- 2) Recursively sort the lists that are Separated
- 3) Merge them.

Example :



Code:

```
* sortList (Node head) {  
    if (head == null || head.next == null) {  
        return head;  
    }  
    Node mid = getMid (head);  
    Node left = head;  
    Node right = mid.next;  
    mid.next = null;  
    // Recursion:  
    left = sortList (left);  
    right = sortList (right);  
    return merge (left, right);  
}  
⇒ For Separating all the list nodes  
do Separate nodes and merging
```

```
* getMid (Node head) {  
    Node slow = head, fast = head;  
    while (fast != null && fast.next != null) {  
        slow = slow.next;  
        fast = fast.next.next;  
    }  
    return slow;  
}
```

⇒ For getting the Middle Node.

\* merge (Node left, Node right) {

    Node merged = new Node (-1);

    Node curr = merged;

    while (left != null && right != null) {

        if (left.data < right.data) {

            curr.next = left;

            left = left.next;

        } else {

            curr.next = right;

            right = right.next;

}

        curr = curr.next;

}

    if (left != null)

        curr.next = left;

    else

        curr.next = right;

    return merged.next;

⇒ Lastly create an empty node and add the values of both lists in such way that which is smaller.

\* Delete an Node without head [The node to be Deleted is given as input]

Delete-Node (Node value) {

    Node tptr = value;

    Node prev = null;

    while (tptr != null && tptr.next != null)

        tptr.val = tptr.next.val;

        prev = tptr;

        tptr = tptr.next;

    prev.next = null;

// when an head is not given, we cannot access the node's [To be Deleted] previous value.

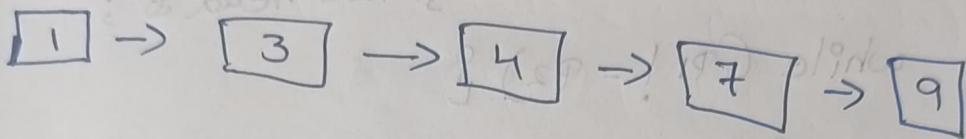
// So, just swap the node's next value to the node and traverse for all elements.

2021 2022 To make add int b2021 2022 in pointers both your node in

★ Find | Delete  $k^{th}$  Node from last.

Since it is an linked-list we can't have access to last element or count of nodes, So, we have to do it manually.

$$\text{Eq. } K = 2$$



Code :

Delete - Nth - node (Node head, int k) {

Node slow = head, fast = head.

for (int i=0; i<k; i++) { }  
for (int i=0; i<k; i++) { }  
for (int i=0; i<k; i++) { }

?  $\text{fast} = \text{fast.next}$ ; ~~head = 9~~

if (fast == null) {

return head.next

3

```
while(fat.next != null) {
```

~~fact = fact.next~~

~~Slow = slow.next~~

۳

~~Slow.next = slow.next.next~~

... - snow.heat.the  
. 273V30iv has best. 273000 ob

## ★ Intersection of Two lists :

Intersection of list (Node head A, Head B)

}

if (head A == null || head B == null) {

    return null;

    Node p1 = head A, p2 = head B;

    while (p1 != p2) {

        if (p1 == null) {

            p1 = head B;

        } else {

            p1 = p1.next;

        if (p2 == null) {

            p2 = head A; best = best

        } else {

            p2 = p2.next;

    return p1;

    // generate through Both the list  
    // till any of Node is equal, If a pvt  
    // reaches end change its direction / pos  
    // to another list and vice versa.  
    // Both the lists will meet at the  
    // intersection point.

★ Rotate a List : ~~using stack~~

rotate-list (Node head) {  
    int k;

    if (head == null || head.next == null  
        || k == 0) {

        return head;

    }

    Node dptr = head;

    int count = 1;

    while (dptr.next != null) {

        count++;

        dptr = dptr.next;

}

    dptr.next = head;

    k = k % count;

    k = count - k;

    while (k > 0) {

        dptr = dptr.next;

        k++

}

    head = dptr.next;

    dptr.next = null;

    return head;

}

|| Read the Code you will understand  
Automatically.