2   x > 2   (co) 2 3 6  Dote Date Page
Matrix Addition int int
Matrix Addition  Noid add Matrices (int m1[3][3], int m2[3][3], int res[3][3], rows, columns) {
Jor (int ) = 0; 1< rows; 1+1/1
for (int j=0; j < columns; j++) {
KS [i][j] = m1[i][j] + m2[i][j];
3
}
}
Acres St. Co. Co. Co. Co. Co. Co. Co. Co. Co. Co
 Matrix Transpase
 void transpase Matrix (int mat [3][3], int rows, int columns) {
int res[columns][rows];
for (int i=0; icrows; i++){
for (int j=0; j < columns; j++) {
AS (1) (1) = M(1 (1) 2013)
7
3
Est file the first a manifestion and 13
Matrix Multiplication
void multiply Matrices (int m1[3][3], int m2[3][3], int res[3][3], int r1, int c1, int c2)
for (inti=0; ic rows rl; itt){
for (int j=0; j < c2; j++){
res (i) (j] = 0;
for (int k=0; k < cl; k++) {
RS[i](j]+= m [i][k] * m2[k][j];
3
2
7
<u>J</u>



Sparse Matrix Addition struct Sparse \* add (struct Sparse \* st. struct Sparse \* s2) { struct Sparse \*sum; sum = (struct Sparse \*) malloc (sizeos (struct Sparse)); sum > e = (struct Element \*) malloc ((stonum + s2 onum) \* size of (struct Element)): int i=0 j=0, k=0; while ( i < sl > e[i].; \$8 j < s2 > e[j].j) { if (sl→e(i].i < s2→e[;].i) { sum→e[k++]= sl→e[i++];} else if (slae[i] i > 52 pe [j] i) { sum = e[k++] = 52 a e [j++]; } if (sl→e[i].; < s2 →e[j].;) { sym+e[k++] = sl →e [i++];} else if (slae[i].; > 52 se[j].j) {sumae[k++] = 52 ae [j++];} else f

sum → e[k] = sl → e[i]:

sum > e(k]. x = s1 → e(i++7.x + s2 → e[i++].x;

3

for (; icslanum; i++) { sum ae[k++] = slae[i];} for (; j < s2 → nym; j++) { sym →e[k++] = s2 →e[j];}

Sum-num= k 11 k Stoks no. of non-zero elements

sum -) m = sl + m. syman = slan;

return sym;

2 12 / (0/ 236 Sparse Matrix Transpose struct Sparse \*transpare (struct Sparse \*s) { struct Sparse \* result = (struct Sparse \*) malloc (size of (struct Sparse)); 18541+ +m = 5 +n result + n = s+m nsult - num = s - num result > c = (struct Element \*) malloc (result > num \* size of (struct Element)); int col Count [son];

for (int i=0; ix s=n==; i++) { col Count[i] = 0;}

for (int i=0; i <s >nym; i++) { col Count[s > eli]; ]++; }

int colposition [s →n];

colPosition[0]=0;

for (int i=1; i< s→n; i++){

colposition [i] = colposition [i-1] + col Count[i-1];

for (int i= 0; i < s > nym; i++) {

int j = s + eli ];

result >e [colposition[j]]. i = s >e [i].j; result -e [ Col Position []]]. j = s -e [].i;

result to [ Colposition (j]]. x = ste [i].x;

colposition GJ++;

return usult;

Sporse Motrix Multiplication struct Sparse \*multiply (struct Sparse \*s), struct Sparse \*s2) { it (slan != s2am) { print("Multiplication not possible"); return NULL; } struct Sparse \* result; result = (struct Sporse +) malloc (size of (struct Sporse)); Hoult -m = slam; result in = s2 in. result + nym = 0. result -e = (struct Element +) mg/loc (result - num \* sizeof (struct Element)); int sum, k=0; for (int i=0; icsl+m; i++) { for (int j=0; j <52-n; j++) { sym=0; for (int l=0; l < sl + num; l++){ for (int m: 0; m < s2 + nym; m++) { if (slae[1].; == 52 ae[m]. i 28 slae[1]. i == i 48 52 ae[m]. j == j) { sym += si>e[l].x + s2>e[m].x; 3 . 10 Ed. 12 x 1 W 1 . 4 : A . 1 1 5 - . . . if (sym!=0) {

result - num ++:

k++;

3

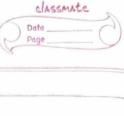
return result:

result  $\rightarrow e[k] \cdot j = j'$ ; result  $\rightarrow e[k] \cdot x = sym;$ 

result -> e = (struct Element x) realloc (result -> e, result -) num\*

 $\text{result} \rightarrow \text{e [k]} \cdot i = i$ ; size of (struct Element)

	2 k 2 1 / co/238	)
	Polynomial Addition	
	struct Poly moded (struct Poly mpl, struct Poly *p2) {	
	struct Poly tsum;	
-	sum = (struct Poly +) mella (sizeaf (struct Poly));	
	sum + terms = (struct Term+) molloc((p)+n+p2+n) * size of (struct Term));	
	int 1:0, 1:0, k:0;	
	while (icplan (1 jcp2 an) [	_
	if (pl-terms li] exp > p2 + terms [i] exp) {sumtkerns (k++] = pl-) terms [i++	5;
	else if (pl->terms [i] exp < p2-> terms [i] exp) { sum > terms [k++]=p2-> terms [j++	ز
	La la company of the first of the company of the co	
	sum > terms[k]. exp = p) -> terms[i]. exp;	
	sym + terms [k++]. coeff = p) + terms [i++]. coeff + p2 -> terms [i++]. coeff;	
	}	_
744	3 - Les Carres Des labordes d'ar les	_
	for (; icpl-n; i+t) { sym-> terms (k++) = pl -> terms (i]; }	
	for (; j  terms (k++] = p ≥ → terms [j]; }	_
	sum → n= k // k stores no. of elements in sum	_
	return sum; feel logarities Shares to the same	_
	3-2 State of the second of the	
	United States	
		_
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p.		_
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Polynomial Multiplication struct Poly \*multiply (struct Poly \*p), struct Poly \*p2) {

struct Poly \* result = (struct Poly \*)malloc (size of (struct Poly));

for (int i=0; i < result =n-1; i++){

for (int j=i+); i < Hsult >n : j++) { if (result - 4 mms [i]. exp = = result - terms [j]. exp) {

result -> termsli]. coest += result -> terms [j]. coest; for (int 1=j , 1 = result =n-1; l+t) { result = terms[s] = result = terms[l+1];} result - n --:

\_j -- ;

return result;

for (int i = 0; jep2 = n; j++) { result > terms [k]. coeff = pl > terms [i]. coeff \* p2 > terms [i]. coeff. Asult -> terms [k]. exp = pl -> terms [i]. exp + p2-> terms [j]. exp;

HS4/+ +n= pl+n+p2+n; result = terms = (struct Term \*) molloc (result +n \* size of (struct Term)); int k = 0 : for (int i=0; ixplan; it) {

1	J	Ċ	0/	2	3	ė

2/22 Using Stack, check for Palindrome string bool is Palindrome (const char \* str) { struct Stack stack: initialize ( & stack); int length = strun (str); For (int ico; ic langth; itt) { push (Astack, str [i]); } for (int i=0; ix length; i++) { is (pap (Estack)! = str [i]) { return false; } return true: Application of Hegps as prisrity queues 1- Initialization: Create empty hap datastructure. This can be bingry min-hap or binary max-heap, depending on whether you want to implement min or max-priority queve.

2. Insertion: To insert on element with a priority value into priority queue

- Re-heapity the heap to maintain heap property (min-heap or max-heap property) 3. Extraction: To extract on element with highest (max-priority queue) or lowest (min-priority queue) priority from queue:

-Add element to hegp

- Remove cast element

- Retyrn root element of Meap

- Re-heapity heap to maintain heap property

priority queue) priority without removing it:

- Replace root with last element in Hegp.

- Retrieve root element of Hugp (one with Mighest or Lowest priority)

4. Peck: To peck of element with highest (max-priority queue) or lowest (min-