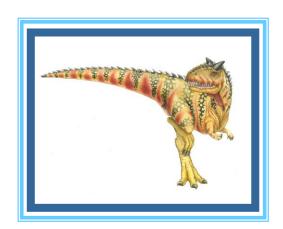
Introduction to Operating Systems





What is an Operating System?

- A program that acts as an intermediary between a user of a computer and the computer hardware
- Operating system goals:
 - Execute user programs and make solving user problems easier
 - Make the computer system convenient to use
 - Use the computer hardware in an efficient manner





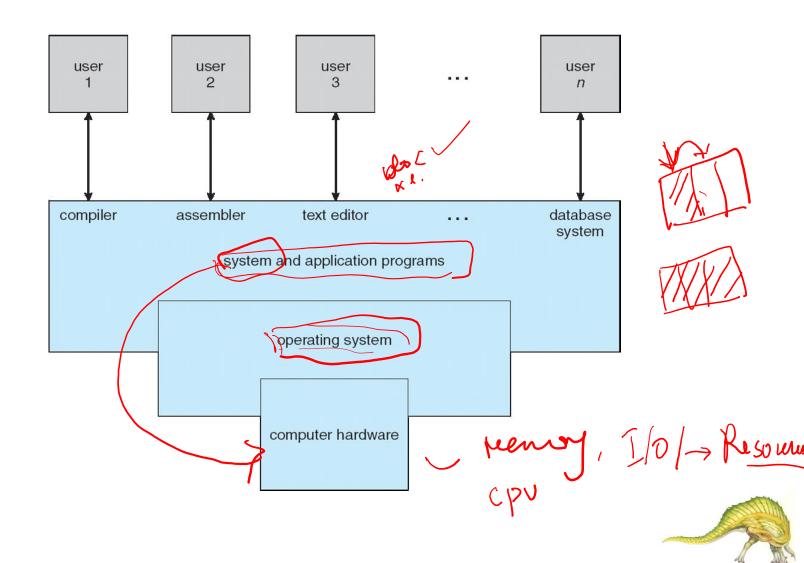
Computer System Structure

- Computer system can be divided into four components:
 - □ Hardware provides basic computing resources
 - ▶ CPU, memory, I/O devices
 - Operating system
 - Controls and coordinates use of hardware among various applications and users
 - Application programs define the ways in which the system resources are used to solve the computing problems of the users
 - Word processors, compilers, web browsers, database systems, video games
 - Users
 - People, machines, other computers





Four Components of a Computer System





What Operating Systems Do

- Depends on the point of view
- Users want convenience, ease of use and good performance
 - Don't care about resource utilization
- But shared computer such as mainframe or minicomputer must keep all users happy
- Users of dedicate systems such as workstations have dedicated resources but frequently use shared resources from servers
- Handheld computers are resource poor, optimized for usability and battery life
- Some computers have little or no user interface, such as embedded computers in devices and automobiles





Operating System Definition

- OS is a resource allocator
 - Manages all resources
 - Decides between conflicting requests for efficient and fair resource use
- OS is a control program
 - Controls execution of programs to prevent errors and improper use of the computer





Operating System Definition (Cont.)

- No universally accepted definition
- "Everything a vendor ships when you order an operating system" is a good approximation
 - But varies wildly
- "The one program running at all times on the computer" is the kernel.
- Everything else is either
 - a system program (ships with the operating system), or
 - an application program.



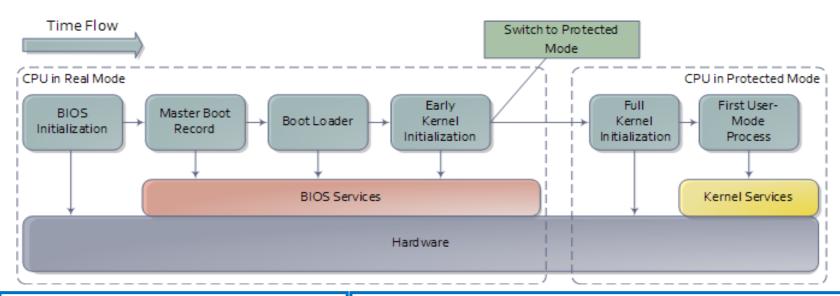


Computer Startup

- bootstrap program is loaded at power-up or reboot
 - Typically stored in ROM or EPROM, generally known as firmware
 - Initializes all aspects of system
 - Loads operating system kernel and starts execution



Starting an Operating System(Booting)



- ✓ Power On Switch sends electricity to the motherboard on a wire called the *Voltage*Good line.
- ✓ If the power supply is good, then the BIOS (Basic Input/Output System) chip takes over.
- ✓ In Real Mode, CPU is only capable of using approximately 1 MB of memory built into the motherboard.
- ✓ The BIOS will do a Power-On Self Test (POST) to make sure that all hardware are

- ✓ BIOS will then look for a small sector at the very beginning of your primary hard disk called MBR.
- The MBR contains a list, or map, of all of the partitions on your computer's hard disk (or disks).
- After the MBR is found the Bootstrap Loader follows basic instructions for starting up the rest of the computer, including the operating system.
- In Early Kernel Initialization stage, a smaller core of the Kernel is activated.
- This core includes the device drivers needed to use computer's RAM chips.



System Boot

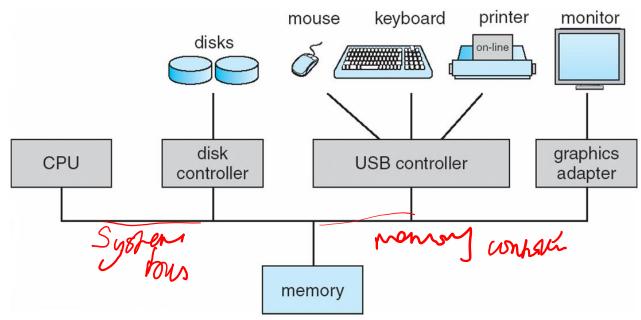
- When power initialized on system, execution starts at a fixed memory location
 - Firmware ROM used to hold initial boot code
- Operating system must be made available to hardware so hardware can start it
 - Small piece of code bootstrap loader, stored in ROM or EEPROM locates the kernel, loads it into memory, and starts it
 - Sometimes two-step process where boot block at fixed location loaded by ROM code, which loads bootstrap loader from disk
- Common bootstrap loader, GRUB, allows selection of kernel from multiple disks, versions, kernel options
- Kernel loads and system is then running





Computer System Organization

- Computer-system operation
 - One or more CPUs, device controllers connect through common bus providing access to shared memory
 - Concurrent execution of CPUs and devices competing for memory cycles







Computer-System Operation

- I/O devices and the CPU can execute concurrently
- □ Each device controller is in charge of a particular device type
- Each device controller has a local buffer
- CPU moves data from/to main memory to/from local buffers
- I/O is from the device to local buffer of controller
- Device controller informs CPU that it has finished its operation by causing an interrupt





Common Functions of Interrupts

- Interrupt transfers control to the interrupt service routine generally, through the interrupt vector, which contains the addresses of all the service routines
- Interrupt architecture must save the address of the interrupted instruction
- A trap or exception is a software-generated interrupt caused either by an error or a user request
- An operating system is interrupt driven





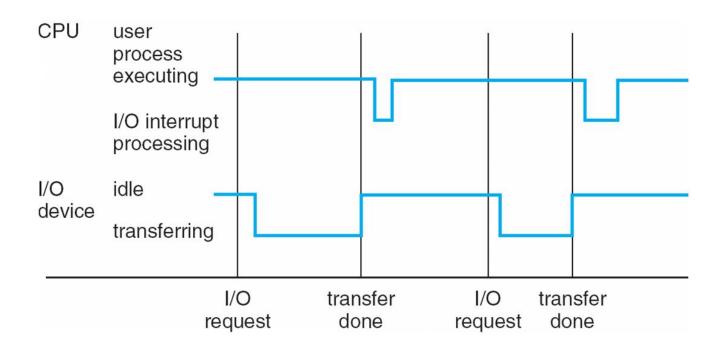
Interrupt Handling

- The operating system preserves the state of the CPU by storing registers and the program counter
- Determines which type of interrupt has occurred:
 - polling
 - vectored interrupt system
- Separate segments of code determine what action should be taken for each type of interrupt





Interrupt Timeline







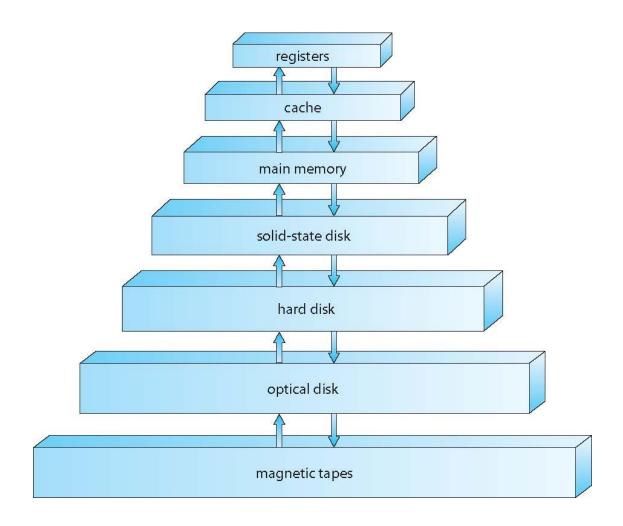
Storage Structure

- Main memory only large storage media that the CPU can access directly
 - Random access
 - Typically volatile
- Secondary storage extension of main memory that provides large nonvolatile storage capacity
- Hard disks rigid metal or glass platters covered with magnetic recording material
 - Disk surface is logically divided into tracks, which are subdivided into sectors
 - The disk controller determines the logical interaction between the device and the computer
- Solid-state disks faster than hard disks, nonvolatile
 - Various technologies
 - Becoming more popular





Storage-Device Hierarchy







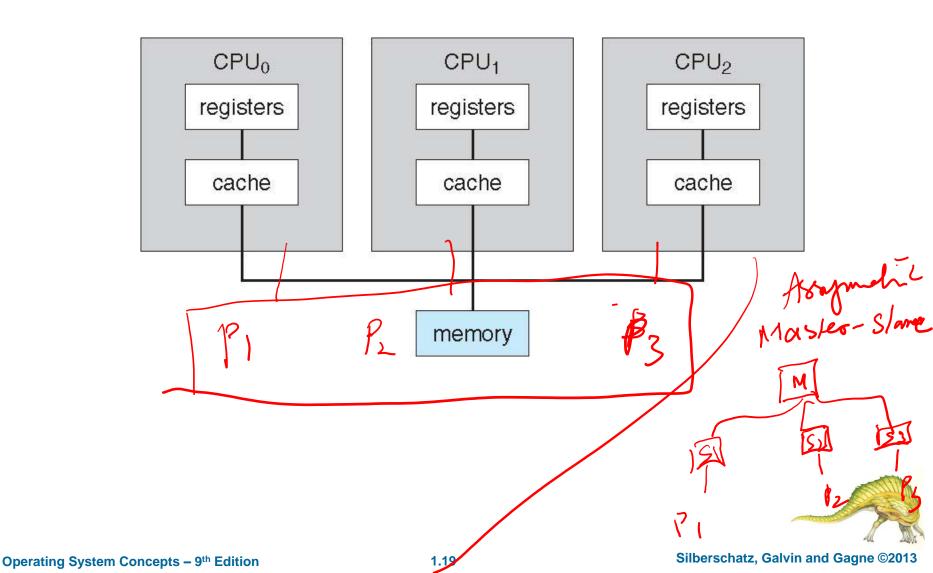
Computer-System Architecture

- Most systems use a single general-purpose processor
 - Most systems have special-purpose processors as well
- Multiprocessors systems growing in use and importance
 - Also known as parallel systems, tightly-coupled systems
 - Advantages include:
 - 1. Increased throughput
 - 2. Economy of scale
 - 3. Increased reliability graceful degradation or fault tolerance
 - Two types:
 - Asymmetric Multiprocessing each processor is assigned a specific task.
 - 2. Symmetric Multiprocessing each processor performs all tasks





Symmetric Multiprocessing Architecture





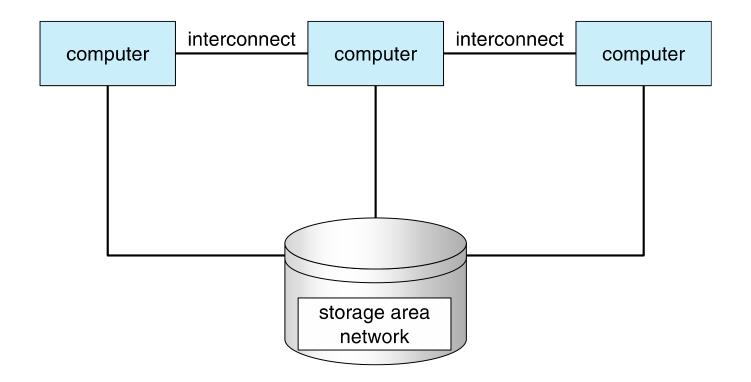
Clustered Systems

- ☐ Like multiprocessor systems, but multiple systems working together
 - Usually sharing storage via a storage-area network (SAN)
 - Provides a high-availability service which survives failures
 - Asymmetric clustering has one machine in hot-standby mode
 - Symmetric clustering has multiple nodes running applications, monitoring each other
 - Some clusters are for high-performance computing (HPC)
 - Applications must be written to use parallelization
 - Some have distributed lock manager (DLM) to avoid conflicting operations





Clustered Systems





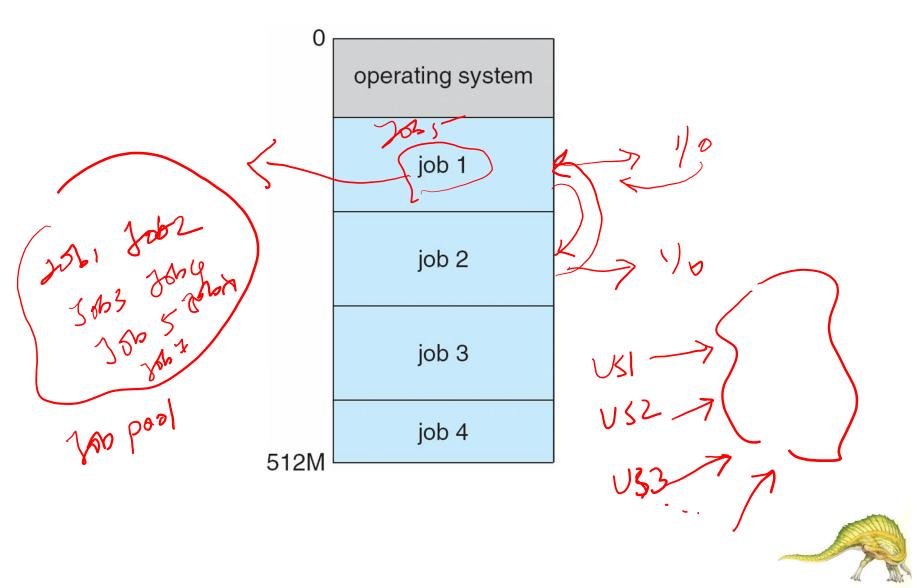


Operating System Structure

- Multiprogramming (Batch system) needed for efficiency
 - □ Single user cannot keep CPU and I/O devices busy at all times
 - Multiprogramming organizes jobs (code and data) so CPU always has one to execute
 - A subset of total jobs in system is kept in memory
 - One job selected and run via job scheduling
 - □ When it has to wait (for I/O for example), OS switches to another job
- ☐ Timesharing (multitasking) is logical extension in which CPU switches jobs so frequently that users can interact with each job while it is running, creating interactive computing
 - □ Response time should be < 1 second</p>
 - Many users share the computer simultaneously
 - Uses CPU scheduling and multiprogramming to provide for each user
 - □ Each user has at least one program executing in memory ⇒ process
 - If several jobs ready to run at the same time ⇒Job and CPU scheduling
 - If processes don't fit in memory, swapping moves them in and out to runge
 - Virtual memory allows execution of processes not completely in memory



Memory Layout for Multiprogrammed System





Operating-System Operations

- Interrupt driven (hardware and software)
 - Hardware interrupt by one of the devices
 - Software interrupt (exception or trap):
 - Software error (e.g., division by zero)
 - Request for operating system service
 - Other process problems include infinite loop, processes modifying each other or the operating system





Operating-System Operations (cont.)

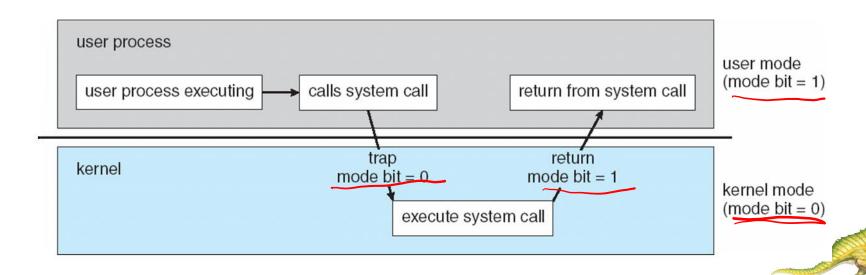
- Dual-mode operation allows OS to protect itself and other system components
 - User mode and kernel mode
 - Mode bit provided by hardware
 - Provides ability to distinguish when system is running user code or kernel code
 - Some instructions designated as privileged, only executable in kernel mode
 - System call changes mode to kernel, return from call resets it to user
- Increasingly CPUs support multi-mode operations
 - i.e. virtual machine manager (VMM) mode for guest VMs





Transition from User to Kernel Mode

- Timer to prevent infinite loop / process hogging resources
 - Timer is set to interrupt the computer after some time period
 - Keep a counter that is decremented by the physical clock.
 - Operating system set the counter (privileged instruction)
 - When counter zero generate an interrupt
 - Set up before scheduling process to regain control or terminate program that exceeds allotted time





Process Management

- A process is a program in execution. It is a unit of work within the system. Program is a passive entity, process is an active entity.
- Process needs resources to accomplish its task
 - CPU, memory, I/O, files
 - Initialization data
- Process termination requires reclaim of any reusable resources
- Single-threaded process has one program counter specifying location of next instruction to execute
 - Process executes instructions sequentially, one at a time, until completion
- Multi-threaded process has one program counter per thread
- Typically system has many processes, some user, some operating system running concurrently on one or more CPUs
 - Concurrency by multiplexing the CPUs among the processes / threads





Process Management Activities

The operating system is responsible for the following activities in connection with process management:

- Creating and deleting both user and system processes
- Suspending and resuming processes
- Providing mechanisms for process synchronization
- Providing mechanisms for process communication
- Providing mechanisms for deadlock handling







Memory Management

- To execute a program all (or part) of the instructions must be in memory
- All (or part) of the data that is needed by the program must be in memory.
- Memory management determines what is in memory and when
 - Optimizing CPU utilization and computer response to users
- Memory management activities
 - Keeping track of which parts of memory are currently being used and by whom
 - Deciding which processes (or parts thereof) and data to move into and out of memory
 - Allocating and deallocating memory space as needed





Storage Management

- OS provides uniform, logical view of information storage
 - Abstracts physical properties to logical storage unit file
 - Each medium is controlled by device (i.e., disk drive, tape drive)
 - Varying properties include access speed, capacity, datatransfer rate, access method (sequential or random)
- File-System management
 - Files usually organized into directories
 - Access control on most systems to determine who can access what
 - OS activities include
 - Creating and deleting files and directories
 - Primitives to manipulate files and directories
 - Mapping files onto secondary storage
 - ▶ Backup files onto stable (non-volatile) storage media





Mass-Storage Management

- Usually disks used to store data that does not fit in main memory or data that must be kept for a "long" period of time
- Proper management is of central importance
- Entire speed of computer operation hinges on disk subsystem and its algorithms
- OS activities
 - Free-space management
 - Storage allocation
 - Disk scheduling
- Some storage need not be fast
 - Tertiary storage includes optical storage, magnetic tape
 - Still must be managed by OS or applications
 - Varies between WORM (write-once, read-many-times) and RW (read-write)



I/O Subsystem

- One purpose of OS is to hide peculiarities of hardware devices from the user
- □ I/O subsystem responsible for
 - Memory management of I/O including buffering (storing data temporarily while it is being transferred), caching (storing parts of data in faster storage for performance), spooling (the overlapping of output of one job with input of other jobs)
 - General device-driver interface
 - Drivers for specific hardware devices





Protection and Security

- Protection any mechanism for controlling access of processes or users to resources defined by the OS
- Security defense of the system against internal and external attacks
 - Huge range, including denial-of-service, worms, viruses, identity theft, theft of service
- Systems generally first distinguish among users, to determine who can do what
 - User identities (user IDs, security IDs) include name and associated number, one per user
 - User ID then associated with all files, processes of that user to determine access control
 - Group identifier (group ID) allows set of users to be defined and controls managed, then also associated with each process, file
 - Privilege escalation allows user to change to effective ID with more rights



End of Chapter 1

