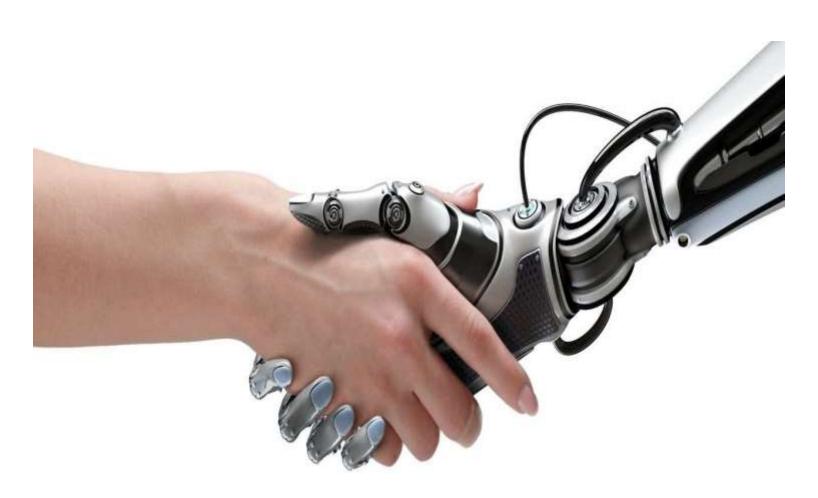
ARTIFICIAL INTELLIGENCE (A.I.)



| BLIND SEARCH SEARCH | INFORMED SEARCH |
|----------------------------------|--|
| BLIND UNINFORMED SEARCH | 1. Search with information |
| à. No knowledge | a. Use knowledge to find steps to solution. |
| 2 —° | steps to solution. |
| 3. Time Consuming. | 3. Quick Solution |
| 4. More Complexity (Time, space) | 4. hess complexity (Time, space) |
| 5. DFS, BFS etc V | 5. A*, Best first search |

HEURISTIC

Hewristic in AI' (Rule of thumbe) [What, Why, How]

The is a technique disigned to solve a problem

quickly.

BLIND SCARCH

Hewristic Search

8 puzzle
(0 (ba))
(320) → Time lomplikity
is more
Such problems are known
as NP (Non polynomial) problem

is Exponential

L. Advantage -> optimal

-> Reduce time

-> Multiple sol

Explore best

-> Quickly solve the problem

-> Best Solution

-> does not gourantee

| Hone to find hewist 1. Euclidian Distance | (straight Line Distance): |
|--|---|
| | (1/2) gird Motance |
| E.D = 1 (3 | (e-21)2+(ya-yi)2- |
| 2. Manhattan dista | nce (used when we want to find Vertical, horizontal distance) |
| 1 3 2 6 5 4 8 7 | 1 2 3 4 5 6 7 8 |
| start state | Goal State |
| In Blind Search (: | 320) -> then only you get optimal sol |
| To find herristic value | -> Use Manhatlan distance |
| H.B =0+1+1+2+ | |
| A District | La Explore all possible position La select convect Manhattan value |
| | les dess misblaced tile |

LE BRANCH AND BOUND SEARCH (UNIFORM COST SEARCH) → In branch and bound search method, cost function (denoted by g(x)) is designed that assigns cumulature expense to the path from start node to the current node x by applying the sequence of operators. -> Branch and bound search extands the least-lost path at each iteration till we reach goal state. - Sometimes et is also called a Uniform Cost search. -> Two parts -> Branch: Several chaices are found > Bound: Setting bound on sol Quality Ly Pruring teimming of branches where I solution quality is poor.

```
Algorithm (Branch and Bound)
Input: START and GOAL states
Local Variables: OPEN. CLOSED. NODE. SUCCs. FOUND:
Output: Yes or No
Method:

    initially store the start node with g(root) = 0 in a OPEN list: CLOSED = φ:

   FOUND = false:

    while (OPEN ≠ o and FOUND = false) do

    remove the top element from OPEN list and call it NODE:

    • if NODE is the goal node. then FOUND = true else

    put NODE in CLOSED list:

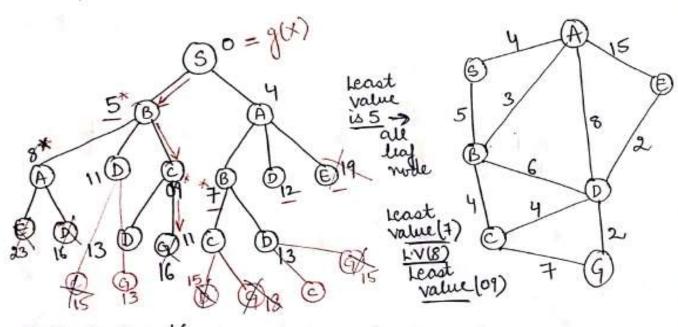
      • find SUCCs of NODE. if any. and compute their 'g' values and store them
        in OPEN list:

    sort all the nodes in the OPEN list based on their cost-function values:

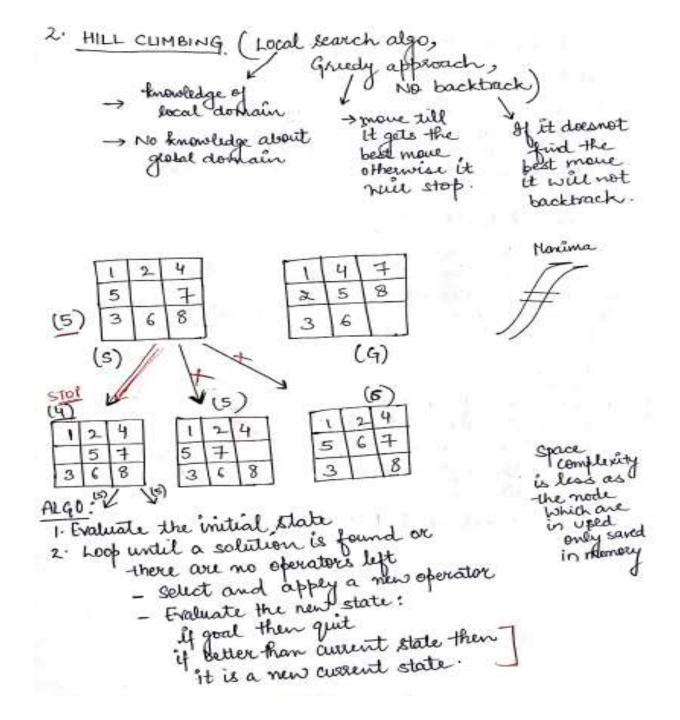
} /* end while */
  • if FOUND = true then return Yes otherwise return No:

    Stop
```

Example



- 1. Bound set = 16 value greater than 16 is prured off.
- 2. Now Explore the path having value less than 16
- 3. Now discard path having value greater than 13.

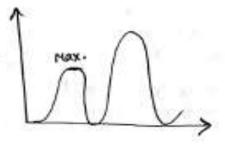


Algorithm (Simple Hill Climbing) Input: START and GOAL states Local Variables: OPEN. NODE, SUCCs. FOUND: Output: Yes or No Method: . store initially the start node in a OPEN list (maintained as stack): FOUND = false: while (OPEN ≠ empty and Found = false) do remove the top element from OPEN list and call it NODE; if NODE is the goal node, then FOUND = true else find SUCCs of NODE, if any: sort SUCCs by estimated cost from NODE to goal state and add them to the front of OPEN list: } /* end while */ if FOUND = true then return Yes otherwise return No: Stop

PROBLEMS IN HILL CLIMBING

Local Maximum:

It is a state better than all "its neighbours but not better than some other states which one far away.



Plateau | Flat Manimum :-

It is a flat area of the search space where all neighbouring states has the same value. Same Reusistic

Ridge :-

marima

It is an area of search space - that is higher than surrounding areas but that cannot be travoused by single moves in any one direction.

It is a special kind of local

After, selectory thus peak le A 14 cannot charge the forestern tomore's (B).

& BEST FIRST SEARCH (INFORMED, HEURISTIC) L> works as Greedy method. ALGORITHM: Let OPEN be a priority queue containing initial state if open' is empty return failure Node ← Remove-first (OPEN) if Node is a Goal then return the path from initial to Node else generate all silcussors of Node and put the newly generated Node into "OPEN" according to their & values L hewistic value END LOOP !

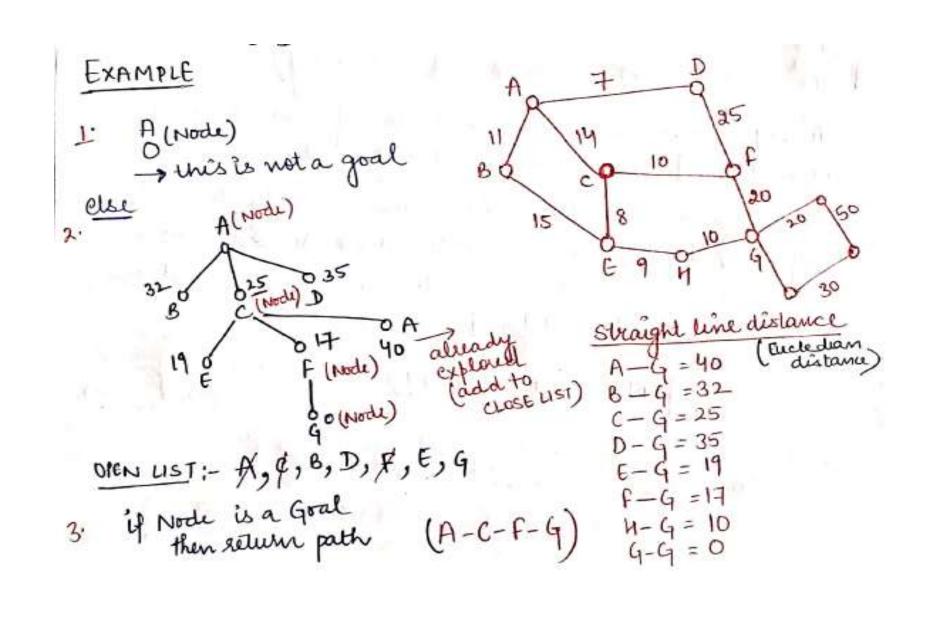
```
Algorithm (Best-First Search)
Input: START and GOAL states
Local Variables: OPEN. CLOSED. NODE. FOUND:
Output: Yes or No
Method:
• initialize OPEN list by root node: CLOSED = o: FOUND = false:

    while (OPEN ≠ oand FOUND = false) do

  • if the first element is the goal node, then FOUND = true else
    remove it from OPEN list and put it in CLOSED list;

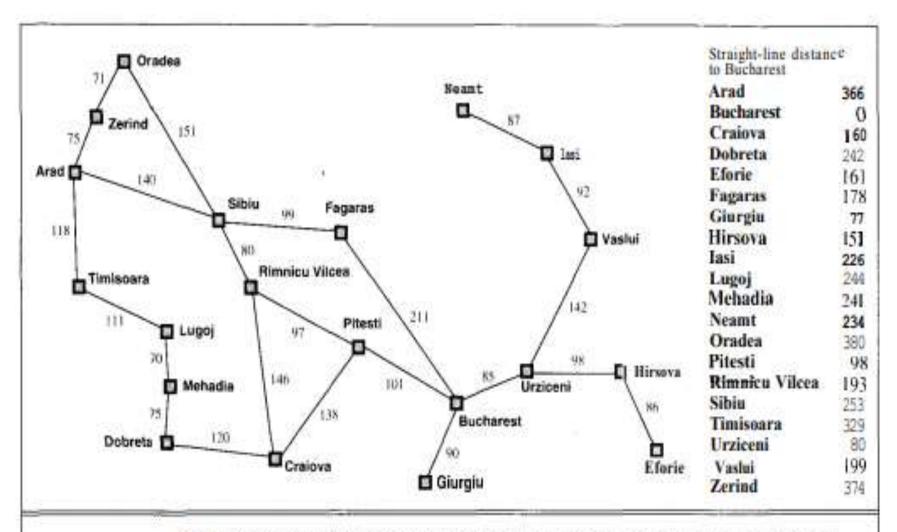
    add its successor, if any, in OPEN list;

  · sort the entire list by the value of some heuristic function that
    assigns to each node, the estimate to reach to the goal node:
. if FOUND - true then return Yes otherwise return No:
Stop
```

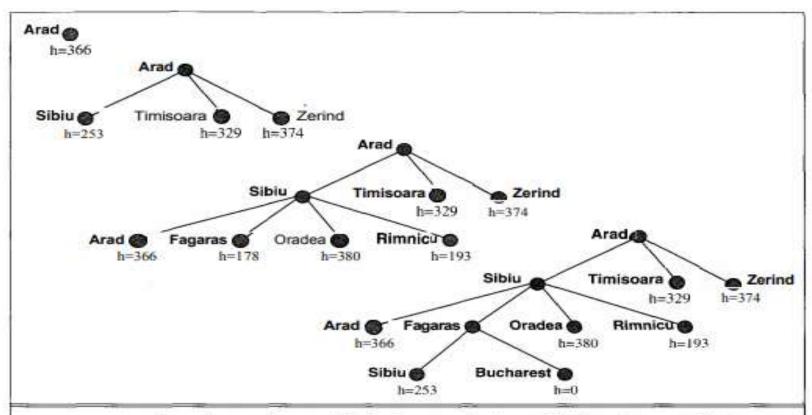


If you are working according to heuristic you will follow the path: A-C-F-G (cost is 44 without heuristic) But if you follow the poth: A-C-E-H-G (cost is 41)
this is less as So if we consider both the huristic and cost then there is changes to get optimal sol. I which is used by A* search.

Assignment



Map of Romania with road distances in km, and straight-line distances to Bucharest.



Stages in a greedy search for Bucharest, using the straight-line distance to Bucharest as the heuristic function h_{SLD}. Nodes are labelled with their h-values.

4. A* Algorithm: (Informed Searching) At algorithm is a combination of branch and bound, and best search methods. It uses a hewristic or evaluation function usually denoted by f(x) to determine the order in which the search for a node N is defined as follows: |f(N) = g(N) + R(N) | function g is a measure of cost of getting from stait function h is an estimate of additional cost of getting from current node N to goal node.

A* algorithm is called OR graph | true search algo.

A* algo incrementally searches all the soutes starting from the start node with it juids the shortest path to a goal.

```
Algorithm (A*) -
Input: START and GOAL states
Local Variables: OPEN, CLOSED, Best Node, SUCCs, OLD. FOUND:
Method:

    initialization OPEN list with start node; CLOSED= $\phi$: $g = 0$, $f = h$.

    while (OPEN ≠ φ and Found = false) do

  . remove the node with the lowest value of f from OPEN list and store
    it in CLOSED list. Call it as a Best_Node;

    if (Best_Node = Goal state) then FOUND = true else

                                                            Control
    · generate the SUCCs of the Best_Mode:
                                                            Litery
    · for each SUCC do
                                                           245(0):0

    establish parent link of SUCC: /* This link will Help to

     recover path once the solution is found */ 3,6 00 est lode compute g(SUCC) - g(Best Node) + cost of get3,6 00 est lode
        to SUCC:

    if SUCC ∈ OPEN then /* already being generated but not

        processed */
       · call the matched node as OLD and add it in the successor list
         of the Best Node:
       . ignore the SUCC node and change the parent of OLD, if required
         as follows:
         . if q(SUCC) < q(OLD) then make parent of OLD to be Best Node
         and change the values of g and f for DLD else ignore:

    If SUCC ∈ CLOSED then /* already processed */

        · call the matched node as OLD and add it in the list of the
         Best_Node successors:
       . ignore the SUCC node and change the parent of OLD, if required
          as follows:
          • if g(SUCC) < g(OLD) then make parent of OLD to be
            Best Node and change the values of g and f for OLD and
            propogate the change to OLD's children using depth first
           search else ignore:

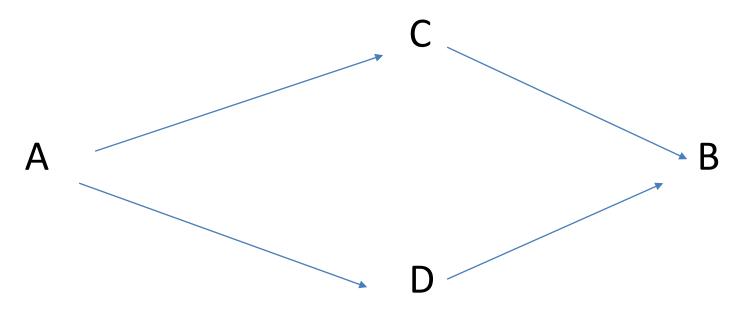
    If SUCC & OPEN or CLOSED

       · add it to the list of Best_Node's successors:

    compute f(SUCC) - g(SUCC) + h(SUCC);

        · put SUCC on OPEN list with its f value
      /* End while */
. if FOUND - true then return Yes otherwise return No:
· Stop
```

Cost fn g(n) -Distance A-C-B – 5 kms (5mns) H(n)-Traffic-5 mns ----10mns



Cost fn g(n) Distance A-D-B – 7 kms(7mns) H(n)-Traffic – 1mnt ---- 8mns Let us consider an example of eight puzzle again and solve it by using A^* algorithm. The simple evaluation function f(x) is defined as follows:

$$f(X) = g(X) + h(X)$$
, where

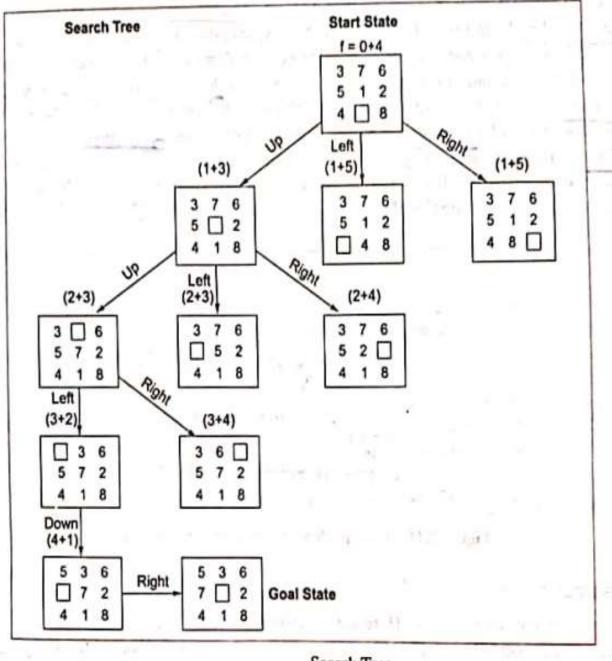
h(X) = the number of tiles not in their goal position in a given state X

g(X) = depth of node X in the search tree

Given

| Start State | | | - |
|-------------|---|-----|---|
| 3 | 7 | 6 | |
| 5 | 1 | · 2 | - |
| 4 | | 8 | |





Search Tree

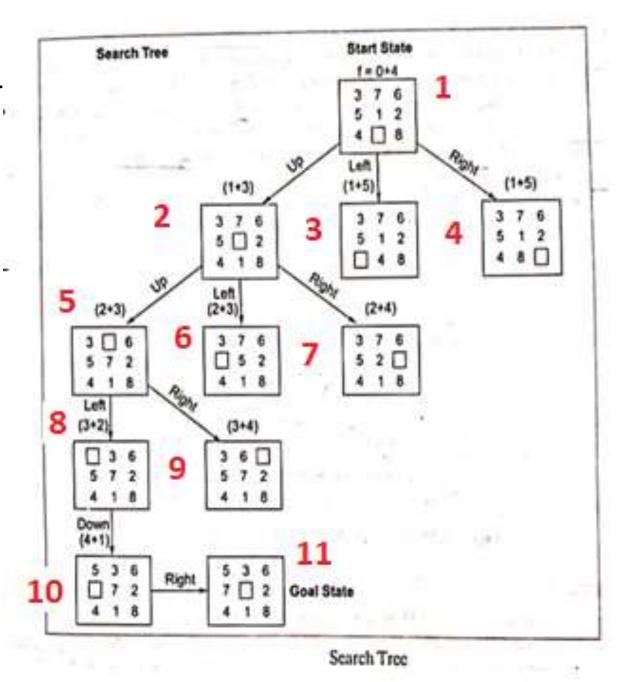
Open List-3-1(6),4-1(6),6-2(5),7-2(6),9-5(7), 12-10(7)

Closed List- 1(4), 2-1(4), 5-2(5), 8-5(5), 10-8(5)

BestNode- 11-10(5)

SUCCs-

Old-



SB→e

4+12+4.

= 20

SB->

4+5+11

HOW TO MAKE A* ADMISSIBLE: Ly acceptable

- A* star search algorithm is admissible, if for any graph it always terminates in an optimal path from start state to goal state, if path exists.
- → If hewistic function "h" underestimates the actual value from current state to goal state, then "it bounds to give an optimal solution & hence is called admissible function.
- -> So, A* always terminates with oftimal both in case h is an admissible hewustic function

estimated actual optimal h(m) < h*(m) Underestimation h(A)=30 h(B) = 20 $A* \rightarrow f(m) = g(m) + h(m)$ f(A) = g(A) + R(A) = 200+ 30 = 230 f(B) = g(B) + h(B) = 200+20 = 220 7(9) = g(9)+ R(9) 230 = 200+50 +0 = 250 After reaching goal you observe that A has I value 230 which is less than G=250, so now explore A f(q)= g(q)+h(q) = 200+40+0 - value is less as compared to seg path So if herestic function h underestimates the actual value from current state to goal state, then it bounds to give an optimal solution

$$h(n) \ge h^*(n)$$

$$f(A) = g(A) + h(A)$$

= 200+80 = 280

$$f(B) = g(B) + R(B)$$

= 200 + 70 = 270

$$f(9) = g(9) + h(9)$$

= $a50 + 0 = a50$