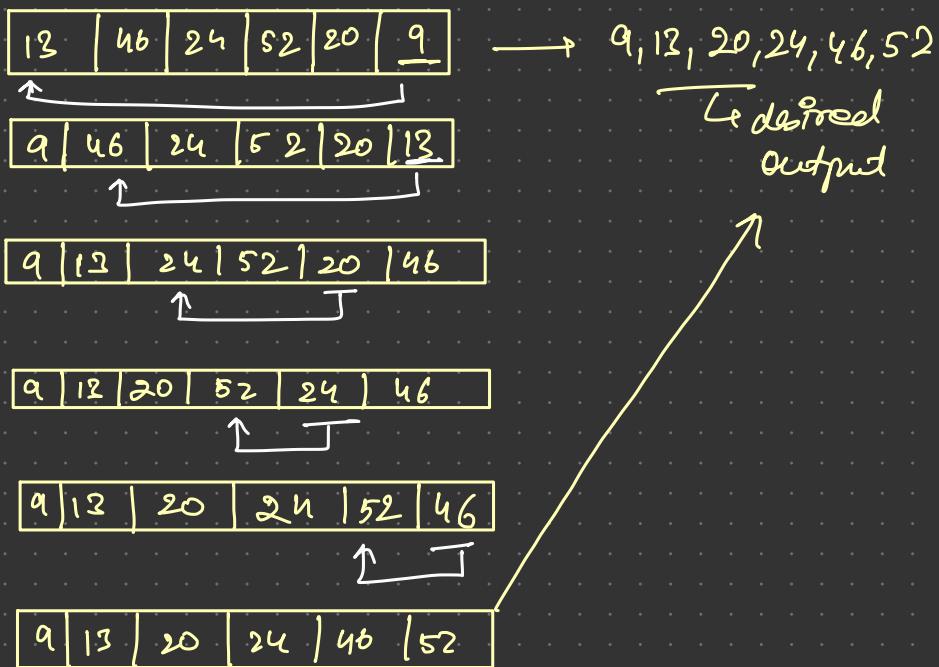


# Sorting Algorithms

## → Selection Sort →

- As name suggest "selection"
- Selection of minimum element
- swapping minimum with desired position in the array.



Time complexity  $\rightarrow O(n^2)$

Auxiliary space  $\rightarrow O(1)$

→ It is not stable as the relative order of same key is not maintained, it can be made stable.

→ It is an in-place algorithm i.e. no extra space is used.

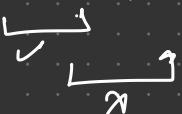
```
void selectionSort(vector<int>& nums)
{
    //THIS ONE IS UNSTABLE
    int k=0;
    for(int i=0;i<nums.size();i++)
    {
        int mini=i;
        for(int j=i+1;j<nums.size();j++)
        {
            if(nums[j]<nums[mini])
                mini=j;
        }
        swap(nums[k++],nums[mini]);
    }
}
void selectionSortStable(vector<int>& nums)
{
    for(int i=0;i<nums.size();i++)
    {
        int min=i;
        for(int j=i+1;j<nums.size();j++)
        {
            if(nums[j]<nums[min])
                min=j;
        }
        int key=nums[min];
        for(int j=min;j>i;j--)
        {
            nums[j]=nums[j-1];
        }
        nums[i]=key;
    }
}
```

## Bubble Sort

— Traverse from left and compare adjacent elements and higher one is placed at right.

— Two adjacent are sorted at one step

13	46	24	52	20	9
----	----	----	----	----	---



13	24	46	52	20	9
----	----	----	----	----	---



13	24	46	20	50	9
----	----	----	----	----	---



13	24	46	20	9	50
----	----	----	----	---	----



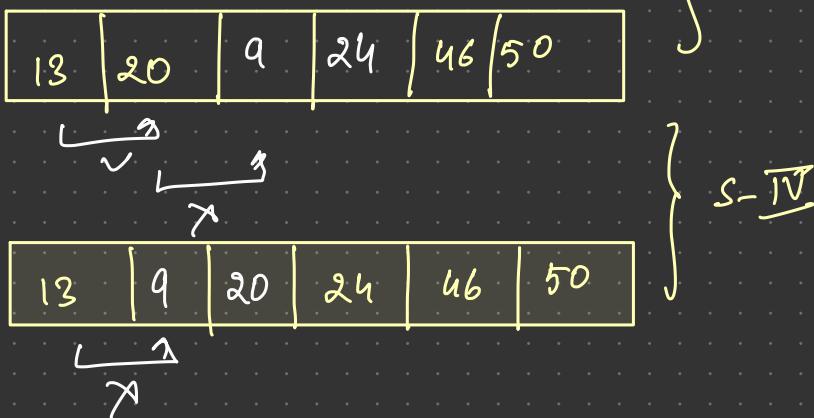
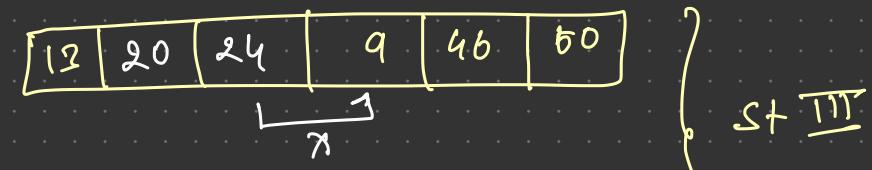
13	24	20	46	9	50
----	----	----	----	---	----



13	24	20	9	46	50
----	----	----	---	----	----



} S-II



[9 | 13 | 20 | 24 | 46 | 50]  $\rightleftharpoons$  sorted at step-IV

Time complexity  $\rightarrow O(N^2)$   
Auxiliary space  $\rightarrow O(1)$

→ Bubble sort is stable

→ Bubble sort is inplace

```

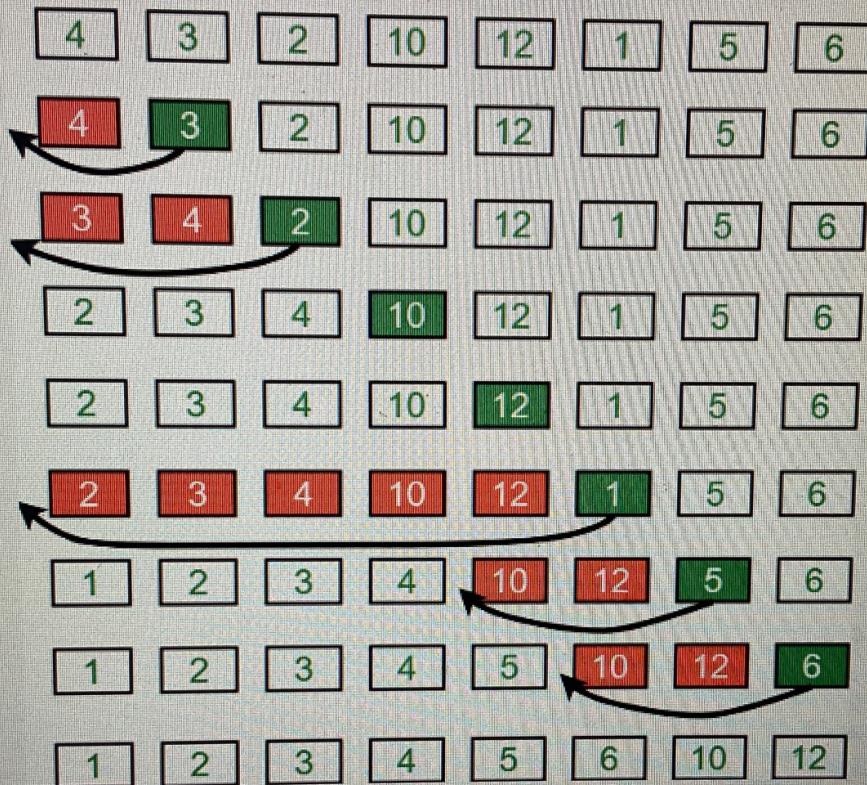
6
7 void bubbleSort(vector<int>& nums)
8 {
9     for(int i=0;i<nums.size();i++)
10    {
11        for(int j=0;j<nums.size()-i-1;j++)
12        {
13            if(nums[j]>nums[j+1])
14            swap(nums[j],nums[j+1]);
15        }
16    }
17 }
```

## Insertion Sort

→ simple sorting algorithm that works similar to the way sorting playing cards in your hands.

→ Array is splitted into sorted and unsorted part

### Insertion Sort Execution Example



## Time Complexity

- Worst case  $\rightarrow O(N^2)$
- Average case  $\rightarrow O(N^2)$
- Best case  $\rightarrow O(N)$

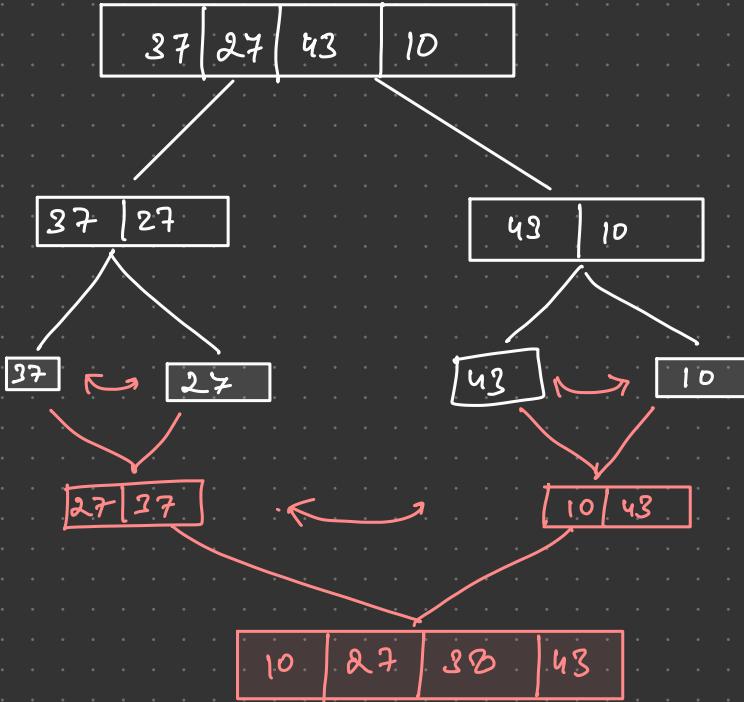
Auxiliary space  $\rightarrow O(1)$

- Follows incremental approach
- Insertion sort is in-place sorting algorithm
- Insertion sort is stable sorting algorithm

```
4 void insertionSort(vector<int>& nums)
5 {
6     for(int i=0;i<=nums.size()-1;i++)
7     {
8         int j=i;
9         while(j>0 && nums[j-1]>nums[j])
10        {
11            swap(nums[j-1],nums[j]);
12            j--;
13        }
14    }
15 }
```

# Merge Sort

- Divide and conquer algorithm
- Divide
  - ↓
  - sort
  - ↓
  - merge



$$T(n) = 2T(n/2) + \Theta(n)$$

Time complexity  $\rightarrow O(n \log(n))$

Auxiliary Space  $\rightarrow O(N)$

- stable sorting algorithm
- guaranteed performance
- parallelizable

— not recommended for small datasets.

```
30 void mergeSort(vector<int>& nums,int left,int right)
31 {
32     if(left>=right)
33         return;
34     int mid=(left+right)/2;
35     mergeSort(nums,left,mid);
36     mergeSort(nums,mid+1,right);
37     merge(nums,left,mid,right);
38 }
39
40 void merge(vector<int>& nums,int left,int mid,int right)
41 {
42     int size=right-left+1;
43     vector<int> temp(size,0);
44     int i=left;
45     int j=mid+1;
46     int k=0;
47     while(i<=mid && j<=right)
48     {
49         if(nums[i]<nums[j])
50             temp[k++]=nums[i++];
51         else
52             temp[k++]=nums[j++];
53     }
54     while(i<=mid)
55         temp[k++]=nums[i++];
56     while(j<=right)
57         temp[k++]=nums[j++];
58
59     for(int i=left;i<=right;i++)
60     {
61         nums[i]=temp[i-left];
62     }
63 }
```

## - Quick Sort -

- Algorithm based on divide and conquer
- Picks a pivot and partition the given array around picked pivot by placing it in its correct position.



## Time Complexity →

Best case  $\rightarrow \mathcal{O}(N \cdot \log(N))$ .

Average Case  $\rightarrow \Theta(N \cdot \log(N))$

Worst Case  $\rightarrow \mathcal{O}(N^2)$ .

Auxiliary space  $\rightarrow \mathcal{O}(1)$

→ Not a stable algorithm

```
22
23 void quickSort(vector<int>& nums,int low,int high)
24 {
25     if(low<high)
26     {
27         int pivot_index=quick(nums,low,high);
28         quickSort(nums,low,pivot_index-1);
29         quickSort(nums,pivot_index+1,high);
30     }
31 }
```

```

4
5  int quick(vector<int>& nums,int low,int high)
6  {
7      int pivot=nums[low];//starting element as pivot
8      int i=low;
9      int j=high;
10     while(i<j)
11     {
12         while(i<high && nums[i]<=pivot)
13             i++;
14         while(j>low && nums[j]>pivot)
15             j--;
16         if(i<j)
17             swap(nums[i],nums[j]);
18     }
19     swap(nums[low],nums[j]); //swap to its correct position
20     return j;
21 }
22

```

→ Heap sort → - sorting technique based on binary heap data structure.

Auxiliary space →  $O(1) \cong O(\log N)$   
 Time complexity →  $O(N \cdot \log(N))$

- Typically not stable
- In place sorting algorithm.

recusive  
stack  
space

```
5 void heapSort(vector<int>& nums)
6 {
7     int n=nums.size();
8     for(int i=1;i<n;i++)
9     {
10         int ch=i;
11         while(ch>0)
12         {
13             int par=(ch-1)/2;
14             if(nums[par]<nums[ch])
15             {
16                 swap(nums[par],nums[ch]);
17             }
18             else
19                 break;
20
21             ch=par;
22         }
23     }
24     int size=n;
25     while(size>1){
26         swap(nums[0],nums[size-1]);
27         size--;
28         int par=0;
29         int lef=par*2+1;
30         int rig=par*2+2;
31         while(lef<size)
32         {
33             int mini=par;
34             if(nums[mini]<nums[lef])
35                 mini=lef;
36             if(rig<size && nums[mini]<nums[rig])
37                 mini=rig;
38             if(mini==par)
39                 break;
40
41             swap(nums[mini],nums[par]);
42             par=mini;
43             lef=2*par+1;
44             rig=2*par+2;
45         }
46     }
47 }
```