DAFNY

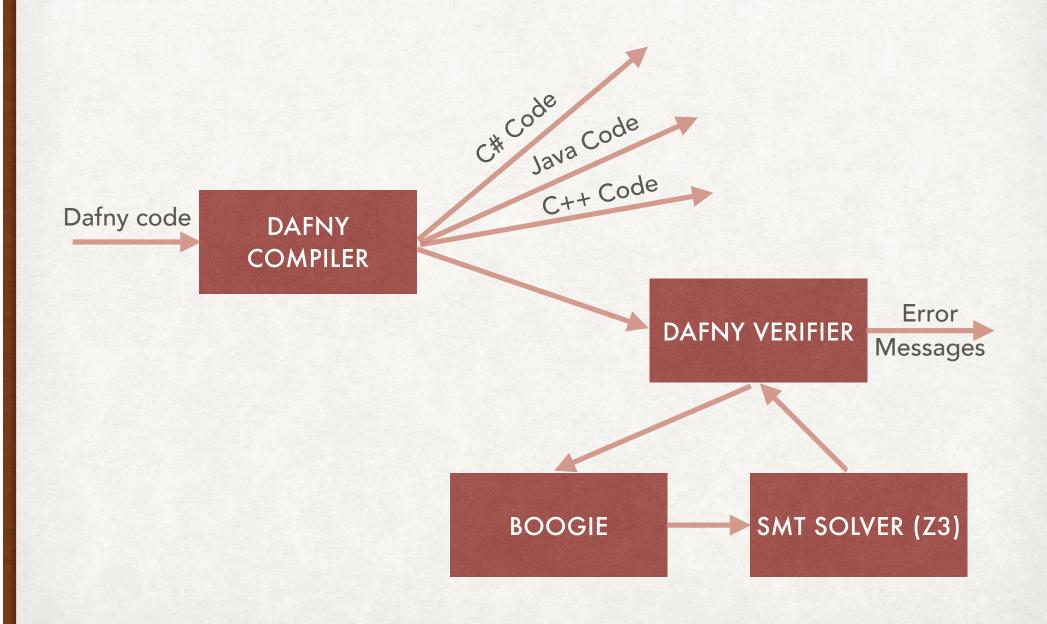
INTRODUCTION

- Dafny is a programming language, verifier and compiler, all rolled together into one.
- Dafny uses SMT solver (Z3) to automatically prove correctness.
- Highly sophisticated specification language
 - Pre-conditions and Post-conditions for methods
 - Loop Invariants
 - Assertions
 - Lemmas and Theorems

DAFNY - PROGRAMMING LANGUAGE

- Dafny hosts two separate sub-languages
- Language for the executable code
 - Imperative, Object-oriented
 - Methods, loops, if-then-else, arrays, classes, inductive datatypes,...
- Specification Language
 - Functional
 - Pure functions, predicates, algebraic datatypes, sets, sequences, lemmas...

DAFNY TOOL



DAFNY BASICS: METHODS

- Parameters and return values must be typed.
- Return values are named.
- Multiple return values are also allowed.

```
method Abs(x: int) returns (y: int)
{
    ...
}
```

```
method MultipleReturns(x: int, y: int)
returns (more: int, less: int)
{
    ...
}
```

DAFNY BASICS: METHODS

```
method MultipleReturns(x: int, y: int)
returns (more: int, less: int)
{
   more := x + y;
   less := x - y;
}
```

- To return a value, assign it to the named return variable.
 - In case of a single return value, the 'return' statement can also be used.
- Assignments use ':='.
- Compound statements such as if..then..else, while loops are available.

DAFNY BASICS: PRE AND POST-CONDITIONS

```
method MultipleReturns(x: int, y: int)
returns (more: int, less: int)
   ensures less < x
   ensures x < more
{
   more := x + y;
   less := x - y;
}</pre>
```

- Post-conditions specified using 'ensures'.
 - Are the post-conditions satisfied in the above example?

DAFNY BASICS: PRE AND POST-CONDITIONS

```
method MultipleReturns(x: int, y: int)
returns (more: int, less: int)
   requires y > 0
   ensures less < x
   ensures x < more
{
   more := x + y;
   less := x - y;
}</pre>
```

- Post-conditions specified using 'ensures'.
 - Are the post-conditions satisfied in the above example?
- Pre-conditions specified using 'requires'
 - Dafny will output 'Verified'

DAFNY BASICS: ASSERTIONS

- Local variables declared using 'var'.
- Methods calls are not allowed inside assert conditions.
 - Will the assert in 'Testing' method succeed?

```
method Abs(x: int) returns (y: int)
   ensures y >= 0
   if x < 0
     { return -x; }
   else
     { return x; }
method Testing()
  var v := Abs(3);
   assert v == 3;
```

DAFNY BASICS: ASSERTIONS

- Local variables declared using 'var'.
- Methods calls are not allowed inside assert conditions.
 - Will the assert in 'Testing' method succeed?

```
method Abs(x: int) returns (y: int)
   ensures y >= 0
   ensures x >= 0 ==> y == x
   ensures x < 0 ==> y == -x
   if x < 0
     { return -x; }
   else
     { return x; }
method Testing()
  var v := Abs(3);
   assert v == 3;
```

DAFFY BASICS: FUNCTIONS

```
function abs(x: int): int
{
   if x < 0 then -x else x
}</pre>
```

- Pure mathematical functions
 - Cannot write to memory
 - Body is a single expression
 - Single return value
 - Not compiled and executed
- Can only be used inside pre-conditions, post-conditions, assertions and loop invariants

DAFNY BASICS: FUNCTIONS

```
function abs(x: int): int
   if x < 0 then -x else x
method Abs(x: int) returns (y: int)
   ensures y == abs(x)
   if x < 0
     { return -x; }
   else
    { return x; }
method Testing()
   var v := Abs(3);
   assert v == 3;
```

DAFFY BASICS: LOOP INVARIANTS

```
var i := 0;
  while i < n
     invariant 0 <= i
  {
     i := i + 1;
}</pre>
```

- Specified using the keyword 'invariant'
 - Must hold upon entering the loop
 - After every iteration of the loop body
 - Dafny will automatically check the conditions

DAFFY BASICS: LOOP INVARIANTS

```
method ComputeFib(n: nat) returns (b: nat) function fib(n: nat): nat
   ensures b == fib(n)
                                                if n == 0 then 0 else
{
                                                if n == 1 then 1 else
   if n == 0 { return 0; }
                                                   fib(n - 1) + fib(n - 2)
   var i: int := 1;
   var a := 0;
       b := 1;
   while i < n
      invariant 0 < i <= n
      invariant a == fib(i - 1)
      invariant b == fib(i)
      a, b := b, a + b;
      i := i + 1;
```

DAFFY BASICS: TERMINATION

- Dafny will try to prove that every loop terminates
 - Reports an error if it can't prove termination.
- Dafny will guess an expression which decreases with every iteration and has a lower bound.
 - This can be explicitly specified using the 'decreases' annotation.
 - Use the annotation 'decreases *' if we don't want Dafny to prove termination.

```
while 0 < i
    invariant 0 <= i
    decreases i
    {
        i := i - 1;
}</pre>
```

DAFFY BASICS: TERMINATION

- Dafny will try to prove that every loop terminates
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- Dafny will guess an expression which decreases with every iteration and has a lower bound.
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 - Use the annotation 'decreases *' if we don't want Dafny to prove termination.

```
while 0 < i
    invariant 0 <= i
    decreases i
    {
        i := i - 1;
}</pre>
```

```
while i < n
    invariant 0 <= i <= n
    decreases n - i
{
       i := i + 1;
}</pre>
```

DAFFY BASICS: ARRAYS

- Built in data type 'array<T>'
 - Also built in field 'Length'.
- Dafny will automatically check that arrays are accessed within the bounds.

```
method M(x: int)
   var a := new int[10];
   var b := a[x]; //Error
   if 0 \le x < 10
     { b := a[x]; //OK }
method M(x: int, c: array<int>)
 requires 0 \ll x \ll 10
   var a := new int[10];
   var b := a[x]; //OK
   If 0 <= x < 10
     \{ b := c[x]; //Error \}
```

DAFNY BASICS: QUANTIFIERS

```
method Find(a: array<int>, key: int) returns (index: int)
   ensures 0 <= index ==> index < a.Length && a[index] == key
   ensures index < 0 ==> forall k :: 0 <= k < a.Length ==> a[k] != key
   index := 0;
  while index < a.Length</pre>
      invariant 0 <= index <= a.Length</pre>
      invariant forall k :: 0 <= k < index ==> a[k] != key
      if a[index] == key { return; }
      index := index + 1;
   index := -1;
```

 Quantifiers can be used in pre/post-conditions, assertions, invariants using the 'forall' keyword.

DAFNY BASICS: LEMMAS

```
method ComputePow2(n: nat) returns (p:nat) function pow2(n: int): int
 ensures p = pow2(n)
                                              requires 0 <= n
                                             {
  if n = 0
                                                  if n == 0 then 1
   { p := 1; }
                                               else
  else if n \% 2 = 0
                                                  2*pow2(n-1)
    \{ p := ComputePow2(n / 2); \}
                                             }
      p := p * p; }
  else
    \{ p := ComputePow2(n-1); \}
      p := 2 * p;
```

DAFNY BASICS: LEMMAS

```
method ComputePow2(n: nat) returns (p:nat) function pow2(n: int): int
ensures p = pow2(n)
                                              requires 0 <= n
 if n = 0
                                                  if n == 0 then 1
   { p := 1; }
                                               else
  else if n \% 2 = 0
                                                  2*pow2(n-1)
    \{ p := ComputePow2(n / 2); \}
                                             }
      p := p * p;
      Lemma(n); }
                                       lemma Lemma(n : nat)
  else
                                          requires n % 2 == 0
    \{ p := ComputePow2(n-1); \}
                                          ensures pow2(n) == pow2(n/2)
      p := 2 * p;
                                       * pow2(n/2)
                                          if n != 0 then Lemma(n-2);
```

DAFNY: RESOURCES

- https://rise4fun.com/Dafny/tutorial/Guide
- https://github.com/dafny-lang/dafny