ABSTRACT FORWARD PROPAGATE

KILDALL'S ALGORITHM

```
AbstractForwardPropagate(Γ_c,P)
  S := {l_0};
  μ̂(l_0) := α(P);
  μ̂(l) := ⊥, for l ∈ L\{l_0};
  while S ≠ ∅ do{
     l := Choose S;
     S := S \ {l};
     foreach (l,c,l′) ∈ T do{
        F := ŝp(μ̂(l),c);
        if ¬(F ≤ μ̂(l′)) then{
           μ̂(l′) := μ̂(l′) ⊔ F;
           S := S ∪ {l′};
        }
     }
  }
```

- Does this algorithm actually calculate the abstract JOP?
- What are the conditions under which it is guaranteed to terminate?

A Unified Approach to Global Program Optimization. Gary Kildall. POPL 1973.

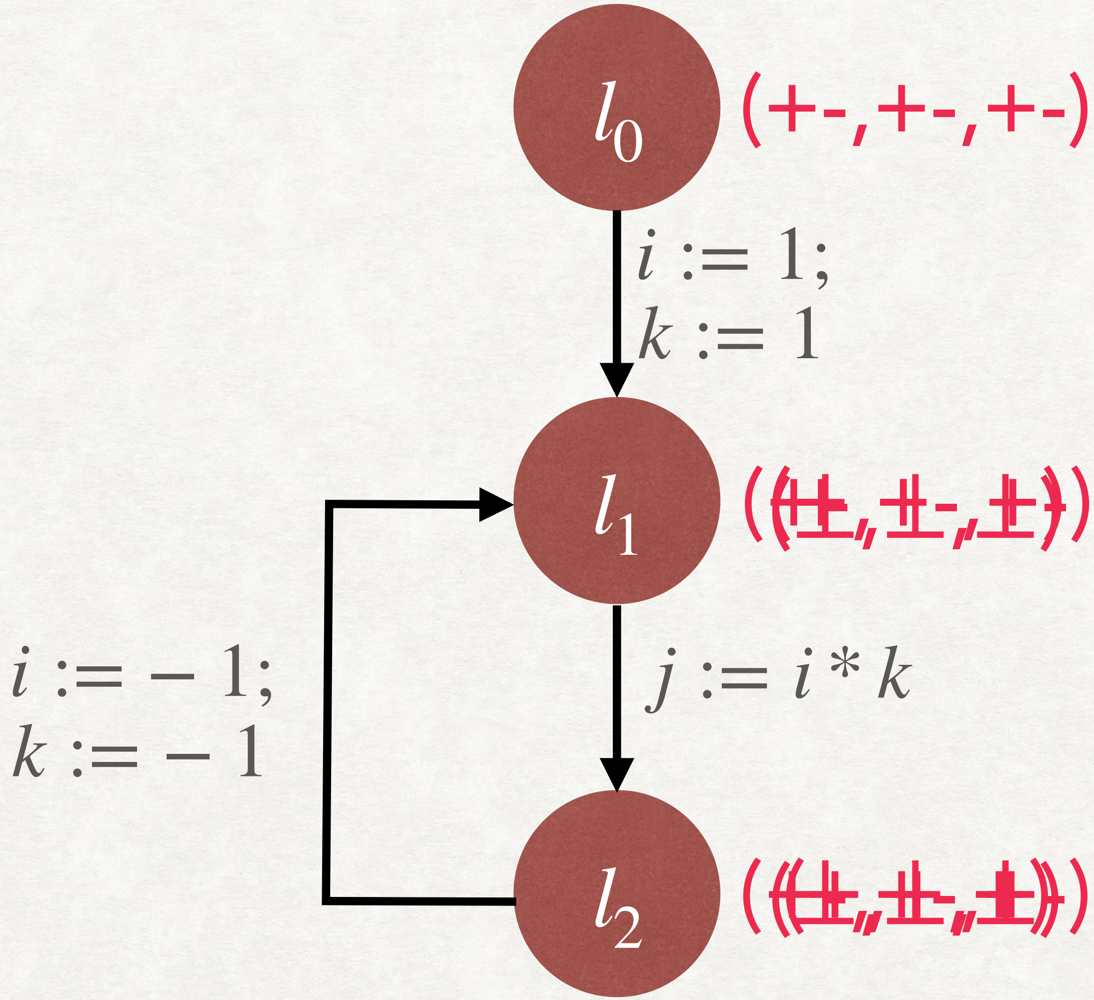
ABSTRACT FORWARD PROPAGATE

KILDALL'S ALGORITHM

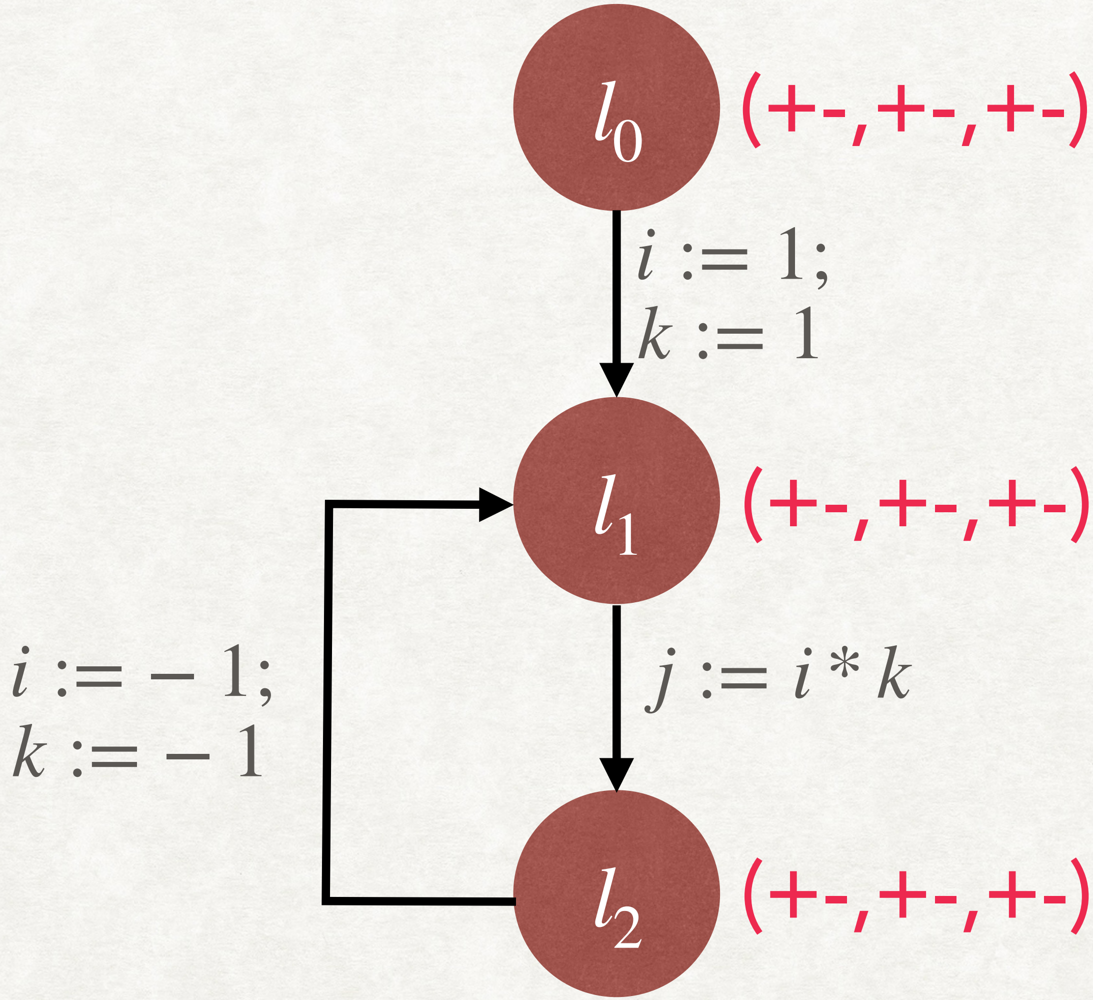
```
AbstractForwardPropagate(Γ_c,P)
  S := {l_0};
  μ̂_K(l_0) := α(P);
  μ̂_K(l) := ⊥, for l ∈ L\{l_0};
  while S ≠ ∅ do{
     l := Choose S;
     S := S \ {l};
     foreach (l, c, l') ∈ T do{
        F := f̂_c(μ̂_K(l));
        if ¬(F ≤ μ̂_K(l')) then{
           μ̂_K(l') := μ̂_K(l') ⊔ F;
           S := S ∪ {l'};
        }
     }
  }
```

A Unified Approach to Global Program Optimization. Gary Kildall. POPL 1973.

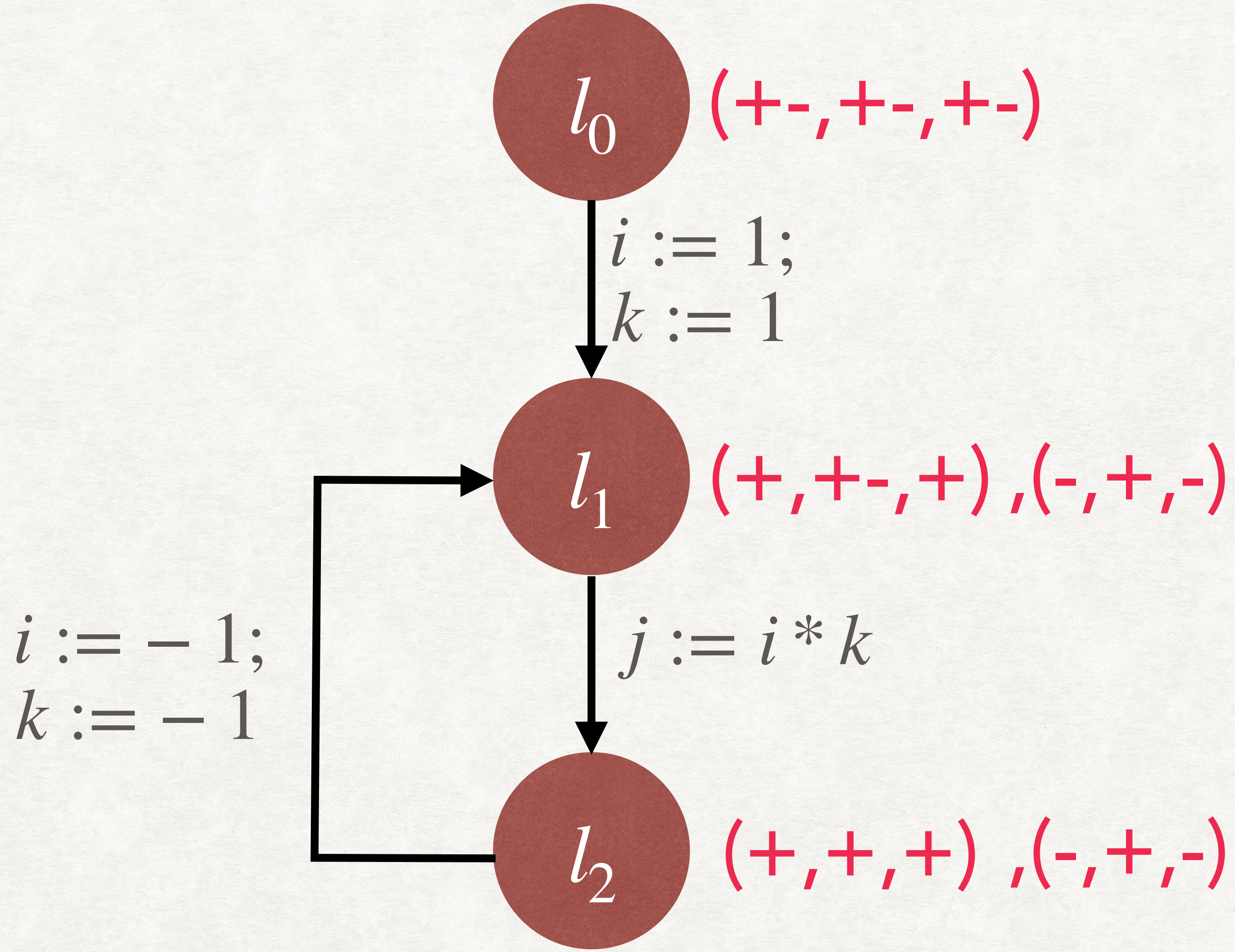
EXAMPLE - KILDALL'S ALGORITHM



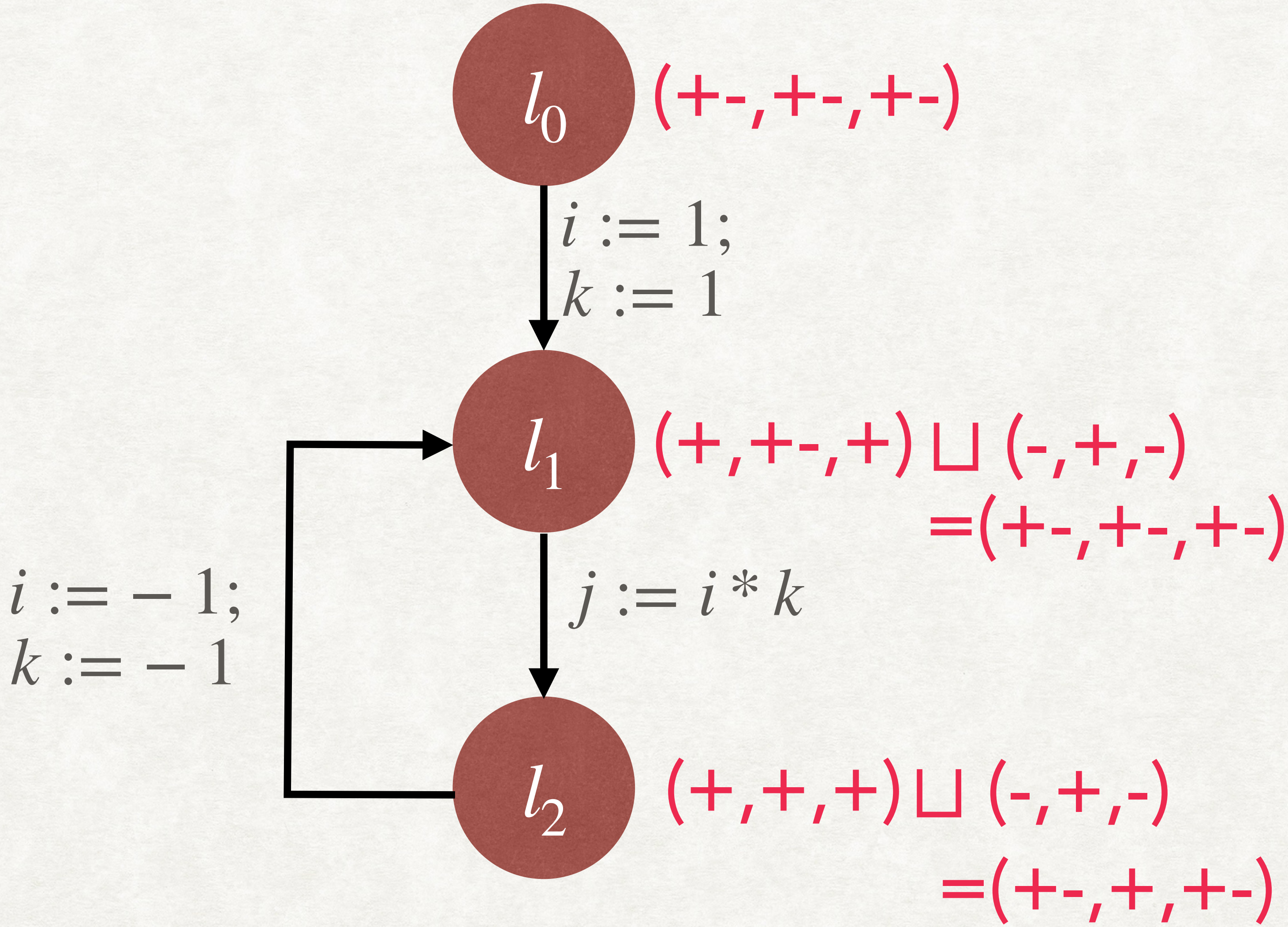
EXAMPLE - KILDALL'S ALGORITHM



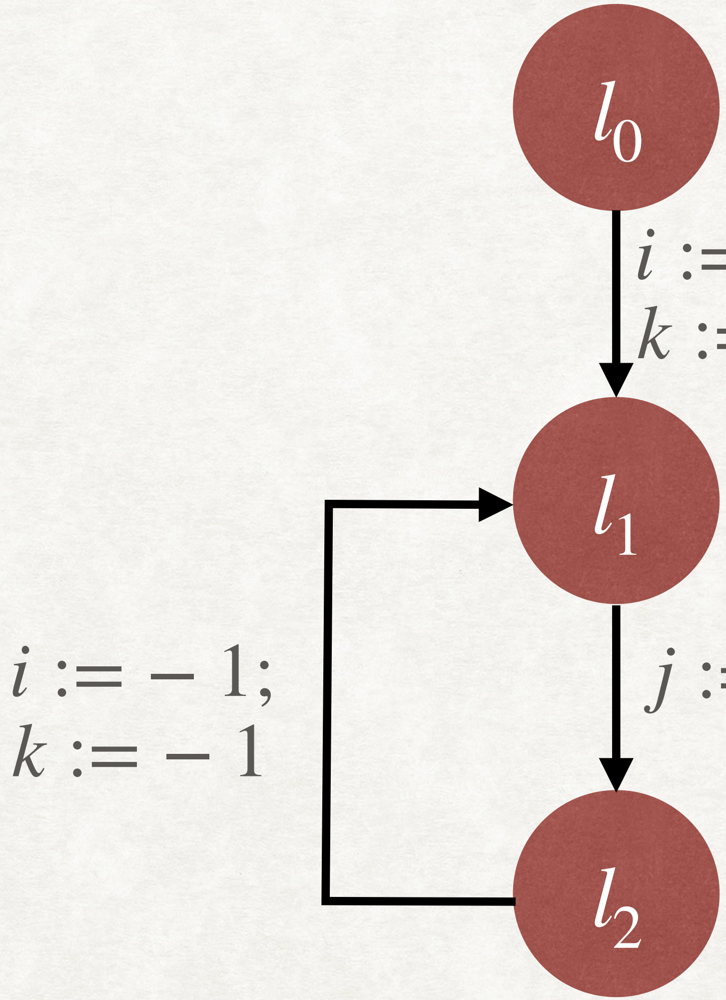
EXAMPLE - ABSTRACT JOP



EXAMPLE - ABSTRACT JOP



EXAMPLE - KILDALL VS ABSTRACT JOP



	$\hat{\mu}_K$	$\hat{\mu}$
l_0	(+-,+-,+-)	(+-,+-,+-)
l_1	(+-,+-,+-)	(+-,+-,+-)
l_2	(+-,+-,+-)	(+-,+,+-)

$\hat{\mu}_K \neq \hat{\mu}$: This is because Kildall's Algorithm applies join eagerly

We will prove that $\hat{\mu}_K \geq \hat{\mu}$

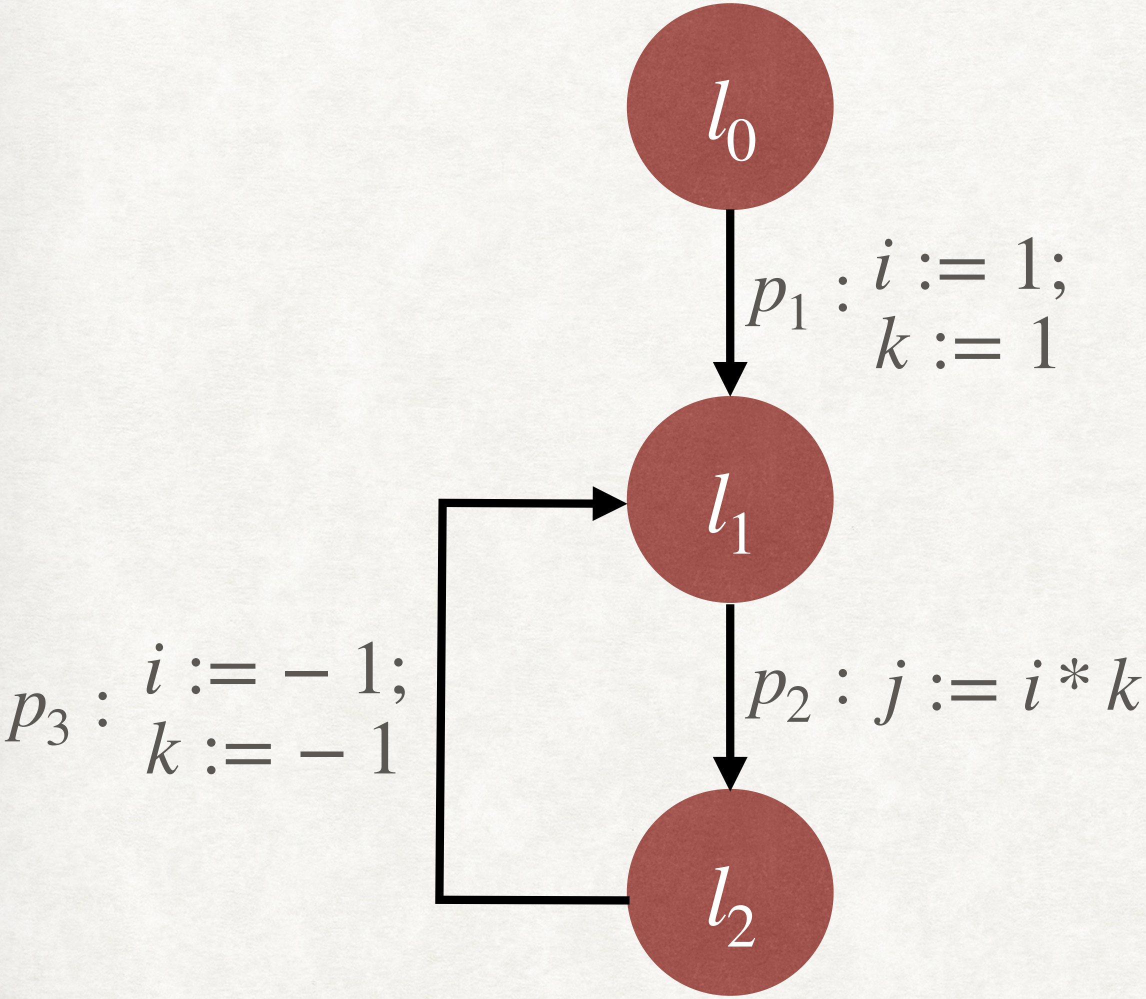
PROPERTIES OF KILDALL'S ALGORITHM

1. The values computed using Kildall's algorithm are an over-approximation of the abstract JOP, if the underlying AI framework is monotonic.
2. In general, Kildall's algorithm computes the least solution to a system of equations.
3. If the abstract domain satisfies the ascending chain condition, then Kildall's algorithm is guaranteed to terminate.

DATAFLOW EQUATIONS

- Program $\Gamma_c = (V, L, l_0, l_e, T)$ induces a system of data flow equations:
  - $X_{l_0} = d_0$
  - For all other locations $l \in L \backslash \{l_0\},\ X_l = \bigsqcup_{(l',c,l) \in T} \hat{f}_c(X_{l'})$
- For collecting semantics, replace d_0 with c_0, $\sqcup$ with $\cup$ and $\hat{f}_c$ with f_c.

EXAMPLE - DATAFLOW EQUATIONS



$$X_{l_0} = d_0$$

$$X_{l_1} = \hat{f}_{p_1}(X_{l_0}) \sqcup \hat{f}_{p_3}(X_{l_2})$$

$$X_{l_2} = \hat{f}_{p_2}(X_{l_1})$$

DATAFLOW EQUATIONS AS FUNCTION

- Consider the 'vectorised' lattice $(\bar{D}, \bar{\leq})$, where $\bar{D} = D^{|L|}$.
  - $\bar{d} \,\bar{\leq}\, \bar{d}' \Leftrightarrow \forall l \in L \,.\, \bar{d}(l) \leq \bar{d}'(l)$
  - Homework: Prove that if $(D, \leq)$ is a complete lattice, then $(\bar{D}, \bar{\leq})$ is also a complete lattice.
- We can view the data flow equations as a function $\bar{f} : \bar{D} \to \bar{D}$:
  - $(\bar{f}(\bar{d}))(l_0) = d_0$
  - $(\bar{f}(\bar{d}))(l) = \bigsqcup_{(l',c,l) \in T} \hat{f}_c(\bar{d}(l'))$