

# WEAKEST-PRECONDITION

## SYMBOLIC EXECUTION IN THE BACKWARD DIRECTION

- Given an error condition or a post-condition, propagate the condition backwards through the program.
- Given a set of states  $S$  and a command  $c$ , the weakest precondition operator  $wp(S, c)$  consists of all states that would always lead to a state in  $S$  after executing  $c$ .

$$wp(S, c) \triangleq \{\sigma \mid \forall \sigma'. (\sigma, c) \hookrightarrow^* (\sigma', \text{skip}) \rightarrow \sigma' \in S\}$$

Equivalently,  $\sigma \in wp(S, c) \Leftrightarrow \forall \sigma'. (\sigma, c) \hookrightarrow^* (\sigma', \text{skip}) \rightarrow \sigma' \in S$

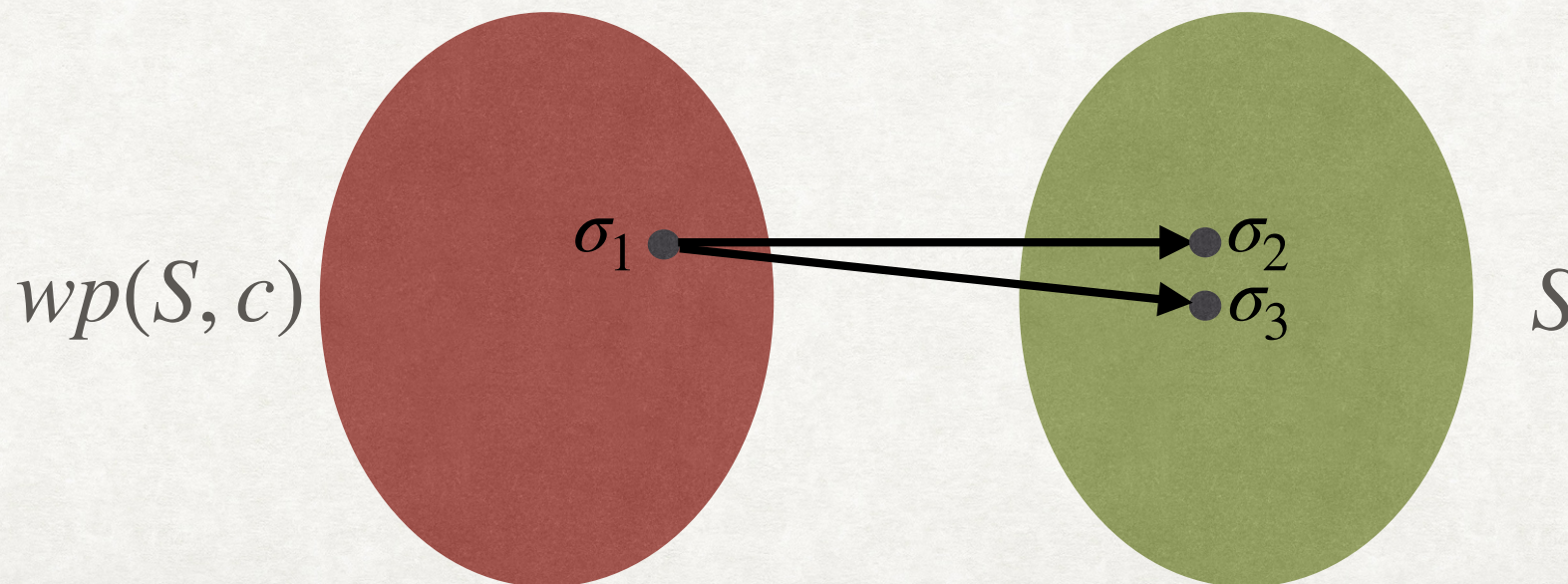


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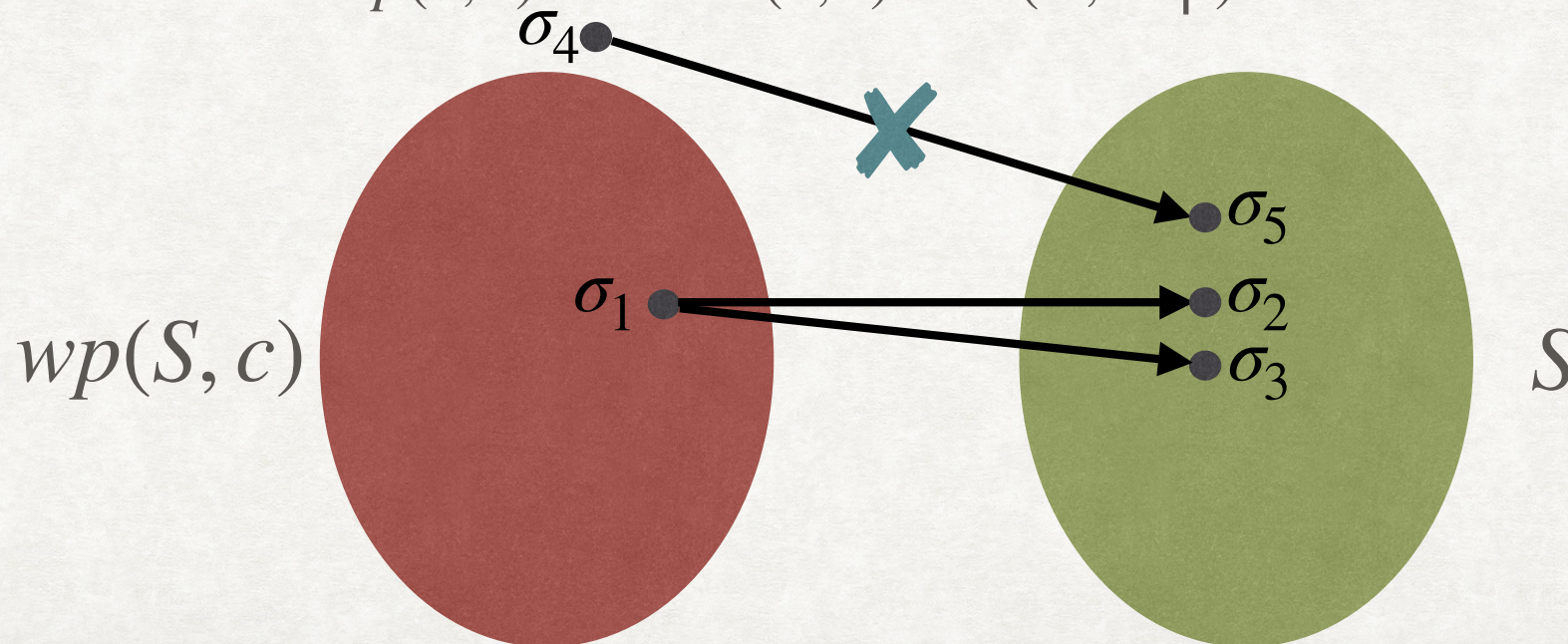


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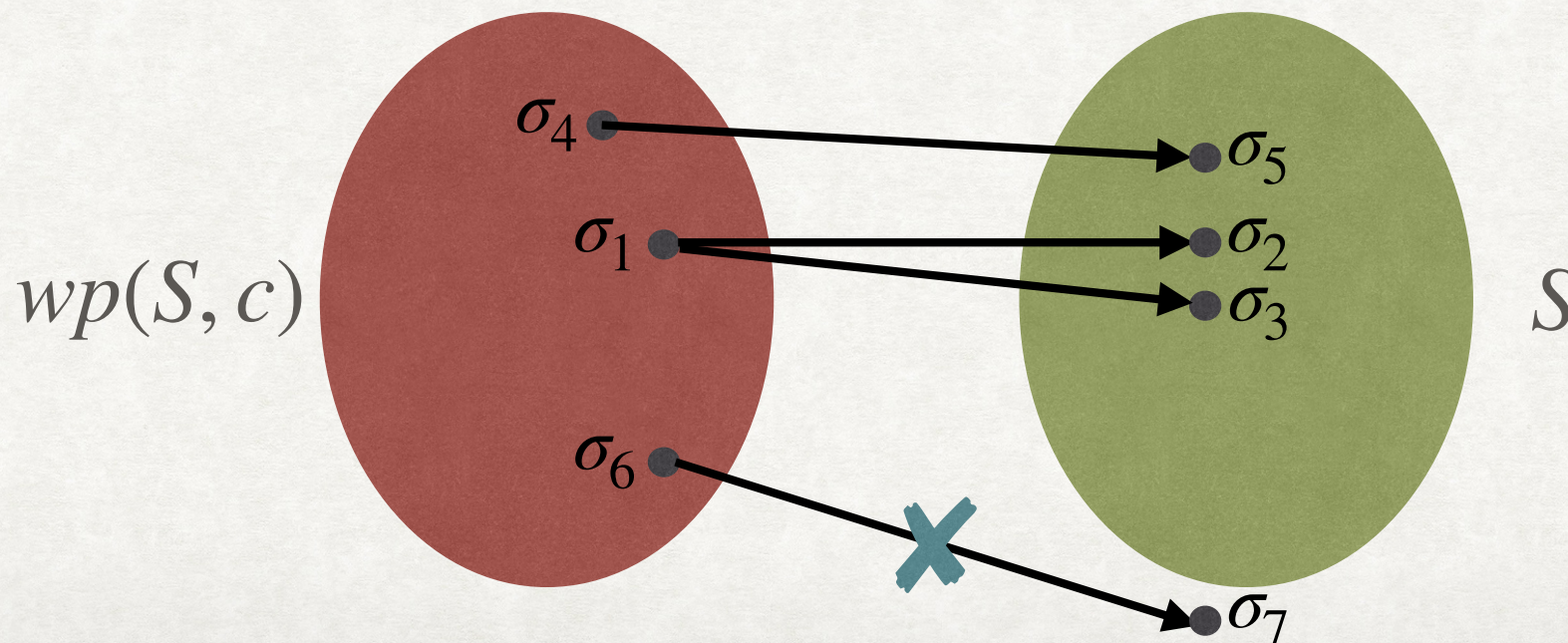


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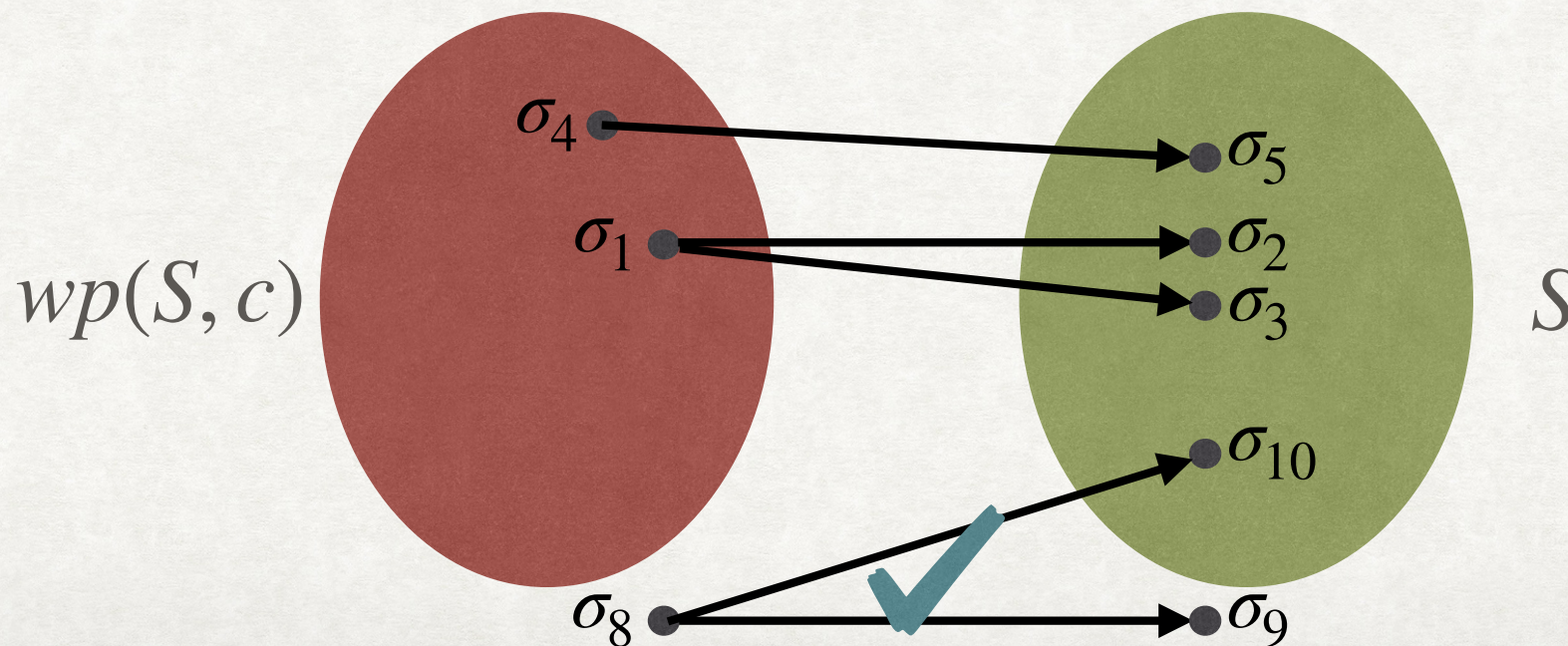


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$$wp(S, c) \triangleq \{\sigma \mid \forall \sigma'. (\sigma, c) \hookrightarrow^* (\sigma', \text{skip}) \rightarrow \sigma' \in S\}$$

- We can use a FOL formula  $F$  to represent a set of states.
- The symbolic weakest pre-condition operator can be defined as:

$$\sigma \models wp(F, c) \Leftrightarrow \forall \sigma'. (\sigma, c) \hookrightarrow^* (\sigma', \text{skip}) \rightarrow \sigma' \models F$$

- We now use the symbolic FOL semantics ( $\rho$ ) for individual commands:

$$wp(F, c) \triangleq \forall V'. \rho(c) \rightarrow F[V'/V]$$



# WEAKEST PRE-CONDITION

## EXAMPLES

$$\begin{aligned} wp(x > 10, x := x + 1) &\triangleq \forall x'. x' = x + 1 \rightarrow x' > 10 \\ &\equiv x + 1 > 10 \equiv x > 9 \end{aligned}$$



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$$\begin{aligned}wp(x > 10, x := x + 1) &\triangleq \forall x'. x' = x + 1 \rightarrow x' > 10 \\ &\equiv x + 1 > 10 \equiv x > 9\end{aligned}$$

$$wp(\top, c) \equiv ???$$



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$$wp(\perp, c) \equiv ???$$



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$$wp(\perp, c) \equiv \text{All states for which } c \text{ does not terminate}$$



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$$wp(\perp, c) \equiv \text{All states for which } c \text{ does not terminate}$$

$$wp(\perp, \text{assume}(x > 0)) \equiv ???$$



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$$wp(\top, c) \equiv \top$$

$$wp(\perp, c) \equiv \text{All states for which } c \text{ does not terminate}$$

$$\begin{aligned}wp(\perp, \text{assume}(x > 0)) &\equiv \forall x'. x > 0 \wedge x' = x \rightarrow \perp \\ &\equiv x \leq 0\end{aligned}$$



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## ASSIGNMENT STATEMENT

- $wp(F, x:=e) \triangleq F[e/x]$



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$$wp(F, x:=e) \triangleq \forall V'. \rho(x:=e) \rightarrow F[V'/V]$$

$$\equiv \forall V'. x' = e \wedge \text{frame}(x') \rightarrow F[V'/V]$$

$$\equiv F[e/x]$$



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$$\equiv F[e/x]$$

### EXAMPLES:

- $wp(x = 5, x:=6) \equiv ???$



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### EXAMPLES:

- $wp(x = 5, x:=6) \equiv \perp$



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### EXAMPLES:

- $wp(x = 5, x:=6) \equiv \perp$
- $wp(x = 5, x:=5) \equiv ???$



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### EXAMPLES:

- $wp(x = 5, x:=6) \equiv \perp$
- $wp(x = 5, x:=5) \equiv \top$



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### EXAMPLES:

- $wp(x = 5, x:=6) \equiv \perp$
- $wp(x = 5, x:=5) \equiv \top$
- $wp(x > 5, x:=y+1) \equiv ???$



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- $wp(F, x:=e) \triangleq F[e/x]$

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### EXAMPLES:

- $wp(x = 5, x:=6) \equiv \perp$
- $wp(x = 5, x:=5) \equiv \top$
- $wp(x > 5, x:=y+1) \equiv x > 5[(y+1)/x] \equiv y > 4$



# WEAKEST PRE-CONDITION

## HAVOC, ASSUME

- $wp(F, x:=havoc) \equiv \forall x. F$

$$\begin{aligned} wp(F, x:=havoc) &\triangleq \forall V'. frame(x) \rightarrow F[V'/V] \\ &\equiv \forall x'. F[x'/x] \equiv \forall x. F \end{aligned}$$

- $wp(F, assume(G)) \equiv ???$



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## HAVOC, ASSUME

- $wp(F, x:=havoc) \equiv \forall x. F$

$$\begin{aligned} wp(F, x:=havoc) &\triangleq \forall V'. frame(x) \rightarrow F[V'/V] \\ &\equiv \forall x'. F[x'/x] \equiv \forall x. F \end{aligned}$$

- $wp(F, assume(G)) \equiv G \rightarrow F$

$$\begin{aligned} wp(F, assume(G)) &\triangleq \forall V'. G \wedge frame(\emptyset) \rightarrow F[V'/V] \\ &\equiv \forall V'. G \rightarrow F \equiv G \rightarrow F \end{aligned}$$