

# INTRODUCTION TO ROUTING PROTOCOLS

(1)

## why ROUTING PROTOCOLS ?

### STATIC ROUTING

- Too much overhead
- Must be configured manually
- Manual configurations on all routers
- No account for link failure
- Not used for large n/w
- Every n/w requires new static route

AUTONOMOUS SYSTEM (AS) - is an internetwork under the control of single org<sup>n</sup>.

- One routing protocol cannot handle the task of updating the routing tables. For this Internet is divided into autonomous systems.

## ROUTING PROTOCOLS

### INTRA DOMAIN

- Routing inside an AS
- Protocols like RIP, EIGRP, OSPF

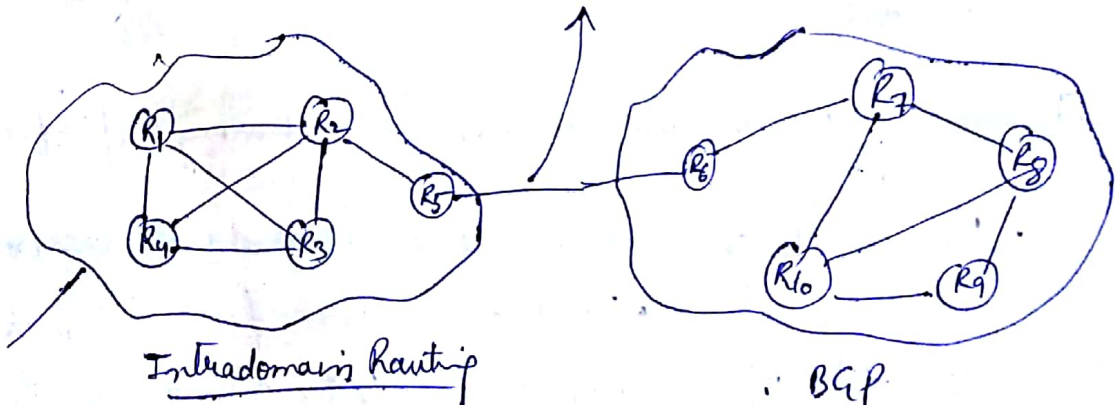
↓  
Interior Gateway Protocol

### INTER DOMAIN

- Routing b/w A.S
- Protocols BGP (Border Gateway Protocol)

↓  
Exterior Gateway Protocol

Interdomain Routing



RIP  
OSPF  
EIGRP  
IGRP

# INTRADOMAIN ROUTING PROTOCOLS

## DISTANCE VECTOR

RIP

## LINK STATE ROUTING

OSPF

## DISTANCE VECTOR ROUTING

- In DVR, the least cost route b/w any two nodes is the route with minimum distance.
- Each node maintains a vector (table) of minimum distances to every node.

### DVR works in 3 STEPS

① Initialization - At the beginning, Each node can know only the distance b/w itself & its immediate neighbours..

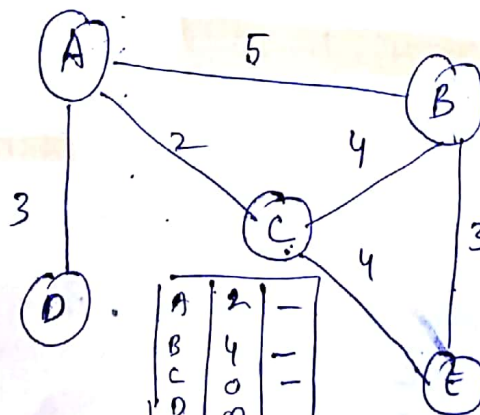
Below Figure shows the initial table for each node.

To	COST	NEXT
A	0	-
B	5	-
C	2	-
D	3	-
E	$\infty$	-

A's TABLE

To	COST	NEXT
A	3	-
B	$\infty$	-
C	$\infty$	-
D	0	-
E	$\infty$	-

D's TABLE



To	COST	NEXT
A	2	-
B	4	-
C	0	-
D	$\infty$	-
E	4	-

To	COST	NEXT
A	5	-
B	0	-
C	4	-
D	$\infty$	-
E	3	-

B's TABLE

To	COST	NEXT
A	$\infty$	-
B	3	-
C	4	-
D	$\infty$	-
E	0	-

- ② SHARING - whole idea of DVR is sharing of info b/w neighbours.
- How much of table must be shared with each neighbour?
  - The best solution for each node is to send its entire table to the neighbour & let the neighbour decide



updating - when a node receives a two column table from a neighbour, it needs to update its routing table.

(d) After receiving updates from C

To	COST
A	2
B	4
C	0
D	$\infty$
E	4

RECEIVED FROM C

To	COST	NEXT
A	4	C
B	6	C
C	2	C
D	$\infty$	C
E	6	C

As modified Table

To	COST	NEXT
A	0	-
B	5	-
C	2	-
D	3	-
E	$\infty$	-

As OLD TABLE

COMPARE

After receiving updates from B

To	COST	NEXT
A	5	B
B	0	-
C	4	B
D	$\infty$	B
E	3	B

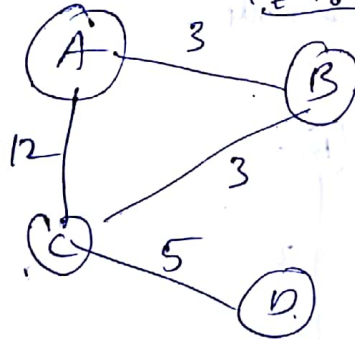
As Modified Table

A	0
B	5
C	2
D	3
E	$\infty$

To	COST	NEXT
A	0	-
B	5	-
C	2	-
D	3	-
E	6	C

As NEW TABLE

EXAMPLE -



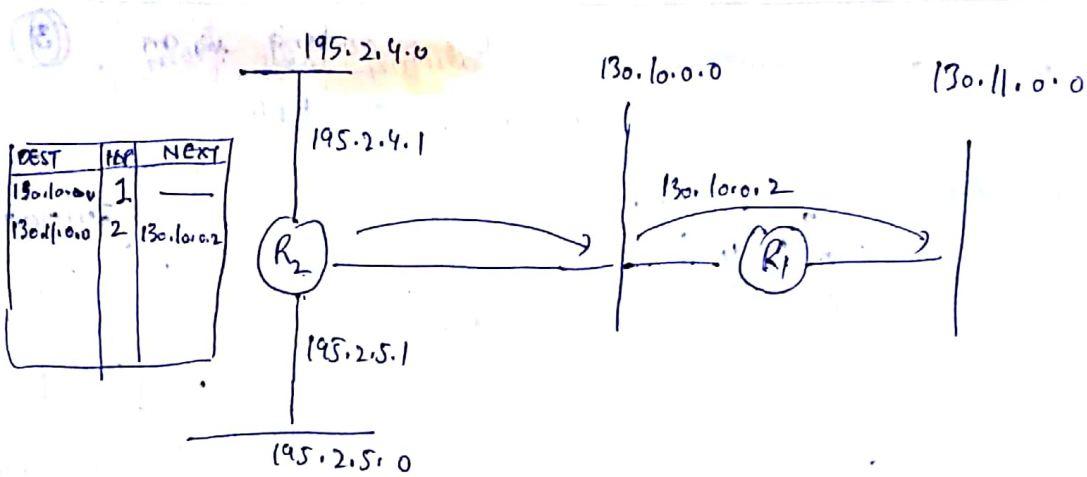
- Draw routing Table of node A
- ① after receiving updates from node C
  - ② after receiving updates from B
  - ③ Final routing table of A

## RIP

is an intradomain routing protocol based on DVR.

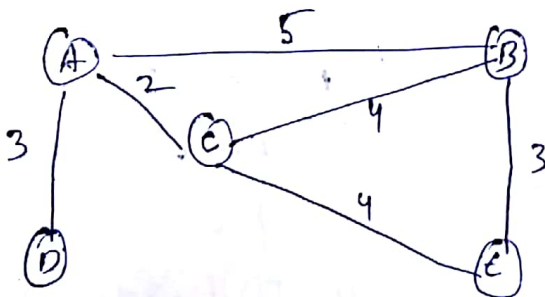
CONSIDERATIONS To keep while implementing DVR in RIP

- Metirc used in RIP is hop count i.e. no. of links to reach the destination.
- Any route in an AS using RIP cannot have more than 15 hops
- The next node column defines address of the router to which packet is to be sent to reach its destination.

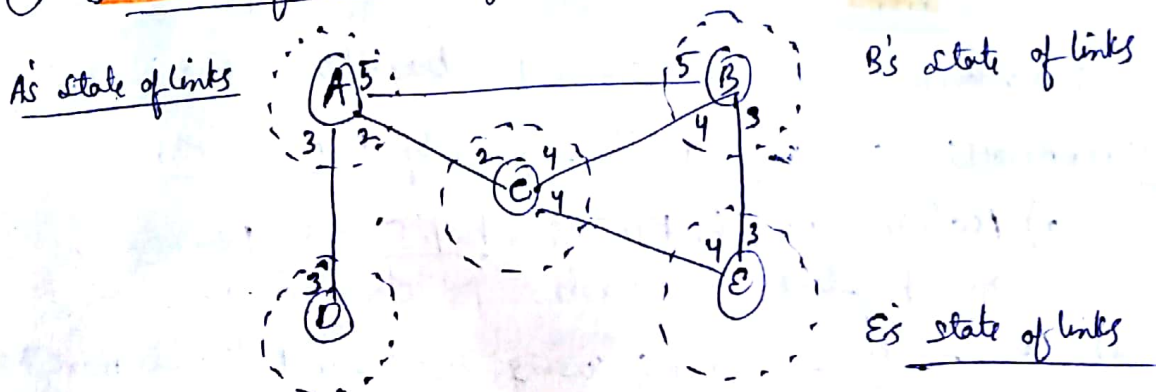


### LINK STATE ROUTING

- In link state routing, if each node in the domain has entire topology of the domain, the node can use Dijkstra's algorithm to build a routing table.
- In link state routing, 4 sets of actions are required to create routing table at each node.

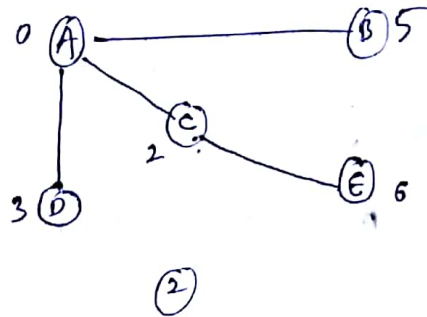
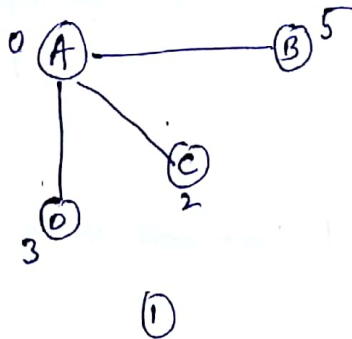


① Creation of states of the links by each node, called link state Packet (LSP)



dissemination of LSP's to every other router called Flooding in an efficient & reliable way (3)

(3) Formation of Shortest Path tree : Dijkstra's algorithm



(4) Calculation of Shortest Routing Table Based on the Shortest Path tree  
Each node uses the shortest path tree protocol to construct its routing table

Routing table for node A

Node	Cost	Next Router
A	0	—
B	5	—
C	2	—
D	3	—
E	6	C

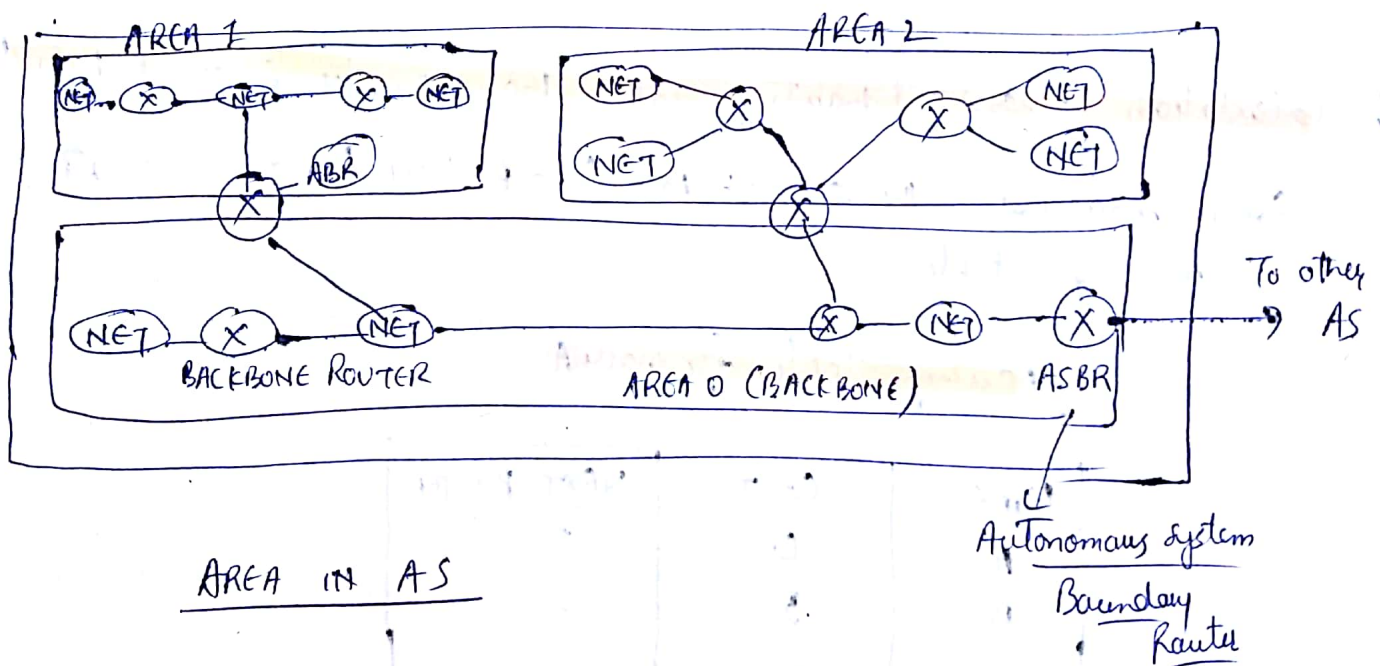
OSPF - is an intradomain routing Protocol based on link state routing.

Area - OSPF divided an autonomous system into areas. An area is a collection of n/w, hosts & routers all within an autonomous system.

• An autonomous system can be divided into different areas



- Routers inside an area flood area with routing inform<sup>n</sup>.
- At the Border of an area, special routers called area border routers disseminate inform<sup>n</sup> to other areas.
- Among the areas inside AS is a special area called backbone. Backbone serves as a Primary area & other areas called Secondary areas.
- Routers inside the backbone area are called Backbone routers.



## How LINK STATE ROUTING IS BETTER THAN DISTANCE VECTOR ROUTING

Distance vector Routing - designed to run on small n/w (usually fewer than 100 routers).

Example RIP & IGRP

Disadvantage - • DVP do not scale well because they require higher CPU & bandwidth utilization.

• They also take longer to converge than do link state protocols

convergence time is amount of time it takes to propagate changes in n/w topology

• Distance vector routing Protocols are great for small environment. but when it comes to enterprise n/w, you must deploy link state protocols

Link state Routing - designed to operate in large enterprise

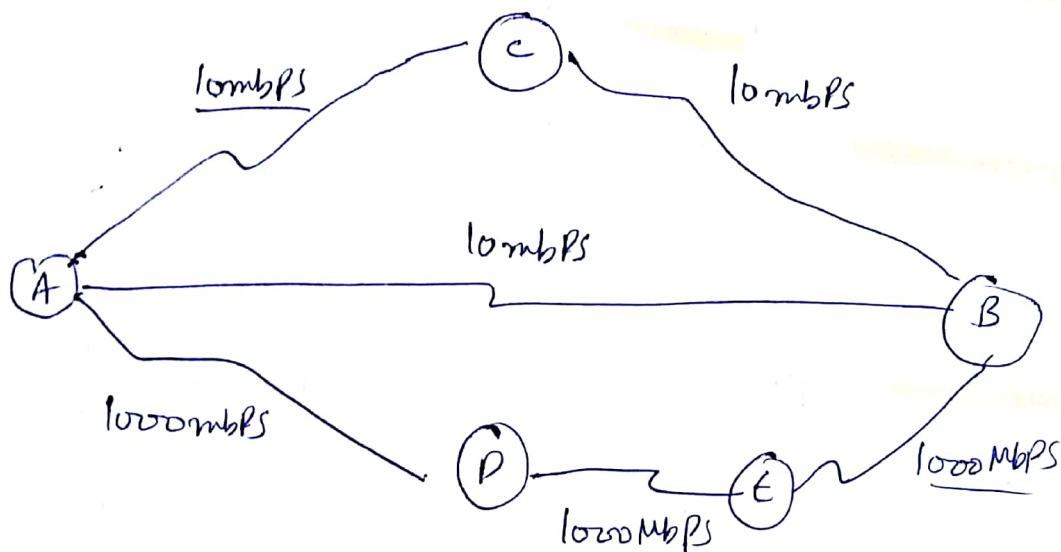
Example - OSPF

• This algorithm takes into account bandwidth as well as other factors when calculating the best path for a packet to traverse a n/w.

• Additionally link state convergence occurs faster than distance vector convergence.

in DVR Protocol - uses hop count as a metric to select the best path

in Link state Routing Protocol - uses bandwidth as a metric & develops a calculate metric to route the decisions



In the above topology, if we follow

a) Distance Vector Routing algorithm -

then optimal path is  $A \xrightarrow{10mbps} B$  because of less hop counts, but if we follow

b) Link state routing (OSPF) - then optimal path is

$A \xrightarrow{1000mbps} D \xrightarrow{1000mbps} E \xrightarrow{1000mbps} B$

because link state routing uses bandwidth as a metric to take routing decisions.