

## **OPERATING SYSTEM LAB (COM -312)**

### **Implementing Solution To Producer Consumer Problem Using Semaphore**

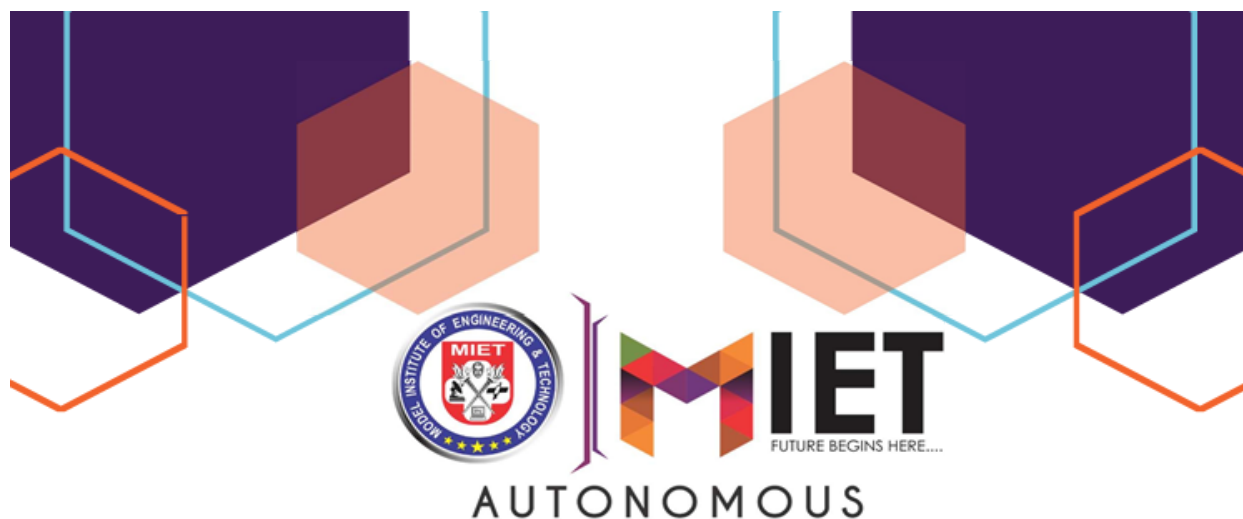
---

**CSE, MODEL INSTITUTE OF ENGINEERING AND  
TECHNOLOGY**

**BACHELOR OF ENGINEERING**  
**In**  
**Computer Science & Engineering**

**SUBMITTED BY**

Kashish Sharma , Harpreet Kour , Shivam Bhardwaj , Yatharth Choudhary  
(2021a1r058) , (2021a1r042) , (2021a1r060), (2021a1r040)



**SUBMITTED TO**

Department of Computer Science & Engineering  
Model Institute of Engineering and Technology (Autonomous)

Jammu, India

## ACKNOWLEDGEMENT

Under CRIE (Center for Research, Innovation and Entrepreneurship) we both worked under the guidance of mentors of CRIE, who guided us through a way leading to professionalism and practical hands on work experience. It is indeed with a great pleasure and immense sense of gratitude that we acknowledge the help of these individuals. We are highly indebted to our Director Ankur Gupta, “MODEL INSTITUTE OF ENGINEERING AND TECHNOLOGY”, for the facilities provided to accomplish this main project. We would like to thank our Prof. Ashok Kumar, Head of the Department of Computer Science and Engineering, MIET. For this constructive criticism throughout our project. We feel elated in manifesting our sense of gratitude for our internal project guide. Asst. Saurabh Sharma, Department of Computer Science and Engineering, MIET. He has been a constant source of inspiration for us and we are very deeply thankful to him.

## FACULTY AND MEMBER

**Dr. Ankur Gupta** - He is the Director at the Model Institute of Engineering and Technology, Jammu, India, besides being a Professor in the Department of Computer Science and Engineering. Prior to joining academia, he worked as Technical Team Lead at Hewlett Packard, developing software in the network management and e-Commerce domains. He has 2 patents to his name and 20 patents pending. He obtained his B.E, Hons. He is a senior member of the ACM, senior member IEEE and a life-member of the Computer Society of India. He has received competitive grants over Rs 2 crores from various funding agencies and faculty awards from IBM and EMC.

**Prof. Ashok Kumar** - He has over 21 years of experience in industry, academics, research and academic administration. He did his Doctorate from IKG Punjab Technical University, Jalandhar; Master of Engineering from Punjab Engineering College, Chandigarh and Bachelor's degree from NIT, Calicut, Kerala. He is Fellow of Institution of Electronics and Telecommunication Engineers (IETE). His research interest areas are Optical Networks, Electronics Product Design, Wireless Sensor Networks and Digital Signal Processing.. He is reviewer of many reputed national and international journals. He has also coordinated many national and international conferences.

**Mr. Saurabh Sharma** - He is working as Assistant Professor in the Department of Computer Science at Model Institute of Engineering and Technology, Jammu and has over 10 years of experience in academics and research. He has done his M. Tech CSE in Web Mining from Maharishi Markandeshwar University, Mullana, Ambala (HR) in the year 2011. His research interest areas are Web Mining, Data Mining, Machine Learning, ANN, Pattern Classification and Object Identification. He has, to his credit, 25 research papers published in journals of national and international repute and he has guided 11 M. Tech Dissertations.

## ABSTRACT

This paper represents the working of the two processes namely producer and consumer Working without and without Semaphores .The main contribution of our work is the facility to run the simulation in a multi-processors environment. This has been deemed appropriate since most of the current applications are distributed in nature and run in multi-processing environment.The. report describes the usage of Semaphores in solving the critical section problems that occur during the producer consumer functioning.

### Keywords-

Bounded Buffer problem,producer,consumer  
Synchronisation ,Semaphores,multiprocesing .

## CONTENTS

Acknowledgement	i
Faculty and member	ii
Abstract	iii
Contents	iv
Introduction	v
Problem	vi
Working	vii
Algorithm	viii
Flowchart	ix

Pseudo Code	x
Implementation	xi
a) Coding	
b) Output	
References	xii

## 1. Introduction

The Producer –Consumer problem is a family of problems described by Edsger W. Dijkstra since 1965. The Producer-Consumer problem is a classic problem this is used for multi-process synchronization, i.e. synchronization between more than one processes. Process Synchronization can be defined as the coordination between two process that have access to common materials such as a common section of code, resources or data etc. In the producer-consumer problem, there is one Producer that is producing something and there is one Consumer that is consuming the products produced by the Producer. a producer and a consumer coupled via a buffer with unbounded capacity .

## 2. Producer-consumer Problem

There is one Producer in the producer-consumer problem , Producer is producing some items, whereas there is one Consumer that is consuming the items produced by the Producer. The same memory buffer is shared by both producers and consumers which is of fixed-size. The Producer produces an item, put it into the memory buffer, and again start producing items. Whereas the task of the Consumer is to consume the item from the memory buffer The producer is a cyclic process and each time it goes through its cycle it produces a certain portion of information, that has to be processed by the consumer. The consumer is also a cyclic process and each time it goes through its cycle, it can process the next portion of information, as has been produced by the producer .We assume the two processes to be connected for a purpose via a buffer with unbounded capacity. The consumer waits when the buffer is empty; Producer waits when the buffer is full. The below flowchart describes the functioning of the whole process.

## 3. Working–

The task of the Producer is to produce the item, put it into the memory buffer, and again start producing items .The working of Producer is explained in the following code--

```

int count=0;
void producer(void)
{
int itemP;
while(1)
{
Producer_item(itemP);
while(count==n);//buffer is full
buffer[in]=itemP;
in=(in+1)mod n;
count=count+1;
}
}

```

In this code, Rp is a register which keeps the value of m[count].

-Rp is incremented (As element has been added to buffer)

-Then, the Incremented value of Rp is stored back to m[count]

Fig 1

<b>0</b>	<b>A</b>
<b>1</b>	<b>B</b>
<b>2</b>	<b>C</b>
<b>3</b>	<b>D</b>
<b>4</b>	<b>E</b>
<b>5</b>	
<b>6</b>	
<b>7</b>	

As we can see from the above figure, Buffer has total 8 spaces out of which the first 5 are filled, in = 5(pointing next empty position) and out = 0(pointing first filled position).

The producer here ,who wanted to produce an item F according to code it will enter into the producer() function, while(1) will always be true, itemP = F will be tried to insert into the buffer, before that while(count == n); will evaluate to be False

Buffer[in] = itemP → Buffer[5] = F. ( F is inserted now)

$in = (in + 1) \bmod n \rightarrow (5 + 1) \bmod 8 \rightarrow 6$ , therefore in = 6; (next empty buffer)

After insertion of F, Buffer looks like this-

<b>0</b>	<b>A</b>
<b>1</b>	<b>B</b>
<b>2</b>	<b>C</b>
<b>3</b>	<b>D</b>
<b>4</b>	<b>E</b>
<b>5</b>	<b>F</b>
<b>6</b>	
<b>7</b>	

Fig 1.1

Since  $\text{count} = \text{count} + 1$ ; is divided into three parts:

Load  $R_p, m[\text{count}] \rightarrow$  will copy count value which is 5 to register  $R_p$ .

Increment  $R_p \rightarrow$  will increment  $R_p$  to 6.

Suppose just after Increment and before the execution of third line (store  $m[\text{count}], R_p$ ), Context Switching occurs and code jumps to consumer code.

## 4. CONSUMER -

The task of the Consumer is to consume the item from the memory buffer.

Data can only be consumed by the consumer if and only if the memory buffer is not empty. In case it is found that the buffer is empty, the consumer is not allowed to use any data from the memory buffer.

Accessing memory buffer should not be allowed to producer and consumer at the same time.

The Consumer code is given below-

```
Void consumer(void)
```

```
{
```

```
int itemc;
```

```
while(1)
{
while(count==0);
itemc=Buffer(out);
out=(out+1)mod n;

count=count+1;
Process_item(itemc);
}
}
```

In the above code,

Rc is a register which keeps the value of m[count]

- Rc is decremented (As element has been removed out of buffer)
- The decremented value of Rc is stored back to m[count].

The Consumer who wanted to consume the first element , according to code it will enter into the consumer() function, while(1) will always be true, while(count == 0); will evaluate to be False( since the count is still 5, which is not equal to 0)

itemC = Buffer[out] → itemC = A ( since out is 0)

out = (out + 1) mod n → (0 + 1) mod 8 → 1, therefore out = 1 ( first filled position)

So, A is removed now. After removal of A, Buffer look like this

Fig 2

<b>0</b>	
<b>1</b>	<b>B</b>
<b>2</b>	<b>C</b>
<b>3</b>	<b>D</b>
<b>4</b>	<b>E</b>
<b>5</b>	<b>F</b>
<b>6</b>	
<b>7</b>	

Since count = count - 1; is divided into three parts:

Load Rc, m[count] → will copy count value which is 5 to register Rp.

Decrement Rc → will decrement Rc to 4.

store m[count], Rc → count = 4.

Now the current value of count is 4

Suppose after this, Context Switch occurs back to the leftover part of producer code.

Since context switch at producer code was occurred after Increment and before the execution of the third line (store m[count], Rp). So we resume from here since Rp holds 6 as incremented value

Hence store m[count], Rp → count = 6

Now the current value of count is 6, which is wrong as Buffer has only 5 elements, Since context switch at producer code was occurred after Increment and before the execution of the third line (store m[count], Rp)

This condition is known as Race Condition and Problem is Producer-Consumer Problem.

## **Problems In Producer-Consumer Working**

When the producer wants to put a new item in the buffer, and it is already full. So here the solution for the producer is to go to sleep and to be awakened when the consumer has removed one or more items. Similarly, if the consumer wants to remove an item from the buffer and sees that the buffer is empty, it goes to sleep until the producer puts something in the buffer and wakes it up. This approach sounds simple enough, but it leads to race conditions. So these problems need to be solved.

### **Solution-**

The producer should produce data only when the buffer is not full. If the buffer is full, then the producer should not be allowed to put any data into the buffer.

The consumer should consume data only when the buffer is not empty. If the buffer is empty, then the consumer should not be allowed to take any data from the buffer.

The producer and consumer should not access the buffer at the same time.

The problems in Producer –Consumer working can

be solved with the help of Semaphores.

### **Semaphores-**

Semaphores are integer variables that are used to solve the critical section problem. Semaphore cannot be implemented in the user mode because race condition may always arise when two or more processes try to access the variable simultaneously. It always needs support from the operating system to be implemented.

According to the demand of the situation, Semaphore can be divided into two categories.



1. Counting Semaphore
2. Binary Semaphore or Mutex

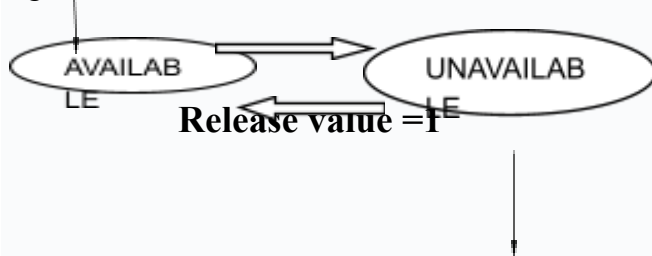
## Binary Semaphores

In these type of Semaphores the integer value of the semaphore can only be either 0 or 1. If the value of the Semaphore is 1, it means that the process can proceed to the critical section (the common section that the processes need to access). However, if the value of binary semaphore is 0, then the process cannot continue to the critical section of the code. When a process is using the critical section of the code, we change the Semaphore value to 0, and when a process is not using it, or we can allow a process to access the critical section, we change the value of semaphore to 1. Binary semaphore is also called mutex lock.

Initial value=0

Initial value=1      Release value =0

Fig3



## Counting Semaphores

Counting semaphores are signaling integers that can take on any integer value. Using these Semaphores we can coordinate access to resources and here the Semaphore count is the number of resources available. If the value of the Semaphore is anywhere above 0, processes can access the critical section, or the shared resources. The number of processes that can access the resources / code is the value of the semaphore. However, if the value is 0, it means that there aren't any resources that are available or the critical section is already being accessed by a number of processes and cannot be accessed by more processes. Counting semaphores are generally used when the number of instances of a resource are more than 1, and multiple processes can access the resource. The Semaphores used in solution to Producer-Consumer problem are-

### Full()-

The full variable is used to track the space filled in the buffer by the Producer process. It is initialized to 0 initially as initially no space is filled by the Producer process.

### Empty()-

The Empty variable is used to track the empty space in the buffer. The Empty variable is initially initialized to the buffer size as initially, the whole buffer is empty.

### Mutex()-

Mutex is used to achieve mutual exclusion. Mutex ensures that at any particular time only the producer or the consumer is accessing the buffer. Mutex is a binary semaphore variable that has a value of 0 or 1.

We will use the Signal() and wait() operation in the above-mentioned semaphores to arrive at a solution to the Producer-Consumer problem.

**Signal()** - The signal function increases the semaphore value by 1.

**Wait()** - The wait operation decreases the semaphore value by 1.

Mutex is used to solve the producer-consumer problem as mutex helps in mutual exclusion. It prevents more than one process to enter the critical section. As mutexes have binary values i.e 0 and 1. So whenever any process tries to enter the critical section code it first checks for the mutex value by using the wait operation.

**wait(mutex)** decreases the value of mutex by 1. So, Suppose a process P1 tries to enter the critical section when mutex value is 1. P1 executes wait(mutex) and decreases the value of mutex. Now, the value of mutex becomes 0 when P1 enters the critical section of the code.

Now, suppose Process P2 tries to enter the critical section then it will again try to decrease the value of mutex. But the mutex value is already 0. So, wait(mutex) will not execute, and P2 will now keep waiting for P1 to come out of the critical section. Now, P2 comes out of the critical section by executing signal(mutex).

**signal(mutex)** increases the value of mutex by 1. mutex value again becomes 1. Now, the process P2 which was in a busy-waiting state will be able to enter the critical section by executing wait(mutex).

So, mutex helps in the mutual exclusion of the processes.

The Producer Code using semaphore-

```
void Producer(){
    while(true){
        // producer produces an item/data
        wait(Empty);
        wait(mutex);
        add();
        signal(mutex);
        signal(Full);
    }
}
```

Understanding the above Producer process code :

**wait(Empty)** - Before producing items, the producer process checks for the empty space in the buffer. If the buffer is full producer process waits for the consumer process to consume items from the buffer. so, the producer process executes wait(Empty) before producing any item.

**wait(mutex)** - Only one process can access the buffer at a time. So, once the producer process enters into the critical section of the code it decreases the value of mutex by executing wait(mutex) so that no other process can access the buffer at the same time.

**add()** - This method adds the item to the buffer produced by the Producer process. once the Producer process reaches add function in the code, it is guaranteed that no other process will be able to access the shared buffer concurrently which helps in data consistency.

**signal(mutex)** - Now, once the Producer process added the item into the buffer it increases the mutex value by 1 so that other processes which were in a busy-waiting state can access the critical section.

**signal(Full)** - When the producer process adds an item into the buffer spaces is filled by one item so it increases the Full semaphore so that it indicates the filled spaces in the buffer correctly.

The Consumer code using Semaphore is -

```
void Consumer() {
    while(true){
        // consumer consumes an item
```

```

wait(Full);
wait(mutex);
consume();
signal(mutex);
signal(Empty);
}}

```

In the above Consumer process code :

`wait(Full)` - Before the consumer process starts consuming any item from the buffer it checks if the buffer is empty or has some item in it. So, the consumer process creates one more empty space in the buffer and this is indicated by the full variable. The value of the full variable decreases by one when the `wait(Full)` is executed. If the Full variable is already zero i.e the buffer is empty then the consumer process cannot consume any item from the buffer and it goes in the busy-waiting state.

`wait(mutex)` - It does the same as explained in the producer process. It decreases the mutex by 1 and restricts another process to enter the critical section until the consumer process increases the value of mutex by 1.

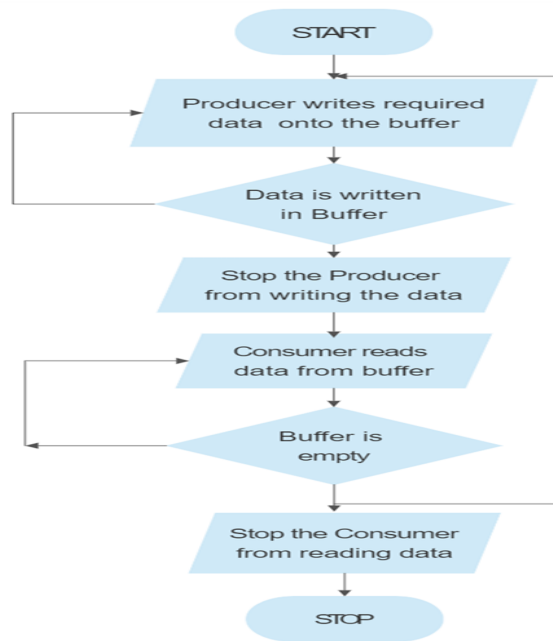
`consume()` - This function consumes an item from the buffer. when code reaches the consuming () function it will not allow any other process to access the critical section which maintains the data consistency.

`signal(mutex)` - After consuming the item it increases the mutex value by 1 so that other processes which are in a busy-waiting state can access the critical section now.

`signal(Empty)` - when a consumer process consumes an item it increases the value of the Empty variable indicating that the empty space in the buffer is increased by 1.

In the above section in both the Producer process code and consumer process code, we have the wait and signal operation on mutex which helps in mutual exclusion and solves the problem of the Producer consumer process.

## 5. Flowchart

**Fig 4**

## 6. Pseudo Code

```

int count=0;
void producer(void)
{
  int itemP;
  while(1)
  {
    Producer_item(itemP);
    while(count==n);//buffer is full
    buffer[in]=itemP;
    in=(in+1)mod n;
    count=count+1;
  }
}

```

Void consumer(void)

```

{
  int itemc;

```

```
while(1)
{
while(count==0);
itemc=Buffer(out);
out=(out+1)mod n;

count=count+1;
Process_item(itemc);
}
}
```

The entire producer consumer process functioning is implemented in the below C code-

```

#include <pthread.h>
#include <semaphore.h>
#include <stdlib.h>
#include <stdio.h>

/*
This program provides a possible solution for producer-consumer problem using mutex and semaphore.
I have used 5 producers and 5 consumers to demonstrate the solution. You can always play with these values.
*/

#define MaxItems 5 // Maximum items a producer can produce or a consumer can consume
#define BufferSize 5 // Size of the buffer

sem_t empty;
sem_t full;
int in = 0;
int out = 0;
int buffer[BufferSize];
pthread_mutex_t mutex;

void *producer(void *pno)
{
    int item;
    for(int i = 0; i < MaxItems; i++) {
        item = rand(); // Produce an random item
        sem_wait(&empty);
        pthread_mutex_lock(&mutex);
        buffer[in] = item;
        printf("Producer %d: Insert Item %d at %d\n", *((int *)pno), buffer[in], in);
        in = (in+1)%BufferSize;
        pthread_mutex_unlock(&mutex);
        sem_post(&full);
    }
}

void *consumer(void *cno)
{
    for(int i = 0; i < MaxItems; i++) {
        sem_wait(&full);

```

Fig5

```

        sem_wait(&full);
        pthread_mutex_lock(&mutex);
        int item = buffer[out];
        printf("Consumer %d: Remove Item %d from %d\n", *((int *)cno), item, out);
        out = (out+1)%BufferSize;
        pthread_mutex_unlock(&mutex);
        sem_post(&empty);
    }
}

int main()
{
    pthread_t pro[5], con[5];
    pthread_mutex_init(&mutex, NULL);
    sem_init(&empty, 0, BufferSize);
    sem_init(&full, 0, 0);

    int a[5] = {1,2,3,4,5}; //Just used for numbering the producer and consumer

    for(int i = 0; i < 5; i++) {
        pthread_create(&pro[i], NULL, (void *)producer, (void *)&a[i]);
    }
    for(int i = 0; i < 5; i++) {
        pthread_create(&con[i], NULL, (void *)consumer, (void *)&a[i]);
    }

    for(int i = 0; i < 5; i++) {
        pthread_join(pro[i], NULL);
    }
    for(int i = 0; i < 5; i++) {
        pthread_join(con[i], NULL);
    }
    pthread_mutex_destroy(&mutex);
    sem_destroy(&empty);
    sem_destroy(&full);
    return 0;
}

```

Output-

```

Consumer 2: Remove Item 1350490027 from 1
Consumer 2: Remove Item 783368690 from 2
Consumer 2: Remove Item 1102520059 from 3
Consumer 2: Remove Item 2044897763 from 4
Producer 1: Insert Item 1967513926 at 0
Consumer 5: Remove Item 1967513926 from 0
Producer 2: Insert Item 1540383426 at 1
Producer 2: Insert Item 304089172 at 2
Producer 2: Insert Item 1303455736 at 3
Consumer 5: Remove Item 1540383426 from 1
Producer 2: Insert Item 35005211 at 4
Consumer 5: Remove Item 304089172 from 2
Producer 1: Insert Item 1365180540 at 0
Consumer 5: Remove Item 1303455736 from 3
Consumer 1: Remove Item 35005211 from 4
Producer 2: Insert Item 521595368 at 1
Consumer 5: Remove Item 1365180540 from 0
Producer 1: Insert Item 294702567 at 2
Producer 1: Insert Item 1726956429 at 3
Producer 1: Insert Item 336465782 at 4
Consumer 1: Remove Item 521595368 from 1
Consumer 1: Remove Item 294702567 from 2
Consumer 1: Remove Item 1726956429 from 3
Consumer 1: Remove Item 336465782 from 4

```

**Fig 6**

## References

<https://www.scaler.com>

<https://www.geeksforgeeks.org>

<https://afteracademy.com>

<https://www.javatpoint.com>

## Github:

<https://github.com/Kashish9086>

<https://github.com/HKour2003>

**KASHISH SHARMA ,HARPREET KOUR**

**EMAIL ID-**

[2021a1r058@mietjammu.in](mailto:2021a1r058@mietjammu.in)

(KASHISH SHARMA)

[2021a1r042@mietjammu.in](mailto:2021a1r042@mietjammu.in)

(HARPREET KOUR)

[2021a1r060@mietjammu.in](mailto:2021a1r060@mietjammu.in)

(SHIVAM BHARDWAJ)

[2021a1r040@mietjammu.in](mailto:2021a1r040@mietjammu.in)

(YATHARTH CHOUDHARY)