

PA4 due tomorrow  
PA5 Late/Resubmit → tomorrow

## Sorting Quickly

Quiz 4 Sort → in-place

```

public class SortQuickly {
    public static void swap(String[] array, int i1, int i2) {
        String temp = array[i1];
        array[i1] = array[i2];
        array[i2] = temp;
    }

    public static int partition(String[] array, int low, int high) {
        int pivotStartIndex = high - 1;
        String pivot = array[pivotStartIndex];
        int smallerBefore = low, largerAfter = high - 2;

        while (smallerBefore <= largerAfter) {
            if (array[smallerBefore].compareTo(pivot) < 0) {
                smallerBefore++;
            } else {
                swap(array, smallerBefore, largerAfter);
                largerAfter--;
            }
        }

        swap(array, smallerBefore, pivotStartIndex);
        return smallerBefore;
    }

    public static void qsort(String[] array, int low, int high) {
        if (high - low <= 1) { return; }
        int splitAt = partition(array, low, high);
        qsort(array, low, splitAt);
        qsort(array, splitAt + 1, high);
    }

    public static void sortD(String[] array) {
        qsort(array, 0, array.length);
    }

    public static void main(String[] args) {
        String[] str = {"f", "b", "a", "a", "d", "c"};
        int[] result = SortQuickly.sortD(str);
        System.out.println(Arrays.deepToString(result));
    }
}

```

Handwritten notes on the code:

- $O(N)$  3 [ (next to swap method)
- $N$  (next to partition method)
- $O(N)$  (next to qsort method)
- $O(N)$  (next to sortD method)
- $O(N)$  (next to main method)
- Annotations:  $N$ ,  $O$ ,  $N$  above partition parameters;  $N$ ,  $O$ ,  $N$  above qsort parameters;  $O$ ,  $6$  above sortD parameters.
- Diagram: A sequence of numbers 1, 2, 3, 4, 5 in boxes. Arrows point from 1, 2, 3, 4 to a box containing 5. A label 'ss' is near the first arrow, and 'la' is near the last arrow.

Draw the picture of sortD()

What is the tight bound of sortD:

Best case:  $O(N * \log_2(N)) \rightarrow$  median value at every levelWorst case:  $O(N^2) \rightarrow$  sorted arraylow = 0  
high = 6

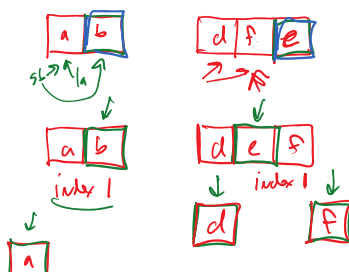
0	1	2	3	4	5
f	b	a	e	d	c

pivotIndex = 5  
pivot = "c"  
sb = 0, 1, 2  
la = 4, 5, 6

d	b	a	e	f	c
e	b	a	d	f	c
a	b	e	d	f	c

a	b	c	d	f	e
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returned index 2

 $N = 6$ 

height = 3

 $\hookrightarrow \log_2(N)$  $O(N * \log_2(N))$ 

worst case → sorted array

1	2	3	4	5
---	---	---	---	---

 $N = 5$ 

1	2	3	4	5
---	---	---	---	---

levels = 5

