

CSE12 - Lecture 11 - A00

Monday, October 17, 2022 8:00 AM

PA3 due tomorrow
Exam1 → Friday

Measuring Runtime

Count how many times each line executes, then say which $\Theta()$ statement(s) is(are) true.

```
int maxDifference(int[] arr){
    max = 0;
    for (int i=0; i<arr.length; i++) {
        for (int j=0; j<arr.length; j++) {
            if (arr[i] - arr[j] > max)
                max = arr[i] - arr[j];
        }
    }
    return max;
}
```

Assume $n = \text{arr.length}$

- A. $f(n) = \theta(2^n)$
- ☒ B. $f(n) = \theta(n^2)$
- C. $f(n) = \theta(n)$
- D. $f(n) = \theta(n^3)$
- E. Other/none/more



Count how many times each line executes, then say which $\Theta()$ statement(s) is(are) true.

```
int sumTheMiddle(int[] arr){
    int range = 100;
    int start = arr.length/2 - range/2;
    int sum = 0;
    for (int i=start; i<start+range; i++) {
        sum += arr[i];
    }
    return sum;
}
```

Assume $n = \text{arr.length}$

- A. $f(n) = \theta(2^n)$
- B. $f(n) = \theta(n^2)$
- C. $f(n) = \theta(n)$
- ☒ D. $f(n) = \theta(1)$
- E. None of these

$$6 + 300 = 306$$

$$f(n) = 306$$

$$C = 306$$

$$N_0 = 0$$

$$g(n) = 1$$

constant time

Big O Upper bound

$f(n) = O(g(n))$, $f(n) \leq c * g(n)$
for all $n \geq n_0$

Big Ω Lower bound

$f(n) = \Omega(g(n))$, $f(n) \geq c * g(n)$
for all $n \geq n_0$

Big Θ Tight bound

$f(n) = \Theta(g(n))$, $f(n) = c * g(n)$
for all $n \geq n_0$

For each function in the list below, it is related to the function below it by O , and the reverse is not true. That is, n is $O(n^2)$ but n^2 is not $O(n)$.

- $f(n) = 1/(n^2)$
- $f(n) = 1/n$
- $f(n) = 1$
- $f(n) = \log(n)$
- $f(n) = \sqrt{n}$
- $f(n) = n$
- $f(n) = n^2$
- $f(n) = n^3$
- $f(n) = n^4$
- ... and so on for constant polynomials ...
- $f(n) = 2^n$
- $f(n) = n!$
- $f(n) = n^n$

$$\text{range} = 100$$

$$\text{start} = \frac{N}{2} - 50$$

$$\text{start} + \text{range} = \frac{N}{2} - 50 + 100 = \frac{N}{2} + 50$$

$$N = 100$$

$$\text{start} = \frac{100}{2} - 50 = 0$$

$$\text{start} + \text{range} = \frac{100}{2} + 50 = 100 \rightarrow 100$$

$$N = 10000$$

$$\text{start} = \frac{10000}{2} - 50 = 4950$$

$$\text{start} + \text{range} = \frac{10000}{2} + 50 = 5050 \rightarrow 100$$

```
void printAllItemsTwice(int arr[], int size)
```

```
{
  for (int i = 0; i < size; i++) {
    printf("%d\n", arr[i]);
  }
  for (int i = 0; i < size; i++) {
    printf("%d\n", arr[i]);
  }
}
```

$$N \left[\begin{array}{l} 1 + (N+1) + N \\ N \end{array} \right] 3N+2 + N \left[\begin{array}{l} 1 + (N+1) + N \\ N \end{array} \right] 3N+2$$

What is the tight bound?

$$2N+4$$

$$C=6$$

$$N_0=4$$

$$g(N) = N$$

$$\Theta(N)$$

```
void printFirstItemThenFirstHalfThenSayHi100Times(int arr[], int size)
```

```
{
  printf("First element of array = %d\n", arr[0]);
  for (int i = 0; i < size/2; i++) {
    printf("%d\n", arr[i]);
  }
  for (int i = 0; i < 100; i++) {
    printf("Hi\n");
  }
}
```

$$\frac{N}{2} \left[\begin{array}{l} 1 + (\frac{N}{2}+1) + \frac{N}{2} \\ \frac{N}{2} \end{array} \right] \frac{3N}{2}+2 + 1 \left[\begin{array}{l} 1 + (100+1) + 100 \\ 100 \end{array} \right] 302$$

What is the tight bound?

$$\frac{3N}{2} + 304$$

$$\Theta(N)$$

```
void printAllNumbersThenAllPairSums(int arr[], int size)
```

```
{
  for (int i = 0; i < size; i++) {
    printf("%d\n", arr[i]);
  }
  for (int i = 0; i < size; i++) {
    for (int j = 0; j < size; j++) {
      printf("%d\n", arr[i] + arr[j]);
    }
  }
}
```

$$N \left[\begin{array}{l} 1 + (N+1) + N \\ N \end{array} \right] 3N+2 + N \left[\begin{array}{l} 1 + (N+1) + N \\ N \end{array} \right] \left[\begin{array}{l} 1 + (N+1) + N \\ N \end{array} \right] \rightarrow 2N+2 + N(3N+2)$$

$$N^2 + N$$

What is the tight bound?

$$5N+4 + 3N^2+2N$$

$$3N^2+7N+4$$

$$\Theta(N^2)$$

Selection Sort

```
import java.util.Arrays;
public class Sort {
    public static void sortA(int[] arr) {
        for(int i = 0; i < arr.length; i += 1) {
            System.out.print(Arrays.toString(arr) + " -> ");
            int minIndex = i;
            for(int j = i; j < arr.length; j += 1) {
                if(arr[minIndex] > arr[j]) { minIndex = j; }
            }
            int temp = arr[i];
            arr[i] = arr[minIndex];
            arr[minIndex] = temp;
            System.out.println(Arrays.toString(arr));
        }
    }
}
```

N

$\frac{N^2}{2}$

Selection Sort – what does it print out?

$N=5$

Sort.sortA(new int[]{ 53, 83, 15, 45, 49 });

[53, 83, 15, 45, 49] ->

$15 \mid 83 \ 53 \ 45 \ 49$ 5
 $15 \ 45 \mid 53 \ 83 \ 49$ 4
 $15 \ 45 \ 49 \mid 83 \ 53$ 3
 $15 \ 45 \ 49 \ 53 \mid 83$ 2
 $15 \ 45 \ 49 \ 53 \ 83 \mid$ 1
 $15 \ 45 \ 49 \ 53 \ 83$ 0

$\frac{N}{2}$

Worst case: reverse sorted array
 83 53 49 45 15

best case: sorted case
 15 45 49 53 83

What is the runtime? Consider the shape of the input array.

Worse case: $\Theta(N^2)$

Best case: $\Theta(N^2)$

is sorted ()
 $\Theta(N)$

Insertion Sort

```
import java.util.Arrays;

public class Sort {
    public static void sortB(int[] arr) {
        for(int i = 0; i < arr.length; i += 1) {
            System.out.print(Arrays.toString(arr) + " -> ");
            for(int j = i; j > 0; j -= 1) {
                if(arr[j] < arr[j-1]) {
                    int temp = arr[j-1];
                    arr[j-1] = arr[j];
                    arr[j] = temp;
                }
            }
            System.out.println(Arrays.toString(arr));
        }
    }
}
```

Handwritten notes: A large red bracket on the left side of the code, spanning from the first for loop to the end of the method, is labeled with a red N . A smaller red bracket on the right side, spanning the inner if loop, is labeled with a red $\frac{N}{2}$.

Insertion Sort – what does it print out?

Sort.sortB(new int[]{ 53, 83, 15, 45, 49 });

[53, 83, 15, 45, 49] -> 53 83 15 45 49
 53 83 15 45 49
 15 53 83 45 49
 15 45 53 83 49
 15 45 49 53 83

15 45 49 53 83

What is the runtime? Consider the shape of the input array.

Worse case: $O(N^2)$

Best case: $O(N^2)$